

F-L-C SEMINAR 3 & 4

- Given grammar $G = (\{S, H\}, \{a, c, d, e\}, \{S \rightarrow b^2 S e^1 | H, H \rightarrow C H d^2 | cd^3, S\})$, find the generated language (+ proof). - 25P

$$(5) \xrightarrow{(2)} H \xrightarrow{(4)} cd$$

$$\xrightarrow{(3)} \text{CHd}^2 \xrightarrow{(4)} \text{C}^2\text{d}^{\text{B}}$$

$$(3) \quad C^3Hd^6 \xrightarrow{(4)} C^4d^7$$

$$(3) \quad C^4 + H^8 \xrightarrow{(4)} C^5 H^8$$

(3) ...

$$? \quad H \xrightarrow{*} c^n d^{2n-1}, \forall n \in N^*$$

I ? P(1) is true

$$H \stackrel{(4)}{\Rightarrow} cd = d^{d^{2 \cdot i - 1}} \Rightarrow P(i) \text{ is true}$$

$$\text{II } p(k) \xrightarrow{?} p(k+1), \forall k \in \mathbb{N}^*$$

$$P(k) \text{ is true} \Rightarrow H \stackrel{*}{\Rightarrow} C^k d^{2k-1} \Rightarrow$$

$$\Rightarrow H \stackrel{?}{=} C c^k d^{2k+1} d^2 = C^{k+1} d^{2(k+1)-1} \stackrel{(3)}{\Rightarrow} p(k+1) \text{ is true}$$

I, II $\Rightarrow P(n)$ is true, $\forall n \in N^*$ \Rightarrow

$\Rightarrow H \stackrel{?}{\Rightarrow} \tilde{o}^{2^{n-1}} d, \forall n \in N^*$

$S \Rightarrow H$

(1) $\Downarrow b^2 Se \stackrel{(2)}{\Rightarrow} b^2 He_2$

(1) $\Downarrow b^4 Se^2 \stackrel{(2)}{\Rightarrow} b^4 He^2$

(1) $\Downarrow b^6 Se^4 \stackrel{(2)}{\Rightarrow} b^6 He^4$

(1) $\Downarrow b^8 Se^6 \dots$

? ~~$S \Rightarrow b^{2^m} H e^{2^m}$~~ $\frac{b^{2^m} H e^{2^m}}{K_m \in N}$

? $S \stackrel{?}{\Rightarrow} b^{2^m} \tilde{o}^m d^{2^{m-1}} e^m, \forall m \in N, \forall a \in N^*$

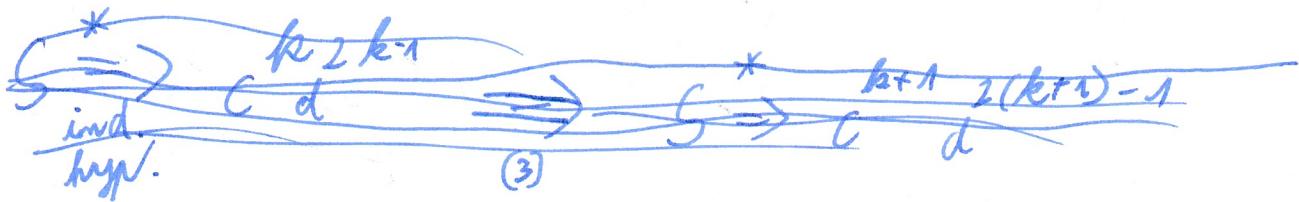
I ? $P(1, 0)$ is true $\Leftrightarrow S \stackrel{?}{\Rightarrow} cd$

$S \Rightarrow H \stackrel{(2)}{\Rightarrow} cd \Rightarrow P(1, 0)$ is true

$k > 1, l > 0$
 $k \in \mathbb{N}^*, l \in \mathbb{N}$

II. $P(k, l) \Rightarrow P(k+1, l), P(k, l+1), P(k+1, l+1)$

$P(k, l)$ is true $\Rightarrow S^* \Rightarrow b^{2l} c^k d^{2k-1} e^l$



$H \Rightarrow cd \xrightarrow{(4)} \xrightarrow{(3)} c^k d^{2(k+1)-1} \xrightarrow{(2)} S^* \Rightarrow c^{k+1} d^{2(k+1)}$

