Understanding SHAMap

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XRPL Core Dev Bootcamp 2025



Core Ledger Team @ Ripple

- Low-level protocols, high-performance code
- Personal focus: memory + throughput scalability
- We're hiring! Engineers welcome
 - Staff Software Engineer -> https://ripple.com/careers/all-jobs/job/6437475/
 - Senior Software Engineer -> https://ripple.com/careers/all-jobs/job/6437475/



Agenda

- What is SHAMap?
- Why SHAMap matters
- Data Structure Fundamentals
- Nodes and Hashing
- Mutability & Snapshots
- Traversal & Iteration
- Synchronization & Proofs
- Storage, Caching, and Thread Safety
- ► Q&A



SHAMap in One Sentence

A hybrid of <u>Merkle tree</u> and <u>radix trie</u> used to store transactions or ledger state

TODO: remind Merkle and Radix tree/trie structure

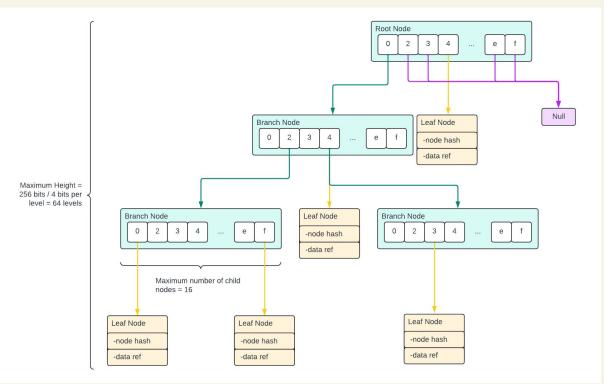


Role of SHAMap in XRPL

- Backbone of ledger structure
- Used to store:
 - Ledger State
 - Transactions



High-Level Diagram





Why Merkle Tree + Patrcia Trie?

- Efficient comparison: O(1) with hashes
- Cryptographic integrity
- Ordered key space



Uniform Leaf Type

SHAMap is homogeneous: leaves are of one type per map

State map vs. Tx map



Types of SHAMap Nodes

- SHAMapTreeNode (base)
- SHAMapInnerNode
- SHAMapLeafNode (+ subclasses)



SHAMapInnerNode Overview

- Up to 16 children
- Bitset + hashes
- Holds no data
- FullBelow optimization

SHAMapLeafNode Overview

- Holds SHAMapItem
- Can be account state, tx, or tx+meta
- Immutable once inserted

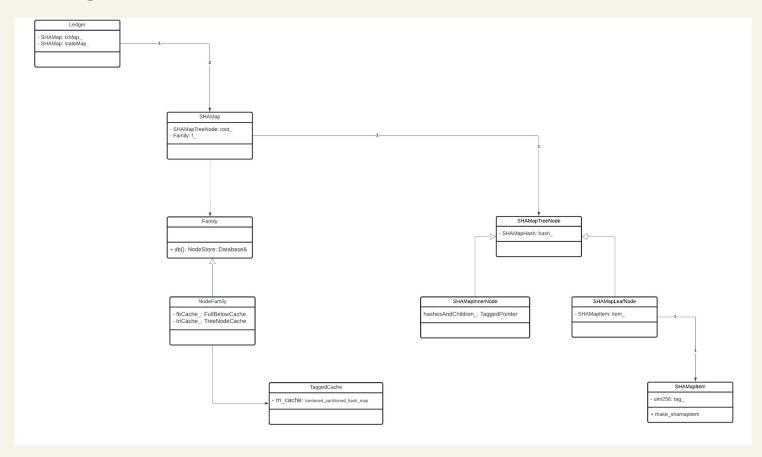


SHAMapItem

- tag
- size
- raw data
- intrusive_ptr based



Class diagram



Hashing: The Foundation

- Each node stores a hash
- Inner = hash of children
- Leaf = hash of content



Hash Update Logic

- `updateHash()` on leaves: compute hash
- Inner node: recompute after all children are updated



