EE6106 Lecture 11 (Date: 5th March 2024) pecal p, Q ~ measures on (D, f). $p(A) + Q(A^c) \ge \frac{1}{2} exp(-D(P, Q))$ EE6106 = 0 Moral: If IP and Q are "close", then P(A) + Q(A') cannot be small . D(Pu, Pu') = \(\subseteq D(\vi, \vi') \mathbb{E}_v[\tau:(n)] . E ~ set of arm distributions ≥ E ~ set of MAB instances Rn(IT, EK) = sup Rn(TT, V) worst case regret of IT on Ek most robust worst case policy . $R_n^*(\varepsilon^k) = \inf_{\Pi \in \Pi_1 \in \mathbb{R}} R_n(\Pi_1 \varepsilon^k)$ $\longrightarrow \min_{\Pi \in \Pi_1 \in \Pi_2} \operatorname{minimax regret}$ Ky (TC, N) = Ry (11, N) 3 Rn (ek) = inf sup Rn(π, ν)
π νε εκομηνη (π) lower bound on minimax regret: means that any policy in the worst case will have at least as much regret THEOREM Let K>1, n> K-1. Then for any policy TT, I pe & Ex (1) Whow the Light on the way i was John W Tino J. W. J.s. 1 (11, M) ≥ 1 √n(K-1) tolding the arginer $\mathcal{E}_{\mathcal{N}}^{k}(1)$: Gaussian distributions with variance 1 for all arms' rewards r Effectively, this theorem gives a lower bound on the minimax regret Note: When we did UCB, there was a NAKLOGK upper bound that we had got Proof Proof Fix policy TI, let △ ∈ [0,1/2] ... open interval Consider $M = (\Delta, 0, ..., 0)$

i = argmin Ex [Tj(n)]

j=1

Note: i is the arm which on average gets pulled the least.

Note:
$$R_{n}(\pi,\mu) \geq P_{\mu}(T_{i}(n) \leq \frac{n}{2}) \frac{n\Delta}{2}$$

This is a bad event since we are pulling the optimal arm less than N_{i} times: \underline{Ida} : lover bound regret in terms of a bad/have N_{i} times: \underline{Ida} : lover bound regret in terms of a bad/have N_{i} times: \underline{Ida} : \underline

: $\underset{j>1}{\operatorname{argmin}} \notin [T_j(n)] = i \implies \mathbb{E}_{M}[T_i(n)] \leq \frac{n}{k-1}$

orm which on average it is pulsed the land

Construct a new arm M.

Then:
$$R_{1}(\Pi_{1},M) + R_{1}(\Pi_{1},M') \geq \sqrt{n(k-1)}$$

$$R_{2}(\Pi_{1},M) + R_{2}(\Pi_{1},M') \geq \sqrt{n(k-1)}$$

$$R_{3}(\Pi_{1},M) + R_{2}(\Pi_{1},M') \geq \sqrt{n(k-1)}$$

$$R_{3}(R_{2}(\Pi_{1},M) + R_{2}(\Pi_{1},M')) \geq \sqrt{n(k-1)}$$

$$R_{4}(\Pi_{1},M) + R_{2}(\Pi_{1},M') \geq \sqrt{n(k-1)}$$

$$R_{5}(R_{2}(\Pi_{1},M) + R_{2}(\Pi_{1},M')) \geq \sqrt{n(k-1)}$$

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$$R_{5}(R_{5}(\Pi_{1},M') + R_{5}(\Pi_{1},M')) \geq \sqrt{n(k-1$$

Note: In terms of estimation jargon, consistency tells that the estimator will converge to true means as the number of samples goes to infinity.

Note: $R_n(\Pi, V) = O(n^q)$ H O is "Little O".

Hence, logarithmic and poly-logarithmic policies are consistent let $\mu^* \in \mathbb{R}$, arm is s.t.

 $\mu(v_i) < \mu^*$ $\dim \left(v_{i,\mu} *, \epsilon \right) = \inf \left\{ \mathcal{D}(v_i, v_i') : \mu(v_i') > \mu_i^* v_i' \in \epsilon \right\} \tag{1}$

Interpretation of (1):

Pick an arm i. μ^* is some greater number. What is the slightest

perturbation required to push the KL divergence higher? to push its

perturbation required to the

mean to higher than μ^* ?

M(Vi) μ^*

Note that the perturbation is in the KLDivergence and not in the mean.

Ans: Because we are working in a general family and hence the distributions are no are parametrized by not just the means, but other parameters as well.

Note: It is possible for a "rich" even family that even a slight perturbation in the KLD can push the means far far away.

THEOREM: Let To be a consistent policy on E". For v E E'

 $\underset{n\to\infty}{\liminf} \frac{R_n(\pi,\nu)}{\log n} \geq \underbrace{\frac{\Delta_i}{\dim f(\nu_i,\mu^*,\varepsilon)}}$

where μ^* is the optimal mean in ν

Note: This tells us that logarithmic regret is the best that we can do

Corner case: If the family of distributions in very rich, then dinf (Vi, M*, E) = 0 Then the regret becomes super-logarithmic & then logarithmic regret becomes impossible.

R. (11, 0) = 0 (ma) + 0 = 0

Jof: Policy IT is Consistent on E' if Aus E'

Announcements

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① HW2 is due on next Tuesday (12th March)

② Extra class on a handful of Saturdays (9th March & the first)

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3 Extra class tomorrow 530 PM (6th March)