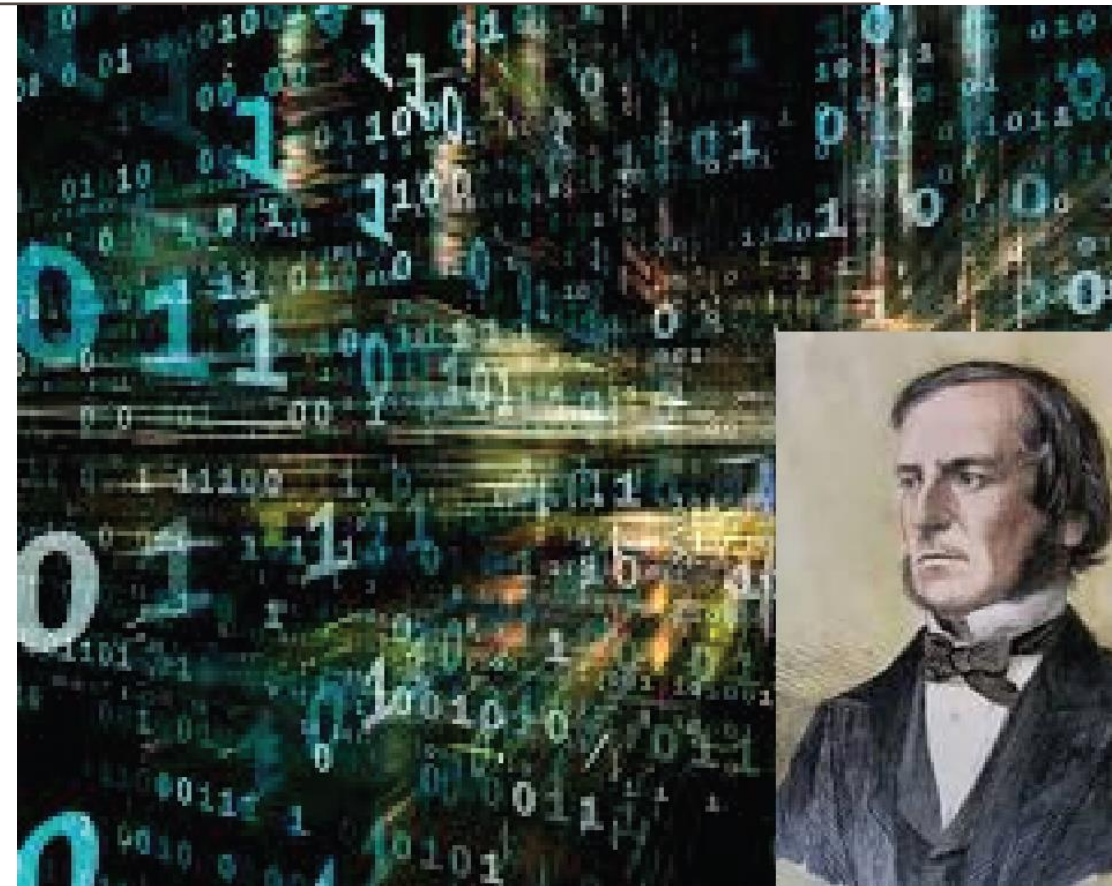




# DIGITAL CIRCUITS

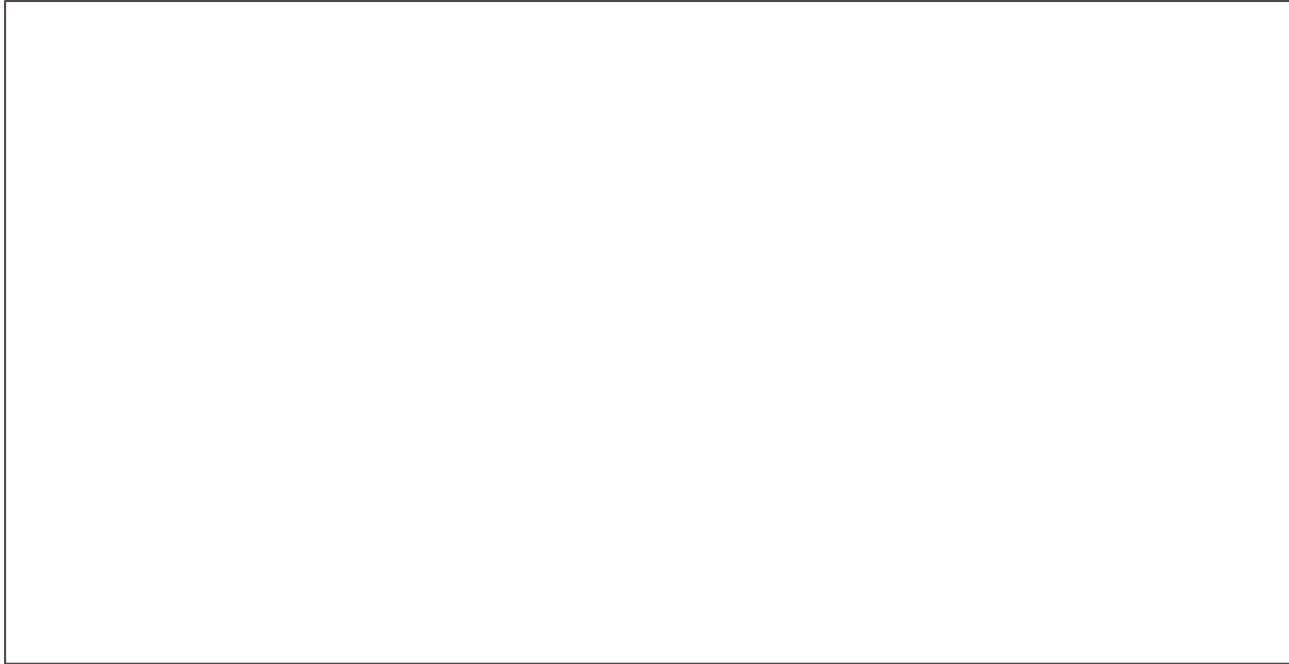
## Week-3, Lecture-1 Boolean Algebra

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14<sup>th</sup> August, 2018



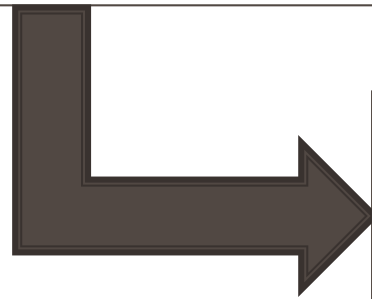
# Digital Circuits: Announcements/Revision

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# Digital Circuits



Boolean Algebra

# Theorems Proof Using *Huntington's Postulates*

Prove using Huntington's Postulates:  $xy + x'z + yz = xy + x'z$

The “.” operator can be omitted between variables to make representation look compact

Solution:

$$\begin{aligned} & xy + x'z + yz \\ &= xy + x'z + yz(x + x') \\ &= xy + x'z + yzx + yzx' \\ &= xy.1 + xyz + x'z.1 + x'zy \\ &= xy.(1 + z) + x'z.(1 + y) \\ &= xy + x'z \end{aligned}$$

Theorem:  $x + 1 = 1$

Proof:

$$\begin{aligned} & x + 1 = 1.(x + 1) \\ &= (x + x').(x + 1) \\ &= x + x'.1 \\ &= x + x' = 1 \end{aligned}$$

