

Assignment - 1

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2) For every $b > 1$ and every constant n_0 ,

$$\log_b^n = O(n^x)$$

Solution: Let us assume this to be true.

$$\therefore f(n) = O(g(n)) \quad n^{(n)} \propto n \cdot (n)^p$$

$$O(g(n)) = \{f(n) : \text{for positive constant } c,$$

$$(f(n) \leq cg(n) \text{ at } n \geq n_0)\}$$

~~$\frac{\log_b^n}{g(n)} \rightarrow 0$~~

And,

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \quad f(n) = \log_b^n$$

$$g(n) = n^x$$

$$\lim_{n \rightarrow \infty} \frac{\log_b^n}{n^x} \quad + b > 1$$

$$x > 1$$

$$f'(n) = \frac{1}{n(\log b)} \quad g'(n) = xn^{x-1}$$

$$\lim_{n \rightarrow \infty} \frac{\left(\frac{1}{n(\log b)}\right)}{xn^{x-1}} \rightarrow \left(\frac{1}{n(\log b)}\right) \lim_{n \rightarrow \infty} \left(\frac{1}{n^{x-1}}\right) \left(\frac{1}{n}\right)$$

$$3 \left(\frac{1}{n \log b} \right) \lim_{n \rightarrow \infty} \frac{\log n}{n^{1-\epsilon}} = \frac{\text{constant}}{\log n}$$

$$\begin{aligned} & \text{if } \log b > 1 \quad (n) \rightarrow \infty \Rightarrow \frac{1}{\log b} \rightarrow 0 \\ & \text{and if } \log b < 1 \quad (n) \rightarrow \infty \end{aligned}$$

$$\text{Hence } \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \quad \text{if } \log b < 1.$$

i. Our assumption is true.

\Rightarrow

3) For every $\epsilon > 0$ and every constant $d > 0$

$$n^d = \Omega(\gamma^n)$$

$$(n)^d > n^d = \frac{(n)^d}{n^d} n^d$$

Solution - Let us assume it is true.

$$f(n) = \Omega(g(n))$$

$$(n^d) \geq n^d \quad \text{and} \quad (n^d) \leq n^d$$

$\Omega(g(n)) = \{ f(n) : \text{there exists positive constant } c \in \mathbb{N}^0$
such that

$$0 \leq cg(n) \leq f(n) \forall n \geq n_0 \}$$

$$\text{And } \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$$

$$0 = \frac{(n^d)}{(n^d)} \text{ will vanish}$$

$$(1) \lim_{n \rightarrow \infty} f(n) = \infty \Rightarrow g(n) = \infty$$

then we can apply L'Hopital's Rule

$$\lim_{n \rightarrow \infty} \frac{m}{\sqrt[n]{e^n}} \Rightarrow \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)}$$

[Applying L'Hopital's Rule]

$$f'(n) = d n^{d-1} = \frac{(n)^d - n^0}{(n)^0}$$

$$g'(n) = r^n (\ln r)$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{d n^{d-1}}{r^n (\ln r)}$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)} = \lim_{n \rightarrow \infty} \frac{f''(n)}{g''(n)},$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{d^{d-1}}{r^n (\ln r)} = \lim_{n \rightarrow \infty} \frac{d(d-1)n^{d-2}}{r^n (\ln r)^2}$$

on left side we are going to take limit of $\frac{d^{d-1}}{r^n (\ln r)^2}$ as $n \rightarrow \infty$

We can observe that as we derivative
for $n \rightarrow \infty$ the terms $(\ln r)^2$ increases -
the numerator decreases and the denominator
increases -

