## SOUTH CHINA UNIVERSITY OF TECHNOLOGY

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# **TEMPLATE**



0 error(s), 0 warning(s)

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## 1 Graph Theory

## 1.1 Shortest Path

## 1.1.1 Dijkstra

```
1 typedef long long LL;
2 const int MAXN = ;
3 const int MAXM = ;
4 const LL DINF = ;
5 typedef pair<LL, int> P;
6 struct Edge {
7
       int to, nxt;
8
       LL w;
9 }e[MAXM];
int head[MAXN], ecnt;
11 LL d[MAXN];
12 priority_queue<P, vector<P>, greater<P> > q;
   inline void addEdge(int x, int y, LL w) {
13
14
        e[++ecnt] = (Edge) \{y, head[x], w\}; head[x] = ecnt;
15
16
   void dijkstra(int st, int n) {
17
        for(int i = 0; i <= n; i++) d[i] = DINF;</pre>
18
        d[st] = 0;
        q.push(make_pair(0, st));
19
20
        while(!q.empty()) {
21
            P x = q.top(); q.pop();
            int u = x.second;
22
            if(d[u] != x.first) continue;
23
24
            for(int i = head[u], v; i; i = e[i].nxt) {
25
                v = e[i].to;
26
                if(d[v] > d[u] + e[i].w) {
27
                    d[v] = d[u] + e[i].w;
28
                    q.push(make_pair(d[v], v));
29
                }
30
           }
31
        }
32
```

## 1.1.2 SPFA

```
1
   struct Edge {
2
       int to, nxt;
3
       LL w;
4
   }e[MAXE];
5
   int head[MAXN], ecnt;
6 LL d[MAXN];
7
   bool exist[MAXN];
8
   queue<int> q;
   inline void addEdge(int x, int y, LL w) {
9
10
      e[++ecnt] = (Edge) {y, head[x], w}; head[x] = ecnt;
   }
11
12
   void SPFA(int st) {
13
       memset(d,0x3f,sizeof(d));
14
       d[st] = 0;
15
        q.push(st);
16
        exist[st] = 1;
17
       while(!q.empty()) {
```

```
18
            int u = q.front(); q.pop();
19
             exist[u] = 0;
20
            for(int i = head[u], v; i; i = e[i].nxt) {
21
                 v = e[i].to;
22
                 if(d[v] > d[u] + e[i].w) {
23
                     d[v] = d[u] + e[i].w;
                     //pre[v] = u;
24
                     if(!exist[v]) {
25
26
                          q.push(v);
27
                          exist[v] = 1;
                     }
28
29
                 }
30
            }
31
        }
32
   }
```

#### 1.1.3 Johnson

```
1
   void johnson() {
2
       //全源带负权最短路,新建超级节点向全部点连权为0的边, 计算D, 利用Dij计算点对距离
3
       for(int i = 1; i <= n; i++) addEdge(0, i, 0);</pre>
4
       spfa(0);
5
       for(int u = 1; u <= n; u++)</pre>
6
           for(int i = head[u]; i; i = e[i].nxt)
7
               e[i].w += D[e[i].from] - D[e[i].to];
8
       dijkstra(s,n);
9
       //ans = d[n] - D[s] + D[n];
10
11
   }
```

## 1.1.4 K Shortest Path (A\*)

```
//可重复走同一条边 利用了反向边表示,ecnt初始化为1
   //调用 dijkstra (ed, n) 跑反向图,注意 if (i lo 1) {松弛}
3
   int shrt[MAXN];
   LL A_star(int st, int ed, int k, int n) {
4
        if(d[st] == d[0]) return -1;
5
6
       if(st == ed) k++;
7
        q.push(make_pair(d[st], st));
8
        while(!q.empty()) {
9
            P x = q.top(); q.pop();
10
            int u = x.second;
            LL xd = x.first;
11
12
            ++shrt[u];
13
            if(u == ed) {
14
                if(shrt[ed] == k) return xd;
15
           for(int i = head[u], v; i; i = e[i].nxt)
16
17
           if(!(i & 1)) {
18
                v = e[i].to;
19
                if(shrt[v] < k)q.push(make_pair(xd-d[u]+e[i].w+d[v], v));</pre>
20
            }
21
        }
22
        return -1;
23
```

## 1.1.5 K Shortest Path (Protractable Heap)

```
//可重复走同一条边
1
   typedef double LD;
3 const int MAXN = ;
   const int MAXM = ;
   const int MAXLT = MAXM * 20;
5
   const LD DINF = ;
6
7
   const LD eps = ;
8
9
   namespace LT{
10
        int tcnt;
11
        int ls[MAXLT], rs[MAXLT], dis[MAXLT];
12
        int to[MAXLT];
13
        LD val[MAXLT];
14
15
        inline int newnode(LD w, int _to) {
16
            ++tcnt;
17
            ls[tcnt] = rs[tcnt] = 0; dis[tcnt] = 1;
18
            val[tcnt] = w; to[tcnt] = _to;
19
            return tcnt;
20
21
        inline int copynode(int id) {
22
            ++tcnt;
23
            ls[tcnt] = ls[id]; rs[tcnt] = rs[id]; dis[tcnt] = dis[id];
24
            val[tcnt] = val[id]; to[tcnt] = to[id];
25
            return tcnt;
26
        void push_up(int x) {
27
28
            if(dis[ls[x]] < dis[rs[x]]) swap(ls[x], rs[x]);</pre>
29
            dis[x] = dis[rs[x]] + 1;
30
31
        int merge(int x, int y) {
32
            if(!x || !y) return x^y;
33
            if(val[x] - val[y] > eps) swap(x, y);
34
            int p = copynode(x);
            rs[p] = merge(rs[p],y);
35
36
            push_up(p);
37
            return p;
38
       }
39
   }
40 int rt[MAXN];
41
42 typedef pair<LD, int> P;
43
   struct Edge {
44
       int to, nxt;
45
        LD w;
   }e[MAXM];
46
47
   int head[MAXN], ecnt;
48 int stan, sta[MAXN], fa[MAXN];
49 int vis[MAXN], cov[MAXM];
50 LD d[MAXN];
51 priority_queue<P, vector<P>, greater<P> > q;
52 inline void addEdge(int x, int y, LD w) {
53
        e[++ecnt] = (Edge) {y, head[x], w}; head[x] = ecnt;
54
   }
55
   void init(int n, int m) {
56
        ecnt = 1; stan = 0;
57
        for(int i = 1; i <= n; i++)</pre>
58
            head[i] = cov[i] = fa[i] = 0;
59
       for(int i = 1; i <= m; i++) {
```

```
60
             int u, v; LD w;
61
             scanf("%d%d%lf", &u,&v,&w);
62
             addEdge(u, v, w);
63
             addEdge(v, u, w);
64
         }
65
66
    void dijkstra(int st, int n) {
         for(int i = 0; i <= n; i++) {d[i] = DINF; vis[i] = 0;}</pre>
67
68
         while(!q.empty()) q.pop();
69
         d[st] = 0;q.push(make_pair(0, st));
70
         while(!q.empty()) {
71
             P x = q.top(); q.pop();
72
             int u = x.second;
73
             if(vis[u]) continue;
74
             vis[u] = 1;
             for(int i = head[u], v; i; i = e[i].nxt)
75
76
                 if(i & 1) {
                     v = e[i].to;
77
                     if(d[v]-(d[u] + e[i].w) > eps) {
78
79
                         d[v] = d[u] + e[i].w;
80
                         q.push(make_pair(d[v], v));
                     }
81
82
                 }
83
        }
84
    }
85
    void buildT(int u) {
86
         sta[++stan] = u;
87
         vis[u] = 1;
88
         for(int i = head[u], v; i; i = e[i].nxt)
89
         if(i & 1) {
90
             v = e[i].to;
             if(fabs(d[v] - (d[u] + e[i].w)) < eps && !vis[v]) {</pre>
91
                 fa[v] = u; cov[i^1] = 1; buildT(v);
92
93
         }
94
95
    }
96
    void buildH(int st, int n) {
         buildT(st);
97
         for(int i = 2, u, v; i <= ecnt; i += 2) {
98
99
             if(!cov[i]) {
100
                 u = e[i^1].to; v = e[i].to;
101
                 if(fabs(d[u] - d[0]) < eps || fabs(d[v] - d[0]) < eps) continue;
102
                 rt[u] = LT::merge(rt[u], LT::newnode(d[v]+e[i].w-d[u], v));
103
             }
104
105
         for(int i = 2, u; i <= n; i++)
106
             if(fa[u = sta[i]]) rt[u] = LT::merge(rt[u], rt[fa[u]]);
107
    //求前k短路径,其和不超过W,问最大k int ans = 0;
108
    //求第 k短路径,LD ans = 0;
109
110
    void getKth(int st, int ed, LD W) {
111
         while(!q.empty()) q.pop();
         //最短路要记入答案,注意d[st]-Weps与W-d[st]<eps
112
         //! if(d/st/ - W > eps) return;
113
         //! else\{W -= d[st]; ++ans;\}
114
         //注意st和ed相同
115
         //\#if(st == ed) k++;
116
         //\#if(d[st] == d[0]) \{ans = -1; return;\}
117
118
         //\#if(--k == 0) \{ans = d | st |; return; \}
119
120
        int u = rt[st], v;
```

```
if(u) q.push(make_pair(LT::val[u], u));
121
122
         while (!q.empty()) {
123
             u = q.top().second; LD cur = q.top().first;
124
             q.pop();
125
             //! if (cur + d[st] - W > eps) break;
126
             //! else \{W = cur + d[st]; ++ans;\}
127
             //\#if (--k == 0) \{ans = cur + d[st]; break;\}
128
             v = LT::ls[u];
129
             if (v) q.push(make_pair(cur - LT::val[u] + LT::val[v], v));
130
             v = LT::rs[u];
             if (v) q.push(make_pair(cur - LT::val[u] + LT::val[v], v));
131
             v = rt[LT::to[u]];
132
             if (v) q.push(make_pair(cur + LT::val[v], v));
133
134
135
    }
136
    void sol() {
137
        int n, m, st = , ed = ;
         //LD W; int k;
138
139
         init(n, m);
140
         dijkstra(ed, n);
141
         for(int i = 0; i <= n; i++) vis[i] = 0;
         buildH(ed, n);
142
143
         getKth(st, ed, W);
         printf("%d\n", ans);
144
145
    }
```

## 1.2 Network Flow

## 1.2.1 ISAP

```
namespace NWF {
1
2
        struct Edge{
3
            int to, nxt;LL f;
4
        }e[MAXM << 1];
5
        int S, T, tot;
        int ecnt, head[MAXN], cur[MAXN], pre[MAXN], num[MAXN], dis[MAXN];
6
7
        queue<int> q;
8
        void init(int _S, int _T, int _tot){
9
            ecnt = 1; S = _S; T = _T; tot = _tot;
10
            memset(num, 0, (tot + 1) * sizeof(int));
            memset(head, 0, (tot + 1) * sizeof(int));
11
12
        inline void addEdge(int u, int v, LL f) {
13
14
            e[++ecnt] = (Edge) {v, head[u], f}; head[u] = ecnt;
15
            e[++ecnt] = (Edge) {u, head[v], 0}; head[v] = ecnt;
16
        void bfs() {
17
            memset(dis, 0, (tot + 1) * sizeof(int));
18
19
            q.push(T);
20
            dis[T] = 1;
21
            while(!q.empty()) {
22
                int u = q.front(), v; q.pop();
23
                num[dis[u]]++;
                for(int i = cur[u] = head[u]; i; i = e[i].nxt) {
24
25
                    if(!dis[v = e[i].to]) {
                        dis[v] = dis[u] + 1;
26
27
                        q.push(v);
28
                    }
29
```

```
30
            }
31
32
        LL augment() {
33
            LL flow = INF;
            for(int i = S; i != T; i = e[cur[i]].to)
34
35
                 flow = min(flow, e[cur[i]].f);
36
            for(int i = S; i != T; i = e[cur[i]].to) {
                 e[cur[i]].f -= flow;
37
38
                 e[cur[i] ^ 1].f += flow;
39
40
            return flow;
41
        LL isap() {
42
43
            bfs();
            int u = S, v;
44
45
            LL flow = 0;
            while(dis[S] <= tot) {</pre>
46
47
                 if(u == T) {
48
                     flow += augment();
49
                     u = S;
50
                 bool fg = 0;
51
52
                 for(int i = cur[u]; i; i = e[i].nxt) {
53
                     if(e[i].f && dis[u] > dis[v = e[i].to]) {
54
                         pre[v] = u;
55
                         cur[u] = i;
56
                         u = v;
                         fg = 1;
57
58
                         break;
59
                     }
60
                 if(fg) continue;
61
62
                 if(!--num[dis[u]]) break;
63
                 int maxDis = tot;
                 for(int i = head[u]; i; i = e[i].nxt) {
64
                     if(e[i].f \&\& maxDis > dis[v = e[i].to]) {
65
66
                         maxDis = dis[v];
                         cur[u] = i;
67
68
                     }
69
                 }
70
                 num[dis[u] = maxDis + 1]++;
71
                 if(u != S) u = pre[u];
72
            }
73
            return flow;
74
75
```

#### 1.2.2 HLPP

```
1
   namespace NWF{
2
        struct Edge{
3
            int to,nxt;LL f;
4
        }e[MAXM << 1];
5
        int S, T, tot;
        int ecnt, head[MAXN], dis[MAXN], num[MAXN];
6
7
        LL sumf[MAXN];
8
        queue<int> q;
9
       list<int> dep[MAXN];
10
       void init(int _S,int _T,int _tot){
11
            ecnt = 1;S = _S;T = _T;tot = _tot;
```

```
memset(num, 0, (tot + 1) * sizeof(int));
12
            memset(head, 0, (tot + 1) * sizeof(int));
13
            memset(sumf, 0, (tot + 1) * sizeof(LL));
14
15
16
        void addEdge(int u,int v,LL f){
            e[++ecnt] = (Edge) {v, head[u], f};head[u] = ecnt;
17
18
            e[++ecnt] = (Edge) {u, head[v], 0}; head[v] = ecnt;
19
20
        void bfs(){
            memset(dis, 0, (tot + 1) * sizeof(int));
21
22
            q.push(T); dis[T] = 1;
23
            while(!q.empty()){
24
                 int u=q.front(), v; q.pop();
25
                 for(int i = head[u]; i; i = e[i].nxt)
                 if(!dis[v = e[i].to]){
26
                     dis[v] = dis[u] + 1;
27
28
                     q.push(v);
29
                 }
30
            }
31
        LL hlpp(){
32
33
            bfs();
34
            dis[S] = tot + 1;
            for(int i = 1;i <= tot; ++i)num[dis[i]]++;</pre>
35
36
            for(int i = tot + 1; ~i; --i)dep[i].clear();
37
            int maxd = dis[S];LL f;
38
            dep[maxd].push_back(S);sumf[S] = INF;
39
            for(;;){
40
                 while(maxd && dep[maxd].empty())maxd--;
41
                 if(!maxd)break;
42
                 int u = dep[maxd].back(), v;dep[maxd].pop_back();
                 int minDis = tot + 1;
43
                 for(int i = head[u]; i;i = e[i].nxt)
44
                 if(e[i].f){
45
                     if(dis[u] > dis[v = e[i].to]){
46
                         f = min(sumf[u], e[i].f);
47
48
                         e[i].f -= f;e[i^1].f += f;
                         if(sumf[u] != INF) sumf[u] -= f;
49
50
                         if(sumf[v] != INF) sumf[v] += f;
51
                         if(v!=S \&\& v!=T \&\& sumf[v] == f){
52
                             maxd = max(maxd, dis[v]);
53
                             dep[dis[v]].push_back(v);
54
55
                         if(!sumf[u])break;
56
                     }else minDis=min(minDis, dis[v] + 1);
57
                 if(sumf[u]){
58
59
                     if(!--num[dis[u]]){
                         for(int i = dis[u];i <= maxd;++i){</pre>
60
61
                             while(!dep[i].empty()){
62
                                  --num[i];
63
                                  dis[dep[i].back()] = tot + 1;
64
                                  dep[i].pop_back();
65
                             }
66
                         }
67
                         maxd = dis[u] - 1; dis[u] = tot + 1;
68
                     }else{
69
                         dis[u] = minDis;
70
                         if(minDis > tot)continue;
71
                         num[minDis]++;
72
                         maxd = max(maxd, minDis);
```

