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Course Code: CSA0389

Course Mame: Data Structure

Assignment - 2

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Describe the Concept of abstract data type (ADT) and how they differ from Concrete data structures. Design an ADT for a stack and implement it using arrays and linked list in c. Anclude operations (ike push, pop, peck, is compty, is full and peck.

Solf-Abstract Data Type (ADT):

An abstract data type (ADI) is a theoletical model that defines a set of operations and the semantics (behavior) of those operations on a datastructure, without specifying how the data structure should be implemented. It provides level description of what operations can be partitioned on the data and what constraints apply to those operations Characteristics of ADTs:

\*Operations: Defines a set of operations that can be performed on the data structure.

\* Semantics: - Specifies the behavior of each operation.

\* Encapsulation: - Hide the implementation details, tocusing on the the interface provided to the user.

ADT For Stack!

A Stack is a fundamental data structure that follows the last In First out (LIFO) Principle. It supports the Hollowing operations!

\*push: Adds an element to the top of the stack.

\* pop: Temoves and Returns the element from the top of stack.

\* peck! Returns—the element—from the top of the stack without senoving it.

in the following in the state of the first the party of t

Capara Harris

\* Is Empty: Checks if the stack is empty

\*Is Full: checks if the Stack is full.

```
Concrete Data Structures!
The implementations using arrays and linked lists are specific ways
of implementing the stack ADT in c.
How ADT differ Iron Concrete Data structures.
ADT tocuses on the operations and their behaviol, while concrete
data Structures focus on how those operations are realited
using specific programming constructs Carrays are (inted lists).
Advantages of ADT:
    Separating the ADT from its implementation, you achieve modulax
ity, encapsulation, and flexibility in designing and using datastructure
in programs. This separation allows tot easier maintanance, code reuse,
and abstraction of the complex operations.
 Implementation in C using Arrays:
#include Letdio.h>
# deline MAX-SS-ZE 100
type det struct s
     int items[MAX_SIZE];
     int toping in the stopped if supported (1111) 300
   I Stack Array;
 int maine > {
     StackArray Stack; qui- ad- 14 thinks
    Stack. top=-1;
      stack. items[++stack.top]=10,
     stack. items [++stack-top]=20;
      stack. "tems [++ stack. top] = 30;
    (stack. top! = -1) {
    printf ("Top clement: "do In", stack. items [stack. top]);
    4 else {
         printf (" stack is empty ! In");
```

if (stack.top!=-1) {
 pointf("poped element: "hd\n", stack.items[stack.top--]);

```
3 else ;
       printf ("Stack underflows in");
   if (stack . top!=1) { 1 } 1 } 1 | all and a state
      Points ("poped element: 1.d \n", stack. "tems (stack. top--]);
   Belse 8
                           - controlle portion de de
    y Print ("Stack underflow: (n");
   if (stack . top ! = -1) {
         print + ("Top element after pops: "led in" stack items [stack-top];
   3 clse 5
      Print ("Stack is empty: 10");
                               " tog ist of
  return o;
 implementation using linked list!
= include < stdio.h>
# Include < stdlib.h>
-typedel struct dode & by some by
     Int data;
     struct Node *next;
 3 Node;
     Node* top = Null; " and show done."
int main () &
Node * new Node = ( Node *) mallo c (Sizeof (Node));
     of (newhode = = Mull) {
         print+ (" memory allocation failed! (n");
         seturn 1;
                                    9 (1) Why = 1 10 + 2 - 17 - 1
     new Mode -> data=10;
     new lode -> next = top;
    -top = newrode;
     newlode = ( Node +) mallec (site of ( Node));
     : f (new Mode = Mull) }
        pointf (" memoly allocation lailed in ");
      yetwin 1;
```

```
newplode -> data=20;
   newplode -> next = top;
  -top= new lode;
   rewrode . (rlode *) nalloc (size of (rlode));
   ? + (new Mode == Mull) [
        printh (" memory allocation wiled ! In");
       return 1;
   new Mode -> data = 30;
   new rlode -> next + top;
   -top = new Mode;
   newhode = ( node + ) malloc (Sized ( node));
   if (top! = rlull) {
       Printh (" Top element : " " top -> data);
   3 else &
        Printf (" stack is empty: In");
   if (top!=Null) &
      dode * temp = top;
      printf (" popped element: lodin", temp-) data)
  to top= top- next;
      -free (temp);
   Jelse s
       Printf ("Stack underflow: In");
   if (-top 1 = Null) &
        print + ("Top clement after pops! 1/6 d in", top -> da ta);
    Jelse s
         printf ("Stack is empty: in");
   while (top! = wall) {
        lode * temp = top;
        -top = -top-> next;
         Free (-temp);
retugn o;
```

