## STA447/2006 (Stochastic Processes) Lecture Notes, Winter 2016

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**Note:** These lecture notes will be posted on the STA447/2006 course web page for your convenience, and will be updated regularly. However, they are just rough, point-form notes, with no guarantee of completeness or accuracy. They should in no way be regarded as a substitute for attending and actively learning from the course lectures.

#### **Introduction:**

- Discuss course web page, outline, evaluation, etc. (www.probability.ca/sta447)
- Schedule: will take 15-minute break if you return promptly!
- Your background knowledge: STA347 last semester? previously? other?
- Your status: undergrad? grad? special? STA specialist? major? Act Sci? other?
- You should already know basic probability theory: probability spaces, random variables, expected value, independence, conditional probability, discrete and continuous distributions, etc. (You do <u>not</u> need to know measure theory.)
- This class considers stochastic <u>processes</u>, i.e. randomness which <u>proceeds in time</u>.
  - Will develop their mathematical theory (with a <u>few</u> applications).

#### Markov chains:

- EXAMPLE (Frog Example):
  - 1000 lily pads arranged in a circle. (diagram)
  - Frog starts at pad #1000.
  - Each minute, she jumps either straight up, or one pad clockwise, or one pad counter-clockwise, each with probability 1/3.
  - (see e.g. www.probability.ca/frogwalk)
- So, P(at pad #1 after 1 step) = 1/3.
  - P(at pad #1000 after 1 step) = 1/3.
  - P(at pad #999 after 1 step) = 1/3.
  - $\mathbf{P}(\text{at pad } \#2 \text{ after } 2 \text{ steps}) = 1/9.$
  - P(at pad #999 after 2 steps) = 2/9.
  - etc.
- What happens in the long run?
  - What is P(frog at pad #428 after 987 steps)?

- What is  $\lim_{k\to\infty} \mathbf{P}(\text{frog at pad } \#428 \text{ after } k \text{ steps})$ ?
- Will the frog necessarily eventually return to pad #1000?
- Will the frog necessarily <u>eventually</u> visit every pad?
- And what happens if we have:
  - different jump probabilities?
  - different arrangement of the pads?
  - more pads?
  - infinitely many pads?
  - etc.
- A (discrete time, discrete space, time homogeneous) <u>Markov chain</u> is specified by three ingredients:
  - A state space S, any non-empty finite or countable set. (e.g. the 1000 lily pads)
  - <u>transition probabilities</u>  $\{p_{ij}\}_{i,j\in S}$ , where  $p_{ij}$  is the probability of jumping to j if you start at i. (So,  $p_{ij} \geq 0$ , and  $\sum_j p_{ij} = 1$  for all i.)
  - <u>initial probabilities</u>  $\{\nu_i\}_{i\in S}$ , where  $\nu_i$  is the probability of starting at i (at time 0). (So,  $\nu_i \geq 0$ , and  $\sum_i \nu_i = 1$ .)
- In the frog example,  $S = \{1, 2, 3, \dots, 1000\}$ , and

$$p_{ij} = \begin{cases} 1/3, & |j-i| \le 1\\ 1/3, & |j-i| = 999\\ 0, & \text{otherwise} \end{cases}$$

and  $\nu_{1000} = 1$  (with  $\nu_i = 0$  otherwise).

- Let  $X_n$  be the Markov chain's state at time n.
  - Then  $\mathbf{P}(X_{n+1} = j \mid X_n = i) = p_{ij}, \forall i, j \in S, n = 0, 1, 2, \dots$  (Doesn't depend on n: time-homogeneous.)
  - Also  $\mathbf{P}(X_{n+1} = j | X_0 = i_0, X_1 = i_1, X_2 = i_2, \dots, X_n = i_n) = p_{i_n j}$ . (Markov property.)
  - Also  $P(X_0 = i, X_1 = j, X_2 = k) = \nu_i p_{ij} p_{jk}$ , etc.
  - More generally,  $\mathbf{P}(X_0 = i_0, X_1 = i_1, \dots, X_n = i_n) = \nu_{i_0} p_{i_0 i_1} p_{i_1 i_2} \dots p_{i_{n-1} i_n}$
  - The random sequence  $\{X_n\}_{n=0}^{\infty}$  "is" the Markov chain.
- In the frog example:
  - $-\mathbf{P}(X_0 = 1000) = 1, \mathbf{P}(X_0 = 972) = 0, \text{ etc.}$
  - $\mathbf{P}(X_1 = 1) = 1/3$ ,  $\mathbf{P}(X_1 = 1000) = 1/3$ ,  $\mathbf{P}(X_2 = 2) = 1/9$ ,  $\mathbf{P}(X_2 = 999) = 2/9$ , etc.

# More Examples of Markov Chains:

- Example: simple random walk (s.r.w.).
  - Let 0 . (e.g. <math>p = 1/2)
  - Suppose you repeatedly bet \$1.
  - Each time, you have probability p of winning \$1, and probability 1-p of losing \$1.
  - Let  $X_n$  be net gain (in dollars) after n bets.
  - Then  $\{X_n\}$  is a Markov chain, with  $S = \mathbf{Z}$ ,  $\nu_0 = 1$ , and

$$p_{ij} = \begin{cases} p, & j = i+1\\ 1-p, & j = i-1\\ 0, & \text{otherwise} \end{cases}$$

- What happens in the long run? Will you necessarily go broke? etc.
- (see e.g. www.probability.ca/longrun)
- Example: Bernoulli process. (e.g. counting sunny days)
  - Let 0 . (e.g. <math>p = 1/2)
  - Repeatedly flip a "p-coin" (i.e., a coin whose probability of heads is p).
  - Let  $X_n = \#$  of heads on first n flips.
  - Then  $\{X_n\}$  is Markov chain, with  $S = \{0, 1, 2, \ldots\}$ ,  $X_0 = 0$  (i.e.  $\nu_0 = 1$ ), and

$$p_{ij} = \begin{cases} p, & j = i+1\\ 1-p, & j = i\\ 0, & \text{otherwise} \end{cases}$$

- Example: Branching process. (e.g. amoebas, infected people)
  - Let  $\rho$  be any prob dist on  $\{0, 1, 2, \ldots\}$ , the "offspring distribution".
  - Let  $X_n$  be the size of a "population" at time n.
  - Each of the  $X_n$  items at time n has a random number of offspring which is i.i.d.  $\sim \rho$ . (diagram)
  - That is,  $X_{n+1} = Z_{n,1} + Z_{n,2} + \ldots + Z_{n,X_n}$ , where  $\{Z_{n,i}\}_{i=1}^{X_n}$  are i.i.d.  $\sim \rho$ .
  - Here  $S = \{0, 1, 2, \ldots\}.$
  - $p_{ij}$  is more complicated; in fact  $p_{ij} = (\rho * \rho * \dots * \rho)(j)$ , a convolution of i copies of  $\rho$ . (In particular,  $p_{00} = 1$ .)

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- Will  $X_n = 0$  for some n? etc.
- Example: simple finite Markov chain.

- Let  $S = \{1, 2, 3\}$ , and  $X_0 = 3$ , and

$$(p_{ij}) = \begin{pmatrix} 0 & 1/2 & 1/2 \\ 1/3 & 1/3 & 1/3 \\ 1/4 & 1/4 & 1/2 \end{pmatrix}.$$

- What happens in the long run? (diagram)
- Example: Ehrenfest's Urn
  - Have d balls in total, divided into two urns.
  - At each time, we choose one of the d balls uniformly at random, and move it to the <u>other</u> urn.
  - Let  $X_n = \#$  balls in Urn 1 at time n.
  - Then  $\{X_n\}$  is Markov chain, with  $S = \{0, 1, 2, \dots, d\}$ , and  $p_{i,i-1} = i/d$ , and  $p_{i,i+1} = (d-i)/d$ , with  $p_{i,i} = 0$  otherwise.
  - What happens in the long run? Does  $X_n$  become uniformly distributed? Does it stay close to  $X_0$ ? to d/2?
- Example: simple discrete-time queue.
  - At each time n, one person (or internet packet or ...) gets "served", and  $Z_n$  new people arrive, where  $\{Z_n\}$  are i.i.d.  $\sim \rho$ , with  $\rho$  an "arrival distribution" on  $\{0, 1, 2, ...\}$ .
  - Let  $X_n = \#$  of people in the queue at time n.
  - Then  $X_{n+1} = X_n \min(1, X_n) + Z_n$ .
  - Here  $\{X_n\}$  is Markov chain, with  $S = \{0, 1, 2, 3, \ldots\}$ , and  $p_{ij} = \rho(j-i+\min(1,i))$ .
  - Important in many applications ...
- Example: human Markov chain!
  - Everyone take out a coin (or borrow one).
  - Then pick out two other students, one for "heads" and one for "tails".
  - Each time the frog comes to you, catch it, and flip your coin. Then
    toss the frog to either your heads or your tails student, depending on
    the result of the flip.
  - What happens in the long run?

#### **Elementary Computations:**

- Let  $\{X_n\}$  be a Markov chain, with state space S, and transition probabilities  $p_{ij}$ , and initial probabilities  $\nu_i$ .
- Recall that:
  - $\mathbf{P}(X_0 = i_0) = \nu_{i_0}$ .
  - $\mathbf{P}(X_0 = i_0, X_1 = i_1) = \nu_{i_0} p_{i_0 i_1}.$

- $\mathbf{P}(X_0 = i_0, X_1 = i_1, \dots, X_n = i_n) = \nu_{i_0} p_{i_0 i_1} p_{i_1 i_2} \dots p_{i_{n-1} i_n}.$
- etc.
- In frog example:  $P(X_0 = 1000, X_1 = 999, X_2 = 1000) = \nu_{1000} p_{1000,999} p_{999,1000} = (1)(1/3)(1/3) = 1/9$ , etc.
- Now, let  $\mu_i^{(n)} = \mathbf{P}(X_n = i)$ .
  - Then  $\mu_i^{(0)} = \nu_i$ .
- What is  $\mu_i^{(1)}$  in terms of  $\nu_i$  and  $p_{ij}$ ?

$$-\mu_j^{(1)} = \mathbf{P}(X_1 = j) = \sum_{i \in S} \mathbf{P}(X_0 = i, X_1 = j) = \sum_{i \in S} \nu_i p_{ij}.$$

- ("Law of Total Probability", "additivity", "partition")
- In matrix form:
  - Write  $\mu^{(n)} = (\mu_1^{(n)}, \mu_2^{(n)}, \mu_3^{(n)}, \ldots)$ . [row vector]
  - And write  $\mathbf{P} = (p_{ij}) = \begin{pmatrix} p_{11} & p_{12} & p_{13} & \dots \\ p_{21} & p_{22} & p_{23} & \dots \\ p_{31} & \vdots & \vdots & \ddots \end{pmatrix}$ . [matrix]
  - And write  $\nu = (\nu_1, \nu_2, \nu_3, \ldots)$ . [row vector]
  - Then  $\mu^{(1)} = \nu P = \mu^{(0)} P$ . [matrix multiplication]
- e.g. if  $S = \{1, 2, 3\}$ , and  $\mu^{(0)} = (1/7, 2/7, 4/7)$ , and

$$(p_{ij}) = \begin{pmatrix} 0 & 1/2 & 1/2 \\ 1/3 & 1/3 & 1/3 \\ 1/4 & 1/4 & 1/2 \end{pmatrix},$$

then  $\mu_2^{(1)} = \mathbf{P}(X_1 = 2) = \mu_1^{(0)} p_{12} + \mu_2^{(0)} p_{22} + \mu_3^{(0)} p_{32} = (1/7)(1/2) + (2/7)(1/3) + (4/7)(1/4) = 13/42.$ 

- Similarly,  $\mu_k^{(2)} = \sum_{i \in S} \sum_{j \in S} \nu_i p_{ij} p_{jk}$ , etc.
  - Matrix form:  $\mu^{(2)} = \mu^{(0)}PP = \mu^{(0)}P^2$
  - By induction:  $\mu^{(n)} = \mu^{(0)} P^n$ , for n = 1, 2, 3, ...
  - Convention:  $P^0 = I$  (identity). Then true for n = 0 too.
  - e.g. in frog example,  $\mu_{999}^{(2)} = \nu_{1000} p_{1000,999} p_{999,999} + \nu_{1000} p_{1000,1000} p_{1000,999} + 0 = (1)(1/3)(1/3) + (1)(1/3)(1/3) + 0 = 2/9.$
- *n*-step transitions:  $p_{ij}^{(n)} = \mathbf{P}(X_{m+n} = j \mid X_m = i)$ .
  - (Again, doesn't depend on m: time-homogeneous.)
  - $p_{ij}^{(1)} = p_{ij}$ . (of course)
  - What about  $p_{ij}^{(2)}$ ?
  - Well,  $p_{ij}^{(2)} = \mathbf{P}(X_2 = j \mid X_0 = i) = \sum_{k \in S} \mathbf{P}(X_2 = j, X_1 = k \mid X_0 = i) = \sum_{k \in S} p_{ik} p_{kj}$ .

- Matrix form:  $P^{(2)} = (p_{ij}^{(2)}) = PP = P^2$ .
- By induction:  $P^{(n)} = P^n$ , i.e. to compute probabilities of *n*-step jumps, you can take  $n^{\text{th}}$  powers of the transition matrix P.
- Convention:  $P^{(0)} = I = \text{identity matrix}$ , i.e.  $p_{ij}^{(0)} = \delta_{ij} = \begin{cases} 1, & i = j \\ 0, & \text{otherwise} \end{cases}$
- Observation:  $p_{ij}^{(m+n)} = \mathbf{P}(X_{m+n} = j \mid X_0 = i) = \sum_{k \in S} \mathbf{P}(X_{m+n} = j, X_m = k \mid X_0 = i) = \sum_{k \in S} p_{ik}^{(m)} p_{kj}^{(n)}$ .
  - Matrix form:  $P^{(m+n)} = P^{(m)}P^{(n)}$ .
  - (Of course, since  $P^{(m+n)} = P^{m+n} = P^m P^n$ .)
  - "Chapman-Kolmogorov equations".
  - Follows that e.g.  $p_{ij}^{(m+n)} \ge p_{ik}^{(m)} p_{kj}^{(n)}$  for any state k.

#### Classification of States:

- Shorthand: write  $\mathbf{P}_i(\cdots)$  for  $\mathbf{P}(\cdots | X_0 = i)$ . And, write  $\mathbf{E}_i(\cdots)$  for  $\mathbf{E}(\cdots | X_0 = i)$ .
- Defn: a state i of a Markov chain is <u>recurrent</u> (or, <u>persistent</u>) if  $\mathbf{P}_i(X_n = i$  for some  $n \geq 1$ ) = 1. Otherwise, i is <u>transient</u>. (Previous examples? Frog? s.r.w.?)
- Let  $N(i) = \#\{n \ge 1 : X_n = i\} = \text{total } \# \text{ times the chain hits } i$ . (Random variable; could be infinite.)
- RECURRENCE THEOREM:
  - *i* recurrent  $\underline{\text{iff}} \sum_{n=1}^{\infty} p_{ii}^{(n)} = \infty \ \underline{\text{iff}} \ \mathbf{P}_i(N(i) = \infty) = 1.$
  - And, i transient iff  $\sum_{n=1}^{\infty} p_{ii}^{(n)} < \infty$  iff  $\mathbf{P}_i(N(i) = \infty) = 0$ .
- To prove this, let  $f_{ij} = \mathbf{P}_i(X_n = j \text{ for } \underline{\text{some}} \ n \ge 1)$ .
- Then i recurrent iff  $f_{ii} = 1$ .
  - And, i transient <u>iff</u>  $f_{ii} < 1$ .
- Also,  $\mathbf{P}_i(N(i) \ge 1) = f_{ii}$ , and  $\mathbf{P}_i(N(i) \ge 2) = (f_{ii})^2$ , etc.
  - In general, for  $k = 0, 1, 2, ..., \mathbf{P}_i(N(i) \ge k) = (f_{ii})^k$ .
- ullet Also, recall that if Z is any non-negative-integer-valued random variable, then

$$\sum_{k=1}^{\infty} \mathbf{P}(Z \ge k) = \mathbf{E}(Z).$$

PROOF OF RECURRENCE THEOREM: First, by continuity of probabilities,

$$\mathbf{P}_{i}(N(i) = \infty) = \lim_{k \to \infty} \mathbf{P}_{i}(N(i) \ge k) = \lim_{k \to \infty} (f_{ii})^{k} = \begin{cases} 1, & f_{ii} = 1\\ 0, & f_{ii} < 1 \end{cases}$$

Second, using countable linearity,

$$\sum_{n=1}^{\infty} p_{ii}^{(n)} = \sum_{n=1}^{\infty} \mathbf{P}_i(X_n = i) = \sum_{n=1}^{\infty} \mathbf{E}_i(\mathbf{1}_{X_n = i})$$

$$= \mathbf{E}_i \left(\sum_{n=1}^{\infty} \mathbf{1}_{X_n = i}\right) = \mathbf{E}_i \left(N(i)\right) = \sum_{k=1}^{\infty} \mathbf{P}_i \left(N(i) \ge k\right)$$

$$= \sum_{k=1}^{\infty} (f_{ii})^k = \begin{cases} \infty, & f_{ii} = 1\\ \frac{f_{ii}}{1 - f_{ii}} < \infty, & f_{ii} < 1 \end{cases} \quad Q.E.D.$$

- EXAMPLE:  $S = \{1, 2, 3, 4\}$ , and  $P = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1/4 & 1/4 & 1/2 & 0 \\ 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 1/3 & 2/3 \end{pmatrix}$ .
  - Here  $f_{11} = 1$ ,  $f_{22} = 1/4$ ,  $f_{33} = 1$ , and  $f_{44} = 1$ .
  - So, states 1, 3, and 4 are recurrent, but state 2 is transient.
  - Also,  $f_{12} = 0 = f_{13} = f_{14} = f_{32} = f_{31}$ .
  - And,  $f_{34} = 1 = f_{43}$ .
  - And,  $f_{21} = 1/3$  [since e.g.  $f_{21} = p_{21} + p_{22}f_{21} + p_{23}f_{31} + p_{24}f_{41} = (1/4) + (1/4)f_{21} + 0 + 0$ , so  $f_{21} = (1/4)/(3/4) = 1/3$ ; alternatively, in this special case only,  $f_{21} = \mathbf{P}_2(X_1 = 1 \mid X_1 \neq 2) = (1/4)/[(1/4) + (1/2)] = 1/3$ ].
  - And,  $f_{23} = 2/3$ , and  $f_{24} = 2/3$ , etc.
  - (Harder example to come on homework!)
- What about e.g. Frog Example? Harder. Later!
- EXAMPLE: Simple random walk (s.r.w.). ( $S = \mathbf{Z}$ , and  $p_{i,i+1} = p$ , and  $p_{i,i-1} = 1 p$ .)
  - Is the state 0 recurrent?
  - Well, if *n* odd, then  $p_{00}^{(n)} = 0$ .
  - If n even, then  $p_{00}^{(n)} = \mathbf{P}(n/2 \text{ heads and } n/2 \text{ tails on first } n \text{ tosses})$ =  $\binom{n}{n/2} p^{n/2} (1-p)^{n/2} = \frac{n!}{[(n/2)!]^2} p^{n/2} (1-p)^{n/2}$ . [binomial distribution]
  - Stirling's approximation: if n large, then  $n! \approx (n/e)^n \sqrt{2\pi n}$
  - So, for n large and even,

$$p_{00}^{(n)} \approx \frac{(n/e)^n \sqrt{2\pi n}}{[(n/2e)^{n/2} \sqrt{2\pi n/2}]^2} p^{n/2} (1-p)^{n/2}$$
  
=  $[4p(1-p)]^{n/2} \sqrt{2/\pi n}$ .

- Now, if p = 1/2, then 4p(1-p) = 1, so  $\sum_{n=1}^{\infty} p_{00}^{(n)} \approx \sum_{n=2,4,6,\dots} \sqrt{2/\pi n} = \infty$ , so state 0 is <u>recurrent</u>.

- But if  $p \neq 1/2$ , then 4p(1-p) < 1, so  $\sum_{n=1}^{\infty} p_{00}^{(n)} \approx \sum_{n=2,4,6,\dots} [4p(1-p)]^{n/2} \sqrt{2/\pi n} < \infty$ , so state 0 is <u>transient</u>.
- (Similarly true for all other states besides 0, too.)

## **Communicating States:**

- Say that state i communicates with state j, written  $i \to j$ , if  $f_{ij} > 0$ , i.e. if it is <u>possible</u> to get from i to j, i.e. if  $\exists m \geq 1$  with  $p_{ij}^{(m)} > 0$ .
  - Write  $i \leftrightarrow j$  if both  $i \to j$  and  $j \to i$ .
- Say a Markov chain is <u>irreducible</u> if  $i \to j$  for all  $i, j \in S$ . (Previous examples?)
- CASES THEOREM: For an <u>irreducible</u> Markov chain, either (a)  $\sum_{n=1}^{\infty} p_{ij}^{(n)} = \infty$  for all  $i, j \in S$ , so <u>all</u> states are recurrent. ("recurrent Markov chain") or (b)  $\sum_{n=1}^{\infty} p_{ij}^{(n)} < \infty$  for all  $i, j \in S$ , so <u>all</u> states are transient. ("transient Markov chain")
- This follows immediately from:
- SUM LEMMA: if  $i \to k$ , and  $\ell \to j$ , and  $\sum_{n=1}^{\infty} p_{k\ell}^{(n)} = \infty$ , then  $\sum_{n=1}^{\infty} p_{ij}^{(n)} = \infty$ .
- PROOF OF SUM LEMMA: Find  $m, r \ge 1$  with  $p_{ik}^{(m)} > 0$  and  $p_{\ell j}^{(r)} > 0$ . Note that  $p_{ij}^{(m+s+r)} \ge p_{ik}^{(m)} p_{k\ell}^{(s)} p_{\ell j}^{(r)}$ . Hence,

$$\sum_{n=1}^{\infty} p_{ij}^{(n)} \ge \sum_{n=m+r+1}^{\infty} p_{ij}^{(n)} = \sum_{s=1}^{\infty} p_{ij}^{(m+s+r)} \ge \sum_{s=1}^{\infty} p_{ik}^{(m)} p_{k\ell}^{(s)} p_{\ell j}^{(r)}$$

$$= p_{ik}^{(m)} p_{\ell j}^{(r)} \sum_{s=1}^{\infty} p_{k\ell}^{(s)} = (positive)(positive)(\infty) = \infty. \quad Q.E.D.$$

- EXAMPLE: simple random walk. Irreducible!
  - p = 1/2: case (a).
  - $p \neq 1/2$ : case (b).
- What about Frog Example? Also irreducible, but which case?? Answer given by:
- FINITE SPACE THEOREM: an irreducible Markov chain on a <u>finite</u> state space always falls into case (a), i.e.  $\sum_{n=1}^{\infty} p_{ij}^{(n)} = \infty$  for all  $i, j \in S$ , and all states are recurrent.
- PROOF OF FINITE SPACE THEOREM:
  - Choose any state  $i \in S$ . Then

$$\sum_{j \in S} \sum_{n=1}^{\infty} p_{ij}^{(n)} = \sum_{n=1}^{\infty} \sum_{j \in S} p_{ij}^{(n)} = \sum_{n=1}^{\infty} 1 = \infty.$$

- Since S is finite, there must be at least one  $j \in S$  with  $\sum_{i=1}^{\infty} p_{ij}^{(n)} = \infty$ .
- So, we must be in case (a).
- So, in Frog Example,  $\mathbf{P}(\exists n \geq 1 \text{ with } X_n = 1000 \mid X_0 = 1000) = 1.$ 
  - But what about  $P(\exists n \ge 1 \text{ with } X_n = 428 \,|\, X_0 = 1000)$ ??
- To continue, define  $T_i = \min\{n \geq 1 : X_n = i\}$ .  $(T_i = \infty \text{ if } \underline{\text{never}} \text{ hit } i.)$
- HIT LEMMA: If  $j \to i$  with  $j \neq i$ , then  $\mathbf{P}_i(T_i < T_i) > 0$ .
  - Intuitively obvious(?). But formal proof is:
  - Since  $j \to i$ , there is some possible path from j to i, i.e. there is  $m \in \mathbf{N}$  and  $x_0, x_1, \dots, x_m$  with  $x_0 = j$  and  $x_m = i$  and  $p_{x_r x_{r+1}} > 0$  for all 0 < r < m - 1.
  - Let  $S = \max\{r : x_r = j\}$  be the last time this path hits j.
  - Then  $x_S, x_{S+1}, \ldots, x_m$  is a possible path which goes from j to i without first returning to j.
  - So,  $P_j(T_i < T_j) \ge P_j(\text{this path}) = p_{x_S x_{S+1}} p_{x_{S+1} x_{S+2}} \dots p_{x_{m-1} x_m} > 0$ , Q.E.D.
- F-LEMMA: If  $j \to i$  and  $f_{ij} = 1$ , then  $f_{ij} = 1$ .
- PROOF OF F-LEMMA:
  - Assume  $i \neq j$  (otherwise trivial).
  - Since  $j \to i$ ,  $\mathbf{P}_i(T_i < T_i) > 0$  by Hit Lemma.
  - But  $1 f_{ij} = \mathbf{P}_i(T_i = \infty) \ge \mathbf{P}_i(T_i < T_i) \ \mathbf{P}_i(T_i = \infty)$  $= \mathbf{P}_{i}(T_{i} < T_{i}) (1 - f_{ij}).$
  - If  $f_{jj} = 1$ , then  $1 f_{jj} = 0$ , so must have  $1 f_{ij} = 0$ , i.e.  $f_{ij} = 1$ . Q.E.D.