#### 1.2.3 Dinic

注意当流为浮点数的时候,要判断 eps 以及不能使用 sumf-=tmpf, 否者 1e18 将不会发生改变

```
namespace NWF {
1
2
        struct Edge {
3
            int to, nxt;LL f;
4
        } e[MAXM << 1];
5
        int S, T, tot;
6
        int ecnt, head[MAXN], cur[MAXN], dis[MAXN];
        queue<int> q;
7
8
        void init(int _S, int _T, int _tot){
9
            ecnt = 1; S = _S; T = _T; tot = _tot;
10
            memset(head, 0, (tot + 1) * sizeof(int));
11
12
        void addEdge(int u, int v, LL f) {
            e[++ecnt] = (Edge) \{v, head[u], f\}; head[u] = ecnt;
13
14
            e[++ecnt] = (Edge) {u, head[v], 0}; head[v] = ecnt;
15
16
        bool bfs() {
            memset(dis, 0, (tot + 1) * sizeof(int));
17
            q.push(S); dis[S] = 1;
18
19
            while (!q.empty()) {
20
                int u = q.front(), v; q.pop();
                for (int i = cur[u] = head[u]; i ; i = e[i].nxt) {
21
22
                    if (e[i].f && !dis[v = e[i].to]) {
23
                         q.push(v);
24
                         dis[v] = dis[u] + 1;
25
                    }
26
                }
27
            }
28
            return dis[T];
29
30
        LL dfs(int u, LL maxf) {
31
            if (u == T) return maxf;
            LL sumf = maxf;
32
            for (int &i = cur[u]; i; i = e[i].nxt) {
33
                if (e[i].f && dis[e[i].to] > dis[u]) {
34
35
                    LL tmpf = dfs(e[i].to, min(sumf, e[i].f));
                    e[i].f -= tmpf; e[i ^ 1].f += tmpf;
36
37
                    sumf -= tmpf;
38
                    if (!sumf) return maxf;
39
                }
40
            }
41
            return maxf - sumf;
42
        LL dinic() {
43
            LL ret = 0;
44
45
            while (bfs()) ret += dfs(S, INF);
46
            return ret;
47
        }
48
        void rebuild(){
49
            //无向图采用e[i].f = e[i^1].f的方式建立图
```

```
for(int i = 2; i <= ecnt; i+=2) e[i].f = e[i^1].f = (e[i].f + e[i^1].f) >> 1;

//有向图

//for(int i = 2; i <= ecnt; i+=2) e[i].f += e[i^1].f, e[i^1].f = 0;

}

}
```

## 1.2.4 Bound Flow

```
1
   namespace NWF{
2
        //在Edge中添加下限, delta_flow[i]:节点i的入流-出流, bound_flow注意要处理除了SS和TT以
        外的点
        int delta_flow[MAXN];
3
        void addEdge(int u, int v, int mxf, int mnf) {
4
5
            addEdge(u, v, mxf-mnf);
6
            delta_flow[u] -= mnf; delta_flow[v] += mnf;
7
        }
8
        void delEdge(int u) {
9
            for(int i = head[u]; i; i = e[i].nxt) e[i].f = e[i^1].f = 0;
10
        }
        int bound_flow() {
11
12
            int SS = ++tot, TT = ++tot, sum = 0;
13
            head[SS] = head[TT] = 0;
14
            for(int i =1; i <= tot;++i) {</pre>
15
                if(delta_flow[i]<0) addEdge(i, TT, -delta_flow[i]);</pre>
16
                if(delta_flow[i]>0){
17
                    sum+=delta_flow[i];
                    addEdge(SS, i, delta_flow[i]);
18
19
                }
20
            }
21
            addEdge(T,S, INF);
            int tS = S, tT = T;
22
            S = SS; T = TT;
23
24
            if (dinic()==sum) {
                delEdge(SS); delEdge(TT);
25
26
                int flow = e[ecnt].f;
                e[ecnt].f = e[ecnt^1].f = 0;
27
                //S = tS; T = tT; // 有上下界有源汇最大流
28
29
                //return flow + dinic();
                S = tT; T = tS;// 有上下界有源汇最小流
30
31
                return flow - dinic();
32
33
                return -1;
34
           }
35
        }
36
   }
```

## 1.2.5 Modeling Optimization

利用分治优化建模,每个点i向j连边,费用为 |ai-aj|

```
int pos[MAXN];
pair<int, int> tmp[MAXN];

void CDQ(int L, int R) {
   if (L == R) return;
   int mid = (L + R) >> 1;
   CDQ(L, mid); CDQ(mid + 1, R);
   inplace_merge(tmp + L, tmp + mid + 1, tmp + R + 1);
   for (int i = L; i <= R; ++i) pos[tmp[i].second] = i;
   for (int i = 2; i <= R - L + 1; ++i) {</pre>
```

```
10
            addEdge(tot + i, tot + i - 1, INF, tmp[L + i - 1].first - tmp[L + i - 2].first);
11
            addEdge(tot + i - 1, tot + i, INF, tmp[L + i - 1].first - tmp[L + i - 2].first);
12
13
        for (int i = L; i <= R; ++i) {
            if (i <= mid)</pre>
14
15
                addEdge(i+i-1, tot + pos[i] - L + 1, 1, 0);
16
                addEdge(tot + pos[i] - L + 1, i+i, 1, 0);
17
18
19
        tot += R - L + 1;
20
```

#### 1.2.6 Gomory-Hu Tree

两点间的割可以转为树上两点的距离

```
1
   namespace NWF{
2
        Edge Te[MAXN];
3
        int Tcnt, Thead[MAXN];
4
        void TaddEdge(int u, int v, LL f) {
5
            Te[++Tcnt] = (Edge) {v, Thead[u], f}; Thead[u] = Tcnt;
6
            Te[++Tcnt] = (Edge) {u, Thead[v], f}; Thead[v] = Tcnt;
7
8
        int node[MAXN], tmp[MAXN];
9
        void build(int 1, int r) {
10
            if (1 == r) return;
11
            S = node[1]; T = node[1+1];
            rebuild();
12
13
            LL cut = dinic();
14
            TaddEdge(S, T, cut);
15
            int tl = 1, tr = r;
16
            for(int i = 1; i <= r; i++) {</pre>
17
                 if(dis[node[i]]) tmp[tl++] = node[i]; else tmp[tr--] = node[i];
18
19
            for(int i=1; i<=r; i++) node[i] = tmp[i];</pre>
20
            build(l,tl-1); build(tr+1,r);
21
        int log2n;
22
        int dep[MAXN], anc[MAXN][MAXS];LL mnl[MAXN][MAXS];//anc: 祖先; mnl: 最小边
23
        void lca_dfs(int u, int _fa) {
24
            for(int i=Thead[u], v; i; i=Te[i].nxt) {
25
26
                 if((v = Te[i].to) == _fa) continue;
27
                dep[v] = dep[u] + 1;
28
                anc[v][0] = u;
29
                mnl[v][0] = Te[i].f;
30
                lca_dfs(v, u);
31
            }
32
33
        void work() {
34
            if(tot == 0) return;
35
            log2n = log2(tot)+1; Tcnt = 1;
36
            for(int i=1; i<=tot; i++) node[i]=i, Thead[i]=0;</pre>
37
            build(1, tot);
            dep[1] = 1; anc[1][0] = 0; mnl[1][0] = INF;
38
39
            lca_dfs(1, -1);
            for(int j = 1; j <= log2n; j++) {</pre>
40
                for(int i = 1; i <= tot; i++) {</pre>
41
                     anc[i][j] = anc[anc[i][j-1]][j-1];
42
43
                     mnl[i][j] = min(mnl[i][j-1], mnl[anc[i][j-1]][j-1]);
44
                }
45
```

```
46
47
        LL get_cut(int u,int v) {
48
            LL res=INF;
49
            if(dep[u] < dep[v]) swap(u, v);</pre>
50
             for(int i = log2n; i >= 0; i--) {
51
                 if(dep[anc[u][i]] >= dep[v]){
52
                     res = min(res, mnl[u][i]);
53
                     u = anc[u][i];
54
                 }
55
            if(u == v) return res;
56
             for(int i = log2n; i>=0 ; i--) {
57
                 if(anc[u][i] != anc[v][i]) {
58
59
                     res = min(res, mnl[u][i]);
60
                     res = min(res, mnl[v][i]);
61
                     u = anc[u][i];
62
                     v = anc[v][i];
63
                 }
64
             }
65
            res = min(res, mnl[u][0]);
66
            res = min(res, mnl[v][0]);
67
            return res;
68
        }
69
   }
```

#### 1.2.7 MCMF

```
1
    namespace NWF{
2
        struct Edge {
3
            int to, nxt;LL f, c;
4
        } e[MAXM << 1];
5
        int S, T, tot;
6
        int ecnt, head[MAXN], cur[MAXN];LL dis[MAXN];
7
        bool exist[MAXN];
8
        queue<int> q;
9
        void init(int _S, int _T, int _tot){
            ecnt = 1; S = _S; T = _T; tot = _tot;
10
            memset(head, 0, (tot + 1) * sizeof(int));
11
12
13
        void addEdge(int u, int v, LL f, LL c) {
14
            e[++ecnt] = (Edge) {v, head[u], f, c}; head[u] = ecnt;
15
            e[++ecnt] = (Edge) \{u, head[v], 0,-c\}; head[v] = ecnt;
16
        bool spfa() {
17
18
            for(int i = 0;i <= tot; ++i){</pre>
19
                dis[i] = INF; cur[i] = exist[i] = 0;
20
21
            q.push(S);dis[S] = 0;exist[S] = 1;
22
            while(!q.empty()) {
23
                 int u = q.front(), v; q.pop();exist[u] = 0;
24
                for(int i = head[u]; i; i = e[i].nxt) {
25
                     if(e[i].f \&\& dis[v = e[i].to] > dis[u] + e[i].c) {
                         dis[v] = dis[u] + e[i].c;
26
27
                         cur[v] = i;
28
                         if(!exist[v]) {
29
                             q.push(v);
30
                             exist[v] = 1;
31
                         }
32
                     }
33
```

```
34
35
            return dis[T] != INF;
36
        LL mcmf() {
37
38
            LL cost = 0;
            //while(spfa() & dis[T] < 0) {//最小费用可行流
39
40
            while(spfa()) {
41
                LL flow = INF;
                for(int i = T; i != S; i = e[cur[i] ^ 1].to)
42
43
                     flow = min(flow, e[cur[i]].f);
                for(int i = T; i != S; i = e[cur[i] ^ 1].to) {
44
                    e[cur[i]].f -= flow;
45
                    e[cur[i] ^ 1].f += flow;
46
47
                }
48
                cost += flow * dis[T];
49
            }
50
            return cost;
51
        }
52
   }
```

## 1.3 Tree Related

## 1.3.1 Union Set

```
int fa[MAXN], rnk[MAXN];
   int Find(int x) { return x == fa[x] ? x : fa[x] = Find(fa[x]); }
   bool same(int x, int y){ return Find(x) == Find(y); }
   void unite(int x, int y)
4
5
   {
       x = Find(x);
6
7
       y = Find(y);
        if(x == y) return;
8
9
        if(rnk[x] < rnk[y]) {</pre>
10
            fa[x] = y;
11
        }
12
        else {
13
            fa[y] = x;
14
            if(rnk[x] == rnk[y]) rnk[x]++;
15
       }
16
   }
```

## 1.3.2 Kruskal

```
namespace MST{
1
2
        struct Edge{
3
            int u,v; LL w;
4
            bool operator < (const Edge& x) const { return w < x.w; }</pre>
5
        }e[MAXM];
6
        int ecnt, fa[MAXN];
7
        void addEdge(int u, int v, LL w) {
8
            e[++ecnt] = (Edge)\{v, u, w\}; //headp[u] = ecnt;
9
        int Find(int x) { return x == fa[x] ? x : fa[x] = Find(fa[x]); }
10
        LL kruskal(int n) {
11
            sort(e + 1, e + ecnt + 1);
12
13
            for(int i = 1; i <= n; i++) fa[i] = i;</pre>
14
            LL sum = 0;
15
            for (int i = 1; i <= ecnt; i++){</pre>
```

```
16
                 int fu = Find(e[i].u), fv = Find(e[i].v);
17
                 if(fu != fv){
18
                     fa[fu] = fv;
19
                      sum += e[i].w;
20
21
             }
22
             return sum;
23
        }
24
```

## 1.3.3 Prim

```
1
    namespace MST {
2
        struct Edge{
3
            int to,nxt; LL w;
4
        }e[MAXM];
5
        int ecnt, head[MAXN], vis[MAXN]; // pre [MAXN];
6
        LL dis[MAXN];
7
        void addEdge(int u, int v, LL w){
            e[++ecnt] = (Edge){v, head[u], w}; head[u] = ecnt;
8
9
            e[++ecnt] = (Edge){u, head[v], w}; head[v] = ecnt;
10
        LL Prim(int n){
11
12
            for (int i = 1; i <= n; i++){
                 //pre[i] = 0;
13
14
                 vis[i] = 0;
                 dis[i] = INF;
15
16
            }
17
            vis[1] = 1;
18
            LL sum = 0;
            for (int i = head[1]; i; i = e[i].nxt)
19
20
                 dis[e[i].to] = min(dis[e[i].to],e[i].w);
21
            for (int j = 1; j < n; j++){
22
                 int u; LL minDis = INF;
23
                 for (int i = 1; i <= n; ++i)</pre>
24
                     if (!vis[i] && dis[i] < minDis){</pre>
25
                         minDis = dis[i];
26
                         u = i;
27
                     }
28
                 if (minDis == INF) return -1;
29
                 vis[u] = 1;
30
                 sum += minDis;
                 for (int i = head[u], v; i; i = e[i].nxt)
31
                 if (!vis[v = e[i].to] && e[i].w < dis[v]){</pre>
32
33
                      //pre[u] = v;
34
                     dis[v] = e[i].w;
35
                 }
36
            }
37
            return sum;
38
        }
39
```

## 1.3.4 Spanning Tree Calculation

关联矩阵 B:n\*m 的矩阵, 其中 ek=(vi,vj), Bik 和 Bjk 一个为 1 一个为 1, 第 k 列其他元素为 0 度数矩阵 D:n\*n 的矩阵, 其中 i!=j 时, D[i][j]=0;i==j 时, D[i][j]=vi 的度邻接矩阵 A:n\*n 的矩阵, vi,vj 有边相连, 为 1, 否则为 0 Kirchhoff 矩阵: B\*BT = D - A 即: 如果 i=j, 那么 aij 为点 i(j) 的度数。如果 i!=j, 那么 Aij 为 i 到 j 的边数的相反数。生成树个数: Kirchhoff 矩阵 n-1 阶主子式的行列式值构造 Kirchhoff 矩阵, 调用 det(n)

```
1
    LL a[MAXN][MAXN];
2
    void getSTC(int n, int m) {
3
        for(int i = 1; i <= n; i++) {</pre>
4
             for(int j = 1; j <= n; j++)</pre>
5
                 a[i][j] = a[j][i] = 0;
6
7
        for(int i = 1,u, v; i <= m; i++) {</pre>
8
             scanf("%d%d", &u,&v);
9
             if(u == v) continue;
10
             a[u][v] = --a[v][u];
11
        for(int i = 1; i <= n; i++) {</pre>
12
            int t = 0;
13
14
             for(int j = 1; j <= n; j++)</pre>
15
                 t += a[i][j];
16
            a[i][i] = -t;
17
        LL ans = det(); // 删掉一行一列以后求行列式的值
18
19
```

