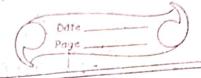
Linked List -> Definition: - A Linked List is collection of object that is stored in list form -> Nodes :- Elements present in linked list called as - contain 2 things 1 data 2 Pointen to the next mode - Linked List is linear data structure in which element are not stored in contiguous memory locations HEAD -> (data mext); -> (data mext) -> (data mext) -> NUL Node - Data holds actual value associated with node. -) Next pointen stories memory address of meat node - LL accessed through head node, which points to the first node. The last mode in the list points to NULL or nollpto indicating end of list. This mode is called as twi mode -> Why II Data structure? Allows dynamic insertion and deletion of elements during oun time. -> Due to Ease of insertion and deletion compared. to arrays, adjustments are minimal With LI there is no requirement to shift elements, only the address meeds updating. -> LL can GROW as clements are added. LL can also SHRINK as elements are removed

-> NO WASTAGE OF MEMORY "because of fixed size

allocation.



Example In a system, maintuining sorted array Such as id[] = [1000, 1010, 1050, 2000, 2040] can such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2040] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 1010, 1050, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 1010, 2000, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 1010, 2000, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 2000, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 2000, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 2000, 2000, 2000, 2000] can Such as id[] = [1000, 20

3 types: Single Linked List

Double Linked List

Circular Linked List

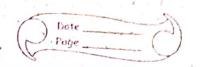
1 Singly Linked List

- -> linear duta staucture that holds value and reference to next node.
- -) This list has head and tail.
- -> We can traverse this LL in forward direction

Application.

O memory management: enabling dynamic memory

Destabase indexing: Implemented in database for swift operations like inscrtion deletion



(3) R	epresentino	Polani	mial	and	spasse	matrix
	PIESMINE	1		-	1	Was to the second

(a) In Operating Systems for tasks such as process scheduling and resource management.

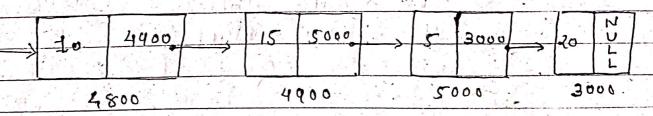
Advantages

- Dynamic memory allocation
- -> Cache friendliness
- -) Space efficient

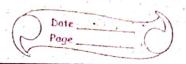
Disadvantage

- Increased memory overhead
- -> If nodes next pointer is lost, accessing elements is difficult.
- I Not suitable for parallel processing
- -) Backward traversing not possible
 Head

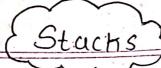
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SINKLY LINKED LIST

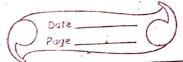


25				
2	Doubly Linked List			
->	A Doubly Linked List CDDD is a special type of			
	linked list in which each mode contains a			
•	pointer to a previous mode as well as the next node			
	of linked list.			
	Head			
×	clata)			
	Previous next			
*	Application			
- <u>- 1</u>	Used by web browsen for backward and forward			
6 E	navigation of web pages.			
	Most recently use or Least recently used cache			
	are constructed using DLL			
_	To maintain under and sedo functionalíties.			
*	Advantages			
-	traverse in both direction.			
W. 1. (-4)	injertion, and deletion of mode is easy!			
ti , ii <u>d</u> .	DDL get the previous mode so we can easily delete			
	the mode.			
504	Corner of the state of the stat			
*	Disadvantage			
. : 1921	requires extra space for previous pointer			
A second				



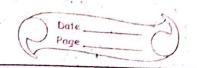
1 - is Empty ()

- is Full)



-> Stack is type of linear list where you can insert and delete elements, but only at one end the man of the same of the -> Cuhen you want to add element to stack it is calledias invest when you want to remove element from strick, it is called POP? -> In a stack, the most accessible element called as top and the least accessible element Called as botitom. - Due to work with one end of the stack, the elements are removed in opp order from how they were added. This behavior known -> Whenever you add a new element to stack.

you increase TOP pointer by "one". When you delete an element, the TOP pointer decrease openations on stuck 1 Paimary Stack operation - PUSH - POP @ Austiliary Stack operation -TOPC)



Application of stack in Ds.

- Recursion: Recursion is like a function that keeps calling itself to solve problem in smaller steps This stacks can remember these steps so you can goback when you are done
- Undal Redo operation stucks
- Expression evaluation: make sude that you are cuiculating moth problem in right order
- Riversing characters
- = servicing Hardware interrupts
- solving combinational problems, using backtrucking

In real life:

- CDIDVD stand
- Stack of books
 - -undolRedo intext editors
- History of web browser
- Call logs, Grooyle photos or Emails

Advantages of Stack Disadvantages - Easy implementation - Limited capacity - No random access

- Efficient memory - Fast access time

Lac Carive Function

- Supports backtracking

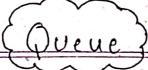
Cull limitations

-used in compiler design. - Stuck overflow

and lunderflow

As	Date Page
48-23 48-23	
*	Operations on stucks
Barry In	@ PUSH OPERATION
The regardens	
11.17.	void push Cint data) &
	1 (SIZE - top > 1)
	ary [top] = data;
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	11
	else &
Silver?	coutec" SPACE NOT AVAILABLE "<< endly
Darloy Mila	The same of the sa
	@ POP OPERATION
	Void Pop() S if (top = = -1) &
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	cout << " Stack underflow "<< endl
31 (37)	100 W 22 Stack Graces 100
	else
W	tup
ive of 2	Frank of the Lond & Party & Arthur
Alexander Control	Y he en that I have a fill of the that he had been the the the second of
Alleria Se	
Antico File	3) PEEK OPERATION
A CHARLES	ANTENNAME OF THE PARTY OF THE P
Barthay 1	int get Top()
Secretary of	if (top = = -1) {
	coutec "There is no element in stack" << endli
And the Williams	A C J C J T J J J J J J J J J J J J J J J

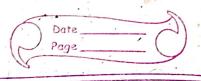
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A., J. Tanga			., .	
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- CPU tusk scheduling



->	Queue is defined as Linear Data structure
."	in which all additions of the elements made at one
	end and all deletions from the list are made at
	other lending
ب	It follows First in First Out (FIFO) rule The first
	item added is first one to be removed.
	The end from where you remove items from the
	queue is called front (sometimes head of queye')
	The end where you add new items to queue is
	called as [Rear (sometimes tail of queue)
	Adding item to queue called Engreque
	Aa, Taking an item out of queue called Doqueur
	The first the second of the se
	CHARACTERISTCS
	= hundler multiple data
T.	- we can access both ends
1	They are fast and flexible
Ł.	Application
	- multiprogramming: Simultaneous execution of programs
	- Metwork: Dato routing and communication infrastrut
	- Job scheduling: sequential execution of scheduled
4.	tasks assigned to processon
1-3	
	Real time application
	-ATM Boothline
	- Ticket Counter Line



A dvantages

- efficient data management
- FIFO operations provide simplicity.
- Multiple consumer service utilization
- swift interprocess communication

Dis advantages

17.510w middle operation

- 5- spare Limitations
- 1-10(N) reach complexity
 - Predefined muximum Jize.