

## EXP 5

### FIRST AND FOLLOW

**AIM:** To write a program for displaying the first and follow of a given production

#### **ALGORITHM:**

For computing the first:

1. If  $X$  is a terminal then  $FIRST(X) = \{X\}$

Example:  $F \rightarrow (E) \mid id$

We can write it as  $FIRST(F) \rightarrow \{ (, id \}$

2. If  $X$  is a non terminal like  $E \rightarrow T$  then to get

$FIRST(E)$  substitute  $T$  with other productions until you get a terminal as the first symbol

3. If  $X \rightarrow \epsilon$  then add  $\epsilon$  to  $FIRST(X)$ .

For computing the follow:

1. Always check the right side of the productions for a non-terminal, whose FOLLOW set is being found. ( never see the left side ).

2. (a) If that non-terminal ( $S, A, B, \dots$ ) is followed by any terminal ( $a, b, \dots, *, +, (, ) \dots$ ) , then add that “terminal” into FOLLOW set.

- (b) If that non-terminal is followed by any other non-terminal then add “FIRST of other nonterminal” into FOLLOW set.

## **CODE:**

```
#include<stdio.h>
#include<string.h>
#include<conio.h>
#define max 20
char prod[max][10];
char ter[10],nt[10];
char first[10][10],follow[10][10];
int eps[10];
int count=0;
int findpos(char ch)
{
    int n;
    for(n=0;nt[n]!='\0';n++)
        if(nt[n]==ch)
            break;
    if(nt[n]=='\0')
        return 1;
    return n;
}

int IsCap(char c)
{
    if(c >= 'A' && c <= 'Z')
        return 1;
    return 0;
}
```

```

void add(char *arr,char c)
{
int i,flag=0;
for(i=0;arr[i]!='\0';i++)
{
if(arr[i] == c)
{
flag=1;
break;
}
}
if(flag!=1)
arr[strlen(arr)] = c;
}
void addarr(char *s1,char *s2)
{
int i,j,flag=99;
for(i=0;s2[i]!='\0';i++)
{
flag=0;
for(j=0;;j++)
{
if(s2[i]==s1[j])
{
flag=1;
break;
}
}
if(j==strlen(s1) && flag!=1)
{

```

```

s1[strlen(s1)] = s2[i];
break;
}
}
}
}

void addprod(char *s)
{
int i;
prod[count][0] = s[0];
for(i=3;s[i]!='\0';i++)
{
if(!IsCap(s[i]))
add(ter,s[i]);
prod[count][i-2] = s[i];
}
prod[count][i-2] = '\0';
add(nt,s[0]);
count++;
}

void findfirst()
{
int i,j,n,k,e,n1;
for(i=0;i<count;i++)
{
for(j=0;j<count;j++)
{
n = findpos(prod[j][0]);
if(prod[j][1] == (char)238)

```

```

eps[n] = 1;
else
{
for(k=1,e=1;prod[j][k]!='\0' && e==1;k++)
{
if(!IsCap(prod[j][k]))
{
e=0;
add(first[n],prod[j][k]);
}
else
{
n1 = findpos(prod[j][k]);
addarr(first[n],first[n1]);
if(eps[n1] == 0)
e=0;
}

}

if(e==1)
eps[n]=1;
}
}
}

void findfollow()
{
int i,j,k,n,e,n1;
n = findpos(prod[0][0]);
add(follow[n],'$');

```

```

for(i=0;i<count;i++)
{
for(j=0;j<count;j++)
{
k = strlen(prod[j])-1;
for(;k>0;k--)
{
if(IsCap(prod[j][k]))
{
n=findpos(prod[j][k]);
if(prod[j][k+1] == '\0')    // A -> aB
{
n1 = findpos(prod[j][0]);
addarr(follow[n],follow[n1]);
}
if(IsCap(prod[j][k+1]))    // A -> aBb
{
n1 = findpos(prod[j][k+1]);
addarr(follow[n],first[n1]);
if(eps[n1]==1)
{
n1=findpos(prod[j][0]);
addarr(follow[n],follow[n1]);
}
}
else if(prod[j][k+1] != '\0')
add(follow[n],prod[j][k+1]);
}
}
}

```

```

}
}
}
void main()
{
char s[max],i;
printf("\nEnter the productions(type 'end' at the last of the production)\n");
scanf("%s",s);
while(strcmp("end",s))
{
addprod(s);
scanf("%s",s);
}
findfirst();
findfollow();
for(i=0;i<strlen(nt);i++)
{
printf("%c\t",nt[i]);
printf("%s",first[i]);
if(eps[i]==1)
printf("%c\t",(char)238);
else
printf("\t");
printf("%s\n",follow[i]);
}
getch();
}

