

18.10.11

Lecture 8 handout

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(7.1)

Definition

A network $N(x,y)$ is a ^{simple} digraph D , a source x and a sink y , and a capacity function $c: E \rightarrow \mathbb{R}$

Intermediate vertices: $V(D) - \{x, y\} =: I$.

A flow is a function $f: E \rightarrow \mathbb{R}$, $0 \leq f(e) \leq c(e)$ for all $e \in E$.

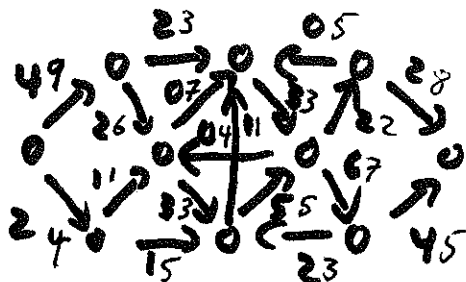
- conservation condition: $f^+(v) = f^-(v)$ for all $v \in I$.

For $X \subseteq V(D)$, $f^+(X) - f^-(X)$ is the net flow of X . $(f^+(X) := \sum_{v \in X} f^+(v))$.

$$\text{val}(f) := f^+(x) - f^-(x)$$

Prop: $\text{val}(f) = f^+(X) - f^-(X)$ for any $X \subseteq V(D)$ with $x \in X$ and $y \in V(D) \setminus X$.

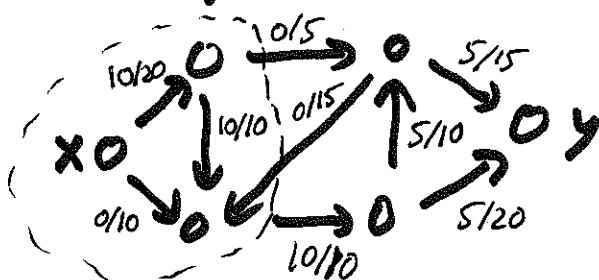
Problem: Find the maximal flow in $N(x,y)$



Cuts

A cut of $N(x,y)$ is a set $X \subseteq V(D)$ s.t.
 $x \in X$ and $y \in V(D) \setminus X$.

The capacity of X is the sum of capacities
 of edges from X to $V(D) \setminus X$.



$$\text{cap}(X) = 15$$

$$\text{val}(X) = 10.$$

Theorem: $\text{val}(f) \leq \text{cap}(X)$. Equality iff edges in $d^+(X)$
 are f -saturated, and edges in $d^-(X)$ are f -zero.
 In this case, f is a max flow and X is a min cut.

Proof: $\text{val}(f) = f^+(x) - f^-(x) = \sum_{v \in X} \left(\sum_w f(v \rightarrow w) - \sum_u f(u \rightarrow v) \right)$

$$= \sum_{v \in X} \left(\sum_{w \in V \setminus X} f(v \rightarrow w) - \sum_{u \in V \setminus X} f(u \rightarrow v) \right) \leq \sum_{v \in X} \sum_{w \in V \setminus X} f(v \rightarrow w)$$

equal if f -saturated
equal if f -zero

$$\leq \sum_{v \in X} \sum_{w \in V \setminus X} c(v \rightarrow w) = \text{cap}(X).$$

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Max flow min cut

The value of the maximal flow equals the cost of the minimal cut.

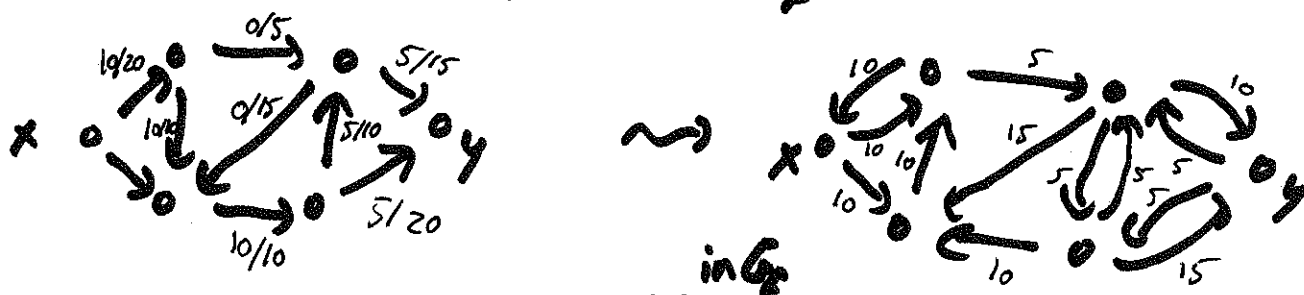
PROOF • one-directional: $o \rightarrow o \rightsquigarrow o \rightleftarrows o$

Residual capacity:

$$c_f(u \rightarrow v) ::= \begin{cases} c(u \rightarrow v) - f(u \rightarrow v) & \text{if } u \rightarrow v \in E \\ f(v \rightarrow u) & \text{if } v \rightarrow u \in E \\ 0 & \text{otherwise.} \end{cases}$$

"how much edge operates below capacity"

Form a residual graph G_f :



If there is no path from x to y , set $X = \{\text{vertices reachable from } x \text{ in } G_f\}$

Then X is a min cut, and L is a max flow from the proposition. Otherwise, augment the flow!

Inclusion-Exclusion



$$|A \cup B \cup C| = |A| + |B| + |C| \\ - |A \cap B| - |A \cap C| - |B \cap C| \\ + |A \cap B \cap C|$$

Example: Integers ≤ 100 divisible by none of 2, 3, 5, 7.

Example: Surjections $f: X \rightarrow Y$.

$$N(i): i \text{ not in } \text{Im}(f) = (n-1)^m$$

$$N(i_1, \dots, i_k) = (n-k)^m$$

$$|\text{Surjections}| = |\text{functions } X \rightarrow Y| - \sum N(i) + \dots$$

$$= n^m - \binom{n}{1} (n-1)^m + \binom{n}{2} (n-2)^m - \dots + (-1)^{n-1} \cdot n$$

$$= \sum_{i=0}^{n-1} (-1)^i \binom{n}{i} (n-i)^m$$

Example: Cayley's formula $T(n) = n^{n-2}$

Next time: Midterm test.