

Escuela de Ingeniería, Ciencia y Tecnología.

## Code Constructs

September 14th, 2020

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- 1 Read the introduction of the section 6 "Recognizing C Code Constructs in Assembly" and explain what means a "Code Construct". What aspects may impact the way as assembly code is generated?

Code construct makes reference to the abstraction of a code implementation. Showing properties and structures but no details. Assembly code abstracted from a C file may change depending on the compiler and OS used by the user. It will be noticed during this lab, the code obtained is in some small aspects different to the one on the book.

- 2 Read the section "Global vs Local Variables" and identify what are the differences in the compilation of a code that employs global vs one that employs local Variables.

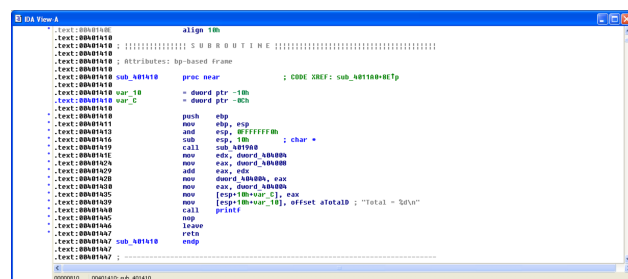
Consider the following two C codes:

```
1  #include<stdio.h>
2
3  int x=1;
4  int y=2;
5
6  void main(){
7      x = x+y;
8      printf("Total = %d\n", x);
9  }
10
```

```
1  #include<stdio.h>
2
3  void main(){
4      int x=1;
5      int y=2;
6
7      x = x+y;
8      printf("Total = %d\n", x);
9  }
10
```

Figure 1: C code for local and global variables.

Now It's possible to analyze the difference between the **assembly code** obtained on *IDA pro* and the one book. Let's analyze **global variable declaration** first:



```

.text:00401000 align 10h
.text:00401000 ; SUBROUTINE
.text:00401000 attributes: bp-based frame
.text:00401000 proc near ; CODE XREF: sub_401000+00p
.text:00401000 var_10 = dword ptr -10h
.text:00401000 var_0 = dword ptr -0Ch
.text:00401000
.text:00401000 push ebp
.text:00401001 mov esp, ebp
.text:00401002 and esp, 0FFFFFFFh ; char *
.text:00401003 sub esp, 10h
.text:00401004 call sub_401000
.text:00401005 mov eax, dword_40CF60
.text:00401006 add eax, dword_40CF60
.text:00401007 mov ecx, dword_40CF60
.text:00401008 mov ecx, dword_40CF60
.text:00401009 push ecx
.text:0040100A push offset aTotalD ; "total = %d\n"
.text:0040100B call printf
.text:0040100C nop
.text:0040100D retn
.text:0040100E sub_401000 endp

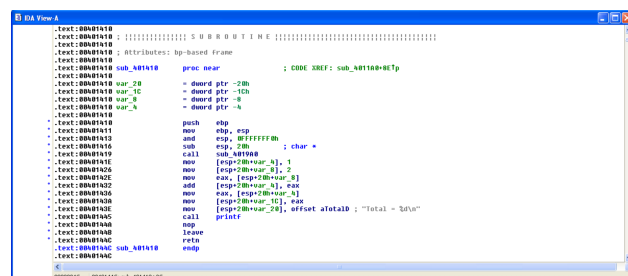
```

00401003 mov eax, dword\_40CF60  
00401008 add eax, dword\_40CF60  
0040100E mov dword\_40CF60, eax  
00401013 mov ecx, dword\_40CF60  
00401019 push ecx  
0040101A push offset aTotalD ; "total = %d\n"  
0040101F call printf

Listing 6-3: Assembly code for the global variable example in Listing 6-1

Figure 2: Comparison between own assembly code and book's for global variables.

