Week 10

Project 3: An Introduction to File Systems

Classes COP4610 / CGS5765 Florida State University

Introduction

- The goal of project 3 is to understand
 - basic file system design and implementation
 - file system testing
 - data serialization/de-serialization
- At the end of the project, you will feel like a file system expert!

Outline

- Background
 - Setting up your environment
 - Mounting file systems
- Project 3
 - Specification
 - Downloading and testing file system image
 - General FAT32 data structures
 - Endian-ness

Environment Setup

Get ready for Project 3!

Project 3 Environment

- Must develop inside Linux environment with root access
 - Use Debian virtual machine or lab machine from project 2
- I will be grading your projects inside of my copy of the class virtual machine

Project Environment

- Kernel version no longer matters
 - Can go back to the default 2.6.26 version if you trust it more
- Entire project will be in userspace, not kernel!
 - Please use debuggers, they will save you time
 - gdb, ddd, others...
- Programming language is still C

Running out of room?

- You must have at least 64MB free plus room for your source code
- To see how much room you have left inside your machine, issue the following command:

df-h

```
    user@cop4610: ~

user@cop4610:~$ df -h
                           Used Avail Use% Mounted on
Filesystem
                     Size
/dev/sda1
                           5.0G 2.2G 70% /
                    7.5G
tmpfs
                    1014M
                              0 1014M 0% /lib/init/rw
udev
                      10M
                           644K 9.4M
                                       7% /dev
tmpfs
                    1014M
                                        0% /dev/shm
                              0 1014M
user@cop4610:~$
```

- /dev/sda is root file system mounted on "/"
 - Has 2.2GB currently available

Cleaning up space

- If you need to clean up some more space
 - Issue 'make clean' inside of your 2.6.32 kernel source
 - Delete the 2.6.26 kernel source

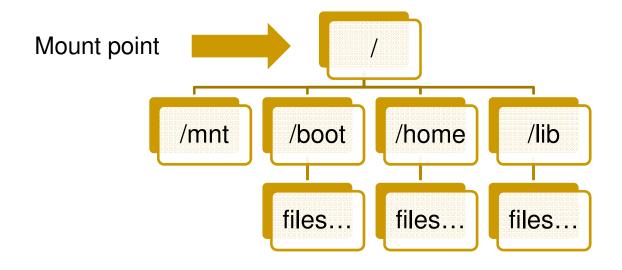
Mounting File Systems

Unix File Hierarchy

- All files accessible in a Unix system are arranged in one big tree
 - Also called the *file hierarchy*
 - Tree is rooted (starts) at /
- These files can be spread out over several devices
- The mount command serves to attach the file system found on some device to the big file tree

'mount' command

- mount
- mount <device> <mount directory>
- Typing 'mount' without arguments shows you what is mounted and where
- Second example attaches a device or partition to a directory
 - Must have root privileges



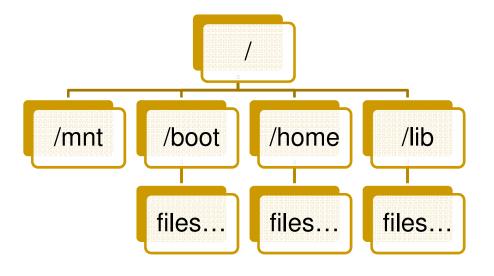


The device sda partition 1 is mounted at "/". All files and dirs below "/" come from this device.

 Type command 'mount' without any arguments to see what is mounted and where

```
user@cop4610:~

user@cop4610:~$ mount
/dev/sda1 on / type ext3 (rw,errors=remount-ro)
tmpfs on /lib/init/rw type tmpfs (rw,nosuid,mode=0755)
proc on /proc type proc (rw,noexec,nosuid,nodev)
sysfs on /sys type sysfs (rw,noexec,nosuid,nodev)
procbususb on /proc/bus/usb type usbfs (rw)
udev on /dev type tmpfs (rw,mode=0755)
tmpfs on /dev/shm type tmpfs (rw,nosuid,nodev)
devpts on /dev/pts type devpts (rw,noexec,nosuid,gid=5,mode=620)
user@cop4610:~$
```