#### END OF WEEK #2 -

- Putting all of the above together, we obtain:
- STRONGER RECURRENCE THEOREM: If chain irreducible, then the following are equivalent (and all correspond to "case (a)"):

  - (1) There are  $k, \ell \in S$  with  $\sum_{n=1}^{\infty} p_{k\ell}^{(n)} = \infty$ . (2) For all  $i, j \in S$ , we have  $\sum_{n=1}^{\infty} p_{ij}^{(n)} = \infty$ .
  - (3) There is  $k \in S$  with  $f_{kk} = 1$ , i.e. with k recurrent.
  - (4) For all  $j \in S$ , we have  $f_{jj} = 1$ , i.e. all states are recurrent.
  - (5) For all  $i, j \in S$ , we have  $f_{ij} = 1$ .
- PROOF:
  - (1)  $\Rightarrow$  (2): Sum Lemma.
  - (2)  $\Rightarrow$  (3): Recurrence Theorem (with i = j = k).

- (3)  $\Rightarrow$  (1): Recurrence Theorem (with  $\ell = k$ ).
- (2)  $\Rightarrow$  (4): Recurrence Theorem (with i = j).
- (4)  $\Rightarrow$  (5): F-Lemma.
- (5)  $\Rightarrow$  (3): Immediate.
- Q.E.D.
- Frog Example:  $P(\exists n \ge 1 \text{ with } X_n = 428 \,|\, X_0 = 1000) = 1$ , etc.
- Simple random walk with p = 1/2:  $\mathbf{P}(\exists n \geq 1 \text{ with } X_n = 1,000,000 \mid X_0 = 0) = 1$ , etc. (And similarly for any conceivable pattern of values, i.e. the chain's values "fluctuate" arbitrarily..)
- Example:  $S = \{1, 2, 3\}$ , and  $(p_{ij}) = \begin{pmatrix} 0 & 1/2 & 1/2 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$ .
  - Then  $\sum_{n=1}^{\infty} p_{12}^{(n)} = \sum_{n=1}^{\infty} (1/2) = \infty$ .
  - And  $f_{22} = 1$ . Recurrent!
  - But  $f_{11} = 0 < 1$ . Transient!
  - Also  $f_{12} = 1/2 < 1$ .
  - Not irreducible!
- Example: Simple random walk with p > 1/2.
  - Irreducible.
  - $-f_{00} < 1$ . (transient)
  - Claim:  $f_{05} = 1$ . Contradiction? No!
  - Indeed, let  $Z_n = X_n X_{n-1}$ .
  - Then  $P(Z_n = +1) = p$ ,  $P(Z_n = -1) = 1 p$ , and  $\{Z_n\}$  i.i.d.
  - So, by Strong Law of Large Numbers, w.p. 1,  $\lim_{n\to\infty} \frac{1}{n}(Z_1 + Z_2 + ... + Z_n) = \mathbf{E}(Z_1) = p(1) + (1-p)(-1) = 2p 1 > 0.$
  - So, w.p. 1,  $\lim_{n\to\infty} (Z_1 + Z_2 + \ldots + Z_n) = +\infty$ .
  - i.e., w.p.  $1, X_n X_0 \to \infty$ , so  $X_n \to \infty$ .
  - Follows that if i < j, then  $f_{ij} = 1$  (since must pass j when going from i to  $\infty$ ).
  - In particular,  $f_{05} = 1$ .
  - (Similarly, if p < 1/2 and i > j, then  $f_{ij} = 1$ .)

## **Stationary Distributions:**

- What about a Markov chain's long-run probabilities?
  - Does  $\lim_{n\to\infty} \mathbf{P}[X_n = i]$  exist?
  - What does it equal?
- Let  $\pi$  be a probability distribution on S, i.e.  $\pi_i \geq 0$  for all  $i \in S$ , and  $\sum_{i \in S} \pi_i = 1$ .
- Defn:  $\pi$  is stationary for a Markov chain  $P = (p_{ij})$  if  $\sum_{i \in S} \pi_i p_{ij} = \pi_j$  for all  $j \in S$ .
  - Matrix notation:  $\pi P = \pi$ .
  - Then by induction,  $\pi P^n = \pi$  for n = 0, 1, 2, ..., i.e.  $\sum_{i \in S} \pi_i p_{ij}^{(n)} = \pi_j$ .
  - Intuition, if chain starts with probabilities  $\{\pi_i\}$ , then chain will keep the same probabilities one time unit later.
  - That is, if  $\mu^{(n)} = \pi$ , i.e.  $\mathbf{P}(X_n = i) = \pi_i$  for all i, then  $\mu^{(n+1)} = \mu^{(n)}P = \pi P = \pi$ , i.e.  $\mu^{(n+1)}$  also equals  $\pi$ .
  - And then, by induction,  $\mu^{(m)} = \pi$  for all  $m \ge n$ . ("stationary")
- Frog Example:
  - Let  $\pi_i = \frac{1}{1000}$  for all  $i \in S$ .
  - Then  $\pi_i \geq 0$  and  $\sum_i \pi_i = 1$ .
  - Also, for all  $j \in S$ ,  $\sum_{i \in S} \pi_i p_{ij} = \frac{1}{1000} (\frac{1}{3} + \frac{1}{3} + \frac{1}{3}) = \frac{1}{1000} = \pi_j$ .
  - So,  $\{\pi_i\}$  is stationary distribution!
- (More generally, if chain is "doubly stochastic", i.e.  $\sum_{i \in S} p_{ij} = 1$  for all  $j \in S$ , and if  $\pi_i = 1/|S|$  for all  $i \in S$ , then  $\{\pi_i\}$  is stationary [check].)
- Ehrenfest's Urn example:  $(S = \{0, 1, 2, \dots, d\}, p_{ij} = i/d \text{ for } j = i-1, p_{ij} = (d-i)/d \text{ for } j = i+1)$ 
  - Does  $\pi_i = \frac{1}{d+1}$  for all i?
  - Well, if e.g. j = 1, then  $\sum_{i \in S} \pi_i p_{ij} = \frac{1}{d+1} (p_{01} + p_{21}) = \frac{1}{d+1} (1 + \frac{2}{d}) \neq \frac{1}{d+1} = \pi_j$ .
  - So, should <u>not</u> take  $\pi_i = \frac{1}{d+1}$  for all i.
  - So,  $\pi_i = ???$
- Defn: a Markov chain is <u>reversible</u> (or <u>time reversible</u>, or satisfies <u>detailed balance</u>) with respect to a probability distribution  $\{\pi_i\}$  if  $\pi_i p_{ij} = \pi_j p_{ji}$  for all  $i, j \in S$ .
- PROPOSITION: if chain is reversible w.r.t.  $\{\pi_i\}$ , then  $\{\pi_i\}$  is a stationary distribution. (Converse false.)
  - PROOF: for  $j \in S$ ,  $\sum_{i \in S} \pi_i p_{ij} = \sum_{i \in S} \pi_j p_{ji} = \pi_j \sum_{i \in S} p_{ji} = \pi_j$ . Q.E.D.

- Frog Example:
  - $-\pi_i = 1/1000$
  - If  $|j-i| \le 1$  or |j-i| = 999, then  $\pi_i p_{ij} = (1/1000)(1/3) = \pi_j p_{ji}$ .
  - Otherwise both sides 0.
  - So, reversible! (easier way to check stationarity)
- Example:  $S = \{1, 2, 3\}$ ,  $p_{12} = p_{23} = p_{31} = 1$ ,  $\pi_1 = \pi_2 = \pi_3 = 1/3$ . Then  $\{\pi_i\}$  stationary (check!), but chain is <u>not</u> reversible w.r.t.  $\{\pi_i\}$ .
- Ehrenfest's Urn:
  - New idea: perhaps each ball is equally likely to be in either Urn.
  - That is, let  $\pi_i = 2^{-d} {d \choose i} = 2^{-d} \frac{d!}{i!(d-i)!}$
  - Then  $\pi_i \geq 0$  and  $\sum_i \pi_i = 1$ .
  - Stationary? Need to check if  $\sum_{i \in S} \pi_i p_{ij} = \pi_j$  for all  $j \in S$ . Possible but messy. (Need the Pascal's Triangle identity that  $\binom{d-1}{j-1} + \binom{d-1}{j} = \binom{d}{j}$ .) Better way?
  - Use reversibility!
  - If j = i + 1, then

$$\pi_i p_{ij} = 2^{-d} {d \choose i} \frac{d-i}{d} = 2^{-d} \frac{d!}{i!(d-i)!} \frac{d-i}{d} = 2^{-d} \frac{(d-1)!}{i!(d-i-1)!}.$$

Also

$$\pi_j p_{ji} = 2^{-d} \binom{d}{j} \frac{j}{d} = 2^{-d} \frac{d!}{j!(d-j)!} \frac{j}{d} = 2^{-d} \frac{(d-1)!}{(j-1)!(d-j)!} = 2^{-d} \frac{(d-1)!}{i!(d-i-1)!} = \pi_i p_{ij}.$$

- If j = i 1, then again  $\pi_i p_{ij} = \pi_i p_{ij}$  [check! or just interchange i and j].
- Otherwise both sides 0.
- So, reversible!
- So,  $\{\pi_i\}$  is stationary distribution!
- Intuitively,  $\pi_i$  is larger when i is close to d/2.
- But does  $\mu_i^{(n)} \to \pi_i$ ? We'll see!

# Obstacles to Convergence:

- If chain has stationary distribution  $\{\pi_i\}$ , does  $\lim_{n\to\infty} \mathbf{P}[X_n=i]=\pi_i$ ?
- Not necessarily!
- Example:  $S = \{1, 2\}$ , and  $\nu_1 = 1$ , and  $(p_{ij}) = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ .
  - If  $\pi_1 = \pi_2 = \frac{1}{2}$  (say), then  $\{\pi_i\}$  stationary (check!).
  - But  $\lim_{n\to\infty} \mathbf{P}[X_n = 1] = \lim_{n\to\infty} 1 = 1 \neq \frac{1}{2} = \pi_1$ .

- Not irreducible! ("reducible")
- Example:  $S = \{1, 2\}$ , and  $\nu_1 = 1$ , and  $(p_{ij}) = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ .
  - Again, if  $\pi_1 = \pi_2 = \frac{1}{2}$ , then  $\{\pi_i\}$  stationary (check!).
  - But  $\mathbf{P}(X_n = 1) = \begin{cases} 1, & n \text{ even} \\ 0, & n \text{ odd} \end{cases}$
  - So,  $\lim_{n\to\infty} \mathbf{P}[X_n=1]$  does not even exist!
  - "periodic"
- Defn: the <u>period</u> of a state i is the greatest common divisor of the set  $\{n \geq 1; p_{ii}^{(n)} > 0\}.$ 
  - e.g. if period of i is 2, this means that it is only possible to get from i to i in an <u>even</u> numbers of steps.
  - If period if each state is 1, say chain is "aperiodic".
- Example:  $S = \{1, 2, 3\}$ , and  $p_{12} = p_{23} = p_{31} = 1$ .
  - Then period of each state is 3.
- Example:  $S = \{1, 2, 3\}$ , and  $(p_{ij}) = \begin{pmatrix} 0 & 1/2 & 1/2 \\ 1/2 & 0 & 1/2 \\ 1/2 & 1/2 & 0 \end{pmatrix}$ .
  - Then period of state 1 is  $gcd\{2,3,\ldots\}=1$ .
  - Aperiodic!
- Observation: if  $p_{ii} > 0$ , then period of i is 1 (since  $gcd\{1, ...\} = 1$ ).
  - (Converse false, as in previous example.)
  - Or, if there is some  $n \ge 1$  with  $p_{ii}^{(n)} > 0$  and  $p_{ii}^{(n+1)} > 0$ , then period of i is 1 (since  $gcd\{n, n+1, \ldots\} = 1$ ).
- Frog Example:  $p_{ii} > 0$ , so chain aperiodic.
- Simple Random Walk: can only return after <u>even</u> number of steps, so period of each state is 2.
- Ehrenfest's Urn: again, can only return after <u>even</u> number of steps, so period of each state is 2.
- EQUAL PERIODS LEMMA: if  $i \leftrightarrow j$ , then the periods of i and of j are equal.
- PROOF:
  - Let the periods of i and j be  $t_i$  and  $t_j$ .
  - Find  $r, s \in \mathbf{N}$  with  $p_{ij}^{(r)} > 0$  and  $p_{ii}^{(s)} > 0$ .
  - Then  $p_{ii}^{(r+s)} \ge p_{ij}^{(r)} p_{ji}^{(s)} > 0$ , so  $t_i$  divides r + s.
  - Also if  $p_{jj}^{(n)} > 0$ , then  $p_{ii}^{(r+n+s)} \ge p_{ij}^{(r)} p_{jj}^{(n)} p_{ji}^{(s)} > 0$ , so  $t_i$  divides r+n+s, hence  $t_i$  divides n.

- So,  $t_i$  is a common divisor of  $\{n \in \mathbf{N}; \ p_{jj}^{(n)} > 0\}.$
- So,  $t_j \ge t_i$  (since  $t_j$  is greatest common divisor).
- Similarly,  $t_i \ge t_j$ , so  $t_i = t_j$ . Q.E.D.
- COR: if chain <u>irreducible</u>, then all states have same period.
- COR: if chain irreducible and  $p_{ii} > 0$  for <u>some</u> state i, then chain is aperiodic.

## Markov Chain Convergence Theorem:

- MARKOV CHAIN CONVERGENCE THEOREM: If a Markov chain is irreducible, and aperiodic, and has a stationary distribution  $\{\pi_i\}$ , then  $\lim_{n\to\infty} p_{ij}^{(n)} = \pi_j$  for all  $i, j \in S$ , and  $\lim_{n\to\infty} \mathbf{P}(X_n = j) = \pi_j$  for any initial probabilities  $\{\nu_i\}$ .
- To prove this (big) theorem, we need some lemmas.
- STATIONARY RECURRENCE LEMMA: If chain irreducible, and has stationary dist, then it is <u>recurrent</u>.
- PROOF:
  - Suppose the chain is <u>not</u> recurrent.
  - Then by Stronger Recurrence Theorem, for all  $i, j \in S$ ,  $\sum_{n=1}^{\infty} p_{ij}^{(n)} < \infty$ .
  - Hence,  $\lim_{n\to\infty} p_{ij}^{(n)} = 0$ .
  - But  $\pi_j = \sum_{i \in S} \pi_i p_{ij}^{(n)}$  for all n.
  - Since  $\sum_{i \in S} \sup_n |\pi_i p_{ij}^{(n)}| \leq \sum_{i \in S} \pi_i = 1 < \infty$ , and each term  $\pi_i p_{ij}^{(n)} \to 0$ , it follows from the "Dominated Convergence Theorem" or "Weierstrass M-test" that as  $n \to \infty$ ,  $\sum_{i \in S} \pi_i p_{ij}^{(n)} \to 0$  as well.
  - This implies that  $\pi_j = 0$  for all  $j \in S$ .
  - But we must have  $\sum_{j\in S} \pi_j = 1$ . Impossible!
  - So, the chain must be recurrent. Q.E.D.
- ASIDE: THE (WEIERSTRASS) M-TEST (our version, anyway; also follows from the Dominated Convergence Theorem).
  - THM: If  $\lim_{n\to\infty} b_{nk} = a_k \ \forall k$ , and  $\sum_{k=1}^{\infty} \sup_{n} |b_{nk}| < \infty$ , then  $\lim_{n\to\infty} \sum_{k=1}^{\infty} b_{nk} = \sum_{k=1}^{\infty} a_k$ .
  - PROOF:
  - Let  $\epsilon > 0$ .
  - Note that  $a_k \leq \sup_n b_{nk}$ , so  $\sum_{k=1}^{\infty} \sup_n |b_{nk} a_k| \leq 2 \sum_{k=1}^{\infty} \sup_n |b_{nk}| < \infty$ .
  - So, can find  $K \in \mathbf{N}$  such that  $\sum_{k=K+1}^{\infty} \sup_{n} |b_{nk} a_k| < \frac{\epsilon}{2}$ .

- Then for  $1 \le k \le K$ , find  $N_k$  with  $|b_{nk} a_k| < \frac{\epsilon}{2K}$  for all  $n \ge N_k$ .
- Let  $N = \max(N_1, \ldots, N_K)$ .
- Then for  $n \geq N$ ,  $\left|\sum_{k=1}^{\infty} b_{nk} \sum_{k=1}^{\infty} a_k\right| \leq \sum_{k=1}^{\infty} \left|b_{nk} a_k\right| < K \frac{\epsilon}{2K} + \frac{\epsilon}{2} = \epsilon$ . Q.E.D.
- NUMBER THEORY LEMMA: If a set A of positive integers is non-empty, and additive (i.e.  $m + n \in A$  whenever  $m \in A$  and  $n \in A$ ), and aperiodic (i.e. gcd(A) = 1), then there is  $n_0 \in \mathbb{N}$  such that  $n \in A$  for all  $n \geq n_0$ .
- (Proof omited; see e.g. Durrett p. 51 / 2nd ed. p. 24, or Rosenthal p. 92.)
- COR: If a state i is aperiodic, and  $f_{ii} > 0$ , then there is  $n_0(i)$  such that  $p_{ii}^{(n)} > 0$  for all  $n \ge n_0(i)$ .
- PROOF: Let  $A = \{n \ge 1 : p_{ii}^{(n)} > 0\}.$ 
  - Then A is non-empty since  $f_{ii} > 0$ .
  - And, A is additive since  $p_{ii}^{(m+n)} \ge p_{ii}^{(m)} p_{ii}^{(n)}$ .
  - And, A is aperiodic by assumption.
  - Hence, the result follows from the Number Theory Lemma. Q.E.D.
- COR: If a chain is irreducible and aperiodic, then for any states  $i, j \in S$ , there is  $n_0(i, j)$  such that  $p_{ij}^{(n)} > 0$  for all  $n \ge n_0(i, j)$ .
- PROOF:
  - Find  $n_0(i)$  as above.
  - Find  $m \in \mathbf{N}$  such that  $p_{ij}^{(m)} > 0$ .
  - Then let  $n_0(i,j) = n_0(i) + m$ .
  - Then if  $n \ge n_0(i,j)$ , then  $n-m \ge n_0(i)$ , so  $p_{ij}^{(n)} \ge p_{ii}^{(n-m)} p_{ij}^{(m)} > 0$ . Q.E.D.
- PROOF OF MARKOV CHAIN CONVERGENCE THEOREM (long!):
- Define a new Markov chain  $\{(X_n, Y_n)\}_{n=0}^{\infty}$ , with state space  $\overline{S} = S \times S$ , and transition probabilities  $\overline{p}_{(ij),(k\ell)} = p_{ik}p_{j\ell}$ .
  - Intuition: new chain has two coordinates, each of which is an independent copy of the original Markov chain. ("coupling")

#### END OF WEEK #3 -

- The new chain has stationary distribution  $\overline{\pi}_{(ij)} = \pi_i \pi_j$  (because of independence).
- Furthermore,  $\overline{p}_{(ij),(k\ell)}^{(n)} > 0$  whenever  $n \ge \max[n_0(i,k), n_0(j,\ell)]$ .
  - So, new chain is <u>irreducible</u> and <u>aperiodic</u>.
- So, by Stationary Recurrence Lemma, new chain is <u>recurrent</u>.

- Choose any  $i_0 \in S$ , and set  $\tau = \inf\{n \geq 0; X_n = Y_n = i_0\}$ .
- By Stronger Recurrence Theorem,  $\overline{f}_{(ij),(i_0i_0)} = 1$ , i.e.  $\mathbf{P}_{(ij)}(\tau < \infty) = 1$ .
- Note also that if  $n \geq m$ , then

$$\mathbf{P}_{(ij)}(\tau = m, X_n = k) = \mathbf{P}_{(ij)}(\tau = m) p_{i_0,k}^{(n-m)} = \mathbf{P}_{(ij)}(\tau = m, Y_n = k).$$

• Hence, for  $i, j, k \in S$ ,

$$\begin{aligned} \left| p_{ik}^{(n)} - p_{jk}^{(n)} \right| &= \left| \mathbf{P}_{(ij)}(X_n = k) - \mathbf{P}_{(ij)}(Y_n = k) \right| \\ &= \left| \sum_{m=1}^{n} \mathbf{P}_{(ij)}(X_n = k, \ \tau = m) + \mathbf{P}_{(ij)}(X_n = k, \ \tau > n) \right| \\ &- \sum_{m=1}^{n} \mathbf{P}_{(ij)}(Y_n = k, \ \tau = m) - \mathbf{P}_{(ij)}(Y_n = k, \ \tau > n) \right| \\ &= \left| \mathbf{P}_{(ij)}(X_n = k, \ \tau > n) - \mathbf{P}_{(ij)}(Y_n = k, \ \tau > n) \right| \\ &\leq 2 \mathbf{P}_{(ij)}(\tau > n), \end{aligned}$$

which  $\to 0$  as  $n \to \infty$  since  $\mathbf{P}_{(ij)}(\tau < \infty) = 1$ .

- (Above factor of "2" not really necessary, since both terms non-negative.)
- Hence, it follows that

$$\left| p_{ij}^{(n)} - \pi_j \right| = \left| \sum_{k \in S} \pi_k \left( p_{ij}^{(n)} - p_{kj}^{(n)} \right) \right| \le \sum_{k \in S} \pi_k \left| p_{ij}^{(n)} - p_{kj}^{(n)} \right|,$$

which  $\to 0$  as  $n \to \infty$  since  $|p_{ij}^{(n)} - p_{kj}^{(n)}| \to 0$  for all  $k \in S$  (using the M-test).

• Finally, for any  $\{\nu_i\}$  (again using the M-test),

$$\lim_{n \to \infty} \mathbf{P}(X_n = j) = \lim_{n \to \infty} \sum_{i \in S} \mathbf{P}(X_0 = i, X_n = j) = \lim_{n \to \infty} \sum_{i \in S} \nu_i \, p_{ij}^{(n)}$$
$$= \sum_{i \in S} \nu_i \, \lim_{n \to \infty} p_{ij}^{(n)} = \sum_{i \in S} \nu_i \, \pi_j = \pi_j.$$

Q.E.D. (phew!)

- So, for Frog Example,  $\lim_{n\to\infty} \mathbf{P}(X_n = 428) = 1/1000$ , regardless of  $\{\nu_i\}$ .
- COR: If chain irreducible and aperiodic, then it has <u>at most one</u> stationary distribution.
  - Proof: If it has at least one, then by the above, each one must be equal to  $\lim_{n\to\infty} \mathbf{P}(X_n=j)$ , so they're all equal.
- Example:  $S = \{1, 2, 3\}$ , and  $(p_{ij}) = \begin{pmatrix} 2/3 & 1/3 & 0 \\ 1/3 & 2/3 & 0 \\ 0 & 0 & 1 \end{pmatrix}$ 
  - Stationary dist #1:  $\pi_1 = \pi_2 = 1/2$  and  $\pi_3 = 0$ .
  - Stationary dist #2:  $\pi_1 = \pi_2 = 0$  and  $\pi_3 = 1$ .
  - Stationary dist #3:  $\pi_1 = \pi_2 = 1/8$  and  $\pi_3 = 3/4$ .

