

Homotopical Algebra

Hand in 2

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Notation :

From now on in this document, we will use the following notation. For $c' \xrightarrow{f} c \in \text{Ob } \mathcal{C}/c$, we will note (c', f) and $(c'', h, c') \in \mathcal{C}((c'', g), (c', f))$ for every morphisms from (c'', g) to (c', f)

$$\begin{array}{ccc} c'' & \xrightarrow{h} & c' \\ & \searrow g \quad \swarrow f & \\ & c & \end{array}$$

This may be a confusing notation since it does not contain the information of f or g , but it will be made clear from context.

Exercise 7 :

Part (a) :

Let $d \in \text{Ob } \mathcal{D}$ and let us define the functor $\bar{G} : \mathcal{D}/d \rightarrow \mathcal{C}/G(d)$. We set $\bar{G}(d', f) = (G(d'), G(f))$ on objects of \mathcal{D}/d and $\bar{G}((d'', h, d')) = (G(d''), G(h), G(d'))$ for the a morphism from (d'', g) to (d', f) . Observe that $(G(d''), G(h), G(d'))$ is indeed a morphism since we have that the commutativity of the diagram

$$\begin{array}{ccc} d'' & \xrightarrow{h} & d' \\ & \searrow g \quad \swarrow f & \\ & d & \end{array}$$

implies the commutativity of the diagram

$$\begin{array}{ccc} G(d'') & \xrightarrow{G(h)} & G(d') \\ & \searrow G(g) \quad \swarrow G(f) & \\ & G(d) & \end{array}$$

Now we need to prove that this defines a functor. Let us compute the image of $\text{Id}_{(d', f)}$ for $(d', f) \in \text{Ob } \mathcal{D}/d$. We have

$$\bar{G}(\text{Id}_{(d', f)}) = (G(d'), G(\text{Id}_{d'}), G(d')) = (G(d'), \text{Id}_{G(d')}, G(d')) = \text{Id}_{\bar{G}(d', f)}.$$

Now let us compute the image of a composition. For (d''', i, d'') and (d'', j, d') , we have

$$\begin{aligned} \bar{G}((d'', j, d')(d''', i, d'')) &= \bar{G}(d''', ji, d') \\ &= (G(d'''), G(ji), G(d')) \\ &= (G(d'''), G(j)G(i), G(d')) \\ &= (G(d''), G(j), G(d'))(G(d'''), G(i), G(d'')) \\ &= \bar{G}(d'', j, d')\bar{G}(d''', i, d''). \end{aligned}$$

This allows to say that \bar{G} is a functor.

Part (b) :

As we know that the pair (F, G) is a pair of adjoint functors, we know there exists a natural isomorphism

$$\tau : \mathcal{C}(-, G(-)) \Longrightarrow \mathcal{D}(F(-), -).$$

We set the following functor $\bar{F} : \mathcal{C}/G(d) \longrightarrow \mathcal{D}/d$ the following way. On the objects, we set $\bar{F}(c, f) = (F(c), \tau_{(c,d)}(f))$ and on morphisms $\bar{F}(c', h, c) = (F(c'), F(h), F(c))$ for a morphism from (c, f) to (c', g) . We need now to prove that $(F(c'), F(h), F(c))$ is indeed a morphism from $(F(c'), \tau_{(c',d)}(g))$ to $(F(c), \tau_{(c,d)}(f))$, i.e. we need to prove that the triangle

$$\begin{array}{ccc} F(c') & \xrightarrow{F(h)} & F(c) \\ & \searrow \tau_{(c',d)}(g) & \swarrow \tau_{(c,d)}(f) \\ & d & \end{array}$$

Note first that since (c', h, c) is a morphism from (c, f) to (c', g) , we have $fh = g$. The fact that τ is a natural transformation gives us that

$$\begin{aligned} \mathcal{D}(F(-), -)(h^{op}, \text{Id}_d) \tau_{(c,d)}(f) &= \text{Id}_d \tau_{(c,d)}(f) F(h) = \tau_{(c,d)}(f) F(h) \\ \tau_{(c',d)} \mathcal{C}(-, G(-))(h^{op}, \text{Id}_d)(f) &= \tau_{(c',d)}(G(\text{Id}_d)fh) = \tau_{(c',d)}(fh) = \tau_{(c',d)}(g) \end{aligned}$$

those two ligns are equal and so we conclude that the last triangle commutes.

