

# Introduction to Logic

## Midterm Examination, Semester 1/2021

2 Oct 2021, 13.30-15.00

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### Problem 1 (5 pts)

The passage below contains an argument.

<sup>1</sup>The government must save Thai airways from bankruptcy. <sup>2</sup>It is the pride of our nation.  
More importantly, <sup>3</sup>its failing would make our tourism industry collapse.

Identify the premises, the conclusion, and the hidden premise(s) (if any).

**Example.** <sup>1</sup>Boxing causes injury, so <sup>2</sup>it is not a sport we should encourage.  
Statement 2 is the conclusion. Statement 1 is a premise. The hidden premise is

We should not encourage a sport that causes injury.

### Problem 2 (10 pts)

Each passage below contains an argument. Draw a diagram showing the inferential relationship among the statements in the passage. If a statement is redundant or plays no role in the argument, do not include them in the diagram.

**2.1** <sup>1</sup>Proteins are discovered not invented. <sup>2</sup>Inventions are patentable but discoveries are not.  
Hence, <sup>3</sup>the patenting of proteins is simply flawed.

**2.2** <sup>1</sup>The Big Bang theory is being regarded as wrong. <sup>2</sup>According to this theory, the universe began with the Big Bang, a huge explosion occurring 20 billion years ago. The problem is <sup>3</sup>astronomers have found a huge cluster of galaxies that is too big to have been formed in 20 billion years. Based on recent data, it is now known that <sup>4</sup>galaxies form vast ribbons stretching billions of light years and <sup>5</sup>are separated by empty spaces spanning hundreds of millions of light years. Because <sup>6</sup>galaxies travel much slower than the speed of light, these facts imply that <sup>7</sup>such a large cluster of galaxies must have taken at least 100 billion years to form, five times as long as the time since the Big Bang presumably occurred.

### Problem 3 (9 pts)

Each passage below contains a compound statement. Write each statement below as a formula in propositional logic using the given propositional letters and their specified meaning.

**Example.** “If you have not paid your tuition fee, you will not be allowed to graduate.”

$p$  = You have paid your tuition fee.

$g$  = You are allowed to graduate.

**Ans.**  $\neg p \rightarrow \neg g$

**3.1** “Our constitution neither acknowledges nor tolerates racisms.”

$a$  = Our constitution acknowledges racisms.

$t$  = Our constitution tolerates racisms.

**3.2** “The defendant will receive probation provided that he/she cooperates with the attorney.”

$p$  = The defendant will receive probation.

$c$  = The defendant cooperates with the attorney.

**3.3** “All of these are equivalent: (a)  $S$  is the empty set; (b)  $\overline{S}$  is the universal set; and (c)  $S$  is a subset of every set.”

$a$  =  $S$  is the empty set.

$b$  =  $\overline{S}$  is the universal set.

$c$  =  $S$  is a subset of every set.

### Problem 4 (5 pts)

Rewrite the following code fragment into an equivalent one without the **else** statement.

```
if(x > 1) {  
    if(y > 1)  
        printf("a");  
    else  
        printf("b");  
} else {  
    printf("c");  
    if(y > 1)  
        printf("d");  
}
```

### Problem 5 (10 pts)

For each formula below, check whether it is satisfiable or not. If the formula is satisfiable, give a truth assignment which makes the formula true. If not, show that it is unsatisfiable.

$$5.1 \quad (p \wedge q \wedge \neg p \wedge r) \vee (\neg p \wedge s \wedge \neg q \wedge \neg s) \vee (r \wedge \neg p \wedge \neg q \wedge p) \vee \neg q$$

$$5.2 \quad (p \vee \neg q \vee r) \wedge (p \vee q) \wedge (r \vee \neg q \vee \neg s) \wedge (\neg p \vee s) \wedge (\neg r \vee \neg q) \wedge (\neg s \vee q)$$

## Problem 6 (10 pts)

For each pair of formulas below, either show that the two formulas are logically equivalent or describe a truth assignment which makes one formula true and the other formula false.

$$6.1 \quad p \leftrightarrow (q \leftrightarrow r) \text{ and } (p \leftrightarrow q) \leftrightarrow r$$

$$6.2 \quad \neg p \vee (q \vee (\neg r \vee s)) \text{ and } (p \wedge r) \rightarrow (q \vee s)$$

## Problem 7 (10 pts)

Draw a reduced OBDD for the formula  $(p \rightarrow q) \rightarrow (p \rightarrow r)$ .

