





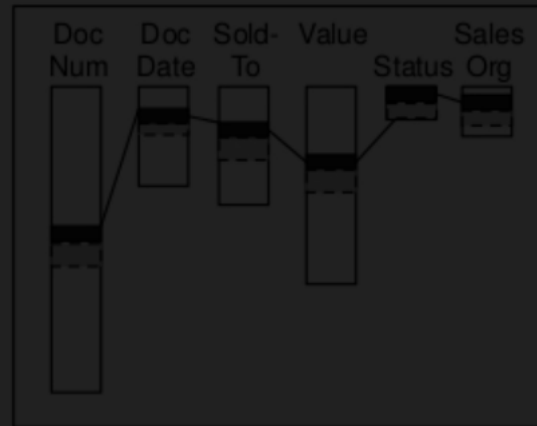






Optimal for row-wise access  
(e.g., SELECT \*)

```
SELECT *  
FROM Sales Orders  
WHERE Document Number = '95779216'  
(OLTP-style query)
```

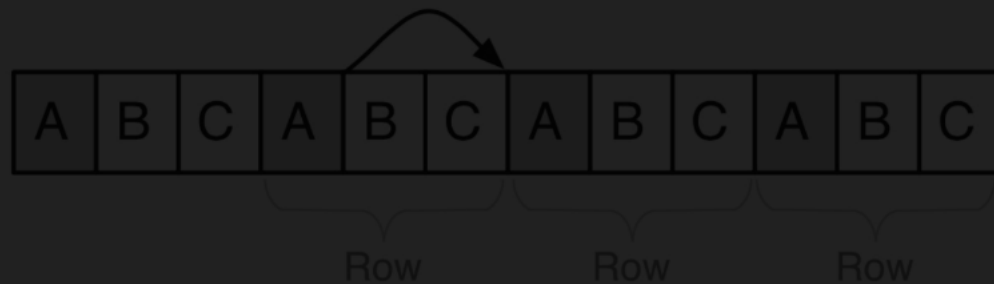


Optimal for attribute focused access  
(e.g., SUM, GROUP BY)

```
SELECT SUM(Value)  
FROM Sales Orders  
WHERE Document Date > 2011-08-28  
(OLAP-style query)
```



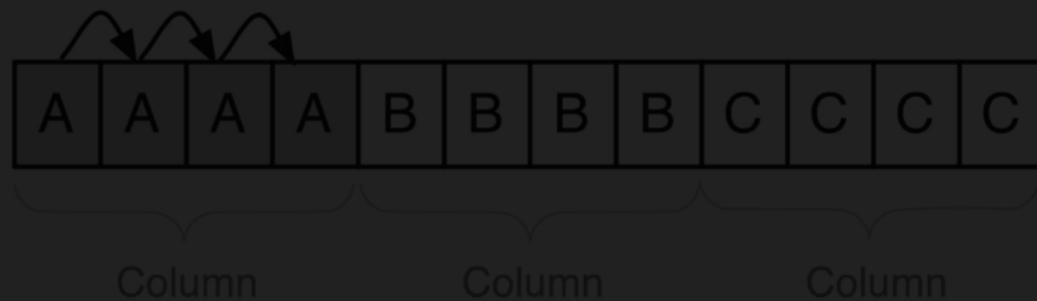
## Column Operation



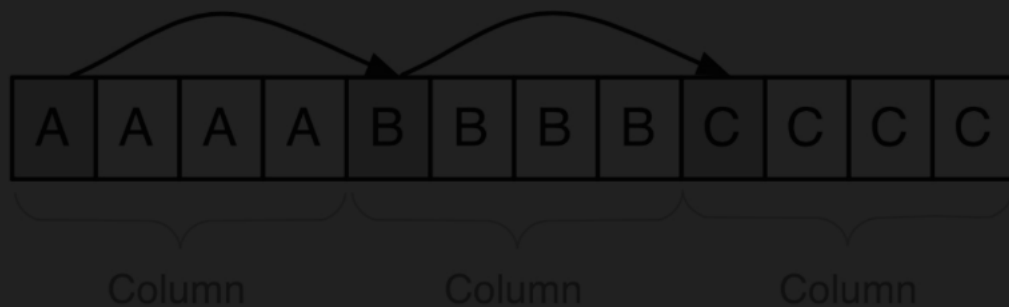
## Row Operation



## Column Operation



## Row Operation



Stride access ~256 secs

## Row Store – Layout

Table: humans

	First Name	Last Name	Gender	Country	City	Birthday
Row 1						
Row 2						
Row 3						
...						
Row $8 \times 10^9$						

