

Department of Electrical and Computer Engineering

The University of Texas at Austin

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Problem Set 4 Solution

Yale N. Patt, Instructor

Stephen Pruet, Emily Bragg, Siavash Zangeneh, TAs

1. The number of occurrences of value x , in the input vector V_0, V_1, \dots, V_{n-1}

2. a) Single processor system - using Horner's rule

$$((((ax + b)x + c)x + d)x + e)x + f)x + g$$

Number of operations = 12 (6 multiplies & 6 adds)

Number of time-steps = 12

b) Using 4 processors: time steps = 5, operations = 15.

One way of doing this is:

Step 1	$x * x$	$b * x$	$d * x$	$f * x$
Step 2	$x^2 * x^2$	$a * x^2$	$dx + e$	$bx + c$
Step 3	$ax^2 * x^4$	$(bx + c) * x^4$	$(dx + e) * x^2$	$fx + g$
Step 4	$ax^6 + bx^5 + cx^4$	$dx^3 + ex^2 + fx + g$		
Step 5	$ax^6 + bx^5 + cx^4 + dx^3 + ex^2 + fx + g$			

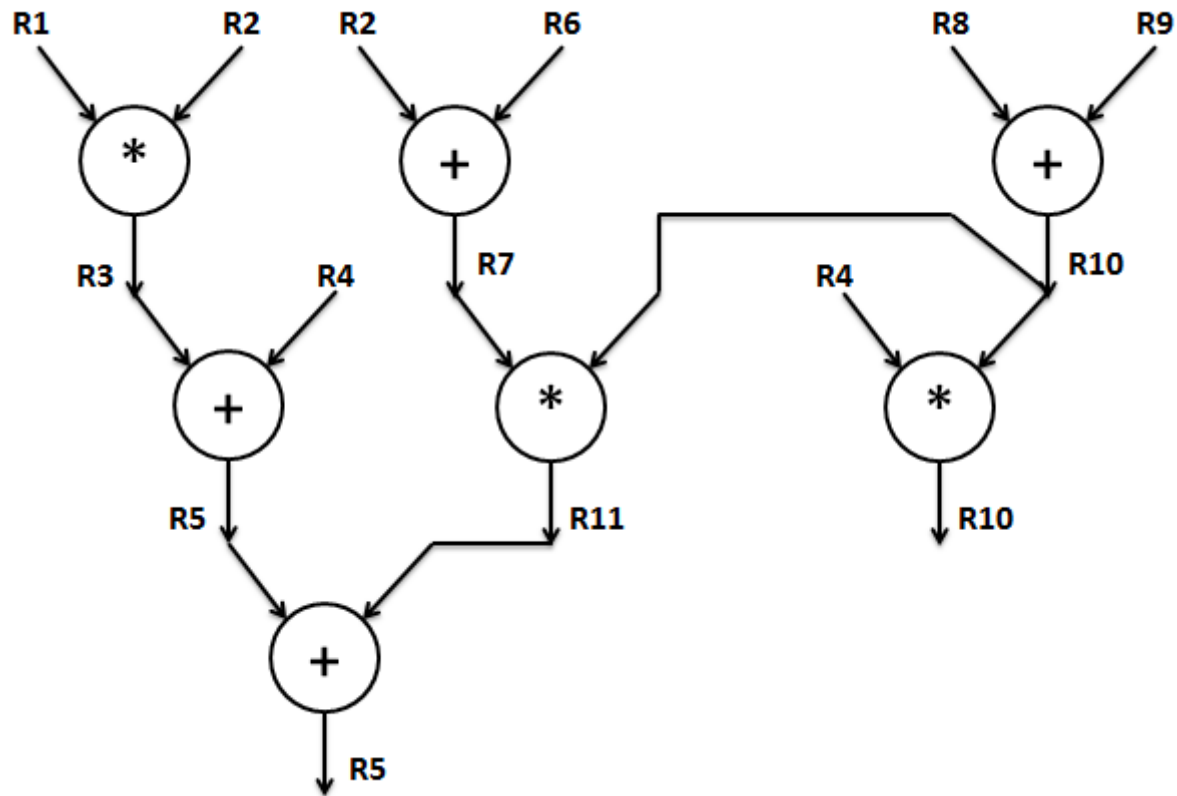
An intuitive proof for why it can't be done in less than 5 time steps:

Lets assume it takes x time steps. In the x th time step, only one operation can be performed (either an add or a multiply) to get the final result. In the $(x-1)$ th time step, at the most two operations can be performed, one each to get each of the operands for the operation in step x . In the $(x-2)$ nd time step & other below, possibly 4 units of work can be found. If $x = 4$, then theres no way to compute all the operands such that only 3 operations are left to be performed after 2 time steps.

c) $\text{Speedup} = T_1 / T_p = 12/5 = 2.4$

3. T_1 should be the time to solve the problem, using **the fastest** sequential algorithm and implementation.

4. The Solution:



5. The table is shown below:

		k = 2	k = 4	k = 8	k = 64
Number of switches in each level	n/k	$n/2$	$n/4$	$n/8$	$n/64$
Number of levels (= latency)	$\log_k n$	$\log_2 n$	$\log_4 n$	$\log_8 n$	$\log_{64} n$
Cost	$k^2 * n/k * \log_k n$ $= nk \log_k n$	$2n \log_2 n$	$4n \log_4 n$	$8n \log_8 n$	$64n \log_{64} n$

Reasonable choice depends on whether lower cost, latency, or contention is desired.

For $n=64$, Using $k=64$ yields the lowest latency, and $k=2$ or $k=4$ yields the lowest cost. If we would like low latency we would choose k to be 64. If we would like lowest cost, we would choose k to be 2 or 4. If we would like to minimize contention, we would choose k to be 64.

Another approach could be to minimize the cost-latency product, in which case we would choose $k=8$ for $n=64$.

6. Initial state of directory for cache line A: 0000

After processor 1 reads A: 1000

After processor 2 writes to A: 0101

After processor 3 reads A: 0110