- Threads and Multithreading
   Each process has an address space and a single thread of control.
- Multithreading allows a process to have other threads of
- Fork is used to create a new (child) process; problems with fork:
  - 1. fork is too expensive mainly because memory has to be copied from parent to the child process (e.g. duplicating all descriptors)
- 2. Since each process has its own address space, IPC is required to pass information between parent and child after the fork.

#### Reference:

 $\underline{https://computing.llnl.gov/tutorials/pthreads/}$ 

Threads solve both these problems:

- Thread creation can be 10-100 times faster than process creation.
- All threads of a process share the same global memory, making information sharing easy. Synchronization is required.

## per **process** items:

- address space global variables
- open files,
- current working directory
- timers
- signals
- semaphores
- accounting information user and group Ids

# per thread items:

- thread ID
- program counterstack (for local variables and return addresses)
- state
- register setchild threads
- priority
- errno

Q: what do threads of same process have in common (share)?

## **Basic Posix Thread Functions**

- In real-time Posix each process can contain several "threads" of execution.
- When a program is started by exec, a single thread is created, called the *initial thread* or the *main thread*.

Additional threads are created by pthread\_create function which takes 4 pointer

( similar to fork in process )
#include <pthread.h>

int pthread\_create ( pthread\_t \*tid, const pthread\_attr\_t \*attr, void \* (\*func) (void \*), void \*arg);

#### where

- tid is the thread ID (returned by the call)
   attr is a set of attributes. Each thread has several attributes: its initial stack size, its priority, whether it should be a daemon thread or not, and so on.

We can specify these attributes by initializing a  $pthread\_attr\_t$  variable that overrides the defaults.

Normally the default values are taken, by specifying the attr argument as null pointer.

- func is the address of thread start function. The thread starts by calling this function and
  then terminates either explicitly (by calling pthread\_exit) or implicitly (by letting the function
  return).
- arg points to a set of parameters to be passed to the function when it is called.

The return value from pthread functions is either 0 if OK or a positive number (Exxx) on an error. (for example, if pthread\_create() cannot create a new thread because system limitations (on number of threads) exceeded, then return value is EAGAIN)

A thread can be aborted by use of pthread\_cancel.

#### Waiting for a thread

- One thread can wait for another thread to terminate by calling the pthread\_join function: (similar to waitpid() for process)

#include <pthread.h>

int pthread\_join (pthread\_t tid, void \*\*status);

#include <pthread.h>

int pthread\_join (pthread\_t tid, void \*\*status);

- tid is thread id of the thread we want to wait on
- \*status, if not null, takes the return value from the thread as a pointer to some object (i.e. return value is stored in the location on which status is

pointing)

#### **Detach**

The term **detaching** means cleaning up after a thread and reclaiming its storage space.

Two ways to achieve this:

- by calling the *pthread\_join* function and waiting for it to terminate, or
- by setting the detach attribute of the thread (either at creation time or dynamically by calling the pthread\_detach function:

(a thread is either joinable or detached)

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#include <pthread.h>

int pthread\_detach (pthread\_t tid);

If the detached attribute is set , the thread is not joinable and its storage space may be reclaimed automatically when the thread terminates.

When a program terminates all its threads terminate. Hence, sometimes it is necessary for a main program to issue a *pthread join*.

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# Other thread functions

• A thread can determine its id by calling *pthread self*:

#include <pthread.h>

pthread\_t pthread\_self ( void);

(similar to getpid() in process)

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• A thread can detach itself by calling

pthread\_detach(pthread\_self());

· One way for a thread to terminate is to call pthread exit:

#include <pthread.h>
void pthread\_exit(void \*status);

If the thread is not detached, its thread ID and exit status are retained for a later *pthread\_join* by some other thread.

- Forking processes (programs) which contain multiple threads is not straightforward, because some threads in the process may hold resources or executing system calls.
- The **POSIX** standard specifies that the child process will only have one thread.

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# Posix mutexes and condition variables

- The thread extension to Posix provides mutexes and condition variables.
- They can be used to synchronize threads of the same process.

They can also be used to synchronize threads of multiple processes, if the mutex or condition variable is stored in memory that is shared between processes.

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• *Mutexes* and condition variables, when combined, provide the functionality of a monitor but with a procedure interface.

Each monitor has an associated *mutex* variable and all operations on the monitor are surrounded by calls to lock and unlock the *mutex*.