- So, here, the stationary distribution is not unique!
- But chain is not irreducible.
- What about periodic chains? (e.g. s.r.w., Ehrenfest)
- PERIODIC CONVERGENCE THM: Suppose chain irreducible, with period  $b \geq 2$ , and stat dist  $\{\pi_i\}$ . Then  $\forall i, j \in S$ ,  $\lim_{n \to \infty} \frac{1}{b} \left[ p_{ij}^{(n)} + \ldots + p_{ij}^{(n+b-1)} \right] = \pi_j$ , and also

$$\lim_{n \to \infty} \frac{1}{b} \mathbf{P}[X_n = j \text{ or } X_{n+1} = j \text{ or } \dots \text{ or } X_{n+b-1} = j] = \pi_j.$$

- (Note: still have  $\lim_{n\to\infty} \frac{1}{b} \left[ p_{ij}^{(n)} + \dots + p_{ij}^{(n+b-1)} \right] = \pi_j$  for aperiodic chains, too.)
- e.g. Ehrenfest's Urn: b=2, so  $\lim_{n\to\infty}\frac{1}{2}\mathbf{P}[X_n=j \text{ or } X_{n+1}=j]=2^{-d}\binom{d}{j}$ .
- PROOF (outline only; optional):
  - Fix  $i \in S$ .
  - For r = 0, 1, 2, ..., b 1, let  $S_r = \{j \in S : p_{ij}^{(bm+r)} > 0 \text{ for some } m \in \mathbb{N}\}.$
  - Then  $S = S_0 \overset{\bullet}{\cup} S_1 \overset{\bullet}{\cup} \dots \overset{\bullet}{\cup} S_{b-1}$ . (disjoint) (partition)
  - Furthermore  $P^{(b)}$  is irreducible and aperiodic when restricted to  $S_0$ .
  - Also  $\pi(S_0) = \pi(S_1) = \ldots = \pi(S_{b-1}) = 1/b$ .
  - And,  $\{b\,\pi_i\}_{i\in S_0}$  is stationary for  $P^{(b)}$  when restricted to  $S_0$ .
  - Follows that  $\lim_{n\to\infty} p_{ij}^{(bn)} = b \,\pi_j$  for all  $j \in S_0$ .
  - Then follows that  $\lim_{n\to\infty} p_{ij}^{(bn+r)} = b \pi_j$  for all  $j \in S_r$ , for  $0 \le r \le b-1$ .
  - Hence,  $\lim_{n\to\infty} \frac{1}{b} [p_{ij}^{(n)} + p_{ij}^{(n+1)} + \dots + p_{ij}^{(n+b-1)}] = \frac{1}{b} [b \pi_j + 0] = \pi_j$  for any  $j \in S$ .
  - Then the second statement follows from the first using the M-test, just like in the main proof. Q.E.D.
- (e.g. for Ehrenfest's Urn, if i = 0, then  $S_0 = \{\text{even } i \in S\}$ , and  $S_1 = \{\text{odd } i \in S\}$ , and  $\lim_{n \to \infty} \frac{1}{2} [p_{ij}^{(n)} + p_{ij}^{(n+1)}] = \pi_j = 2^{-d} \binom{d}{j}$ .)
- COROLLARY: If Markov chain P is irreducible (not necessarily aperiodic), then it has <u>at most one</u> stationary distribution (just like before).
- What about simple random walk? Does it have a stationary dist?
  - No!
  - Know that  $p_{ii}^{(n)} \approx [4p(1-p)]^{n/2} \sqrt{2/\pi n}$ , so  $p_{ii}^{(n)} \leq \sqrt{2/\pi n} \to 0$ .
  - Then for any  $i, j \in S$ , find  $m \in \mathbf{N}$  with  $p_{ji}^{(m)} > 0$ , then  $p_{ii}^{(n+m)} \ge p_{ij}^{(n)} p_{ji}^{(m)}$ , so we must have  $p_{ij}^{(n)} \le p_{ii}^{(n+m)} / p_{ji}^{(m)} \to 0$  as well.
  - Then, if had stat dist  $\{\pi_i\}$ , then  $\forall j \in S$ ,  $\pi_j = \sum_{i \in S} \pi_i p_{ij}^{(n)} \to 0$  (using M-test).

- [Or, alternatively, would have  $\frac{1}{2}[p_{ij}^{(n)}+p_{ij}^{(n+1)}]\to\pi_j$  and also  $\frac{1}{2}[p_{ij}^{(n)}+p_{ij}^{(n+1)}]\to0.$ ]
- So, would have  $\pi_j = 0$  for all j, so  $\sum_j \pi_j = 0$ . Impossible!
- [Aside: here  $\sum_{j} p_{ij}^{(n)} = 1$  for all n, even though  $\sum_{j} \lim_{n \to \infty} p_{ij}^{(n)} = 0$ . So, M-test conditions are <u>not</u> satisfied.]
- If S is infinite, can there ever be a stationary distribution? Yes!
- Example:  $S = \mathbf{N} = \{1, 2, 3, \ldots\}$ , and for  $i \geq 2$ ,  $p_{i,i} = p_{i,i+1} = 1/4$  and  $p_{i,i-1} = 1/2$ , and  $p_{1,1} = 3/4$  and  $p_{1,2} = 1/4$ .
  - Let  $\pi_i = 2^{-i}$ , so  $\pi_i \ge 0$  and  $\sum_i \pi_i = 1$ .
  - Then for any  $i \in S$ ,  $\pi_i p_{i,i+1} = 2^{-i} (1/4) = 2^{-i-2}$ .
  - Also,  $\pi_{i+1}p_{i+1,i} = 2^{-(i+1)}(1/2) = 2^{-i-2}$ . Equal!
  - And  $\pi_i p_{i,j} = 0$  if  $|j i| \ge 2$ .
  - So reversible! So,  $\{\pi_i\}$  is stationary dist.
  - Also irreducible and aperiodic (easy).
  - So,  $\lim_{n\to\infty} P(X_n=j) = \pi_j = 2^{-j}$  for all  $j \in S$ .

# Application – Metropolis Algorithm (Markov Chain Monte Carlo) (MCMC):

- Let  $S = \mathbf{Z}$ , and let  $\{\pi_i\}$  be <u>any</u> prob dist on S. Assume  $\pi_i > 0$  for all i.
- Can we <u>create</u> Markov chain transitions  $\{p_{ij}\}$  so that  $\lim_{n\to\infty} p_{ij}^{(n)} = \pi_j$ .
- Yes! Let  $p_{i,i+1} = \frac{1}{2} \min[1, \frac{\pi_{i+1}}{\pi_i}], p_{i,i-1} = \frac{1}{2} \min[1, \frac{\pi_{i-1}}{\pi_i}],$  and  $p_{i,i} = 1 p_{i,i+1} p_{i,i-1}$ , with  $p_{ij} = 0$  otherwise.
- Equivalent algorithmic version: Given  $X_{n-1}$ , let  $Y_n$  equal  $X_{n-1} \pm 1$  (prob 1/2 each), and  $U_n \sim \text{Uniform}[0,1]$  (indep.), and

$$X_n = \begin{cases} Y_n, & U_n < \frac{\pi_{Y_n}}{\pi_{X_{n-1}}} & \text{("accept")} \\ X_{n-1}, & \text{otherwise} & \text{("reject")} \end{cases}$$

- Then  $\pi_i p_{i,i+1} = \pi_i \frac{1}{2} \min[1, \frac{\pi_{i+1}}{\pi_i}] = \frac{1}{2} \min[\pi_i, \pi_{i+1}].$
- Also  $\pi_{i+1}p_{i+1,i} = \pi_{i+1} \frac{1}{2} \min[1, \frac{\pi_i}{\pi_{i+1}}] = \frac{1}{2} \min[\pi_{i+1}, \pi_i].$
- So  $\pi_i p_{ij} = \pi_j p_{ji}$  if j = i + 1, hence for all  $i, j \in S$ .
- So, chain is <u>reversible</u> w.r.t.  $\{\pi_i\}$ , so  $\{\pi_i\}$  stationary.
- Also irreducible and aperiodic (easy).
- So,  $\lim_{n \to \infty} p_{ij}^{(n)} = \pi_j$ , i.e.  $\lim_{n \to \infty} \mathbf{P}[X_n = j] = \pi_j$ . Q.E.D.
- Widely used to <u>sample from</u> complicated distributions  $\{\pi_i\}$ , and thus estimate their probability / expected values / etc.
  - [ \*\* Animated version available at: www.probability.ca/met \*\* ]
- Also works on continuous state spaces, with  $\pi$  a density function (e.g. the

Bayesian posterior density).

- "markov chain monte carlo" gives 635,000 hits in Google!

## Application - Random Walks on Graphs:

- Let V be a non-empty finite or countable set.
- Let  $w: V \times V \to [0, \infty)$  be a symmetric weight function (i.e. w(u, v) = w(v, u)).
  - Usual (unweighted) case: w(u, v) = 1 if there is an edge between u and v, otherwise w(u, v) = 0. (diagram)
  - Or can have other weights, multiple edges, self-loops (w(u, u) > 0), etc.
- Let  $d(u) = \sum_{v \in V} w(u, v)$ . ("degree" of vertex u)
- Define a Markov chain on S = V by  $p_{uv} = \frac{w(u,v)}{d(u)}$ .
  - Check:  $\sum_{v \in V} p_{uv} = \frac{\sum_{v \in V} w(u,v)}{\sum_{v \in V} w(u,v)} = 1.$
  - "(simple) random walk on the weighted undirected graph (V, w)"
- Other examples: Irreducible? Aperiodic? Stationarity distribution?
- Example:  $V = \{1, 2, 3, 4, 5\}$ , with w(i, i + 1) = w(i + 1, i) = 1 for i = 1, 2, 3, 4, and w(5, 1) = w(1, 5) = 1, with w(i, j) = 0 otherwise. ("ring") (diagram) Irreducible? Aperiodic? Stationarity distribution?
- [Reminder: HW#1 due next class at 6:10 sharp!]

#### END OF WEEK #4 -

- Example:  $V = \{0, 1, 2, ..., K\}$ , with w(i, 0) = w(0, i) = 1 for i = 1, 2, 3, with w(i, j) = 0 otherwise. ("star") (diagram)
- Example:  $V = \{1, 2, ..., K\}$ , with w(i, i + 1) = w(i + 1, i) = 1 for  $1 \le i \le K 1$ , with w(i, j) = 0 otherwise. ("stick") (diagram)
- Example:  $V = \mathbf{Z}$ , with w(i, i + 1) = w(i + 1, i) = 1 for all  $i \in V$ , and w(i, j) = 0 otherwise.
  - Random walk on this graph corresponds to simple random walk with p = 1/2.
- Example:  $V = \{1, 2, ..., 1000\}$ , with w(i, i) = 1 for  $1 \le i \le 1000$ , and w(i, i+1) = w(i+1, i) = 1 for  $1 \le i \le 999$ , and w(1000, 1) = w(1, 1000) = 1, and w(i, j) = 0 otherwise.
  - Random walk on this graph corresponds to the Frog Example!
- Let  $Z = \sum_{u \in V} d(u) = \sum_{u,v \in V} w(u,v)$ .
  - In unweighted case,  $Z = 2 \times (\text{number of edges})$ .
  - Assume that Z is <u>finite</u> (it might not be, if V is infinite).

- And, assume that d(u) > 0 for all  $u \in V$  (so any isolated point has a self-loop), to make  $p_{uv} = \frac{w(u,v)}{d(u)}$  well-defined.
- Let  $\pi_u = \frac{d(u)}{Z}$ , so  $\pi_u \ge 0$  and  $\sum_u \pi_u = 1$ .
  - Then  $\pi_u p_{uv} = \frac{d(u)}{Z} \frac{w(u,v)}{d(u)} = \frac{w(u,v)}{Z}$ .
  - And,  $\pi_v p_{vu} = \frac{d(v)}{Z} \frac{w(v,u)}{d(v)} = \frac{w(v,u)}{Z} = \frac{w(u,v)}{Z}$ . Same!
  - So, chain is reversible w.r.t.  $\{\pi_u\}$ .
  - So,  $\{\pi_u\}$  is stationary dist.
- If graph is <u>connected</u>, then chain is irreducible.
- If graph is <u>bipartite</u> (i.e., can be divided into two subsets s.t. all links go from one to the other), then the chain has period 2.
  - Otherwise, the chain is aperiodic (since can return to u in 2 steps).
  - (i.e., 1 and 2 are the only possible periods)
- This proves: THM: for random walk on a connected non-bipartite graph, if  $Z < \infty$ , then  $\lim_{n \to \infty} p_{uv}^{(n)} = \pi_v = \frac{d(v)}{Z}$  for all  $u, v \in V$ .
  - i.e.,  $\lim_{n\to\infty} \mathbf{P}[X_n = v] = \frac{d(v)}{Z}$ .
- What about bipartite graphs? Use Periodic Convergence Thm!
  - THM: for random walk on any connected graph with  $Z < \infty$  (whether bipartite or not),  $\lim_{n\to\infty} \frac{1}{2}[p_{uv}^{(n)} + p_{uv}^{(n+1)}] = \frac{d(v)}{Z}$ .
- Example:  $V = \{1, 2, ..., K\}$ , with w(i, i + 1) = w(i + 1, i) = 1 for  $1 \le i \le K 1$ , with w(i, j) = 0 otherwise. ("stick")
  - Connected, but bipartite.
  - $p_{12} = 1$ , and  $p_{K,K-1} = 1$ , and  $p_{i,i+1} = p_{i,i-1} = 1/2$  for  $2 \le i \le K 1$ .
  - $\pi_i = \frac{1}{2K-2}$  for i = 1, K, and  $\pi_i = \frac{2}{2K-2}$  for  $2 \le i \le K-1$ .
  - Then, know that  $\lim_{n\to\infty} \frac{1}{2} [p_{ij}^{(n)} + p_{ij}^{(n+1)}] = \pi_j$  for all  $j \in V$ .
- What about star? Or, star with an extra edge between 0 and 0?

## Application – Gambler's Ruin:

- Let 0 < a < c be integers, and let 0 .
- Suppose player A starts with a dollars, and player B starts with c-a dollars.
- At each bet, A wins \$1 with prob p, or loses \$1 with prob 1-p.
- Let  $X_n$  be the amount of money A has at time n.
  - So,  $X_0 = a$ .
- Let  $T_i = \inf\{n \geq 0 : X_n = i\}$  be the first time A has i dollars.
- QUESTION: what is  $\mathbf{P}_a(T_c < T_0)$ , i.e. the prob that A reaches c dollars before reaching 0 (i.e., before losing all their money)?

- (see e.g. www.probability.ca/gamone)
- Example: What does it equal if c = 10,000, a = 9,700, and p = 0.49?
- Example: Is it higher if c = 8, a = 6, p = 1/3 ("born rich"), or if c = 8, a = 2, p = 2/3 ("born lucky")?
- Here  $\{X_n\}$  is a Markov chain (good), but there's no limit to how long the game might take (bad).
  - So, how to solve it??
- Key: write  $P_a(T_c < T_0)$  as s(a), and consider it to be a <u>function</u> of a.
  - Can we related the different unknown s(a) to each other?
- Clearly s(0) = 0, and s(c) = 1.
- Furthermore, on the <u>first</u> bet, A either wins or loses \$1.
  - So, for  $1 \le a \le c 1$ ,

$$s(a) = \mathbf{P}_a(T_c < T_0)$$

$$= \mathbf{P}_a(T_c < T_0, \ X_1 = X_0 + 1) + \mathbf{P}_a(T_c < T_0, \ X_1 = X_0 - 1)$$

$$= \mathbf{P}(X_1 = X_0 + 1) \ \mathbf{P}_a(T_c < T_0 \mid X_1 = X_0 + 1)$$

$$+ \mathbf{P}(X_1 = X_0 - 1) \ \mathbf{P}_a(T_c < T_0 \mid X_1 = X_0 - 1)$$

$$= p s(a+1) + (1-p) s(a-1).$$

- This gives c-1 equations for the c-1 unknowns.
  - Can solve using simple algebra!
- Re-arranging, p s(a) + (1-p) s(a) = p s(a+1) + (1-p) s(a-1).
  - Hence,  $s(a+1) s(a) = \frac{1-p}{p} [s(a) s(a-1)].$
  - Let x = s(1) (unknown).
  - Then s(1) s(0) = x, and  $s(2) s(1) = \frac{1-p}{p}[s(1) s(0)] = \frac{1-p}{p}x$ .
  - Then  $s(3) s(2) = \frac{1-p}{p}[s(2) s(1)] = \left(\frac{1-p}{p}\right)^2 x$ .
  - In general, for  $1 \le a \le c 1$ ,  $s(a+1) s(a) = \left(\frac{1-p}{p}\right)^a x$ .
  - Hence, for  $1 \le a \le c 1$ ,

$$s(a) = s(a) - s(0)$$

$$= (s(a) - s(a-1)) + (s(a-1) - s(a-2)) + \dots + (s(1) - s(0))$$

$$= \left( \left( \frac{1-p}{p} \right)^{a-1} + \left( \frac{1-p}{p} \right)^{a-2} + \dots + \left( \frac{1-p}{p} \right) + 1 \right) x$$

$$= \begin{cases} \left( \frac{\left( \frac{1-p}{p} \right)^a - 1}{\left( \frac{1-p}{p} \right) - 1} \right) x, & p \neq 1/2 \\ ax, & p = 1/2 \end{cases}$$

- But s(c) = 1, so we can solve for x:

$$x = \begin{cases} \frac{\left(\frac{1-p}{p}\right)-1}{\left(\frac{1-p}{p}\right)^{c}-1}, & p \neq 1/2\\ 1/c, & p = 1/2 \end{cases}$$

• We then obtain our <u>final Gambler's Ruin fo</u>rmula:

$$s(a) = \begin{cases} \frac{\left(\frac{1-p}{p}\right)^{a} - 1}{\left(\frac{1-p}{p}\right)^{c} - 1}, & p \neq 1/2 \\ a/c, & p = 1/2 \end{cases}$$

• Example: If c = 10,000, a = 9,700, p = 0.49, then

$$s(a) = \frac{\left(\frac{0.51}{0.49}\right)^{9,700} - 1}{\left(\frac{0.51}{0.49}\right)^{10,000} - 1} \doteq 0.000006134 \doteq 1/163,000.$$

• Example: If c = 8, a = 6, p = 1/3 ("born rich"),

$$s(a) = \frac{\left(\frac{2/3}{1/3}\right)^6 - 1}{\left(\frac{2/3}{1/3}\right)^8 - 1} = 63/255 \doteq 0.247.$$

• Example: If c = 8, a = 2, p = 2/3 ("born lucky"),

$$s(a) = \frac{\left(\frac{1/3}{2/3}\right)^2 - 1}{\left(\frac{1/3}{2/3}\right)^8 - 1} = (3/4) / (255/256) \doteq 0.753.$$

- So, it is better to be born lucky than rich!
- Check: is s(a) continuous as a function of p, as  $p \to 1/2$ ?

#### Martingales:

- MOTIVATION: Gambler's ruin with p = 1/2.
  - Let  $T = \inf\{n \geq 0 : X_n = 0 \text{ or } c\} = \text{time when game ends.}$
  - Then  $\mathbf{E}(X_T) = c \mathbf{P}(X_T = c) + 0 \mathbf{P}(X_T = 0) = c s(a) + 0 (1 s(a)) = c (a/c) + 0 (1 a/c) = a.$
  - So  $\mathbf{E}(X_T) = \mathbf{E}(X_0)$ , i.e. "on average it stays the same".
  - Makes sense since  $\mathbf{E}(X_{n+1} | X_n = i) = (1/2)(i+1) + (1/2)(i-1) = i$ .
  - Reverse logic: If we knew that  $\mathbf{E}(X_T) = \mathbf{E}(X_0) = a$ , then could compute that a = c s(a) + 0 (1 s(a)), so must have s(a) = a/c. (Easier solution!)
- DEFN: A sequence  $\{X_n\}_{n=0}^{\infty}$  of random variables is a <u>martingale</u> if  $\mathbf{E}|X_n| < \infty$  for each n, and also  $\mathbf{E}(X_{n+1}|X_0,...,X_n) = X_n$  (i.e., it stays same on average).

- SPECIAL CASE: If  $\{X_n\}$  is a Markov chain (with  $\mathbf{E}|X_n| < \infty$ ), then  $\mathbf{E}[X_{n+1}|X_0,...,X_n] = \sum_j j P[X_{n+1} = j|X_0,...,X_n] = \sum_j j p_{X_n,j}$ , so martingale if  $\sum_j j p_{ij} = i$  for all i.
- EXAMPLE: Let  $\{X_n\}$  be simple random walk with p = 1/2 (i.e., "simple symmetric random walk", or s.s.r.w.).
  - Martingale, since  $\sum_{j} j p_{ij} = (i+1)(1/2) + (i-1)(1/2) = i$ .
- (Optional aside: in defn of martingale, suffices to check that  $\mathbf{E}|X_n| < \infty$  for all n, and also  $\mathbf{E}(X_{n+1}|\mathcal{F}_n) = X_n$ , where  $\{\mathcal{F}_n\}$  is any nested <u>filtration</u> for  $\{X_n\}$ , i.e.  $\mathcal{F}_n$  is a sub- $\sigma$ -algebra with  $\sigma(X_0, X_1, \ldots, X_n) \subseteq \mathcal{F}_n \subseteq \mathcal{F}_{n+1}$ .)
- If  $\{X_n\}$  martingale, then it follows from "double-expectation formula" that

$$\mathbf{E}(X_{n+1}) = \mathbf{E} \big[ \mathbf{E}(X_{n+1} \mid X_0, X_1, \dots, X_n) \big] = \mathbf{E}(X_n),$$

i.e. that  $\mathbf{E}(X_n) = \mathbf{E}(X_0)$  for all n.