**2(a)**  $x + 0 = 0 + x = x$

2(b)  $x.1 = 1.x = x$

3(a)  $x + y = y + x$

3(b)  $x.y = y.x$

4(a)  $x.(y + z) = x.y + x.z$

4(b)  $x + (y.z) = (x + y).(x + z)$

5(a)  $x + x' = 1$

5(b)  $x.x' = 0$

# Duality Principle: Explanation

The dual of a true **statement** is also true.

$$x + x' = 1$$

Dual Statement  
 $x \cdot x' = 0$

$$x \cdot (y + z) = x \cdot y + x \cdot z$$

Dual Statement  
 $x + (y \cdot z) = (x + y) \cdot (x + z)$

$$(x_1 + x_2)' = x_1' x_2'$$

Dual Statement  
 $(x_1 \cdot x_2)' = x_1' + x_2'$

$$x + y$$

Dual Expression:  $x \cdot y$

They are not equal.

$(x + \bar{z}) \cdot y$  Dual Expression:  
 $x \cdot \bar{z} + y$

They are not equal (Verify using Truth Table)

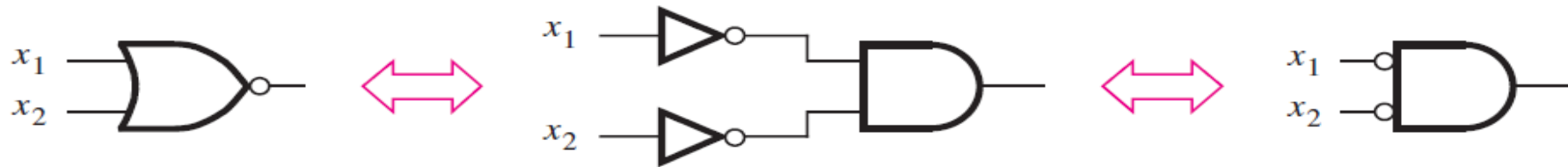
- Statement has LHS and RHS
- Dual has to be taken on both the sides
- Principle of duality applies to a **statement** and **NOT to expression**
- In general, **expression** and its **dual expression** are not EQUAL

