

Project 1. MIPS Assembler

Due 23:59, March. 27th

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1. Introduction

This project aims to develop a MIPS Instruction Set Architecture (ISA) assembler. The assembler will translate assembly language codes into a binary file. This project is intended to help you understand the MIPS better.

The scope of this project includes the development of a simplified assembler that omits the linking process, and consequently, **symbol and relocation tables are not required**.

The assembler must support a subset of the MIPS instruction set, particularly handling labels for jump/branch targets and labels for the static data section.

Instruction Set

Refer to the green card page from the textbook, also attached at the last pages of the current document. The assembler should implement the following instructions, focusing on signed operations and handling sign extension for immediate and offset fields where applicable:

ADDI	ADD	AND	ANDI	BEQ	BNE	J
JAL	JR	LUI	LW	<u>LA</u> *	NOR	OR
ORI	SLTI	SLT	SLL	SRL	SW	SUB

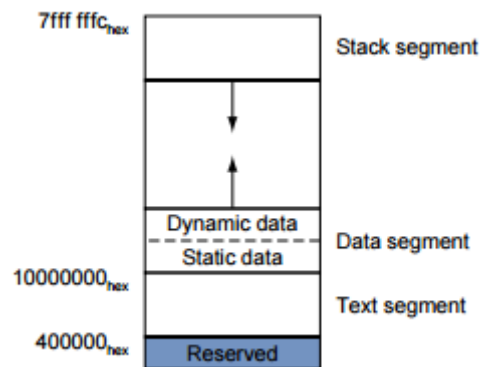
- Implement only signed operation instructions
- Implement sign extension for immediate fields and offset fields in specific instructions (addiu, sltiu, beq, bne, lw, sw)
- Implement only 4B words loads and stores
- The assembler must support decimal and hexadecimal (0x prefix) for immediate fields and the .data section
- Use the format "\$n" for register names, where n ranges from 0 to 31.
- The la (load address) instruction is a pseudo instruction and thus should be translated into lui or ori instructions

la \$2, VAR1 // VAR1 is a label in the data section // It should be converted to lui and ori instructions. lui \$reg, upper 16bit address ori \$reg, \$reg, lower 16bit address // Skipped if the lower 16bit address is 0x0000
Case 1) load address is 0x1000 0000 lui \$2, 0x1000
Case 2) load address is 0x1000 0004 lui \$2, 0x1000 ori \$2, \$2, 0x0004

Directives

- .text: Indicates instructions stored in the user text segment, starting from 0x400000.
- .data: Indicates data stored in the data segment, starting from 0x10000000.
- .word: Stores n 32-bit quantities in consecutive memory words.
- You may assume that the .data and .text directives appear only once, and the .data appear before .text directive.

Memory Layout



Execution command

> ./runfile <assembly file>

The program should produce a single output file (*.o) from the input assembly file (*.s).

Input format

```
1      .data
2  array: .word 3
3         .word 123
4         .word 4346
5  array2: .word 0x11111111
6      .text
7  main:
8      addi    $2, $0, 1024
9      add     $3, $2, $2
10     or      $4, $3, $2
11     sll     $6, $5, 16
12     addi    $7, $6, 9999
13     sub     $8, $7, $2
14     nor     $9, $4, $3
15     ori     $10, $2, 255
16     srl     $11, $6, 5
17     la      $4, array2
18     and     $13, $11, $5
19     andi    $14, $4, 100
20     lui     $17, 100
21     addi    $2, $0, 0xa
```

Output Format

The output format is an ASCII string of '0's and '1's representing the binary code. The ASCII string follows a simplified custom format.

- The first two words (32bits) are the size of text section and data section.
- The next bytes are the instructions in binary. The length must be equal to the specified text section length.
- After the text section, the rest of bytes are the initial values of the data section.

```
<text section size> // There should be no newlines in the actual output
<data section size>
<instruction 1>
...
<instruction n>
<value 1>
...
<value m>
```

- Please refer to the example output file inside the project files for a better idea of this format.

Program Language

You should use C or C++ for this project, considering the continuity with subsequent projects.

GitLab Repository Setup

Create a new repository by forking the project prepared by the TAs at

cs311.kaist.ac.kr/handout/project1.