- But what about  $\mathbf{E}(X_T)$  for a random time T?
- DEFN: A non-negative-integer-valued random variable T is a stopping time for  $\{X_n\}$  if the event  $\{T=n\}$  is determined by  $X_0, X_1, \ldots, X_n$ .
  - i.e., can't look into <u>future</u> before deciding to stop.
  - e.g.  $T = \inf\{n \ge 0 : X_n = 5\}$  is a valid stopping time. (=  $\infty$  if never hit 5)
  - e.g.  $T = \inf\{n \ge 0 : X_n = 0 \text{ or } X_n = c\}$  is a valid stopping time.
  - e.g.  $T = \inf\{n \ge 2 : X_{n-2} = 5\}$  is a valid stopping time.
  - e.g.  $T = \inf\{n \geq 2 : X_{n-1} = 5, X_n = 6\}$  is a valid stopping time.
  - e.g.  $T = \inf\{n \ge 0 : X_{n+1} = 5\}$  is <u>not</u> a valid stopping time (since it looks into the future).
- Do we always have  $\mathbf{E}(X_T) = \mathbf{E}(X_0)$ , if T is a stopping time?
  - At least if  $P(T < \infty) = 1$ ?
- Not necessarily!
  - e.g. let  $\{X_n\}$  be s.s.r.w. with  $X_0 = 0$ . Martingale!
  - Let  $T = T_{-5} = \inf\{n \ge 0 : X_n = -5\}$ . Stopping time!
  - And,  $P(T < \infty) = 1$  since s.s.r.w. is recurrent.
  - But  $X_T = -5$ , so  $\mathbf{E}(X_T) = -5 \neq 0 = \mathbf{E}(X_0)$ .
  - What went wrong? Need some <u>boundedness</u> conditions!
- OPTIONAL STOPPING LEMMA: If  $\{X_n\}$  martingale, with stopping time T which is <u>bounded</u> (i.e.,  $\exists M < \infty$  with  $\mathbf{P}(T \leq M) = 1$ ), then  $\mathbf{E}(X_T) = \mathbf{E}(X_0)$ .
- PROOF: Using the double-expectation formula, and then the fact that

" $1 - \mathbf{1}_{T \leq k-1}$ " is completely determined by  $X_0, X_1, \ldots, X_{k-1}$  (and thus can be treated as a constant in the conditional expectation; this fact is optional), we have:

$$\mathbf{E}(X_{T}) - \mathbf{E}(X_{0}) = \mathbf{E}(X_{T} - X_{0}) = \mathbf{E}\left[\sum_{k=1}^{T} (X_{k} - X_{k-1})\right]$$

$$= \mathbf{E}\left[\sum_{k=1}^{M} (X_{k} - X_{k-1}) \mathbf{1}_{k \le T}\right] = \sum_{k=1}^{M} \mathbf{E}[(X_{k} - X_{k-1}) \mathbf{1}_{k \le T}]$$

$$= \sum_{k=1}^{M} \mathbf{E}[(X_{k} - X_{k-1}) (1 - \mathbf{1}_{T \le k-1})]$$

$$= \sum_{k=1}^{M} \mathbf{E}(\mathbf{E}[(X_{k} - X_{k-1}) (1 - \mathbf{1}_{T \le k-1}) \mid X_{0}, X_{1}, \dots, X_{k-1}])$$

$$= \sum_{k=1}^{M} \mathbf{E}(\mathbf{E}[(X_{k} - X_{k-1}) \mid X_{0}, X_{1}, \dots, X_{k-1}] (1 - \mathbf{1}_{T \le k-1}))$$

$$= \sum_{k=1}^{M} \mathbf{E}((0) (1 - \mathbf{1}_{T \le k-1})) = 0, \quad Q.E.D.$$

- Question: How does this proof break down if  $M = \infty$ ?
- Example: s.s.r.w., with  $X_0 = 0$ , and let  $T = \min \left( 10^{12}, \inf\{n \ge 0 : X_n = -5\} \right)$ .
  - Then  $T \leq 10^{12}$ , so T bounded, so  $\mathbf{E}(X_T) = \mathbf{E}(X_0) = \mathbf{E}(0) = 0$ .
  - But nearly always have  $X_T = -5$ . Contradiction??
  - No, since by the Law of Total Expectation,  $0 = \mathbf{E}(X_T) = \mathbf{P}(X_T = -5)\mathbf{E}(X_T | X_T = -5) + \mathbf{P}(X_T \neq 5)\mathbf{E}(X_T | X_T \neq 5)$ , and  $\mathbf{E}(X_T | X_T = -5) = -5$ , and  $\mathbf{P}(X_T = -5) \approx 1$ , and  $\mathbf{P}(X_T \neq 5) \approx 0$ , but the equation still holds since  $\mathbf{E}(X_T | X_T \neq 5)$  is <u>huge</u>.
- Can we apply this to the Gambler's Ruin problem?
  - No, since there T is not bounded!
  - Need something more general!
- OPTIONAL STOPPING THM: If  $\{X_n\}$  is martingale with stopping time T, and  $\mathbf{P}(T < \infty) = 1$ , and  $\mathbf{E}|X_T| < \infty$ , and  $\lim_{n \to \infty} \mathbf{E}(X_n \mathbf{1}_{T>n}) = 0$ , then  $\mathbf{E}(X_T) = \mathbf{E}(X_0)$ .
- PROOF:
  - Let  $S = \min(T, n)$ . Stopping time! Bounded!
  - Then by Optional Stopping Lemma,  $\mathbf{E}(X_S) = \mathbf{E}(X_0)$  (for any n).
  - But  $X_S = X_{\min(T,n)} = X_T X_T \mathbf{1}_{T>n} + X_n \mathbf{1}_{T>n}$ .
  - So,  $X_T = X_S + X_T \mathbf{1}_{T>n} X_n \mathbf{1}_{T>n}$ .

- So,  $\mathbf{E}(X_T) = \mathbf{E}(X_S) + \mathbf{E}(X_T \mathbf{1}_{T>n}) \mathbf{E}(X_n \mathbf{1}_{T>n})$ . (three terms to consider)
- First term =  $\mathbf{E}(X_0)$  from above.
- Second term  $\to 0$  as  $n \to \infty$  by Dominated Convergence Thm (optional), since  $\mathbf{E}|X_T| < \infty$  and  $\mathbf{1}_{T>n} \to 0$  (since  $\mathbf{P}(T < \infty) = 1$ ).
- Third term  $\to 0$  as  $n \to \infty$  by assumption.
- So,  $\mathbf{E}(X_T) \to \mathbf{E}(X_0)$ , i.e.  $\mathbf{E}(X_T) = \mathbf{E}(X_0)$ . Q.E.D.
- OPTIONAL STOPPING COROLLARY: If  $\{X_n\}$  is martingale with stopping time T, which is "bounded up to time T" (i.e.,  $\exists M < \infty$  with  $\mathbf{P}(|X_n|\mathbf{1}_{n\leq T} \leq M) = 1$  for all n), and  $\mathbf{P}(T < \infty) = 1$ , then  $\mathbf{E}(X_T) = \mathbf{E}(X_0)$ .

## • PROOF:

- It follows that  $\mathbf{P}(|X_T| \leq M) = 1$ . [Formally, this holds since  $\mathbf{P}(|X_T| > M) = \sum_n \mathbf{P}(T = n, |X_T| > M) = \sum_n \mathbf{P}(T = n, |X_n| \mathbf{1}_{n \leq T} > M) \leq \sum_n \mathbf{P}(|X_n| \mathbf{1}_{n \leq T} > M) = \sum_n (0) = 0$ .]
- Hence,  $\mathbf{E}|X_T| \leq M < \infty$ .
- Also,  $|\mathbf{E}(X_n\mathbf{1}_{T>n})| \leq \mathbf{E}(|X_n|\mathbf{1}_{T>n}) \leq \mathbf{E}(M\mathbf{1}_{T>n}) = M\mathbf{P}(T>n)$ , which  $\to 0$  as  $n \to \infty$  since  $\mathbf{P}(T<\infty) = 1$ .
- Hence, result follows from Optional Stopping Theorem. Q.E.D.
- Example: Gambler's Ruin with p = 1/2, and  $T = \inf\{n \ge 0 : X_n = 0 \text{ or } X_n = c\}$ .
  - Then  $\mathbf{P}(T < \infty) = 1$  (game must eventually end). [Formally:  $\mathbf{P}(T > mc) \le (1 p^c)^m \to 0$  as  $m \to \infty$ , since if win c times in a row then game over.]
  - Also,  $|X_n|\mathbf{1}_{n\leq T}\leq c<\infty$  for all n.
  - So, by Optional Stopping Corollary,  $\mathbf{E}(X_T) = \mathbf{E}(X_0) = a$ .
  - Hence, as before, a = c s(a) + 0 (1 s(a)), so must have s(a) = a/c. (Easier solution!)
- What about Gambler's Ruin with  $p \neq 1/2$ ?
  - Here  $\{X_n\}$  is <u>not</u> a martingale:  $\sum_j j p_{ij} = p(i+1) + (1-p)(i-1) = i + 2p 1 \neq i$ .
  - Trick: let  $Y_n = \left(\frac{1-p}{p}\right)^{X_n}$ .
  - Then  $\mathbf{E}(Y_{n+1} | Y_0, Y_1, \dots, Y_n) = p \left[ Y_n \left( \frac{1-p}{p} \right) \right] + (1-p) \left[ Y_n / \left( \frac{1-p}{p} \right) \right] = Y_n (1-p) + Y_n(p) = Y_n.$
  - So,  $\{Y_n\}$  is a martingale!
  - And,  $P(T < \infty) = 1$  as before (with the same T).

- And,  $|Y_n|\mathbf{1}_{n\leq T} \leq \max\left(\left(\frac{1-p}{p}\right)^0, \left(\frac{1-p}{p}\right)^c\right) < \infty$  for all n.
- Hence,  $\mathbf{E}(Y_T) = \mathbf{E}(Y_0) = \left(\frac{1-p}{p}\right)^a$
- But  $Y_T = \left(\frac{1-p}{p}\right)^c$  if win, or  $Y_T = \left(\frac{1-p}{p}\right)^0 = 1$  if lose.
- Hence,  $\left(\frac{1-p}{p}\right)^a = \mathbf{E}(Y_T) = s(a) \left(\frac{1-p}{p}\right)^c + [1-s(a)](1) = 1 + s(a) \left[\left(\frac{1-p}{p}\right)^c 1\right].$
- Solving,  $s(a) = \frac{\left(\frac{1-p}{p}\right)^a 1}{\left(\frac{1-p}{p}\right)^c 1}$ . (Again, easier solution!)

Reminders: No class Feb 18 (Reading Week). Midterm on Feb 25 in HA401 [last name A-P] and HA410 [last name Q-Z]) – BRING STUDENT CARD.

## END OF WEEK #5 -

(Midterm Test)

## END OF WEEK #6

- WALD'S THM: Suppose  $X_n = a + Z_1 + \ldots + Z_n$ , where  $\{Z_i\}$  are iid, with finite mean m. Let T be a stopping time for  $\{X_n\}$  which has finite mean, i.e.  $\mathbf{E}(T) < \infty$ . Then  $\mathbf{E}(X_T) = a + m \mathbf{E}(T)$ .
- Special case: if m = 0, then  $\{X_n\}$  is a martingale, and Wald's Thm says that  $\mathbf{E}(X_T) = a = \mathbf{E}(X_0)$ , as usual.
- Example:  $\{X_n\}$  is s.s.r.w. with  $X_0 = 0$ , and  $T = \inf\{n \ge 0 : X_n = -5\}$ .
  - Then  $\mathbf{E}(X_T) = -5 \neq 0 = \mathbf{E}(X_0)$ .
  - Contradiction?? No; it turns out that here  $\mathbf{E}(T) = \infty$ .
- PROOF: We compute (using that  $Z_i$  indep of  $\{T \geq i\} = \{T \leq i 1\}^C$ ) that

$$\mathbf{E}(X_T) - a = \mathbf{E}(X_T - a) = \mathbf{E}(Z_1 + \dots + Z_T)$$

$$= \mathbf{E}\left[\sum_{i=1}^T Z_i\right] = \mathbf{E}\left[\sum_{i=1}^\infty Z_i \mathbf{1}_{T \ge i}\right] = \sum_{i=1}^\infty \mathbf{E}\left[Z_i \mathbf{1}_{T \ge i}\right]$$

$$= \sum_{i=1}^\infty \mathbf{E}[Z_i] \mathbf{E}[\mathbf{1}_{T \ge i}] = m \sum_{i=1}^\infty \mathbf{P}[T \ge i] = m \mathbf{E}(T), \quad Q.E.D.$$

- (Optional aside: the above calculation uses the Dominated Convergence Thm; indeed, setting  $Y = \sum_{i=1}^{\infty} |Z_i| \mathbf{1}_{T \geq i}$ , we have  $\mathbf{E}(Y) = \mathbf{E}[\sum_{i=1}^{\infty} |Z_i| \mathbf{1}_{T \geq i}] = \sum_{i=1}^{\infty} \mathbf{E}[|Z_i| \mathbf{1}_{T \geq i}] = \sum_{i=1}^{\infty} \mathbf{E}[Z_i| \mathbf{E}[\mathbf{1}_{T \geq i}]] = \mathbf{E}|Z_1| \sum_{i=1}^{\infty} \mathbf{P}[T \geq i] = \mathbf{E}|Z_1| \mathbf{E}(T) < \infty.$
- EXAMPLE: Gambler's Ruin with  $p \neq 1/2$ , and  $T = \inf\{n \geq 0 : X_n = 0 \text{ or } c\}$ . (see e.g. www.probability.ca/gamone)
  - What is  $\mathbf{E}(T)$  = expected number of bets in the game?

- Well, here  $m = \mathbf{E}(Z_i) = p(1) + (1-p)(-1) = 2p 1$ .
- Also,  $\mathbf{E}(X_T) = c \, s(a) + 0 \, (1 s(a)) = c \, \frac{\left(\frac{1-p}{p}\right)^a 1}{\left(\frac{1-p}{p}\right)^c 1}.$
- And,  $\mathbf{E}(T) < \infty$ . [For example, this follows since  $\mathbf{P}(T \geq cn) \leq (1 p^c)^n$  so  $\mathbf{E}(T) = \sum_{i=1}^{\infty} \mathbf{P}(T \geq i) \leq \sum_{j=0}^{\infty} c \, \mathbf{P}(T \geq cj) \leq \sum_{j=0}^{\infty} c \, (1 p^c)^j = c/[1 (1 p^c)] = c/p^c < \infty$ .]
- Hence, by Wald's Thm,  $\mathbf{E}(X_T) = a + m \mathbf{E}(T)$ .
- So,  $\mathbf{E}(T) = \frac{1}{m} \left( \mathbf{E}(X_T) a \right) = \frac{1}{2p-1} \left( c \frac{\left(\frac{1-p}{p}\right)^a 1}{\left(\frac{1-p}{p}\right)^c 1} a \right).$
- e.g. p = 0.49, a = 9,700, c = 10,000:  $\mathbf{E}(T) = 484,997$ . (large!)
- But what about  $\mathbf{E}(T)$  when p = 1/2??
  - Then m=0, so the above method does not work.
- LEMMA: Let  $X_n = a + Z_1 + \ldots + Z_n$ , where  $\{Z_i\}$  i.i.d. with mean 0 and variance  $v < \infty$ . Let  $Y_n = (X_n a)^2 nv = (Z_1 + \ldots + Z_n)^2 nv$ . Then  $\{Y_n\}$  is a martingale.
- PROOF:
  - Check:  $\mathbf{E}|Y_n| \leq \mathbf{Var}(X_n) + nv = 2nv < \infty$ .
  - Also, since  $Z_{n+1}$  indep of  $Z_1, \ldots, Z_n, Y_0, \ldots, Y_n$ , we have (optional)

$$\mathbf{E}[Y_{n+1} \mid Y_0, Y_1, \dots, Y_n] = \mathbf{E}[(Z_1 + \dots + Z_n + Z_{n+1})^2 - (n+1)v \mid Y_0, Y_1, \dots, Y_n]$$

$$= \mathbf{E}[(Z_1 + \dots + Z_n)^2 + (Z_{n+1})^2 + 2Z_{n+1}(Z_1 + \dots + Z_n) - nv - v \mid Y_0, Y_1, \dots, Y_n]$$

$$= \mathbf{E}[Y_n + (Z_{n+1})^2 - v + 2Z_{n+1}(Z_1 + \dots + Z_n) \mid Y_0, Y_1, \dots, Y_n]$$

$$= Y_n + v - v + 2\mathbf{E}(Z_{n+1})\mathbf{E}[Z_1 + \dots + Z_n \mid Y_0, Y_1, \dots, Y_n] = Y_n + 0, \ Q.E.D.$$

- COR: If  $\{X_n\}$  is Gambler's Ruin with p=1/2, and  $T=\inf\{n\geq 0: X_n=0 \text{ or } c\}$ , then  $\mathbf{E}(T)=\mathbf{Var}(X_T)=a(c-a)$ .
- PROOF:
  - Let  $Y_n = (X_n a)^2 n$  (since here v = 1). Martingale (by Lemma)!
  - Choose M > 0, and let  $S_M = \min(T, M)$ . Stopping time! Bounded!
  - Hence, by Optional Stopping Lemma,  $\mathbf{E}[Y_{S_M}] = \mathbf{E}[Y_0] = (a-a)^2 0 = 0$ .
  - But  $Y_{S_M} = (X_{S_M} a)^2 S_M$ , so  $\mathbf{E}(S_M) = \mathbf{E}[(X_{S_M} a)^2]$ .
  - As  $M \to \infty$ ,  $S_M \to T$  (obviously). This implies that  $\mathbf{E}(S_M) \to \mathbf{E}(T)$  [optional: by Monotone Convergence Thm], and  $\mathbf{E}[(X_{S_M} a)^2] \to \mathbf{E}[(X_T a)^2]$  [optional: by Bounded Convergence Thm, since for any n,  $(X_{S_M} a)^2 \le \max(a^2, (c a)^2) < \infty$ ].
  - Hence,  $\mathbf{E}(T) = \mathbf{E}[(X_T a)^2] = \mathbf{Var}(X_T)$  (since  $\mathbf{E}(X_T) = a$ ).

- But  $\mathbf{Var}(X_T) = (a/c)(c-a)^2 + (1-a/c)a^2 = (a/c)(c^2+a^2-2ac) + (a^2-a^3/c) = ac+a^3/c-2a^2+a^2-a^3/c = ac-a^2 = a(c-a), Q.E.D.$
- e.g. c = 10,000, a = 9,700, p = 1/2:  $\mathbf{E}(T) = a(c-a) = 2,910,000$ . (even larger!)

## Martingale Convergence Theorem:

- EXAMPLE: Let  $\{X_n\}$  be a Markov chain on  $S = \{2^m : m \in \mathbf{Z}\}$ , with  $X_0 = 1$ , and  $p_{i,2i} = 1/3$  and  $p_{i,i/2} = 2/3$  for  $i \in S$ .
  - Martingale, since  $\sum_{j} j p_{ij} = (2i)(1/3) + (i/2)(2/3) = i$ .
  - What happens in the long run?
  - Trick: let  $Y_n = \log_2 X_n$ . Then  $Y_0 = 0$ , and  $\{Y_n\}$  is s.r.w. with p = 1/3, so  $Y_n \to -\infty$  w.p. 1.
  - Hence,  $X_n = 2^{Y_n} \to 2^{-\infty} = 0$  w.p. 1.
- EXAMPLE: Let  $\{X_n\}$  be Gambler's Ruin with p=1/2. Then  $X_n \to X$  w.p. 1, where  $\mathbf{P}(X=c)=a/c$  and  $\mathbf{P}(X=0)=1-a/c$ .
- MARTINGALE CONVERGENCE THM: Any non-negative martingale  $\{X_n\}$  converges w.p. 1 to some random variable X (e.g.  $X \equiv 0$ ).
  - Intuition: since it's non-negative (i.e., bounded on one side), it can't "spread out" forever. And since it's a martingale, it can't "drift" anywhere. So eventually it has to stop somewhere.
  - Proof omitted; see e.g. Rosenthal, p. 169.
- Example: s.s.r.w. martingale, but <u>not</u> non-negative, does <u>not</u> converge.
- Example: s.s.r.w. stopped at zero martingale, non-negative, converges to zero.
- Example: s.r.w. with p = 2/3 stopped at zero non-negative, does not converge (might increase to infinity), but not a martingale.

# Application – Branching Processes:

- Let  $\mu$  be any prob dist on  $\{0, 1, 2, \ldots\}$ . ("offspring distribution")
- Have  $X_n$  individuals at time n. (e.g., people with colds)
- Start with  $X_0 = a$  individuals. Assume  $0 < a < \infty$ .
- Each of the  $X_n$  individuals at time n has a random number of offspring which is i.i.d.  $\sim \mu$ , i.e. has i children with probability  $\mu\{i\}$ . (diagram)
- That is,  $X_{n+1} = Z_{n,1} + Z_{n,2} + \ldots + Z_{n,X_n}$ , where  $\{Z_{n,i}\}_{i=1}^{X_n}$  are i.i.d.  $\sim \mu$ .
- Then  $\{X_n\}$  is Markov chain, on state space  $S = \{0, 1, 2, \ldots\}$ .
- $-p_{00}=1.$
- $p_{ij}$  is more complicated; in fact (optional),  $p_{ij} = (\mu * \mu * \dots * \mu)(j)$ , a convolution of i copies of  $\mu$ .

- Will  $X_n = 0$  for some n?
  - How can martingales help?
- Let  $m = \sum_{i} i \mu\{i\} = \text{mean of } \mu$ . ("reproductive number")
  - Assume  $0 < m < \infty$ .
  - Then  $\mathbf{E}(X_{n+1} | X_0, \dots, X_n) = \mathbf{E}(Z_{n,1} + Z_{n,2} + \dots + Z_{n,X_n} | X_0, \dots, X_n) = m X_n$ . So, by induction,  $\mathbf{E}(X_n) < \infty$  for all n.
- Let  $Y_n = X_n/m^n$ .
  - Then since  $Y_n \leftrightarrow X_n$  is one-to-one function,

$$\mathbf{E}(Y_{n+1} | Y_0, \dots, Y_n) = \mathbf{E}(\frac{X_{n+1}}{m^{n+1}} | Y_0, \dots, Y_n)$$

$$= \mathbf{E}(\frac{X_{n+1}}{m^{n+1}} | X_0, \dots, X_n) = \frac{m X_n}{m^{n+1}} = \frac{X_n}{m^n} = Y_n.$$

- And, must have each  $\mathbf{E}|Y_n| < \infty$  (since  $\mathbf{E}|X_n| < \infty$  and m > 0).
- Hence,  $\{Y_n\}$  is martingale.
- So,  $\mathbf{E}(Y_n) = \mathbf{E}(Y_0) = a$  for all n, i.e.  $\mathbf{E}(X_n/m^n) = a$ , so  $\mathbf{E}(X_n) = a \, m^n$ .
  - (This also follows from the "induction" above.)
- If m < 1, then  $\mathbf{E}(X_n) = a \, m^n \to 0$ .
  - But  $\mathbf{E}(X_n) = \sum_{k=0}^{\infty} k \mathbf{P}(X_n = k) \ge \sum_{k=1}^{\infty} \mathbf{P}(X_n = k) = \mathbf{P}(X_n \ge 1)$ .
  - Hence,  $\mathbf{P}(X_n \ge 1) \le \mathbf{E}(X_n) = a \, m^n \to 0$ , i.e.  $\mathbf{P}(X_n = 0) \to 1$ .
  - Certain extinction!
- If m > 1, then  $\mathbf{E}(X_n) \to \infty$ .
  - In this case, it turns out that  $\mathbf{P}(X_n \to \infty) > 0$ . ("flourishing")
  - But assuming  $\mu\{0\} > 0$ , still have  $\mathbf{P}(X_n \to \infty) < 1$ , indeed  $\mathbf{P}(X_n \to 0) > 0$  (e.g., if no one has any offspring at all on the first iteration: prob =  $(\mu\{0\})^a > 0$ ).
  - So, have possible extinction, but also possible flourishing.
- But what if m = 1?
  - Then  $\mathbf{E}(X_n) = \mathbf{E}(X_0) = a$  for all n.
  - In fact,  $\{X_n\}$  is a martingale, and non-negative.
  - So, by Martingale Convergence Thm, must have  $X_n \to X$  w.p. 1, for some random variable X.
  - But how can  $\{X_n\}$  converge w.p. 1? Either (a)  $\mu\{1\} = 1$ , or (b) X = 0.
  - (In all other cases,  $\{X_n\}$  would continue to fluctuate, i.e. <u>not</u> converge w.p. 1.)
  - So, if non-degenerate (i.e.,  $\mu\{1\} < 1$ ), then  $X \equiv 0$ , i.e.  $\{X_n\} \rightarrow 0$  w.p. 1.

- Certain extinction, even when m=1!

