

* EEG106 Lecture 2 (Date : 16 Jan 2024)

WMA : $w_{i,1} = 1$

At time $t \geq 1$

$$- w_{0,t} = \sum w_{i,t} \mathbb{1}_{\{x_{i,t} = 0\}}$$

$$- w_{1,t} = \dots$$

$$- A_t = \mathbb{1}_{\{w_{1,t} \geq w_{0,t}\}}$$

$$- w_{i,t+1} = \begin{cases} w_{i,t} & \text{if } y_t = x_t \\ w_{i,t}(1-\beta) & \text{if } y_t \neq x_t \end{cases}$$

$$W_t = \sum_{i=1}^K w_{i,t}$$

$$W_{T+1} \geq (1-\beta)^{L_{i,T}}$$

$$W_{T+1} \leq \left(\frac{1-\beta}{2}\right)^{L_T K}$$

Principle : Every time we make a mistake, the weight shrinks by at least a factor of $\left(\frac{1-\beta}{2}\right)$

Note : L_T = errors by algorithm

$L_{i,T}$ = error made by each expert

$$\left(\frac{1-\beta}{2}\right)^{L_T K} \geq (1-\beta)^{L_{i,T}}$$

$$\therefore L_T \log\left(\frac{1-\beta}{2}\right) + \log K \geq L_{i,T} \log(1-\beta)$$

Just for "aesthetic sense" :

$$\boxed{\frac{1-\frac{1}{x}}{x} \leq \log x \leq x-1} \quad \text{for } x > 0$$

all agree at $x=1$

$$\therefore L_T \left(\frac{-\beta}{2}\right) + \log K \geq L_{i,T} \left(\frac{-\beta}{1-\beta}\right)$$

$$\Rightarrow \boxed{L_T \leq \left(\frac{2}{1-\beta}\right) L_{i,T} + \frac{2}{\beta} \log K}$$

expert's error (choose best expert!)

$$\therefore L_T \leq \frac{2}{(1-\beta)} \left(\min_i L_{i,T} \right) + \frac{2}{\beta} \log K$$

\downarrow
 ≥ 2

Define regret $R_T := L_T - \left(\min_i L_{i,T} \right)$ i.e. $\lim_{T \rightarrow \infty} \frac{R_T}{T} = 0$

should ideally be sublinear wrt $t(T) \rightarrow$ Want $R_T = o(T)$

Note that for a finite horizon setting, it is also possible to incur negative regret

Note: If we know that there is a best expert

$$L_T \leq \frac{2}{\beta} \log K$$

If we don't know there is a best expert, but they exist anyway

$$L_T \leq 2 \log K$$

WMA

$$R_T \leq \left(\frac{1+\beta}{1-\beta} \right) \left(\min_i L_{i,T} \right) + \frac{2}{\beta} \log K$$

- ① Not guaranteed to have sublinear regrets
- ② Functional form of R_T depends on functional form of $L_{i,T}$

Claim: Possibility of "linear" regret is due to two aspects (issues)

- Algorithm is deterministic
- The setting is an "All or nothing" loss

This means that if there are these two issues, NO DETERMINISTIC ALGORITHM can guarantee sub-linear regret

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Simple Idea: The adversary knows the algorithm, and hence knowing the history, the adversary can choose rewards in a way that we get bad rewards through determinization

Example

Consider any deterministic algorithm

- ① Pick expert 1 \rightarrow always 1
- ② Pick expert 2 \rightarrow always 0
- ③ Algorithm chooses whether to choose expert 1 or 2
- ④ Adversary set $X_t \neq A_t$
- ⑤ Loss = T
- ⑥ $R_T \geq \frac{T}{2}$ (More than half are 0s or more than half are 1s)

Formulation 2 : Predicting a "REAL" sequence with expert advice in an adversarial setting

Protocol :

- ① Receive advice $(Y_{i,t} ; 1 \leq i \leq k)$ such that $Y_{i,t} \in [0,1]$
- ② Play $A_t \in [0,1]$
- ③ X_t revealed

Loss function $l(x,y)$
 $l(\cdot, \cdot) : [0,1]^2 \rightarrow [0,1]$
 $l(x, \cdot)$ is convex (this assumption stops loss to extremes)