## 1.3.5 Minimum Spanning Tree Calculation

```
1
   typedef long long LL;
2 const int MAXN = ;
   const int MAXM = ;
3
4
   int sum,ans1,ans2=1, Mod=;
5
   int fa1[MAXN],fa2[MAXN];
   bool vis[MAXN];
7
    struct Edge{int v,u,val;}e[MAXM];
   bool cmp(Edge A, Edge B){return A.val<B.val;}</pre>
9
   int getfa1(int *fa,int x){return fa[x]=fa[x]==x?x:getfa1(fa,fa[x]);}
   int getfa2(int *fa,int x){return fa[x]==x?x:getfa2(fa,fa[x]);}
10
    void dfs(int tot,int l,int r){
11
12
        if(tot==0){++sum; return;}
        for(int i=1,fx,fy;i<r;++i)</pre>
13
14
            if(!vis[i]){
                 vis[i]=true;
15
16
                 fx=getfa2(fa2,e[i].u);fy=getfa2(fa2,e[i].v);
17
                 if(fx!=fy){
                     fa2[fx]=fy;
18
19
                     dfs(tot-1,i+1,r);
20
                     fa2[fx]=fx;
21
                 vis[i]=false;
22
23
            }
24
   }
   void sol(){
25
        int n,m;
26
        scanf("%d %d",&n,&m);
27
28
        for(int i=1;i<=n;++i)fa1[i]=fa2[i]=i;</pre>
29
        for(int i=1;i<=m;++i)</pre>
30
            scanf("%d %d %d",&e[i].v,&e[i].u,&e[i].val);
31
        std::sort(e+1,e+m+1,cmp);e[m+1].val=-1;
32
        for(int i=1,j=1,fx,fy,tot;i<=m;++i){</pre>
33
            for(;e[i].val==e[j].val;++j);
34
            tot=0;
            for(int k=i;k<j;++k){</pre>
35
36
                 fx=getfa1(fa1,e[k].u);fy=getfa1(fa1,e[k].v);
37
                 if(fx!=fy){++tot;ans1++;fa1[fx]=fy;}
38
```

```
39
             if(!tot)continue;
40
             sum=0;dfs(tot,i,j);
41
             (ans2*=sum)%=Mod;
42
             for(int k=i;k<j;++k){</pre>
43
                 fx=getfa1(fa2,e[k].u);fy=getfa1(fa2,e[k].v);
44
                 if(fx!=fy)fa2[fx]=fy;
45
             }
46
        if(ans1!=n-1)puts("0");else printf("%d",ans2);
47
48
```

#### 1.3.6 Steiner Tree

```
const int MAXH = 128;
1
   const int MAXW = 128;
    const int MAXST = 1256;
4
    namespace SteinerTree{
        const int dx[4] = \{0,0,1,-1\};
5
6
        const int dy[4] = \{1,-1,0,0\};
7
        int n, m, k, stn;
8
        bool vis[MAXH][MAXW];
9
        LL G[MAXH][MAXW],f[MAXST][MAXH][MAXW];
10
        queue<pair<int, int> > q;
11
        struct PRE{int s, x, y;} pre[MAXST][MAXH][MAXW];
12
13
        void spfa(int st) {
14
            while(!q.empty()) {
15
                 int vx = q.front().first,vy = q.front().second;
16
                 q.pop();
17
                 vis[vx][vy] = 0;
18
                 for (int i = 0, ux,uy; i < 4; ++i) {
19
                     ux = vx + dx[i];
                     uy = vy + dy[i];
20
                     if (ux==0||uy==0||ux==n+1||uy==m+1)continue;
21
22
                     if (f[st][vx][vy]+G[ux][uy]<f[st][ux][uy]) {</pre>
23
                          f[st][ux][uy] = f[st][vx][vy] + G[ux][uy];
                          pre[st][ux][uy] = (PRE) {st, vx, vy};
24
25
                          if (!vis[ux][uy]) {
26
                              vis[ux][uy]=1;
27
                              q.push(make_pair(ux,uy));
28
                          }
29
                     }
                 }
30
            }
31
32
33
        LL sum = 0;
        void init() {
34
35
            k = 0;
36
             for(int i = 1; i <= n; ++i)</pre>
37
                 for(int j = 1; j <= m; ++j) {</pre>
                     scanf("%d", &G[i][j]);
38
39
                     if(G[i][j] == 0) k++;
40
             stn = 1 << k;
41
             for (int st = 0; st < stn; ++st)</pre>
42
                 for (int i = 1; i <= n; ++i)</pre>
43
                     for (int j = 1; j <= m; ++j)
44
45
                          f[st][i][j] = INF;
46
             int tk = 0;
47
             for(int i = 1; i <= n; ++i)</pre>
```

```
for(int j = 1; j <= m; ++j) {</pre>
48
49
                     vis[i][j] = 0;
50
                     if(G[i][j] == 0) \{f[1 << tk][i][j] = 0; tk++;\}
51
52
        void dfs(int st, int x, int y) {
53
54
            vis[x][y] = 1;
             PRE tmp = pre[st][x][y];
55
56
             if (tmp.x == 0 && tmp.y == 0) return ;
57
             dfs(tmp.s, tmp.x, tmp.y);
58
             if (tmp.x == x && tmp.y == y) dfs(st - tmp.s, tmp.x, tmp.y);
59
        void sol(int _n, int _m) {
60
61
             n = _n; m = _m;
            init();
62
63
             for (int st = 0; st < stn; ++st) {</pre>
                 for (int i = 1; i <= n; ++i)
64
                     for (int j = 1; j <= m; ++j) {
65
                          for (int s = st&(st-1); s; s = st&(s-1))
66
67
                          if(f[st-s][i][j]+f[s][i][j]-G[i][j] < f[st][i][j]){</pre>
68
                              f[st][i][j] = f[st-s][i][j]+f[s][i][j]-G[i][j];
                              pre[st][i][j] = (PRE) {s, i, j};
69
70
                          if (f[st][i][j]!=INF) {
71
72
                              q.push(make_pair(i,j));
73
                              vis[i][j]=1;
74
                          }
75
76
                 spfa(st);
77
             }
78
             int ansx, ansy, fg = 0;
             for(int i = 1; i <= n && !fg; ++i)</pre>
79
                for(int j = 1; j <= m; ++j)</pre>
80
                if(!G[i][j]) {ansx = i; ansy = j; fg = 1; break;}
81
             printf("%d\n", f[stn-1][ansx][ansy]);
82
            memset(vis, 0, sizeof vis);
83
84
             dfs(stn-1, ansx, ansy);
             for(int i = 1; i <= n; i++, puts("")) {</pre>
85
86
                 for(int j = 1; j <= m; j++) {</pre>
87
                     if(G[i][j] == 0) putchar('x');
88
                     else if(vis[i][j]) putchar('o');
89
                     else putchar('_');
90
                 }
91
            }
92
        }
93
```

## 1.3.7 Tree Divide and Conquer

```
struct Edge {
1
2
       int to, nxt, w;
3
   }e[MAXM];
   int head[MAXN], ecnt;
4
5
   int sz[MAXN];
   int d[MAXN], t[5], ans;
6
7
   bool vis[MAXN];
   inline void add_edge(int u, int v, int w) {
8
9
        e[++ecnt] = (Edge) {v, head[u], w}; head[u] = ecnt;
10
        e[++ecnt] = (Edge) {u, head[v], w}; head[v] = ecnt;
11
   }
```

```
int getsz(int x, int fa) {
12
13
        sz[x] = 1;
14
        for(int i = head[x]; i; i = e[i].nxt) {
15
            int y = e[i].to;
16
            if(vis[y] || y == fa) continue;
17
            sz[x] += getsz(y, x);
18
        }
19
        return sz[x];
20
21
    int getrt(int x) {
22
        int tot = getsz(x, 0) >> 1;
23
        while(1) {
24
            int u = -1;
25
            for(int i = head[x]; i; i = e[i].nxt) {
26
                 int y = e[i].to;
27
                 if(vis[y] || sz[y] > sz[x]) continue;
28
                 if(u == -1 || sz[y] > sz[u]) u = y;
29
30
            if(\sim u \&\& sz[u] > tot) x = u;
31
            else break;
32
33
        return x;
34
   }
35
   void getdep(int x, int fa) {
36
        t[d[x]]++;
37
        for(int i = head[x]; i; i = e[i].nxt) {
38
            int y = e[i].to;
            if(vis[y] || y == fa) continue;
39
40
            d[y] = (d[x] + e[i].w) % 3;
41
            getdep(y, x);
42
43
    int cal(int x, int v) {
44
        t[0] = t[1] = t[2] = 0;
45
        d[x] = v \% 3;
46
47
        getdep(x, 0);
        return t[0] * t[0] + t[1] * t[2] * 2;
48
49
   }
50
   void solve(int x) {
51
        vis[x] = 1;
52
        ans += cal(x, 0);
53
        for(int i = head[x]; i; i = e[i].nxt) {
54
            int y = e[i].to;
55
            if(vis[y]) continue;
56
            ans -= cal(y, e[i].w);
57
            solve(getrt(y));
58
        }
59
60
    int main() {
61
        solve(getrt(1));
62
```

## 1.3.8 Dominator Tree

```
#define LL long long
#define FILE "dagch"
using namespace std;

const int N = 200010;
struct Node{int to,next;}E[N<<1];</pre>
```

```
7
   int n,m,q,head[N],tot,dfn[N],clo,rev[N],fa[N],semi[N],Ans[N];
8
    vector<int>G[N];
9
    struct Union_Merge_Set{
10
        int fa[N],Mi[N];
11
        inline void init(){
12
             for(int i=0;i<=n;++i)</pre>
13
                 fa[i]=Mi[i]=semi[i]=i;
14
        inline int find(int x){
15
16
             if(x==fa[x])return x;
             int fx=fa[x],y=find(fa[x]);
17
             if(dfn[semi[Mi[fx]]]<dfn[semi[Mi[x]]])Mi[x]=Mi[fx];</pre>
18
19
             return fa[x]=y;
20
21
    }uset;
22
    inline void tarjan(int u) {
23
        rev[dfn[u] = ++tarjan_time] = u;
24
        for(auto v : G[u])
25
26
        if(!dfn[v]) {
27
             fa[v] = u;
28
             tarjan(v);
29
        }
30
    }
31
    inline void get_semi() {
32
        for(int i = tarjan_time; i >= 2; i--) {
33
             int u = rev[i], tsemi = n;
34
             for(auto v : rG[u]) {
                 if(!dfn[v]) continue;
35
36
                 if(dfn[v] < dfn[u]) tsemi = min(tsemi, dfn[x]);</pre>
37
                 else{
38
                      uset.find(x);
                      tsemi = min(tsemi, dfn[semi[uset.Mi[x]]]);
39
                 }
40
             }
41
             uset.fa[y] = fa[y];
42
             semi[y] = rev[tsemi];
43
            Ans[rev[tsemi]]++;
44
45
        }
46
    }
47
48
    inline void solve() {
        scanf("%d %d %d", &n, &m, &q);
49
50
        fa[1]=1;
51
        for(int i = 1, u, v; i <= m; ++i){
52
             scanf("%d%d", &u,&v);
             link(v,u);
53
54
            G[u].push_back(v);
55
56
        for(int i = 1; i <= n; i++)</pre>
57
             if(G[i].size())
58
                 sort(G[i].begin(), G[i].end());
59
        uset.init();
60
61
62
        tarjan(1);
63
        build();
64
        for(int i=1;i<=q;++i)</pre>
             printf("%d ",Ans[gi()]);
65
66
        printf("\n");
67
        for(int i=0;i<=n;++i){</pre>
```

```
68
            G[i].clear();head[i]=0;
69
            Ans[i]=semi[i]=fa[i]=0;
70
        }
71
        clo=tot=0;
72
    }
73
74
    int main() {
        int T; scanf("%d", &T);
75
76
        while(T--) solve();
77
        return 0;
78
    }
```

## 1.4 LCA

## 1.4.1 Tree Decomposition LCA

见树链剖分

## 1.4.2 Tarjan LCA

```
vector< pair<int,int> > G[MAXN],ask[MAXN];
   int fa[MAXN], ans[MAXN], vis[MAXN] ,dis[MAXN];
3
   int Find(int x){
4
        return x == fa[x] ? x : fa[x] = Find(fa[x]);
5
   }
6
   void init(int n){
7
        memset(ans, 0,sizeof ans);
8
        memset(vis, 0,sizeof vis);
9
        for(int i = 0; i <= n; i++){</pre>
10
            G[i].clear();
11
            ask[i].clear();
12
        }
13
    void LCA(int u){
14
15
        int v;
16
        fa[u] = u;
17
        vis[u] = true;
        for(auto it : ask[u])
18
            if(vis[v = it.first])
19
20
                ans[it.second] = dis[u] + dis[v] - 2 * dis[Find(it.first)];
21
        for(auto it : G[u])
22
        if(!vis[v = it.first]){
23
            dis[v] = dis[u] + it.second;
24
            LCA(v);
25
            fa[v] = u;
26
        }
27
    }
```

## 1.5 Tarjan

## 1.5.1 SCC

```
namespace SCC{
vector<int> G[MAXN];
int dfs_clock, scc_cnt, dfn[MAXN], low[MAXN], sccno[MAXN];
stack<int> S;
void addEdge(int u, int v) {
```

```
6
            G[u].push_back(v);
7
        }
8
        void tarjan(int u) {
9
             dfn[u] = low[u] = ++dfs_clock;
            S.push(u);
10
             for(auto v : G[u]) {
11
12
                 if(!dfn[v]) {
13
                     tarjan(v);
14
                     low[u] = min(low[u], low[v]);
15
                 }else if(!sccno[v]) {
16
                     low[u] = min(low[u], dfn[v]);
17
18
19
            if(dfn[u] == low[u]) {
20
                 scc_cnt++;
21
                 for(;;) {
                     int v = S.top(); S.pop();
22
23
                     sccno[v] = scc_cnt;
24
                     if(v == u) break;
25
                 }
26
            }
27
        void findSCC(int n) {
28
29
            for(int i = 1; i <= n; i++)
30
                 if(!dfn[i]) tarjan(i);
31
32
        void init(int n){
33
            dfs_clock = scc_cnt = 0;
34
             for(int i = 0;i <= n;++i){</pre>
35
                 dfn[i] = low[i] = sccno[i] = 0;
36
                 G[i].clear();
37
            }
        }
38
39
```

## 1.5.2 BCC

```
1
    namespace BCC{
2
        struct Edge {
3
            int to, nxt;
4
        }e[MAXM << 1];
        int ecnt, head[MAXN];
5
6
        int dfs_clock, dfn[MAXN], low[MAXN];
7
8
        int is_vertex[MAXN], vbcc_cnt, vbccno[MAXN];
9
        vector<int> vbcc[MAXN];
10
        stack<int> vS;
11
12
        int ebcc_cnt, ebccno[MAXN];
13
        stack<int> eS;
14
15
        inline void addEdge(int u, int v) {
            e[++ecnt] = (Edge) {v, head[u]}; head[u] = ecnt;
16
            e[++ecnt] = (Edge) {u, head[v]}; head[v] = ecnt;
17
18
        inline void init(int n) {
19
20
            ecnt = 1;
21
            dfs_clock = 0;
22
            vbcc cnt = 0;
23
            ebcc_cnt = 0;
```

```
24
            for(int i = 1; i <= n; ++i){
25
                 head[i] = dfn[i] = low[i] = 0;
26
                 is_vertex[i] = 0;
27
                 vbccno[i] = 0;
28
                 ebccno[i] = 0;
29
            }
30
            while(!vS.empty()) vS.pop();
31
32
        //root's edge = -1;
        void tarjan(int u, int edge) {
33
34
            dfn[u] = low[u] = ++dfs_clock;
35
            int ch = 0;
36
            vS.push(u);
37
            eS.push(u);
38
            for(int i = head[u], v; i; i = e[i].nxt) {
39
                 if(!dfn[v = e[i].to]) {
                     tarjan(v, i ^ 1);
40
41
                     low[u] = min(low[u], low[v]);
                     if(low[v] >= dfn[u]) {
42
43
                         if(edge > 0 || ch > 1) is_vertex[u] = 1;
44
45
                         vbcc[++vbcc_cnt].clear();
46
                         vbcc[vbcc_cnt].push_back(u);
                         for(int x;;){
47
48
                              x = vS.top();vS.pop();
49
                              vbcc[vbcc_cnt].push_back(x);
                              vbccno[x] = vbcc_cnt;
50
51
                              if(x == v)break;
52
                         }
53
                     if(low[v] > dfn[u]) {
54
                     // i && i ^ 1 is bridge
55
56
57
                 else if(dfn[v] < dfn[u] && i != edge)</pre>
58
59
                     low[u] = min(low[u], dfn[v]);
60
            if(dfn[u] == low[u]) {
61
62
                 ebcc cnt++;
63
                 for(int v;;) {
64
                     v = eS.top(); eS.pop();
65
                     ebccno[v] = ebcc_cnt;
66
                     if(v == u) break;
67
                 }
68
            }
69
        void findBCC(int n){
70
            for(int i = 1; i <= n; i++)</pre>
71
72
                 if(!dfn[i]) tarjan(i, -1);
73
74
            //findBridge
            for(int u = 1; u <= n; u++) {</pre>
75
                 for(int i = head[u], v; i; i = e[i].nxt)
76
77
                 if(ebccno[u] != ebccno[v = e[i].to]) {
78
                     //is bridge
79
                 }
80
            }
81
        }
82
    }
```

