

HEURISTIC OPTIMIZATION

Generalised Local Search Machines

adapted from slides for SLS:FA, Chapter 3

Outline

1. The Basic GLSM Model
2. State and Transition Types
3. Modelling SLS Methods Using GLSMs

The Basic GLSM Model

Many high-performance SLS methods are based on combinations of *simple (pure) search strategies* (e.g., ILS, MA).

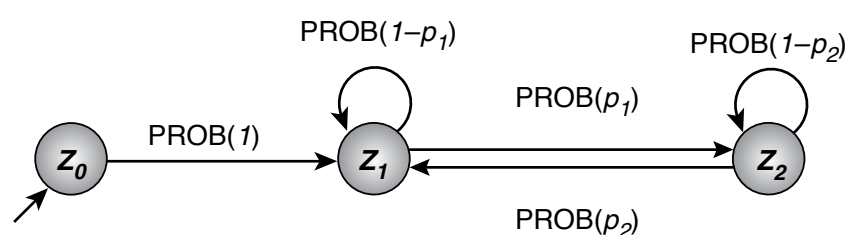
These hybrid SLS methods operate on two levels:

- ▶ **lower level:** execution of underlying simple search strategies
- ▶ **higher level:** activation of and transition between lower-level search strategies.

Key idea underlying Generalised Local Search Machines:

Explicitly represent higher-level search control mechanism in the form of a *finite state machine*.

Example: Simple 3-state GLSM



- ▶ States z_0, z_1, z_2 represent simple search strategies, such as Random Picking (for initialisation), Iterative Best Improvement and Uninformed Random Walk.
- ▶ $\text{PROB}(p)$ refers to a probabilistic state transition with probability p after each search step.

Generalised Local Search Machines (GLSMs)

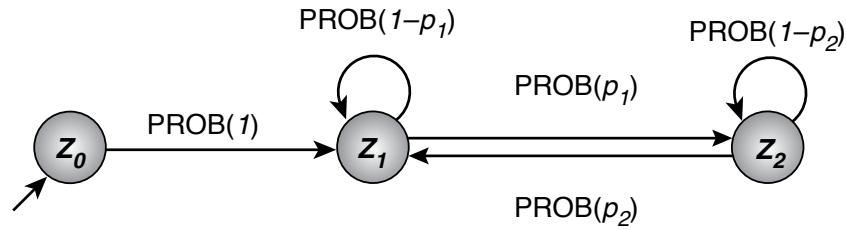
- ▶ States \cong simple search strategies.
- ▶ State transitions \cong search control.
- ▶ GLSM \mathcal{M} starts in initial state.
- ▶ In each iteration:
 - ▶ \mathcal{M} executes one search step associated with its current state z ;
 - ▶ \mathcal{M} selects a new state (which may be the same as z) in a probabilistic manner.
- ▶ \mathcal{M} terminates when a given termination criterion is satisfied.

Formal definition of a GLSM

A *Generalised Local Search Machine* is defined as a tuple $\mathcal{M} := (\mathcal{Z}, z_0, M, m_0, \Delta, \sigma_Z, \sigma_\Delta, \tau_Z, \tau_\Delta)$ where:

- ▶ \mathcal{Z} is a set of states;
- ▶ $z_0 \in \mathcal{Z}$ is the *initial state*;
- ▶ M is a set of *memory states* (as in SLS definition);
- ▶ m_0 is the *initial memory state* (as in SLS definition);
- ▶ $\Delta \subseteq \mathcal{Z} \times \mathcal{Z}$ is the *transition relation*;
- ▶ σ_Z and σ_Δ are sets of *state types* and *transition types*;
- ▶ $\tau_Z : \mathcal{Z} \mapsto \sigma_Z$ and $\tau_\Delta : \Delta \mapsto \sigma_\Delta$ associate every state z and transition (z, z') with a state type $\sigma_Z(z)$ and transition type $\tau_\Delta((z, z'))$, respectively.

