Egght: $R_n = n\mu^* - E\left[\sum X_E\right]$ We Algorithm

First, recall:

For $E \ge 0$, $P(M_{i,n} \le M_i - E)$ $M_{i,n} = M_{i,n} + \sqrt{\frac{2\log(1/8)}{n}}$ $P(N_{i,n} \le M_{i,n} + \sqrt{\frac{2\log(1/8)}{n}})$

$$U(B_{i}(t-1)) = \begin{cases} \infty & \text{if } T_{i}(t-1) = 0 \\ \mu_{i}(t-1) + \sqrt{\frac{2\log(1/8)}{T_{i}(t-1)}} \end{cases}$$
empirical estimate

Ti(t): number of times the arm has been pulled up to time t

Algo:
At time
$$t \ge 1$$
,
$$-A_t = \underset{i}{\operatorname{argmax}} UCB_i(t-1)$$

- Observe rewards, update UCBs

Idea: UCB is an optimistic estimate of the mean of arm i.

We are pulling the arm with the highest optimistic estimate

"Optimism in the face of Uncertainty"

 $T_i(t-1)$ is a random variable and not deterministic, unlike the n in the earlier setting.

: Ti(t-1) is a random quantity dependent on the past

we cannot blindly say that this UCB is a perfectly bound

Intuition: In the line
$$\mu i \leq \mu_{i,n} + \sqrt{\frac{2\log(1/8)}{\eta}}$$

if the n was dependent on μ_i , then this bound will not be valid.

Theorem: With $\delta = \frac{1}{n^2}$, UCB has regret

$$R_n \leq 3 \sum \Delta i + \sum_{i:\Delta_i \geq 0} \frac{16 \log(n)}{\Delta i}$$

. The algorithm is not "ANYTIME", i.e. it requires knowledge of the horizon 'n' Note: The regret is logarithmic wit to the horizon 'n' Note: The bound is instance dependent bound (i.e. depends on Di) Recall: $R_n = \sum \triangle_i \mathbb{E}[T_i(n)]$ Proof: Fix suboptimal arm i. WLOGI, arm 1 is optimal arm. $Gi = \begin{cases} G \mu_1 < \min_{t \in [n]} UCB_1(t) \end{cases} \cap \begin{cases} \mu_{i,ui} + \sqrt{\frac{2\log(1/8)}{u_i}} < \mu_i \end{cases}$ goodevent associated with the UCB of aum 1 is to be specified asm i never "violated" later (think of it like a constant) Mi M1 (UCB lies here) EVENT 2 (UCB; at the end of uipulls) $\mathbb{E}\left[T_{i}(n)\right] = \mathbb{P}\left(G_{i}\right)\mathbb{E}\left[T_{i}(n)/G_{i}\right] + \mathbb{P}\left(G_{i}^{c}\right)\mathbb{E}\left[T_{i}(n)/G_{i}\right]$ < 1 E[T;(n) | G;] + P(G; c) n (#) Claim: Under Gi, Ti(n) < Ui Proof: Suppose Gi; occurs and Ti(n) > ui ⇒ ∃ t s.t Ti(t) = U; , A; (t+1) = i Now, Mi,ui + \frac{2\log(1/8)}{ui} < \mui < UCBI(6) ... under Gi → Arm i not pulled (#) $\mathbb{E}[T_i(n)] \leq u_i + n P(G_i^c)$ Bounding IP(Gic) Gic = { M = min UCB, (t) } U { Mi, ui + \frac{2log(1/5)}{ui} > Mi} ... (DeMorgan Laws

$$\frac{P(T)}{P(T)}: T \subseteq \left\{ M_1 \gtrsim \min_{S \in [n]} M_1 \lesssim + \sqrt{\frac{2\log(V_S)}{n \cdot S}} \right\}$$

$$= \bigcup_{S \in [n]} \left\{ \mu_1 \gtrsim \mu_{1,S} + \sqrt{\frac{2\log(V_S)}{s}} \right\}$$

$$\Rightarrow P(T) \leq nS \qquad \text{(union bound, further } P(\text{sub-event}) \leq S)$$

$$\frac{P(K)}{P(K)}: P\left(M_{i,u_i} - M_i \geqslant \Delta_i - \sqrt{\frac{2\log(V_S)}{u_i}}\right)$$
Set $u_i \leq t$.
$$\Delta_i - \sqrt{\frac{2\log(V_S)}{u_i}} \geqslant \Delta_i$$

$$\left\{ u_i = \int_{\frac{3}{2}}^{\frac{3}{2}\log(V_S)} \right\}$$

$$\left\{ u_i = \int_{\frac{3}{2}}^{\frac{3}{2}\log(V_S)} \right\}$$

$$P(K) \leq P\left(\mu_{i,u_i} - \mu_i \geqslant \Delta_i \right)$$

$$\leq e^{\frac{3}{8}} \leq n^{-2} \qquad \text{(Note: } e^{\frac{3}{8}} \leq S \text{, here } S = \frac{1}{n},$$

$$E\left[T_i(n)\right] \leq \int_{\frac{3}{2}}^{\frac{3}{2}\log(V_S)} + n\left(nS + \frac{1}{n^2}\right)$$

$$E\left[T_i(n)\right] \leq 3 + \frac{16\log n}{\Delta_i^2} \qquad \dots \text{(Regret Decomposition Function)}$$

$$\Rightarrow R_n \leq 3 \leq \Delta_i + \sum_{i: \Delta_i \geq 0} \frac{16\log(n)}{\Delta_i} \qquad \dots \text{(Regret Decomposition Function)}$$

E: Can we achieve the Laid Robbins' Bound by choosing Ui even more smartly?

 $\underline{\beta}$: Is $\delta = \frac{1}{n^2}$ the most optimal choice?

If δ is not chosen to be powerlaw (n) and something like e^{-n} , we shall end up exploring too much.

Morals

i)
$$E[T_i(n)] = O\left(\frac{\log(n)}{\Delta_i^2}\right)$$

2) Note that the way we have derived this bound tries to suggest that for a modest (n), IRn increases very fast when Di decreases. This is not very intuitive & logical.

Thm For
$$S = \frac{1}{n^2}$$
, for UCB,

$$R_n \leq 8\sqrt{n K \log(n)} + 3 \sum \Delta i$$

(Removing inverse-dependence on the sub-optimality gap) This bound is also sometimes called the Worst-Case bound or the instance - independent bound, though it is not completely fitting to say so

$$\frac{\text{Proof}}{\text{i: } \Delta_{i} < \Delta_{i}} : R_{n} = \sum_{i: \Delta_{i} < \Delta_{i}} \Delta_{i} \text{ E[T_{i}(n)]} + \sum_{i: \Delta_{i} \geq \Delta_{i}} \Delta_{i} \text{ E[T_{i}(n)]}$$

$$i: \Delta_{i} < \Delta_{i} = 0$$

$$i: \Delta_{i} \geq \Delta_{i} \text{ exp}$$

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$$\leq n\Delta + \sum_{i:\Delta_i \geq \Delta} \left(3\Delta_i + \frac{6\log n}{\Delta_i} \right) - \frac{6\log 2n}{\Delta_i} = \left(\frac{1}{2} \right)$$

$$\leq 3 \sum \Delta i + n \Delta + \frac{16 K \log(n)}{\Delta}$$

If we set
$$1 \Delta = \sqrt{\frac{16 \, \text{K log(n)}}{n}}$$
, we get

254MB-64 = -4-1.5(X+ A-B) -0.5(X+A-B)