## Row Store – Full Table Scan

Table: humans

	First Name	Last Name	Gender	Country	City	Birthday
Row 1						
Row 2						
Row 3						
...						
Row $8 \times 10^9$						

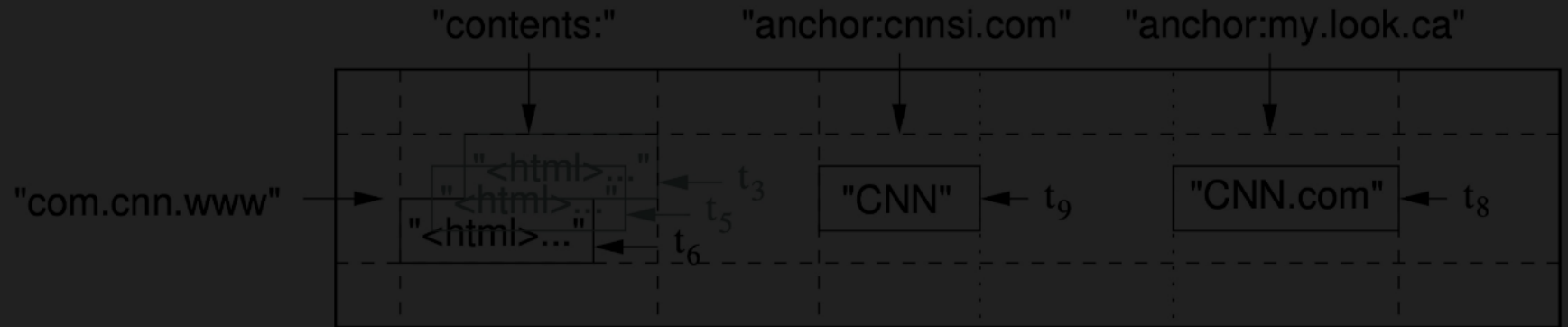
~0.5 secs

## Column Store – Layout

Table: humans					
First Name	Last Name	Gender	Country	City	Birthday

## Column Store – Full Column

Table: humans					
First Name	Last Name	Gender	Country	City	Birthday





하나의 테이블은 메타 태블릿 조각을 가지고 있음.  
(색인 데이터)