Now the **global variables**:



```

.text:00401000 ; SUBROUTINE
.text:00401000 attributes: bp-based frame
.text:00401000 proc near ; CODE XREF: sub_401000+00p
.text:00401000 var_20 = dword ptr -20h
.text:00401000 var_1C = dword ptr -1Ch
.text:00401000 var_8 = dword ptr -8
.text:00401000 var_4 = dword ptr -4
.text:00401000
.text:00401000 push ebp
.text:00401001 mov esp, ebp
.text:00401002 and esp, 0FFFFFFFh ; char *
.text:00401003 sub esp, 20h
.text:00401004 call sub_401000
.text:00401005 mov [esp+20h+var_8], 1
.text:00401006 mov [esp+20h+var_8], 2
.text:00401007 add [esp+20h+var_8], eax
.text:00401008 mov [esp+20h+var_8], eax
.text:00401009 mov [esp+20h+var_8], eax
.text:0040100A call printf
.text:0040100B nop
.text:0040100C retn
.text:0040100D sub_401000 endp

```

00401006 mov dword ptr [ebp-4], 0  
0040100D mov dword ptr [ebp-8], 1  
00401014 mov eax, [ebp-4]  
00401017 add eax, [ebp-8]  
0040101A mov [ebp-4], eax  
0040101D mov ecx, [ebp-4]  
00401020 push ecx  
00401021 push offset aTotalD ; "total = %d\n"  
00401026 call printf

Listing 6-4: Assembly code for the local variable example in Listing 6-2, without labeling

Figure 3: Comparison between own assembly code and book's for local variables.

The difference on the compilation of a code with *global* and *local* variables is the way they are referenced. Global variables are referenced by *memory address*, the local ones are referenced by *stack address*. It means stacks are used to store all the variables and values created inside a function. Those stacks are used only during the execution of the function that's why this variables doesn't get an actual memory space.

### 3 Read the section "Disassembling Arithmetic Operations" and explain to your classmates how the operations (addition, subtraction, increment, decrement and modulo) are represented in assembly code.

Consider the following C code.

```

void main(){
    int a = 0;
    int b = 1;
    a = a + 11;
    a = a - b;
    b = a % 3;
}

```

Now Analyze and compare the own *assembly code* vs. the one provided by the book.

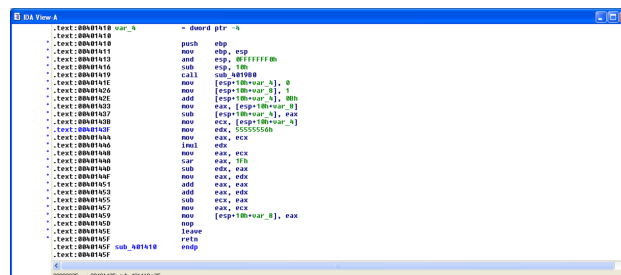


Figure 4: Assembly code obtained on *IDA pro*.

```

00401006    mov     [ebp+var_4], 0
0040100D    mov     [ebp+var_8], 1
00401014    mov     eax, [ebp+var_4]
00401017    add     eax, 0Bh
0040101A    mov     [ebp+var_4], eax
0040101D    mov     ecx, [ebp+var_4]
00401020    sub     ecx, [ebp+var_8]
00401023    mov     [ebp+var_4], ecx
00401026    mov     edx, [ebp+var_4]
00401029    sub     edx, 1
0040102C    mov     [ebp+var_4], edx
0040102F    mov     eax, [ebp+var_8]
00401032    add     eax, 1
00401035    mov     [ebp+var_8], eax
00401038    mov     eax, [ebp+var_4]
0040103B    cdq
0040103C    mov     ecx, 3

```

```

00401041    idiv    ecx
00401043    mov     [ebp+var_8], edx

```

Listing 6-7: Assembly code for the arithmetic example in Listing 6-6

Figure 5: Assembly code provided by the book.

*Addition* is placed on the following 2 lines, doing the addition first, then reorganizing the memory, associating the value obtained to the stack address of the variable `a`. The process is analog for the **subtraction** on the lines 5 and 6.