$\xrightarrow{(1)} S^* \Rightarrow b^{2l} c^{k+1} d^{2(k+1)-1} \Rightarrow P(k+1, l)$ is true

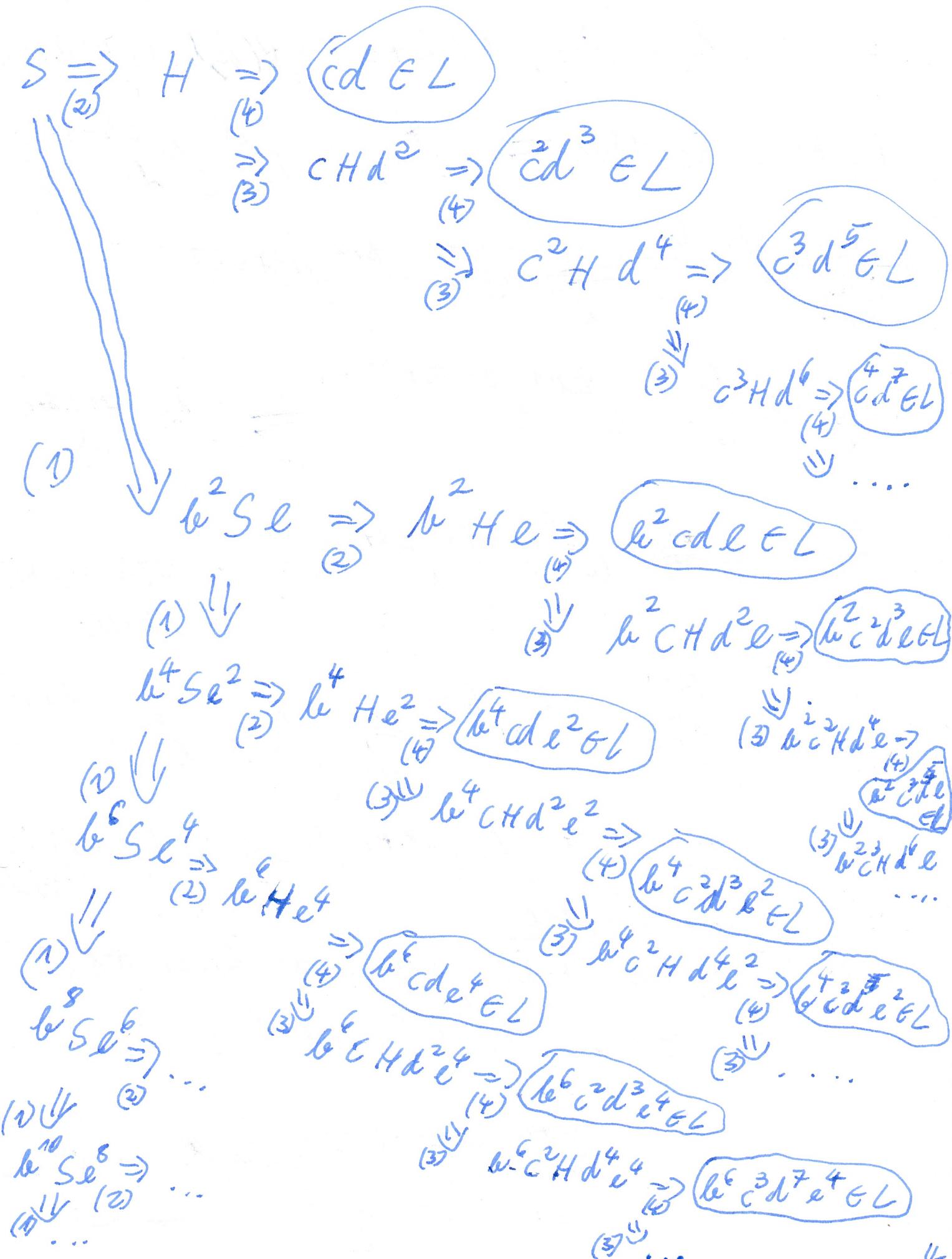
$H \Rightarrow (cd)^{k-1} \xrightarrow{(4)} \xrightarrow{(3)} c^k d^{2k-1} \xrightarrow{(2)} S^* \Rightarrow c^k d^{2k+1}$

