

Category Theory (Continued)

Example 61.2.2 (Free and forgetful functors)

Note that these are both informal terms, and don't have a rigid definition.

- (a) We talked about a **forgetful functor** earlier, which takes the underlying set of a category like Vect_k . Let's call it $U : \mathsf{Vect}_k \to \mathsf{Set}$.
 - Now, given a map $T: V_1 \to V_2$ in Vect_k , there is an obvious $U(T): U(V_1) \to U(V_2)$ which is just the set-theoretic map corresponding to T.
 - Similarly there are forgetful functors from Grp , CRing , etc., to Set . There is even a forgetful functor $\mathsf{CRing} \to \mathsf{Grp}$: send a ring R to the abelian group (R,+). The common theme is that we are "forgetting" structure from the original category.
- (b) We also talked about a **free functor** in the example. A free functor $F : \mathsf{Set} \to \mathsf{Vect}_k$ can be taken by considering F(S) to be the vector space with basis S. Now, given a map $f : S \to T$, what is the obvious map $F(S) \to F(T)$? Simple: take each basis element $s \in S$ to the basis element $f(s) \in T$.

Similarly, we can define $F:\mathsf{Set}\to\mathsf{Grp}$ by taking the free group generated by a set S.

Remark 61.2.3 — There is also a notion of "injective" and "surjective" for functors (on arrows) as follows. A functor $F: \mathcal{A} \to \mathcal{B}$ is **faithful** (resp. **full**) if for any $A_1, A_2, F: \operatorname{Hom}_{\mathcal{A}}(A_1, A_2) \to \operatorname{Hom}_{\mathcal{B}}(FA_1, FA_2)$ is injective (resp. surjective).

We can use this to give an exact definition of concrete category: it's a category with a faithful (forgetful) functor $U: A \to Set$.

之角很多次提到3 concrete Category 都没去理它。直港上就是比较符合认识的对象组成的 Cat. 像 Grp. Cring, 首先的j都是基于集合的对象,然后Mor就是一般脑子里想到的mapping。 标准定义就是说,对于 concrete cat. 中两个 ob) 之间的 arrow,都唯一对应着 Set 中这两个obj 的 underlying sets 之间的态射;这是显然的 (从例子看),很显然无论也可元素对应关系 preserve. 至于forgetful,这更显然。

Example 61.2.4 (Functors from G)

Let G be a group and $\mathcal{G} = \{*\}$ be the associated one-object category.

(a) Consider a functor $F: \mathcal{G} \to \mathsf{Set}$, and let S = F(*). Then the data of F corresponds to putting a group action of G on S.

^aAgain, experts might object that $\operatorname{Hom}_{\mathcal{A}}(A_1, A_2)$ or $\operatorname{Hom}_{\mathcal{B}}(FA_1, FA_2)$ may be proper classes instead of sets, but I am assuming everything is locally small.

g =	$ \begin{array}{ccc} & & & & & & & \\ & & & & & & & \\ & & & & $	() F(e)
	a. C* a. Fra	FCay
ed & n_ 1, 5	74.00 11 11 11 11 11 11	
	作用。 Homg L*,*) = G's underlying set = {(5) = {FCe), FCa,1, { > 映射的疑合	e, a,, a ₂ S
	お紹介(日日) · · · · · · · · · · · · · · · · · ·	:a (xes.aeg), Jusp. H :
	(Fa)(x),又易和Fe)=ids,那么(x,e) H	
_	杂件, Cresp. 左群作刚,因此下可看作一个释作	•
,	H be a group and construct \mathcal{H} the same spond to homomorphisms $G \to H$.	ne way. Then functors $\mathcal{G} \to \mathcal{H}$
监上问	l还是很直觉的。唯一要想的就是保证 ♂ca.	X6 a2) = 5(a.) X4 5(a2)
gond, st	突然想起来 functor 的这义:	
The fu	这些提起来 functor 的这之: nctor respects composition: if $A_1 \stackrel{f}{ o} A$	$_2 \xrightarrow{g} A_3$ are arrows in \mathcal{A} , then
The function $F(g \circ g)$	文然提起来 functor 的定义: $\operatorname{nctor}\ \operatorname{respects}\ \operatorname{composition}\colon \ \operatorname{if}\ A_1 \xrightarrow{f} A_0 = F(g) \circ F(f).$	
The function $F(g \circ g)$	这些提起来 functor 的这之: nctor respects composition: if $A_1 \stackrel{f}{ o} A$	
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	文然提起来 functor 的定义: $\operatorname{nctor}\ \operatorname{respects}\ \operatorname{composition}\colon \ \operatorname{if}\ A_1 \xrightarrow{f} A_0 = F(g) \circ F(f).$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。
The fu F(g o j	突然想起来 functor 的这丈: \mathbf{x} $$	unctor 导出的 J 肯定是同态。