Goal : minimize the total loss $L_T = \sum_1^T l(X_t, A_t)$

Note : It is still possible that A_t & X_t are binary

Exponential WMA

- $w_{i,1} = 1$ (initial weights)
- At time $t \geq 1$, $A_t = \frac{\sum w_{i,t} Y_{i,t}}{\sum w_{i,t}}$ \rightarrow weighted average of the experts' predictions
- $w_{i,t+1} = w_{i,t} e^{-\beta l(X_t, Y_{i,t})}$
 (can be visualized as \swarrow the loss which expert i would've incurred)

Analysis of EWMA

$$W_t = \sum_{i=1}^k w_{i,t}$$

$$W_{T+1} \geq e^{-\beta L_{i,T}} \rightarrow \text{total loss of expert } i \text{ upto time } T \dots \text{(Lower bound)}$$

$$\frac{W_{t+1}}{W_t} = \frac{\sum \tilde{w}_{i,t} e^{-\beta l_{i,t}}}{\sum \tilde{w}_{i,t}} \quad \begin{array}{l} l_{i,t} \rightarrow \text{loss of expert } i \text{ AT time } t \\ \rightarrow \text{looks like a pmf} \end{array}$$

Jensen: $f(\mathbb{E}(X)) \leq \mathbb{E}(f(X))$
 f is convex

Define Z (r.v.) s.t.

$$Z = Y_{i,t} \text{ w.p. } \frac{w_{i,t}}{\sum w_{i,t}}$$

Note: $A_t = \mathbb{E}[Z]$

RHS then is $\mathbb{E}[e^{-\beta \ell(Z, X_t)}]$? ... actually $\mathbb{E}[e^{-\beta \ell(X_t, Z)}]$

Hoeffding Lemma

$$X \sim [a, b]$$

$$\mathbb{E}[e^{s(X - \mathbb{E}(X))}] \leq e^{\frac{s^2(b-a)^2}{8}}$$

Centred MGF of X

① replace s by β

② $\ell(Z, X_t)$ is an r.v. $\in [0, 1]$

we have, $\mathbb{E}[e^{-\beta \ell(X_t, Z)}] \leq e^{-\beta \mathbb{E}(\ell(X_t, Z))} e^{\beta^2/8}$

Further using Jensen's inequality,

$$\begin{aligned} \frac{W_{t+1}}{W_t} &\leq e^{-\beta \ell(X_t, \mathbb{E}(Z))} e^{\beta^2/8} \\ &= e^{-\beta \ell(X_t, A_t)} e^{\beta^2/8} \end{aligned}$$

$$\therefore \frac{W_{T+1}}{W_1} = \frac{W_{T+1}}{K} \leq \left(e^{-\beta \ell(X_1, A_1)} e^{\beta^2/8} \right) \dots \left(e^{-\beta \ell(X_t, A_t)} e^{\beta^2/8} \right)$$

$$\therefore \frac{W_{T+1}}{K} \leq e^{-\beta L_T} e^{T\beta^2/8}$$

$$\boxed{\therefore W_{T+1} \leq K e^{-\beta L_T} e^{T\beta^2/8}}$$

Combining bounds,

$$K e^{-\beta L_T} e^{T\beta^2/8} \geq e^{-\beta L_{i,T}}$$

Note that the exponents for L_T & $L_{i,T}$ match, which is good

$$\therefore L_T \leq L_{i,T} + \frac{\beta T}{8} + \frac{\log K}{\beta}$$

$$\therefore L_T \leq \min_i L_{i,T} + \frac{\beta T}{8} + \frac{\log K}{\beta}$$

This is still not sublinear regret, but we can CHOOSE β !!

Now, $\beta \in (0, \infty)$, so use $\beta \propto \sqrt{\frac{1}{T}}$!!

requires the knowledge of the horizon T

$$\frac{\beta T}{8} \propto \sqrt{T} \text{ and } \frac{\log K}{\beta} \propto \sqrt{T}$$

choosing optimal β , ($\beta = \sqrt{\frac{8 \log K}{T}}$)

$$R_T \leq \sqrt{\frac{T \log K}{2}}$$

HW

Can you use an adaptive learning rate $\beta = \sqrt{\frac{8 \log K}{t}}$?

Randomization :

- No assumption is made on the loss function (general case)
- Play $\rightarrow Y_{i,t}$ w.p. $\frac{w_{i,t}}{\sum w_{i,t}}$

Now, the actions & the losses are both random variables

In next class: Bound $\mathbb{E}[L_T]$ from above