

Lecture 18
March 19th, 2015

$\overbrace{\quad}^{n \text{ digits}}$

$$\begin{array}{r} x_0 \\ y_0 \end{array} \quad \begin{array}{r} x_1 \\ y_1 \end{array}$$

$$X = 10^{\frac{n}{2}} x_0 + x_1$$

$$Y = 10^{\frac{n}{2}} y_0 + y_1$$

e.g.

$$\begin{array}{r} \overbrace{\quad}^8 \\ 1234 \quad 5678 \\ 3141 \quad 5926 \end{array} = 10^4 \cdot 1234 + 5678$$

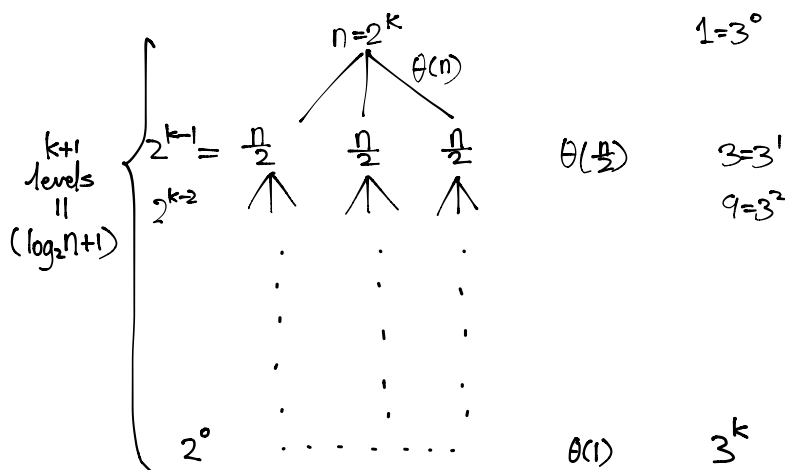
$$= 10^4 \cdot 3141 + 5926$$

multiples of size $\Theta(n)$: 3

$$XY = 10^n x_0 y_0 + 10^{\frac{n}{2}} (x_0 y_1 + x_1 y_0) + x_1 y_1$$

$$= 10^n x_0 y_0 + 10^{\frac{n}{2}} [(x_0 + x_1)(y_0 + y_1) - x_0 y_0 - x_1 y_1] + x_1 y_1$$

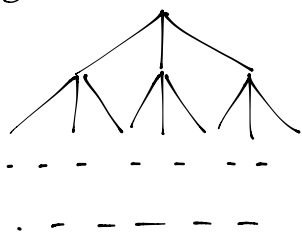
3 multiplications



arity of tree of calls
 $a=3$
division of problems of
 $\frac{1}{b} = \frac{1}{2}$ the size

use these 2 $\begin{cases} \log_2 3 < \log_2 4 = 2 \\ x^{\log_y z} = z^{\log_y x} \end{cases} \Rightarrow 3^k = 3^{(\log_2 n)} = (2^{\log_2 3})^{\log_2 n} = 2^{(\log_2 3)(\log_2 n)} = n^{\log_2 3} \in O(n^2)$

Why count leaves?



for binary tree:
 $3^0 + 3^1 + \dots + 3^{k-1} > 3^k - 1$
 $2^0 + 2^1 + \dots + 2^{k-1} = 2^k - 1$

3^k leaves at bottom

For n a power of 2

$$T(n) = \begin{cases} 3T(\frac{n}{2}) + n, & n > 1 \\ 1, & n = 1 \end{cases}$$

$$\begin{aligned} T(n) &= 3T(\frac{n}{2}) + n = 3[3T(\frac{n}{4}) + \frac{n}{2}] + n \\ &= 3[3[3T(\frac{n}{8}) + \frac{n}{4}] + \frac{n}{2}] + n \\ &= n + \frac{3}{2}n + \frac{3^2}{2^2}n + \frac{3^3}{2^3}n + \dots + \frac{3^k}{2^k}n \end{aligned}$$

