

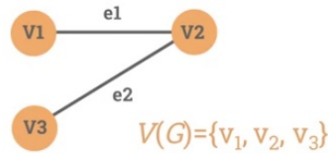
Cheatsheet - Graphs

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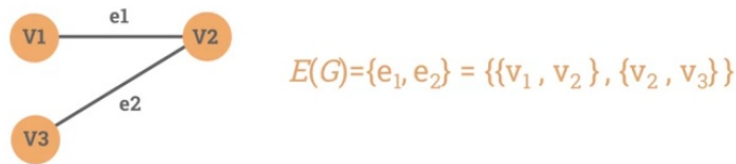
1. Intro

Graphs are **discrete** structures consisting of **vertices (nodes)** and **edges** connecting them. Graph theory is an area of discrete mathematics which studies these types of discrete structures.

The **graph** G can be represented as an ordered pair $G = (V, E)$, where V is a set of nodes/vertices and E is a set of edges, lines or connections. A **vertex** (singular of "vertices") is a basic element of a graph, usually drawn as a node or a dot. The set of vertices of G is usually denoted by $V(G)$ or V .



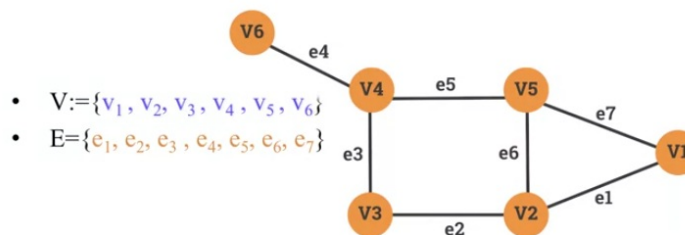
An **edge** is a link between 2 vertices, usually drawn as a line connecting two vertices. The set of edges in a graph G is usually denoted by $E(G)$ or E .



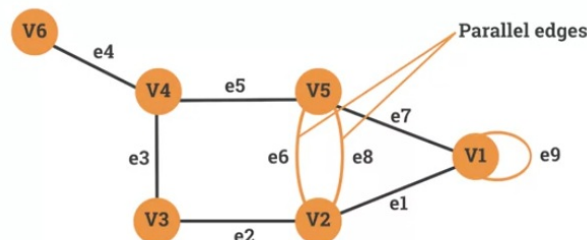
Two vertices are said to be **adjacent** if they are endpoints of the same edge. Two edges are said to be **adjacent** if they share the same vertex. If a vertex v is an endpoint of an edge e , then we say that e and v are **incident**.

A **directed graph**, also called a **digraph**, is a graph in which the edges have a direction. This is usually indicated with an arrow on the edge.

1.1. Examples



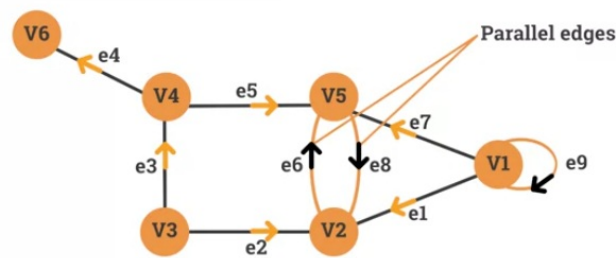
v_1 and v_2 are endpoints of the edge e_1 . We say that v_1 and v_2 are **adjacent**.
The edges e_1 and e_7 share the same vertex v_1 . We say that e_1 and e_7 are **adjacent**.
The vertex v_2 is an endpoint of the edge e_1 . We say that e_1 and v_2 are **incident**.



v_2 and v_5 are linked with two edges (e_6 and e_8).
 e_6 and e_8 are called **parallel** edges.

v_1 is linked to itself by e_9 . The edge e_9 is called a **loop**.

And an example of a directed graph:



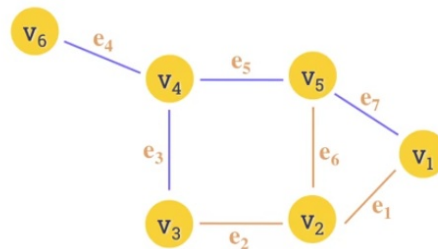
e_1 is a connection from v_1 to v_2 but not from v_2 to v_1

e_6 is a connection from v_2 to v_5 whereas e_8 is a connection from v_5 to v_2

2. Concepts

2.1. Walk

A **walk** is a sequence of vertices and edges of a graph where vertices and edges can be repeated. A **walk of length k** in a graph is a succession of k (not necessarily different) edges of the form uv, vw, wx, \dots, yz .