Example: Simple 3-state GLSM (formal definition)



- ▶ $Z := \{z_0, z_1, z_2\}$; z_0 = initial machine state
- ▶ no memory ($M := \{m_0\}$; m_0 = initial and only memory state)
- ▶ $\Delta := \{(z_0, z_1), (z_1, z_2), (z_1, z_1), (z_2, z_1), (z_2, z_2)\}$
- ▶ $\sigma_Z := \{z_0, z_1, z_2\}$
- ▶ $\sigma_\Delta := \{\text{PROB}(p) \mid p \in \{1, p_1, p_2, 1 - p_1, 1 - p_2\}\}$
- ▶ $\tau_Z(z_i) := z_i, \quad i \in \{0, 1, 2\}$
- ▶ $\tau_\Delta((z_0, z_1)) := \text{PROB}(1), \tau_\Delta((z_1, z_2)) := \text{PROB}(p_1), \dots$

Example: Simple 3-state GLSM (semantics)

- ▶ Start in initial state z_0 , memory state m_0 (never changes).
 - ▶ Perform one search step according to search strategy associated with state type z_0 (e.g., random picking).
 - ▶ With probability 1, switch to state z_1 .
 - ▶ Perform one search step according to state z_1 ; switch to state z_2 with probability p_1 , otherwise, remain in state z_1 .
 - ▶ In state z_2 , perform one search step according to z_2 ; switch back to state z_1 with probability p_2 , otherwise, remain in state z_2 .
- ~> After one z_0 step (initialisation), repeatedly and probabilistically switch between phases of z_1 and z_2 steps until termination criterion is satisfied.

Note:

- ▶ *State types* formally represent (subsidiary) search strategies, whose definition is not part of the GLSM definition.
- ▶ *Transition types* formally represent mechanisms used for switching between GLSM states.
- ▶ Multiple states / transitions can have the same type.
- ▶ σ_Z, σ_Δ should include only state and transition types that are actually used in given GLSM ('no junk').
- ▶ Not all states in Z may actually be reachable when running a given GLSM.
- ▶ *Termination condition* is not explicitly captured in GLSM model, but considered part of the execution environment.

GLSM Semantics

Behaviour of a GLSM is specified by *machine definition* + *run-time environment* comprising specifications of

- ▶ state types,
- ▶ transition types;
- ▶ problem instance to be solved,
- ▶ search space,
- ▶ solution set,
- ▶ neighbourhood relations for subsidiary SLS algorithms;
- ▶ termination predicate for overall search process.

Run GLSM \mathcal{M} :

set *current machine state* to z_0 ; set *current memory state* to m_0 ;

While *termination criterion* is not satisfied:

 perform *search step* according to type of current machine state;
 this results in a new *search position*

 select *new machine state* according to *types of transitions*
 from *current machine state*, possibly depending on
 search position and *current memory state*; this may
 change the *current memory state*

Note:

- ▶ The *current search position* is only changed by the subsidiary search strategies associated with states, *not* as side-effect of machine state transitions.
- ▶ The *machine state* and *memory state* are only changed by state-transitions, *not* as side-effect of search steps.
(Memory state is viewed as part of higher-level search control.)
- ▶ The operation of \mathcal{M} is uniquely characterised by the evolution of *machine state*, *memory state* and *search position* over time.

GLSMs are factored representations of SLS strategies:

- ▶ Given GLSM represents the way in which *initialisation* and *step function* of a hybrid SLS method are composed from respective functions of *subsidiary component SLS methods*.
- ▶ When modelling hybrid SLS methods using GLSMs, *subsidiary SLS methods* should be as simple and pure as possible, leaving *search control* to be represented explicitly at the GLSM level.
- ▶ *Initialisation* is modelled using *GLSM states* (advantage: simplicity and uniformity of model).
- ▶ *Termination of subsidiary search strategies* are often reflected in *conditional transitions* leaving respective GLSM states.

State and Transition Types

In order to completely specify the search method represented by a given GLSM, we need to define:

- ▶ the GLSM model (states, transitions, ...);
- ▶ the search method associated with each *state type*, *i.e.*, step functions for the respective subsidiary SLS methods;
- ▶ the semantics of each *transition type*, *i.e.*, under which conditions respective transitions are executed, and how they effect the memory state.