If $n = 2^k$, the last term is $3^k = 3^{\log_2 n}$

$$\underbrace{n + \frac{3}{2}n + \dots + 3^{\log_2 n}}_{(1 + \log_2 n) \text{ terms}}$$

$$= n \left[1 + 3 \cdot \frac{1}{2} + 3^2 \cdot \frac{1}{2^2} + 3^3 \cdot \frac{1}{2^3} + \dots + \frac{3^{\log_2 n}}{2^{\log_2 n}} \right]$$

$$= n \left[1 + 3 \cdot \frac{1}{2} + 3^2 \cdot \frac{1}{2^2} + 3^3 \cdot \frac{1}{2^3} + \dots + \frac{3^{\log_2 n}}{2^{\log_2 n}} \right]$$

$$= \Theta(n \left(\frac{3}{2}\right)^{\log_2 n})$$

$$= \Theta(n \cdot n^{\log_2 \frac{3}{2}})$$

$$= \Theta\left(\frac{n(n^{\log_2 3})}{n^{\log_2 2}}\right)$$

$$= \Theta(n^{\log_2 3})$$

So doing work at each division
= doing work at these leaves

$$\begin{aligned} S &= 1 + r + r^2 + \dots + r^n \\ rS &= r + r^2 + r^3 + \dots + r^{n+1} \\ rS - S &= r^{n+1} - 1 \\ S &= \frac{r^{n+1} - 1}{r - 1} = \Theta(r^n) \text{ if } r > 1 \end{aligned}$$

→ comparing with fixed r

for n a power of b .

$$T(n) = \begin{cases} aT(\frac{n}{b}) + n^k, & n > 1 \\ 1, & n = 1 \end{cases}$$

$$\sum_{i=0}^{\log_b n} a_i \left(\frac{n}{b^i}\right)^k = \left[\sum_{i=0}^{\log_b n} \left(\frac{a}{b^k}\right)^i \right] n^k$$

$$\left[\frac{a}{b^k} > 1 \right]$$

$$\frac{a}{b^k} = 1$$

$$\left[\frac{a}{b^k} < 1 \right]$$

To be continued.

Tutorial

For n a power of b

$$T(n) = \begin{cases} aT(\frac{n}{b}) + n^k, & n > 1 \\ 1, & n = 1 \end{cases}$$

• $f \in \Theta(n^k)$

$$\exists B, \exists c_1, c_2. \forall n \geq B, c_1 \cdot n^k \leq f(n) \leq c_2 n^k$$

\downarrow n sufficiently large, $f(n)$ is bounded & sandwiched

$$\left. \begin{aligned} T(n) &\leq c_2 \cdot n^k \sum_{i=0}^{\log_b n} \left(\frac{a}{b^k}\right)^i \\ T(n) &\geq c_1 \cdot n^k \sum_{i=0}^{\log_b n} \left(\frac{a}{b^k}\right)^i \end{aligned} \right\} \Rightarrow O^k, \text{ just take } f(n) = n^k$$

The assignment asks you to compute this part:

$$\left(\frac{a}{b^k}\right)^i$$

$$L = \{0^n 1^n \mid n \in \mathbb{N}\}$$

X defined by:

• $\varepsilon \in X$

$$\forall s, t \in \{0, 1\}^* \cdot s \in X \Rightarrow 00s11 \in X$$

$$\cdot s01t \in X \Rightarrow st \in X$$

$X \subseteq L$ Intuition for $L \subseteq X$

0.1 Apply rule 3 in reverse: 0011

Then apply rule 2 in reverse to get ε .

Proof of $X \subseteq L$

By structure induction on def'n of X

Let $s \in X$ be arbitrary

WTS $s \in L$

① $s, \varepsilon \in X$ by def'n

② s is generated by 2nd Rule.

So $s = 00s'11$ for some $s' \in X$

s' generated before, $s \Rightarrow$ IH applies $s' \in L$

$$\exists n \in \mathbb{N}, s' = 0^n 1^n$$

$$\text{Then } s = 0^{n+2} 1^{n+2} \in L \quad \checkmark$$

③ s is generated by 3rd rule, so

$$\exists s', t' \in \{0, 1\}^* \text{ s.t. } s = s'01t'$$

IH applies to $s'01t'$

$$s'01t' \in L \text{ so } \exists n \in \mathbb{N} \ s'01t' = 0^n 1^n$$

$$\text{Then } s't' = 0^{n-1} 1^{n-1} \in L \quad (\text{and } n \geq 1 \text{ since } \underbrace{|s'01t'|}_{\text{length}} \geq 2)$$