## IMPORTANT

- Given a function, DO NOT realize the dual of that function.
- In general, they will be different functions

# DeMorgan's Theorems (1)

$$(x_1 + x_2)' = x_1'x_2'$$



$$(b) \overline{x_1 + x_2} = \bar{x}_1 \bar{x}_2$$

$x_1$	$x_2$	$(x_1 + x_2)'$
0	0	<b>1</b>
0	1	<b>0</b>
1	0	<b>0</b>
1	1	<b>0</b>

$x_1$	$x_2$	$x_1'$	$x_2'$	$x_1'x_2'$
0	0	1	1	<b>1</b>
0	1	1	0	<b>0</b>
1	0	0	1	<b>0</b>
1	1	0	0	<b>0</b>

# DeMorgan's Theorems (1)

$$(x_1 + x_2)' = x_1'x_2'$$

Proof:

Let  $A = x_1 + x_2$  and  $B = x_1'x_2'$

$$B + A$$

$$= x_1'x_2' + x_1 + x_2 = x_1 + x_1'x_2' + x_2$$

$$= (x_1 + x_1')(x_1 + x_2') + x_2$$

$$= x_1 + x_2' + x_2 = x_1 + 1$$

$$= 1(\text{should prove } x + 1 = 1)$$

$$B.A$$

$$= x_1'x_2'(x_1 + x_2)$$

$$= x_1'x_2'x_1 + x_1'x_2'x_2$$

$$= 0 + 0 = 0$$

$$\mathbf{2(a)} \quad x + 0 = 0 + x = x$$

$$2(b) \quad x.1 = 1.x = x$$

$$3(a) \quad x + y = y + x$$

$$3(b) \quad x.y = y.x$$

$$4(a) \quad x.(y + z) = x.y + x.z$$

$$4(b) \quad x + (y.z) = (x + y).(x + z)$$

$$5(a) \quad x + x' = 1$$

$$5(b) \quad x.x' = 0$$

# DeMorgan's Theorems (1)

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$$(x_1 + x_2)' = x_1' x_2'$$

Proof:

Let  $A = x_1 + x_2$  and  $B = x_1' x_2'$

$$B + A = 1$$

$$B.A = 0$$

## Postulate 5: Existence of a complement

For every element  $x \in B$ , there exists  $x' \in B$  such that:

a.  $x + x' = 1$

b.  $x.x' = 0$

Postulate 5 imply that  $A' = B$

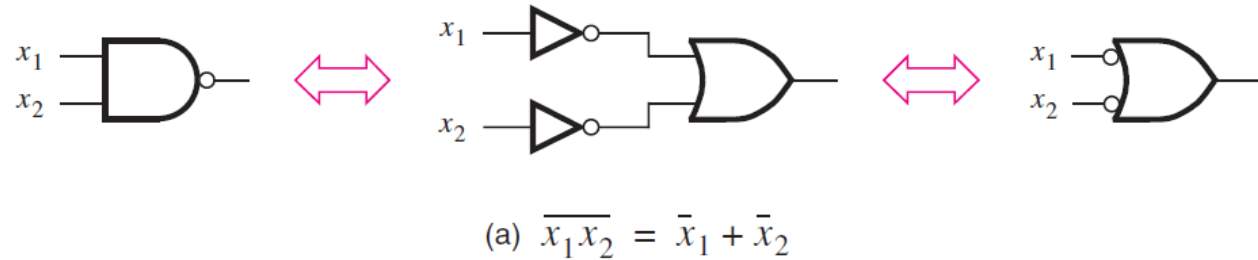
$$\Rightarrow (x_1 + x_2)' = x_1' x_2'$$



# DeMorgan's Theorems

$$(x_1 + x_2)' = x_1' x_2'$$

$$(x_1 x_2)' = x_1' + x_2' \text{ (Dual)}$$



$x_1$	$x_2$	$(x_1 x_2)'$
0	0	<b>1</b>
0	1	<b>1</b>
1	0	<b>1</b>
1	1	<b>0</b>

$x_1$	$x_2$	$x_1'$	$x_2'$	$x_1' + x_2'$
0	0	1	1	<b>1</b>
0	1	1	0	<b>1</b>
1	0	0	1	<b>1</b>
1	1	0	0	<b>0</b>

# DeMorgan's Theorems (More variables)

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$$(x_1 + x_2)' = x_1'x_2'$$

$$(x_1 + x_2 + x_3)' = x_1'x_2'x_3'$$

Proof:

Let  $x_2 + x_3 = A$

$$(x_1 + x_2 + x_3)'$$

$$= (x_1 + A)'$$

$$= x_1'A' = x_1'(x_2 + x_3)'$$

$$= x_1'(x_2'x_3') = x_1'x_2'x_3'$$

# DeMorgan's Theorems (More variables)

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$$(x_1 + x_2)' = x_1' x_2'$$

$$(x_1 x_2)' = x_1' + x_2' \text{ (Dual)}$$

$$(x_1 + x_2 + x_3)' = x_1' x_2' x_3'$$

$$(x_1 x_2 x_3)' = x_1' + x_2' + x_3'$$

$$(x_1 + x_2 + x_3 + x_4)' = x_1' x_2' x_3' x_4'$$

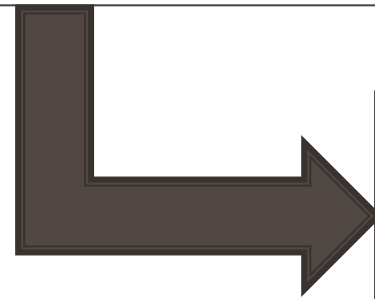
$$(x_1 x_2 x_3 x_4)' = x_1' + x_2' + x_3' + x_4'$$

$$(x_1 + x_2 + \cdots + x_N)' = x_1' x_2' \cdots x_N'$$

$$(x_1 x_2 \cdots x_N)' = x_1' + x_2' + \cdots + x_N'$$

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# Digital Circuits



Synthesis using Logic  
Gates

# Basic Definitions (1)

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**Boolean variable:** A variable which can take only '0' or '1' values.

Example:  $x, y, a, b, \dots$

**Literal:** A Boolean variable or its complement.

Example:  $x, x', y, y', a, a', b, b', \dots$

**Product Term:** A single literal or a logical AND of two or more literals.

Example:  $x, x', xy, yy', abc, a'bc'd, \dots$

**Sum Term:** A single literal or a logical OR of two or more literals.

Example:  $x, x', x + y, y + y', a + b + c, a' + b + c' + d, \dots$

# Basic Definitions (2)

**Normal Term:** A product or sum term in which no variable appears more than once.

Example:

$x, a + b, abc', \dots$  are Normal Terms

$a + a', abcb', a + b + b'$  are NOT Normal Terms

**In the following expression, identify:**

$$a + a'b + (c + c') + dd'$$

- a) Variables and Literals
- b) Product terms and Sum Terms
- c) Normal Terms and Non-normal terms

- a) Variables:  $a, b, c, d$  and Literals:  $a, a', b, c, c', d, d'$
- b) Product terms:  $a, a'b, c, c', dd'$  and Sum Terms:  $a, (c + c')$
- c) Normal Terms:  $a, a'b$  and Non-normal terms:  $(c + c'), dd'$

# Basic Definitions (3)

**Sum of Product (SOP) Expression:** A *logical OR* of a set of *product terms*.

Examples:

$$ab + a'bc + ac'$$

$$abc + a'bc + abc' + ab'c'$$

$$a + b$$

**Product of Sum (POS) Expression:** A *logical AND* of a set of *sum terms*.

Examples:

$$(a + b)(a' + b + c)(a + c')$$

$$(a + b + c)(a' + b + c)(a + b + c')(a + b' + c')$$

$$ab$$

# Synthesis: Using Truth Table (1)

## Problem:

Design a network using **logic gates** that takes two inputs  $x_1$  and  $x_2$ . Assume that  $x_1$  and  $x_2$  represent the states of two switches, either of which may be open (0) or closed (1). The output of the network is 1 when  $(x_1, x_2)$  are in the states (0,0), (0,1) or (1,1). In the state (1,0) the output should be 0.

$x_1$	$x_2$	$f(x_1, x_2)$
0	0	1
0	1	1
1	0	0
1	1	1

**A sum of product (SOP) expression is formed using truth table.**



# Synthesis: Truth Table to SOP

$x_1$	$x_2$	$f(x_1, x_2)$	Product Term
0	0	1	$x_1'x_2'$
0	1	1	$x_1'x_2$
1	0	0	$x_1x_2'$
1	1	1	$x_1x_2$

$$f(x_1, x_2) = x_1'x_2' + x_1'x_2 + x_1x_2$$

## How to find the product term:

- For each row find the product term
  - The variable that has value 0 is complemented and the variable that has value 1 is taken without complement

## How to find the sum:

- Take the sum of product terms that result in function being 1

# Synthesis: SOP to Network

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$$f(x_1, x_2) = x_1'x_2' + x_1'x_2 + x_1x_2$$