$$\text{Hence } \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$$

But, this is a contradiction to our assumption

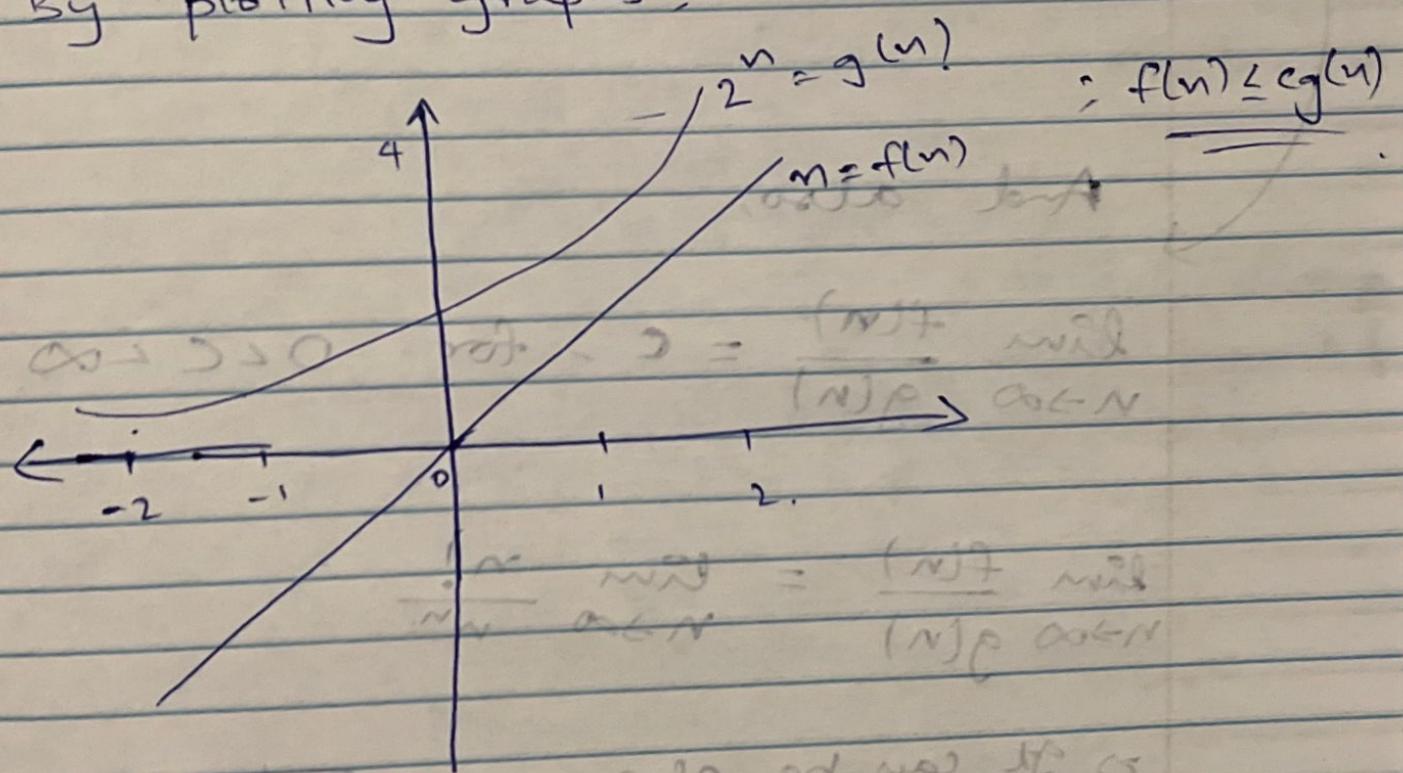
Hence, it is false.

$$n^d = \Omega(r^n) \quad r > 1 \quad \therefore d > 0$$

$$n^d = \Omega(2^n)$$

$$r=2; \quad d=1$$

By plotting graphs; for better understanding



6) $n! = \Theta(n^n)$

Solution:- Let us assume it is true.

$$\text{For } f(n) = \Theta(g(n))$$

$\Theta(g(n)) = \{f(n) : \text{there exists positive const} - c_1, c_2 \in \mathbb{R} \text{ such that}$

$$c_1 g(n) \leq f(n) \leq c_2 g(n)$$

$$f(n) = n!$$

$$g(n) = n^n$$

$$= n \times (n-1) \times \dots \times 1 \quad \Rightarrow \quad n \times n \times \dots \times n$$

And also,

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = c > 0 \quad \text{for } 0 < c < \infty$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{n!}{n^n}$$

\Rightarrow It can be also represented as

$$\lim_{n \rightarrow \infty} \left[\frac{n(n-1)(n-2)\dots 1}{n \times n \times \dots \times n} \right]^{(n)}$$

$$\Rightarrow \lim_{n \rightarrow \infty} \left[\left(1 - \frac{1}{n}\right) \left(1 - \frac{2}{n}\right) \dots \left(1 - \frac{1}{n}\right) \right]$$