The guidelines for forking, cloning and pushing using Git have been uploaded to the notice tab on KLMS.

Grading Policy

The grading is based on five open cases and two hidden cases. The two hidden cases are included to prevent hardcoding. The five open cases can be found in the forked repository. Each open case is worth 10 points, and each hidden case is worth 15 points, making a total of 80 points. For each case, you can get points only if every digit of the output binary matches the expected answer.

Submission

1. Commit and Push: Ensure that your code and Makefile are committed and pushed to your forked GitLab repository.
2. Tagging for Submission: Follow these commands in your working directory to tag your submission:

```
git tag -a submit -m 'your_custom_message'
git push origin submit
```

Important notes

- The absence of Makefile in submission will result in your work not being graded.
- The absence of a "submit" tag will result in your work not being graded.
- Always double-check to confirm the submission tag has been properly applied.

Late Policy and Academic Integrity

- A penalty of 30% will be applied to submissions made within the first day after the deadline (March 28th 23:59). Submissions beyond this period will not be accepted.
- Engaging with peers and referencing external materials is encouraged for educational purposes. However, direct copying from other students or publicly available code is strictly prohibited and will be actively monitored. Detected plagiarism will result in severe penalties.
- The TAs will employ code comparison tools against open-source repositories and submissions from the previous semesters to ensure the originality of your work.

For any further questions or if you encounter any administrative issues (e.g., late submission due to extenuating circumstances, technical difficulties with GitLab), please reach out on Piazza.

MIPS Reference Data

①



CORE INSTRUCTION SET

NAME, MNEMONIC	FOR-MAT	OPERATION (in Verilog)	OPCODE / FUNCT (Hex)
Add	add R	$R[rd] = R[rs] + R[rt]$	(1) 0 / 20 _{hex}
Add Immediate	addi I	$R[rt] = R[rs] + \text{SignExtImm}$	(1,2) 8 _{hex}
Add Imm. Unsigned	addiu I	$R[rt] = R[rs] + \text{SignExtImm}$	(2) 9 _{hex}
Add Unsigned	addu R	$R[rd] = R[rs] + R[rt]$	0 / 21 _{hex}
And	and R	$R[rd] = R[rs] \& R[rt]$	0 / 24 _{hex}
And Immediate	andi I	$R[rt] = R[rs] \& \text{ZeroExtImm}$	(3) c _{hex}
Branch On Equal	beq I	if($R[rs] == R[rt]$) $PC = PC + 4 + \text{BranchAddr}$	(4) 4 _{hex}
Branch On Not Equal	bne I	if($R[rs] != R[rt]$) $PC = PC + 4 + \text{BranchAddr}$	(4) 5 _{hex}
Jump	j J	$PC = \text{JumpAddr}$	(5) 2 _{hex}
Jump And Link	jal J	$R[31] = PC + 8; PC = \text{JumpAddr}$	(5) 3 _{hex}
Jump Register	jr R	$PC = R[rs]$	0 / 08 _{hex}
Load Byte Unsigned	lbu I	$R[rt] = \{24'b0, M[R[rs] + \text{SignExtImm}(7:0)]\}$	(2) 24 _{hex}
Load Halfword Unsigned	lhu I	$R[rt] = \{16'b0, M[R[rs] + \text{SignExtImm}(15:0)]\}$	(2) 25 _{hex}
Load Linked	ll I	$R[rt] = M[R[rs] + \text{SignExtImm}]$	(2,7) 30 _{hex}
Load Upper Imm.	lui I	$R[rt] = \{\text{imm}, 16'b0\}$	f _{hex}
Load Word	lw I	$R[rt] = M[R[rs] + \text{SignExtImm}]$	(2) 23 _{hex}
Nor	nor R	$R[rd] = \sim (R[rs] R[rt])$	0 / 27 _{hex}
Or	or R	$R[rd] = R[rs] R[rt]$	0 / 25 _{hex}
Or Immediate	ori I	$R[rt] = R[rs] \text{ZeroExtImm}$	(3) d _{hex}
Set Less Than	slt R	$R[rd] = (R[rs] < R[rt]) ? 1 : 0$	0 / 2a _{hex}
Set Less Than Imm.	slti I	$R[rt] = (R[rs] < \text{SignExtImm}) ? 1 : 0$	(2) a _{hex}
Set Less Than Imm. Unsigned	sltiu I	$R[rt] = (R[rs] < \text{SignExtImm}) ? 1 : 0$	(2,6) b _{hex}
Set Less Than Unsig.	sltu R	$R[rd] = (R[rs] < R[rt]) ? 1 : 0$	(6) 0 / 2b _{hex}
Shift Left Logical	sll R	$R[rd] = R[rt] \ll \text{shamt}$	0 / 00 _{hex}
Shift Right Logical	srl R	$R[rd] = R[rt] \gg \text{shamt}$	0 / 02 _{hex}
Store Byte	sb I	$M[R[rs] + \text{SignExtImm}(7:0)] = R[rt](7:0)$	(2) 28 _{hex}
Store Conditional	sc I	$M[R[rs] + \text{SignExtImm}] = R[rt];$ $R[rt] = (\text{atomic}) ? 1 : 0$	(2,7) 38 _{hex}
Store Halfword	sh I	$M[R[rs] + \text{SignExtImm}(15:0)] = R[rt](15:0)$	(2) 29 _{hex}
Store Word	sw I	$M[R[rs] + \text{SignExtImm}] = R[rt]$	(2) 2b _{hex}
Subtract	sub R	$R[rd] = R[rs] - R[rt]$	(1) 0 / 22 _{hex}
Subtract Unsigned	subu R	$R[rd] = R[rs] - R[rt]$	0 / 23 _{hex}