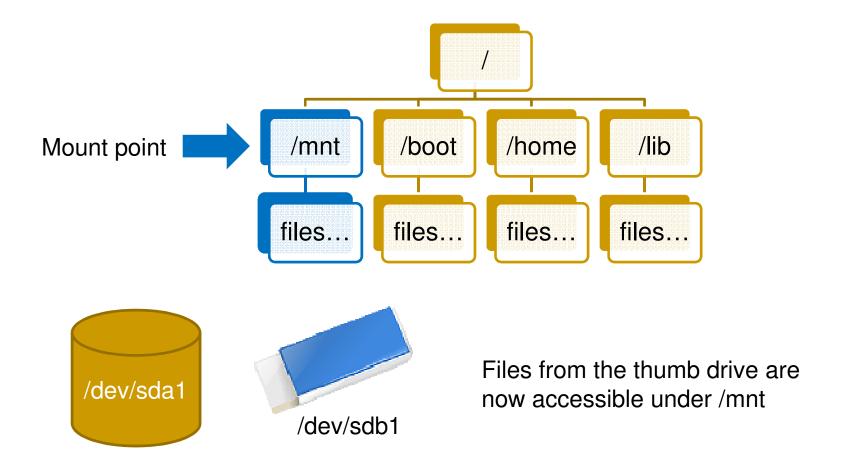






Now suppose we attach a thumb drive and want our thumb drive files accessible under /mnt...

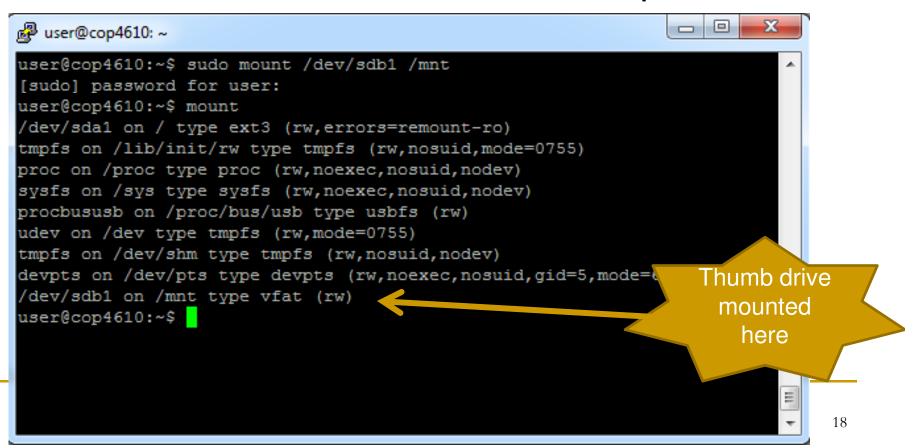
File Hierarchy Example



 The 'mount' command can dynamically attach new devices to new mount points

```
user@cop4610: ~
user@cop4610:~$ sudo mount /dev/sdb1 /mnt
[sudo] password for user:
user@cop4610:~$ mount
/dev/sda1 on / type ext3 (rw,errors=remount-ro)
tmpfs on /lib/init/rw type tmpfs (rw,nosuid,mode=0755)
proc on /proc type proc (rw,noexec,nosuid,nodev)
sysfs on /sys type sysfs (rw,noexec,nosuid,nodev)
procbususb on /proc/bus/usb type usbfs (rw)
udev on /dev type tmpfs (rw,mode=0755)
tmpfs on /dev/shm type tmpfs (rw,nosuid,nodev)
devpts on /dev/pts type devpts (rw,noexec,nosuid,gid=5,mode=620)
/dev/sdb1 on /mnt type vfat (rw)
user@cop4610:~$
```

 The 'mount' command can dynamically attach new devices to new mount points



Un-mount Command

■ umount <dir>

- In our example where the thumb drive was mounted at /mnt, we can issue
 - □ \$> umount /mnt
 - Must have root privileges

Figuring out names of devices

- /etc/fstab Has list of devices and file systems that get auto-mounted on boot
- 'dmesg' command shows output when plugging in a dynamic device

Project 3

More than you wanted to know about FAT32..

Project 3

- You will create a user-space utility to manipulate a FAT32 file system image
 - No more kernel programming!
- Utility must understand a few basic commands to allow simple file system manipulation
- Utility must not corrupt the file system and should be robust

FAT32 Manipulation Utility

Utility only recognizes the following built-in commands:

- open
- close
- create
- rm
- size

- cd
- Is
- mkdir
- rmdir
- read
- write

File System Image

- Manipulation utility will work on a preconfigured FAT32 file system image
 - Actually a file
- File system image will have raw FAT32 data structures inside
 - Just like looking at the raw bytes inside of a disk partition