```

- Consider the grammar

$$S \rightarrow ( L ) \mid a$$

$$L \rightarrow L , S \mid S$$

- Compute the function FIRST for all non terminals

$$\text{FIRST}(S) = \{ (, a )$$

$$\text{FIRST}(L) = \text{FIRST}(S) = \{ (, a )$$

$$S \rightarrow ( L ) \mid a$$

$$L \rightarrow L , S \mid S$$

- Compute the function FOLLOW for all non terminals

$$\text{FOLLOW}(S) = \{ \$, \text{FOLLOW}(L) \}$$

$$= \{ \$, ), , \}$$

$$\text{FOLLOW}(L) = \{ ), , \}$$



- Consider the grammar

$$S \rightarrow L = R \mid R$$

$$L \rightarrow * R \mid \mathbf{id}$$

$$R \rightarrow L$$

- Compute the function FIRST for all non terminals

$$\text{FIRST}(S) = \text{FIRST}(L) \text{ and } \text{FIRST}(R)$$

$$\text{FIRST}(L) = \{ *, \mathbf{id} \}$$

$$\text{FIRST}(R) = \text{FIRST}(L) = \{ *, \mathbf{id} \}$$

$$\text{Therefore } \text{FIRST}(S) = \{ *, \mathbf{id} \}$$

- Consider the grammar

$$S \rightarrow L = R \mid R$$

$$L \rightarrow * R \mid \mathbf{id}$$

$$R \rightarrow L$$

Compute the function FOLLOW for all non terminals

$$\text{FOLLOW}(S) = \{ \$ \}$$

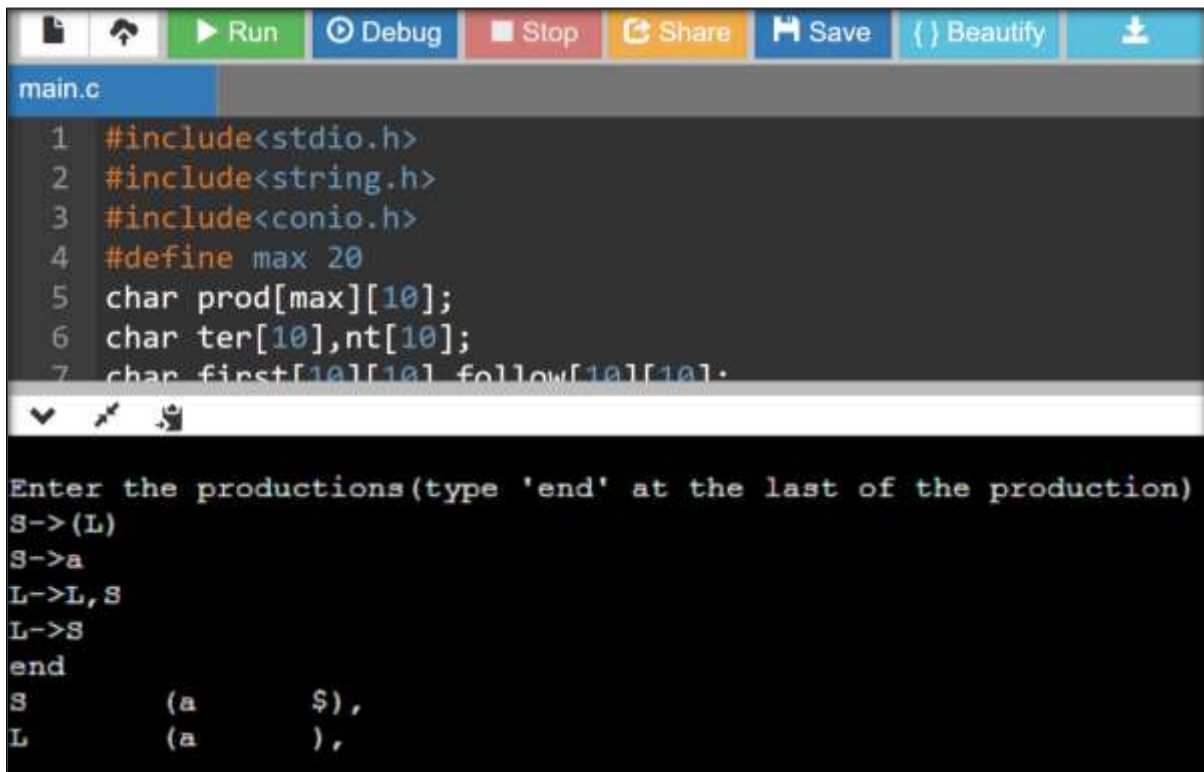
$$\text{FOLLOW}(L) = \{ =, \text{FOLLOW}(R) \}$$