- (1) May cause overflow exception
 (2) $\text{SignExtImm} = \{16\{\text{immediate}[15]\}, \text{immediate}\}$
 (3) $\text{ZeroExtImm} = \{16\{1b'0\}, \text{immediate}\}$
 (4) $\text{BranchAddr} = \{14\{\text{immediate}[15]\}, \text{immediate}, 2'b0\}$
 (5) $\text{JumpAddr} = \{PC + 4[31:28], \text{address}, 2'b0\}$
 (6) Operands considered unsigned numbers (vs. 2's comp.)
 (7) Atomic test&set pair; $R[rt] = 1$ if pair atomic, 0 if not atomic

BASIC INSTRUCTION FORMATS

R	opcode	rs	rt	rd	shamt	funct
	31	26 25	21 20	16 15	11 10	6 5
	0					
I	opcode	rs	rt	immediate		
	31	26 25	21 20	16 15		
	0					
J	opcode	address				
	31	26 25				n

ARITHMETIC CORE INSTRUCTION SET

②

NAME, MNEMONIC	FOR-MAT	OPERATION	OPCODE / FUNCT (Hex)
Branch On FP True	belt FI	if($FPcond$) $PC = PC + 4 + \text{BranchAddr}$	(4) 11/8/1/--
Branch On FP False	belf FI	if(! $FPcond$) $PC = PC + 4 + \text{BranchAddr}$	(4) 11/8/0/--
Divide	div R	$Lo = R[rs]/R[rt]; Hi = R[rs]\%R[rt]$	0/--/--1a
Divide Unsigned	divu R	$Lo = R[rs]/R[rt]; Hi = R[rs]\%R[rt]$	(6) 0/--/--1b
FP Add Single	add.s FR	$F[fd] = F[fs] + F[ft]$	11/10/--/0
FP Add Double	add.d FR	$\{F[fd], F[fd+1]\} = \{F[fs], F[fs+1]\} + \{F[ft], F[ft+1]\}$	11/11/--/0
FP Compare Single	c.x.s* FR	$FPcond = (F[fs] \text{ op } F[ft]) ? 1 : 0$	11/10/--/y
FP Compare Double	c.x.d* FR	$FPcond = (\{F[fs], F[fs+1]\} \text{ op } \{F[ft], F[ft+1]\}) ? 1 : 0$	11/11/--/y
FP Divide Single	div.s FR	$F[fd] = F[fs] / F[ft]$	11/10/--/3
FP Divide Double	div.d FR	$\{F[fd], F[fd+1]\} = \{F[fs], F[fs+1]\} / \{F[ft], F[ft+1]\}$	11/11/--/3
FP Multiply Single	mul.s FR	$F[fd] = F[fs] * F[ft]$	11/10/--/2
FP Multiply Double	mul.d FR	$\{F[fd], F[fd+1]\} = \{F[fs], F[fs+1]\} * \{F[ft], F[ft+1]\}$	11/11/--/2
FP Subtract Single	sub.s FR	$F[fd] = F[fs] - F[ft]$	11/10/--/1
