Lecture 7

A binary tree is

· empty . or

· a left binary tree and

a right binary tree.

The simplest one is empty.

A binary tree is full iff its left & right subtrees are empty, or both are full binary tree.

Prove that all full binary trees have odd # of nodes.

For heN, let P(h) be All FULL binary trees of height h have an odd # of nodes

Proof of the IN, Ph)

By complete induction

Inductive step. Let $h \in \mathbb{N}$, assume all full binary trees of height < h have odd # nodes. Let t be a full binary tree of height h.

Case: left & right subtrees of t are empty

Then tis just the single node tree, has one node, which is odd.

Case: left&right subtrees of t are full binary trees

Lot's call them to & tr, their heights are less than the height of t

By 14, the f.b.t.s th & tr. # of nodes in the is odd, # of nodes in the is odd. The # of nodes in t is 1 (the root) + # nodes in the + # of nodes in the which is 1 + odd + odd = odd

Well ordering Principle

Every non-empty set of notural numbers has a minimal element.

Ø: No Claim (no min)

 \mathbb{N} : Min is \mathbb{O}

(1.3.4.5.7) min is 1

Primes, min is 2.

Consider sets of real #s
(0.1)

Z No
[0.1) min is 0

we like the remainder minimal, say 9 > ra= bg +r

consider set of all natural numbers

let R= (r=1N, 3 g=N, 3 (415926=535898+r) 31415926 eR

not real number set!

By WOP, M has a minimum, call it ro. Let g. E.N be st. 3/415926=5358980+10 to EN since to EMCIN

if ro >53589, ro-53589 \in IN and 314/5926=53589(g+1)+(ro-53589) Then ro is not the minimum. contradiction So ro <53589