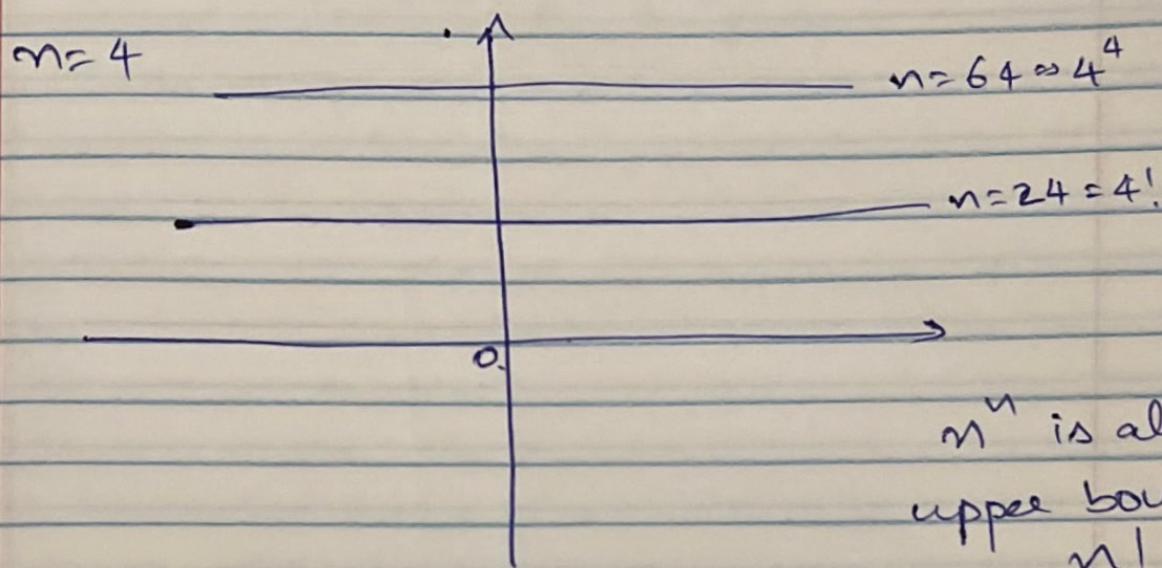
$$\lim_{n \rightarrow \infty} \left(1 - \frac{1}{n}\right) = 1$$

β should be greater than 0.

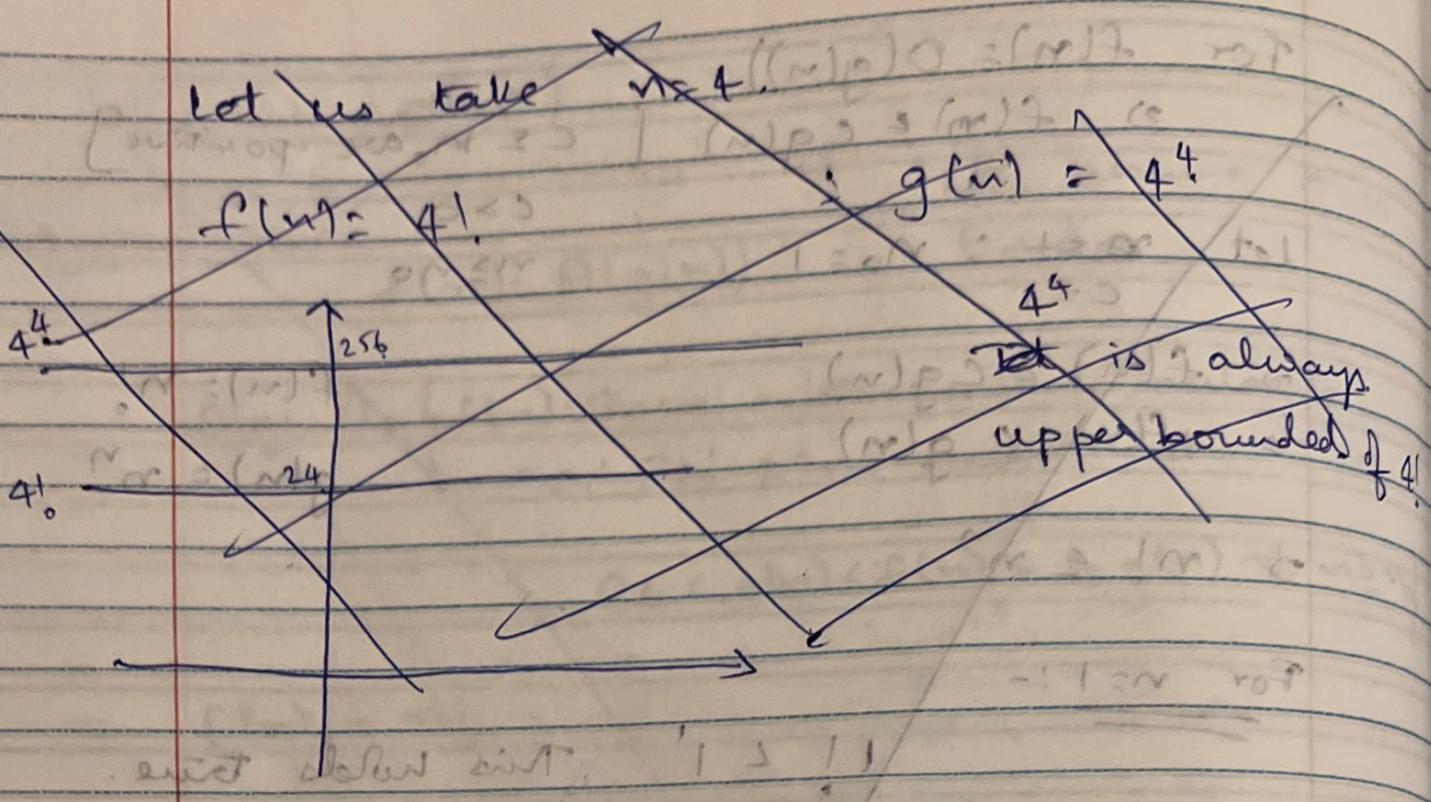
Hence it is a contradiction to our assumption.