FP Subtract Double	sub.d FR	$\{F[fd], F[fd+1]\} = \{F[fs], F[fs+1]\} - \{F[ft], F[ft+1]\}$	11/11/--/1
Load FP Single	lwc1 I	$F[rt] = M[R[rs] + \text{SignExtImm}]$	(2) 31/--/--
Load FP Double	ldc1 I	$F[rt] = M[R[rs] + \text{SignExtImm}];$ $F[rt+1] = M[R[rs] + \text{SignExtImm} + 4]$	(2) 35/--/--
Move From Hi	mfhi R	$R[rd] = Hi$	0/--/--10
Move From Lo	mflo R	$R[rd] = Lo$	0/--/--12
Move From Control	mfc0 R	$R[rd] = CR[rs]$	10/0/--/0
Multiply	mult R	$\{Hi, Lo\} = R[rs] * R[rt]$	0/--/--18
Multiply Unsigned	multu R	$\{Hi, Lo\} = R[rs] * R[rt]$	(6) 0/--/--19
Shift Right Arith.	sra R	$R[rd] = R[rt] \gg \text{shamt}$	0/--/--3
Store FP Single	swc1 I	$M[R[rs] + \text{SignExtImm}] = F[rt]$	(2) 39/--/--
Store FP Double	sdc1 I	$M[R[rs] + \text{SignExtImm}] = F[rt];$ $M[R[rs] + \text{SignExtImm} + 4] = F[rt+1]$	(2) 3d/--/--

FLOATING-POINT INSTRUCTION FORMATS

FR	opcode	fmt	ft	fs	fd	funct
	31	26 25	21 20	16 15	11 10	6 5
	0					
FI	opcode	fmt	ft	immediate		
	31	26 25	21 20	16 15		
	0					

PSEUDOINSTRUCTION SET

NAME	MNEMONIC	OPERATION
Branch Less Than	blt	if($R[rs] < R[rt]$) $PC = \text{Label}$
Branch Greater Than	bgt	if($R[rs] > R[rt]$) $PC = \text{Label}$
Branch Less Than or Equal	btle	if($R[rs] \leq R[rt]$) $PC = \text{Label}$
Branch Greater Than or Equal	bge	if($R[rs] \geq R[rt]$) $PC = \text{Label}$
Load Immediate	li	$R[rd] = \text{immediate}$
Move	move	$R[rd] = R[rs]$

REGISTER NAME, NUMBER, USE, CALL CONVENTION

NAME	NUMBER	USE	PRESERVED ACROSS A CALL?
\$zero	0	The Constant Value 0	N.A.
\$at	1	Assembler Temporary	No
\$v0-\$v1	2-3	Values for Function Results and Expression Evaluation	No
\$a0-\$a3	4-7	Arguments	No
\$t0-\$t7	8-15	Temporaries	No
\$s0-\$s7	16-23	Saved Temporaries	Yes
\$t8-\$t9	24-25	Temporaries	No
\$k0-\$k1	26-27	Reserved for OS Kernel	No
\$gp	28	Global Pointer	Yes
\$sp	29	Stack Pointer	Yes
\$fp	30	Frame Pointer	Yes
\$ra	31	Return Address	Yes

