

## Seminar 8 - LL(1)

$\overbrace{LL(K)}$  uses leftmost derivations  
↑ predictions of length  $K$   
sequence read from left to right

LL( $K$ ) parser: at each moment of parsing an action is uniquely determined

### LL(1) Parser

→ uses predictions of length 1

Steps

- 1) build FIRST & FOLLOW
- 2) build LL(1) parse table
- 3) analyse sequence based on moves and table

#### FIRST

let  $\alpha \in (N \cup \Sigma)^*$ , then  $\text{FIRST}_K$  means first  $K$  terminal symbols that can be generated by  $\alpha$

Notation  $\text{FIRST}_1 = \text{FIRST}$

Example:

$$S \rightarrow Bb \mid Cd$$

$$B \rightarrow aB \mid \epsilon$$

$$C \rightarrow cC \mid \epsilon$$

	$F_0$	$F_1$	$F_2$
S	$\emptyset$	$a, \epsilon, c, d$	$a, \epsilon, c, d$
B	$a, \epsilon$	$a, \epsilon$	$a, \epsilon$
C	$c, \epsilon$	$c, \epsilon$	$c, \epsilon$

(Ex1) Compute FIRST for:

$$\begin{aligned} S &\rightarrow BA^1 \\ A &\rightarrow +BA^2 \\ A &\rightarrow \epsilon^3 \\ B &\rightarrow DC^4 \\ C &\rightarrow *DC^5 \\ C &\rightarrow \epsilon^6 \\ D &\rightarrow (S)^7 \\ D &\rightarrow a^8 \end{aligned}$$

	$F_0$	$F_1$	$F_2$	$F_3$
S	$\emptyset$	$\emptyset$	{, a}	{, a}
A	$+,\epsilon$	$+,\epsilon$	$+,\epsilon$	$+,\epsilon$
B	$\emptyset$	{, a}	{, a}	{, a}
C	* $,\epsilon$	* $,\epsilon$	* $,\epsilon$	* $,\epsilon$
D	{, a}	{, a}	{, a}	{, a}

$$FIRST(S) = \{ (, a) \}$$

$$FIRST(A) = \{ +, \epsilon \}$$

$$FIRST(B) = \{ (, a) \}$$

$$FIRST(C) = \{ *, \epsilon \}$$

$$FIRST(D) = \{ (, a) \}$$

### FOLLOW

→ the next symbol generated after/following A

Rules (FOLLOW(A) for example)

we look for A in the right side of all productions

• if A is at the end of production we add ( union ) FOLLOW(left-side of prod.)

FOLLOW  
Follow

• if after A is a terminal , that terminal is the

! • if after A is another non-terminal , let's say B ,  
then  $FOLLOW(A) = FIRST(B)$  and if  $FIRST(B)$  contains  $\epsilon$

then we also add FOLLOW(~~left-side~~) !

Example

$$S \rightarrow Bb \mid Cd$$

$$B \rightarrow ab \mid \epsilon$$

$$C \rightarrow cc \mid \epsilon$$

	$L_0$	$L_1$	$L_2$
S	$\epsilon$	$\epsilon$	$\epsilon$
B	$\emptyset$	b	b
C	$\emptyset$	d	d

# Algorithm FOLLOW

**INPUT:** G, FIRST(X),  $\forall X \in N \cup \Sigma$

**OUTPUT:** FOLLOW(A),  $\forall A \in N$

**for**  $A \in N - \{S\}$  **do**

{init}

$L_0(A) = \Phi;$

**endFor;**

$L_0(S) = \{\epsilon\};$

{init}

$i = 0;$

**repeat**

$i = i + 1;$

**for**  $B \in N$  **do**

**for**  $A \rightarrow \alpha B \gamma \in P$  **do**

**for**  $\forall a \in FIRST(y)$  **do**

**if**  $a = \epsilon$  **then**  $F_i(B) = F_i(B) \cup F_{i-1}(A)$

**else**  $F_i(B) = F_{i-1}(B) \cup FIRST(y)$

**endif**

**endFor**

**endFor**

**endFor**

**until**  $F_i(X) = F_{i-1}(X), \forall X \in N$

$FOLLOW(X) = F_i(X), \forall X \in N$

1. Rule for  $L_0$

$$L_0(S) = \{\epsilon\}$$

$$L_0(X) = \emptyset \text{ } \forall X \in N - \{S\}$$

Ex 2

Compute FOLLOW for the grammar from Ex 1

	$L_0$	$L_1$	$L_2$	$L_3$	$L_4$
S	$\epsilon$	$\epsilon, )$	$\epsilon, )$	$\epsilon, )$	$\epsilon, )$
A	$\emptyset$	$\epsilon$	$\epsilon, )$	$\epsilon, )$	$\epsilon, )$
B	$\emptyset$	$+, \epsilon$	$+, \epsilon, )$	$+, \epsilon, )$	$+, \epsilon, )$
C	$\emptyset$	$\emptyset$	$+, \epsilon$	$+, \epsilon, )$	$+, \epsilon, )$
D	$\emptyset$	$*, \epsilon$	$*, \epsilon, +$	$*, \epsilon, +, )$	$*, \epsilon, +, )$