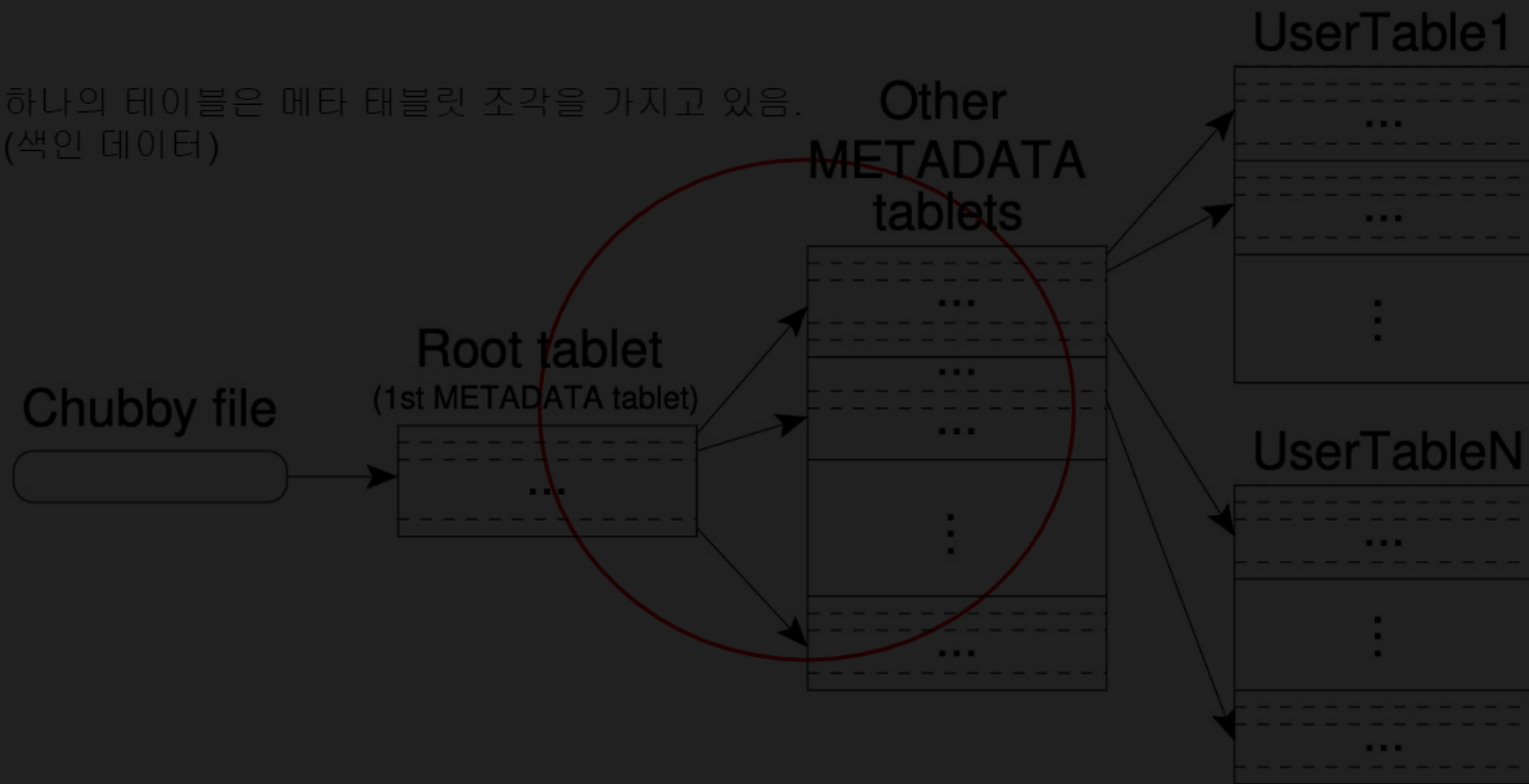


Figure 4: Tablet location hierarchy.





