The extreme point theorem (Theorem 1.7)

Let S be the feasible region of a linear programming problem.

() If S is non-empty and bounded, the problem has an optimal solution at an extreme point of S.

This is the contrepositive of 3 If the public has an optimal solution then either Sis empty or S is unbounded.

(2 (correction): If S has an extreme point, and S is unbounded but the problem has an optimal solution, then it has an optimal solution at an extreme point.

\$1.5 Basic Solutions

Consider the system Ax = b where A is an $m \times n$ motrix having columns A_1, \cdots, A_n , $X \in \mathbb{R}^n$, and $b \in \mathbb{R}^m$

There is a basic solution having basic variables Xi. -. . Xim if and only if Ai, ... Aim are linearly independent.

Defn: The basic solution having basic variables Xi,....Xim. is (provided Ai,... Aim are linearly independent). The unique solution of Ax=b where Xj=0 when $j \neq i_1, \dots, j \neq i_m$.

Defn: The xj, when j ≠ i, ..., j ≠ im are non-basic variable.

Remark: Non-basic variables are always 0.

Basic variables can be positive, negetire, or 0.

Definition: Consider the canonical constraints Ax=b

A basic, positive solution of this system is a basic solution of Ax=6, where x>0 €1R"

Theorem: Let S be the solution set of Ax=b

X is extreme in S if and only if X is a basic feasible positive Solution for the constraints.

Note: Theorem 1.8: "if"

Theorem 1.9: "only if"