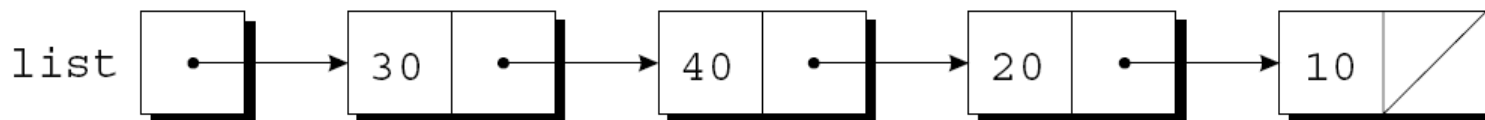


Lecture 17 - Advanced Uses of Pointers

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Dynamic Storage Allocation

- C's data structures, including arrays, are normally fixed in size.
- Fixed-size data structures can be a problem, since we're forced to choose their sizes before executing a program.
- Fortunately, C supports **dynamic storage allocation**: the ability to allocate storage during program execution.
- Using dynamic storage allocation, we can design data structures that grow (and shrink) as needed.

Dynamic Storage Allocation (cont.)

- Dynamic storage allocation is **used most often for strings, arrays, and structures.**
- Dynamically allocated structures **can be linked together to form lists, trees, and other data structures.**
- Dynamic storage allocation **is done by calling a memory allocation function.**

Memory Allocation Functions

- The `<stdlib.h>` header declares three memory allocation functions:

`malloc`—Allocates a block of memory but **doesn't initialize it**.

`calloc`—Allocates a block of memory and **clears it**.

`realloc`—**Resizes** a previously allocated block of memory.

- **These functions return** a value of type `void *` (a “generic” pointer).

Null Pointers

- If a memory allocation function **can't locate** a memory block of the requested size, it **returns a *null pointer***.
- A **null pointer is a special value** that can be distinguished from all valid pointers.
- After we've **stored the function's return value** in a pointer variable, we must test **to see if it's a null pointer**.

Null Pointers (cont.)

- An example of testing `malloc`'s return value:

```
p = malloc(10000);  
if (p == NULL) {  
    /* allocation failed; take appropriate action */  
}
```

- `NULL` is a macro (defined in various library headers) that represents the null pointer.
- Some programmers combine the call of `malloc` with the `NULL` test:

```
if ((p = malloc(10000)) == NULL) {  
    /* allocation failed; take appropriate action */  
}
```

Null Pointers (cont.)

- Pointers test true or false in the same way as numbers.
- All non-null pointers test true; only null pointers are false.

- Instead of writing

```
if (p == NULL) ...  
we could write
```

```
if (!p) ...
```

- Instead of writing

```
if (p != NULL) ...  
we could write
```

```
if (p) ...
```

Dynamically Allocated Strings

- Dynamic storage allocation is often useful for working with strings.
- **Strings** are stored in character arrays, and it **can be hard to anticipate how long these arrays need to be.**
- **By allocating strings dynamically, we can postpone the decision until the program is running.**

Using `malloc` to Allocate Memory for a String

- Prototype for the `malloc` function:

```
void *malloc(size_t size);
```

- `malloc` **allocates** a block of **size bytes** and **returns a pointer** to it.
- `size_t` is an **unsigned integer** type defined in the library.

Using `malloc` to Allocate Memory for a String (cont.)

- A call of `malloc` that allocates memory for a string of `n` characters:

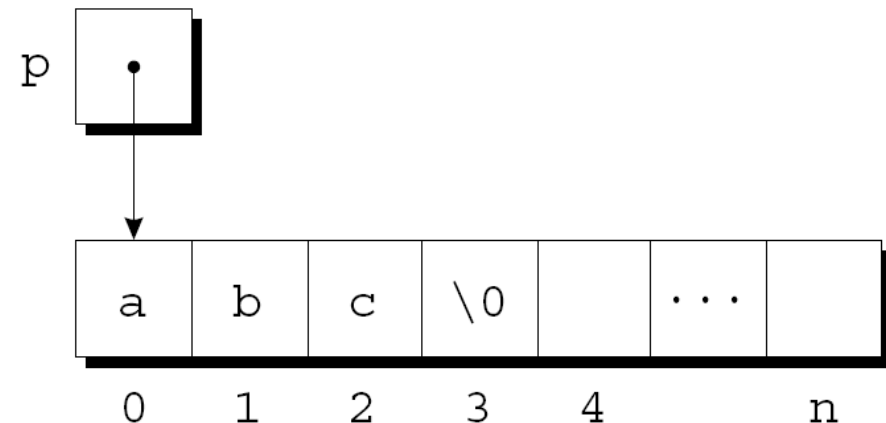
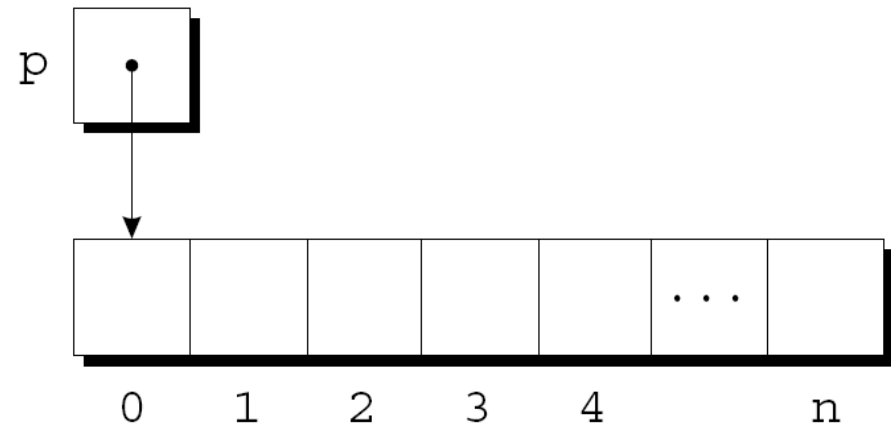
```
char * p;  
p = malloc(n + 1);
```

- Each character requires one byte of memory; adding 1 to `n` leaves room for the null character.
- Some programmers prefer to cast `malloc`'s return value, although the cast is not required:

```
p = (char *) malloc(n + 1);
```

Using `malloc` to Allocate Memory for a String (cont.)