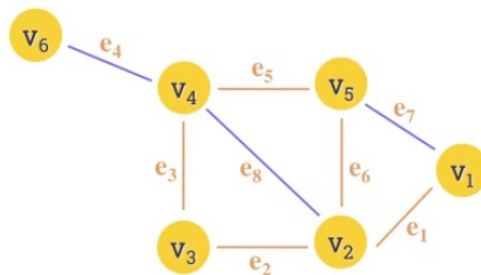


$$v_1v_2, v_2v_3, v_3v_2, v_2v_5 = e_1, e_2, e_2, e_6 = v_1v_2v_3v_2v_5$$

A walk of **length 4** from v_1 to v_5 (passes twice through the edge e_2)

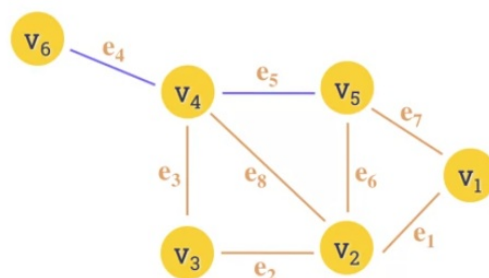
2.2. Trail

A **trail** is a walk in which no edge is repeated. In a trail, vertices can be repeated but no edge is ever repeated. For example, e_1, e_2, e_3, e_5, e_6 is a trail:



2.3. Circuit

A **circuit** is a closed trail. Circuits can have repeated vertices only. For example, $e_7, e_6, e_8, e_3, e_2, e_1$ is a circuit:

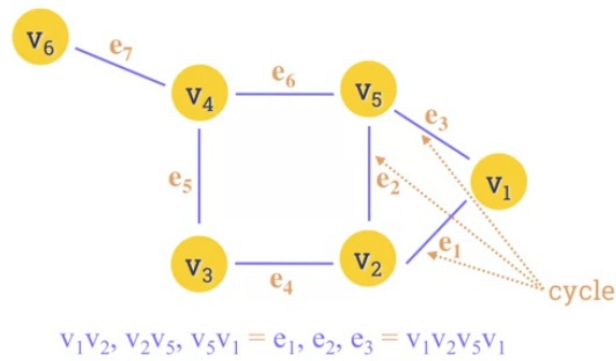


2.4. Path

A **path** is a trail in which neither vertices nor edges are repeated.

2.5. Cycle

A **cycle** is a closed path, consisting of edges and vertices where a vertex is reachable from itself.

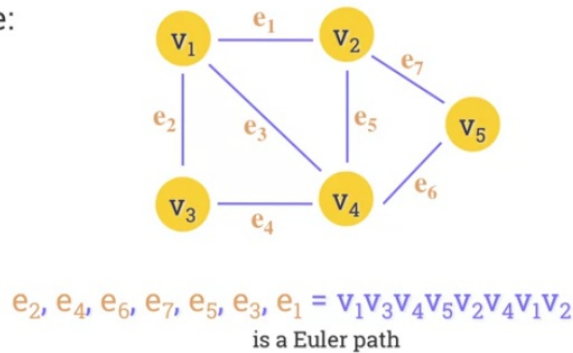


A walk of **length 3** from v_1 to v_1 = closed path = cycle

2.6. Eulerian Path

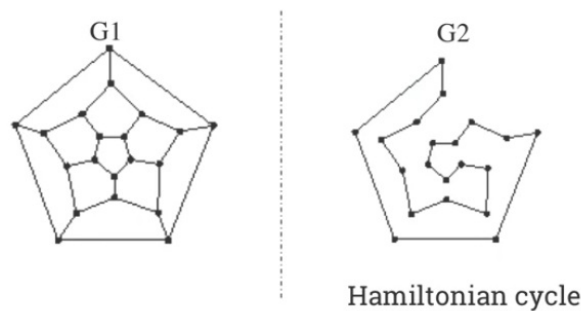
A **Eulerian path** in a graph is a path that uses each edge precisely once. If such a path exists, the graph is called **traversable**.

Example:



2.7. Hamiltonian Path, Cycle & Graph

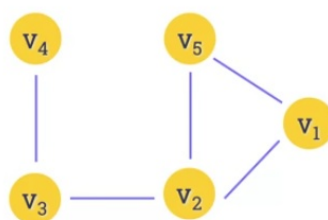
A **Hamiltonian path** (also called a *traceable path*) is a path that visits each vertex exactly once. A **Hamiltonian cycle** is a cycle that visits each vertex exactly once (except for the starting vertex, which is visited once at the start and once again at the end).



A graph that contains a Hamiltonian cycle is called a **Hamiltonian graph**. Any Hamiltonian cycle can be converted to a Hamiltonian path by removing one of its edges.