#### **Brownian Motion:**

- Let  $\{X_n\}_{n=0}^{\infty}$  be s.s.r.w., with  $X_0 = 0$ .
- Represent this as  $X_n = Z_1 + Z_2 + \ldots + Z_n$ , where  $\{Z_i\}$  are i.i.d. with  $\mathbf{P}(Z_i = +1) = \mathbf{P}(Z_i = -1) = 1/2$ .
  - That is,  $X_0 = 0$ , and  $X_{n+1} = X_n + Z_{n+1}$ .
  - Here  $\mathbf{E}(Z_i) = 0$  and  $\mathbf{Var}(Z_i) = 1$ .
- Let M be a large integer, and let  $\{Y_t^{(M)}\}$  be like  $\{X_n\}$ , except with time speeded up by a factor of M, and space shrunk down by a factor of  $\sqrt{M}$ .
  - That is,  $Y_0^{(M)} = 0$ , and  $Y_{i+1 \atop M}^{(M)} = Y_{i \atop M}^{(M)} + \frac{1}{\sqrt{M}} Z_{i+1}$ . (diagram)
  - Fill in  $\{Y_t^{(M)}\}_{t\geq 0}$  by linear interpolation. (file www.probability.ca/sta447/Rbrownian)
- Brownian motion  $\{B_t\}_{t\geq 0}$  is (intuitively) the limit as  $M\to\infty$  of  $\{Y_t^{(M)}\}$ .
- But  $Y_t^{(M)} = \frac{1}{\sqrt{M}}(Z_1 + Z_2 + \ldots + Z_{tM})$  (at least, if  $tM \in \mathbf{Z}$ ; otherwise get errors of order  $O(1/\sqrt{M})$ , which don't matter when  $M \to \infty$ ).
  - Thus,  $\mathbf{E}(Y_t^{(M)}) = 0$ , and  $\mathbf{Var}(Y_t^{(M)}) = \frac{1}{M}(tM) = t$ .
  - So, as  $M \to \infty$ , by the Central Limit Theorem,  $Y_t^{(M)} \to \text{Normal}(0,t)$ .
  - CONCLUSION:  $B_t \sim \text{Normal}(0, t)$ . ("normally distributed")
- Also, if 0 < t < s, then  $Y_s^{(M)} Y_t^{(M)} = \frac{1}{\sqrt{M}} (Z_{tM+1} + Z_{tM+2} + \ldots + Z_{sM})$  (at least, if  $tM, sM \in \mathbf{Z}$ ; otherwise get  $O(1/\sqrt{M})$  errors).
  - So,  $Y_s^{(M)} Y_t^{(M)} \to \text{Normal}(0, s t)$ , and it is <u>independent</u> of  $Y_t^{(M)}$ .
  - CONCLUSION:  $B_s B_t \sim \text{Normal}(0, s t)$ , and it's <u>independent</u> of  $B_t$ .
  - MORE GENERALLY: if  $0 \le t_1 \le s_1 \le t_2 \le s_2 \le \ldots \le t_k \le s_k$ , then  $B_{s_i} B_{t_i} \sim \text{Normal}(0, s_i t_i)$ , and  $\{B_{s_i} B_{t_i}\}_{i=1}^k$  are all independent. ("independent normal increments")
- Finally, if  $0 < t \le s$ , then  $\mathbf{Cov}(B_t, B_s) = \mathbf{E}(B_t B_s) = \mathbf{E}(B_t [B_s B_t + B_t]) = \mathbf{E}(B_t [B_s B_t]) + \mathbf{E}((B_t)^2) = \mathbf{E}(B_t) \mathbf{E}(B_s B_t) + \mathbf{E}((B_t)^2) = (0)(0) + t = t$ .
  - In general,  $\mathbf{Cov}(B_t, B_s) = \min(t, s)$ . ("covariance structure")
- DEFINITION: <u>Brownian motion</u> is a process  $\{B_t\}_{t\geq 0}$  satisfying the above properties, and with continuous sample paths (i.e., the mapping  $t \to B_t$  is continuous).
  - FACT: Such a process exists! (The above construction is <u>intuitive</u>, but a formal proof of existence requires measure theory.)
- Example: What is  $\mathbf{E}[(B_2 + B_3 + 1)^2]$ ?
  - Well,  $\mathbf{E}[(B_2 + B_3 + 1)^2] = \mathbf{E}[(B_2)^2] + \mathbf{E}[(B_3)^2] + 1^2 + 2\mathbf{E}[B_2B_3] +$

$$2\mathbf{E}[B_2(1)] + 2\mathbf{E}[B_3(1)] = 2 + 3 + 1 + 2(2) + 2(0) + 2(0) = 10.$$

- Example: What is  $Var[B_3 + B_5 + 7]$ ?
  - Well,  $\mathbf{Var}[B_3 + B_5 + 7] = \mathbf{E}[(B_3 + B_5)^2] = \mathbf{E}[(B_3)^2] + \mathbf{E}[(B_5)^2] + 2\mathbf{E}[B_3B_5] = 3 + 5 + 2(3) = 14.$
- Aside: w.p. 1, the function  $t \mapsto B_t$  is continuous everywhere, but differentiable nowhere.
- Example: Let  $\alpha > 0$ , and let  $W_t = \alpha B_{t/\alpha^2}$ .
  - Then  $W_t \sim \text{Normal}(0, \alpha^2(t/\alpha^2)) = \text{Normal}(0, t)$ . (same as for  $B_t$ )
  - Also for 0 < t < s,  $\mathbf{E}(W_t W_s) = \alpha^2 \mathbf{E}(B_{t/\alpha^2} B_{s/\alpha^2}) = \alpha^2 (t/\alpha^2) = t$ .
  - In fact,  $\{W_t\}$  has all the same properties as  $\{B_t\}$ .
  - That is,  $\{W_t\}$  "is" Brownian motion, too. ("transformation")
- If 0 < t < s, then given  $B_r$  for  $0 \le r \le t$ , what is the <u>conditional</u> distribution of  $B_s$ ?
  - Similar to above,  $B_s | B_t = B_t + (B_s B_t) | B_t = B_t + \text{Normal}(0, s t) \sim \text{Normal}(B_t, s t)$ . (i.e., given  $B_t$ ,  $B_s$  is normal with mean  $B_t$ , variance s t.)
  - So, in particular,  $\mathbf{E}[B_s | \{B_r\}_{0 \le r \le t}] = B_t$ .
  - Hence,  $\{B_t\}$  is a (continuous-time) martingale!
  - So, similar results apply just like for discrete-time martingales.
- Example: let a, b > 0, and let  $\tau = \min\{t \ge 0 : B_t = -a \text{ or } b\}$ .
  - What is  $p \equiv \mathbf{P}(B_{\tau} = b)$ ?
  - Well, here  $\{B_t\}$  is martingale, and  $\tau$  is stopping time.
  - Furthermore,  $\{B_t\}$  is bounded up to time  $\tau$ , i.e.  $|B_t|\mathbf{1}_{t\leq \tau}\leq \max(|a|,|b|)$ .
  - So, just like for discrete martingales, must have  $\mathbf{E}(B_{\tau}) = \mathbf{E}(B_0) = 0$ .
  - Hence, p(b) + (1-p)(-a) = 0, so  $p = \frac{a}{a+b}$ . (as expected)
  - But what is  $e \equiv \mathbf{E}(\tau)$ ?

#### END OF WEEK #7 -

- To continue, let  $Y_t = B_t^2 t$ .
  - Then for 0 < t < s,  $\mathbf{E}[Y_s \mid \{B_r\}_{r \le t}] = \mathbf{E}[B_s^2 s \mid \{B_r\}_{r \le t}]$ =  $\mathbf{Var}[B_s \mid \{B_r\}_{r \le t}] + (\mathbf{E}[B_s \mid \{B_r\}_{r \le t}])^2 - s$ =  $(B_t)^2 + (s - t) - s = Y_t$ .
  - By the law of iterated expectations (optional; e.g. Rosenthal, Prop 13.2.7),  $\mathbf{E}[Y_s \mid \{Y_r\}_{r \leq t}] = \mathbf{E}\left[\mathbf{E}[Y_s \mid \{B_r\}_{r \leq t}] \mid \{Y_r\}_{r \leq t}\right] = \mathbf{E}[Y_t \mid \{Y_r\}_{r \leq t}] = Y_t.$
  - So,  $\{Y_t\}$  is also a martingale!
- Back to  $\tau = \min\{t \geq 0 : B_t = -a \text{ or } b\}$ . What is  $e \equiv \mathbf{E}(\tau)$ ?

- Well, with  $Y_t = B_t^2 t$ , have  $\mathbf{E}(Y_\tau) = \mathbf{E}(B_\tau^2 \tau) = \mathbf{E}(B_\tau^2) \mathbf{E}(\tau) = pb^2 + (1-p)(-a)^2 e = \frac{a}{a+b}b^2 + \frac{b}{a+b}a^2 e = ab e$ .
- Assuming  $\mathbf{E}(Y_{\tau}) = 0$ , solve to get e = ab. (like for discrete Gambler's Ruin)
- But  $\tau$  is not bounded ...
- To justify this argument, i.e. show that  $\mathbf{E}(Y_{\tau}) = 0$ , let  $\tau_M = \min(\tau, M)$ .
  - Then  $\tau_M$  is bounded, so  $\mathbf{E}(Y_{\tau_M}) = 0$ .
  - But  $Y_{\tau_M} = B_{\tau_M}^2 \tau_M$ , so  $\mathbf{E}(\tau_M) = \mathbf{E}(B_{\tau_M}^2)$ .
  - As  $M \to \infty$ ,  $\mathbf{E}(\tau_M) \to \mathbf{E}(\tau)$  by the Monotone Convergence Thm, and  $\mathbf{E}(B_{\tau_M}^2) \to \mathbf{E}(B_{\tau}^2)$  by the Bounded Convergence Thm.
  - Therefore,  $\mathbf{E}(\tau) = \mathbf{E}(B_{\tau}^2)$ , i.e.  $\mathbf{E}(Y_{\tau}) = 0$  as above.
- Example: Suppose  $X_t = 2 + 5t + 3B_t$  for  $t \ge 0$ .
  - What are  $\mathbf{E}(X_t)$  and  $\mathbf{Var}(X_t)$  and  $\mathbf{Cov}(X_t, X_s)$ ?
  - Well,  $\mathbf{E}(X_t) = 2 + 5t$ , and  $\mathbf{Var}(X_t) = 3^2 \mathbf{Var}(B_t) = 9t$ .
  - Follows that  $X_t \sim \text{Normal}(2+5t, 9t)$ .
  - Also for 0 < t < s,  $\mathbf{Cov}(X_t, X_s) = \mathbf{E}[(X_t 5t 2)(X_s 5s 2)] = \mathbf{E}[(3B_t)(3B_s)] = 9\mathbf{E}[B_t B_s] = 9t$ .
  - Fancy notation:  $dX_t = 5 dt + 3 dB_t$ . ("diffusion")
- More generally, could have  $X_t = x_0 + \mu t + \sigma B_t$ . (file "Rbrownian")
  - Then  $dX_t = \mu dt + \sigma dB_t$ . ( $\mu = \text{"drift"}; \sigma = \text{"volatility"}; \sigma \geq 0$ )
  - Then  $\mathbf{E}(X_t) = x_0 + \mu t$ , and  $\mathbf{Var}(X_t) = \sigma^2 t$ , and  $\mathbf{Cov}(X_t, X_s) = \sigma^2 \min(s, t)$ .
  - Optional: Even more generally, could have  $dX_t = \mu(X_t) dt + \sigma(X_t) dB_t$ , where  $\mu$  and  $\sigma$  are functions, i.e. non-constant drift and volatility.

# **Application** – **Financial Modeling:**

- Common model for stock price:  $X_t = x_0 \exp(\mu t + \sigma B_t)$ .
  - i.e. if  $Y_t = \log(X_t)$ , then  $Y_t = y_0 + \mu t + \sigma B_t$ , i.e.  $dY_t = \mu dt + \sigma dB_t$ .
  - That is, changes occur <u>proportional</u> to total price (makes sense).
  - So,  $Y_t = \log(X_t)$  is a <u>diffusion</u>.
- Also assume a risk-free interest rate r, so that \$1 today is worth  $\$e^{rt}$  a time t later.
  - Equivalently, \$1 at a future time t > 0 is worth  $\$e^{-rt}$  at time 0 (i.e. "today").
  - So, "discounted" stock price (in "today's dollars") is

$$D_t \equiv e^{-rt} X_t = e^{-rt} x_0 \exp(\mu t + \sigma B_t) = x_0 \exp((\mu - r)t + \sigma B_t).$$

- Defn: A "European call option" is the <u>option</u> to buy the stock for some amount K at some fixed future time S > 0?
  - At time S, this is worth  $\max(0, X_S K)$ .
  - At time 0, it's worth only  $e^{-rS} \max(0, X_S K)$ .
  - But at time 0,  $X_S$  is unknown (random).
- QUESTION: what is the "fair price" of this option?
  - This means the fair "no-arbitrage" price, i.e. a price such that you cannot make a guaranteed profit by buying or selling the option, combined with buying and selling the stock.
  - Note: this assumes the ability to buy/sell arbitrary amounts of stock at any time, infinitely often, including going negative (i.e., "shorting" the stock), with <u>no</u> transaction fees.
  - So, what is the fair price at time 0?
  - Is it simply the expected value,  $\mathbf{E}[e^{-rS} \max(0, X_S K)]$ ?
  - No! This would allow for arbitrage!
- FACT: the fair price for the option equals  $\mathbf{E}[e^{-rS} \max(0, X_S K)]$ , but only after replacing  $\mu$  by  $r \frac{\sigma^2}{2}$ .
  - i.e., such that  $X_S = x_0 \exp([r \frac{\sigma^2}{2}]S + \sigma B_S)$ , where  $B_S \sim \text{Normal}(0, S)$ .
  - WHY?? Well, if  $\mu = r \frac{\sigma^2}{2}$ , then  $\{D_t\}$  becomes a martingale (HW#3), and this <u>turns out</u> to be a key fact. (finance/actuarial classes . . . )
- So, fair price is now just an <u>integral</u> (with respect to a normal density).
  - After some computation (HW#3), this fair price becomes:

$$x_0 \Phi\left(\frac{(r+\frac{\sigma^2}{2})S - \log(K/x_0)}{\sigma\sqrt{S}}\right) - e^{-rS}K\Phi\left(\frac{(r-\frac{\sigma^2}{2})S - \log(K/x_0)}{\sigma\sqrt{S}}\right),$$

where  $\Phi(u) = \int_{-\infty}^{u} \frac{1}{\sqrt{2\pi}} e^{-v^2/2} dv$  is the cdf of a standard normal distribution. ["Black-Scholes formula". Do <u>not</u> have to memorise!]

- Note: this price does <u>not</u> depend on the drift ("appreciation rate")  $\mu$ . [Surprising! Intuition: if  $\mu$  large, then can make good money from stock, so don't need the option.]
- However, it is an increasing function of the volatility  $\sigma$ . [Makes sense.]

## **Application – Sequence Waiting Times:**

- Suppose we repeatedly flip a fair coin. Let  $\tau$  be the first time the sequence "HTH" is completed. What is  $\mathbf{E}(\tau)$ ?
  - And, is the answer the same for "THH"?
  - Try it out: file www.probability.ca/sta447/Rseqwait
- One solution: use Markov chains!

- Suppose an irreducible Markov chain on a (discrete) state space S has a stationary distribution  $\pi$ . Then if the chain starts at some  $i \in S$ , how long until it returns to i, on average? That is, what is  $m_i \equiv \mathbf{E}_i(T_i)$ ?
  - Well, whatever  $m_i$  is, by the SLLN, over the long run, the chain will spend a fraction  $1/m_i$  of iterations at i:  $\lim_{n\to\infty} \frac{1}{n} \sum_{k=1}^n \mathbf{1}_{X_k=i} = 1/m_i$ .
  - Hence, in particular, by the Bounded Convergence Theorem,  $\lim_{n\to\infty} \mathbf{E}\left(\frac{1}{n}\sum_{k=1}^n \mathbf{1}_{X_k=i}\right) = 1/m_i$ .
  - But in the long run, the chain must spend an expected fraction  $\pi_i$  of iterations at i. That is,  $\lim_{n\to\infty} \mathbf{E}\left(\frac{1}{n}\sum_{k=1}^n \mathbf{1}_{X_k=i}\right) = \lim_{n\to\infty} \frac{1}{n}\sum_{k=1}^n \mathbf{E}_i(\mathbf{1}_{X_k=i}) = \lim_{n\to\infty} \frac{1}{n}\sum_{k=1}^n \mathbf{P}_i(X_k=i) = \lim_{n\to\infty} \mathbf{P}_i(X_n=i) = \lim_{n\to\infty} p_{ii} = \pi_i.$
  - So, we must have  $1/m_i = \pi_i$ , i.e.  $m_i = 1/\pi_i$ !
  - This is the RETURN TIME THEOREM: If an irreducible Markov chain on a discrete state space S has a stationary distribution  $\pi$ , then for any state  $i \in S$ , the mean return time satisfies  $m_i \equiv \mathbf{E}_i(T_i) = 1/\pi_i$ .
  - (For more details, see e.g. Rosenthal, Theorem 8.4.9.)
  - Let's apply this to sequence waiting times!
- Let  $X_n$  be the amount of the desired sequence (HTH) that the chain has "achieved so far". (For example, if the flips begin with HHTTHT, then  $X_1 = 1$ ,  $X_2 = 1$ ,  $X_3 = 2$ ,  $X_4 = 0$ ,  $X_5 = 1$ , and  $X_6 = 2$ .) Take  $X_0 = 0$ .
  - So,  $S = \{0, 1, 2, 3\}$ , with  $X_0 = 0$ , and  $X_{\tau} = 3$ .
  - Then  $p_{01} = p_{12} = p_{23} = 1/2$ . (Probability of continuing the sequence.)
  - Also  $p_{00} = p_{20} = 1/2$ . But instead of  $p_{10} = 1/2$ , have  $p_{11} = 1/2$ . Key!
  - (That is, if you fail to match the second flip, T, then you've already matched the first flip, H, for the next try.)
  - For completeness, assume we "start over" as soon as we win, so  $p_{3j} = p_{0j}$  for all j, i.e.  $p_{31} = p_{30} = 1/2$ .

- Thus, 
$$P = \begin{pmatrix} 1/2 & 1/2 & 0 & 0 \\ 0 & 1/2 & 1/2 & 0 \\ 1/2 & 0 & 0 & 1/2 \\ 1/2 & 1/2 & 0 & 0 \end{pmatrix}$$
.

- Compute(!) that the stationary distribution is (0.3, 0.4, 0.2, 0.1).
  - So, mean time to return from state 3 to state 3 is 1/0.1 = 10.
  - But returning from state 3 to state 3 has the same probabilities as going from state 0 to state 3.
  - Hence, the mean time to go from state 0 to state 3 is 10.
  - That is, mean waiting time for HTH is 10.
  - Solved it!
  - Try it out: file www.probability.ca/sta447/Rseqwait

- What about THH? Is it the same?
  - Here we compute similarly (check) that  $P = \begin{pmatrix} 1/2 & 1/2 & 0 & 0 \\ 0 & 1/2 & 1/2 & 0 \\ 0 & 1/2 & 0 & 1/2 \\ 1/2 & 1/2 & 0 & 0 \end{pmatrix}$ .
  - Compute (check) that the stationary distribution is (1/8, 1/2, 1/4, 1/8).
  - So, mean time to return to state 3 is 1/(1/8) = 8. Smaller!
  - Try it out: file www.probability.ca/sta447/Rseqwait
- ANOTHER APPROACH (to HTH), USING MARTINGALES:
- Suppose that at each time n, a new "player" appears, and bets \$1 on heads, then if they win they bet \$2 on tails, then if they win again they bet \$4 on heads. (Each player stops betting as soon as they either lose once or win three bets in a row.)
  - Let  $S_n$  be the total amount won by <u>all</u> the betters by time n.
  - Then  $\{S_n\}$  is a martingale with stopping time  $\tau$ .
  - Then have(!) that  $S_{\tau} = -(\tau 3) + (-1) + (1) + (7) = -\tau + 10$ .
  - It follows (optional: by Dominated Convergence Thm, since  $|S_n S_{n-1}| \le 7$ , and  $\mathbf{E}(\tau) < \infty$ ) that  $\mathbf{E}(S_\tau) = \mathbf{E}(S_0) = 0$ .
  - Hence,  $0 = \mathbf{E}(S_{\tau}) = -\mathbf{E}(\tau) + 10$ , whence  $\mathbf{E}(\tau) = 10$ . Same as before!
- Similarly, for THH, get that  $S_{\tau} = -(\tau 3) + (-1) + (-1) + (7) = -\tau + 8$ , whence  $\mathbf{E}(\tau) = 8$ . Same as before!

#### **Poisson Processes:**

- MOTIVATING EXAMPLE:
  - Suppose an average of  $\lambda = 2.5$  fires in Toronto per day.
  - Intuitively, this is caused by a very <u>large</u> number n of buildings, each of which has a very <u>small</u> probability p of having a fire.
  - Then mean =  $np = \lambda$ , so  $p = \lambda/n$ .
  - Then # fires today is Binomial $(n, \lambda/n) \approx \text{Poisson}(\lambda) = \text{Poisson}(2.5)$ .
  - [That is,  $\mathbf{P}(\# \text{ fires} = k) \approx e^{-2.5} \frac{(2.5)^k}{k!}$ , for k = 0, 1, 2, 3, ...]
  - And, # fires today and tomorrow combined  $\approx \text{Poisson}(2 * \lambda) = \text{Poisson}(5)$ , etc.
  - Full distribution?  $\mathbf{P}$ (fire within next hour)? etc.
- Let  $\{Y_n\}_{n=1}^{\infty}$  be i.i.d.  $\sim \text{Exp}(\lambda)$ , for some  $\lambda > 0$ .
  - So,  $Y_n$  has density function  $\lambda e^{-\lambda y}$  for y > 0.
  - And,  $\mathbf{P}(Y_n > y) = e^{-\lambda y}$  for y > 0.
  - And,  $\mathbf{E}(Y_n) = 1/\lambda$ .
- Let  $T_0 = 0$ , and  $T_n = Y_1 + Y_2 + \ldots + Y_n$  for  $n \ge 1$ . ("nth arrival time")

- [e.g.  $T_n = \text{time of } n^{\text{th}} \text{ fire.}$ ]
- Let  $N(t) = \max\{n \ge 0 : T_n \le t\} = \#\{n \ge 1 : T_n \le t\} = \#$  arrivals up to time t.
  - "Counting process". (Counts number of arrivals.)
  - [e.g. N(t) = # fires between times 0 and t.]
  - "Poisson process with intensity  $\lambda$ "
- What is distribution of N(t), i.e.  $\mathbf{P}(N(t)=m)$ ?
  - Well, N(t) = m iff both  $T_m \le t$  and  $T_{m+1} > t$ , which is iff there is  $0 \le s \le t$  with  $T_m = s$  and  $T_{m+1} T_m > t s$ .
  - Recall that  $Y_n \sim \text{Exp}(\lambda) = \text{Gamma}(1,\lambda)$ , so  $T_m := Y_1 + Y_2 + \ldots + Y_m \sim \text{Gamma}(m,\lambda)$ , with density function  $f_{T_m}(s) = \frac{\lambda^m}{\Gamma(m)} s^{m-1} e^{-\lambda s} = \frac{\lambda^m}{(m-1)!} s^{m-1} e^{-\lambda s}$ .
  - Also  $\mathbf{P}(T_{m+1} T_m > t s) = \mathbf{P}(Y_{m+1} > t s) = e^{-\lambda(t-s)}$ . So,

$$\mathbf{P}(N(t) = m) = \mathbf{P}(T_m \le t, T_{m+1} > t) = \mathbf{P}(\exists \ 0 \le s \le t : T_m = s, Y_{m+1} > t - s)$$

$$= \int_0^t f_{T_m}(s) \mathbf{P}(Y_{m+1} > t - s) ds = \int_0^t \frac{\lambda^m}{(m-1)!} s^{m-1} e^{-\lambda s} e^{-\lambda(t-s)} ds$$
$$= \frac{\lambda^m}{(m-1)!} e^{-\lambda t} \int_0^t s^{m-1} ds = \frac{\lambda^m}{(m-1)!} e^{-\lambda t} [\frac{t^m}{m}] = \frac{(\lambda t)^m}{m!} e^{-\lambda t}.$$

- Hence,  $N(t) \sim \text{Poisson}(\lambda t)$ .
- Thus,  $\mathbf{P}(N(t) = k) = e^{-\lambda t} \frac{(\lambda t)^k}{k!}$  for k = 0, 1, 2, ...
- Hence also  $\mathbf{E}(N(t)) = \lambda t$ , and  $\mathbf{Var}(N(t)) = \lambda t$ .

## END OF WEEK #8 -

- Now, recall the "memoryless" (or "forgetfulness") property of the  $\text{Exp}(\lambda)$  distribution: for a, b > 0,  $\mathbf{P}(Y_n > b + a \mid Y_n > a) = \mathbf{P}(Y_n > b) = e^{-\lambda b}$ .
  - This means the process  $\{N(t)\}$  "starts over" in each new time interval.
  - It follows that  $N(t+s) N(s) \sim N(t) \sim \text{Poisson}(\lambda t)$ .
  - Also follows that if  $0 \le a < b \le c < d$ , then N(d) N(c) indep. of N(b) N(a), and similarly for multiple non-overlapping time intervals. ("independent increments")
  - MORE GENERALLY: if  $0 \le t_1 \le s_1 \le t_2 \le s_2 \le \ldots \le t_k \le s_k$ , then  $N(s_i) N(t_i) \sim \text{Poisson}(\lambda(s_i t_i))$ , and  $\{N(s_i) N(t_i)\}_{i=1}^k$  are all independent. ("independent Poisson increments")
- DEFN: A <u>Poisson processes</u> with intensity  $\lambda > 0$  is a collection  $\{N(t)\}_{t\geq 0}$  of random non-decreasing integer counts N(t), satisfying: (a) N(0) = 0; (b)  $N(t) \sim \text{Poisson}(\lambda t)$  for all  $t \geq 0$ ; and (c) independent Poisson increments (as above).