State types

- ▶ State type semantics are often most conveniently specified procedurally (see algorithm outlines for 'simple SLS methods' from Chapter 2).
- ▶ *initialising state type* = state type τ for which search position after one τ step is independent of search position before step.
initialising state = state of initialising type.
- ▶ *parametric state type* = state type τ whose semantics depends on memory state.
parametric state = state of parametric type.

Transitions types (1)

- ▶ *Unconditional deterministic transitions* – type *DET*:
 - ▶ executed always and independently of memory state or search position;
 - ▶ every GLSM state can have at most one outgoing DET transition;
 - ▶ frequently used for leaving initialising states.
- ▶ *Probabilistic transitions* – type *PROB(p)*:
 - ▶ executed with probability p , independently of memory state or search position;
 - ▶ probabilities of PROB transitions leaving any given state must sum to one.

Note:

- ▶ DET transitions are a special case of PROB transitions.
- ▶ For a GLSM \mathcal{M} any state that can be reached from initial state z_0 by following a chain of $\text{PROB}(p)$ transitions with $p > 0$ will eventually be reached with arbitrarily high probability in any sufficiently long run of \mathcal{M} .
- ▶ In any state z with a $\text{PROB}(p)$ self-transition (z, z) with $p > 0$, the number of GLSM steps before leaving z is distributed geometrically with mean and variance $1/p$.

Transitions types (2)

- ▶ *Conditional probabilistic transitions* – type $\text{CPROB}(C, p)$:
 - ▶ executed with probability proportional to p iff *condition predicate* C is satisfied;
 - ▶ all CPROB transitions from the current GLSM state whose condition predicates are not satisfied are *blocked*, i.e., cannot be executed.

Note:

- ▶ Special cases of $\text{CPROB}(C, p)$ transitions:
 - ▶ $\text{PROB}(p)$ transitions;
 - ▶ *conditional deterministic transitions*, type $\text{CDET}(C)$.
- ▶ Condition predicates should be efficiently computable (ideally: \leq linear time w.r.t. size of given problem instance).

Commonly used simple condition predicates:

\top	always true
$\text{count}(k)$	total number of GLSM steps $\geq k$
$\text{countm}(k)$	total number of GLSM steps modulo $k = 0$
$\text{scount}(k)$	number of GLSM steps in current state $\geq k$
$\text{scountm}(k)$	number of GLSM steps in current state modulo $k = 0$
lmin	current candidate solution is a local minimum w.r.t. the given neighbourhood relation
$\text{evalf}(y)$	current evaluation function value $\leq y$
$\text{noimpr}(k)$	incumbent candidate solution has not been improved within the last k steps

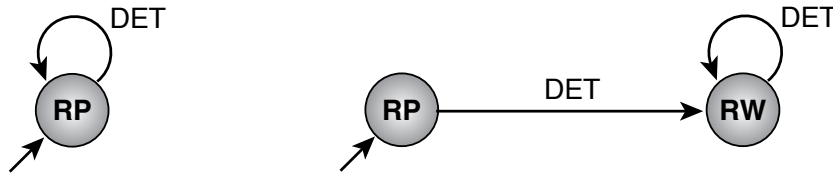
All based on local information; can also be used in negated form.

Transition actions:

- ▶ Associated with individual transitions; provide mechanism for modifying current memory states.
- ▶ Performed whenever GLSM executes respective transition.
- ▶ Modify memory state only, *cannot* modify GLSM state or search position.
- ▶ Have read-only access to search position and can hence be used, e.g., to memorise current candidate solution.
- ▶ Can be added to any of the previously defined transition types.

