

1. Design a context-free grammar for $\{a^i b^j c^k \mid i \neq j \text{ or } j \neq k\}$, that is, the set of strings of a 's followed by b 's followed by c 's, such that there are either a different number of a 's and b 's or a different number of b 's and c 's, or both.

$$\begin{aligned} S &\rightarrow A_1 C \mid A_2 C \mid AB_1 \mid AB_2 \\ A_1 &\rightarrow aA_1 b \mid aA_1 \mid a \\ A_2 &\rightarrow aA_2 b \mid A_2 b \mid b \\ C &\rightarrow Cc \mid \varepsilon \\ B_1 &\rightarrow bB_1 c \mid bB_1 \mid b \\ B_2 &\rightarrow bB_2 c \mid B_2 c \mid c \\ A &\rightarrow Aa \mid \varepsilon \end{aligned}$$

(注意: $Cc \mid \varepsilon$ 若为 $Cc \mid c$ 则不能产生 a, c 同时为 0 个, 或 b, c)

2. Design a context-free grammar for the set of all strings with twice as many 0's as 1's.

$$\begin{aligned} S &\rightarrow 0S0S1S \mid 0S1S0S \mid 1S0S0S \mid \varepsilon \text{ 或} \\ S &\rightarrow 0S0S1S \mid 0S1S0S \mid 1S0S0S \mid 001 \mid 010 \mid 110 \text{ (但此文法无法产生 } \varepsilon \text{) 或} \\ S &\rightarrow 0S0S1 \mid 0S1S0 \mid 1S0S0 \mid SS \mid \varepsilon \end{aligned}$$

若仅为 $S \rightarrow 0S0S1 \mid 0S1S0 \mid 1S0S0 \mid \varepsilon$ (产生式中缺少 S 的), 则无法产生: 001100 或开头结尾都是 1 的串.

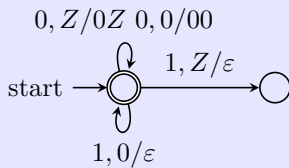
3. Design a context-free grammar for the language consisting of all strings over $\{a, b\}$ that are **not** of the form ww , for some string w . Explain how your grammar works. You needn't prove it's correctness formally.

如果串长为奇数, 显然不是 ww 形式 (对应下面文法中的 A 或 B)。而对于长度为偶数 ($2k$) 的串, 一定有在 j ($1 \leq j \leq k$) 位置和 $(k+j)$ 位置不相同, 且只要有一处即可。可以将偶数串拆分成两个奇数串, 只要中间位置不同, 分别通过 A 和 B 生成。

$$\begin{aligned} S &\rightarrow A \mid B \mid AB \mid BA \\ A &\rightarrow XAX \mid a \\ B &\rightarrow XBX \mid b \\ X &\rightarrow a \mid b \end{aligned}$$

任何偶数串, 拆成两个奇数串, 两个位于中间的字符不相同即可, 如 $aaaabbbb = \underline{aa\check{a}ab} \underline{bbb}$ 或 $aabaaa = \underline{aa\check{b}aa} \underline{\check{a}}$ 。
但 ww 形式无法生成, 拆成任何两个奇数串, 中间部分都相等, 比如 $baabbaab = \underline{baa\hat{b}baa} \underline{\hat{b}} = \underline{b\hat{a}a} \underline{bb\hat{a}ab}$ 。

4. Design a PDA to accept the set of all strings of 0's and 1's such that no prefix has more 1's than 0's.



5. Design a PDA to accept: $\{0^n 1^m \mid n < m < 2n\}$. You may accept either by final state or by empty stack, whichever is more convenient.