# DATABASE STORAGE ENGINES

## B-TREE

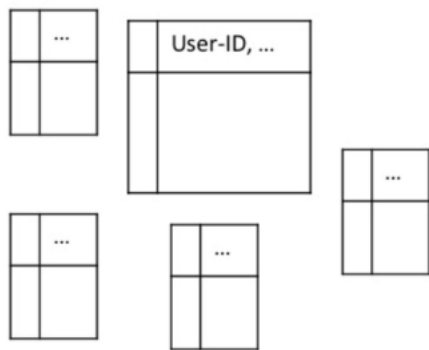


## LSM TREE



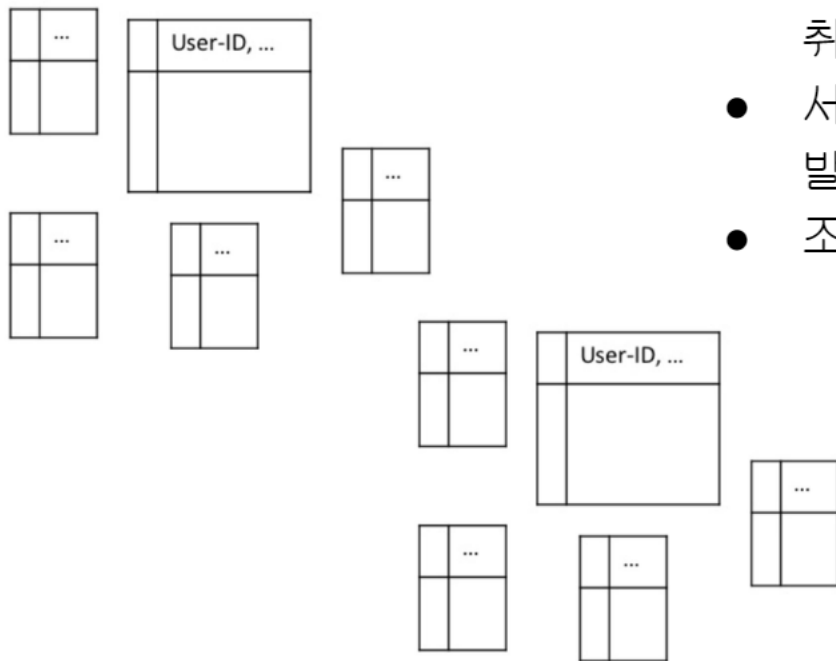


# 관계형 데이터베이스의 문제점



- **User profile**을 저장하는 테이블이 너무 커지거나, 쓰기연산이 **Heavy** 한 경우 취할 수 있는 전략?
- 서비스 정책 변경에 따른 **Scheme** 변화가 발생할 때 해야하는 일들?
- 조인 쿼리가 너무 무겁다면?

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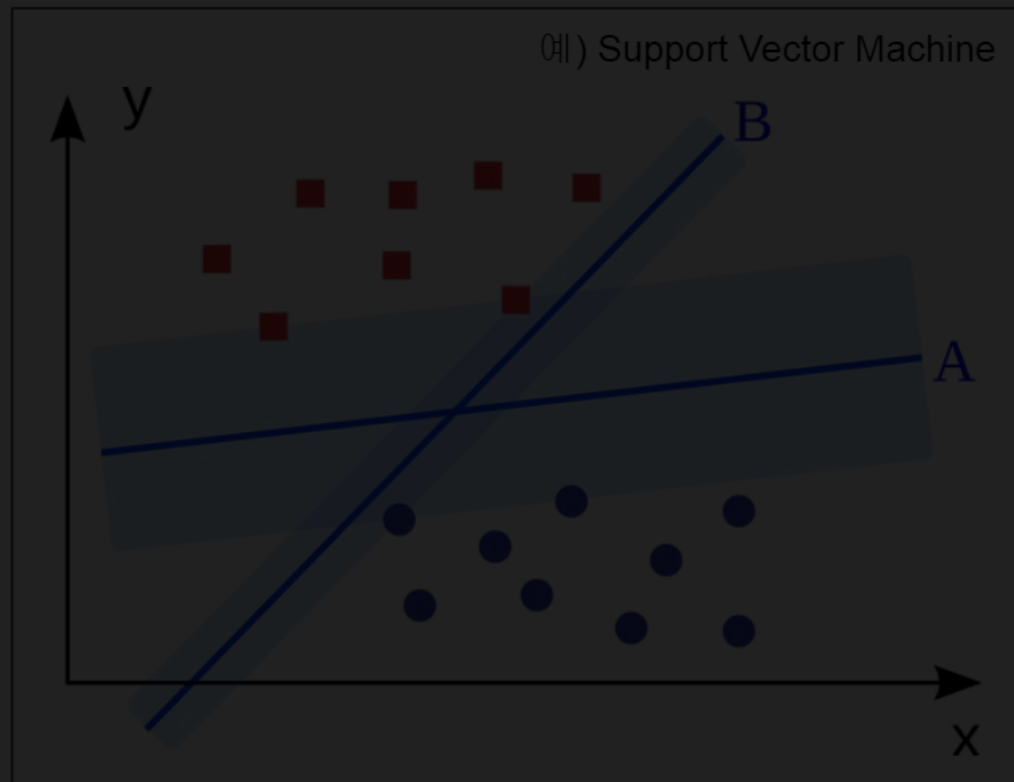








