	THE RESIDENCE OF STREET OF
EE6106	Orline Learning and Optemization Lecture 1 (12/01/2024)
	A policy of All robots these
Gradin	q -> Class participation (5:/)
10-27	Quizzes/Homework (20%)
-	Midterm Exam (30%)
	Endterm Exam (45%)
7	which is the second
Learning	g from Experts: Adversarial setting
	· ?X+? is a binary sequence. (arbitrary). (Finite
	horizon will be generally discussed. {0,1}
	horizon will be generally discussed. {0,1}  At time t > 1, predictions (Yi, t, 1 < i < k) of k experts
	are revealed.
	· algorithm takes action A > your prediction
1	· Finally, Xt is revealed once the (:: Yi, t \in \{0,1\})
	. Total barizon: T & Note that algorithm makes mistake
	prediction has been made by us  Total horizon: T { Note that algorithm makes mistake }  when $A_t \neq X_t$
	Call it to it is the total number of mistakes
	· Gloat a to minimize the actal T
	minimize T 1 = LT Loss  Minimize T 1 = LT Total
	t = 1

Here, we shall see that T being apriori known/unknown does not matter
Take 1: Assume 3 perfect expert (You don't know who it is)
be more than 1
MATORITY ALCORITUM
MAJORITY ALGORITHM > log_K bound
$T_1 = \{1, 2, \dots, K\}$ $T_1 = \{1, 2, \dots, K\}$
$\frac{1}{t} = \text{set of trusted experts at time t}$ $At t \ge 1,$
At = majority action from { Yit : it Tt}
$T = T_{t+1} = T_{t} \setminus \{i : Y_{i,t} \neq X_{t}\}$
E+1 + 1 + 1/E + 1/
Whenever we make a mistake, the cardinality of the set of trusted
experts reduces by half a minimum has principal and a minimum has been applied to the minimum has been applied
experts reduces by half  Claim: Under MA, LT \le \log_2 K  [UB]
050057
9: What if the set {0,13 (alphabet) was something other than binary?
17 1/t e 1/1/2/11/11 ! (   09 K : / as   ong as K > m
S. What if an expert is allowed to make upto n mistakes?
UB 1 n can be a function of T
the Experis Adversaria satting
log_K " " compos tounid a far contril and
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
Diose K Lime t 3 1 prodictions ( Y & desert) of L see a
Laboration of the state of the
Take 2: No assumption about the experts' existence
- Idea: Given some hind of a priority order to the experts by means
of weights und show and you contribute
and a second a carlow continuous tests stalls to a second good . The
WEIGHTED MAJORITY ALGORITHM
- Initialize $\omega_{i,1} = 1 + i$
- At time t = 1, what is the weight of all the experts

	Page No.
	Oate
San Maria	that are predicting 1?
and the same of th	$W_{\alpha} = \sum_{i} \omega_{i}$ , 11.
S. R. S.	$W_{0,E} = \sum \omega_{n,e}  \mathcal{I}_{\{Y_{i,E} = 0\}}$
And believed many productions.	のでは、これでは、これでは、これでは、これでは、これでは、これでは、これでは、これ
graph manufacture and a	$W_{ij}e = \sum \omega_{1,t}  11_{\{Y_{i,t}=1\}}$
Secret Land Control	$A_{t} = 11_{\{W_{V} \in W_{0}, e\}}$
STREET, THE PROPERTY AND	EWVE > Mole}
A CONTRACTOR OF THE PARTY OF TH	to it named a
Unce	f is revealed.
A STATE OF THE PARTY OF THE PAR	$\omega_{i,t+1} = \omega_{i,t}$ if $Y_{i,t} = X_t$
	$\omega_{i,t+1} = \omega_{i,t} (1-\beta)$ if $Y_{i,t} \neq X_{t}$
and the second second	B ∈ (0,1) is called the learning rate
and	weights -> post-facto the weights which were correct
1000	is equal to WMA (1)
IVICA	$\beta = 1$
ise p	elential functions or lyapunov function
	E Comment of the second of the
Anal	ysis of WMA: TDEA
1 2-4	$W_t = \sum_{i=1}^{\infty} w_{i,t} \longrightarrow \text{can be bound } W_T \text{ from above & below?}$
	CY
	$W_4 = K$ (by construction)
Lower	Bound: Let the number of mistakes of expert i until time t
	be Lit
For	any i, we have the following trivial bound
	$W_{T+1} \ge (1-\beta)^{L_{ij}T}$
	T+1   weight expert i will
	we are have at the end
	also why
	also why write the inequality not saying that if you can compute an we know the "i" to choose
1:	
UPI	per Bound:
	look at the observation every time up make an error.
	look at the observation every time us make an error.