The university announced the selected conditions agreed muchus -for placement training. The student 441, 299,00, EDWEDTO SIGHE - to Check whether his name?s listed or not. The list is not saled in any Order . Sidentify - the seasoning technique - that can be opplied and explain the searching technique that can be opplied and explain the searching steps with the suitable procedure. List includes 20142015, 20142033, 20142011, 20142017, 20142010, 20142056, व्याप व्यवहरू

Great Seatch Cinear Search walks by checking each element in the list one by one until the desired element is found or the end of the list reached. Sit's a sample scorching technique that doesn't aquive any paux solling of the data.

Stops for Gnear Scorpch!

1. Start from the first element. 2. Check if the Current element is aqual to the target element.

3. If the current element is not the target, move to the next element in

the list. 4. Continue - this process until oither target element is found it you

reached the end of the list.

5. It -the -larget is found, returns it's position, It-the end of the list is reached and the element has not been found, indicate that element is not present.

Proceduje:

Given the list!

- 2014 2015, 20142033, 20142011, 20142017, 20142016, 20142056, 20142025.

1. Start at the first element of the list.

2. Compare '20142010' with '20142015' (first element), '2014 2033' (second element), 'adiuso 11' (this d'element), 'aduso 17' (fougth element) these are not aqual.

3. Confare solubolo' with 'solubolo' Chifth element). They are equal. 4. The plannent apivalous' is found at the Pasition (index 11) in the list.

```
C code -177 Cheax Seozch!
 # Include < stolio. h>
 Pot mosnes &
     305 pollog (105 plas, 105 plas, 2505 plas, 2105 plas)= [ ]230 plas)
     Int taget = 2042010;
     But D= Size of (signlumbers) /size of (signlumbers [0]);
     int found = 0;
    -13 (1=0; izn; 1++) {
          if (reg humbers[i] == toxget) {
              Printf (" Pogidsolien humber 1.1 hound at index bd. in", toxget, i);
             -found = 1;
           2 break;
       pentf ("Regionalism dumber and not found in list. In", target);
    31 ( found) }
 zebojn o;
 Exploration of the Code?
The "engroundress" exercy contains the list of registration numbers.
a garget is the registration number use are searching for.
2 0 % the total numbers of elements in orray
es Throse through each element of the oxfac.
5. I the Current element matches the 'target', Print its index and set
  the "found" flog to "!
5 It the loop completes without finding the target, point that the
   registration number is not found.
The program will print the "order of the found registration number or
   3 raticale that the registration is not present.
 output - Degistration number 20142010 Lound at "index 4.
```

```
isite pseudocade 137 stack operations.
 1. Instialize stack ():
      Initialize necessary ubgliable of structures to represent the stack.
2. push (clements):
      it stack is full:
           Print "stack overflow"
       else:
          add element to the top of the stack
          incorment top pointer
3. Pop ():
       if Stack is empty:
           point "stack underflow"
           Jeturn null (3) appropriate error value)
       else:
           semove and setup clement from top of the stack
          decrement end pointer.
4. peck():
                               in with the top in
        it stack is empty:
             point "stack is empty".
            setupo null (8) appropriate error value)
            setupo element at the top of the stack (without semoving it)
5. is Empty ():
       seturn true of top is -1 (stack is empty)
        otherwise, vetupo false
       vetupo true if top is equal to massite -1 (stack is full)
6. is full ():
        otherwise, return talse.
```

## -Explanation of the Pseudocock!

- + Inthalize—the necessary variables as obtain standingers—to represent
- + Adds an element to the top of the stock. Checks it the stack is full before pushing.
- \* Removes and between the element from the top of the stack Checks if the stack is empty before period.
- \* Returns the element at the top of the stack estimat remaining ?t. Checks if the stack is empty before pecking.
- -x Checks if the stack is empty by inspecting the top pointer or equivalent variable.
- + Checks if the Stack is full by composing the top pointer or equivalent utiliable to the modimum site of the stack.

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