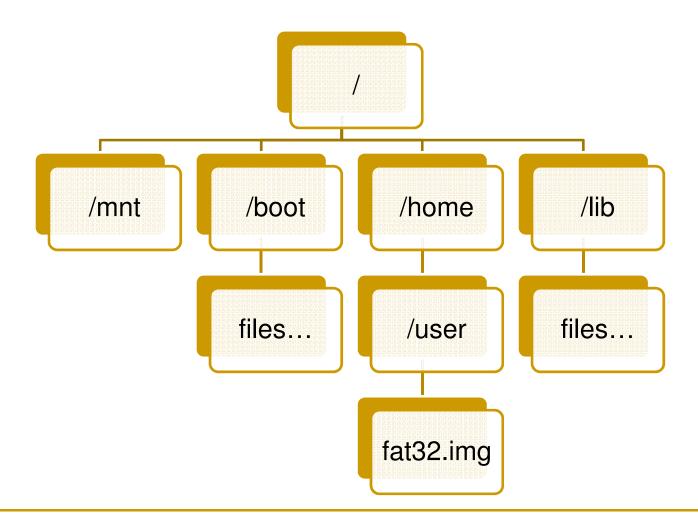
File System Image

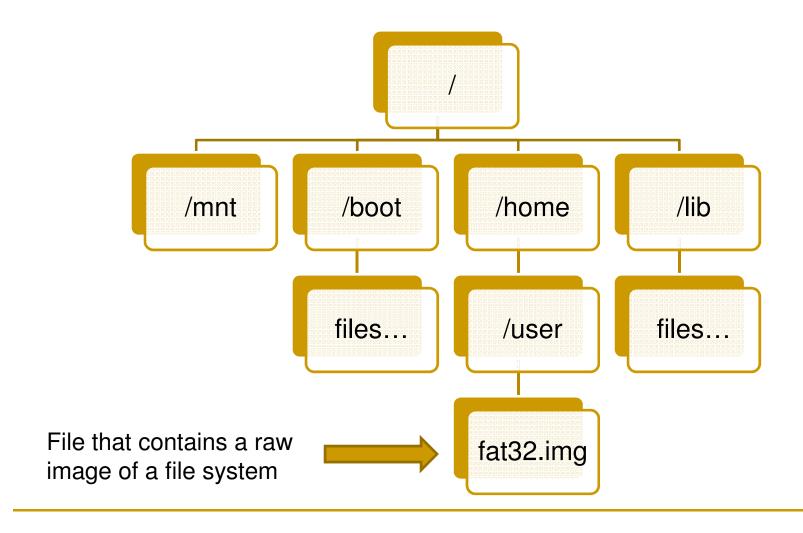
- Your FAT32 manipulation utility will have to
 - Open the FAT32 file system image
 - Read parts of the FAT32 file system image and interpret the raw bytes inside to service your utility's file system commands...

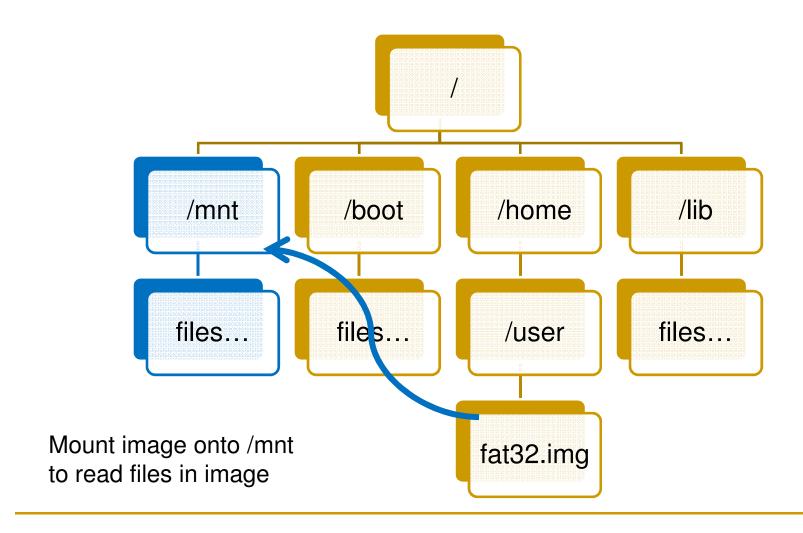
...just like a file system!

File System Image

- Sometimes you may want to check that you haven't corrupted your file system image, or that you can add or write files successfully
 - Mount your file system image with the OS FAT32 driver
 - Just like the file system image is a device







```
$> sudo mount -o loop fat32.img /mnt
$> cd /mnt
```

- fat32.img is your image file
- /mnt is your mounting directory
- Once the file is mounted, you can go into the /mnt directory and issue all your normal file system commands like:
 - □ ls, cat, cd, ...

General FAT32 Data Structures

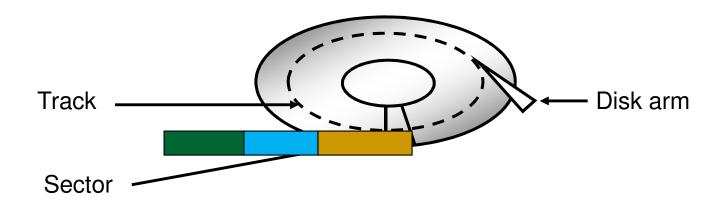
Terminology

- Byte 8 bits of data, the smallest addressable unit in modern processors
- Sector Smallest addressable unit on a storage device. Usually this is 512 bytes
- Cluster FAT32-specific term. A group of sectors representing a chunk of data
- FAT Stands for file allocation table and is a map of files to data