#### 1.6 Cactus

#### 1.6.1 Circle-Square Tree

```
1 #include <bits/stdc++.h>
2 using namespace std;
3 typedef pair<int, int> P;
4 const int MAXN = 2e4 + 5;
5 const int S = 15;
6
   namespace Tree {
7
        struct Edge {
8
            int to, nxt, w;
9
        }e[MAXN << 1];
10
        int ecnt, head[MAXN];
11
        int rt, isrt[MAXN], fa[MAXN][S + 3];
12
        int sz[MAXN];
13
        inline void addEdge(int u, int v, int w) {
14
            e[++ecnt] = (Edge) {v, head[u], w}; head[u] = ecnt;
15
            fa[v][0] = u;
16
        }
   }
17
18
   int n, m, Q;
   namespace BCC {
19
20
        struct Edge {
            int to, nxt, w;
21
22
        }e[MAXN << 1];
        int ecnt, head[MAXN];
23
24
        int dfs clock, dfn[MAXN], low[MAXN];
25
        int is vertex[MAXN], vbcc cnt, vbccno[MAXN];
26
        vector<P> vbcc[MAXN];
27
        stack<P> vs;
28
        int tag[MAXN];
29
        inline void addEdge(int u, int v, int w) {
30
            e[++ecnt] = (Edge) {v, head[u], w}; head[u] = ecnt;
31
            e[++ecnt] = (Edge) {u, head[v], w}; head[v] = ecnt;
32
33
        inline void init(int n) {
34
            ecnt = 1;
35
            dfs_clock = 0;
36
            vbcc cnt = 0;
            for(int i = 0; i \le 2 * n; i++){
37
                head[i] = dfn[i] = low[i] = 0;
38
39
                vbccno[i] = 0;
40
                tag[i] = 0;
41
42
            while(!vs.empty()) vs.pop();
43
44
        //root's edge = -1;
        void tarjan(int u, int edge) {
45
46
            dfn[u] = low[u] = ++dfs clock;
47
            vs.push(P(u, e[edge ^ 1].w));
48
            for(int i = head[u], v; i; i = e[i].nxt) {
49
                if(!dfn[v = e[i].to]) {
                    tarjan(v, i ^ 1);
50
                    low[u] = min(low[u], low[v]);
51
52
                    if(low[v] >= dfn[u]) {
53
                         if(vs.top().first == v) {
54
                             Tree::addEdge(u, v, vs.top().second);
55
                             vs.pop();
56
                             continue;
```

```
57
58
                          vbcc[++vbcc_cnt].clear();
59
                          vbcc[vbcc_cnt].push_back(P(u, 0));
60
                          Tree::isrt[u] = 1;
61
                          int &sz = Tree::sz[n + vbcc_cnt];
62
                          tag[vs.top().first] = n + vbcc_cnt;
63
                          //Tree::addEdge(u, rt, 0);
                          for(P x;;) {
64
65
                               x = vs.top(); vs.pop();
66
                               sz += x.second;
67
                               //Tree :: addEdge(rt, x.first, sz);
                               vbcc[vbcc_cnt].push_back(x);
68
                               vbccno[x.first] = vbcc_cnt;
69
                               if(x.first == v) break;
70
71
                          }
72
                      }
73
74
                 else if(dfn[v] < dfn[u] && i != edge)</pre>
75
                      low[u] = min(low[u], dfn[v]);
76
77
             for(int i = head[u], v; i; i = e[i].nxt) {
78
                 if(tag[v = e[i].to]) {
79
                      int r = tag[v]; Tree::sz[r] += e[i].w;
80
                      tag[v] = 0;
81
                 }
82
             }
83
         void findBCC(int n) {
84
85
             for(int i = 1; i <= n; i++)</pre>
86
                 if(!dfn[i]) tarjan(i, -1);
87
    }
88
89
    namespace Tree {
90
         int dis[MAXN], dep[MAXN], len[MAXN];
         inline void init(int n) {
91
92
             BCC::init(n);
93
             rt = n;
             ecnt = 1;
94
95
             for(int i = 0; i <= 2 * n; i++) {
96
                 head[i] = 0;
97
                 fa[i][0] = isrt[i] = dis[i] = dep[i] = len[i] = 0;
98
             }
99
         void dfs(int x) {
100
101
             for(int i = head[x], y; i; i = e[i].nxt) {
102
                 if(!dep[y = e[i].to]) {
103
                      dep[y] = dep[x] + 1;
104
                      dis[y] = dis[x] + e[i].w;
105
                      dfs(y);
106
                 }
107
             }
108
109
         void pre() {
             for(int k = 1; k <= BCC::vbcc_cnt; k++) {</pre>
110
111
                 vector<P> &E = BCC::vbcc[k];
112
113
                 addEdge(E[0].first, rt, 0);
                 int cnt = 0;
114
                 for(int i = E.size() - 1; i >= 1; i--) {
115
116
                      cnt += E[i].second;
117
                      len[E[i].first] = cnt;
```

```
118
                      addEdge(rt, E[i].first, min(cnt, sz[rt] - cnt));
119
                  }
120
             }
121
             for(int k = 1; k <= S; k++) {</pre>
122
                  for(int i = 1; i <= rt; i++) {</pre>
123
                      fa[i][k] = fa[fa[i][k - 1]][k - 1];
124
125
             }
126
             dep[1] = 1;
127
             dfs(1);
128
         int up(int x, int d) {
129
             for(int i = S; i >= 0; i--) {
130
131
                  if(dep[fa[x][i]] >= d) x = fa[x][i];
132
             }
             return x;
133
134
135
         int lca(int u, int v) {
136
             if(dep[u] > dep[v]) swap(u, v);
137
             v = up(v, dep[u]);
138
             if(u == v) return u;
             for(int i = S; i >= 0; i--) {
139
                  if(fa[u][i] != fa[v][i]) {
140
141
                      u = fa[u][i], v = fa[v][i];
142
143
             }
144
             return fa[u][0];
145
146
         int query(int u, int v) {
147
             int 1 = 1ca(u, v);
             if(1 <= n) return dis[u] + dis[v] - 2 * dis[1];</pre>
148
149
             int x = up(u, dep[1] + 1), y = up(v, dep[1] + 1);
             int res = dis[u] - dis[x] + dis[v] - dis[y];
150
             int tmp = abs(len[x] - len[y]);
151
             return res + min(tmp, sz[1] - tmp);
152
153
         }
154
    }
155
156
    int main() {
157
         ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
158
         using namespace Tree;
159
         cin >> n >> m >> Q;
160
         init(n);
161
         for(int i = 1, u, v, w; i <= m; i++) {
162
             cin >> u >> v >> w;
163
             BCC::addEdge(u, v, w);
164
         BCC::findBCC(n);
165
         pre();
166
167
         int u, v;
168
         while(Q--) {
169
             cin >> u >> v;
170
             cout << query(u, v) << endl;</pre>
171
172
         return 0;
173
```

## 2 Data Structures

## 2.1 Basic Structures

#### 2.1.1 RMQ

```
struct RMQ {
1
2
        int d[MAXN][S + 2];
        inline void init(int *a, int n) {
3
4
            for(int i = 1; i <= n; i++) d[i][0] = a[i];</pre>
5
            for(int k = 1; (1 << k) <= n; k++)
6
                for(int i = 1; i + (1 << k) - 1 <= n; i++)
7
                     d[i][k] = min(d[i][k - 1], d[i + (1 << (k - 1))][k - 1]);
8
9
        inline int query(int 1, int r) {
10
            if(1 > r) swap(1, r);
            int k = 0;
11
12
            while ((1 << (k + 1)) <= r - 1 + 1) k++;
            return min(d[1][k], d[r - (1 << k) + 1][k]);
13
14
        }
15
   }rmq;
16
   const int MAXM = 2e5 + 5, MAXN = 3e6 + 5, S = 22;
17
    const LL INF = 1e18;
18
   #define belong(x) (x / S + 1)
19
   #define pos(x) (x % S + 1)
   int Log[MAXN], sz;
20
21
   struct RMQ {
        LL a[MAXN];
22
        LL d[MAXM][S + 2];
23
24
        LL pre[MAXM][S + 2], aft[MAXM][S + 2];
        inline void init(int n) {
25
26
            sz = n / S + 1;
27
            Log[0] = -1; for(int i = 1; i <= n; i++) Log[i] = Log[i / 2] + 1;
28
            for(int i = 1; i <= sz; i++) {</pre>
29
                pre[i][0] = aft[i][S + 1] = INF;
30
            }
            for(int i = 1; i <= n; i++) {</pre>
31
32
                pre[belong(i)][pos(i)] = min(pre[belong(i)][pos(i) - 1], a[i]);
33
            }
            for(int i = n; i >= 1; i--) {
34
35
                aft[belong(i)][pos(i)] = min(aft[belong(i)][pos(i) + 1], a[i]);
36
37
            for(int i = 1; i <= sz; i++) {
38
                d[i][0] = aft[i][1];
39
            for(int k = 1; k <= S; k++)</pre>
40
                for(int i = 1; i + (1 << k) <= sz; i++)
41
                     d[i][k] = min(d[i][k-1], d[i+(1 << (k-1))][k-1]);
42
43
        inline LL ask(int 1, int r) {
44
            assert(1 <= r);</pre>
45
46
            LL res = INF;
            if(belong(1) == belong(r)) {
47
                for(int i = 1; i <= r; i++) res = min(res, a[i]);</pre>
48
49
50
            res = min(aft[belong(1)][pos(1)], pre[belong(r)][pos(r)]);
51
            int k = Log[belong(r) - belong(l) - 1];
52
53
            if(~k) {
```

## 2.1.2 Divide Blocks

```
int belong[MAXN], 1[MAXN], r[MAXN];
    int sz, num;
3
   void build(int n) {
4
        sz = sqrt(n);
        num = n / sz; if(n % sz) num++;
5
6
        for(int i = 1; i <= num; i++) {</pre>
7
            l[i] = (i - 1) * sz + 1;
            r[i] = i * sz;
8
9
        }
10
        r[num] = n;
11
        for(int i = 1; i <= n; i++) {</pre>
12
            belong[i] = (i - 1) / sz + 1;
13
14
   }
```

## 2.2 Heap Structures

## 2.2.1 Leftist Tree

```
const int MAXN = ;
2
   namespace LeftistTree{
3
        int ls[MAXN], rs[MAXN];
4
        int dis[MAXN];
        int fg[MAXN], sfa[MAXN], rt[MAXN];//利用rt得到堆根节点
5
6
        int val[MAXN];
7
        void push_down(int x) {};
8
9
        void push_up(int x) {
10
            if(dis[ls[x]] < dis[rs[x]]) swap(ls[x], rs[x]);</pre>
11
            dis[x] = dis[rs[x]] + 1;
12
        int merge(int x, int y) {
13
            if(!x || !y) return x^y;
14
            if(val[x] > val[y] \mid | (val[x] == val[y] && x > y)) swap(x, y);
15
16
            push_down(x);
17
            rs[x] = merge(rs[x], y);
18
            push_up(x);
            return x;
19
20
21
        int getSfa(int x) {return sfa[x] == x ? x : sfa[x] = getSfa(sfa[x]);}
22
        int uni(int x, int y) {//返回合并后的根
            if(!x || !y) return x^y;
23
            if(fg[x] || fg[y]) return;
24
25
            x = getSfa(x); y = getSfa(y);
26
            if(x == y) return;
27
            int z = merge(x, y);
28
            return sfa[x] = sfa[y] = sfa[z] = z;
29
        }
       void uni2(int a, int b) {
```

```
31
            //val[a] \gg = 1;
32
           int c = merge(ls[a], rs[a]);
33
           ls[a] = rs[a] = dis[a] = 0;
34
           int a1 = merge(c, a);
35
            //val[b] \gg = 1;
36
           c = merge(ls[b], rs[b]);
37
           ls[b] = rs[b] = dis[b] = 0;
38
            int b1 = merge(c, b);
39
            c = merge(a1, b1);
40
           sfa[a] = sfa[b] = sfa[c] = c;
            printf("%d\n", val[c])
41
42
       int pop(int x) {//返回堆顶值,也可以用于返回根
43
           if(!x || fg[x]) return -1;
44
45
           x = getSfa(x); fg[x] = 1;
            push_down(x);//在删除堆顶时要下传标记
46
           int y = merge(ls[x], rs[x]);
47
           sfa[x] = sfa[y] = y;
48
49
           return val[x];
50
       void init(int n) {
51
           for(int i = 1; i <= n; i++) {
52
53
                sfa[i] = i;
                ls[i] = rs[i] = dis[i] = fg[i] = 0;
54
55
            }
56
       }
57
   //可持久化版本见k短路
```

## 2.3 Sequence Structures

## 2.3.1 Cartesian Tree

```
struct CartesianTree{
1
2
        int rt, fa[MAXN], ls[MAXN], rs[MAXN];
3
        int top, st[MAXN];
        int cnt[MAXN];
4
5
        void build(LL *a,int n) {
            top = rt = 0;
6
7
            for(int i = 1; i <= n; i++) {
8
                ls[i] = rs[i] = fa[i] = 0;
9
                while(top && a[st[top]] > a[i]) ls[i] = st[top--];
10
                fa[i] = st[top];
                if(ls[i]) fa[ls[i]] = i;
11
12
                if(fa[i]) rs[fa[i]] = i; else rt = i;
13
                st[++top] = i;
14
            }
15
        void dfs(int x) {
16
17
            cnt[x] = 1;
18
            if(ls[x]) {dfs(ls[x]); cnt[x] += cnt[ls[x]];}
19
            if(rs[x]) {dfs(rs[x]); cnt[x] += cnt[rs[x]];}
20
        LL getAns(LL *a, int n) {
21
22
            //dfs(rt);
23
24
            return res;
25
        }
   }T;
26
```

#### 2.3.2 TreeArray

```
1 //树状数组上二分
2 int BS(int x) {
3    int res = 0;
4    for (int i = 1 << 18; i; i >>= 1)
5        if ((res | i) <= Tn && T[res | i] <= x)
             x -= T[res |= i];
7    return res;
8 }</pre>
```

## 2.3.3 Segment Tree

```
1
   #define Ls(x) (x << 1)
2
3
   #define Rs(x) (x << 1 | 1)
   struct Tree {
5
        int 1, r, lazy;
6
        LL sum, mx;
7
   }tree[MAXN << 2];</pre>
   int A[MAXN];
8
9
   void push_up(int x) {
10
       tree[x].sum = tree[Ls(x)].sum + tree[Rs(x)].sum;
        tree[x].mx = max(tree[Ls(x)].mx, tree[Rs(x)].mx);
11
12
   }
13
   void push_down(int x) {
        if(tree[x].lazy) {
14
15
            tree[Ls(x)].sum += tree[x].lazy * (tree[Ls(x)].r - tree[Ls(x)].l + 1);
16
            tree[Rs(x)].sum += tree[x].lazy * (tree[Rs(x)].r - tree[Rs(x)].l + 1);
17
            tree[Ls(x)].mx += tree[x].lazy;
            tree[Rs(x)].mx += tree[x].lazy;
18
            tree[Ls(x)].lazy += tree[x].lazy;
19
20
            tree[Rs(x)].lazy += tree[x].lazy;
21
            tree[x].lazy = 0;
22
23
   }
   void build(int x, int L, int R) {
24
25
       tree[x].lazy = 0;
26
        tree[x].1 = L; tree[x].r = R;
27
        if(L == R) {
28
            tree[x].sum = A[L];
29
            tree[x].mx = A[L];
30
31
            return;
32
33
        int mid = (L + R) \gg 1;
34
        build(Ls(x), L, mid);
35
        build(Rs(x), mid + 1, R);
36
        push_up(x);
37
38
   void update(int x, int L, int R, LL val) {
        if(tree[x].1 >= L && tree[x].r <= R) {</pre>
39
40
            tree[x].lazy += val;
            tree[x].sum += val * (tree[x].r - tree[x].l + 1);
41
            tree[x].mx += val;
42
43
            return;
44
45
        push_down(x);
        int mid = (tree[x].l + tree[x].r) >> 1;
46
```

```
if(L <= mid) update(Ls(x), L, R, val);</pre>
47
48
        if(R > mid) update(Rs(x), L, R, val);
49
        push_up(x);
50
    }
    LL query(int x, int L, int R) {
51
        if(tree[x].l >= L \&\& tree[x].r <= R)
52
53
             return tree[x].sum;
54
        push down(x);
55
        int mid = (tree[x].l + tree[x].r) >> 1;
56
        LL res = 0;
57
        if(L <= mid) res += query(Ls(x), L, R);</pre>
58
        if(R > mid) res += query(Rs(x), L, R);
59
60
        return res;
61
62
   LL query2(int x, int L, int R) {
        if(tree[x].l >= L \&\& tree[x].r <= R)
63
64
            return tree[x].mx;
        push down(x);
65
66
        int mid = (tree[x].l + tree[x].r) >> 1;
67
        LL res = -INF;
68
        if(L <= mid) res = max(res, query2(Ls(x), L, R));</pre>
69
        if(R > mid) res = max(res, query2(Rs(x), L, R));
70
        return res;
71
    }
```