- Memory allocated using `malloc` **isn't cleared**, so `p` will point to an **uninitialized array of $n + 1$ characters**:
- **Calling `strcpy` is one way to initialize** this array:
`strcpy(p, "abc");`
- The first four characters in the array will now be `a`, `b`, `c`, and `\0`:



Using Dynamic Storage Allocation in String Functions

- Dynamic storage allocation makes it possible to write functions that return a pointer to a “new” string.
- Consider the problem of **writing a function that concatenates two strings without changing either one.**
- The function will **measure the lengths of the two strings** to be concatenated, then **call `malloc` to allocate the right amount of space** for the result.

Using Dynamic Storage Allocation in String Functions (cont.)

```
char *concat(const char *s1, const char *s2)
{
    char *result;

    result = malloc(strlen(s1) + strlen(s2) + 1);
    if (result == NULL) {
        printf("Error: malloc failed in concat\n");
        exit(EXIT_FAILURE);
    }
    strcpy(result, s1);
    strcat(result, s2);
    return result;
}
```

Using Dynamic Storage Allocation in String Functions (cont.)

- A call of the `concat` function:

```
p = concat("abc", "def");
```

- After the call, **p will point to the string "abcdef"**, which is stored in a dynamically allocated array.

Using Dynamic Storage Allocation in String Functions (cont.)

- Functions such as `concat` that dynamically allocate storage **must be used with care**.
- **When the string that `concat` returns is no longer needed, we'll want to call the `free` function to release the space that the string occupies.**
- **If we don't, the program may eventually run out of memory.**

Program: Printing a One-Month Reminder List (Revisited)

- The `remind2.c` program is based on the `remind.c` program of Lecture 13, which **prints a one-month list of daily reminders**.
- The **original `remind.c` program stores reminder strings in a two-dimensional array of characters**.
- **In the new program, the array will be one-dimensional; its elements will be pointers to dynamically allocated strings**.

Program: Printing a One-Month Reminder List (Revisited) (cont.)

- **Advantages** of switching to dynamically allocated strings:
 - Uses space more efficiently by allocating the exact number of characters needed to store a reminder.
 - Avoids calling `strcpy` to move existing reminder strings in order to make room for a new reminder.
- Switching from a two-dimensional array to an array of pointers requires **changing only eight lines** of the program (shown in **bold**).

Program: Printing a One-Month Reminder List (Revisited) (cont.)

remind2.c

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

#define MAX_REMIND 50    /* maximum number of reminders */
#define MSG_LEN 60      /* max length of reminder message */

int read_line(char str[], int n);
int main(void)
{
    char *reminders[MAX_REMIND];
    char day_str[3], msg_str[MSG_LEN+1];
    int day, i, j, num_remind = 0;
```

```
char reminders[MAX_REMIND][MSG_LEN+3];
```

Program: Printing a One-Month Reminder List (Revisited) (cont.)

```
for (;;) {
    if (num_remind == MAX_REMIND) {
        printf("-- No space left --\n");
        break;
    }
    printf("Enter day and reminder: ");
    scanf("%2d", &day);
    if (day == 0)
        break;
    sprintf(day_str, "%2d", day);
    read_line(msg_str, MSG_LEN);

    for (i = 0; i < num_remind; i++)
        if (strcmp(day_str, reminders[i]) < 0)
            break;
    for (j = num_remind; j > i; j--)
        reminders[j] = reminders[j-1];
```

```
strcpy(reminders[j],
       reminders[j-1]);
```

Program: Printing a One-Month Reminder List (Revisited) (cont.)

```
reminders[i] = malloc(2 + strlen(msg_str) + 1);  
if (reminders[i] == NULL) {  
    printf("-- No space left --\n");  
    break;  
}
```

```
strcpy(reminders[i], day_str);  
strcat(reminders[i], msg_str);
```

```
num_remind++;  
}
```

```
printf("\nDay Reminder\n");  
for (i = 0; i < num_remind; i++)  
    printf(" %s\n", reminders[i]);
```

```
return 0;
```

Dynamically Allocated Arrays

- The **close relationship between arrays and pointers** makes a **dynamically allocated array** as **easy to use** as an ordinary array.
- Suppose a program **needs an array of n integers**, where n is computed during program execution.
- We'll **first declare a pointer variable**:

```
int *a;
```
- **Once the value of n is known**, the program can **call `malloc` to allocate space** for the array:

```
a = malloc(n * sizeof(int));
```
- **Always use the `sizeof` operator** to calculate the amount of space required for each element.

Using `malloc` to Allocate Storage for an Array (cont.)

- We can now ignore the fact that `a` is a pointer and use it instead as an array name, thanks to the relationship between arrays and pointers in C.
- For example, we could use the following loop to initialize the array that `a` points to:

```
for (i = 0; i < n; i++)  
    a[i] = 0;
```
- We also have the option of using pointer arithmetic instead of subscripting to access the elements of the array.

The `calloc` Function

- The `calloc` function is an **alternative to** `malloc`.
- Prototype for `calloc`:

```
void *calloc(size_t nmemb, size_t size);
```
- Properties of `calloc`:
 - Allocates space for an array with `nmemb` **elements**, **each of which** is `size` **bytes long**.
 - **Returns a null pointer if** the requested **space isn't available**.
 - **Initializes** allocated memory by setting **all bits to 0**.

The `calloc` Function (cont.)

- A call of `calloc` that allocates space for **an array of `n` integers**:

```
a = calloc(n, sizeof(int));
```

- By calling `calloc` with 1 as its first argument, we can **allocate space for a data item of any type**:

```
struct point { int x, y; } *p;
```

```
p = calloc(1, sizeof(struct point));
```


The `realloc` Function

- The `realloc` function can resize a dynamically allocated array.
- Prototype for `realloc`:

```
void *realloc(void *ptr, size_t size);
```
- `ptr` must point to a memory block obtained by a previous call of `malloc`, `calloc`, or `realloc`.
- `size` represents the new size of the block, which may be larger or smaller than the original size.