예) Support Vector Machine



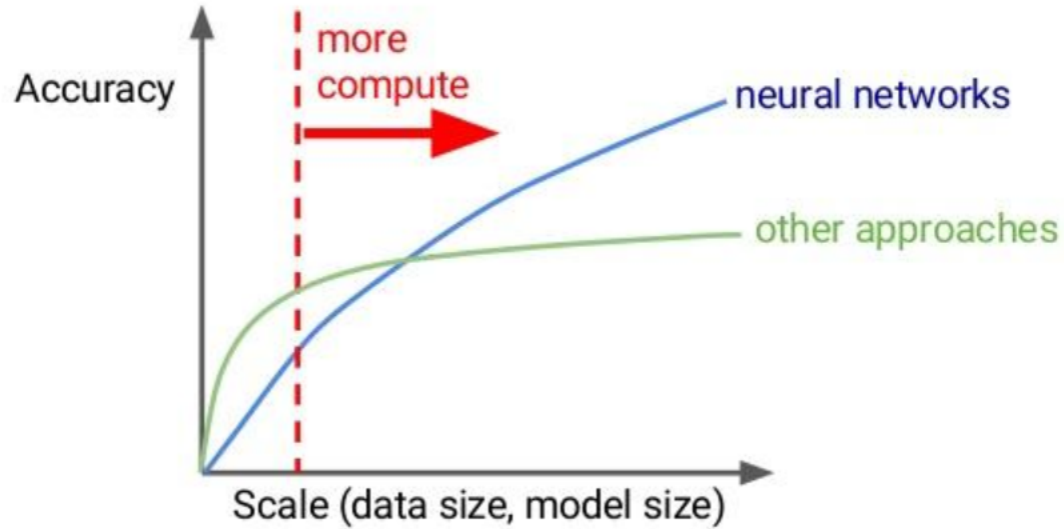
# 락 장르





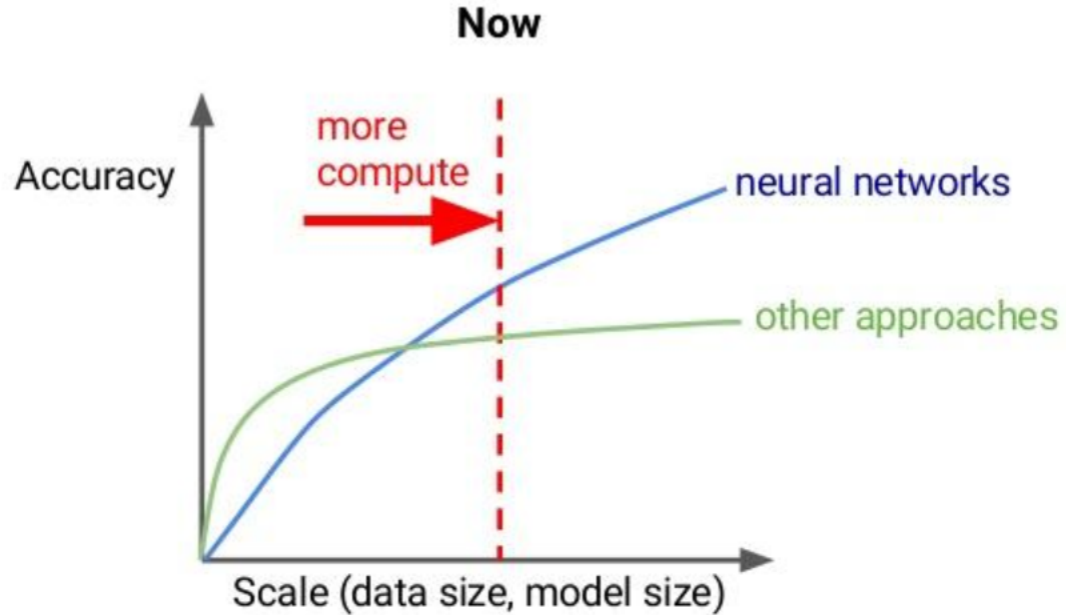


## 1980s and 1990s



과거 스케일에서는 다른 모델이 우수한 성능을 보여줬다면,





현재의 스케일에서는 딥 뉴럴 네트워크가 가장 우수하다.