OPCODES, BASE CONVERSION, ASCII SYMBOLS

MIPS opcode (31:26)	(1) MIPS funct (5:0)	(2) MIPS funct (5:0)	Binary	Deci- mal	Hexa- decim- al	ASCII Char- acter	Deci- mal	Hexa- decim- al	ASCII Char- acter
(1)	sll	add.f	00 0000	0	0	NUL	64	40	@
j	srl	sub.f	00 0001	1	1	SOH	65	41	A
j	sra	mul.f	00 0010	2	2	STX	66	42	B
j	sra	div.f	00 0011	3	3	ETX	67	43	C
beq	sllv	sqrt.f	00 0100	4	4	EOT	68	44	D
bne		abs.f	00 0101	5	5	ENQ	69	45	E
blez	srlv	mov.f	00 0110	6	6	ACK	70	46	F
bgtz	sra	neg.f	00 0111	7	7	BEL	71	47	G
addi	jr		00 1000	8	8	BS	72	48	H
addiu	jalr		00 1001	9	9	HT	73	49	I
slli	movz		00 1010	10	a	LF	74	4a	J
slltu	movn		00 1011	11	b	VT	75	4b	K
andi	syscall	round.w.f	00 1100	12	c	FF	76	4c	L
ori	break	trunc.w.f	00 1101	13	d	CR	77	4d	M
xori		ceil.w.f	00 1110	14	e	SO	78	4e	N
lui	sync	floor.w.f	00 1111	15	f	SI	79	4f	O
(2)	mfhi		01 0000	16	10	DLE	80	50	P
mtli			01 0001	17	11	DC1	81	51	Q
mflo	movz.f		01 0010	18	12	DC2	82	52	R
mtlo	movn.f		01 0011	19	13	DC3	83	53	S
			01 0100	20	14	DC4	84	54	T
			01 0101	21	15	NAK	85	55	U
			01 0110	22	16	SYN	86	56	V
			01 0111	23	17	ETB	87	57	W
mult			01 1000	24	18	CAN	88	58	X
multu			01 1001	25	19	EM	89	59	Y
div			01 1010	26	1a	SUB	90	5a	Z
divu			01 1011	27	1b	ESC	91	5b	[
			01 1100	28	1c	FS	92	5c	\
			01 1101	29	1d	GS	93	5d]
			01 1110	30	1e	RS	94	5e	^
			01 1111	31	1f	US	95	5f	_
lb	add	cvt.s.f	10 0000	32	20	Space	96	60	`
lh	addu	cvt.d.f	10 0001	33	21	!	97	61	a
lwl	sub		10 0010	34	22	"	98	62	b
lwr	subu		10 0011	35	23	#	99	63	c
lbu	and	cvt.w.f	10 0100	36	24	\$	100	64	d
lhu	or		10 0101	37	25	%	101	65	e
lwr	xor		10 0110	38	26	&	102	66	f
	nor		10 0111	39	27	'	103	67	g
sb			10 1000	40	28	(104	68	h
sh			10 1001	41	29)	105	69	i
swl	slt		10 1010	42	2a	*	106	6a	j
sw	sltu		10 1011	43	2b	+	107	6b	k
			10 1100	44	2c	,	108	6c	l
			10 1101	45	2d	-	109	6d	m
			10 1110	46	2e	.	110	6e	n
swr			10 1111	47	2f	/	111	6f	o
cache									
ll	tge	c.f.f	11 0000	48	30	0	112	70	p
lwc1	tgeu	c.un.f	11 0001	49	31	1	113	71	q
lwc2	tlit	c.eq.f	11 0010	50	32	2	114	72	r
pref	tlit	c.ueq.f	11 0011	51	33	3	115	73	s
	teq	c.olt.f	11 0100	52	34	4	116	74	t
ldc1		c.ult.f	11 0101	53	35	5	117	75	u
ldc2	tne	c.ole.f	11 0110	54	36	6	118	76	v
		c.ule.f	11 0111	55	37	7	119	77	w
sc		c.sf.f	11 1000	56	38	8	120	78	x
swc1		c.ngle.f	11 1001	57	39	9	121	79	y
swc2		c.seq.f	11 1010	58	3a	:	122	7a	z
		c.ngl.f	11 1011	59	3b	;	123	7b	{
		c.lt.f	11 1100	60	3c	<	124	7c	
sdcl		c.ngf.f	11 1101	61	3d	=	125	7d	}
sdcl		c.le.f	11 1110	62	3e	>	126	7e	~
		c.ngt.f	11 1111	63	3f	?	127	7f	DEL