4 Read the section "Recognizing if Statements" and explain to your classmates how to recognize an if/else structure in assembly code.

```
1 #include<stdio.h>
2
3 void main(){
4     int x = 1;
5     int y = 2;
6     if(x==y){
7         printf("x equals y.\n");
8     }else{
9         printf("x is not equal to y.\n")
10     }
11 }
```

```

00401006 mov     [ebp+var_8], 1
0040100D mov     [ebp+var_4], 2
00401014 mov     eax, [ebp+var_8]
00401017 cmp     eax, [ebp+var_4] ❶
0040101A jnz     short loc_40102B ❷
0040101C push    offset aXEqualsY_ ; "x equals y.n"
00401021 call    printf
00401026 add     esp, 4
00401029 jmp     short loc_401038 ❸
0040102B loc_40102B:
0040102B push    offset aXIsNotEqualToY ; "x is not equal to y.n"
00401030 call    printf

```

First we make sure that the beginning and the end of the program is clear. It starts on *text:0040141E*(line 1) and ends on *00401454*.

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doesn't occur the program continues with the *if* section of the code (prints *x is not equal to y*).

Both of those options end with *jmp* instructions that redirect the process to the code after the if statement.

## 5 Read the section "Recognizing Nested if Statements" and explain to your classmates how to recognize a "Nested IF" structure in assembly code.

Consider the following *C* code:

```

1  #include<stdio.h>
2
3  void main (){
4      int x = 1;
5      int y = 1;
6      int z = 2;
7      if(x==y){
8          if(z==0){
9              printf("x equals y.\n");
10             }else{
11                 printf("x is not equal to y.\n");
12             }
13         }else{
14             if(z==0){
15                 printf("z zero and x != y.\n");
16             }else{
17                 printf("z non-zero and x != y.\n");
18             }
19         }
20     }
  
```

Now Analyze and compare the own *assembly code* vs. the one provided by the book. item[R].] Consider the following *C* code.

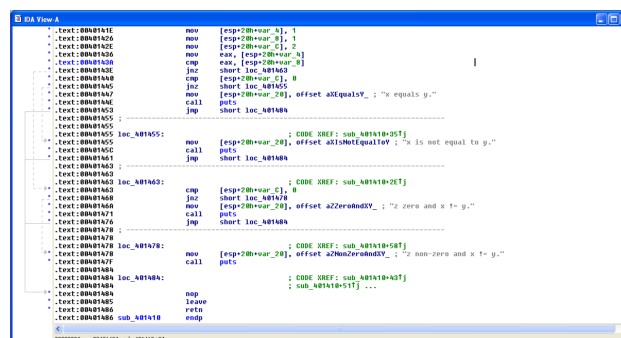


Figure 7: Assembly code obtained on *IDA pro*.

```

00401006    mov     [ebp+var_8], 0
0040100D    mov     [ebp+var_4], 1
00401014    mov     [ebp+var_C], 2
0040101B    mov     eax, [ebp+var_8]
0040101E    cmp     eax, [ebp+var_4]
00401021    jnz     short loc_401047
00401023    cmp     [ebp+var_C], 0
00401027    jnz     short loc_401038
00401029    push    offset aZIsZeroAndXY_ ; "z is zero and x = y.\n"
0040102E    call    printf
00401033    add     esp, 4
00401036    jmp     short loc_401045
00401038 loc_401038:
00401038    push    offset aZIsNonZeroAndX ; "z is non-zero and x = y.\n"
0040103D    call    printf
00401042    add     esp, 4

```

```

00401045 loc_401045:
00401045    jmp     short loc_401069
00401047 loc_401047:
00401047    cmp     [ebp+var_C], 0
0040104B    jnz     short loc_40105C
0040104D    push    offset aZZeroAndXY_ ; "z zero and x != y.\n"
00401052    call    printf
00401057    add     esp, 4
0040105A    jmp     short loc_401069
0040105C loc_40105C:
0040105C    push    offset aZNonZeroAndXY_ ; "z non-zero and x != y.\n"
00401061    call    printf
00401061

```

Listing 6-11: Assembly code for the nested if statement example shown in Listing 6-10

Figure 8: Assembly code provided by the book.

The behavior of this assembly is similar to the one on the "single if" code. The difference is the use of the *jnz*'s one inside the others according to the intersection of the conditions.

Just like on the exercise before, all of the functions produced according to the conditions end with *jmp* instructions that lead the code to what's after the nested if's.

## 6 Read the section "Recognizing Loops" and explain to your classmates how to recognize a FOR structure in assembly code.