#### 2.3.4 LiChao Tree

```
const double eps = 1e-12;
2
    namespace LiT{
3
        const int MLIMIT = 40000;
4
        typedef double LD;
        struct line{LD k,b;int l,r,id;} T[MAXN << 2];</pre>
5
6
        //inline\ LD\ calc(line\ \&a,int\ pos)\ \{return\ a.k*vec[pos]+a.b;\}
7
        inline LD calc(line &a,int pos) {return a.k*pos+a.b;}
8
        inline double cross(line &a,line &b) {
9
            if(b.k == a.k) return -1e9;
10
            return (double)(a.b-b.b)/(b.k-a.k);
11
        void build(int v, int l, int r) {
12
13
            T[v].k = 0;T[v].b = -1e18;
14
            T[v].1 = 0;T[v].r = MLIMIT;
15
            T[v].id = 0;
16
            if(1 == r)return;
17
            int mid = (1+r) >> 1;
            build(v<<1,1,mid);</pre>
18
19
            build(v << 1|1, mid+1, r);
20
21
        void ins(int v,int l,int r, line k) {
22
            if(k.1 <= 1 && r <= k.r) {</pre>
23
                 LD fl = calc(k, 1), fr = calc(k, r);
24
                 LD gl = calc(T[v], 1), gr = calc(T[v], r);
25
                 if(fl - gl > eps \&\& fr - gr > eps) T[v] = k;
                 else if(fl - gl > eps || fr - gr > eps) {
26
                     int mid = (1+r)>>1;
27
28
                     if(calc(k, mid) - calc(T[v], mid) > eps) swap(k, T[v]);
                     //if(vec[mid] - cross(k, T[v]) > eps)
29
30
                     if(mid - cross(k, T[v]) > eps)
31
                         ins(v << 1, l, mid, k); else ins(v << 1|1, mid+1, r, k);
32
```

```
33
                 return;
34
            }
35
            int mid=(l+r)>>1;
36
            if(k.l <= mid) ins(v<<1, l, mid, k);</pre>
37
            if(mid < k.r) ins(v<<1|1, mid+1, r, k);</pre>
38
39
        LD ans; int ansid;
40
        void que(int v, int l, int r, int x) {
41
             LD tmp = calc(T[v], x);
             if(tmp > ans || (tmp == ans && T[v].id < ansid)) {</pre>
42
43
                 ans = tmp;
                 ansid = T[v].id;
44
45
46
            if(1 == r) return;
47
            int mid = (l+r)>>1;
48
            if(x <= mid) que(v<<1,1,mid,x); else que(v<<1|1,mid+1,r,x);
49
        }
50
   };
   //左闭右闭
```

## 2.3.5 Splay Tree

```
1
    namespace splay{
2
        int n, m, sz, rt;
3
        int val[MAXN], id[MAXN];
        int tr[MAXN][2], size[MAXN], fa[MAXN], rev[MAXN], s[MAXN], lazy[MAXN];
4
        void push_up(int x) {
5
6
            int 1 = tr[x][0], r = tr[x][1];
7
            s[x] = max(val[x], max(s[1], s[r]));
8
            size[x] = size[l] + size[r] + 1;
9
10
        void push_down(int x) {
            int 1 = tr[x][0], r = tr[x][1];
11
            if(lazy[x]) {
12
13
                if(1) {
                    lazy[1] += lazy[x];
14
                     s[1] += lazy[x];
15
                     val[1] += lazy[x];
16
17
                if(r) {
18
19
                    lazy[r] += lazy[x];
20
                     s[r] += lazy[x];
21
                    val[r] += lazy[x];
22
23
                lazy[x] = 0;
24
            }
            if(rev[x]) {
25
26
                 rev[x] = 0;
27
                 rev[1] ^= 1; rev[r] ^= 1;
                swap(tr[x][0], tr[x][1]);
28
29
            }
30
31
        void rotate(int x, int &k) {
32
            int y = fa[x];
            int z = fa[y];
33
            int 1, r;
34
            if(tr[y][0] == x) 1 = 0;
35
36
            else l = 1;
37
            r = 1 ^1;
38
            if(y == k) k = x;
```

```
39
            else {
40
                if(tr[z][0] == y) tr[z][0] = x;
41
                else tr[z][1] = x;
42
43
            fa[x] = z; fa[y] = x; fa[tr[x][r]] = y;
            tr[y][1] = tr[x][r]; tr[x][r] = y;
44
45
            push_up(y); push_up(x);
46
47
        void splay(int x, int &k) {
48
            int y, z;
            while(x != k) {
49
                y = fa[x];
50
51
                z = fa[y];
                if(y != k) {
52
53
                     if((tr[y][0] == x) ^ (tr[z][0] == y)) rotate(x, k);
54
55
                     else rotate(y, k);
56
57
                 rotate(x, k);
58
            }
59
        int find(int x, int rank) {
60
61
            push_down(x);
62
            int 1 = tr[x][0], r = tr[x][1];
63
64
            if(size[l] + 1 == rank) return x;
65
            else if(size[1] >= rank) return find(1, rank);
66
            else return find(r, rank - size[l] - 1);
67
        void update(int 1, int r, int v) {
68
            int x = find(rt, 1), y = find(rt, r + 2);
69
70
            splay(x, rt); splay(y, tr[x][1]);
            int z = tr[y][0];
71
            lazy[z] += v;
72
73
            val[z] += v;
74
            s[z] += v;
75
        }
        void reverse(int 1, int r) {
76
77
            int x = find(rt, 1), y = find(rt, r + 2);
78
            splay(x, rt); splay(y, tr[x][1]);
79
            int z = tr[y][0];
80
            rev[z] ^= 1;
81
        void query(int 1, int r) {
82
            int x = find(rt, 1), y = find(rt, r + 2);
83
84
            splay(x, rt); splay(y, tr[x][1]);
            int z = tr[y][0];
85
            printf("%d\n", s[z]);
86
87
        void build(int 1, int r, int f) {
88
89
            if(1 > r) return;
90
            int now = id[1], last = id[f];
            if(1 == r) {
91
                fa[now] = last; size[now] = 1;
92
                if(1 < f) tr[last][0] = now;
93
                else tr[last][1] = now;
94
95
                return;
96
            }
            int mid = (l + r) \gg 1; now = id[mid];
97
            build(1, mid - 1, mid); build(mid + 1, r, mid);
98
99
            fa[now] = last;
```

```
100
             push_up(now);
101
             if(mid < f) tr[last][0] = now;</pre>
102
             else tr[last][1] = now;
103
         void init() {
104
105
             s[0] = -INF;
             scanf("%d%d", &n, &m);
106
107
             for(int i = 1; i <= n + 2; i++) id[i] = ++sz;
108
             build(1, n + 2, 0); rt = (n + 3) >> 1;
109
110
    }
111
    namespace splay{
112
         //内存回收池见fhq_treap
         int tcnt, root;
113
114
        int sz[MAXN];
115
         int tr[MAXN][2], fa[MAXN];
116
         int val[MAXN];
117
         int newnode(int w) {
             ++tcnt;
118
119
             sz[tcnt] = 1;
120
             fa[tcnt] = tr[tcnt][0] = tr[tcnt][1] = 0;
121
             //val[tcnt] = w;
122
             return tcnt;
123
         void push_up(int v) {
124
125
             int 1 = tr[v][0], r = tr[v][1];
126
             sz[v] = sz[1] + 1 + sz[r];
127
128
         void push down(int v) {
129
             if(!v) return;
130
         void init() {
131
132
             tcnt = 2;
             tr[root = fa[1] = 2][0] = 1;
133
             sz[1] = 1; sz[2] = 2;
134
             //val[1] = -INF; val[2] = INF;//权值平衡树
135
             //val[1] = val[2] = 0;//位置平衡树
136
             //1,2为哨兵节点,根据题意也可设置为n+1,n+2或1,n+1
137
138
        }
139
         void rotate(int x) {
140
             int y = fa[x], z = fa[y];
141
             push_down(y);push_down(x);
142
             int lr = tr[y][1] == x;
143
             if(z) tr[z][tr[z][1]==y] = x;
144
             fa[x] = z;
145
             fa[tr[y][lr] = tr[x][lr^1]] = y;
             fa[tr[x][lr^1] = y] = x;
146
147
             push_up(y); push_up(x);
148
149
         void splay(int x, int k) {
             for(int y, z; (y = fa[x]) != k; rotate(x)) {
150
                 if((z = fa[y]) != k) {
151
                     if((tr[y][0] == x) ^ (tr[z][0] == y))
152
153
                         rotate(x); else rotate(y);
154
                 }
155
             if(!k) root = x;
156
157
158
         int find(int x, int rank) {
159
             push down(x);
160
             int l = tr[x][0], r = tr[x][1];
```

```
161
             if(sz[1] + 1 == rank) return x;
162
             if(sz[1] >= rank) return find(1, rank);
163
             return find(r, rank - sz[1] - 1);
164
165
         int build(int 1, int r) {
166
             if(1 > r) return 0;
             if(1 == r) {
167
                 int num; scanf("%d", &num);
168
169
                 return newnode(num);
170
             }
             int mid = (1 + r) >> 1;
171
             int ls = build(1, mid-1);
172
             int num; scanf("%d", &num);
173
174
             int v = newnode(num);
             int rs = build(mid+1, r);
175
             if(ls) fa[ls] = v;
176
177
             tr[v][0] = 1s;
             if(rs) fa[rs] = v;
178
             tr[v][1] = rs;
179
180
             push_up(v);
181
             return v;
182
         void insert(int pos, ...) {
183
             int x = find(root, pos+1), y = find(root, pos+2);
184
185
             splay(x, 0); splay(y, x);
186
             //int z = newnode(w); //插入一个节点
187
             //int z = build(1, n); //插入n个节点
188
             fa[tr[y][0] = z] = y;
189
             splay(z,0);
190
         void modifyOrQuery(int 1, int r, int v) {
191
             int x = find(root, 1), y = find(root, r + 2);
192
             splay(x, 0); splay(y, x);
193
             int z = tr[y][0];
194
             if(!z) return;
195
             //标记对本身无效,处理时将z点重新计算
196
197
             splay(z,0);
198
         }
         void display(int v) {
199
200
             if(!v) return;
201
             push_down(v);
202
             display(tr[v][0]);
             if(val[v]) printf("%d ", val[v]);
203
204
             display(tr[v][1]);
205
         /*int findValue(int v) {
206
207
             int res = root;
             for(int \ cur = root; cur; \ res = cur, \ cur = tr[cur][val[cur] <= v]);
208
209
             return res;
210
211
         void insert(int w) {
212
             int y = find Value(w);
213
             int z = newnode(w);
214
             fa[tr[y][val[y] \le w] = z] = y;
215
             splay(z,0);
         }*/
216
         /*void split(int v) {//splay维护区间[l,r],区间分裂为[l,k-1],[k,k],[k+1,r];
217
             //ump查看标号是否出现,mp维护子区间左端点
218
             if(ump.find(v) == ump.end()) {
219
                 auto\ it = mp.upper\_bound(v); --it;
220
221
                 int z = it -> second;
```

```
222
                      splay(z, 0);
223
                      int pos = sz[tr[z]/0]];
                      int \ x \, = \, find \, (\, root \, , \, pos) \, , \ y \, = \, find \, (\, root \, , \, pos + rc \, [\, z] - l\, c \, [\, z] + 2) \, ;
224
225
                      splay(x, 0); splay(y, x);
226
                      z = tr[y][0];
                      if(lc[z] != v)  {
227
                           tr[z][0] = newnode(lc[z], v-1);
228
229
                           fa[tr[z][0]] = z;
230
                           mp/lc/z// = tr/z//0/;
231
                      if(rc[z] != v)  {
232
                           tr[z][1] = newnode(v+1, rc[z]);
233
234
                           fa / tr / z / / 1 / / = z;
235
                          mp[v+1] = tr[z][1];
236
                     lc[z] = rc[z] = v;
237
238
                     splay(z,0);
239
                     mp / v / = z;
                     ump / v / = z;
240
241
           }*/
242
243
244
     }
```

#### 2.3.6 Scapegoat Tree

```
1
    struct ScapegoatTree{
2
        int Tsn; queue<int> q;
3
        int val[MAXM], ext[MAXM];
4
        int sz[MAXM], tsz[MAXM];
5
        int fa[MAXM], tr[MAXM][2];
6
        int root;
7
        double alp;
8
        void init() {
9
            root = 0;
10
            alp = 0.7;
11
        int newnode(int x) {
12
13
            if(q.empty()) q.push(++Tsn);
14
            int tcnt = q.front(); q.pop();
15
            val[tcnt] = x; ext[tcnt] = 1;
16
            fa[tcnt] = tr[tcnt][0] = tr[tcnt][1] = 0;
17
            sz[tcnt] = tsz[tcnt] = 0;
18
            return tcnt;
19
20
        void push_up(int v) {
21
            sz[v] = ext[v];tsz[v] = 1;
22
            if(tr[v][0]) {
23
                sz[v] += sz[tr[v][0]];
24
                tsz[v] += tsz[tr[v][0]];
25
            if(tr[v][1]) {
26
27
                 sz[v] += sz[tr[v][1]];
28
                tsz[v] += tsz[tr[v][1]];
29
            }
30
        bool isBad(int v) {
31
32
            return (double(tsz[ tr[v][0] ]) > double(tsz[v]) * alp) ||
33
                 (double(tsz[ tr[v][1] ]) > double(tsz[v]) * alp) ||
34
                 (sz[v] * 2 < tsz[v]);
```

```
35
36
        vector<int> vec;
37
        void rRecycle(int v) {
38
            if(tr[v][0]) rRecycle(tr[v][0]);
39
            if(ext[v]) vec.push_back(v); else q.push(v);
40
            if(tr[v][1]) rRecycle(tr[v][1]);
41
42
        int rBuild(int 1, int r) {
            int mid = (l + r) \gg 1, v = vec[mid];
43
            tr[v][0] = (1 \le mid-1) ? rBuild(1, mid - 1) : 0;
44
45
            if(tr[v][0]) fa[tr[v][0]] = v;
            tr[v][1] = (mid+1 <= r) ? rBuild(mid + 1, r) : 0;
46
47
            if(tr[v][1]) fa[tr[v][1]] = v;
48
            push_up(v);
49
            return v;
50
        void rebuild(int x) {
51
52
            int v = 0;
53
            for(;x; x= fa[x]) {
54
                 push_up(x);
55
                 if(isBad(x)) v = x;
56
            if(v && isBad(v)){
57
                 vec.clear();
58
59
                 int u = fa[v], lr = tr[u][1] == v;
60
                 rRecycle(v);
61
                 if(vec.size()) v = rBuild(0, vec.size() - 1); else v = 0;
62
                 if(u == 0) fa[root = v] = 0;
63
                 else{
64
                     tr[u][lr] = v;
65
                     if(v) fa[v] = u;
66
                 }
            }
67
        }
68
        void ins(int x) {
69
            int p = root, q = root;
70
71
            for(;p && val[p] != x; q = p, p = tr[p][ val[p] < x]);</pre>
72
            if(!q) {
73
                 p = root = newnode(x);
74
            }else if(p) {
75
                 ext[p]++;
76
            }else{
77
                 fa[p = tr[q][val[q] < x] = newnode(x)] = q;
            }
78
79
            rebuild(p);
80
        void del(int x) {
81
82
            int p = root;
83
            for(;p && val[p] != x; p = tr[p][ val[p] < x]);</pre>
            if(p && ext[p]){
84
85
                 --ext[p];
86
                 rebuild(p);
87
            }
88
89
        int get_rank(int x) {
            int ret = 0;
90
91
            for(int p = root;p;) {
                 if(val[p] < x) {</pre>
92
                     ret += sz[tr[p][0]] + ext[p];
93
                     p = tr[p][1];
94
95
                 }else p = tr[p][0];
```

```
96
             }
97
             return ret + 1;
98
99
         int get_Kth(int p, int k) {
100
             if(sz[tr[p][0]] >= k) return get_Kth(tr[p][0] ,k);
101
             k -= sz[tr[p][0]];
102
             if(ext[p] >= k) return val[p];
103
             k -= ext[p];
104
             return get_Kth(tr[p][1], k);
105
         int pre(int x) {
106
             int id = get_rank(x);
107
             return get_Kth(root, id - 1);
108
109
110
         int nxt(int x) {
             int id = get_rank(x + 1);
111
             return get_Kth(root ,id);
112
113
         void display(int v) {
114
115
             if(tr[v][0]) display(tr[v][0]);
             cerr<<val[v]<<" ";
116
117
             if(tr[v][1]) display(tr[v][1]);
         }
118
119
    }T;
```