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# **Mutexes: Locking and Unlocking**

- Mutexes are used to provide mutually exclusive access to critical sections.
  - (mutex status is locked or unlocked)
- · Posix mutexs are variables declared to be of type

# pthread\_mutex\_t

 Statically allocated mutexes can be initialized to the constant PTHREAD\_MUTEX\_INITIALIZER:

```
    Dynamically allocated mutexes can be initialized at run time by
calling the <a href="pthread_mutex_init">pthread_mutex_init</a> function.
```

· Lock and unlock functions:

```
#include <pthread.h>
```

```
int pthread_mutex_lock(pthread_mutex_t *mptr);
```

\* If the mutex is locked by some other thread, then we are blocked until the mutex is unlocked.

```
int pthread_mutex_trylock (pthread_mutex_t *mptr);
```

\* \_trylock returns EBUSY if the mutex is already locked.

int pthread\_mutex\_unlock (pthread\_mutex\_t \*mptr);

#### Example:

The Multiple-Producers/Single-Consumer Problem:

- Producers and consumer are represented by threads,
- buff is an integer array to hold the shared data,
- producers set buff[i] to i,
- consumer verifies that each entry is correct,
- consumer starts after all producers terminate.

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```
Example: mutex-producer-consumer
#include
               "unpipc.h"
#define
              MAXNITEMS
                                     1000
#define
              MAXNTHREADS
                                     100
/* globals shared by threads */
             nitems;
                                    /* read-only by
producer and consumer */
struct {
  pthread_mutex_t
                             mutex;
              buff[MAXNITEMS];
  int
int     nput; /* next index to store */
int     nval; /* next value to store */
} shared = { PTHREAD_MUTEX_INITIALIZER };
/* end globals */
```

```
void *produce(void *), *consume(void *);

/* include main */
int main(int argc, char **argv)
{
  int     i, nthreads, count[MAXNTHREADS];
  pthread_t tid_produce[MAXNTHREADS], tid_consume;

if (argc != 3)
  err_quit("usage: prodcons <#items> <#threads>");
  nitems = min(atoi(argv[1]), MAXNITEMS);
  nthreads = min(atoi(argv[2]), MAXNTHREADS);
```

#### The produce and consume functions are

```
void * produce(void *arg)
{
    for ( ; ; ) {
        Pthread_mutex_lock(&shared.mutex);
        if (shared.nput >= nitems) {
            /*** buffer is full ***/
            pthread_mutex_unlock(&shared.mutex);
            return(NULL);/* we're done */
            } /** if **/
            /*** buffer is not full ***/
            shared.buff[shared.nput] = shared.nval;
            shared.nput++;
            shared.nval++;
            pthread_mutex_unlock(&shared.mutex);
            *((int *) arg) += 1;
            } /** for **/
            **
```

## **Condition Variables:** Waiting and Signaling

- A mutex is for locking and a condition variable is for waiting.
- A condition variable is a variable of type: pthread\_cond\_t
- The associated functions are:

int pthread\_cond\_signal( pthread\_cond\_t \*cptr);
(both return zero on success, nonzero on error)

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A condition variable is always associated with a mutex.

The condition variable is a signaling mechanism associated with a mutex.

When a thread waits on a condition variable, its lock on the associated mutex is released.

When the thread successfully returns from the wait, it again holds the lock.

However, because more than one thread may be released, the programmer must again test for the condition which caused it to wait initially.

General structure of the code that waits on a condition variable:

```
struct {
    pthread_mutex_t mutex;
    pthread_cond_t cond;
    /* variables for maintaining the condition */
    } varx = {
        PTHREAD_MUTEX_INITIALIZER,
        PTHREAD_COND_INITIALIZER,....
    };

pthread_mutex_lock(&varx.mutex);
    while (condition is false)
    pthread_mutex_wait(&varx.cond, &varx.mutex);
    ... modify condition ...
    pthread_mutex_unlock(&varx.mutex);
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```

General structure of the code that signals a condition variable:

```
pthread_mutex_lock(&var.mutex);

set condition true

pthread_cond_signal(&var.cond);

pthread_mutex_unlock(&var.mutex);
```

The producers/consumer example where the consumer thread starts right after all producer threads have <u>started</u>:

The consumer must synchronize with the producers.

A non-busy-waiting solution.

```
Example: producer consumer
#include
               "unpipc.h"
               MAXNITEMS
                                      1000000
#define
               MAXNTHREADS
                                     100
/* globals shared by threads */
               nitems;     /* read-only by producer and consumer */
buff[MAXNITEMS];
int
int
int
               Nsignals;
struct {
  pthread_mutex_t
int nput; /* next index to store */
int nval; /* next value to store */
} put = { PTHREAD_MUTEX_INITIALIZER };
  /** struct put is used by producer only ***/
 struct{
   pthread_mutex_t
   pthread_cond_t
                         cond:
                        nready
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```