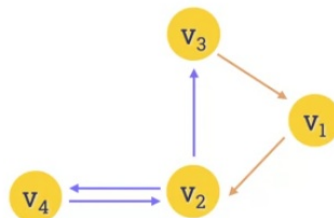
2.8. Connectivity

An **undirected** graph is **connected** if you can get from **any node to any other** by following a **sequence of edges**. Or, **any two nodes** are **connected** by a path.



Connected graph

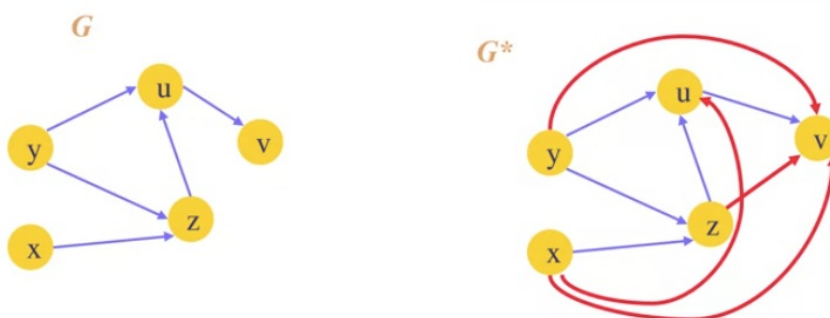
A directed graph is **strongly connected** if there is a **directed path** from any node to any other node.



Strongly connected directed graph

2.9. Transitive Closure

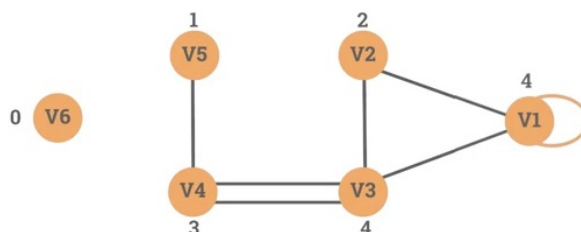
Given a digraph G , the transitive closure of G is the digraph G^* such that G^* has the same vertices as G . If G has a directed path from u to v ($u \neq v$), G^* has a directed edge from u to v .



The transitive closure provides reachability information about a digraph.

3. Degree of a Vertex

The degree of a vertex is the number of edges incident on v . A loop contributes **twice** to the degree. An **isolated vertex** has a degree of 0.

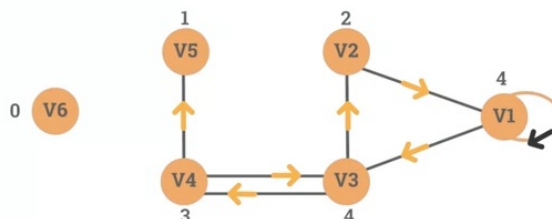


In the case of directed graphs, $\text{In-deg}(v)$ is the number of edges for which v is the terminal vertex. $\text{Out-deg}(v)$ is the number of edges for which v is the initial vertex.

And the degree $\text{deg}(v)$ is:

$$\text{deg}(v) = \text{Out-deg}(v) + \text{In-deg}(v)$$

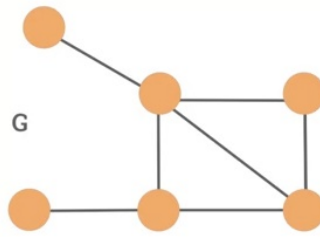
A loop contributes **twice** to the degree as it contributes 1 to both in-degree and out-degree.



$$\begin{aligned} \text{deg}(v_1) &= \text{in-deg}(v_1) + \text{out-deg}(v_1) = 2 + 2 = 4 \\ \text{deg}(v_2) &= \text{in-deg}(v_2) + \text{out-deg}(v_2) = 1 + 1 = 2 \\ \text{deg}(v_3) &= \text{in-deg}(v_3) + \text{out-deg}(v_3) = 2 + 2 = 4 \\ \text{deg}(v_4) &= \text{in-deg}(v_4) + \text{out-deg}(v_4) = 1 + 1 = 2 \\ \text{deg}(v_5) &= \text{in-deg}(v_5) + \text{out-deg}(v_5) = 1 + 0 = 1 \\ \text{deg}(v_6) &= \text{in-deg}(v_6) + \text{out-deg}(v_6) = 0 + 0 = 0 \end{aligned}$$

3.1. Degree Sequence

Given an undirected graph G , a **degree sequence** is a **monotonic non-increasing** sequence of the vertex degrees of all the vertices of G . The **sum of the degree sequence** of a graph is always **even**.