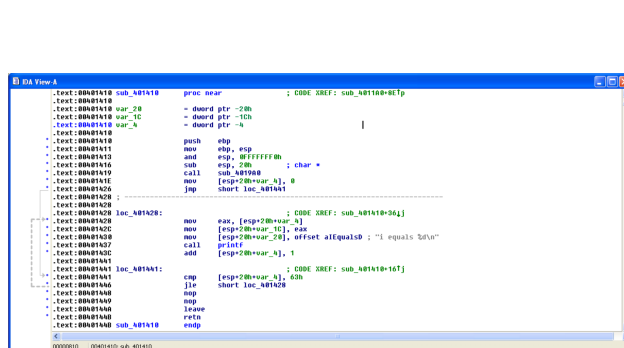
Consider the following C code:

```

1  #include<stdio.h>
2
3  void main(){
4      int i;
5      for(i=0; i<100;i++){
6          printf("i equals %d\n", i);
7      }
8  }

```

Now Analyze and compare the own *assembly code* vs. the one provided by the book.



```

00401004    mov     [ebp+var_4], 0
0040100B    jmp     short loc_401016
0040100D loc_40100D:
0040100D    mov     eax, [ebp+var_4]
00401010    add     eax, 1
00401013    mov     [ebp+var_4], eax
00401016 loc_401016:
00401016    cmp     [ebp+var_4], 64h
0040101A    jge     short loc_40102F
0040101C    mov     ecx, [ebp+var_4]
0040101F    push    ecx
00401020    push    offset aID ; "i equals %d\n"
00401025    call    printf
0040102A    add     esp, 8
0040102D    jmp     short loc_40100D

```

Listing 6-13: Assembly code for the for loop example in Listing 6-12

Figure 9: Assembly code provided by the book.

Code starts creating the local variable *i*. After it it executes a *jmp* instruction that leads to a comparison, in this case it's a *jle* instruction. This checks if the first value is less or equal that the second. This condition determines if the for loop executes or not. If it does the *jle* instructions redirects the code to the section in which prints the value of *i* and increases it by one.

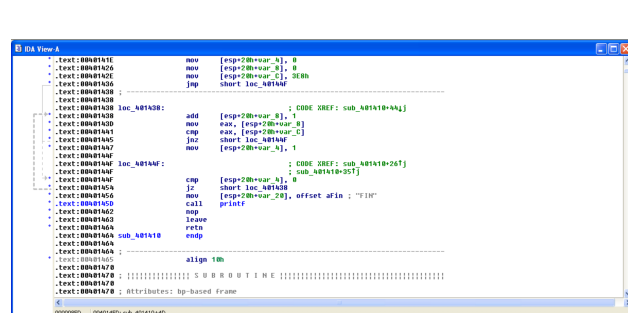
Then it goes back to the "condition checking" section to determine if this sequence of instructions should be repeated.

## 7 Read the section "Recognizing Loops" and explain to your classmates how to recognize a WHILE structure in assembly code.

Consider the following C code:

```
1  #include<stdio.h>
2
3  void main(){
4      int status = 0;
5      int result = 0;
6
7      while(status == 0){
8          result = performAction();
9          status = chechResult(result);
10     }
11 }
```

Now let's analyze and compare It's *assembly codes* with the one provided by the book:



00401036	mov	[ebp+var_4], 0
0040103D	mov	[ebp+var_8], 0
00401044	loc_401044:	
00401044	cmp	[ebp+var_4], 0
00401048	jnz	short loc_401063
0040104A	call	performAction
0040104F	mov	[ebp+var_8], eax
00401052	mov	eax, [ebp+var_8]
00401055	push	eax
00401056	call	checkResult
0040105B	add	esp, 4
0040105E	mov	[ebp+var_4], eax
00401061	jmp	short loc_401044

Listing 6-15: Assembly code for the while loop example in Listing 6-14

Figure 10: Assembly codes.

The while loop is a sort of for loop with out an **explicit** iteration variable. This loops has almost the same behavior as the for loop. The difference is the existence of an accessible iteration variable and the breaking point of the loop. Breaking point is also given by a comparison, in this case a **Boolean** one.

## 8 BIBLIOGRAPHY:

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