$\xrightarrow{(1)} S^* \Rightarrow b^{2(l+1)} c^k d^{2k+1} e^l \Rightarrow P(k, l+1)$ is true
applying the

By same logic, we can deduce that
 $P(k, l) \Rightarrow P(k+1, l+1)$

$\Rightarrow \{ b^{2^m} c^m d^{2^{m-1}} e^m \mid m \in \mathbb{N}^*, m \in \mathbb{N} \} \subseteq L(G)$

? $L(G) \subseteq L$



All grammar productions have been used in all possible combinations \Rightarrow valid justif.

$$L \subseteq L(G), L(G) \subseteq L \Rightarrow L(G) = L$$

2. Find grammars that generate the following languages

A. $L_1 = \{x^n y^n \mid n \in \mathbb{N}\}$ + proof

It is enough to find G_1 s.t. $L_1 \subseteq L(G_1)$ and $L(G_1) \subseteq L_1$
 Let $G_1 = (N, \Sigma, P, S)$

$$N = \{S\}$$

$$\Sigma = \{x, y\}$$

$$P: S \xrightarrow{(1)} xSxy \mid \xrightarrow{(2)} \epsilon$$

I $L_1 \subseteq L(G_1)$

$$L_1 \subseteq L(G_1) (\Leftrightarrow) \forall n \in \mathbb{N}: x^n y^n \in L(G_1)$$

Get $n \in \mathbb{N}$ - fixed

$$S \xrightarrow{(2)} \epsilon = xy^0 \xrightarrow{(1)} xxy \xrightarrow{(1)} x^n y^n \Rightarrow$$

$$\Rightarrow S \xrightarrow{*} x^n y^n \Rightarrow x^n y^n \in L(G_1) \Rightarrow L_1 \subseteq L(G_1)$$

$$L_1 = \{x^n y^n \mid n \in \mathbb{N}\}$$

II $L(G_1) \subseteq L_1 \Leftrightarrow ?S \text{ generates } x^n y^n, n \in \mathbb{N}$

$$S \stackrel{(2)}{\Rightarrow} S = \{x^0 y^0 \in L_1\}$$

$$\stackrel{(1)}{\Rightarrow} xSyz \stackrel{(2)}{\Rightarrow} \{xy \in L_1\}$$

$$\stackrel{(1)}{\Downarrow} x^2Syz^2 \stackrel{(2)}{\Rightarrow} \{x^2y^2 \in L_1\}$$

$$\stackrel{(1)}{\Downarrow} x^3Syz^3 \stackrel{(2)}{\Rightarrow} \{x^3y^3 \in L_1\}$$

$\Downarrow \dots$

$\Rightarrow S, \text{ only generates } x^n y^n, \forall n \in \mathbb{N}$

$$\Rightarrow L(G_1) \subseteq L_1$$

$$I, II \Rightarrow L(G_1) = L_1$$

B. $L_2 = \{a^n b^{2n} \mid n \in \mathbb{N}^*\} + \text{proof}$

$$\text{Get } G_2 = (N, \Sigma, P, S)$$

$$N = \{S\}$$

$$\Sigma = \{a, b\}$$

$$P: S \xrightarrow{(1)} a^5 b^2 / a^2 b^2$$

I $L_2 \subseteq L(G_2) \Leftrightarrow \forall n \in \mathbb{N} \exists: a^n b^{2n} \in L(G_2)$