$$= \{ =, \$ \}$$

$$\text{FOLLOW}(R) = \{ \text{FOLLOW}(S) \text{ \& } \text{FOLLOW}(L) \}$$

$$= \{ \$, = \}$$

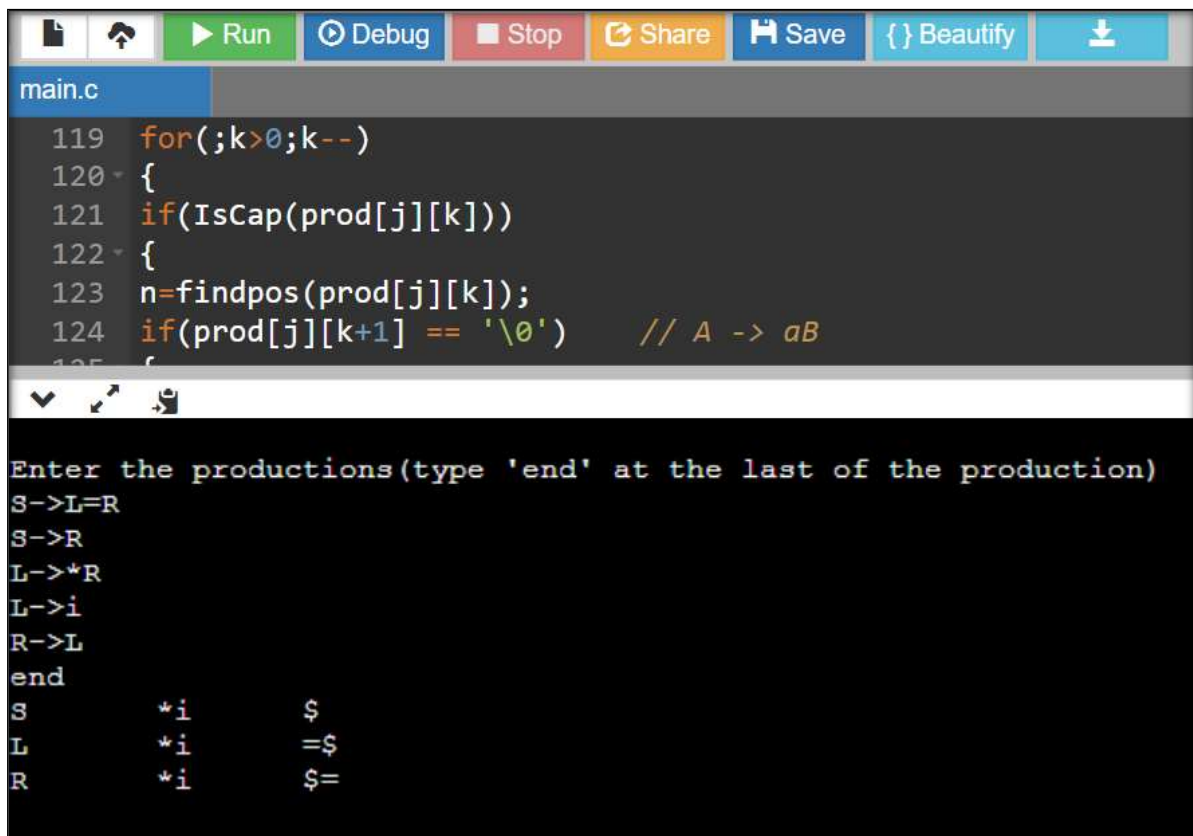
## OUTPUT:



The screenshot shows a C program in a code editor. The code defines a grammar with three non-terminals: S, L, and R. S is the start symbol. The grammar rules are: S → (L S), L → L S, and L → S. The program prompts the user to enter productions, and the user input is shown in the terminal window below the code.

```
main.c
1  #include<stdio.h>
2  #include<string.h>
3  #include<conio.h>
4  #define max 20
5  char prod[max][10];
6  char ter[10],nt[10];
7  char first[10][10],follow[10][10];

Enter the productions(type 'end' at the last of the production)
S->(L)
S->a
L->L,S
L->S
end
S      (a      $),
L      (a      ),
```



The screenshot shows the continuation of the C program. It defines the LR(0) item sets and transitions for the grammar. The code defines the LR(0) item sets and transitions. The program prompts the user to enter productions, and the user input is shown in the terminal window below the code.

```
main.c
119  for(;k>0;k--)
120  {
121  if(IsCap(prod[j][k]))
122  {
123  n=findpos(prod[j][k]);
124  if(prod[j][k+1] == '\\0')    // A -> aB
125  {
126  if(ter[n] == '\\0')
127  {
128  strcpy(ter+n,prod[j][k]);
129  }
130  }
131  }
132  }

Enter the productions(type 'end' at the last of the production)
S->L=R
S->R
L->*R
L->i
R->L
end
S      *i      $
L      *i      =$
R      *i      $=
```

**RESULT:** The program for finding first and follow is successfully compiled and executed

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