

Lecture 13  
Feb. 26th, 2015

## PROGRAM CORRECTNESS

$b^n$  for  $b > 0, n \in \mathbb{N}$

$$236^{23} = 236 \cdots 236 = (236^6)(236^6) \cdot 236 = (236^6)^2 \cdot 236 = ((236^3)^2 \cdot 236)^2 \cdot 236 = (((236^1)^2)^2 \cdot 236)^2 \cdot 236$$

in python:

# pre:  $n \in \mathbb{N}, b \in \mathbb{R}$  and  $b \neq 0$  in case  $0^0$  bothers you.  
# post: return  $b^n$ .

```
def pow(b, n):
    if n >= 1:
        p = pow(b, n // 2)
        if n % 2 == 1:
            return p * p * b
        else:
            return p * p
    else:
        return 1
```

For  $n \in \mathbb{N}$ , Let  $Q(n)$  be: IF pre( $n, b$ ), THEN post( $n, b$ )

If  $b \in \mathbb{R}$  and  $b \neq 0$  then  $\text{pow}(b, n)$  returns  $b^n$ .

Base Case:  $n = 0$ . suppose  $b \in \mathbb{R}, b \neq 0$ .

From code,  $\text{pow}(b, n)$  strings 1st branch, returns  $1 = b^0 = b^n$

IS: let  $n \in \mathbb{N}$ , assume  $n \geq 1$

IH: Assume  $Q$  for all natural numbers less than  $n$ .

Suppose  $b \in \mathbb{R}, b \neq 0$

From code and  $n \geq 1$ : calls  $\text{pow}(b, \lfloor \frac{n}{2} \rfloor)$

Since  $n > 0$ :  $\frac{n}{2} < n$ , so  $\lfloor \frac{n}{2} \rfloor < n$

also  $\lfloor \frac{n}{2} \rfloor \in \mathbb{Z}$  and  $\frac{n}{2} \geq \frac{1}{2} \geq 0$

So  $\lfloor \frac{n}{2} \rfloor \in \mathbb{N}$

So  $Q(\lfloor \frac{n}{2} \rfloor)$  by IH  $\wedge (n \in \mathbb{N}, b \in \mathbb{R}, b \neq 0)$

So PRE is true for  $\text{pow}(b, \lfloor \frac{n}{2} \rfloor)$ , so  $\text{pow}(b, \lfloor \frac{n}{2} \rfloor)$  returns  $b^{\lfloor \frac{n}{2} \rfloor}$  so  $p = b^{\lfloor \frac{n}{2} \rfloor}$

If  $n$  is odd: 1st inner branch of First Branch

returns  $p * p * b = b^{\lfloor \frac{n}{2} \rfloor} b^{\lfloor \frac{n}{2} \rfloor} b = b^{\frac{n-1}{2}} b^{\frac{n-1}{2}} b = b^{n-1} b = b^n$

If  $n$  is even: returns  $p * p = b^{\lfloor \frac{n}{2} \rfloor} b^{\lfloor \frac{n}{2} \rfloor} = b^{\frac{n}{2}} b^{\frac{n}{2}} = b^n$

```
def pow(b, n):
    m = 0
    r = 1
    while m < n:
        r = r * b
        m = m + 1
    return r
```

For  $i \in \mathbb{N}$ , let  $r_i, m_i$  be the values of  $r, m$  after  $i$  iteration.

$r_0 = 1, m_0 = 0$ .

$r_{i+1} = r_i \cdot b, m_{i+1} = m_i + 1$  for each  $i \in \mathbb{N}$

So finally  $m_i = i, r_i = b^i$

For  $i \in \mathbb{N}$ , let  $I(i)$  be : if there are more than  $i$  iterations, then  $m_i = i, r_i = b^i$ .

recursive  $\Rightarrow$  complete induction  
loops  $\Rightarrow$  simple induction

Proof by Simple Induction

$I(0): m_0 = 0, r_0 = 1 = b^0$

IS: let  $i \in \mathbb{N}$ , assume  $I(i)$

Assume more than  $i+1$  iterations, so more than  $i$  iterations.

So by  $I(i): m_i = i, r_i = b^i$ .

So  $r_{i+1} = b \cdot r_i = b^i \cdot b = b^{i+1}$  and  $m_{i+1} = m_i + 1 = i + 1$