Get $n \in \mathbb{N}^*$ -fixed

$$S \stackrel{(2)}{\Rightarrow} ab^2 \stackrel{(1)}{\Rightarrow} a^2 b^4 \stackrel{n-2}{\Rightarrow} a^n b^{2n}$$

$$S \stackrel{*}{\Rightarrow} a^n b^{2n} \Rightarrow a^n b^{2n} \in L(G_2) \Rightarrow L_2 \subseteq L(G_2)$$

II $L(G_2) \stackrel{?}{\subseteq} L_2$

$$S \xrightarrow{(2)} [a^2 b^2 \epsilon L_2]$$

$$\xrightarrow{(1)} a S b^2 \xrightarrow{(2)} [a^2 b^4 \epsilon L_2]$$

$$\xrightarrow{(1)} a^2 S b^4 \xrightarrow{(2)} [a^3 b^6 \epsilon L_2]$$

$$\xrightarrow{(1)} a^3 S b^6 \xrightarrow{(2)} [a^4 b^8 \epsilon L_2]$$

$\Rightarrow S$ only generates $a^n b^{2n}, n \in N^*$...

$$\Rightarrow L(G_2) \subseteq L_2$$

$$I, II \Rightarrow L(G_2) = L_2$$

C. $L_3 = \{a^n b^m \mid n, m \in N^*\}$ - regular grammar required + proof

Let $G_3 = (N, \Sigma, P, S)$

$$N = \{S, H\}$$

$$\Sigma = \{a, b\}$$

$$P: \begin{aligned} S &\xrightarrow{(1)} a S \quad | \quad a H \\ &\xrightarrow{(2)} b H \quad | \quad b \end{aligned} \quad \text{right-linear}$$

I $L_3 \stackrel{?}{\subseteq} L(G_3) \Leftrightarrow \forall n, m \in N^* : a^n b^m \in L(G_3)$

Let $n, m \in N^*$ fixed

$$H \xrightarrow{(4)} b \xrightarrow{(m-1)} b^m \quad (i)$$

$$S \xrightarrow{(2)} aH \xrightarrow{(1)} a^n H \xrightarrow{(2)} a^n b^m \Rightarrow$$

$$\Rightarrow S \xrightarrow{*} a^n b^m \Rightarrow a^n b^m \in L(G_3)$$

$$\Rightarrow L_3 \subseteq L(G_3)$$

II $L(G_3) \subseteq L_3 \Leftrightarrow ? S \text{ only generates } a^n b^m, n, m \in \mathbb{N}^*$

a) ? H only generates $b^m, m \in \mathbb{N}^*$

$$H \xrightarrow{(4)} b$$

$$\xrightarrow{(3)} bH \xrightarrow{(4)} b^2$$

$$\xrightarrow{(3)} b^2 H \xrightarrow{(4)} b^3$$

$$\xrightarrow{(3)} b^3 H \xrightarrow{(4)} b^4$$

$$\dots$$

$\Rightarrow H \text{ only generates } b^m, m \in \mathbb{N}^*$

b) ? S only generates $a^n H, n \in \mathbb{N}^*$

$$S \xrightarrow{(2)} aH$$

$$\xrightarrow{(1)} aS \xrightarrow{(2)} a^2 H$$

$$\xrightarrow{(1)} a^2 S \xrightarrow{(2)} a^3 H$$

$$\xrightarrow{(1)} a^3 S \xrightarrow{(2)} a^4 H$$

$$\dots$$

$\Rightarrow S \text{ only generates } a^n H, n \in \mathbb{N}^*$

a), b), (2) \Rightarrow G_3 only generates $a^n b^n$, $n \in \mathbb{N}$

(Considering that the other prod. rule of S besides (2) is (1), and that applying (1) still generates only results in L_3 , then this should be sound.) I realized late that I could have used $N = \{S, t, BY\}$, $S = \{a, b\}$, $P = \{S \rightarrow AB\}$, (not regular)

$$\Rightarrow L(G_3) \subseteq L_3$$

$$I, II \Rightarrow L(G_3) = L_3$$