```
q = realloc(p, 20000);
```

The `realloc` Function (cont.)

- Properties of `realloc`:
 - When it expands a memory block, `realloc` doesn't initialize the bytes that are added to the block.
 - If `realloc` can't enlarge the memory block as requested, it returns a null pointer; the data in the old memory block is unchanged.
 - If `realloc` is called with a null pointer as its first argument, it behaves like `malloc`.
 - If `realloc` is called with 0 as its second argument, it frees the memory block.

The `realloc` Function (cont.)

- We expect `realloc` to be reasonably efficient:
 - When asked to reduce the size of a memory block, `realloc` should shrink the block “in place.”
 - `realloc` should always attempt to expand a memory block without moving it.
- If it can't enlarge a block, `realloc` will allocate a new block elsewhere, then copy the contents of the old block into the new one.
- Once `realloc` has returned, be sure to update all pointers to the memory block in case it has been moved.

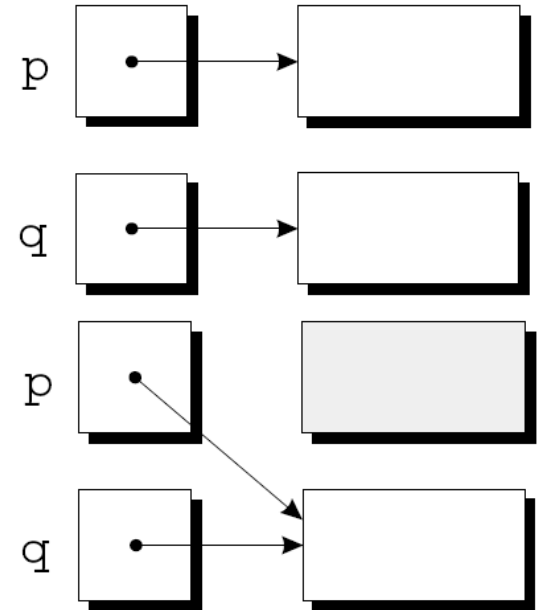
Deallocating Storage

- `malloc` and the other memory allocation functions obtain memory blocks from a storage pool known as the **heap**.
- Calling these functions too often—or asking them for large blocks of memory—can exhaust the heap, causing the functions to return a null pointer.
- To make matters worse, a program may allocate blocks of memory and then lose track of them, thereby wasting space.

Deallocating Storage (cont.)

```
p = malloc(...);  
q = malloc(...);  
p = q;
```

- A snapshot **after the first two statements** have been executed:
- **After q is assigned to p , both variables now point to the second memory block.**
- There are **no pointers to the first block**, so we'll never be able to use it again.



Deallocating Storage (cont.)

- A block of memory that's **no longer accessible** to a program is **said to be *garbage***.
- A program that **leaves garbage behind** has a ***memory leak***.
- Some languages provide a ***garbage collector*** that automatically locates and recycles garbage, **but C doesn't**.
- Instead, **each C program is responsible for recycling its own garbage by calling the `free` function** to release unneeded memory.

The **free** Function

- Prototype for `free`:

```
void free(void *ptr);
```

- `free` will be passed a pointer to an unneeded memory block:

```
p = malloc(...);
```

```
q = malloc(...);
```

```
free(p);
```

```
p = q;
```

- Calling `free` releases the block of memory that `p` points to.

The “Dangling Pointer” Problem

- Using `free` leads to a new problem: ***dangling pointers***.
- `free(p)` deallocates the memory block that `p` points to, but doesn't change `p` itself.
- If we forget that `p` no longer points to a valid memory block, chaos may ensue:

```
char *p = malloc(4);  
...  
free(p);  
...  
strcpy(p, "abc");    /** WRONG **/
```

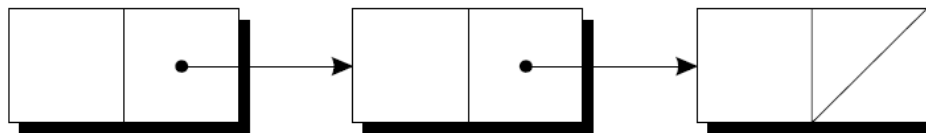
- Modifying the memory that `p` points to is a **serious error**.

The “Dangling Pointer” Problem (cont.)

- Dangling pointers can be hard to spot, since several pointers may point to the same block of memory.
- When the block is freed, all the pointers are left dangling.

Linked Lists

- Dynamic storage allocation is especially useful for building lists, trees, graphs, and other linked data structures.
- A **linked list** consists of a chain of structures (called **nodes**), with each node containing a pointer to the next node in the chain:



- The last node in the list contains a null pointer.

Linked Lists (cont.)

- A linked list is more flexible than an array: we can easily insert and delete nodes in a linked list, allowing the list to grow and shrink as needed.
- On the other hand, we lose the “random access” capability of an array:
 - Any element of an array can be accessed in the same amount of time.
 - Accessing a node in a linked list is fast if the node is close to the beginning of the list, slow if it's near the end.

Declaring a Node Type

- To set up a linked list, we'll **need a structure** that **represents a single node**.
- A node structure will contain data (an integer in this example) plus a pointer to the next node in the list:

```
struct node {  
    int value;           /* data stored in the node */  
    struct node *next;  /* pointer to the next node */  
};
```

Declaring a Node Type (cont.)

- Next, **we'll need a variable** that **always points to the first node** in the list:
`struct node *first = NULL;`
- **Setting `first` to `NULL`** indicates that **the list is initially empty**.

Creating a Node

- As we construct a linked list, we'll **create nodes one by one**, adding each to the list.
- Steps involved in creating a node:
 1. **Allocate memory** for the node.
 2. **Store data** in the node.
 3. **Insert the node into the list.**
- We'll concentrate on the first two steps for now.

Creating a Node (cont.)

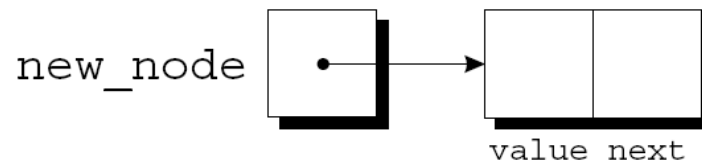
- When we create a node, **we'll need a variable** that can **point to the node temporarily**:

```
struct node *new_node;
```

- We'll use `malloc` to allocate memory for the new node, saving the return value in `new_node`:

```
new_node = malloc(sizeof(struct node));
```

- `new_node` now points to a block of memory just large enough to hold a node structure:

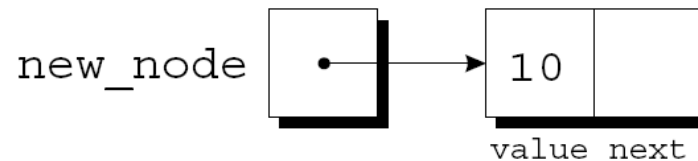


Creating a Node (cont.)

- Next, we'll **store data in the** `value` member of the new node:

```
(*new_node).value = 10;
```

- The resulting picture:



- The **parentheses around** `*new_node` **are mandatory** because the `.` operator would otherwise take precedence over the `*` operator.

The -> Operator

- Accessing a member of a structure using a pointer is so common that C provides a special operator for this purpose.
- This operator, known as **right arrow selection**, is a minus sign followed by >.
- Using the -> operator, we can write

```
new_node->value = 10;
```

instead of

```
(*new_node).value = 10;
```

The -> Operator (cont.)

- The **-> operator produces an lvalue**, so we can use it wherever an ordinary variable would be allowed.
- A `scanf` example:

```
scanf ("%d", &new_node->value) ;
```
- The **& operator is still required**, even though `new_node` is a pointer.

Inserting a Node at the Beginning of a Linked List

- Suppose that `new_node` is pointing to the node to be inserted, and `first` is pointing to the first node in the linked list.
- The first step is to modify the new node's `next` member to point to the node that was previously at the beginning of the list:

```
new_node->next = first;
```
- The second step is to make `first` point to the new node:

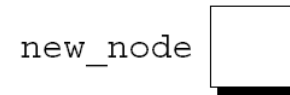
```
first = new_node;
```
- These statements work even if the list is empty.

Inserting a Node at the Beginning of a Linked List (cont.)

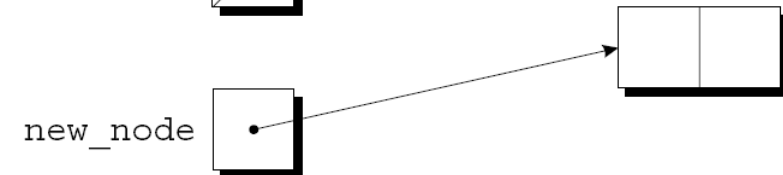
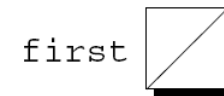
- Let's trace the process of inserting two nodes into an empty list.
- We'll insert a node containing the number 10 first, followed by a node containing 20.

Inserting a Node at the Beginning of a Linked List (cont.)

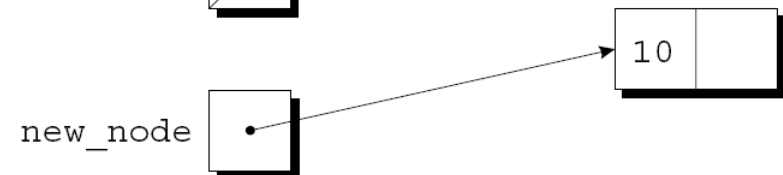
```
first = NULL;
```



```
new_node =  
    malloc(sizeof(struct node));
```

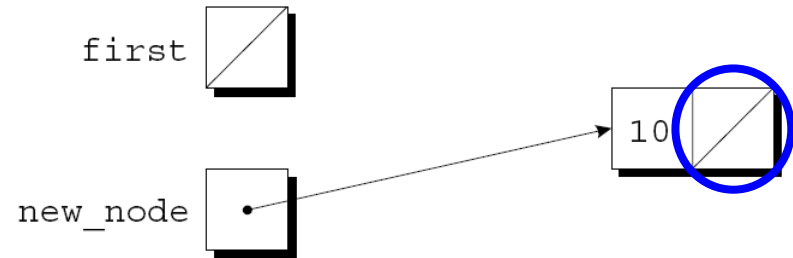


```
new_node->value = 10;
```

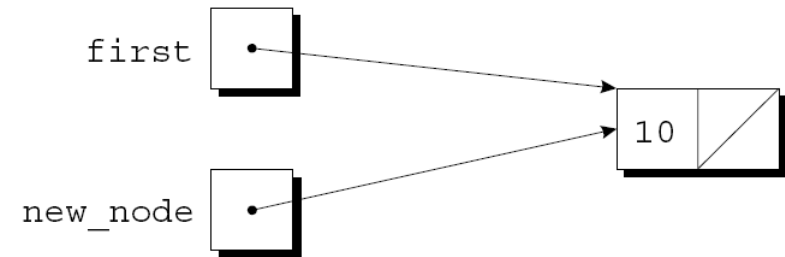


Inserting a Node at the Beginning of a Linked List (cont.)

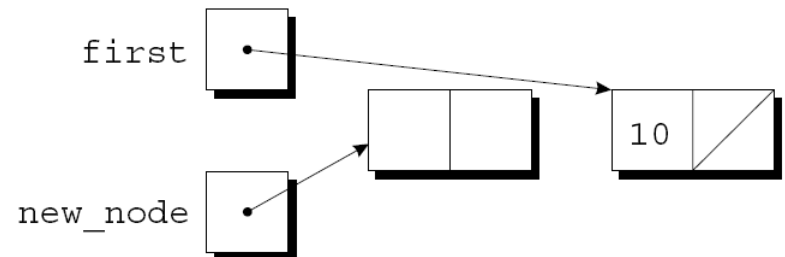
```
new_node->next = first;
```