Hence, it is false.

for example,



n^n is always
upper bound of
 $n!$.



7). $2^n = \Theta(4^n)$

Solution:- $f(n) = \Theta(g(n))$

lets assume it is true.

$\Theta(g(n)) = \{ f(n) : \text{There exist positive constants } c_1, c_2 \text{ & } n_0 \text{ such that}$

$$0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \quad \forall n \geq n_0$$

Here $f(n) = 2^n \therefore g(n) = 4^n = 2^{2n}$

And $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \Theta(1) ; 0 < C < \infty$

$$\lim_{n \rightarrow \infty} \frac{2^n}{2^{2n}} \Rightarrow \lim_{n \rightarrow \infty} 2^{n-2n} = \lim_{n \rightarrow \infty} 2^{-n}$$

\Rightarrow Hence. $\lim_{n \rightarrow \infty} 2^{-n} = 0$

Again, ' c ' has to be greater than ' d '.

Hence, $f(n) \neq \Theta(g(n))$. Hence, our assumption is false.

But by the theorem,

As $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \Rightarrow f(n) = O(g(n)) \neq f(n) \neq \Omega(g(n))$

$\Rightarrow 2^n \neq \Theta(4^n)$
 $2^n = O(4^n)$
 $2^n \neq \Omega(4^n)$

For values of n :

$$\Rightarrow 2^n \leq c \cdot 4^n$$

$$\Rightarrow 2^n \leq c \cdot 2^{2n} \quad \text{Holds true for all values of } n$$

$$\Rightarrow 1 \leq c \cdot 2^n$$

$$5) n^{1.01} = O(n \log^2 n)$$

Let's assume it is true

Solution: $f(n) = O(g(n))$

$$O(g(n)) = \{ f(n); \text{ for some positive constants } c, \\ f(n) \leq cg(n) \forall n \geq n_0 \}$$

And .

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0.$$

$$f(n) = n^{1.01} \quad \left(= n \cdot \frac{(n)^{0.01}}{(n)^{0.01}} \text{ as } n \rightarrow \infty\right)$$

$$g(n) = n \log^2 n$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{n^{1.01}}{n \log^2 n} \Rightarrow \text{Applying L'Hospital Rule}$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)}$$

$$f(n) = n^{1.01}$$

$$f'(n) = (1.01)(n^{1.01-1}) \Rightarrow \underline{\underline{(1.01)(n^{0.01})}}$$

$$g(n) = n \log^2 n$$

$$g'(n) = \frac{d}{dn} (n \log^2 n)$$

Now, using product rule,

$$\frac{d}{dn}(uv) = u \frac{dv}{dn} + v \frac{du}{dn}$$

Here, $u = n$

$$\therefore \frac{d}{dn}(n \log^2 n) = n \cdot \frac{d}{dn}(\log^2 n) + \log^2 n \left(\frac{dn}{dn} \right)$$

$$\Rightarrow (n \times 2 \log n \times 1) + (\log^2 n \times 1)$$

$$\Rightarrow \underline{\underline{2 \log n [1 + \log n]}} \quad \underline{\underline{\log n [2 + \log n]}}$$

$$g'(n) = g'(n \log^2 n) = \log n (2 + \log n)$$

$$\therefore \lim_{n \rightarrow \infty} \frac{(1.01)^n}{\log n (2 + \log n)} \quad \text{Here as we}$$

observe, the

numerator grows

$$\lim_{n \rightarrow \infty} \frac{(1.01)^n}{\log n (2 + \log n)} = \infty \text{ than the denominator}$$

as $n \rightarrow \infty$.