FAT32 Disk Layout

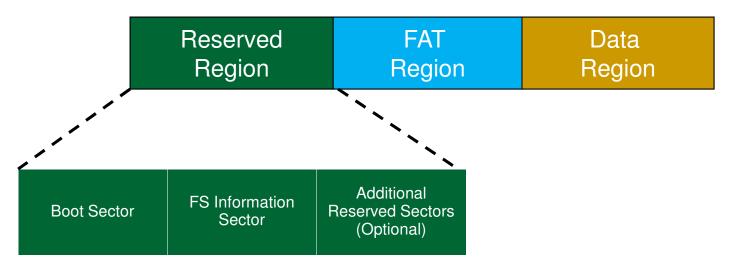
3 main regions...





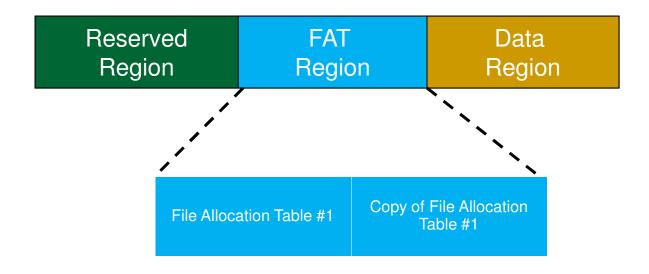
Reserved Region

Reserved Region – Includes the boot sector, the extended boot sector, the file system information sector, and a few other reserved sectors



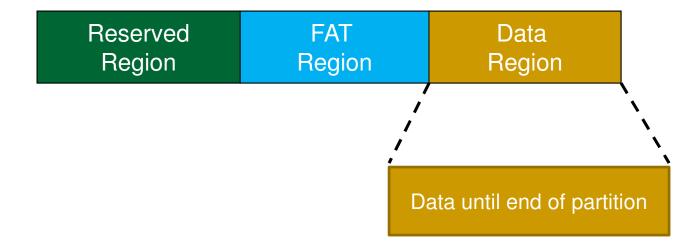
FAT Region

 FAT Region – A map used to traverse the data region. Contains mappings from cluster locations to cluster locations



Data Region

Data Region – Using the addresses from the FAT region, contains actual file/directory data



Endian

Big or little?

Machine Endianness

- The endianness of a given machine determines in what order a group of bytes are handled (ints, shorts, long longs)
 - Big-endian most significant byte first
 - Little-endian least significant byte first
- This is important to understand for this project, since FAT32 is always formatted as little-endian

FAT32 Endianness

- The following are a few cases where endianness matters in your project:
 - Reading in integral values from the FAT32 image
 - Reading in shorts from a FAT32 image
 - Combining multiple shorts to form a single integer from the FAT32 image
 - Interpreting directory entry attributes

Endian Example (English Version)

- Imagine you can only communicate three letters at a time, and your word is "RAPID"
- Big-endian
 - 1. RAP
 - □ 2. ID
 - Word = RAPID
- Little-endian
 - 1. PID
 - 2. RA
 - Word = PIDRA (come again?)

Endian Example (data version)

- short value = 15; /* 0x000F */
- char bytes[2];
- memcpy(bytes, &value, sizeof(short));
- In little-endian:
 - \Box bytes[0] = 0x0F
 - \Box bytes[1] = 0x00
- In big-endian:
 - \Box bytes[0] = 0x00
 - bytes[1] = 0x0F

Endian Example (data version 2)

- int value = 13371337; /* 0x00CC07C9 */
- char bytes[4];
- memcpy(bytes, &value, sizeof(int));
- In little-endian:
 - \Box bytes[0] = 0xC9
 - bytes[1] = 0x07
 - \Box bytes[2] = 0xCC
 - \Box bytes[3] = 0x00

In big-endian:

- \Box bytes[0] = 0x00
- \Box bytes[1] = 0xCC
- \Box bytes[2] = 0x07
- \Box bytes[3] = 0x09

Visualizing Example 2 Value = 13371337 (0x00CC07C9)

index	0	1	2	3
little endian	0xC9	0x07	0xCC	0x00
big endian	0x00	0xCC	0x07	0xC9

Additional Project 3 Information

- Like other projects, may work in teams or alone
- Project deadline is December 3rd, 2010 at 11:59:59pm (Friday before finals week)
 - No extensions
 - Can use remainder of slack days
 - Cannot be more than 3 days late
 - Monday, December 6th is absolute last day to turn in

Project 3 Partners

- Raise your hand if you are looking for a partner...
- Also use the discussion board

Until Next Time

- Set up your environment
- Download the image file
- Practice mounting the image file with the OS FAT32 drivers
 - Make sure you can cd into /mnt and read/write to the files
- Read over the FAT32 Specification