Now we need to prove that \bar{F} is a functor. First consider $\text{Id}_{(c,f)}$ and let us compute its image. We have

$$\bar{F}(\text{Id}_{(c,f)}) = (F(c), F(\text{Id}_{F(c)}, F(c))) = \text{Id}_{F(c)} = \text{Id}_{\bar{F}(c,f)}.$$

Let us compute the image of a composition. Suppose, we have (c'', i, c') and (c', j, c) . Then

$$\begin{aligned} \bar{F}((c', j, c)(c'', i, c')) &= \bar{F}(c'', ji, c) \\ &= (F(c''), F(ji), F(c)) \\ &= (F(c''), F(j)F(i), F(c)) \\ &= (F(c''), F(j), F(c'))(F(c'))(F(c'), F(i), F(c)) \\ &= \bar{F}(c'', j, c')\bar{F}(c', i, c) \end{aligned}$$

and thus we conclude that \bar{F} is a functor. What is left is to prove that the pair (\bar{F}, \bar{G}) is a pair of adjoint functors. Let $(c, f) \in \text{Ob } \mathcal{C}/G(d)$ and $(d', g) \in \text{Ob } \mathcal{D}$. We would like to find a map from $\mathcal{C}/G(d)((c, f), (G(d'), G(g)))$ to $\mathcal{D}/d((F(c), \tau_{(c,d)}(f)), (d', g))$, that is if we have the commutative triangle

$$\begin{array}{ccc} c & \xrightarrow{h} & G(d') \\ & \searrow f & \swarrow G(g) \\ & G(d) & \end{array}$$

we would like to send h to a map \bar{h} such that the triangle

$$\begin{array}{ccc} F(c) & \xrightarrow{\bar{h}} & d' \\ & \searrow \tau_{(c,d)}(f) & \swarrow g \\ & d & \end{array}$$

commutes. Recall that since τ is a natural isomorphism, the square

$$\begin{array}{ccc} \mathcal{C}(G(d'), G(d')) & \xrightarrow{\tau_{(G(d'), d')}} & \mathcal{D}(FG(d'), d') \\ \bar{\mathcal{C}}(h^{op}, g) \downarrow & & \downarrow \bar{\mathcal{D}}(h^{op}, g) \\ \mathcal{C}(c', G(d)) & \xrightarrow{\tau_{(c, d)}} & \mathcal{D}(F(c'), d) \end{array}$$

commutes, where $\bar{\mathcal{C}} = \mathcal{C}(-, G(-))$ and $\bar{\mathcal{D}} = \mathcal{D}(F(-), -)$. Note that we then have that

$$\begin{aligned} \tau_{(c, d)}(f) &= \tau_{(c, d)}(G(g)h) = \tau_{(c, d)}\bar{\mathcal{C}}(h^{op}, g)(\text{Id}_{G(d')}) \\ &= \bar{\mathcal{D}}(h^{op}, g)\tau_{(G(d'), d')}(\text{Id}_{G(d')}) \\ &= g\tau_{(G(d'), d')}(\text{Id}_{G(d')})F(h) \end{aligned}$$

and so our guess is $\bar{h} = \tau' F(h)$, where $\tau' = \tau_{(G(d'), (d'))}$. Let us denote the family of maps obtained by $\theta_{((c, f), (d', g))}$ and let us denote by θ the natural transformation given by this family.