## Problem 8 (20 pts)

Each passage below contains an argument. For each passage, please do the following:

- Write the underlined statements in the passage in propositional logic using the given propositional letters and its specified meaning.
- From the formulas you obtained in (a), determine which formulas are the premises and which formula is the conclusion of the argument in the passage.
- Based on what you identified as the premises and the conclusion in (b), determine whether the argument is valid or not. If so, provide a derivation of the conclusion from the premises using natural deduction rules. If not, give a truth assignment which makes all the premises true but the conclusion false.

**Example.** <sup>1</sup>John must not be at home at the moment. <sup>2</sup>If he were at home, his car must be in the garage. But from what I can see, <sup>3</sup>his car is currently not in the garage.

$h$  = John is at home at the moment.

$g$  = John's car is currently in the garage.

**Ans.**

- Statement 1 =  $\neg h$   
Statement 2 =  $h \rightarrow g$   
Statement 3 =  $\neg g$

- Premises:  $h \rightarrow g, \neg g$   
Conclusion:  $\neg h$

- The argument is valid.

1 : $h \rightarrow g$	premise
2 : $\neg g$	premise
3 : $\neg h$	MT, 1, 2

**8.1**    “<sup>1</sup>The victim was right-handed. <sup>2</sup>If the victim committed suicide and was right-handed, she would not have wounds on the left of her head. <sup>4</sup>Hence, if there are wounds on the left of the victim’s head, she did not commit suicide.

$r$  = The victim was right-handed.

$w$  = There are wounds on the left of the victim’s head.

$s$  = The victim committed suicide.

**8.2**    “<sup>1</sup>You should not stay up all night to study for the exam. <sup>2</sup>If you stay up all night to study for the exam, you will be tired in the morning. And <sup>3</sup>if you are tired in the morning and the exam is difficult, you will not be able to do well on the exam. Obviously, <sup>4</sup>if you stay up all night to study for the exam and still not be able to do well on the exam, then you should not do that.

$s$  = You should stay up all night to study for the exam.

$u$  = You stay up all night to study for the exam.

$t$  = You are tired in the morning.

$d$  = The exam is difficult.

$w$  = You are able to do well on the exam.

## Problem 9 (20 pts)

Imagine a fictional island where two types of inhabitants, called the *knight*s and the *knave*s, are living. A knight always tells the truth, whereas a knave always tells lies (i.e. the opposite of the truth). Each inhabitant is of one of these two types, but unfortunately it is not clear which type he/she is. When you visited this island, you met 5 inhabitants on the island, namely  $A$ ,  $B$ ,  $C$ ,  $D$ , and  $E$ . Below is the transcript from your conversation with some of these inhabitants.

- $A$  said “Both  $C$  and  $D$  are knights.”
- $B$  said “If  $E$  is a knight, then so is  $A$ .”
- $C$  said “Either  $B$  or  $E$  or both are knaves.”
- $D$  said “ $E$  is a knave if and only if  $C$  is.”

You are then asked to determine whether each of the 5 inhabitants is a knight or a knave. Luckily, you are in possession of a highly-efficient SAT solver program, which can determine whether a formula in CNF is satisfiable or not. Explain in detail how you can utilize your SAT solver to solve this.

**Hint:** Introduce the following propositional symbols  $a$ ,  $b$ ,  $c$ ,  $d$ , and  $e$  which mean that  $A$ ,  $B$ ,  $C$ ,  $D$ , and  $E$ , respectively, are *knight*s.