(1) opcode(31:26) == 0

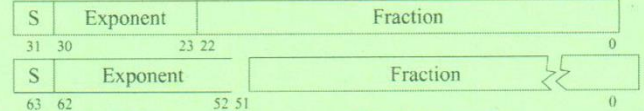
(2) opcode(31:26) == 17_{ten} (11_{hex}); if fmt(25:21) == 16_{ten} (10_{hex}) f = s (single);
if fmt(25:21) == 17_{ten} (11_{hex}) f = d (double)

IEEE 754 FLOATING-POINT STANDARD

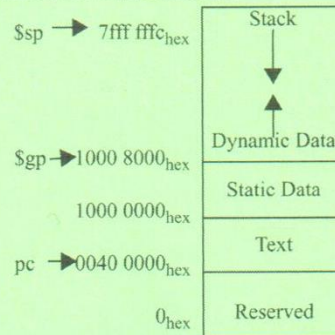
$$(-1)^S \times (1 + \text{Fraction}) \times 2^{(\text{Exponent} - \text{Bias})}$$

where Single Precision Bias = 127,
Double Precision Bias = 1023.

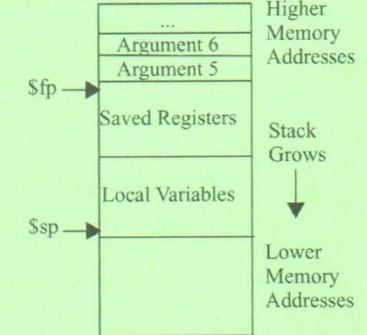
IEEE Single Precision and Double Precision Formats:



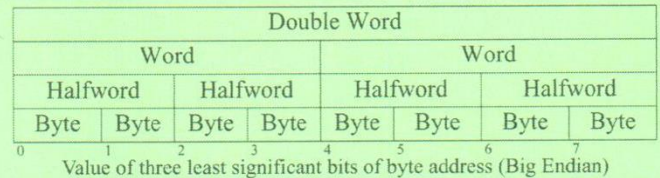
MEMORY ALLOCATION



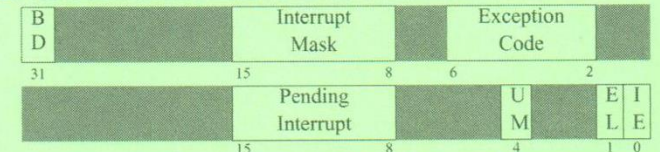
STACK FRAME



DATA ALIGNMENT



EXCEPTION CONTROL REGISTERS: CAUSE AND STATUS



BD = Branch Delay, UM = User Mode, EL = Exception Level, IE = Interrupt Enable

EXCEPTION CODES

Number	Name	Cause of Exception	Number	Name	Cause of Exception
0	Int	Interrupt (hardware)	9	Bp	Breakpoint Exception
4	AdEL	Address Error Exception (load or instruction fetch)	10	Rl	Reserved Instruction Exception
5	AdES	Address Error Exception (store)	11	CpU	Coprocessor Unimplemented
6	IBE	Bus Error on Instruction Fetch	12	Ov	Arithmetic Overflow Exception
7	DBE	Bus Error on Load or Store	13	Tr	Trap
8	Sys	Syscall Exception	15	FPE	Floating Point Exception

SIZE PREFIXES

	PREFIX	SYMBOL	SIZE	PREFIX	SYMBOL	SIZE	PREFIX	SYMBOL	SIZE	PREFIX	SYMBOL	SIZE
10 ³	Kilo-	K	2 ¹⁰	Kibi-	Ki	10 ¹⁵	Peta-	P	2 ⁵⁰	Pebi-	Pi	
10 ⁶	Mega-	M	2 ²⁰	Mebi-	Mi	10 ¹⁸	Exa-	E	2 ⁶⁰	Exbi-	Ei	
10 ⁹	Giga-	G	2 ³⁰	Gibi-	Gi	10 ²¹	Zetta-	Z	2 ⁷⁰	Zebi-	Zi	
10 ¹²	Tera-	T	2 ⁴⁰	Tebi-	Ti	10 ²⁴	Yotta-	Y	2 ⁸⁰	Yobi-	Yi	