#### 2.3.7 FHQ Treap

```
1
   namespace fhq_treap{
       int Tsz; queue<int> q; //内存回收池
2
3
       int tcnt, root;
       //int rt [MAXN]; //可持久化时使用rt,维护版本号(int &root), 空间开大
4
5
       int sz[MAXN], rnd[MAXN];
6
       int tr[MAXN][2];
        //int fa [MAXN];//维护fa时除了在pushup更新v节点左右孩子父节点信息还要在 split和merge结
7
       東时维护root的fa信息,fa[root]=0;
8
       int val[MAXN], rev[MAXN];
9
       void init() {
10
           srand(time(0));
11
           Tsz = tcnt = root = 0;
12
13
       int newnode(int v) {
14
           if(q.empty()) q.push(++Tsz);
15
           tcnt = q.front(); q.pop();
           sz[tcnt] = 1;
16
17
           rnd[tcnt] = rand();
           tr[tcnt][0] = tr[tcnt][1] = 0;
18
19
           //val[tcnt] = v;
20
           return tcnt;
21
22
       /*int copynode(int id) {
           //++tcnt;//获取一个新的节点编号
23
24
           sz/tcnt/ = sz/id/;
           rnd[tcnt] = rnd[id];
25
           tr[tent][0] = tr[id][0];
26
           tr[tcnt][1] = tr[id][1];
27
           //val[tcnt] = val[id];
28
29
           return tent;
30
       }*/
31
       void push_up(int v) {
32
           int 1 = tr[v][0], r = tr[v][1];
```

```
33
            sz[v] = sz[1] + 1 + sz[r];
34
        }
35
        void push_down(int v) {
36
            if(!v) return;
37
            int l = tr[v][0], r = tr[v][1];
38
            //if(l) ;
39
            //if(r)
40
            //swap时候注意交换tr[v][0]和tr[v][1],而不是l和r;
41
        /* //可持久化在push_down是要新建节点,否者历史版本有可能被多次下传,以rev为例
42
        void push_down(int v) {
43
            if(!v \mid / | !rev[v]) return;
44
            int \ \mathcal{E}l = tr[v][0], \ \mathcal{E}r = tr[v][1];
45
46
            if(l) {
47
                l = copynode(l);
                rev[l]^{=1};
48
                swap(tr[l][0], tr[l][1]);
49
50
            };
            if(r) {
51
52
                r = copynode(r);
53
                rev[r]^=1;
                swap(tr[r][0], tr[r][1]);
54
            };
55
56
            rev[v] = 0;
57
        }*/
58
        void split(int v,int k,int &x,int &y) {
59
            if(!v) {x=y=0;return;}
60
            push_down(v);
            //v = copynode(v); //可持久化时复制节点
61
62
            /*if(k > sz[tr[v][0]]) {
63
                split(tr[v]/1], k-sz[tr[v]/0]/-1, tr[v]/1], y);
64
            }else{
65
66
67
                split(tr[v][0], k, x, tr[v][0]);
68
69
            if(val[v] <= k) {
70
                x = v:
71
                split(tr[v][1], k, tr[v][1], y);
72
            }else{
73
74
                split(tr[v][0], k, x, tr[v][0]);
75
            }
76
            push_up(v);
77
        int merge(int x, int y) \{//x堆所有值均小于y堆
78
79
            if(!x || !y) return x|y;
80
            push_down(x); push_down(y);
81
            if(rnd[x]<rnd[y]){</pre>
                //x = copynode(x); //可持久化时复制节点,可不写
82
83
                tr[x][1] = merge(tr[x][1],y);
84
                push_up(x);
85
                return x;
            }else{
86
                //y = copynode(y); //可持久化时复制节点,可不写
87
88
                tr[y][0] = merge(x,tr[y][0]);
89
                push_up(y);
90
                return y;
91
            }
92
93
        void insert(int k) {
```

```
94
             int x,y;
95
             split(root,k,x,y);
96
             root = merge(merge(x,newnode(k)),y);
97
         void recycle(int v) {//回收一颗 treap上所有节点
98
99
             if(!v) return;
100
             q.push(v);
101
             recycle(tr[v][0]); recycle(tr[v][1]);
102
103
         void erase(int k) {
104
             int x,y,z;
             split(root,k,x,y);
105
106
             split(x,k-1,x,z);
107
             z = merge(tr[z][0],tr[z][1]);
108
             root = merge(x,merge(z,y));
109
         void krank(int k) {
110
             int x,y;
111
             split(root,k-1,x,y);
112
113
             printf("%d\n",sz[x]+1);
114
             root = merge(x,y);
115
         int find(int v,int k) {
116
             if(sz[tr[v][0]]==k-1) return val[v];
117
118
             if(sz[tr[v][0]]>=k) return find(tr[v][0],k);
119
             return find(tr[v][1],k-sz[tr[v][0]]-1);
120
121
         void pre(int k) {
122
             int x,y;
123
             split(root,k-1,x,y);
             printf("%d\n",find(x,sz[x]));
124
125
             root=merge(x,y);
126
         void nxt(int k){
127
128
             int x,y;
129
             split(root,k,x,y);
             printf("%d\n",find(y,1));
130
131
             root=merge(x,y);
132
         }
133
         void reverse(int l,int r){
134
             int x,y,z;
135
             split(root, r, x, y);
136
             split(x, l-1, x, z);
             //rev[z] ^= 1;标记对本身无效,处理时将z点重新计算
137
             root = merge(merge(x,z),y);
138
139
140
         /*int getRank(int S) {
141
             int res = sz[tr[S][0]]+1;
             for(;fa[S];S = fa[S])
142
143
                 if(tr[fa[S]]/1] == S) res += sz[tr[fa[S]]/0] + 1;
144
            return res;
145
         }*/
146
         void display(int v) {
147
             if(!v) return;
             push_down(v);
148
149
             display(tr[v][0]);
             printf("%d ",val[v]);
150
             display(tr[v][1]);
151
152
         }
153
    }
```

154 //一种可持久化平衡树的替代(非强制在线),由历史版本向当前版本连边,在 dfs 遍历中利用权值树状数组,普通平衡树等获取答案

## 2.4 Persistent Data Structures

#### 2.4.1 Chairman Tree

```
struct Node {
2
        int 1, r;
3
        LL sum;
   }t[MAXN * 40];
4
   int cnt, n;
5
   int rt[MAXN];
7
   void update(int pre, int &x, int 1, int r, int v) {
        x = ++cnt; t[x] = t[pre]; t[x].sum++;
9
       if(1 == r) return;
10
       int mid = (1 + r) >> 1;
11
        if(v <= mid) update(t[pre].1, t[x].1, 1, mid, v);</pre>
12
       else update(t[pre].r, t[x].r, mid + 1, r, v);
13
   }
   int query(int x, int y, int 1, int r, int v) {
14
15
       if(1 == r) return 1;
16
        int mid = (1 + r) >> 1;
17
        int sum = t[t[y].1].sum - t[t[x].1].sum;
18
        if(sum >= v) return query(t[x].1, t[y].1, 1, mid, v);
19
        else return query(t[x].r, t[y].r, mid + 1, r, v - sum);
20
```

## 2.4.2 Unite Chairman Tree

```
//Q x到 y路 径 第 k大
2
   //L link(x, y)
3 #include <bits/stdc++.h>
4 using namespace std;
5 typedef int LL;
6 const int MAXN = 8e4 + 5;
7 const int S = 18;
   struct Node {
9
       int 1, r;
10
       LL sum;
11 }t[MAXN * 800]; //2 * log^2(n)
12 int n, m, Q;
13 int cnt;
   int rt[MAXN], sz[MAXN];
14
15
   void update(int pre, int &x, int 1, int r, int v) {
16
       x = ++cnt; t[x] = t[pre]; t[x].sum++;
       if(1 == r) return;
17
18
       int mid = (1 + r) >> 1;
19
       if(v <= mid) update(t[pre].1, t[x].1, 1, mid, v);</pre>
20
       else update(t[pre].r, t[x].r, mid + 1, r, v);
21
   int query(int x, int y, int z, int w, int l, int r, int v) {
22
       if(1 == r) return 1;
23
       int mid = (1 + r) >> 1;
24
25
       int sum = t[t[x].1].sum + t[t[y].1].sum - t[t[z].1].sum - t[t[w].1].sum;
26
       if(sum >= v) return query(t[x].1, t[y].1, t[z].1, t[w].1, 1, mid, v);
27
       return query(t[x].r, t[y].r, t[z].r, t[w].r, mid + 1, r, v - sum);
   }
```

```
int fa[MAXN][S + 3], dep[MAXN];
30 int val[MAXN];
31 vector<int> G[MAXN];
32 inline void addEdge(int x, int y) {
33
        G[x].push_back(y);
34
        G[y].push_back(x);
35
   }
   inline void upd(int x) {
36
37
        update(rt[fa[x][0]], rt[x], 1, n, val[x]);
38
        for(int i = 1; i <= S; i++) fa[x][i] = fa[fa[x][i - 1]][i - 1];
39
    inline void Go(int &x, int step) {
40
        for(int i = S; i \ge 0; i--) if(step >> i & 1) x = fa[x][i];
41
42
   }
43
   int lca(int x, int y) {
        if(dep[x] < dep[y]) swap(x, y);</pre>
44
        Go(x, dep[x] - dep[y]);
45
46
        if(x == y) return x;
47
        for(int i = S; i >= 0; i--) if(fa[x][i] != fa[y][i]) {
48
            x = fa[x][i], y = fa[y][i];
49
50
        return fa[x][0];
51
   }
52
   int get_rt(int x) {
        for(int i = S; i >= 0; i--) if(fa[x][i]) x = fa[x][i];
53
54
        return x;
55
   }
56
   void dfs(int x, int f) {
57
        fa[x][0] = f;
58
        dep[x] = dep[f] + 1;
59
        upd(x);
60
        for(auto y : G[x]) {
            if(y == f) continue;
61
62
            dfs(y, x);
        }
63
   }
64
   void unite(int x, int y) {
65
66
        int rx = get_rt(x), ry = get_rt(y);
        if(sz[rx] > sz[ry]) swap(x, y), swap(rx, ry);
67
68
        addEdge(x, y);
69
        dfs(x, y);
70
        sz[ry] += cnt - rt[x] + 1;
71
   }
   void init() {
72
        cnt = 0;
73
74
        for(int i = 1; i <= n; i++) {</pre>
            rt[i] = sz[i] = 0;
75
76
            G[i].clear();
77
            dep[i] = 0;
78
        }
79
80
    int main() {
        ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
81
        int T; cin >> T;
82
        while(T--) {
83
84
            cin >> n >> m >> Q;
85
            init();
86
            vector<int> b;
87
            map<int, int> mp;
            for(int i = 1; i <= n; i++) {</pre>
88
89
                cin >> val[i];
```

```
90
                  b.push_back(val[i]);
91
             }
92
             sort(b.begin(), b.end());
93
             b.erase(unique(b.begin(), b.end()), b.end());
94
             for(int i = 1, tmp; i <= n; i++) {</pre>
95
                  tmp = val[i];
96
                  val[i] = lower_bound(b.begin(), b.end(), val[i]) - b.begin() + 1;
97
                  mp[val[i]] = tmp;
98
             for(int i = 1, u, v; i <= m; i++) {
99
                  cin >> u >> v;
100
                  addEdge(u, v);
101
102
103
             for(int i = 1; i <= n; i++) if(!dep[i]) {</pre>
104
                  dep[i] = 1;
105
                  dfs(i, 0);
                  sz[i] = cnt - rt[i] + 1;
106
107
108
             char s[3]; int x, y, z, k, ans = 0;
109
             while(Q--) {
110
                  cin >> s >> x >> y;
                  x \sim ans; y \sim ans;
111
                  if(s[0] == 'Q') {
112
                      cin >> k; k ^= ans;
113
                      z = lca(x, y);
114
115
                      ans = query(rt[x], rt[y], rt[z], rt[fa[z][0]], 1, n, k);
116
                      ans = mp[ans];
117
                      cout << ans << endl;</pre>
118
                  }
119
                  else {
120
                      unite(x, y);
121
             }
122
123
124
         return 0;
125
```

## 2.4.3 Persistent Trie

```
//区间异或最值查询
2 const int N=5e4+10;
3 int t[N];
4 int ch[N*32][2],val[N*32];
5 int cnt;
6
   void init(){
7
       mem(ch,0);
8
        mem(val,0);
9
       cnt=1;
10
   }
11
   int add(int root,int x){
12
        int newroot=cnt++,ret=newroot;
13
        for(int i=30;i>=0;i--){
14
            ch[newroot][0]=ch[root][0];
15
            ch[newroot][1]=ch[root][1];
16
            int now=(x>>i)&1;
            root=ch[root][now];
17
18
19
            ch[newroot][now]=cnt++;
20
            newroot=ch[newroot][now];
21
            val[newroot]=val[root]+1;
```

```
22
23
24
        return ret;
25
   }
26
   int query(int lt,int rt,int x){
27
        int ans=0;
        for(int i=30;i>=0;i--){
28
29
             int now=(x>>i)&1;
30
             if(val[ch[rt][now^1]]-val[ch[lt][now^1]]){
31
                 ans = (1<<i);
                 rt=ch[rt][now^1];
32
33
                 lt=ch[lt][now^1];
34
                 } else{
35
                 rt=ch[rt][now];
36
                 lt=ch[lt][now];
37
            }
38
        }
39
        return ans;
40
```