Modelling SLS Methods Using GLSMs

Uninformed Picking and Uninformed Random Walk



procedure *step-RP*(π, s)

input: *problem instance* $\pi \in \Pi$,
candidate solution $s \in S(\pi)$

output: *candidate solution* $s \in S(\pi)$

$s' := \text{selectRandom}(S);$

return s'

end *step-RP*

procedure *step-RW*(π, s)

input: *problem instance* $\pi \in \Pi$,
candidate solution $s \in S(\pi)$

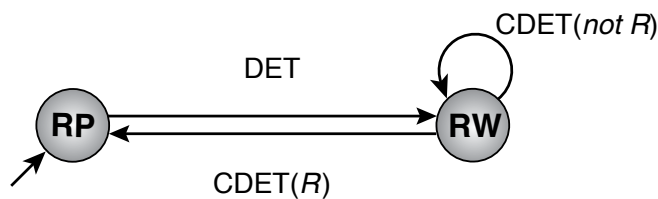
output: *candidate solution* $s \in S(\pi)$

$s' := \text{selectRandom}(N(s));$

return s'

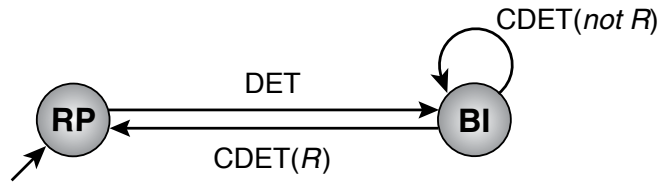
end *step-RW*

Uninformed Random Walk with Random Restart



$R = \text{restart predicate, e.g., countm}(k)$

Iterative Best Improvement with Random Restart



procedure *step-BI*(π, s)

input: *problem instance* $\pi \in \Pi$, *candidate solution* $s \in S(\pi)$

output: *candidate solution* $s \in S(\pi)$

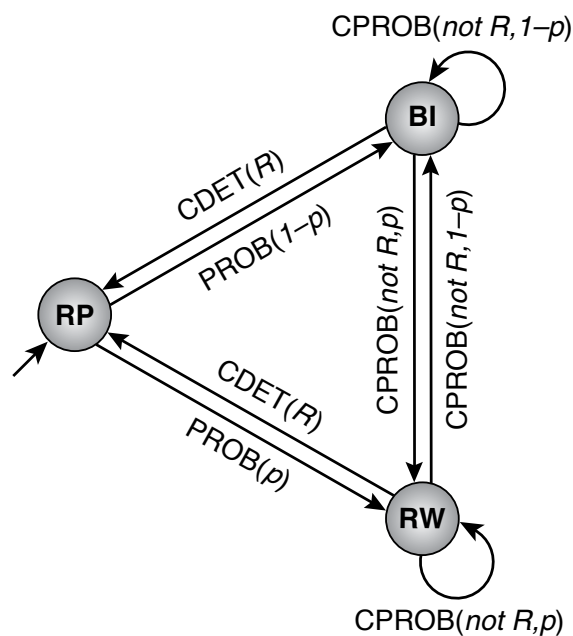
$g^* := \min\{g(s') \mid s' \in N(s)\};$

$s' := \text{selectRandom}(\{s' \in N(s) \mid g(s') = g^*\});$

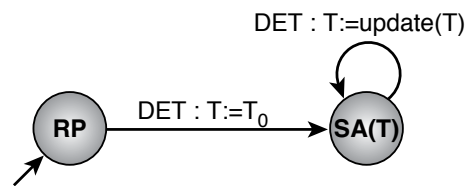
return s'