$$FOLLOW(S) = \{ ) \}$$

$$FOLLOW(A) = FOLLOW(S) \cup FOLLOW(A)$$

$$FOLLOW(B) = FIRST(A) \cup FOLLOW(S) \cup FOLLOW(B)$$

because  $FIRST(A)$  contains  $\epsilon$

$$FOLLOW(C) = FOLLOW(B) \cup FOLLOW(C)$$

$$FOLLOW(D) = FIRST(C) \cup FOLLOW(B) \cup FOLLOW(C)$$

because  $FIRST(C)$  contains  $\epsilon$

## LL(1) TABLE

$M(A, \alpha) = (\alpha, i)$  if  $\alpha \in \text{FIRST}(\alpha)$ ,  $\alpha \neq \epsilon$ ,  $A \rightarrow \alpha$  production with number  $i$

$M(A, b) = (\alpha, i)$  if  $\epsilon \in \text{FIRST}(\alpha)$ ,  $b \in \text{FOLLOW}(A)$ ,  $A \rightarrow \alpha$  production with number  $i$

$M(a, a) = \text{pop}$ ,  $a \in \Sigma$  ( $\Rightarrow$  terminal + terminal = pop)

$M(\$, \$) = \text{acc}$

$M(x, a) = \text{err}$

Example: for example FIRST & example FOLLOW

	a	b	c	d	\$
S	Bb, 1	Bb, 1	Cd, 2	Cd, 2	.
B	ab, 3	ε, 4	.	.	.
C	.	.	cc, 5	ε, 6	.
a	pop	.	.	.	.
b	.	pop	.	.	.
c	.	.	pop	.	.
d	.	.	.	pop	.
\$	.	.	.	.	acc

conf  $(\alpha, \beta, \pi)$   
 input ↑ ↑ output  
 stack      work  
 stack

inst:  $(w\$, \$, \epsilon)$

end:  $(\$, \$, \pi)$

$w = aab \in L(G)$ ?

$(aab\$, \$, \epsilon) \xrightarrow{\text{push}} (aab\$, Bb\$, 1) \xrightarrow{\text{push}} (aab\$, aBb\$, 13)$   
 $\xrightarrow{\text{pop}} (ab\$, Bb\$, 13) \xrightarrow{\text{push}} (ab\$, aBb\$, 133) \xrightarrow{\text{pop}} (b\$, Bb\$, 133)$   
 $\xrightarrow{\text{pop}} (\$, \$, 133) \xrightarrow{\text{pop}} (\$, \$, 133) \checkmark$

Ex 3

Build the LL(1) table for ex 1 & 2

	+	*	(	)	a	\$
S		BA,1		BA,1		
A	+BA,2		E,3			
B		DC,4		DC,4		
C	E,6 $\Rightarrow$ DC,5		E,6			
D		(S),7		a,8		
+	pop					
*		pop				
(		pop				
)			pop			
a				pop		
\$					acc	

w = a\*(a+a) ? L(G)

$D \rightarrow a^8$

- >  $(a^*(a+a)$, S$, \varepsilon) \vdash (a^*(a+a)$, BA$, 1) \vdash (a^*(a+a)$, DCA$, 14)$   
 $\vdash (a^*(a+a)$, aCA$, 148) \vdash (*^*(a+a)$, CA$, 148) \vdash ((^*(a+a)$, *DCA$, 1485)$   
 $\vdash ((a + a)$, DCA$, 1485) \vdash ((a+a)$, (S)CA$, 14857)$   
 $\vdash (a+a)$, S)CA$, 14857) \vdash (a+a)$, BA)CA$, 148571) \vdash (a+a)$, DCA)CA$, 1485714)$   
 $\vdash (a+a)$, aCA)CA$, 14857148) \vdash (+a)$, CA)CA$, 14857148)$   
 $\vdash (+a)$, \varepsilon A)CA$, 148571486) \vdash (+a)$, +BA)CA$, 1485714862) \vdash (a)$, BA)CA$, 1485714862)$   
 $\vdash (a)$, DCA)CA$, 14857148624) \vdash (a)$, aCA)CA$, 148571486248)$   
 $\vdash ()$, CA)CA$, 148571486248) \vdash ()$, \varepsilon A)CA$, 1485714862486)$   
 $\vdash ()$, \varepsilon)CA$, 14857148624863) \vdash ($, CA$, 14857148624863)$   
 $\vdash ($, \varepsilon A$, 148571486248636) \vdash ($, \varepsilon $, 1485714862486363)$   
 $\vdash ($, $, 1485714862486363) \vdash acc$