```
first = new_node;
```

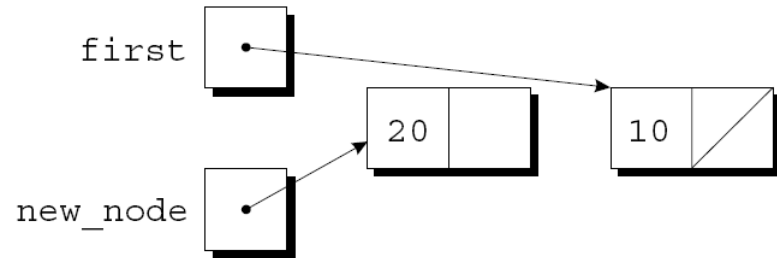


```
new_node =  
    malloc(sizeof(struct node));
```

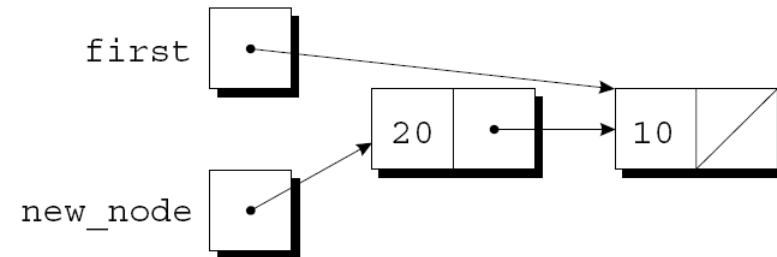


Inserting a Node at the Beginning of a Linked List (cont.)

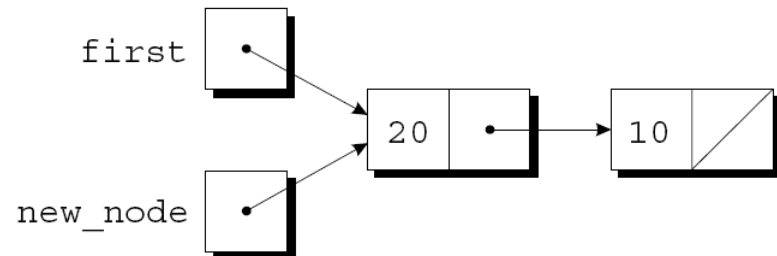
```
new_node->value = 20;
```



```
new_node->next = first;
```



```
first = new_node;
```



Inserting a Node at the Beginning of a Linked List (cont.)

- A function that inserts a node containing n into a linked list, which pointed to by `list`:

```
struct node *add_to_list(struct node *list, int n)
{
    struct node *new_node;

    new_node = malloc(sizeof(struct node));
    if (new_node == NULL) {
        printf("Error: malloc failed in add_to_list\n");
        exit(EXIT_FAILURE);
    }
    new_node->value = n;
    new_node->next = list;
    return new_node;
}
```


Inserting a Node at the Beginning of a Linked List (cont.)

- Note that `add_to_list` returns a pointer to the newly created node (now at the beginning of the list).
- When we call `add_to_list`, we'll need to store its return value into `first`:

```
first = add_to_list(first, 10);  
first = add_to_list(first, 20);
```

Inserting a Node at the Beginning of a Linked List (cont.)

- A function that **uses add to list to create a linked list containing numbers entered by the user:**

```
struct node *read_numbers(void)
{
    struct node *first = NULL;
    int n;

    printf("Enter a series of integers (0 to terminate): ");
    for (;;) {
        scanf("%d", &n);
        if (n == 0)
            return first;
        first = add_to_list(first, n);
    }
}
```

- The **numbers will be in reverse order** within the list.

Searching a Linked List

- A loop that visits the nodes in a linked list, using a pointer variable `p` to keep track of the “current” node:

```
for (p = first; p != NULL; p = p->next)
```

...

- A loop of this form can be used in a function that **searches a list for an integer `n`**.

```
struct node *search_list(struct node *list, int n)
{
    struct node *p;