end *step-BI*

Randomised Iterative Best Improvement with Random Restart

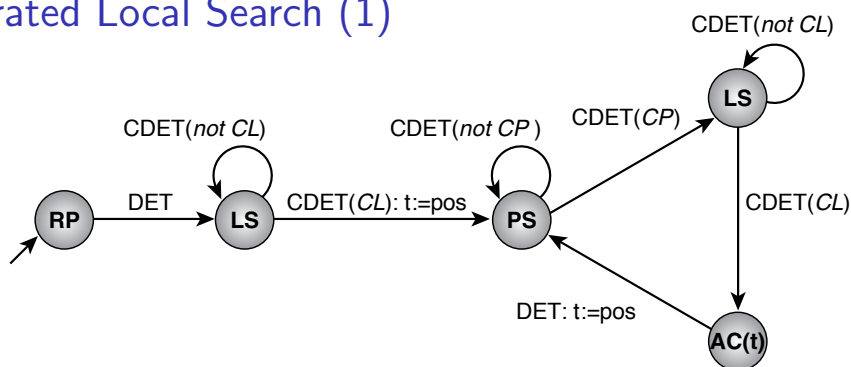


Simulated Annealing



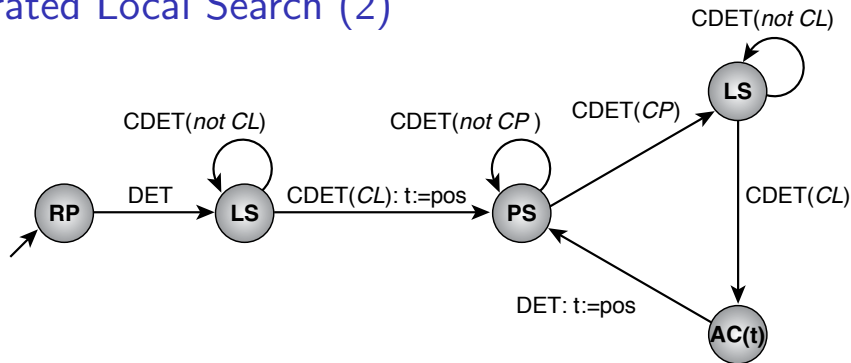
- ▶ Note the use of transition actions and memory for temperature T .
- ▶ The parametric state $SA(T)$ implements probabilistic improvement steps for given temperature T .
- ▶ The initial temperature T_0 and function *update* implement the annealing schedule.

Iterated Local Search (1)



- ▶ The acceptance criterion is modelled as a state type, since it affects the search position.
- ▶ Note the use of transition actions for memorising the current candidate solution (*pos*) at the end of each local search phase.
- ▶ Condition predicates *CP* and *CL* determine the end of perturbation and local search phases, respectively; in many ILS algorithms, $CL := lmin$.

Iterated Local Search (2)



```

procedure step-AC( $\pi, s, t$ )
  input: problem instance  $\pi \in \Pi$ ,
           candidate solution  $s \in S(\pi)$ 
  output: candidate solution  $s \in S(\pi)$ 
  if  $C(\pi, s, t)$  then
    return  $s$ 
  else
    return  $t$ 
  end
end step-AC

```

Ant Colony Optimisation (1)

- General approach for modelling population-based SLS methods, such as ACO, as GLSMs:

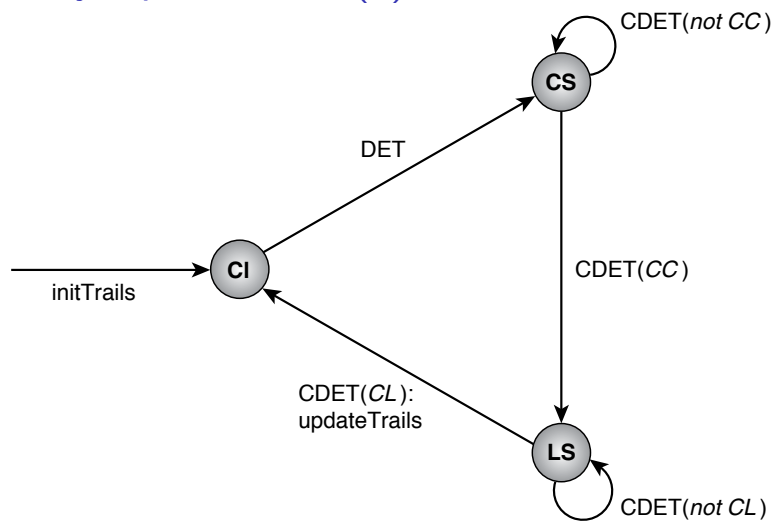
Define search positions as *sets of candidate solutions*; search steps manipulate some or all elements of these sets.

Example: In this view, Iterative Improvement (II) applied to a population sp in each step performs one II step on each candidate solution from sp that is not already a local minimum.

(Alternative approaches exist.)

- Pheromone levels are represented by memory states and are initialised and updated by means of transition actions.

Ant Colony Optimisation (2)



- ▶ The condition predicate CC determines the end of the construction phase.
- ▶ The condition predicate CL determines the end of the local search phase; in many ACO algorithms, $CL := lmin$.