The degree sequence of G is: 4,3,3,2,1,1

Sum of the degree sequence = $1+1+2+3+3+4 = 14$

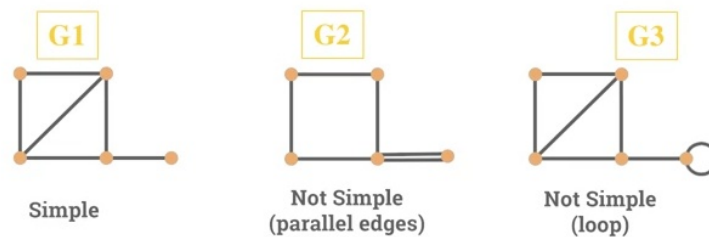
Given a graph G , the sum of the degree sequence of G is twice the number of edges in G .

$$\text{Number of edges of } G = \frac{\text{sum of degree sequences of } G}{2}$$

4. Special Graphs

4.1. Simple Graphs

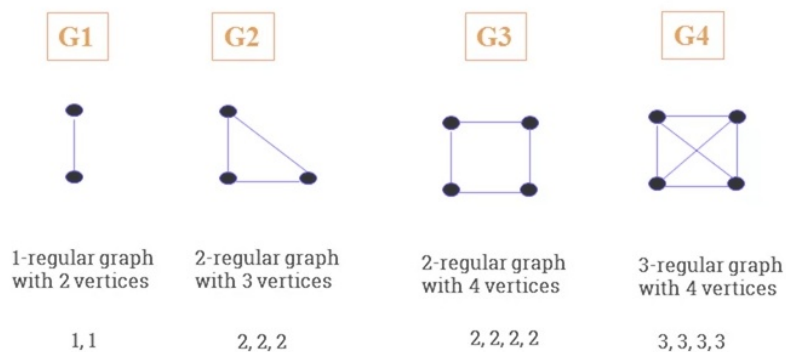
A **simple graph** is a graph without **loops** and **parallel** edges.



Given a **simple** graph G with n vertices, then the degree of each vertex of G is at most equal to $n - 1$.

4.2. Regular Graphs

A graph is said to be **regular** of degree if all local degrees are the same number. A graph G where all the vertices are of the same degree, r , is called an **r -regular** graph.

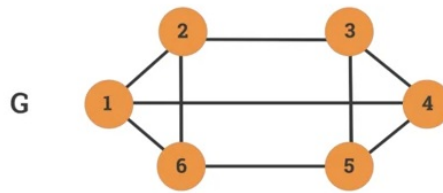


Given a **r -regular** graph G with n vertices, then the following is true:

Degree sequence of $G = r, r, r, \dots, r$ (n times)

Sum of degree sequence of $G = r \times n$

Number of edges in $G = r \times \frac{n}{2}$

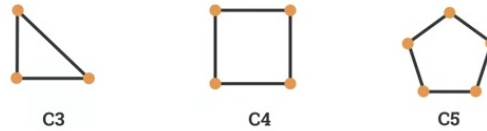


Degree Sequence = 3,3,3,3,3,3

Sum of degree sequence = $3 \times 6 = 18$

Number of edges = $18/2 = 9$

4.3. Special Regular Graphs: Cycles



C_3 is 2-regular graph with 3 vertices

C_4 is 2-regular graph with 4 vertices

C_5 is 2-regular graph with 5 vertices

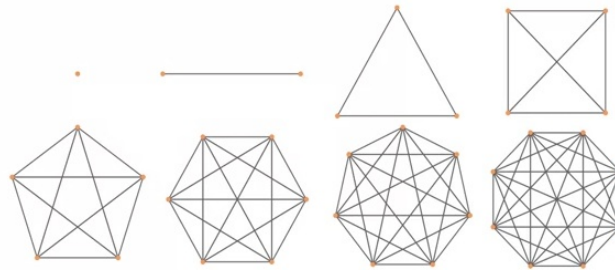
deg seq. of C_3 = 2,2,2

deg seq. of C_4 = 2,2,2,2

deg seq. of C_5 = 2,2,2,2,2

4.4. Complete Graphs

A **complete graph** is a **simple graph** where **every pair of vertices** are **adjacent** (linked with an edge). We represent a complete graph with n vertices using the symbol K_n .



A complete graph with n vertices, K_n , has the following properties:

Every vertex has a degree $(n - 1)$

Sum of degree sequence = $n(n - 1)$

Number of edges = $\frac{n(n - 1)}{2}$



There are 5 vertices

Degree of each vertex = $(5-1) = 4$

Sum of deg. Seq. = $5(5-1) = 20$

Number of edges = $5(5-1)/2 = 20/2 = 10$