SHAMapHash Type

Custom hash type for storage

Used pervasively across nodes



Use Cases for Hashes

- Fast comparison
- Merkle proofs
- Data integrity in DB serialization



Copy-On-Write Model

- Nodes shareable via shared_ptr
- Modified via clone (copy-on-write)
- Each SHAMap has a `cowid`



Mutability & Snapshots

- Immutable: for validation/sync
- Mutable: for building ledgers
- Snapshots: efficient cloning



Traversal

visitNodes: DFS traversal

visitLeaves: filter + callback

Iteration via const_iterator

- Ordered iteration
- STL-style API: `begin()`, `end()`, `lower_bound()`



walkMap vs walkMapParallel

- walkMap: missing node detection
- walkMapParallel: multithreaded sync



Syncing - getMissingNodes

- Traverses tree
- Uses deferred async reads
- Marks fullBelow on success



Proof Paths

- `getProofPath()` builds Merkle path
- `verifyProofPath()` checks it
- Used for light clients



Canonicalization

- Ensures one node per hash
- Avoids duplicates in memory



Node Identification (SHAMapNodeID)

- Encodes depth + path as uint256
- Methods: `getChildNodelD`, `createlD`, etc.

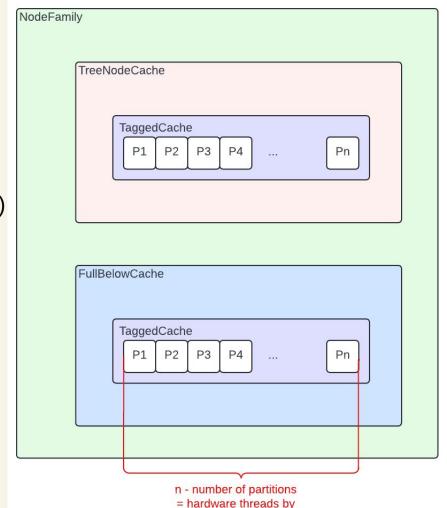
Inner Node Serialization

- Full format: all 16 branches
- Compressed: skip empty branches
- Chosen automatically



Caching Architecture

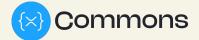
- TreeNodeCache (hash-indexed)
- FullBelowCache (fully synced subtrees)
- Family abstraction (cache+DB mgmt)



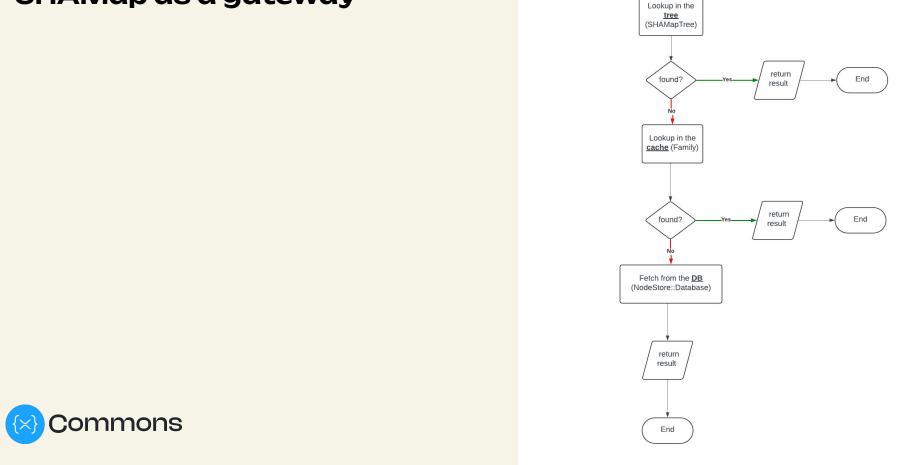


Thread Safety Mechanisms

- Atomic access in SHAMapInnerNode
- canonicalize() logic
- Cache-level synchronization



SHAMap as a gateway



find item in SHAMap

Design Summary

- Immutable-by-default
- Radix-16 + Merkle
- Fast sync, proof, comparison



What's Next?

- Memory footprint optimisations
- Transactions throughput optimisations

Thank you!

- Questions?
- Exercise (TODO)