- MOTIVATING EXAMPLE (cont'd): average of  $\lambda = 2.5$  fires per day.
  - Here, fires approximately follow a Poisson process with intensity 2.5.
  - So, **P**(9 fires today and tomorrow combined)  $\approx e^{-2*2.5} \frac{(2*2.5)^9}{9!} = e^{-5} (\frac{5^9}{9!}) \doteq 0.036$ .
  - **P**(at least one fire in next hour) =  $1 \mathbf{P}$ (no fires in next hour) =  $1 \mathbf{P}(N(1/24) = 0) = 1 e^{-2.5/24} \frac{(2.5/24)^0}{0!} \doteq 1 0.90 = 0.10$ .
  - **P**(exactly 3 fires in next hour) =  $e^{-2.5/24} \frac{(2.5/24)^3}{3!} \doteq 0.00017 \doteq 1/5891$ , etc.
- EXAMPLE: Let  $\{N(t)\}$  be a Poisson process with intensity  $\lambda = 2$ . Then

$$\mathbf{P}[N(3) = 5, \ N(3.5) = 9] = \mathbf{P}[N(3) = 5, \ N(3.5) - N(3) = 4]$$

$$= \mathbf{P}[N(3) = 5] \ \mathbf{P}[N(3.5) - N(3) = 4]$$

$$= \left[e^{-\lambda 3} \frac{(\lambda 3)^5}{5!}\right] \left[e^{-\lambda 0.5} \frac{(\lambda 0.5)^4}{4!}\right]$$

$$= \left(e^{-6} \frac{6^5}{120}\right) \left(e^{-1} \frac{1^4}{24}\right) = e^{-7}(2.7) \doteq 0.0025 \doteq 1/400.$$

- EXAMPLE: Let  $\{N(t)\}$  be a Poisson process with intensity  $\lambda$ .
  - Then for 0 < t < s,

$$\mathbf{P}(N(t) = 1 \mid N(s) = 1) = \frac{\mathbf{P}(N(t) = 1, N(s) = 1)}{\mathbf{P}(N(s) = 1)}$$

$$= \frac{\mathbf{P}(N(t) = 1, N(s) - N(t) = 0)}{\mathbf{P}(N(s) = 1)}$$

$$= \frac{e^{-\lambda t} \frac{(\lambda t)^1}{1!} e^{-\lambda (s-t)} \frac{(\lambda (s-t))^0}{0!}}{e^{-\lambda s} \frac{(\lambda s)^1}{1!}} = t/s.$$

- That is, conditional on N(s) = 1, the first event is <u>uniform</u> over [0, s]. (Distribution does not depend on  $\lambda$ .)
- Also, e.g.

$$\mathbf{P}(N(4) = 1 \mid N(5) = 3) = \frac{\mathbf{P}(N(4) = 1, N(5) = 3)}{\mathbf{P}(N(5) = 3)}$$

$$= \frac{\mathbf{P}(N(4) = 1, N(5) - N(4) = 2)}{\mathbf{P}(N(5) = 3)}$$

$$= \frac{(e^{-4\lambda}(4\lambda)^{1}/1!)(e^{-\lambda}\lambda^{2}/2!)}{e^{-5\lambda}(5\lambda)^{3}/3!} = \frac{(4)^{1}/1!)(1/2!)}{(5)^{3}/3!}$$

$$= \frac{4/2}{125/6} = 24/250 = 12/125.$$

- This also does not depend on  $\lambda$ . [And equals  $\binom{3}{1}(4/5)^1(1/5)^2$ . Why?]

- ALTERNATIVE APPROACH: Given N(t), as  $h \searrow 0$ ,
  - $P(N(t+h) N(t) = 1) = \lambda h + o(h).$
  - $\mathbf{P}(N(t+h) N(t) \ge 2) = o(h).$
  - This (together with independent increments) is another way to characterise Poisson processes.
- NOTE: the  $\{T_i\}$  tend to "clump up" in various patterns just by chance alone.
  - Doesn't "mean" anything at all: they're <u>independent</u>. ("Poisson clumping")
  - But it "seems" like it does have meaning!
  - See e.g. www.probability.ca/pois
- APPLICATION: pedestrian deaths example (true story).
  - 7 pedestrian deaths in Toronto (14 in GTA) in January 2010.
  - Media hype, friends concerned, etc.
  - Facts: Toronto averages about 31.9 per year, i.e.  $\lambda = 2.66$  per month.
  - So, **P**(7 or more) =  $\sum_{j=7}^{\infty} e^{-2.66} \frac{(2.66)^j}{j!} \doteq 1.9\%$ , about once per 52 months, i.e. about once per 4.4 years.
  - Not so rare! doesn't "mean" anything! (Though tragic.) "Poisson clumping"
  - See e.g. www.probability.ca/ped
  - Later, just two in Feb 1 Mar 15, 2010; less than expected (4), but no media!

#### • RELATED APPROACH:

- Suppose have  $\lambda$  buses per hour, i.e. about n buses every  $n/\lambda$  hours.
- Suppose the arrival times are completely <u>random</u>.
- Model this as  $T_1, T_2, \ldots, T_n \sim \text{Uniform}[0, n/\lambda]$ , i.i.d.
- Then for 0 < a < b, as  $n \to \infty$ ,

$$\#\{i: T_i \in [a,b]\} \sim \text{Binomial}(n, \frac{b-a}{n/\lambda})$$

$$= \text{Binomial}(n, \frac{\lambda(b-a)}{n}) \approx \text{Poisson}(\lambda(b-a)).$$

- Like a Poisson process!
- APPLICATION: Waiting Time Paradox.
  - Suppose there are an average of  $\lambda$  buses per hour. (e.g.  $\lambda = 5$ )
  - You arrive at the bus stop at a random time.
  - What is your expected waiting time until the next bus?

- If buses are completely <u>regular</u>, then waiting time is  $\sim$  Uniform[0,  $\frac{1}{\lambda}$ ], so mean =  $\frac{1}{2\lambda}$  hours. (e.g.  $\lambda = 5$ , mean =  $\frac{1}{10}$  hours = 6 minutes)
- If buses are <u>completely random</u>, then they form a Poisson process, so (by memoryless property) waiting time is  $\sim \text{Exp}(\lambda)$ , so mean  $= \frac{1}{\lambda}$  hours. Twice as long! (e.g.  $\lambda = 5$ , mean  $= \frac{1}{5}$  hours = 12 minutes)
- But same number of buses! Contradiction??
- No: you're more likely to arrive during a longer gap.
- Aside: What about streetcars?
  - They can't <u>pass</u> each other, so they sometimes clump up even <u>more</u> than do (independent) buses. (e.g. Spadina streetcar)
- SUPERPOSITION: Suppose  $\{N_1(t)\}_{t\geq 0}$  and  $\{N_2(t)\}_{t\geq 0}$  are two independent Poisson processes, with rates  $\lambda_1$  and  $\lambda_2$  respectively. Let  $N(t) = N_1(t) + N_2(t)$ .
  - Then  $\{N(t)\}_{t\geq 0}$  is a Poisson process with rate  $\lambda_1 + \lambda_2$ .
  - Proof? Sum of two independent Poissons is Poisson!

#### • EXAMPLE:

- Suppose undergrads arrive for office hours according to a Poisson process with intensity  $\lambda_1 = 5$  (i.e. one every 12 minutes on average).
- And, grads arrive independently according to their own Poisson process with intensity  $\lambda_2 = 3$  (i.e. one every 20 minutes on average).
- Then, what is expected number of minutes until first student arrives?
- Well, total # arrivals N(t) is Poisson process with  $\lambda = \lambda_1 + \lambda_2 = 5 + 3 = 8$ .
- Let A = time of first arrival.
- Then,  $\mathbf{P}(A > t) = \mathbf{P}(N(t) = 0) = e^{-\lambda t} \frac{(\lambda t)^0}{0!} = e^{-\lambda t}$ ; so  $A \sim \text{Exp}(\lambda)$ .
- Hence,  $\mathbf{E}(A) = 1/\lambda = 1/8$  hours, i.e. 7.5 minutes.
- THINNING: Let  $\{N(t)\}_{t\geq 0}$  be a Poisson process with rate  $\lambda$ .
  - Suppose each arrival is independently of "type i" with probability  $p_i$ , for  $i = 1, 2, 3, \ldots$  (e.g. bus or streetcar, male or female, undergrad or grad, etc.)
  - Let  $N_i(t)$  be number of arrivals of type i up to time t.
  - THM: The  $\{N_i(t)\}$  are independent Poisson processes, with rates  $\lambda p_i$ .
  - PROOF: "independent increments" is obvious.
  - For the distribution, suppose for notational simplicity that there are just two types, with  $p_1 + p_2 = 1$ .
  - Need to show:  $\mathbf{P}(N_1(t) = j, N_2(t) = k)$ =  $\left(e^{-(\lambda p_1 t)}(\lambda p_1 t)^j / j!\right) \left(e^{-(\lambda p_2 t)}(\lambda p_2 t)^k / k!\right)$ .

- But  $\mathbf{P}(N_1(t) = j, N_2(t) = k)$ =  $\mathbf{P}(j+k \text{ arrivals up to time } t, \text{ of which } j \text{ of type 1 and } k \text{ of type 2})$ =  $\left(e^{-\lambda t}(\lambda t)^{j+k}/(j+k)!\right) \times {j+k \choose j}(p_1)^j(p_2)^k$ . Equal! (Check.)
- EXAMPLE: If students arrive for office hours according to a Poisson process, and each student is independently either undergrad (prob  $p_1$ ) or grad (prob  $p_2$ ), then # undergrads is <u>independent</u> of # grads (and each follows a Poisson distribution).
- ASIDE: Can also have <u>time-inhomogeneous</u> Poisson processes, where  $\lambda = \lambda(t)$ , and  $N(b) N(a) \sim \text{Poisson}(\int_a^b \lambda(t) dt)$ .
- ASIDE: Can also have Poisson processes on other <u>regions</u>, e.g. in two dimensions, etc., cf. <u>www.probability.ca/pois</u>

# Continuous-Time, Discrete-Space Markov Processes:

- Recall: Markov chains  $\{X_n\}_{n=0}^{\infty}$  defined in <u>discrete</u> (integer) time.
  - But Brownian motion  $\{B_t\}_{t\geq 0}$ , and Poisson processes  $\{N(t)\}_{t\geq 0}$ , both defined in <u>continuous</u> (real) time.
  - Can we define Markov processes in continuous time? Yes!
- DEFN: a continuous-time (time-homogeneous, non-explosive) Markov process, on a countable (discrete) state space S, is a collection  $\{X(t)\}_{t\geq 0}$  of random variables such that

$$\mathbf{P}(X_0 = i_0, X_{t_1} = i_1, X_{t_2} = i_2, \dots, X_{t_n} = i_n) = \nu_{i_0} p_{i_0 i_1}^{(t_1)} p_{i_1 i_2}^{(t_2 - t_1)} \dots p_{i_{n-1} i_n}^{(t_n - t_{n-1})},$$

for some <u>initial distribution</u>  $\{\nu_i\}_{i\in S}$  (with  $\nu_i \geq 0$ , and  $\sum_{i\in S} \nu_i = 1$ ), and <u>transition probabilities</u>  $\{p_{ij}^{(t)}\}_{i>0}^{i,j\in S}$  (with  $p_{ij}^{(t)} \geq 0$ , and  $\sum_{j\in S} p_{ij}^{(t)} = 1$ ).

- Just like for discrete-time chains, except need to keep track of the elapsed time (t) too.
- As with discrete chains,  $p_{ij}^{(0)} = \delta_{ij} = \begin{cases} 1, & i = j \\ 0, & i \neq j \end{cases}$
- Let  $P^{(t)} = \left(p_{ij}^{(t)}\right)_{i,j \in S} = \text{matrix version}.$ 
  - Then  $P^{(0)} = I = identity matrix.$
  - Also  $p_{ij}^{(s+t)} = \sum_{k \in S} p_{ik}^{(s)} p_{kj}^{(t)}$ , i.e.  $P^{(s+t)} = P^{(s)} P^{(t)}$ . ("Chapman-Kolmogorov equations", just like for discrete time)
  - If  $\mu_i^{(t)} = \mathbf{P}(X(t) = i)$ , and  $\mu^{(t)} = \left(\mu_i^{(t)}\right)_{i \in S} = \text{row vector, and } \nu = (\nu_i)_{i \in S} = \text{row vector, then } \mu_j^{(t)} = \sum_{i \in S} \nu_i p_{ij}^{(t)}$ , and  $\mu^{(t)} = \nu P^{(t)}$ , and  $\mu^{(t)} P^{(s)} = \mu^{(t+s)}$ , etc.
- Expect that  $\lim_{t\searrow 0} p_{ij}^{(t)} = p_{ij}^{(0)} = \delta_{ij}$ .
  - <u>Assume</u> this is true. ("standard" Markov process)
- Then can compute the process's generator as  $g_{ij} = \lim_{t \searrow 0} \frac{p_{ij}^{(t)} \delta_{ij}}{t} = p'_{ij}(0)$ .

(right-handed derivative)

- So, if  $G = (g_{ij})_{i,j \in S} = \text{matrix}$ , then  $G = P'^{(0)} = \lim_{t \searrow 0} \frac{P^{(t)} I}{t}$ . (righthanded derivative)
- Here  $g_{ii} \leq 0$ , while  $g_{ij} \geq 0$  for  $i \neq j$ .
- In fact, usually (e.g. if S is finite), have

$$\sum_{j \in S} g_{ij} = \sum_{j \in S} \lim_{t \searrow 0} \frac{p_{ij}(t) - \delta_{ij}}{t} = \lim_{t \searrow 0} \frac{\sum_{j \in S} p_{ij}(t) - \sum_{j \in S} \delta_{ij}}{t} = \lim_{t \searrow 0} \frac{1 - 1}{t} = 0.$$

- Furthermore, if t > 0 is small, then  $G \approx \frac{P^{(t)} I}{t}$ , so  $P^{(t)} \approx I + tG$ , i.e.  $p_{ii}^{(t)} \approx \delta_{ii} + t q_{ii}$
- RUNNING EXAMPLE:  $S = \{1, 2\}$ , and  $G = \begin{pmatrix} -3 & 3 \\ 6 & -6 \end{pmatrix}$ .
  - Then for small t > 0,  $P^{(t)} \approx I + tG = \begin{pmatrix} 1 3t & 3t \\ 6t & 1 6t \end{pmatrix}$ .
  - So  $p_{11}^{(t)} \approx 1 3t$ ,  $p_{12}^{(t)} \approx 3t$ , etc.
  - e.g. if t = 0.02, then  $p_{11}^{(0.02)} \doteq 1 3(0.02) = 0.94$ ,  $p_{12}^{(0.02)} \doteq 3(0.02) = 0.06$ ,  $p_{21}^{(0.02)} \doteq 6(0.02) = 0.12$ , and  $p_{22}^{(0.02)} \doteq 1 6(0.02) = 0.88$ , i.e.  $P^{(0.02)} \doteq \begin{pmatrix} 0.94 & 0.06 \\ 0.12 & 0.88 \end{pmatrix}$ .
- What about for larger t?
  - Well, by Chapman-Kolmogorov eqn, for any  $m \in \mathbb{N}$ ,

$$P^{(t)} = [P^{(t/m)}]^m = \lim_{n \to \infty} [P^{(t/n)}]^n = \lim_{n \to \infty} [I + (t/n)G]^n$$
$$= \exp(tG) := I + tG + \frac{t^2G^2}{2!} + \frac{t^3G^3}{3!} + \dots$$

(matrix equation; similar to how  $\lim_{n\to\infty} (1+\frac{c}{n})^n = e^c$ ).

- (Makes sense so that e.g.  $P^{(s+t)} = \exp((s+t)G) = \exp(sG) \exp(tG) =$  $P^{(s)} P^{(t)}$ , etc.)
- So, in principle, the generator G tells us  $P^{(t)}$  for all  $t \geq 0$ .
- Can we actually compute  $P^{(t)} = \exp(tG)$  this way? Yes!
- Method #1: Compute the infinite matrix sum on a computer, numerically and approximately.
- Method #2: Note that in above example, if  $\lambda_1 = 0$  and  $\lambda_2 = -9$ , and  $w_1 = (2,1)$  and  $w_2 = (1,-1)$ , then  $w_1G = \lambda_1w_1 = 0$ , and  $w_2G = 0$  $\lambda_2 w_2 = -9w_2$ . That is,  $\{\lambda_i\}$  are the <u>eigenvalues</u> of G, with corresponding left-eigenvectors  $\{w_i\}$ .
  - Now, if  $w_i$  is a left-eigenvector with corresponding eigenvalue  $\lambda_i$ , then  $w_i \exp(tG) = e^{t\lambda_i} w_i$ . (Check.) Easy!

– So, if initial distribution is (say)  $\nu = (1,0)$ , then first compute that  $\nu = \frac{1}{3}w_1 + \frac{1}{3}w_2$ . Then,

$$\mu^{(t)} = \nu P^{(t)} = \nu \exp(tG) = \left(\frac{1}{3}w_1 + \frac{1}{3}w_2\right) \exp(tG)$$

$$= \frac{1}{3}e^{t\lambda_1}w_1 + \frac{1}{3}e^{t\lambda_2}w_2 = \frac{1}{3}e^{0t}(2,1) + \frac{1}{3}e^{-9t}(1,-1) = \left(\frac{2+e^{-9t}}{3}, \frac{1-e^{-9t}}{3}\right).$$

- So,  $\mathbf{P}[X_t = 1] = p_{11}^{(t)} = \frac{2+e^{-9t}}{3}$ , and  $\mathbf{P}[X_t = 2] = p_{12}^{(t)} = \frac{1-e^{-9t}}{3}$ .
- Check:  $p_{11}^{(0)} = 1$ ,  $p_{12}^{(0)} = 0$ , and  $p_{11}^{(t)} + p_{12}^{(t)} = 1$ . (Phew.)
- (Or, by instead choosing  $\nu = (0,1)$ , could compute  $p_{21}^{(t)}$  and  $p_{22}^{(t)}$ .)

# END OF WEEK #9

• Method #3: Note that

$$p'_{ij}^{(t)} = \lim_{h \searrow 0} \frac{p_{ij}^{(t+h)} - p_{ij}^{(t)}}{h} = \lim_{h \searrow 0} \frac{\left(\sum_{k \in S} p_{ik}^{(t)} p_{kj}^{(h)}\right) - p_{ij}^{(t)}}{h}$$

$$= \lim_{h \searrow 0} \frac{\left(\sum_{k \in S} p_{ik}^{(t)} \left[\delta_{kj} + h g_{kj}\right]\right) - p_{ij}^{(t)}}{h}$$

$$= \lim_{h \searrow 0} \frac{\left(p_{ij}^{(t)} + h \sum_{k \in S} p_{ik}^{(t)} g_{kj}\right) - p_{ij}^{(t)}}{h} = \sum_{k \in S} p_{ik}^{(t)} g_{kj},$$

i.e.  $P'^{(t)} = P^{(t)} G$ . ("forward equations")

- (Makes sense since  $P^{(t)} = \exp(tG)$ , so  $P'^{(t)} = \exp(tG) G = P^{(t)} G$ .)
- So, in above example,

$$p'_{11}^{(t)} = p_{11}^{(t)} g_{11} + p_{12}^{(t)} g_{21} = (-3)p_{11}^{(t)} + (6)p_{12}^{(t)} = (-3)p_{11}^{(t)} + (6)(1 - p_{11}^{(t)})$$
$$= (-9)p_{11}^{(t)} + 6 = (-9)(p_{11}^{(t)} - \frac{2}{3}).$$

- But  $p'_{ij}^{(t)} = \frac{d}{dt}(p_{11}^{(t)}) = \frac{d}{dt}(p_{11}^{(t)} \frac{2}{3}).$
- So,  $\frac{d}{dt}(p_{11}^{(t)} \frac{2}{3}) = (-9)(p_{11}^{(t)} \frac{2}{3}).$
- So,  $p_{11}^{(t)} \frac{2}{3} = Ke^{-9t}$ , i.e.  $p_{11}^{(t)} = \frac{2}{3} + Ke^{-9t}$ .
- But  $p_{11}^{(0)} = 1$ , so  $K = \frac{1}{3}$ , so  $p_{11}^{(t)} = \frac{2}{3} + \frac{1}{3}e^{-9t} = \frac{2+e^{-9t}}{3}$ .
- And then  $p_{12}^{(t)} = 1 p_{11}^{(t)} = \frac{1 e^{-9t}}{3}$ .
- Same answers as before. (Phew.)
- What about LIMITING PROBABILITIES?
- In above example,  $\mu^{(t)} = (\frac{2+e^{-9t}}{3}, \frac{1-e^{-9t}}{3})$ , so  $\lim_{t\to\infty} \mu^{(t)} = (\frac{2}{3}, \frac{1}{3}) =: \pi$ .
  - Note that  $\sum_{i \in S} \pi_i g_{i1} = \frac{2}{3}(-3) + \frac{1}{3}(6) = 0$ , and  $\sum_{i \in S} \pi_i g_{i2} = \frac{2}{3}(3) + \frac{1}{3}(-6) = 0$ .

- i.e.,  $\sum_{i \in S} \pi_i g_{ij} = 0$  for all  $j \in S$ , i.e.  $\pi G = 0$ .
- Does this make sense?
  - Well, as in discrete case,  $\{\pi_i\}$  should be stationary.
  - i.e.  $\sum_{i \in S} \pi_i p_{ij}^{(t)} = \pi_j$  for all  $j \in S$  and <u>all</u>  $t \ge 0$ .
  - In particular, for small t > 0,

$$\pi_j = \sum_{i \in S} \pi_i p_{ij}^{(t)} \approx \sum_{i \in S} \pi_i [\delta_{ij} + t g_{ij}] = \pi_j + t \sum_{i \in S} \pi_i g_{ij}.$$

- So,  $\sum_{i \in S} \pi_i g_{ij} = 0$ .
- So, can check if  $\{\pi_i\}$  is stationary by checking if  $\sum_{i \in S} \pi_i g_{ij} = 0$  for all  $j \in S$ .
- What about reversibility?
  - Well, if  $\pi_i g_{ij} = \pi_j g_{ji}$  for all  $i, j \in S$ , then  $\sum_i \pi_i g_{ij} = \sum_i \pi_j g_{ji} = \pi_j \sum_i g_{ji} = \pi_j \cdot 0 = 0$ , so  $\pi$  is stationary.
  - So, again, reversibility (in the above sense) implies stationary!
  - In above example,  $\pi_1 g_{12} = (2/3)(3) = 2$ , while  $\pi_2 g_{21} = (1/3)(6) = 2$ , so it's reversible.
- Is convergence to  $\{\pi_i\}$  guaranteed?
- CONTINUOUS-TIME MARKOV CONVERGENCE THEOREM: If a continuous-time M.C. is <u>irreducible</u>, and has a <u>stationary distribution</u>  $\pi$ , then  $\lim_{t\to\infty} p_{ij}^{(t)} = \pi_j$  for all  $i,j \in S$ .
  - Like discrete case, but don't need aperiodicity (i.e., in continuous time, it is <u>automatically</u> aperiodic).
  - Proof omitted here; similar to discrete time.
  - See e.g. Durrett, 2nd ed., Theorem 4.4, p. 128.
- CONNECTION TO DISCRETE-TIME MARKOV CHAINS:
  - Let  $\{\hat{p}_{ij}\}_{i,j\in S}$  be the transition probabilities for a <u>discrete-time</u> Markov chain  $\{\hat{X}_n\}_{n=0}^{\infty}$ .
  - Let  $\{N(t)\}_{t\geq 0}$  be a Poisson process with intensity  $\lambda>0$ .
  - Then let  $X_t = \hat{X}_{N(t)}$ .
  - Then  $\{X_t\}$  is just like  $\{\hat{X}_n\}$  except that it jumps at Poisson process event times, not integer times. ("Exponential holding times")
  - In particular,  $\{X_t\}$  is a continuous-time Markov process!
  - That is, we can "create" a continuous-time Markov process from a discrete-time Markov chain.
- What is the generator of this Markov process  $\{X_t\}$ ?
  - Well, here  $p_{ij}^{(t)} = \mathbf{P}_i[\hat{X}_{N(t)} = j].$

- So, 
$$p_{ij}^{(t)} = \sum_{n=0}^{\infty} \mathbf{P}_i[N(t) = n, \ \hat{X}_n = j].$$

- So, 
$$p_{ij}^{(t)} = \sum_{n=0}^{\infty} \mathbf{P}[N(t) = n] \, \hat{p}_{ij}^{(n)} = \sum_{n=0}^{\infty} \left[ e^{-\lambda t} \frac{(\lambda t)^n}{n!} \right] \hat{p}_{ij}^{(n)}$$

- But for small t > 0,  $\mathbf{P}[N(t) = n] = e^{-\lambda t} (\lambda t)^n / n! \approx t^n \lambda^n / n!$ .
- So, for small t > 0 and  $i \neq j$ ,

$$p_{ij}^{(t)} = \sum_{n=0}^{\infty} \mathbf{P}[N(t) = n] \, \hat{p}_{ij}^{(n)}$$

$$= \mathbf{P}[N(t) = 0] \, \hat{p}_{ij}^{(0)} + \mathbf{P}[N(t) = 1] \, \hat{p}_{ij}^{(1)} + \mathbf{P}[N(t) = 2] \, \hat{p}_{ij}^{(2)} + \dots$$

$$= \mathbf{P}[N(t) = 0] \, (0) + \mathbf{P}[N(t) = 1] \, \hat{p}_{ij} + \mathbf{P}[N(t) = 2] \, \hat{p}_{ij}^{(2)} + \dots$$

$$\approx 0 + [t\lambda] \, \hat{p}_{ij} + [t^2 \lambda^2 / 2!] \, \hat{p}_{ij}^{(2)} + \dots$$

$$= [t\lambda] \hat{p}_{ij} + O(t^2) \approx [t\lambda] \hat{p}_{ij},$$

to first order in t, as  $t \searrow 0$ .