Hence, our assumption is contradicting.

Hence,

$n^{1.01} \neq O(n \log^2 n)$ and

$$n^{1.01} = \Omega(n \log^2 n) & n^{1.01} \neq \Omega(n \log^2 n)$$

- i) Let $f(n)$ be a polynomial running time - function of n and ' g ' be an exponential running - time function of n . Then $f(n) = O(g(n))$.

Solution :- Let us assume it is true.

$f(n) =$ polynomial running - time function

$$\Rightarrow a_0 n^m + a_1 n^{m-1} + \dots + a_n$$

$g(n) =$ exponential running - time function
 $\Rightarrow b^n$

Now,

By definition of $O(g(n))$

~~O(g(n)) = { f(n) : for a positive constant 'c'~~

$$f(n) \leq c g(n) \quad \forall n \geq n_0$$

And

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0.$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)}$$

$$f'(n) = n a_0 n^{n-1} + (n-1) a_1 n^{n-2} + \dots +$$

$$((n/p))_0 = (n)_0$$

$$g'(n) = b^n \ln b$$

Hence,

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)} = \lim_{n \rightarrow \infty} \left[\frac{n a_0 n^{n-1}}{b^n \ln b} + \dots \right] = 0$$

let us consider a monomial

$$f(n) = n$$

$$g(n) = b^n$$

$$\lim_{n \rightarrow \infty} \frac{n}{b^n} = \lim_{n \rightarrow \infty} \frac{1}{b^n (\ln b)} \stackrel{\infty}{\rightarrow} 0 = (n)_0$$

for a binomial.

$$f(n) = an^2 + bn + c$$

$$g(n) = b^n \Rightarrow (n)_0 = b^n \ln b$$

$$\lim_{n \rightarrow \infty} \frac{an^2 + bn + c}{b^n} \stackrel{(n/p)^2}{\rightarrow} \lim_{n \rightarrow \infty} \frac{[2an + b]}{b^n (\ln b)} \stackrel{\infty}{\rightarrow} \lim_{n \rightarrow \infty} \frac{2a}{b^n (\ln b)^2}$$

Therefore even for a grade binomial

$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$. [As the denominator increases enormously than the numerator as $n \rightarrow \infty$]

It holds true

for all polynomials.

Hence, $f(n) = O(g(n))$

Where, $f(n)$ = polynomial function

$g(n)$ = exponential function,

Hence, it ~~is~~ base assumption is true.

8) $\sum_{i=1}^n i^k = O(n^{k+1})$; where 'k' is a constant

Solution:- Let us assume it is true.

$$f(n) = O(g(n))$$

$$O(g(n)) = \begin{cases} f(n) : \text{for positive constant } c, \\ f(n) \leq c g(n) \forall n \geq n_0 \end{cases}$$

Here, And $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$.

L-transposition

Here, $f(n) = \sum_{i=1}^n i^n$ where k is a constant.

$$g(n) = n^{k+1}$$

$$f(n) = 1^n + 2^n + \dots + n^n$$

$$g(n) = n^{k+1} \Rightarrow (n^k)n = (n^k \cdot n^k) = n^{2k}$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{(1^n + 2^n + \dots + n^n)}{n^{2k}}$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{\left[\frac{1}{n^k} + \frac{2}{n^k} + \dots + \frac{n}{n^k} \right]}{n}$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{\left[\left(\frac{1}{n}\right)^k + \left(\frac{2}{n}\right)^k + \dots + \left(\frac{n}{n}\right)^k \right]}{n}$$

$$\text{Let } \left(\frac{k}{n}\right)^n = \underset{n \rightarrow \infty}{\cancel{0}} \quad \text{as } n \rightarrow \infty$$

$$\text{Hence, } \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \underset{(d.p.t)/n}{\cancel{1}} = (n)^{1/2}$$

Hence, our assumption is true.

$$\left(\frac{1}{n}\right)^n \text{ and } \left(\frac{1}{n}\right)^n \cancel{\rightarrow \frac{(d.p.t)^n}{n^n}}$$