

Manufactorying Problem——装配线呀嘛装配线,我家有条装配线

```
FASTEST-WAY(a, t, e, x, n)
   f1[1] = e1 + a(1,1)
   f2[1] = e2 + a(2,1)
   for j = 2 to n
       do if f1[j-1] + a(1,j) <= f2[j-1] + t(2,j-1) + a(1,j)
        then f1[j] = f1[j-1] + a(1,j)
             11[j] = 1
        else
             f1[j] = f2[j-1] + t(2,j-1) + a(1,j)
             11[j] = 2
       if f2[j-1] + a(2,j) <= f1[j-1] + t(1,j-1) + a(2,j)
        then f2[j] = f2[j-1] + a(2,j)
             12[j] = 2
        else
             f2[j] = f1[j-1] + t(1,j-1) + a(2,j)
             12[j] = 1
   if f1[n] + x1 <= f2[n] + x2
       then final = f1[n] + x1
             1 fin = 1
        else
             final = f2[n] + x2
             l_fin = 2
```

MATRIX-CHAIN—矩阵链乘! 乘! 乘! 乘!

```
m[i,j] = \begin{cases} 0 & \text{if } i = j \\ \min_{i \le k \le j} \{ m[i,k] + m[k+1,j] + p_{i-1}p_kp_j \} & \text{if } i \le j \end{cases}
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```
MATRIX-CHAIN-ORDER(p)

n = length[p] - 1

for i=1 to n

do m[i,i] = 0 #最底层首先初始化为0

for l=2 to n #1代表链长

do for i=1 to n-l+1

do j = i+l-1

m[i,j] = MAX #无穷

for k=1 to j-1

do q = m[i,k]+m[k+1,j]+p[i-1]*p[k]*p[j]

if q<m[i,j]

then m[i,j] = q

s[i,j] = k

return m and s
```