- But 
$$p_{ij}^{(t)} \approx \delta_{ij} + tg_{ij} = tg_{ij}$$
, so  $tg_{ij} = [t\lambda]\hat{p}_{ij}$ , so  $g_{ij} = \lambda \hat{p}_{ij}$ .

- Also 
$$p_{ii}^{(t)} \approx \mathbf{P}[N(t) = 0] + \mathbf{P}[N(t) = 1] \hat{p}_{ii} + O(t^2) \approx [1 - t\lambda] + [t\lambda] \hat{p}_{ii}$$
.

- But 
$$p_{ii}^{(t)} \approx \delta_{ii} + tg_{ii} = 1 + tg_{ii}$$
, so  $[1 - t\lambda] + [t\lambda]\hat{p}_{ii} = 1 + tg_{ii}$ , so  $g_{ii} = \lambda (\hat{p}_{ii} - 1)$ .

- Check: for  $i \neq j$ ,  $g_{ij} \geq 0$ , and  $g_{ii} \leq 0$ . Good.
- Also,  $\sum_{j \in S} g_{ij} = g_{ii} + \sum_{j \neq i} g_{ij} = \lambda(\hat{p}_{ii} 1) + \sum_{j \neq i} (\lambda \hat{p}_{ij}) = -\lambda + \sum_{j \in S} (\lambda \hat{p}_{ij}) = -\lambda + \lambda \sum_{j \in S} \hat{p}_{ij} = -\lambda + \lambda(1) = 0$ , as it must.
- SPECIAL CASE: if  $\hat{X}_0 = 0$ , and  $\hat{p}_{i,i+1} = 1$  for all i, then  $\hat{X}_n = n$  for all n, so  $X_t = \hat{X}_{N(t)} = N(t) = \text{Poisson process.}$ 
  - And, the Poisson process  $\{N(t)\}_{t\geq 0}$  itself has generator (check):

$$G = \begin{pmatrix} -\lambda & \lambda & 0 & 0 & \dots \\ 0 & -\lambda & \lambda & 0 & \dots \\ 0 & 0 & -\lambda & \lambda & \dots \end{pmatrix}.$$

# Application – Queueing Theory:

- Consider a queue (i.e., a line of customers) with just one server.
  - Let  $T_n$  = time of <u>arrival</u> of  $n^{th}$  customer. (And set  $T_0 = 0$ .)
  - Let  $Y_n = T_n T_{n-1} = \underline{\text{inter-arrival}}$  time between  $(n-1)^{\text{st}}$  and  $n^{\text{th}}$  customers.
  - Let  $S_n$  = time it takes to <u>serve</u> the  $n^{th}$  customer.
  - Let Q(t) = number of customers in the system (i.e., waiting in the queue or being served) at time  $t \ge 0$ . (Assume Q(0) = 0.)
- What happens as  $t \to \infty$ ?
- M/M/1 QUEUE:  $T_n T_{n-1} \sim \text{Exp}(\lambda)$ , and  $S_n \sim \text{Exp}(\mu)$ , all indep.,  $\lambda, \mu > 0$ . (So  $\{T_n\}$  are arrival times of a Poisson process with intensity  $\lambda$ .)

- Then by memoryless property,  $\{Q(t)\}\$  is a Markov process!

## • GENERATOR?

- Well, for  $n \ge 0$ ,  $\mathbf{P}[Q(t) = n + 1 | Q(0) = n]$ 
  - =  $\mathbf{P}$ [one arrival and zero served by time t] +  $\mathbf{P}$ [two arrivals and one served by time t] + . . .
  - $= \mathbf{P}[\text{one arrival and zero served by time } t] + O(t^2)$
  - $\approx \mathbf{P}[\text{one arrival and zero served by time } t].$
- Hence,, for  $n \ge 0$ , to first order as  $t \searrow 0$ ,

$$g_{n,n+1} = \lim_{t \searrow 0} \frac{\mathbf{P}[Q(t) = n+1 \mid Q(0) = n]}{t}$$

$$= \lim_{t \searrow 0} \frac{\mathbf{P}[\text{one arrival and zero served by time } t]}{t}$$

$$= \lim_{t \searrow 0} \frac{\left[e^{-\lambda t} \frac{(\lambda t)^{1}}{1!}\right] \left[e^{-\mu t} \frac{(\mu t)^{0}}{0!}\right]}{t} = \lim_{t \searrow 0} e^{-\lambda t} \lambda e^{-\mu t} = \lambda.$$

- Similarly, for  $n \geq 1$ ,

$$g_{n,n-1} = \lim_{t \searrow 0} \frac{\mathbf{P}[Q(t) = n - 1 \mid Q(0) = n]}{t}$$

$$= \lim_{t \searrow 0} \frac{\mathbf{P}[\text{zero arrivals and one served by time } t]}{t}$$

$$= \lim_{t \searrow 0} \frac{\left[e^{-\lambda t} \frac{(\lambda t)^0}{0!}\right] \left[e^{-\mu t} \frac{(\mu t)^1}{1!}\right]}{t} = \mu.$$

- Also if  $|n-m| \ge 2$  then  $\mathbf{P}[Q(t) = m | Q(0) = n] = O(t^2)$ , so  $g_{n,m} = 0$ .
- But  $\sum_{m=0}^{\infty} g_{n,m} = 0$ , so the generator must be given by:

$$G = \begin{pmatrix} -\lambda & \lambda & 0 & 0 & 0 & \dots \\ \mu & -\lambda - \mu & \lambda & 0 & 0 & \dots \\ 0 & \mu & -\lambda - \mu & \lambda & 0 & \dots \\ 0 & 0 & \ddots & \ddots & \ddots & \end{pmatrix}$$

i.e.  $g_{00} = -\lambda$  and  $g_{nn} = -\lambda - \mu$  for  $n \ge 1$ . (This corresponds to zero arrivals and zero served by time t; check.)

- So we have solved for the queue generator matrix G.

# • STATIONARY DISTRIBUTION $\{\pi_i\}$ ?

- Need  $\sum_{i \in S} \pi_i g_{ij} = 0$  for all  $j \in S$ . (Or, can use <u>reversibility</u>: check.)
- j = 0:  $\pi_0(-\lambda) + \pi_1(\mu) = 0$ , so  $\pi_1 = (\frac{\lambda}{\mu})\pi_0$ .
- $j = 1: \pi_0(\lambda) + \pi_1(-\lambda \mu) + \pi_2(\mu) = 0,$ so  $\pi_2 = (\frac{\lambda}{-\mu})\pi_0 + (\frac{-\lambda - \mu}{-\mu})\pi_1 = (-\frac{\lambda}{\mu})\pi_0 + (1 + \frac{\lambda}{\mu})(\frac{\lambda}{\mu})\pi_0 = (\frac{\lambda}{\mu})^2\pi_0.$
- Then by induction:  $\pi_i = (\frac{\lambda}{\mu})^i \pi_0$ , for  $i = 0, 1, 2, \dots$

- So if  $\lambda < \mu$ , i.e.  $\frac{1}{\mu} < \frac{1}{\lambda}$ , i.e.  $\mathbf{E}(S_n) < \mathbf{E}(T_n - T_{n-1})$ , then since  $\sum_i \pi_i = 1$ ,

$$\pi_0 = \frac{1}{\sum_{i=0}^{\infty} (\frac{\lambda}{\mu})^i} = 1 - (\frac{\lambda}{\mu}),$$

and the stationary distribution is

$$\pi_i = \left(\frac{\lambda}{\mu}\right)^i \left(1 - \frac{\lambda}{\mu}\right), \qquad i = 0, 1, 2, 3, \dots$$

(geometric distribution).

• Furthermore, since the process is clearly irreducible,

$$\lim_{n \to \infty} \mathbf{P}[Q(t) = i] = \pi_i = \left(\frac{\lambda}{\mu}\right)^i \left(1 - \frac{\lambda}{\mu}\right).$$

- By contrast, if  $\lambda > \mu$ , then  $Q(t) \to \infty$  w.p. 1. (see below)
- If  $\lambda = \mu$ , then  $Q(t) \to \infty$  in probability, but not w.p. 1. (see below)

## General (G/G/1) Queue:

- What if we don't assume Exponential distributions, just that  $\{T_n T_{n-1}\}$  i.i.d., and  $\{S_n\}$  i.i.d. (all indep.)?
- Then Q(t) is not Markovian! Have to use "cruder" methods.
- Let  $D_n = \text{time of } \underline{\text{departure}}$  of the  $n^{\text{th}}$  customer.
- Roughly speaking,  $D_n = D_{n-1} + S_n$ .
  - But no one served while queue is <u>empty</u>.
  - So, actually,  $D_n = \max(T_n, D_{n-1}) + S_n$ .
- Let  $W_n = \max(0, D_{n-1} T_n) =$ the amount of time that the  $n^{\text{th}}$  customer has to wait. (With  $W_0 = 0$ .)
- LINDLEY'S EQUATION: For  $n \ge 1$ ,  $W_n = \max(0, W_{n-1} + S_{n-1} Y_n)$ .
- PROOF:
  - The  $(n-1)^{st}$  customer is in the system for a total time  $W_{n-1} + S_{n-1}$ .
  - But the  $n^{\text{th}}$  customer arrives a time  $Y_n$  after the  $(n-1)^{\text{st}}$  customer.
  - If  $W_{n-1} + S_{n-1} \leq Y_n$ , then the  $n^{\text{th}}$  customer doesn't have to wait at all, so  $W_n = 0$ .
  - If  $W_{n-1} + S_{n-1} \ge Y_n$ , then  $W_n = [\text{time the } (n-1)^{\text{st}} \text{ customer is in the system}] [\text{amount of that time that the } n^{\text{th}} \text{ customer was } \underline{\text{not}} \text{ present for}] = (W_{n-1} + S_{n-1}) Y_n$ , Q.E.D.
- LINDLEY'S COROLLARY:  $W_n = \max_{0 \le k \le n} \sum_{i=k+1}^n (S_{i-1} Y_i)$ . (Here, if k = n, then the sum equals zero.)
  - PROOF #1. Write it out:  $W_0 = 0$ ,  $W_1 = \max(0, S_0 Y_1)$ ,  $W_2 = \max(0, W_1 + S_1 Y_2) = \max(0, \max(0, S_0 Y_1) + S_1 Y_2) = \max(0, \max(0, S_0 Y_1) + S_1 Y_2) = \max(0, \max(0, S_0 Y_1) + S_1 Y_2)$

$$\begin{aligned} & \max(0,\ S_1-Y_2,\ S_0-Y_1+S_1-Y_2),\\ & W_3=\max(0,W_2+S_2-Y_3)=\max(0,\max(0,\ S_1-Y_2,\ S_0-Y_1+S_1-Y_2)+\\ & S_2-Y_3)=\max(0,S_2-Y_3,\ S_1-Y_2+S_2-Y_3,\ S_0-Y_1+S_1-Y_2+S_2-Y_3),\\ & \text{etc., each corresponding to the claimed formula.} \end{aligned}$$

- PROOF #2: Induction on n. When n = 0, both sides are zero. If n increases to n + 1, then by Lindley's Equation, each possible value of the "max" gets  $S_n Y_{n+1}$  added to it. And the "max with zero" is covered by allowing for the possibility k = n + 1, Q.E.D.
- THEOREM: For a general (G/G/1) single-server queue:
  - (a) if  $\mathbf{E}(Y_n) < \mathbf{E}(S_n)$ , then  $\lim_{n \to \infty} W_n = \infty$  w.p. 1. (Hence, also  $\lim_{n \to \infty} W_n = \infty$  in probability, so for any  $M < \infty$ ,  $\lim_{n \to \infty} \mathbf{P}(W_n > M) = 1$ .)
  - (b) if  $\mathbf{E}(Y_n) > \mathbf{E}(S_n)$ , then  $\{W_n\}$  is "bounded in probability", i.e. for any  $\epsilon > 0$  there is  $M < \infty$  such that  $\mathbf{P}(W_n > M) < \epsilon$  for all  $n \in \mathbf{N}$ .
  - (c) if  $\mathbf{E}(Y_n) = \mathbf{E}(S_n)$ , and  $S_{n-1}$  and  $Y_n$  are not both constant (i.e.,  $\mathbf{Var}(S_{n-1} Y_n) > 0$ ), then  $W_n \to \infty$  in probability (but <u>not</u> w.p. 1). ("Borderline" case; similar to branching processes with m = 1.)
- PROOF OF (a):
  - By Lindley's Equation,  $W_{n+1} \ge W_n + S_n Y_{n+1}$ .
  - Here the sequence  $\{S_n Y_{n+1}\}$  is i.i.d., with mean > 0.
  - So, by the SLLN,  $\liminf_{n\to\infty} \frac{W_n}{n} \geq \mathbf{E}(S_n Y_{n+1}) > 0$ , w.p. 1.
  - It follows that  $\liminf_{n\to\infty} W_n \ge \infty$ , w.p. 1, Q.E.D.
- PROOF OF (b):
  - By Lindley's Corollary,  $\mathbf{P}(W_n > M) = \mathbf{P}\left(\max_{0 \le k \le n} \sum_{i=k+1}^n (S_{i-1} - Y_i) > M\right).$
  - But  $\{S_{i-1} Y_i\}$  are i.i.d., so this is equivalent to  $\mathbf{P}(W_n > M) = \mathbf{P}\left(\max_{0 \le k \le n} \sum_{i=1}^{n-k} (S_{i-1} Y_i) > M\right).$
  - This is the probability that i.i.d. partial sums with negative mean (since  $\mathbf{E}(S_{i-1} Y_i) < 0$ ) will ever be larger than M.
  - i.e., it is the probability that the  $\underline{\text{maximum}}$  of a sequence of i.i.d. partial sums with negative mean will be larger than M.
  - But by the SLLN, i.i.d. partial sums with negative mean will eventually become <u>negative</u>, w.p. 1.
  - So, w.p. 1, only a <u>finite</u> number of the partial sums will have positive values.
  - So, w.p. 1, the maximum value of the partial sums will be <u>finite</u>.
  - So, as  $M \to \infty$ , the probability that the maximum value will be > M must converge to zero.

- So, for any  $\epsilon > 0$ , there is  $M < \infty$  such that the probability that its maximum value is be > M is  $< \epsilon$ , Q.E.D.
- PROOF OF (c):
  - Trickier! Omitted! For details see e.g. Grimmett & Stirzaker, 2nd ed., Theorem 11.5(4), pp. 432–435.

# Application: Quantum Mechanics (simplified view!):

- According to quantum mechanics, the universe on a fundamental level behaves according to probabilities(!).
  - More specifically, it has a complex-valued wave function  $\Psi$  which evolves according to various rules, and whose squared absolute value  $|\Psi|^2$  gives the probability of observing a given state.
- Here is a very simplified example ...
- Suppose an elementary particle has (just) two "eigenstates",  $E_1$  and  $E_2$ , with "energies"  $\lambda_1 \neq \lambda_2$ , and an "observable state"  $B = \frac{1}{\sqrt{2}}E_1 + \frac{1}{\sqrt{2}}E_2$ .
  - State of system is described by a complex-valued "wave function"  $\Psi_t$ .
  - Then, the projection of the wave function onto each of the two eigenstates  $E_1$  and  $E_2$  evolves by the formulas:  $\langle \Psi_t | E_1 \rangle = \langle \Psi_0 | E_1 \rangle e^{i\lambda_1 t}$ , and  $\langle \Psi_t | E_2 \rangle = \langle \Psi_0 | E_2 \rangle e^{i\lambda_2 t}$ , where  $i = \sqrt{-1}$ .
  - (This is a special case of the "Schrödinger Wave Equation".)
  - Then, by linearity, the projection of the wave function onto the observable state B is given by:  $\langle \Psi_t | B \rangle = \frac{1}{\sqrt{2}} \langle \Psi_t | E_1 \rangle + \frac{1}{\sqrt{2}} \langle \Psi_t | E_2 \rangle$ . ("superposition")
- Now, suppose we "observe" the system at some time t.
  - Then, **P**(observe state *B* at time *t*) =  $|\langle \Psi_t | B \rangle|^2 = |\frac{1}{\sqrt{2}} \langle \Psi_t | E_1 \rangle + \frac{1}{\sqrt{2}} \langle \Psi_t | E_2 \rangle|^2$ .
- QUESTION: Suppose the system is in state B at time 0, so  $\langle \Psi_0 | E_1 \rangle = \langle \Psi_0 | E_2 \rangle = \frac{1}{\sqrt{2}}$ . Suppose we then observe the system at some later time t > 0. Then what is  $\mathbf{P}(\text{observe state } B \text{ at time } t)$ ?

### END OF WEEK #10 -

- SOLUTION:
  - We know that  $\langle \Psi_0 | E_1 \rangle = \langle \Psi_0 | E_2 \rangle = \frac{1}{\sqrt{2}}$ .
  - So, at time t,  $\langle \Psi_t | E_j \rangle = \langle \Psi_0 | E_j \rangle e^{i\lambda_j t} = \frac{1}{\sqrt{2}} e^{i\lambda_j t}$ .
  - **P**(observe state *B* at time *t*) =  $|\langle \Psi_t | B \rangle|^2 = |\frac{1}{\sqrt{2}} \langle \Psi_t | E_1 \rangle + \frac{1}{\sqrt{2}} \langle \Psi_t | E_2 \rangle|^2$ .
  - This is the absolute square of a complex number.
  - We can compute it, using rules of complex numbers. e.g., for  $a, b \in \mathbf{R}$ ,  $|a+ib|^2 = a^2 + b^2$ , and  $e^{ia} = \cos(a) + i\sin(a)$ , and  $|e^{ia}| = 1$ .

- Compute that **P**(observe state *B* at time t) =  $|\frac{1}{2}e^{i\lambda_1t} + \frac{1}{2}e^{i\lambda_2t}|^2 = \frac{1}{4}|e^{i\lambda_1t}| |1 + e^{i(\lambda_2 \lambda_1)t}|^2 = \frac{1}{4}(1) |1 + \cos((\lambda_2 \lambda_1)t) + i\sin((\lambda_2 \lambda_1)t)|^2 = \frac{1}{4} [(1 + \cos((\lambda_2 \lambda_1)t))^2 + (\sin((\lambda_2 \lambda_1)t))^2] = \frac{1}{4} [1 + 1 + 2\cos((\lambda_2 \lambda_1)t))] = \frac{1}{2} [1 + \cos((\lambda_2 \lambda_1)t))].$
- That is, the probability fluctuates between 1 (at times 0,  $2\pi/(\lambda_2 \lambda_1)$ , etc.), and 0 (at times  $\pi/(\lambda_2 \lambda_1)$ , etc.). (Probability "phase".)
  - This is, apparently, actually observed in physics labs!
- Aside: If we observe the system in state B at time s, then the system immediately resets to B, so that  $\langle \Psi_s | E_1 \rangle = \langle \Psi_s | E_2 \rangle = \frac{1}{\sqrt{2}}$ .
- Note: Actual quantum mechanics is more complicated: the wave function  $\Psi(x,t)$  can depend on all spatial points x, not just on two different eigenstates corresponding to one single observable state.
  - Then the more general Schrödinger Wave Equation is the partial differential equation:  $i\frac{h}{2\pi}\frac{\partial}{\partial t}\Psi(x,t) = \hat{H}\Psi(x,t)$ , where  $i=\sqrt{-1}$ , h is Planck's constant (a very small constant), and  $\hat{H}$  is a "Hamiltonian" energy operator (complicated).
  - The above example is the very special case where x takes on just two values (for the two states), and  $\hat{H}$  is a 2 × 2 diagonal matrix with diagonal entries  $\lambda_1$  and  $\lambda_2$ .

# Renewal Theory:

- Like for Poisson Processes, have "arrival times"  $\{T_n\}$ .
- Here  $T_0 = 0$ , and  $T_n = Y_1 + Y_2 + \ldots + Y_n$ , where  $\{Y_n\}_{n=1}^{\infty}$  are independent interarrival times, with  $\{Y_n\}_{n=2}^{\infty}$  i.i.d.
- Then  $N(t) = \max\{n \geq 0; T_n \leq t\} = \#\{n \geq 1; T_n \leq t\}$  is a "renewal process", on the state space  $S = \{0, 1, 2, \ldots\}$ .
  - If  $\{Y_n\}_{n=1}^{\infty}$  are <u>all</u> i.i.d., then the process is <u>zero-delayed</u> (or <u>pure</u> or <u>ordinary</u>).
  - If  $\{Y_n\}$  are i.i.d.  $\sim \text{Exp}(\lambda)$ , then  $\{N(t)\}$  is Markovian (because of the "memoryless property"), and in fact  $\{N(t)\}$  is a Poisson process (already studied).
  - But for other distributions of the  $Y_i$ , usually  $\{N(t)\}$  is <u>not</u> Markovian, e.g. perhaps  $\mathbf{P}[N(11) = 4 \mid N(10) = 3] \neq \mathbf{P}[N(11) = 4 \mid N(10) = 3, N(5) = 3].$
  - In fact, Poisson processes are the <u>only</u> Markovian renewal processes;
     see Grimmett & Stirzaker, 2nd ed., Theorem 8.3(5).
- EXAMPLE: Suppose we replace a light-bulb whenever it burns out.
  - Let  $T_n$  be the  $n^{\text{th}}$  time we replace it.  $(T_0 = 0)$
  - Then  $Y_n = T_n T_{n-1}$  is the <u>lifetime</u> of the  $n^{\text{th}}$  bulb.

