Approximation Algorithm

An approximation algorithm is the way of dealing with NP-completeness for optimization problem.

The goal of approximation algorithm is to give solution closer to optimal solution in polynomial time.

C) cost of solution given by Approximation Sh.

C' > cost of optimal solution

S(1) > Approximation Ration

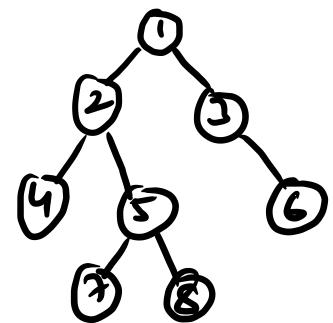
n > input size

- \mathbb{O} Maximization: $\frac{C}{C} \leqslant \mathcal{S}(n)$
- 2) Minimization: C & p(n): S(n) XI

Vertre Cover Problem

A verten cover of a graph G is a subset of verties which cover every elges.

So, verten cover problem is the orinimum Size verten cover.

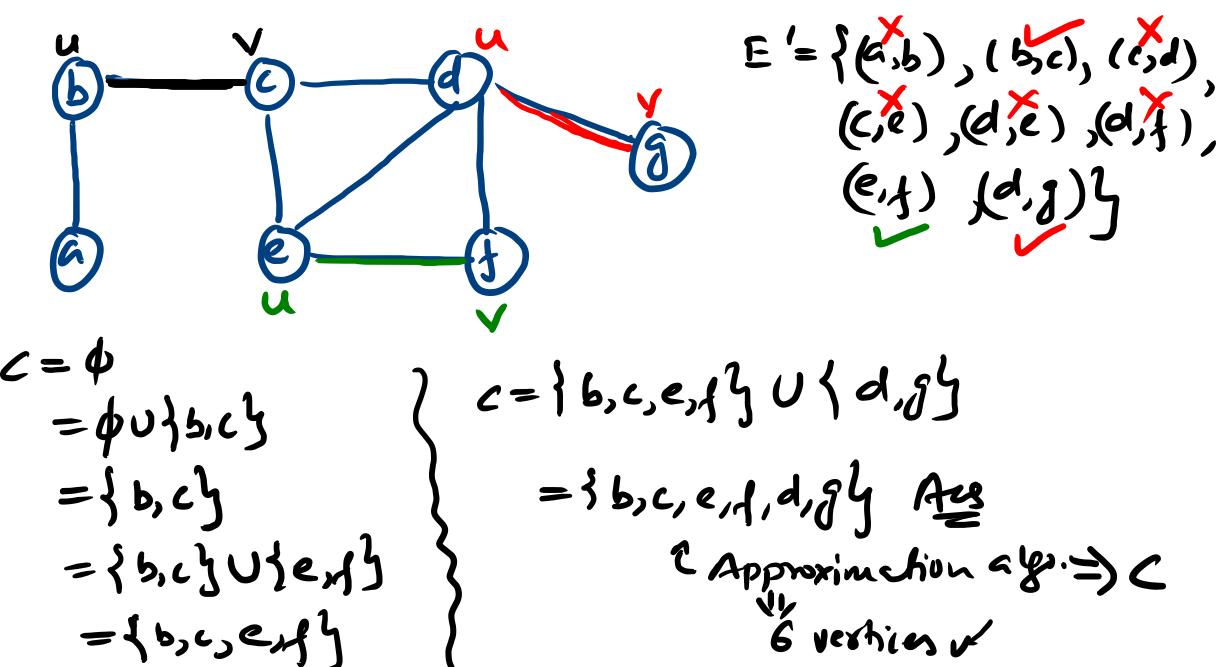


Vertex cover =
$$\{1,3,2,5\}$$

= $\{5,6,2,1\}$
= $\{2,3,5\}$ = $\frac{1}{2}$

approximation - Vertex Cover (6) J. C← Φ 2. E = E(G) 3. Wile E # \$ do let (u,v) an arbitrary edge in E' ce c U { u, v} 4. Remove from E' every ege incident either on

bretun c.



but the optimal Solution for above graph Gis

$$\frac{1}{C} = \frac{C}{C} = \frac{1}{2} = \frac{1}$$

Randomized Algorithm An algorithm that uses random number to decide what to do neut anywhere in its Logic called the Randomized Algorithm. to reduce space and time - this algorithm is used complex?ty Pandom

1 Las Vegas Algo - output is always Correct. e.g. Randomized Quick Sort ej Algosearch Repeat (A) ju (i=0; i<n-1; i++) すかく(デールーン、うてからりナナ) Y (A[i]==A(j)) 474 return True Alp Search Report-LV(A) 1 Wile (True) do i= randomi) med n+1 J= random() mod n+1 4 ((i+j) aw (a(i) = = a(j)) y return True

2) Monte Carlo Algo - output may be incomet e.g. Randamized median e.j. Algo-Jearch (A,K) fr (i=0; i≤n-1; i++) '4 (A[i)==K) 4 5 return True

Algo Monte (etho-Search (A, K, N))

I for (i=0', $i \in \mathbb{R}$; i+1) j=rondom() and n+1 j=rondom() and n+1 j=rondom() and j=rondom()