Lecture 2 RPC & Threads
c++ vs Go non-GC, typesate, memory sate, GC, stroughtforward requires book keeping to keep a track of
10n-GC, typesate, memory sate, GC, straightforward
requires "defective Coo?
book keeping to
Reep a track of
threads 4 free
objects, evor messages are a nightmare
3 J C 0
· Threads
→ manage I/o concurrency of one program talks to ? "go routines" many diff servers I
1 Thread - 1 PC, 1 stack,
1 Thread - 1 PC, 1 stack, 1 address space 1 address space diff threads could refer
1 address space
diff threads could reference other threads' stacks
→ I/O Concurrency: Can have one program launch RPC's to many server 4 wou't for replies.
RPC's to many server &
wou't tor replies.
→ Parallelism: One program can use multiple wres multiple ores multiple of CPU, truly in parallel
→ Convenience: sleep → periodic thing → sleep
event-driven: single thread in a single loop sitting and waiting for any kind of input.
and waiting for any kind of input.

· Threads vs pro	ocess		
	inside a p	process -> mu	.Utiple threads
-, OS keeps prov	resses L		ram suns one
-, OS keeps prov separate.		ONIX proces	s which then
		all go so	s which then utines are
· Thread challens	yes .	inside.	that process
-> Showing men	nary (same address	space)	
but very	yes navy (same address early to get bug	۷ أ	
n=n+1	-> multiple threa	ids might be	executing the
" RACE"	> multiple threa same code		7
solution: locks	(mutex in go)		
	MIL, DERLY,	mu. Unlock()	
-> should locks	be private?		
- not always	a good - embedde	ed inside DS	methode not
idea	0 known	to perogramm	ver 4
- might not	be private? .a good is embedde known want to use I	ocks at all	ج ک
> Problems with Th	reads		
1) Race conditions			
D Co-ordination:	when locking, dif	f threads dont '	know others
	exist, but we diff threads to a	might intention	enally want
	diff threads to	interact	U
- "channels" in	go to send data	between threads	
- sync.cond:	go to send data - condition variables	(it other thro	ed is waiting
0		tor you, g	iveir a RICK /
- wait group!	lauching an unkn for them to finis	own no ot so	utines 4 waiting
	for them to fini	sh	O
- deadlock:	$T_2 \longrightarrow $	T2	

-> Webcrawler demo - final challenge: knowing when to stop the crawl? - social crawler · DFS on web graph, keeps mapped (set) to remember what it fetched.

passed by ref and not copy

(in go all calls share pointer to obj in memory) - Concurrent Crawler · Shared data objects & locks · Non-shared but synced by channels