## Problem 10 (10 pts)

Suppose  $A$ ,  $B$ ,  $C$ ,  $D$ , and  $E$  are the sets given by:

$$A = \{0, 1, 2\}$$

$$B = \{-5, 1, 3, 6, 10\}$$

$$C = \{x \in \mathbb{Z} \mid 0 < x \leq 20 \text{ and } x \text{ is even}\}$$

$$D = \{x \in \mathbb{Z} \mid x = y - z \text{ for some } y \text{ and } z \text{ in } A\}$$

$$E = \{2x + 1 \in \mathbb{Z} \mid x \in A\}$$

List all the members of each of the following sets.

10.1  $B \cup C$

10.2  $\wp(A)$

10.3  $D$

10.4  $A \times E$

10.5  $\wp(\wp(A \cap B))$

## Problem 11 (10 pts)

Suppose  $A = \{x \in \mathbb{Z} \mid -25 \leq x \leq 25\}$ . Let  $P$  be the following binary relation:

$$P = \{(x, y) \in A \times A \mid y = x^2\}$$

11.1 List all the members of  $P$ .

11.2 List all the members of  $P \circ P$ .

## Problem 12 (10 pts)

A binary relation  $R$  on a non-empty set  $A$  is said to be transitive if and only if

$$xRy \text{ and } yRz \text{ implies } xRz, \text{ for all } x, y, z \in A$$

The *transitive closure* of a binary relation  $R$  on  $A$  is the *smallest* transitive relation on  $A$  that includes  $R$ .

Find the transitive closure of the following relation on  $\mathbb{N}$ :

$$R = \{(1, 3), (2, 1), (3, 4), (4, 2), (4, 5)\}.$$

————— This is the end of the exam paper. —————

Table 1: Some Logical Equivalences		
	Equivalences	Name
E1	$\phi \wedge \top \equiv \phi$	Identity Laws
E2	$\phi \vee \perp \equiv \phi$	
E3	$\phi \wedge \perp \equiv \perp$	Domination Laws
E4	$\phi \vee \top \equiv \top$	
E5	$\phi \wedge \neg\phi \equiv \perp$	Complement Laws
E6	$\phi \vee \neg\phi \equiv \top$	
E7	$\phi \wedge \phi \equiv \phi$	Idempotent Laws
E8	$\phi \vee \phi \equiv \phi$	
E9	$\neg(\neg\phi) \equiv \phi$	Double Negation Law
E10	$\phi \wedge \psi \equiv \psi \wedge \phi$	Commutative Laws
E11	$\phi \vee \psi \equiv \psi \vee \phi$	
E12	$\phi \wedge (\psi \wedge \chi) \equiv (\phi \wedge \psi) \wedge \chi$	Associative Laws
E13	$\phi \vee (\psi \vee \chi) \equiv (\phi \vee \psi) \vee \chi$	
E14	$\phi \wedge (\psi \vee \chi) \equiv (\phi \wedge \psi) \vee (\phi \wedge \chi)$	Distributive Laws
E15	$\phi \vee (\psi \wedge \chi) \equiv (\phi \vee \psi) \wedge (\phi \vee \chi)$	
E16	$\neg(\phi \wedge \psi) \equiv \neg\phi \vee \neg\psi$	De Morgan's Laws
E17	$\neg(\phi \vee \psi) \equiv \neg\phi \wedge \neg\psi$	
E18	$\phi \wedge (\phi \vee \psi) \equiv \phi$	Absorption Laws
E19	$\phi \vee (\phi \wedge \psi) \equiv \phi$	
E20	$\phi \rightarrow \psi \equiv \neg\phi \vee \psi$	
E21	$\phi \leftrightarrow \psi \equiv (\phi \rightarrow \psi) \wedge (\psi \rightarrow \phi)$	

Table 1 lists some well-known logical equivalences in propositional logic.