    for (p = list; p != NULL; p = p->next)
        if (p->value == n)
            return p;
    return NULL;
}
```

Deleting a Node from a Linked List

- A big advantage of storing data in a linked list is that we can easily delete nodes.
- Deleting a node involves three steps:
 1. Locate the node to be deleted.
 2. Alter the previous node so that it “bypasses” the deleted node.
 3. Call `free` to reclaim the space occupied by the deleted node.
- Step 1 is harder than it looks, because step 2 requires changing the *previous* node.

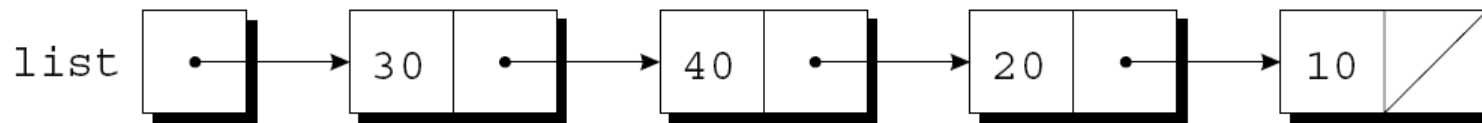
Deleting a Node from a Linked List (cont.)

- We can keep a pointer to the previous node (`prev`) as well as a pointer to the current node (`cur`).
- Assume that `list` points to the list to be searched and `n` is the integer to be deleted.
- A loop that implements step 1:

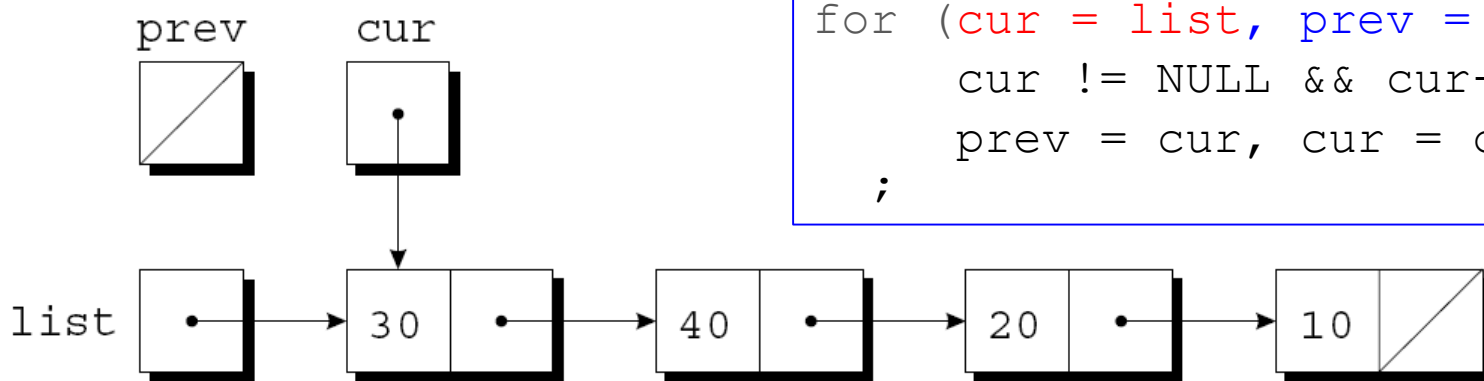
```
for (cur = list, prev = NULL;  
    cur != NULL && cur->value != n;  
    prev = cur, cur = cur->next)  
    ;
```
- When the loop terminates, `cur` points to the node to be deleted and `prev` points to the previous node.

Deleting a Node from a Linked List (cont.)

- Assume that `list` has the following appearance and `n` is 20:



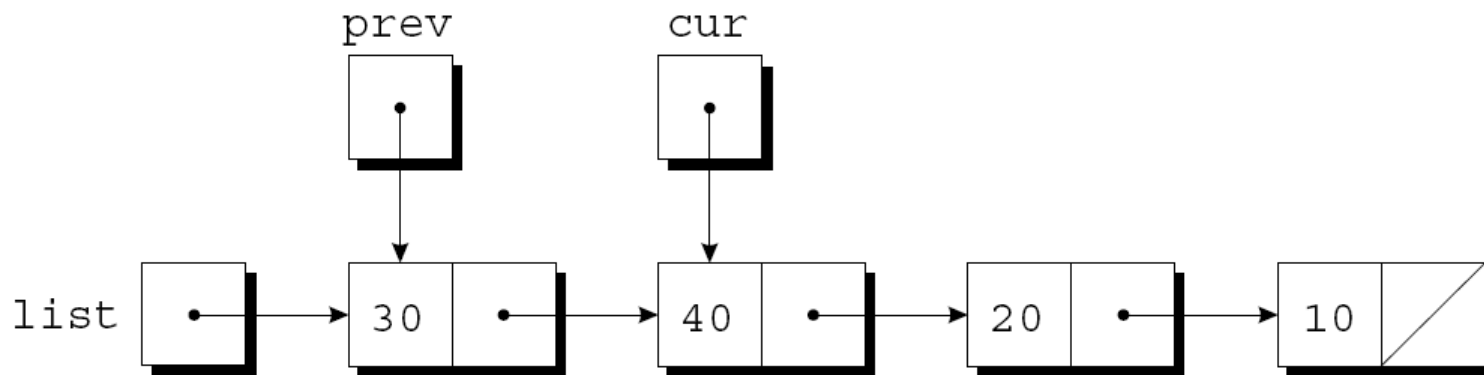
- After `cur = list, prev = NULL` has been executed:



```
for (cur = list, prev = NULL;  
    cur != NULL && cur->value != n;  
    prev = cur, cur = cur->next)  
;
```

Deleting a Node from a Linked List (cont.)

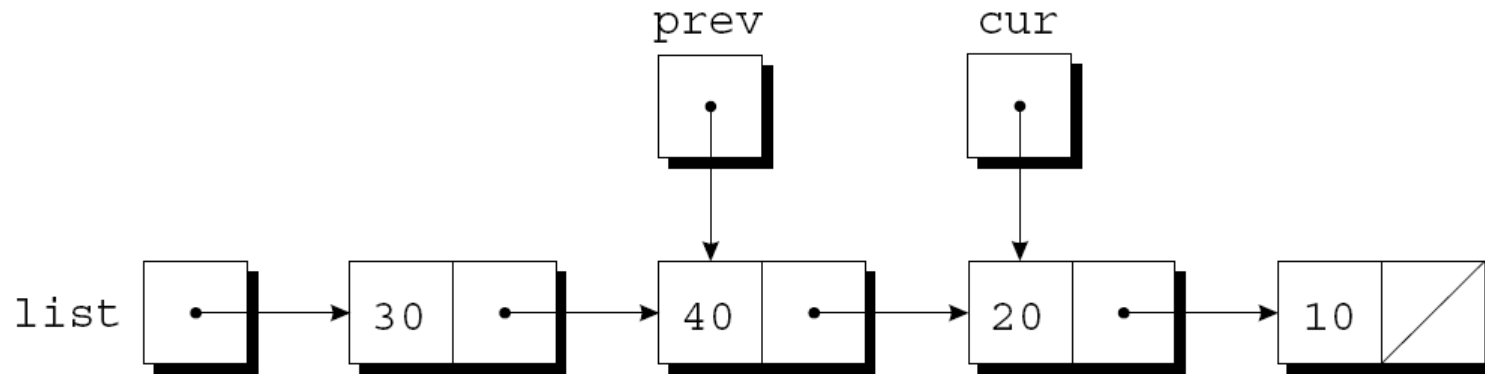
- The test `cur != NULL && cur->value != n` is true, since `cur` is pointing to a node and the node doesn't contain 20.
- After `prev = cur, cur = cur->next` has been executed:



```
for (cur = list, prev = NULL;  
    cur != NULL && cur->value != n;  
    prev = cur, cur = cur->next)  
;
```

Deleting a Node from a Linked List (cont.)

- The test `cur != NULL && cur->value != n` is again true, so `prev = cur, cur = cur->next` is executed once more:



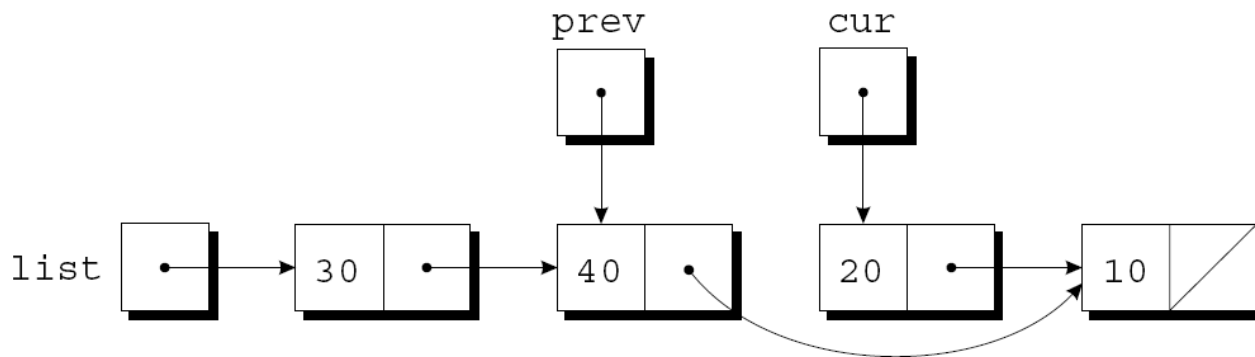
- Since `cur` now points to the node containing 20, the condition `cur->value != n` is false and the loop terminates.

Deleting a Node from a Linked List (cont.)

- At step 2, the following statement

```
prev->next = cur->next;
```

makes the pointer in the previous node point to the node *after* the current node:



- Step 3 is to **release the memory** occupied by the current node:

```
free (cur) ;
```

Deleting a Node from a Linked List (cont.)

- The `delete_from_list` function uses the strategy just outlined.
- When **given a list** and **an integer n** , the function **deletes the first node containing n** .
- If **no node contains n** , `delete_from_list` **does nothing**.
- In either case, the function **returns a pointer to the list**.
- **Deleting the first node** in the list **is a special case** that requires a different bypass step.

Deleting a Node from a Linked List (cont.)

```
struct node *delete_from_list(struct node *list, int n)
{
    struct node *cur, *prev;

    for (cur = list, prev = NULL;
         cur != NULL && cur->value != n;
         prev = cur, cur = cur->next)
        ;
    if (cur == NULL)
        return list;                /* n was not found */
    if (prev == NULL)
        list = list->next;          /* n is in the first node */
    else
        prev->next = cur->next;     /* n is in some other node */
    free(cur);
    return list;
}
```

Ordered Lists

- When the **nodes** of a list are **kept in order**—sorted by the data stored inside the nodes—we say that the list is **ordered**.
- **Inserting a node into an ordered list is more difficult**, because the node won't always be put at the beginning of the list.
- **However, searching is faster**: we can stop looking after reaching the point at which the desired node would have been located.

Program: Maintaining a Parts Database (Revisited)

- The `inventory2.c` program is a modification of the parts database program of Lecture 16, with **the database stored in a linked list** this time.
- **Advantages** of using a linked list:
 - **No need to put a limit on the size** of the database.
 - Database can **easily be kept sorted** by part number.
- In the original program, the database wasn't sorted.

Program: Maintaining a Parts Database (Revisited) (cont.)

- The `part` structure will contain **an additional member** (a pointer to the next node):

```
struct part {  
    int number;  
    char name[NAME_LEN+1];  
    int on_hand;  
    struct part *next;  
};
```

```
struct part {  
    int number;  
    char name[NAME_LEN+1];  
    int on_hand;  
};
```

- `inventory` will **point to the first node** in the list:

```
struct part *inventory = NULL;
```

Program: Maintaining a Parts Database (Revisited) (cont.)

- `find_part` and `insert` will be more complex, however, since we'll keep the nodes in the inventory list sorted by part number.
- In the original program, `find_part` returns an index into the `inventory` array.
- In the new program, `find_part` will return a pointer to the node that contains the desired part number.
- If it doesn't find the part number, `find_part` will return a null pointer.

Program: Maintaining a Parts Database (Revisited) (cont.)

- Since the list of parts is sorted, `find_part` can stop when it finds a node containing a part number that's greater than or equal to the desired part number.
- `find_part`'s search loop:

```
for (p = inventory;  
    p != NULL && number > p->number;  
    p = p->next)  
    ;
```

- When the loop terminates, we'll need to test whether the part was found:

```
if (p != NULL && number == p->number)  
    return p;
```


Program: Maintaining a Parts Database (Revisited) (cont.)

- The **original** version of `insert` stores a new part in the **next available array element**.
- The **new version must determine where the new part belongs** in the list and insert it there.
- It will **also check** whether the part number is **already present** in the list.
- A loop that accomplishes both tasks:

```
for (cur = inventory, prev = NULL;  
    cur != NULL && new_node->number > cur->number;  
    prev = cur, cur = cur->next)  
;
```

Program: Maintaining a Parts Database (Revisited) (cont.)

- Once the loop terminates, `insert` will check whether `cur` isn't NULL and whether `new_node->number` equals `cur->number`.
- If both are true, the part number is already in the list.
- Otherwise, `insert` will insert a new node between the nodes pointed to by `prev` and `cur`.
- This strategy works even if the new part number is larger than any in the list.

Program: Maintaining a Parts Database (Revisited) (cont.)

inventory2.c

```
#include <stdio.h>
#include <stdlib.h>
#include "readline.h"
#define NAME_LEN 25

struct part {
    int number;
    char name[NAME_LEN+1];
    int on_hand;
    struct part *next;
};
```

```
struct part {
    int number;
    char name[NAME_LEN+1];
    int on_hand;
} inventory[MAX_PARTS];
```

```
int num_parts = 0;    /* number of parts */
```

```
struct part *inventory = NULL;    /* points to first part */
```

```
struct part *find_part(int number);
void insert(void);
void search(void);
void update(void);
void print(void);
```

```
int find_part(int number);
```

Program: Maintaining a Parts Database (Revisited) (cont.)

```
int main(void)
{
    char code;

    for (;;) {
        printf("Enter operation code: ");
        scanf(" %c", &code);
        while (getchar() != '\n')    /* skips to end of line */
            ;
        switch (code) {
            case 'i': insert();
                       break;
            case 's': search();
                       break;
            case 'u': update();
                       break;
            case 'p': print();
                       break;
            case 'q': return 0;
            default: printf("Illegal code\n");
        }
        printf("\n");
    }
}
```

Program: Maintaining a Parts Database (Revisited) (cont.)

```
struct part *find_part(int number)
{
    struct part *p;

    for (p = inventory;
         p != NULL && number > p->number;
         p = p->next)
        ;
    if (p != NULL && number == p->number)
        return p;
    return NULL;
}
```

```
int find_part(int number)
{
    int i;

    for (i = 0; i < num_parts; i++)
        if (inventory[i].number == number)
            return i;
    return -1;
}
```

Program: Maintaining a Parts Database (Revisited) (cont.)

```
void insert(void)
{
    struct part *cur, *prev, *new_node;

    new_node = malloc(sizeof(struct part));
    if (new_node == NULL) {
        printf("Database is full; can't add more parts.\n");
        return;
    }

    printf("Enter part number: ");
    scanf("%d", &new_node->number);

    if (num_parts == MAX_PARTS) {
        printf("Database is full; can't add more parts.\n");
        return;
    }
```

Program: Maintaining a Parts Database (Revisited) (cont.)

```
for (cur = inventory, prev = NULL;  
    cur != NULL && new_node->number > cur->number;  
    prev = cur, cur = cur->next)  
;
```

```
if (cur != NULL && new_node->number == cur->number) {  
    printf("Part already exists.\n");  
    free(new_node);  
    return;  
}
```

```
printf("Enter part name: ");  
read_line(new_node->name, NAME_LEN);  
printf("Enter quantity on hand: ");  
scanf("%d", &new_node->on_hand);
```

```
new_node->next = cur;  
if (prev == NULL)  
    inventory = new_node;  
else  
    prev->next = new_node;
```

```
if (find_part(part_number) >= 0) {  
    printf("Part already exists.\n");  
    return;  
}
```

```
// insert at the end  
inventory[num_parts].number = part_number;  
num_parts++;
```



Program: Maintaining a Parts Database (Revisited) (cont.)

```
void search(void)
{
    int number;
    struct part *p;

    printf("Enter part number: ");
    scanf("%d", &number);
    p = find_part(number);
    if (p != NULL) {
        printf("Part name: %s\n", p->name);
        printf("Quantity on hand: %d\n", p->on_hand);
    } else
        printf("Part not found.\n");
}
```

```
i = find_part(number);
if (i >= 0) {
```


Program: Maintaining a Parts Database (Revisited) (cont.)

```
void update(void)
{
    int number, change;
    struct part *p;

    printf("Enter part number: ");
    scanf("%d", &number);
    p = find_part(number);
    if (p != NULL) {
        printf("Enter change in quantity on hand: ");
        scanf("%d", &change);
        p->on_hand += change;
    } else
        printf("Part not found.\n");
}
```

```
i = find_part(number);
if (i >= 0) {
```

```
    inventory[i].on_hand += change;
```

Program: Maintaining a Parts Database (Revisited) (cont.)

```
void print(void)
{
    struct part *p;
    printf("Part Number    Part Name                "
           "Quantity on Hand\n");
    for (p = inventory; p != NULL; p = p->next)
        printf("%7d          %-25s%11d\n", p->number, p->name,
               p->on_hand);
}
```

```
int i;
for (i = 0; i < num_parts; i++)
```

A Quick Review to This Lecture

- The `<stdlib.h>` header declares three memory allocation functions which return **void pointers** (`void *`) to allocated space or **return null pointers** if allocation is failed:

- `malloc` **allocates** a block of **size bytes** but **doesn't initialize it**.

```
void *malloc(size_t size);
```

```
p = malloc(10000);  
if (p == NULL){ /* failed*/ }
```

- `calloc` allocates an array with `nmemb` **elements**, **each of which** is **size bytes long**, and **clears it**.

```
a = calloc(n, sizeof(int));
```

```
void *calloc(size_t nmemb, size_t size);
```

- `realloc` **resizes** a block (where `ptr` points to) to a new size `size`

```
void *realloc(void *ptr, size_t size);
```

A Quick Review to This Lecture (cont.)

- Each C program is responsible for recycling its own garbage by calling the `free` function to release unneeded memory.

```
void free(void *ptr);
```

```
p = malloc(...);  
free(p);
```

- When the block is freed, all the pointers are left dangling.