## 2.4.4 SGT in BBST

```
1 #include <bits/stdc++.h>
2 using namespace std;
3
   const int MAXN = 1e5;
4
   const int MAXM = 2e7;
5
   const int LM = 0;
   const int RM = 70005;
6
7
    namespace T{
8
        int Tsz; queue<int> q;
9
        int ls[MAXM], rs[MAXM], val[MAXM];
10
        int Tan, Ta[105], Tbn, Tb[105];
11
        int newnode() {
12
            if(q.empty()) q.push(++Tsz);
13
            int x = q.front(); q.pop();
            ls[x] = rs[x] = val[x] = 0;
14
15
            return x;
16
        }
        void insert(int &x, int 1, int r, int k, int f) {
17
            if(!x) x = newnode();
18
19
            val[x] += f;
20
            if(1 == r) return;
21
            int mid = (1 + r) >> 1;
22
            if(k <= mid) {</pre>
23
                insert(ls[x], l, mid, k,f);
24
            }else{
25
                 insert(rs[x], mid+1, r, k,f);
26
            }
27
        }
28
        int query(int 1, int r, int y) {
29
            if(1 == r) return 1;
30
            int mid = (1 + r) >> 1;
31
            int sum = 0;
            for(int i = 0; i <= Tan; i++) sum -= val[ls[Ta[i]]];</pre>
32
            for(int i = 0; i <= Tbn; i++) sum += val[ls[Tb[i]]];</pre>
33
34
35
            if(y <= sum) {
36
                 for(int i = 0; i <= Tan; i++) Ta[i] = ls[Ta[i]];</pre>
37
                 for(int i = 0; i <= Tbn; i++) Tb[i] = ls[Tb[i]];</pre>
38
                 return query(1, mid, y);
```

```
39
40
             for(int i = 0; i <= Tan; i++) Ta[i] = rs[Ta[i]];</pre>
41
             for(int i = 0; i <= Tbn; i++) Tb[i] = rs[Tb[i]];</pre>
42
            return query(mid+1, r, y - sum);
43
44
        void recycle(int v) {
45
            if(!v) return;
46
            q.push(v);
47
             recycle(ls[v]);
48
             recycle(rs[v]);
49
    };
50
51
    namespace TT{
52
        int Tsn;
53
        int val[MAXN], Trt[MAXN];
54
        int sz[MAXN];
        int fa[MAXN], tr[MAXN][2];
55
56
57
        int root;
58
        double alp;
59
        int newnode(int x) {
60
61
            int tcnt = ++Tsn;
            val[tcnt] = x; Trt[tcnt] = 0;
62
63
            fa[tcnt] = tr[tcnt][0] = tr[tcnt][1] = 0;
64
             sz[tcnt] = 0;
65
            return tcnt;
66
67
        bool isBad(int v) {
            return (double(sz[ tr[v][0] ]) > double(sz[v]) * alp) ||
68
69
                 (double(sz[ tr[v][1] ]) > double(sz[v]) * alp);
70
71
        vector<int> vec;
        void rRecycle(int v) {
72
73
            if(tr[v][0]) rRecycle(tr[v][0]);
            vec.push_back(v);
74
            T::recycle(Trt[v]);
75
76
            if(tr[v][1]) rRecycle(tr[v][1]);
77
        }
78
        int rBuild(int 1, int r) {
79
             int mid = (l + r) \gg 1, v = vec[mid];
80
             Trt[v] = 0; sz[v] = r - l + 1;
             for(int i = 1; i <= mid; i++) {</pre>
81
                 T::insert(Trt[v], LM, RM, val[vec[i]], 1);
82
83
             }
            tr[v][0] = (1 \leftarrow mid-1) ? rBuild(1, mid - 1) : 0;
84
85
            if(tr[v][0]) fa[tr[v][0]] = v;
            tr[v][1] = (mid+1 \leftarrow r) ? rBuild(mid + 1, r) : 0;
86
87
             if(tr[v][1]) fa[tr[v][1]] = v;
             return v;
88
89
        void rebuild(int v) {
90
91
             if(isBad(v)) {
92
                 vec.clear();
93
                 int u = fa[v], lr = tr[u][1] == v;
94
                 rRecycle(v);
95
                 if(vec.size()) v = rBuild(0, vec.size() - 1); else v = 0;
96
                 if(u == 0) fa[root = v] = 0;
97
                 else{
98
                     tr[u][lr] = v;
99
                     if(v) fa[v] = u;
```

```
100
101
             }
102
         }
103
         int find(int x, int k) {
104
             int 1 = tr[x][0], r = tr[x][1];
105
             if(sz[1] + 1 == k) return x;
106
             if(sz[1] >= k) return find(l, k);
107
             return find(r, k - sz[l] - 1);
108
         void ins(int x, int y) {
109
             int v = find(root, x);
110
             int p = tr[v][1], q = v;
111
112
             if(p) {
113
                 for(;p; q = p, p = tr[p][0]);
114
                 fa[p = tr[q][0] = newnode(y)] = q;
115
             }else {
                 fa[p = tr[q][1] = newnode(y)] = q;
116
117
             }
             int fg = 0;
118
119
             T::insert(Trt[p], LM, RM, y, 1);
120
             sz[p] = 1;
             for(;fa[p]; p = fa[p]) {
121
122
                 if(tr[fa[p]][0] == p)
123
                     T::insert(Trt[fa[p]], LM, RM, y, 1);
                 sz[fa[p]]++;
124
125
                 if(isBad(fa[p])) fg = fa[p];
126
             }
127
             rebuild(fg);
128
129
         void upd(int x, int y) {
130
             int p = find(root, x+1);
131
             int ty = val[p]; val[p] = y;
             T::insert(Trt[p], LM, RM, ty, -1);
132
             T::insert(Trt[p], LM, RM, y, 1);
133
             for(;fa[p]; p = fa[p]) {
134
                 if(tr[fa[p]][0] == p) {
135
                     T::insert(Trt[fa[p]], LM, RM, ty, -1);
136
137
                     T::insert(Trt[fa[p]], LM, RM, y, 1);
138
                 }
139
             }
140
141
         int que(int x, int y,int z) {
142
             x = find(root, x);
             T::Ta[T::Tan = 0] = Trt[x];
143
144
             for(;fa[x]; x = fa[x])
145
                 if(tr[fa[x]][1] == x) {
146
                     T::Ta[++T::Tan] = Trt[fa[x]];
147
                 }
148
             y = find(root, y+1);
             T::Tb[T::Tbn = 0] = Trt[y];
149
150
             for(;fa[y]; y = fa[y]) {
151
                 if(tr[fa[y]][1] == y) {
152
                      T::Tb[++T::Tbn] = Trt[fa[y]];
153
154
155
             return T::query(LM, RM, z);
156
         void init(int n) {
157
158
             alp = 0.7;
             vec.clear();
159
160
             vec.push_back(newnode(RM));
```

```
for(int i = 1, a; i <= n; i++) {</pre>
161
162
                  scanf("%d", &a);
163
                  vec.push_back(newnode(a));
164
              }
165
             root = rBuild(0, vec.size() - 1);
166
             fa[root] = 0;
167
         void display(int v) {
168
169
              if(tr[v][0]) display(tr[v][0]);
              cerr<<val[v]<<" ";
170
171
              if(tr[v][1]) display(tr[v][1]);
172
173
    };
```

## 2.5 Tree Structures

## 2.5.1 dsu on tree

```
const int MAXN = 1e5 + 7;
1
2 vector<int> G[MAXN];
int bgison, dfs_clock, sz[MAXN], st[MAXN], bt[MAXN], et[MAXN];
   int fg[MAXN], col[MAXN];
4
   long long ans[MAXN];
5
6
   void dfs1(int u, int fa) {
7
        sz[u] = 1;
8
        st[bt[u] = ++dfs\_clock] = u;
9
        for(auto v : G[u])
        if(v != fa) {
10
11
            dfs1(v, u);
12
            sz[u] += sz[v];
13
        et[u] = dfs_clock;
14
   }
15
   int maxx = 0;
16
   void dfs2(int u, int fa, int keep) {
17
        int mx = -1, bigson = -1;
18
19
        for(auto &v : G[u])
20
        if(v != fa) {
21
            if(sz[v] > mx)
22
                mx = sz[v], bigson = v;
23
24
        for(auto &v : G[u])
25
        if(v != fa && v != bigson)
26
            dfs2(v,u,0);
        if(bigson != -1) {
27
28
            dfs2(bigson, u, 1);
29
            ans[u] = ans[bigson];
30
            for(int &v : G[u])
                if(v != fa && v != bigson)
31
32
                for(int i = bt[v]; i <= et[v]; i++) {
                    ++fg[col[st[i]]];
33
                    if(fg[ col[st[i]] ] > maxx) maxx=fg[col[st[i]]], ans[u] = 0;
34
                    if(fg[ col[st[i]] ] == maxx) ans[u] += col[st[i]];
35
36
                }
37
38
        ++fg[col[u]];
        if(fg[col[u]] > maxx) maxx = fg[col[u]], ans[u] = 0;
39
40
        if(fg[col[u]] == maxx) ans[u] += col[u];
41
        if(keep == 0) {
```

## 2.5.2 Vitural Tree

```
const int MAXN = ;
   const int MAXM = ;
3
   const LL INF = ;
   const int S = 19;
4
5 int ecnt, head[MAXN];
   struct Edge{int to, nxt; LL w;} e[MAXM], ve[MAXM];
   inline void addEdgeT(int x, int y, LL w) {
7
8
        e[++ecnt] = (Edge) \{y, head[x], w\}; head[x] = ecnt;
9
   }
   int dep[MAXN], dfn_time, dfn[MAXN], fa[MAXN][S+1];
10
   LL dis[MAXN][S+1];
12
   void dfs(int v,int _fa) {
13
        dfn[v] = ++dfn_time;
14
        dep[v] = dep[_fa] + 1;
15
        fa[v][0] = _fa;
16
        for(int i = 1; i <= S; i++) {
17
            fa[v][i] = fa[fa[v][i-1]][i-1];
18
            dis[v][i] = min(dis[v][i-1], dis[fa[v][i-1]][i-1]);
19
        for(int i = head[v], u; i; i = e[i].nxt)
20
        if((u = e[i].to) != _fa) {
21
22
            dis[u][0] = e[i].w;
23
            dfs(u, v);
24
25
   }
26
   int getLca(int u, int v) {
27
        if(dep[u] < dep[v]) swap(u, v);</pre>
28
        for(int i = S; i >= 0; i--)
29
            if(dep[fa[u][i]] >= dep[v]) u = fa[u][i];
30
        if(u == v) return u;
        for(int i = S; i >= 0; i--)
31
32
            if(fa[u][i] != fa[v][i])
33
                u = fa[u][i], v = fa[v][i];
34
        return fa[u][0];
35
   }
   LL getDis(int u, int v) {
36
37
        if(dep[u] < dep[v]) swap(u, v);</pre>
38
        LL res = INF;
39
        for(int i = S; i >= 0; i--)
40
            if(dep[fa[u][i]] >= dep[v]) {
41
                res = min(res, dis[u][i]);
42
                u = fa[u][i];
43
            }
44
        return res;
45
   namespace VituralTree{
46
47
        int hn, h[MAXN];
        int vecnt, vhead[MAXN];
48
49
        Edge ve[MAXM];
50
        int top, st[MAXN], cln, cl[MAXN];
51
        int fgn, fg[MAXN]; //利用 fg[i] == fgn判断是否为当前有效点
        void addEdgeVT(int x, int y, LL w) {
```

```
53
            ve[++vecnt] = (Edge) {y, vhead[x], w}; vhead[x] = vecnt;
54
55
        inline void link(int u, int v) {
56
            LL w = getDis(u, v);
            addEdgeVT(u, v, w);
57
            addEdgeVT(v, u, w);
58
59
60
        inline bool cmp(int a, int b) {return dfn[a] < dfn[b];}</pre>
61
        void build() {
62
            ++fgn;
            for(int i = 1; i <= hn; i++) fg[h[i]] = fgn;</pre>
63
            sort(h + 1, h + hn + 1, cmp);
64
            cl[cln = 1] = st[top = 1] = 1;
65
            for(int i = 1; i <= hn; i++) {</pre>
66
                int rem = getLca(st[top], h[i]);
67
                if(rem == st[top]) {
68
69
                     if(rem != st[top]) cl[++cln] = st[++top] = h[i];
70
                     continue;
71
72
                while(top > 1 && dep[st[top - 1]] >= dep[rem]) {
73
                     link(st[top - 1], st[top]); top--;
74
                if(dep[st[top]] > dep[rem]) {
75
76
                    link(rem, st[top]), top --;
77
78
                if(rem != st[top]) cl[++cln] = st[++top] = rem;
79
                if(h[i] != st[top]) cl[++cln] = st[++ top] = h[i];
80
81
            while(top > 1) {
82
                link(st[top - 1], st[top]), top --;
83
84
        void clear() {
85
            vecnt = 0;
86
            for(;cln; --cln) vhead[cl[cln]] = 0;
87
88
        void sol() {
89
            build();
90
            //注意一号节点可能在虚树外面
91
92
            clear();
93
        }
94
```