- If the bulbs are identical, then  $\{Y_n\}_{n=2}^{\infty}$  are i.i.d.
- Let N(t) be the number of bulbs replaced by time t.
- Then  $\{N(t)\}$  is a renewal process.
- If we <u>started</u> with a <u>fresh</u> bulb, then  $\{Y_n\}_{n=1}^{\infty}$  are all i.i.d., so  $\{N(t)\}$  is a "zero-delayed" renewal process. Otherwise probably not.
- Similarly for fixing a machine, etc.
- EXAMPLE: Let  $\{Q(t)\}$  be a single-server queue.
  - Let  $T_0 = 0$ , and let  $T_n$  be the  $n^{\text{th}}$  time the queue <u>empties</u>. (i.e., the  $n^{\text{th}}$  time s such that Q(s) = 0 but  $\lim_{t \nearrow s} Q(t) > 0$ )
  - Let  $Y_n = T_n T_{n-1}$  be the time between queue emptyings.
  - Let  $N(t) = \#\{n \ge 1; T_n \le t\}$  be the number of queue emptyings by time t.
  - Then  $\{N(t)\}$  is a renewal process.
  - (Probably not zero-delayed ... unless start <u>at</u> an emptying time ... or start empty <u>and</u> have exponential interarrival times ...)
- EXAMPLE: Let  $\{X_n\}_{n=0}^{\infty}$  be an irreducible recurrent discrete-time Markov chain on a discrete state space S.
  - Let  $i, j \in S$ , and assume  $X_0 = j$ .
  - Let  $T_0 = 0$ , and for  $n \ge 2$ , let  $T_n = \min\{t > T_{n-1} : X(t) = i\}$  be the time of the  $n^{\text{th}}$  visit to the state i. (integer-valued)
  - Let  $Y_n = T_n T_{n-1}$ .
  - Then for  $y = 1, 2, \ldots$ , we have  $\mathbf{P}(Y_1 = y) = f_{ji}^{(y)}$ , where

$$f_{ji}^{(n)} := \mathbf{P}_j(X_n = i, \text{ but } X_m \neq i \text{ for } 1 \leq m \leq n-1),$$

so 
$$\sum_{y=1}^{\infty} f_{ji}^{(y)} = f_{ji} = 1.$$

- Similarly for  $n \geq 2$ ,  $\mathbf{P}(Y_n = y) = f_{ii}^{(y)}$ .
- So, if  $N(t) = \#\{n \ge 1; T_n \le t\}$  is the number of visits to i by time t, then  $\{N(t)\}$  is a renewal process.
- If j = i, then the process is zero-delayed, otherwise probably not.
- EXAMPLE: Let  $\{X_n\}_{n=0}^{\infty}$  be an irreducible recurrent <u>continuous</u>-time Markov process on a discrete state space S.
  - Let  $i, j \in S$ , and assume X(0) = j.
  - Let  $T_0 = 0$ , and for  $n \geq 2$ , let  $T_n = \min\{t > T_{n-1} : X(t) = i\}$  be the time of the  $n^{\text{th}}$  arrival at the state i, i.e. the  $n^{\text{th}}$  time s such that Q(s) = i but there is a sequence of times  $t_1, t_2, \ldots \nearrow s$  with  $Q(t_m) \neq i$ .
  - Let  $Y_n = T_n T_{n-1}$  be the time between arrivals at i.
  - Let  $N(t) = \#\{n \ge 1; T_n \le t\}$  be the number of arrivals to i by time t.

- Then  $\{N(t)\}$  is a renewal process.
- If j = i, then the process is zero-delayed, otherwise probably not.
- ELEMENTARY RENEWAL THEOREM: For a renewal process as above, if  $\mathbf{P}(Y_1 < \infty) = 1$ , and if the "mean interarrival time"  $\mu := \mathbf{E}(Y_2) < \infty$ , then (a)  $\lim_{t \to \infty} \frac{N(t)}{t} = 1/\mu$  w.p. 1, and (b)  $\lim_{t \to \infty} \frac{\mathbf{E}[N(t)]}{t} = 1/\mu$ .
- PROOF of (a):
  - By the Strong Law of Large Numbers (SLLN),  $\lim_{n\to\infty} \frac{1}{n} T_n = \lim_{n\to\infty} \frac{1}{n} (Y_1 + Y_2 + \ldots + Y_n) = \mu$  w.p. 1.
  - Also,  $Y_n < \infty$  w.p. 1 (for n = 1 by assumption,  $n \ge 2$  since  $\mu < \infty$ ).
  - Therefore,  $\lim_{t\to\infty} N(t) = \infty$  w.p. 1.
  - Hence,  $\lim_{t\to\infty} \frac{1}{N(t)} T_{N(t)} = \lim_{n\to\infty} \frac{1}{n} T_n = \mu$  w.p. 1.
  - But  $N(t) = \max\{n \geq 0; T_n \leq t\}$ . So,  $T_{N(t)} \leq t < T_{N(t)+1}$ . Hence,

$$\frac{T_{N(t)}}{N(t)} \ \leq \ \frac{t}{N(t)} \ < \ \frac{T_{N(t)+1}}{N(t)} \ = \ \frac{T_{N(t)+1}}{N(t)+1} \ \frac{N(t)+1}{N(t)} \, .$$

- As  $t \to \infty$ , w.p. 1,  $\frac{T_{N(t)}}{N(t)} \to \mu$ , and  $\frac{T_{N(t)+1}}{N(t)+1} \to \mu$ , and  $\frac{N(t)+1}{N(t)} \to 1$ .
- So,  $\frac{t}{N(t)} \rightarrow \mu$  w.p. 1 (by "sandwich theorem").
- So,  $\frac{N(t)}{t} \to 1/\mu$  w.p. 1, Q.E.D.
- PROOF of (b):
  - This is similar to part (a), but it's a bit more subtle (since N(t)/t is neither bounded nor monotone).
  - So, we won't prove it here. Instead see e.g. Grimmett & Stirzaker,
     2nd ed., p. 397, or Resnick, p. 191.
- EXAMPLE: Let  $\{X_n\}_{n=0}^{\infty}$  be an irreducible recurrent discrete-time Markov chain on a discrete state space S.
  - Again let N(t) be the number of visits to the state i by time t.
  - Then  $\{N(t)\}$  is a renewal process.
  - Here  $\mu = m_i$  = the mean return time for the state i.
  - Hence, by the Elementary Renewal Theorem part (a),  $N(t)/t \to 1/m_i$ .
  - But here if t = n, then  $N(n) = \#\{k \le n : X_k = i\} = \sum_{k=1}^n \mathbf{1}_{X_k = i}$ .
  - Hence,  $\lim_{n\to\infty} \frac{1}{n} \sum_{k=1}^n \mathbf{1}_{X_k=i} = 1/m_i$ .
  - Same as we already showed previously!
  - (And, since  $\lim_{n\to\infty}\frac{1}{n}\sum_{k=1}^n\mathbf{1}_{X_k=i}=\pi_i$ , this showed that  $\pi_i=1/m_i$ .)
  - Furthermore, by the Elementary Renewal Theorem part (b), or by the Bounded Convergence Theorem since in this case  $N(t) \leq t$  so

 $N(t)/t \leq 1$ , we also have  $\mathbf{E}[N(t)]/t \to 1/m_i$ .

• BLACKWELL RENEWAL THEOREM: For a renewal process as above, suppose  $\mathbf{P}(Y_1 < \infty)$ , and  $\mu := \mathbf{E}(Y_2) < \infty$ , and  $Y_2$  is "not arithmetic", i.e. there is  $\underline{\mathbf{no}} \ \lambda > 0$  such that  $\mathbf{P}(Y_2 = k\lambda \text{ for some } k \in \mathbf{Z}) = 1$ . (Similar to "aperiodicity".) Then for any fixed h > 0,

$$\lim_{t\to\infty} \mathbf{E}[N(t+h) - N(t)] = h/\mu.$$

- This is a "more refined" theorem, since it doesn't just consider the overall average N(t)/t, but rather considers the specific number of renewals between times t and t + h.
- For a PROOF, see e.g. Grimmett & Stirzaker, 2nd ed., pp. 408–409, or Resnick, Section 3.10.3.
- PRACTICE PROBLEM: Let  $Y_1, Y_2, \ldots$  be i.i.d.  $\sim$  Uniform[0, 10]. Let  $T_0 = 0$ , and  $T_n = Y_1 + Y_2 + \ldots + Y_n$  for  $n \geq 1$ . Let  $N(t) = \max\{n \geq 0 : T_n < t\}$ . (a) Compute (with explanation)  $\lim_{t \to \infty} N(t)/t$ .
  - (b) Approximate (with explanation)  $\mathbf{E}(\#\{n \geq 1: 1234 < T_n < 1236\})$ .

#### Renewal Reward Processes:

- Consider a renewal process as above, with  $P(Y_1 < \infty) = 1$ .
- Suppose the mean interarrival time  $\mu := \mathbf{E}(Y_2) < \infty$ .
- Suppose at the  $k^{\text{th}}$  renewal time  $T_k$ , you receive an "reward" (or cost)  $R_k$ .
  - Assume the  $\{R_k\}$  are i.i.d.
- Let  $R(t) = \sum_{k=1}^{N(t)} R_k$  be the total reward received by time t.
- RENEWAL REWARD THEOREM:  $\lim_{t\to\infty}\frac{R(t)}{t}=\frac{\mathbf{E}[R_1]}{\mu}$  w.p. 1.
  - PROOF: Here

$$\frac{R(t)}{t} = \frac{\sum_{k=1}^{N(t)} R_k}{t} = \frac{\sum_{k=1}^{N(t)} R_k}{N(t)} \frac{N(t)}{t}.$$

- As  $t \to \infty$ , w.p. 1,  $\frac{\sum_{k=1}^{N(t)} R_k}{N(t)} \to \mathbf{E}[R_1]$  by SLLN, and  $\frac{N(t)}{t} \to 1/\mu$  by the Elementary Renewal Theorem part (a).
- Hence,  $\frac{R(t)}{t} \to \mathbf{E}[R_1] \times (1/\mu) = \mathbf{E}[R_1]/\mu$ , Q.E.D.

# **Application – Car Purchases:**

- Suppose each new car's lifetime L has distribution Uniform[0,10] (in years), after which it breaks down.
- Suppose your strategy is to buy a new car as soon as your old car breaks down, or after S years if the car hasn't broken down by then.  $(0 \le S \le 10)$
- Suppose a new car costs 30 (in thousands of dollars).

- Suppose further that if your car breaks down <u>before</u> you sell it, then that costs you an extra 5 (in thousands of dollars).
- QUESTION: On average, how many cars will you buy each year?
- SOLUTION:
  - Here  $\mu = \mathbf{E}(Y_2) = \mathbf{E}[\min(L, S)].$
  - So,  $\mu = S \mathbf{P}(L > S) + \mathbf{E}[L \mid L < S] \mathbf{P}(L < S)$ =  $S[(10 - S)/10] + (S/2)(S/10) = S - (S^2/20)$ .
  - So, by the Elementary Renewal Theorem part (a),  $\lim_{t\to\infty} N(t)/t = 1/\mu = 1/[S (S^2/20)]$ .
  - $-\,$  So, you will buy an average of  $1/[S-(S^2/20)]$  cars per year, i.e. one car every  $[S-(S^2/20)]$  years.
  - e.g. if S=10, then you will buy an average of one car every 5 years. (Of course.)
  - e.g. if S = 9, then you will buy an average of one car every 99/20 = 4.95 years.
  - Or, if S = 1, then you will buy an average of one car every 19/20 = 0.95 years.
- QUESTION: About how many cars will you buy between times 562 and 566 (in years)?

### • SOLUTION:

- To apply the Blackwell Renewal Theorem, we want t = 562, and t + h = 566 so h = 4.
- Hence, by the Blackwell Renewal Theorem, since t=562 is reasonably large, the expected number of purchases between times 562 and 566 is approximately  $h/\mu=4/[S-(S^2/20)]$ .
- e.g. if S = 9, then this equals about 0.808.
- $\bullet$  QUESTION: What is your long-run average car  $\underline{\cos t}$  per year?
  - (And, what choice of S minimises this?)

#### • SOLUTION:

- Let  $T_k$  be the time (in years) of the purchase of your  $k^{\text{th}}$  car.  $(T_0 = 0)$
- Let  $Y_k = T_k T_{k-1}$  be the  $k^{\text{th}}$  interarrival time.
- Let R(t) be your total car cost by time t.
- Then this is a renewal reward process!
- Here  $E[R_1] = 30 + 5 \mathbf{P}(L \le S) = 30 + 5 (S/10) = 30 + S/2.$
- Also  $\mu = \mathbf{E}[Y_2] = S (S^2/20)$  as above.

- Hence, w.p. 1, by the Renewal Reward Theorem,

$$\lim_{t \to \infty} \frac{R(t)}{t} \ = \ \frac{\mathbf{E}[R_1]}{\mu} \ = \ \frac{30 + S/2}{S - (S^2/20)} \,.$$

- If S = 10 (i.e., never sell early), this equals  $\frac{35}{5} = 7$ .
- If S = 9, this equals  $\frac{34.5}{99/20} \doteq 6.970$ . Less!
- Minimised when S = 9.282, then this equals about 6.964.
- Plot: www.probability.ca/sta447/Rcar
- CONCLUSION: Your best policy is to buy a new car as soon as your old car breaks, or is 9.282 years old, whichever comes first.

## Discrete-time Markov Chains on Continuous State Spaces:

- Suppose instead of a discrete set S, our state space  $\mathcal{X}$  is now <u>any</u> non-empty (perhaps uncountable) set, e.g.  $\mathbf{R}$ .
- The (one-step) transition probabilities are then given by P(x, A), for each  $x \in \mathcal{X}$  and each (measurable)  $A \subseteq \mathcal{X}$ .
  - So, for each fixed state  $x \in \mathcal{X}$ ,  $P(x,\cdot)$  is a <u>probability measure</u>.
  - (Also, for each fixed measurable  $A \subseteq \mathcal{X}$ , P(x, A) is a measurable function of  $x \in \mathcal{X}$ .)
  - Here P(x, A) is the probability, if the chain is at a point x, that it will jump to <u>somewhere</u> in the subset A at the next step.
- If  $\mathcal{X}$  is countable, then  $P(x, \{i\})$  corresponds to the transition probability  $p_{xi}$  of the discrete Markov chains as before.
  - But on a general state space, might have  $P(x, \{i\}) = 0$  for all  $i \in \mathcal{X}$ .
- We also require an initial distribution  $\nu$ , which is any probability distribution on  $(\mathcal{X}, \mathcal{F})$ .
- We then have a (discrete-time, general state space, time-homogeneous)  $Markov\ chain\ X_0, X_1, X_2, \ldots$ , where

$$\mathbf{P}(X_0 \in A_0, X_1 \in A_1, \dots, X_n \in A_n)$$

$$= \int_{x_0 \in A_0} \nu(dx_0) \int_{x_1 \in A_1} P(x_0, dx_1) \dots$$

$$\dots \int_{x_{n-1} \in A_{n-1}} P(x_{n-2}, dx_{n-1}) \int_{x_n \in A_n} P(x_{n-1}, dx_n).$$

- If we want, we can simplify  $\int_{x_n \in A_n} P(x_{n-1}, dx_n) = P(x_{n-1}, A_n)$ .
- But cannot further simplify e.g.  $\int_{x_{n-1} \in A_{n-1}} P(x_{n-2}, dx_{n-1}) P(x_{n-1}, A_n)$ .

## END OF WEEK #11 -

- EXAMPLE: Consider the Markov chain on the real line (i.e. with  $\mathcal{X} = \mathbf{R}$ ), where  $P(x, \cdot) = N(\frac{x}{2}, \frac{3}{4})$  for each  $x \in \mathcal{X}$ .
  - Simulation: www.probability.ca/sta447/RconMCex
  - Equivalently,  $X_{n+1} = \frac{1}{2}X_n + Z_{n+1}$ , where  $\{Z_n\}$  are i.i.d. with  $Z_n \sim N(0, \frac{3}{4})$ .
  - Stationary distribution? Convergence? (Soon!)
- Similar to before, write  $\mathbf{P}_x(\cdots)$  for the probability of an event conditional on  $X_0 = x$ , i.e. under the assumption that the initial distribution  $\nu$  is a point-mass at the single state x.
- And, define higher-order transition probabilities inductively by  $P^1(x, A) = P(x, A)$ , and  $P^{n+1}(x, A) = \int_{z \in \mathcal{X}} P(x, dz) P^n(z, A)$  for  $n \geq 1$ .
- A <u>stationary distribution</u> for such a Markov chain is a probability measure  $\pi(\cdot)$  on  $\mathcal{X}$ , such that  $\pi(A) = \int_{\mathcal{X}} \pi(dx) P(x, A)$  for all  $A \subseteq \mathcal{X}$ .
  - (This generalises our earlier definition  $\pi_j = \sum_{i \in S} \pi_i p_{ij}$ .)
  - Like in the discrete case, Markov chains on general state spaces may or may not have stationary distributions.
- How to define a concept like "irreducible"?
- Problem: often (e.g. in the above example),  $p_{ij}^{(n)} = 0$  for all  $i, j \in \mathcal{X}$  and all  $n \geq 1$ . What to do?
- Let  $\tau_A = \inf\{n \geq 0; \ X_n \in A\}$  be the first hitting time of the subset A.
  - Thus,  $\tau_A < \infty$  iff the chain eventually hits the subset A.
- DEFN: a Markov chain on a general state space  $\mathcal{X}$  is  $\phi$ -irreducible if there is a non-zero ( $\sigma$ -finite) measure  $\psi$  on  $\mathcal{X}$  such that if  $\psi(A) > 0$ , then  $\mathbf{P}_x(\tau_A < \infty) > 0$  for all  $x \in \mathcal{X}$ .
  - That is, the chain has positive probability of eventually hitting any subset A of positive  $\psi$  measure.
  - Common choice:  $\psi$  = Lebesgue (length) measure on  $\mathbf{R}$ .
- And, how to define "period"? "aperiodicity"? Can't use "gcd" defn!
- DEFN: the <u>period</u> of a general-state-space Markov chain is the largest (finite) positive integer d such that there are non-empty disjoint subsets  $\mathcal{X}_1, \ldots, \mathcal{X}_d \subseteq \mathcal{X}$ , with  $P(x, \mathcal{X}_{i+1}) = 1$  for all  $x \in \mathcal{X}_i$  ( $1 \le i \le d-1$ ) and  $P(x, \mathcal{X}_1) = 1$  for all  $x \in \mathcal{X}_d$ . ("forced cycle") (diagram)
- DEFN: The chain is <u>aperiodic</u> if its period (as above) equals 1.
- GENERAL STATE SPACE MARKOV CONVERGENCE THEOREM: If a discrete-time Markov chain on a general state space is  $\phi$ -irreducible and aperiodic, and has a stationary distribution  $\pi(\cdot)$ , then for  $\pi$ -almost every  $x \in \mathcal{X}$ ,

$$\lim_{n \to \infty} \sup_{A \in \mathcal{F}} |P^n(x, A) - \pi(A)| = 0.$$

- That is, the Markov chain converges to its stationary distribution in probability (and in "total variation distance").
- For a proof see e.g. my review paper, or the Meyn and Tweedie book.

# Application – Convergence of the Normal Example:

- Consider again the Markov chain from the above Example, with  $\mathcal{X} = \mathbf{R}$ , and with  $P(x,\cdot) = N(\frac{x}{2}, \frac{3}{4})$  for each  $x \in \mathcal{X}$ . Let  $\pi(\cdot) = N(0,1)$  be the standard normal distribution.
- CLAIM:  $\pi(\cdot)$  is a stationary distribution for this chain.
  - PROOF: One way to think of the chain is, to get from  $X_{n-1}$  to  $X_n$ , we first <u>divide</u>  $X_{n-1}$  by 2, and then <u>add</u> an independent  $N(0, \frac{3}{4})$  random variable.
  - That is,  $X_n = \frac{X_{n-1}}{2} + Z_n$ , where  $\{Z_n\}$  are i.i.d.  $\sim N(0, \frac{3}{4})$ .
  - Now, if  $X_{n-1} \sim \pi = N(0,1)$ , then  $\frac{X_{n-1}}{2} \sim N(0,\frac{1}{4})$ , so  $X_n = \frac{X_{n-1}}{2} + Z_n \sim N(0,\frac{1}{4}+\frac{3}{4}) = N(0,1) = \pi$ .
  - Hence,  $\pi$  is stationary for this chain, Q.E.D.
- CLAIM: This Markov chain is  $\phi$ -irreducible, where  $\psi$  is the usual Lebesgue (length) measure on  $\mathbf{R}$ .
  - PROOF: Indeed, if  $x \in \mathbf{R}$ , and if  $A \subseteq \mathbf{R}$  has positive Lebesgue measure (essentially, has positive length), then P(x, A) > 0.
  - Hence, chain is  $\phi$ -irreducible [in fact, with  $\mathbf{P}(\tau_A = 1) > 0$ ], Q.E.D.
- CLAIM: This Markov chain is aperiodic.
  - PROOF (outline): If instead it had period  $d \geq 2$ , then for  $x_1 \in \mathcal{X}_1$  and  $x_2 \in \mathcal{X}_2$ , we would have  $P(x_1, \mathcal{X}_2) = 1$  while  $P(x_2, \mathcal{X}_2) = 0$ .
  - However, if a subset  $A \subset \mathbf{R}$  has positive Lebesgue measure (i.e., positive length), then as above, P(x, A) > 0 for all  $x \in \mathbf{R}$ .
  - Or, if a subset  $A \subset \mathbf{R}$  has zero Lebesgue measure (i.e., zero length), then P(x, A) = 0 for all  $x \in \mathbf{R}$ .
  - So, there is <u>no</u> subset  $A \subset \mathbf{R}$  with  $P(x_1, A) > 0$  but  $P(x_2, A) = 0$ .
  - Contradiction! So, period=1, i.e. it is aperiodic, Q.E.D.
- HENCE, by the General State Space Markov Chain Convergence Theorem, for  $\pi$ -almost every  $x \in \mathcal{X}$ ,  $\lim_{n\to\infty} \sup_{A\in\mathcal{F}} |P^n(x,A) \pi(A)| = 0$ .
  - In particular,  $P^n(x, A) \to \pi(A)$ .
  - So, e.g.,  $\lim_{n\to\infty} \mathbf{P}(X_n < 2) = \pi\{(-\infty, 2)\} = \Phi(2)$ , where  $\Phi$  is the standard normal c.d.f.

• PRACTICE PROBLEM: Consider a discrete-time Markov chain with state space  $\mathcal{X} = \mathbf{R}$ , and with transition probabilities such that  $P(x,\cdot)$  is uniform on the interval [x-1, x+1]. Determine whether or not this chain is  $\phi$ -irreducible.

#### • PRACTICE PROBLEM:

- (a) Prove that a Markov chain on a <u>countable</u> state space  $\mathcal{X}$  is  $\phi$ irreducible if and only if there is  $j \in \mathcal{X}$  such that  $\mathbf{P}_i(\tau_j < \infty) > 0$  for all  $i \in \mathcal{X}$ , i.e. such that j can be reached from any state i.
- (b) Give an example of a Markov chain on a countable state space which is  $\phi$ -irreducible as above, but which is <u>not</u> irreducible according to our previous (discrete state space) definition.

## Convergence From Where?:

- QUESTION: Why does Theorem say " $\pi$ -almost every"  $x \in \mathcal{X}$ ?
- EXAMPLE:
  - State space  $\mathcal{X} = \{1, 2, 3, \ldots\}$  (actually discrete). (diagram)
  - Transitions  $P(1,\{1\}) = 1$ , and for  $x \ge 2$ ,  $P(x,\{1\}) = 1/x^2$  and  $P(x,\{x+1\}) = 1 (1/x^2)$ .
  - Stationary distribution?  $\pi\{1\} = 1$  (of course).
  - $\phi$ -irreducible? Yes! If  $\pi(S) > 0$ , then  $1 \in S$ , so  $P(x, S) \ge P(x, \{1\}) > 0$  for all  $x \in \mathcal{X}$ . So, can take  $\psi = \pi$ .
  - Aperiodic? Yes, since e.g.  $P(1, \{1\}) > 0$ .
  - So, by Theorem, for  $\pi$ -a.e.  $x \in \mathcal{X}$ , have  $\lim_{n \to \infty} P(x, S) = \pi(S)$ .
  - That is,  $\lim_{n\to\infty} \mathbf{P}_x(X_n=1)=1$ .
  - But if  $X_0 = x \ge 2$ , then  $\mathbf{P}_x[X_n = x + n \text{ for all } n] = \prod_{j=x}^{\infty} (1 (1/j^2))$ .
  - This infinite product is positive, since  $\sum_{j=x}^{\infty} (1/j^2) < \infty$ . (optional)
  - This means that  $\lim_{n\to\infty} \mathbf{P}_x(X_n=1) \neq 1$ . Contradiction??
  - No! Here, convergence holds if x = 1, which is  $\pi$ -a.e. since  $\pi(1) = 1$ , but <u>not</u> if  $x \ge 2$ .
- So, the convergence can be subtle! But it <u>usually</u> holds from <u>any</u>  $x \in \mathcal{X}$  ("Harris recurrent"); see e.g. my Harris recurrence paper.

### END OF WEEK #12

- REMINDER: FINAL EXAM! (And office hours.)
- Good luck and best wishes! J.R.