Matrix Chain Multiplication -

Given a Sequence of matorices, find the most efficient way to multiply these matorices together. The Boblem is not actually to perform multiplication but merely to decide in which order we should perform the multiplication.

as the matrix multiplication is associative.

> Two matrices can be multiplied if the no of columns in first is equal to no of rows of 2nd matrix

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \end{bmatrix} \quad B = \begin{bmatrix} b_{11} & b_{12} \\ b_{21} & b_{23} \\ b_{31} & b_{32} \end{bmatrix}$$

$$2 \times 3 \quad 3 \times 2$$

$$C = \begin{bmatrix} a_{11} \times b_{11} + a_{12} \times b_{21} + a_{13} \times b_{31} & a_{11} \times b_{12} + a_{12} \times b_{22} + a_{13} \times b_{32} \\ a_{21} \times b_{11} + a_{22} \times b_{21} + a_{23} \times b_{31} & a_{21} \times b_{12} + a_{22} \times b_{22} + a_{23} \times b_{32} \end{bmatrix}$$

 $A B C 2 \times 3 \times 2 = 2 \times 2$

Total no of multiplication = 2×3×2=12

$$C[u][z] = mm \left\{ C[u,K] + C[k+1,z] + d_{i-1} \times d_k \times d_z \right\}$$

$$i \leq k \leq z$$

How many parenthesis way is possible for 4 matrices ?)

$$\frac{2n_{Cn}}{n+1}$$
 (Take $n=n-1$)

$$= 2(n-1)$$

$$-(n-1)$$

$$\frac{9e^{n=4}}{4} = \frac{6c_3}{4} = \frac{8x5x4}{4}$$

A1 X A2 X A3 X A4

3 2 3 4 4 2 9 5 5

do di di de de de da da da di vala

$$C[1,2] = Mm (C[1,T] + C[2,2] + do \times di \times da)$$

$$[\leq k \leq 2] (C[2,2] + C[3,3] + di \times da \times da)$$

$$2 \leq k \leq 3 (C[2,2] + C[3,3] + di \times da \times da)$$

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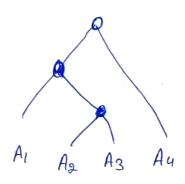
$$2 \leq k \leq 4 (C[2,2] + C[2,4] + da$$

1.e cost of multiplying these AI to Au matrices is (58) number of multiplications. How the parenthasisation should be done

A, A, A, A, A4

(A, A, A3) A4

((A1) (A2 A3))(A4)



@ for Bactice

A1 A2 A3 A 5x4 4x6 6x2 2x7

Using K Matoria

A, Az A3 A4

(A1 A2 A3) A4

((A1)(A2A3)) A4. Ang

AI AD AZ AU

Ang

1	2_	3	4	
0	120	88	158	
	0	US	104	
		0	84	
			0	
	0	1 10	0 120 88	0 120 88 158

1,				
K	1	2	3	4
1		l	1	3
2			2	3
3	1	1		3/
4	1			1

```
Time Complexity
                  Number of elements generated = n(n-1)
        To fund each element, we are calculating
        all and finding the min
                           n^2 \times n
                            =n^3
                            \Rightarrow O(n^3)
                                       of multiplication
            main ()
            € int n=5,.
              int P= [5, 4, 6, 2, 73].
             int m[5][5]=503,2
             int s[s][s] = 203,
             Int toming 2,0
             for (int d=1 g. d<n-1g. d++)
               For (int 1=1, 1<n-d, 1++)
                    for (mt k=1, . K<=1-1, 0 K++)
```

for (Int d=1, d<n-1, d++)

E

For (Int d=1, d<n-d, d++)

E

J=d+d, o

min=32767

For (Int k=1, K<=J-l, k++)

E

2=m[u][k]+m[k+1][J]+P[u-1] * P[k] * P[J], o

3f(2<min)

E

min=2

S[u][J]=k,

3

cart (m[J][n-1], o