## 2.5.3 Tree Decomposition

```
int sz[MAXN], dep[MAXN], top[MAXN], fa[MAXN], son[MAXN], num[MAXN], totw;
1
2
   struct Edge {
3
       int to, nxt;
4
   }e[MAXN << 1];
   int head[MAXN], ecnt;
5
6
   int n, m, Q;
7
   #define Ls(x) (x << 1)
   #define Rs(x) (x << 1 | 1)
   struct Tree {
9
        int 1, r, lazy;
10
        LL sum, mx;
11
   }tree[MAXN << 2];</pre>
12
13 int A[MAXN], B[MAXN];
14 void push_up(int x) {
       tree[x].sum = tree[Ls(x)].sum + tree[Rs(x)].sum;
```

```
16
       tree[x].mx = max(tree[Ls(x)].mx, tree[Rs(x)].mx);
17
   }
18
   void push_down(int x) {
19
        if(tree[x].lazy) {
            tree[Ls(x)].sum += tree[x].lazy * (tree[Ls(x)].r - tree[Ls(x)].l + 1);
20
            tree[Rs(x)].sum += tree[x].lazy * (tree[Rs(x)].r - tree[Rs(x)].l + 1);
21
22
            tree[Ls(x)].mx += tree[x].lazy;
23
            tree[Rs(x)].mx += tree[x].lazy;
24
            tree[Ls(x)].lazy += tree[x].lazy;
25
            tree[Rs(x)].lazy += tree[x].lazy;
26
            tree[x].lazy = 0;
27
28
   }
29
   void build(int x, int L, int R) {
30
        tree[x].lazy = 0;
31
        tree[x].l = L; tree[x].r = R;
        if(L == R) {
32
            tree[x].sum = B[L];
33
34
            tree[x].mx = B[L];
35
            return;
36
        int mid = (L + R) \gg 1;
37
        build(Ls(x), L, mid);
38
39
        build(Rs(x), mid + 1, R);
40
        push_up(x);
41
42
   void update(int x, int L, int R, LL val) {
43
        if(tree[x].1 >= L && tree[x].r <= R) {</pre>
44
            tree(x).lazy += val;
45
            tree[x].sum += val * (tree[x].r - tree[x].l + 1);
46
            tree[x].mx += val;
47
            return;
        }
48
49
        push_down(x);
50
        int mid = (tree[x].l + tree[x].r) >> 1;
        if(L <= mid) update(Ls(x), L, R, val);</pre>
51
        if(R > mid) update(Rs(x), L, R, val);
52
53
        push_up(x);
54
   }
55
   LL query(int x, int L, int R) {
56
        if(tree[x].1 >= L \&\& tree[x].r <= R)
57
            return tree[x].sum;
58
        push_down(x);
59
        int mid = (tree[x].l + tree[x].r) >> 1;
60
        LL res = 0;
61
        if(L <= mid) res += query(Ls(x), L, R);</pre>
62
        if(R > mid) res += query(Rs(x), L, R);
63
        return res;
64
    LL query2(int x, int L, int R) {
65
66
        if(tree[x].1 >= L \&\& tree[x].r <= R)
67
            return tree[x].mx;
68
        push_down(x);
69
        int mid = (tree[x].l + tree[x].r) >> 1;
        LL res = -INF;
70
71
        if(L <= mid) res = max(res, query2(Ls(x), L, R));</pre>
        if(R > mid) res = max(res, query2(Rs(x), L, R));
72
73
        return res;
74
   }
   inline void add edge(int x, int y) {
75
        e[++ecnt] = (Edge) \{y, head[x]\}; head[x] = ecnt;
```

```
77
    }
    void dfs1(int x) {
78
79
         sz[x] = 1; son[x] = 0;
80
         for(int i = head[x]; i; i = e[i].nxt) {
81
             int v = e[i].to;
82
             if(v == fa[x]) continue;
             fa[v] = x;
83
             dep[v] = dep[x] + 1;
84
85
             dfs1(v);
86
             sz[x] += sz[v];
87
             if(sz[v] > sz[son[x]]) son[x] = v;
88
    }
89
    void dfs2(int x) {
90
91
         B[num[x]] = A[x];
92
         if(son[x]) {
93
             top[son[x]] = top[x];
             num[son[x]] = ++totw;
94
95
             dfs2(son[x]);
96
         for(int i = head[x]; i; i = e[i].nxt) {
97
98
             int v = e[i].to;
99
             if(v == fa[x] || v == son[x]) continue;
100
             top[v] = v;
101
             num[v] = ++totw;
102
             dfs2(v);
103
         }
104
    }
105
    void up(int a, int b, int c) {
106
         int f1 = top[a], f2 = top[b];
         while(f1 != f2) {
107
108
             if(dep[f1] < dep[f2]) { swap(a, b); swap(f1, f2); }</pre>
109
             update(1, num[f1], num[a], c);
110
             a = fa[f1];
             f1 = top[a];
111
112
         if(dep[a] > dep[b]) swap(a, b);
113
114
         update(1, num[a], num[b], c);
115
    }
116
    int qsum(int a, int b) {
117
         if(a == b) return query(1, num[a], num[a]);
118
         int f1 = top[a], f2 = top[b];
119
         int res = 0;
         while(f1 != f2) {
120
121
             if(dep[f1] < dep[f2]) { swap(a, b); swap(f1, f2); }</pre>
122
             res += query(1, num[f1], num[a]);
123
             a = fa[f1];
             f1 = top[a];
124
125
126
         if(dep[a] > dep[b]) swap(a, b);
127
         res += query(1, num[a], num[b]);
128
         return res;
129
    int qmax(int a, int b) {
130
         if(a == b) return query2(1, num[a], num[a]);
131
         int f1 = top[a], f2 = top[b];
132
         int res = -10000000000;
133
         while(f1 != f2) {
134
135
             if(dep[f1] < dep[f2]) { swap(a, b); swap(f1, f2); }</pre>
136
             res = max(res, query2(1, num[f1], num[a]));
137
             a = fa[f1];
```

```
138
            f1 = top[a];
139
         }
140
         if(dep[a] > dep[b]) swap(a, b);
141
         res = max(res, query2(1, num[a], num[b]));
142
         return res;
143
    }
     inline void init() {
144
145
         memset(head, 0, sizeof(head)); ecnt = 0;
146
         fa[1] = 0; dep[1] = 1; top[1] = 1; num[1] = 1; totw = 1;
147
148
     inline void pre() {
         dfs1(1); dfs2(1); build(1, 1, totw);
149
    }
150
151
     //-
152
    const int MAXN = ;
153 const int INF = ;
int A[MAXN], B[MAXN], C[MAXN];
155 struct SGT{
         int sL, sR;
156
157
         LL sW, mx[MAXN << 2];
158
         void push_down(int v) {
159
160
         void push_up(int v) {
161
             mx[v] = max(mx[v << 1], mx[v << 1]);
162
163
         void build(int v, int l, int r) {
             if(1 == r) {
164
165
                  mx[v] = B[1];
166
                  return;
167
              }
              int mid = (1 + r) >> 1;
168
169
              build(v << 1, 1, mid);
170
              build(v << 1|1, mid+1,r);
171
              push_up(v);
172
         void upd(int v, int l, int r) {
173
              if(sL <= 1 && r <= sR) {</pre>
174
                  //sum[v] = sW *(r - l + 1);
175
                  mx[v] = sW;
176
177
                  return;
178
              }
179
              push_down(v);
180
              int mid = (1 + r) >> 1;
181
             if(sL <= mid) upd(v << 1, 1, mid);</pre>
              if(mid \langle sR \rangle upd(v \langle \langle 1|1, mid+1, r \rangle;
182
183
              push_up(v);
184
185
         void qmax(int v, int l, int r) {
              if(sL <= 1 && r <= sR) {
186
187
                  sW = max(sW, mx[v]);
188
                  return;
189
190
              push_down(v);
              int mid = (1 + r) >> 1;
191
192
             if(sL <= mid) qmax(v << 1, 1, mid);</pre>
             if(mid < sR) qmax(v<<1|1, mid+1, r);</pre>
193
194
    }T;
195
196
    namespace TD{
         struct Edge {int to, nxt;}e[MAXN << 1];</pre>
197
198
         int ecnt, head[MAXN];
```

```
199
         inline void addEdge(int x, int y) {
200
             e[++ecnt] = (Edge) {y, head[x]}; head[x] = ecnt;
201
             e[++ecnt] = (Edge) \{x, head[y]\}; head[y] = ecnt;
202
203
         int fa[MAXN], sz[MAXN], dep[MAXN], son[MAXN], top[MAXN];
204
         int dfn_time, dfn[MAXN], rnk[MAXN];
205
         void dfs1(int u) {
206
             sz[u] = 1; son[u] = -1;
207
             for(int i = head[u], v; i; i = e[i].nxt) {
208
                  v = e[i].to;
209
                  if(v == fa[u]) continue;
                  fa[v] = u;
210
211
                  dep[v] = dep[u] + 1;
212
                  dfs1(v);
213
                  sz[u] += sz[v];
214
                  if(son[u] == -1 \mid \mid sz[v] > sz[son[u]]) son[u] = v;
215
             }
216
217
         void dfs2(int u) {
218
             dfn[u] = ++dfn_time;
219
              //rnk [dfn\_time] = u;
220
             B[dfn[u]] = A[u];
221
             if(son[u] == -1) return;
222
             top[son[u]] = top[u];
223
             dfs2(son[u]);
224
             for(int i = head[u], v; i; i = e[i].nxt) {
225
                  v = e[i].to;
226
                  if(v == fa[u] || v == son[u]) continue;
227
                  top[v] = v;
228
                  dfs2(v);
             }
229
230
         inline void init(int n) {
231
232
             ecnt = 1;
233
             for(int i = 0; i <= n; i++) head[i] = 0;</pre>
234
235
         inline void pre(int rt) {
236
             fa[rt] = -1; dep[rt] = 1;
237
             top[rt] = rt; dfn time = 0;
238
             dfs1(rt); dfs2(rt);
239
             // edge \rightarrow point
             //for(int \ i = 2; \ i <= ecnt; \ i += 2)  {
240
241
                    if(e[i]. to == fa[e[i^1]. to])  {
                        C[i / 2] = e[i^1]. to;
242
                    }else{
243
                        C[i / 2] = e[i].to;
244
245
246
                    B[dfn[C[i / 2]]] = e[i].w;
247
248
             T.build(1, 1, dfn time);
249
250
         int lca(int u, int v) {
251
             while(top[u] != top[v]) {
252
                  if(dep[top[u]] > dep[top[v]]) swap(u, v);
253
                  v = fa[top[v]];
254
255
             if(dep[u] > dep[v]) swap(u, v);
256
             return u;
257
258
         void upd(int a, int b, int c) {
259
             int ta = top[a], tb = top[b];
```

```
260
             while(ta != tb) {
261
                 if(dep[ta] < dep[tb]) { swap(a, b); swap(ta, tb); }</pre>
262
                 T.sL = dfn[ta]; T.sR = dfn[a]; T.sW = c;
263
                 T.upd(1, 1, dfn_time);
264
                 a = fa[ta]; ta = top[a];
265
             }
             if(dep[a] > dep[b]) swap(a, b);
266
                                                              // point
                                                                ^{\prime} edge
             //if(a == b) return;
267
268
             //if(dep[a] > dep[b]) swap(a, b);
                                                              // edge
                                                              // edge
269
             //a = son[a];
             T.sL = dfn[a]; T.sR = dfn[b]; T.sW = c;
270
             T.upd(1, 1, dfn_time);
271
272
273
         //更新子树,由于dfs2中有if(son[u] == -1) return;小心使用end[u];
274
         void upd2(int a, int c) {
             T.sL = dfn[a]; T.sR = dfn[a] + sz[a] - 1; T.sW = c;
275
276
             T.upd(1, 1, dfn_time);
277
278
         int qmax(int a, int b) {
279
             int ta = top[a], tb = top[b]; int res = - INF;
280
             while(ta != tb) {
281
                 if(dep[ta] < dep[tb]) { swap(a, b); swap(ta, tb); }</pre>
282
                 T.sL = dfn[ta]; T.sR = dfn[a]; T.sW = - INF;
283
                 T.qmax(1, 1, dfn_time); res = max(T.sW, res);
                 a = fa[ta];ta = top[a];
284
285
                                                              // point
286
             if(dep[a] > dep[b]) swap(a, b);
                                                              // edge
287
             //if(a == b) return res;
             //if(dep[a] > dep[b]) swap(a, b);
288
                                                              // edge
289
             //a = son[a];
                                                              // edge
             T.sL = dfn[a]; T.sR = dfn[b]; T.sW = -INF;
290
291
             T.qmax(1, 1, dfn_time); res = max(T.sW, res);
292
             return res;
293
         }
294
```

## 2.5.4 Link-Cut Tree

```
1
   namespace LCT {
2
        int tcnt; //动态开点
3
        int fa[MAXN], tr[MAXN][2], rev[MAXN];
4
        //int val[MAXN], sval[MAXN];
        //void clear(int id) { //或者改为int newnode() {}
5
6
             fa[id] = tr[id][0] = tr[id][1] = 0;
7
             rev[id] = 0; //sz[id] = 1;
        //}
8
        void Rev(int x) {
9
10
           rev[x] ^= 1; swap(tr[x][0], tr[x][1]);
11
12
        void push_up(int x) {
13
            //int l = tr[x][0], r = tr[x][1];
            //sval[x] = sval[l] + val[x] + sval[r];
14
15
        void push_down(int x) {
16
           int 1 = tr[x][0], r = tr[x][1];
17
18
            if(rev[x]) {
19
                if(1) Rev(1);
20
                if(r) Rev(r);
21
                rev[x] = 0;
22
            }
```

```
23
24
        bool isroot(int x) {
25
            return tr[fa[x]][0] != x && tr[fa[x]][1] != x;
26
27
        void pre(int x) {
28
            if(!isroot(x)) pre(fa[x]);
29
            push_down(x);
30
31
        void rotate(int x) {
32
            int y = fa[x], z = fa[y], lr = tr[y][1] == x;
            if(!isroot(y)) tr[z][tr[z][1] == y] = x;
33
            fa[x] = z;
34
35
            fa[tr[y][lr] = tr[x][lr^1]] = y;
36
            fa[tr[x][lr^1] = y] = x;
37
            push_up(y);
38
        inline void splay(int x) {
39
40
            pre(x);
41
            for (int y, z; !isroot(x); rotate(x)) {
42
                y = fa[x]; z = fa[y];
                if (!isroot(y)) rotate((tr[z][0] == y) ^ (tr[y][0] == x) ? x : y);
43
44
            push_up(x);
45
46
        inline int access(int x) {
47
48
            int y = 0;
49
            for (; x; y = x, x = fa[x]) {
50
                splay(x);
                //sz2[x] += sz[tr[x][1]] - sz[y]; //subtree
51
                tr[x][1] = y;
52
53
                push_up(x);
54
            }
            return y; //不求LCA不必
55
56
        inline void makeroot(int x) {
57
58
            access(x); splay(x); Rev(x);
59
        }
        inline int findroot(int x) {
60
61
            access(x); splay(x);
62
            for(;tr[x][0]; x = tr[x][0]) push_down(x);
63
            splay(x);
64
            return x;
65
        inline void lnk(int x, int y) {
66
67
            makeroot(x);
68
            if(findroot(y) != x) fa[x] = y;
69
            //sz2[y] += sz[x]; //subtree
70
        inline void cut(int x,int y) {
71
72
            makeroot(x); //access(y); splay(x);
73
            if(findroot(y) == x&&fa[y] == x&&!tr[y][0]){
74
                fa[y] = tr[x][1] = 0;
75
                push_up(x);
            }
76
77
        inline void cut(int y) { //有根树断开与父节点连边
78
79
            access(y); splay(y);
80
            fa[tr[y][0]] = 0;
            tr[y][0] = 0;
81
82
            push_up(y);
83
```

```
inline int lca(int u, int v) {
84
85
           access(u);
86
           return access(v);
87
       void split(int x, int y) {
88
89
           makeroot(x); access(y); splay(y);
90
       //维护节点或者维护路径
91
       //例如:染色(注意tag\_rev)、tag\_add的区间信息、splay维护连续端最远位置
92
       void upd(int x, int y) {
93
           makeroot(x); val[x] = y; push_up(x);
94
95
       int que(int x, int y) {
96
97
           split(x, y);
98
           //return sval[y];
99
       //维护边权(y \rightarrow eid \rightarrow x),
                                  需要初始化 vcnt, ecnt, 可用map(注意双向维护)维护eid信息
100
101
       //注意 if(x == y) continue;
       struct LCTEdge{int u, v; int w;} e[MAXN];
102
103
       void addEdge(int eid) { //e[eid = ++ecnt] = (Edge)\{x, y, w\};
104
           lnk(e[eid].u, vcnt + eid);
105
           lnk(vcnt + eid, e[eid].v);
106
       void delEdge(int eid) {
107
108
           cut(e[eid].u, vcnt + eid);
           cut(vcnt + eid, e[eid].v);
109
110
       }
       //维护边双连通分量
111
112
       //并查集,所有的fa[x]改为Find(fa[x]),public调用函数使用前Find(x),需要保证只在shink进
       行过合并
       //维护边双时, 节点自身信息在 unite 中维护, 路径等信息在 push_up和 push_down中维护
113
114
       int Rt[MAXN];
       int findroot(int x) {return Rt[x] == x ? x : Rt[x] = findroot(Rt[x]);}
115
116
       void unite_dfs(int x) {
117
           push_down(x);
           if (tr[x][0]) unite_dfs(tr[x][0]), unite(tr[x][0], x);
118
           if (tr[x][1]) unite_dfs(tr[x][1]), unite(tr[x][1], x);
119
120
       void shrink(int x, int y) {
121
122
           split(x, y);
123
           unite_dfs(y);
124
           int z = Find(y);
125
           fa[z] = fa[y]; tr[z][0] = tr[z][1] = 0;
126
           push_up(z);
127
128
       void addEdge(int x, int y) {
129
           x = Find(x); y = Find(y);
           if(findroot(x) != findroot(y)) {
130
131
               lnk(x, y); Rt[findroot(x)] = Rt[findroot(y)];
132
           }else shrink(x, y);
133
       }
134
    };
    //未连成树、初始化节点时,需要调用LCT::push_up 维护节点其余信息
135
    //维护的信息要有 可减性 , 如子树结点数, 子树权值和, 但不能直接维护子树最大最小值, 因为在
136
       将一条虚边变成实边时要排除原先虚边的贡献。
    //新建一个附加值存储虚子树的贡献, 在统计时将其加入本结点答案, 在改变边的虚实时及时维护。
137
    //其余部分同普通 LCT, 在统计子树信息时一定将其作为根节点。
138
    //如果维护的信息没有可减性,如维护区间最值,可以对每个结点开一个平衡树维护结点的虚子树中
139
       的最值。
   if(x == y) continue;
140
   if(Find(x) != Find(y)) {
```

```
142
         unite(x, y);
143
         addEdge(mp[{x, y}]);
144
    }else{
         int eid = que(x, y);
145
         int id = mp[{x, y}];
146
         if(val[eid] > e[id].w) {
147
148
             delEdge(eid - vcnt);
149
             addEdge(id);
150
         }
151
    }
```

#### 2.5.5 Divide Combine Tree

1.[i,i+1] 构造依赖, 利用线段树辅助建图,利用 tarjan 求 scc,利用 rmq 求最左最右边界 2. 定义 (i,j) 为一个好二元组,当且仅当 a[i]-a[j]=1 这样的两项的二元组在 [l,r] 中恰好有 r-l 个线段树维护 val+l=r,其中 val 是区间 [l,r] 中好二元组的个数离线 3. 析合树维护 mx-mn=r-l <=> fx = (mx-mn) - (r-l)

```
namespace DCT{
1
2
        struct RMQ {
3
            int lg[MAXN], mn[MAXN][S+1], mx[MAXN][S+1];
4
            inline void init(int *a, int n) {
5
                for (int i = 2; i \le n; i++) lg[i] = lg[i >> 1] + 1;
6
                for (int i = 1; i <= n; i++) mn[i][0] = mx[i][0] = a[i];
7
                for (int k = 1; (1 << k) <= n; k++)
8
                     for (int i = 1; i + (1 << k) - 1 <= n; i++) {
9
                         mn[i][k] = min(mn[i][k - 1], mn[i + (1 << (k - 1))][k - 1]);
10
                         mx[i][k] = max(mx[i][k - 1], mx[i + (1 << (k - 1))][k - 1]);
11
                     }
12
13
            inline int Min(int 1, int r) {
14
                int len = lg[r - l + 1];
15
                return min(mn[l][len], mn[r - (1 << len) + 1][len]);</pre>
16
17
            inline int Max(int 1, int r) {
                int len = lg[r - l + 1];
18
                return max(mx[1][len], mx[r - (1 << len) + 1][len]);</pre>
19
20
        } D;
21
22
        struct SEG {
23
24
            int setL, setR, setW;
25
            int mn[MAXN << 2], tag[MAXN << 2];</pre>
26
            inline void pushup(int x) {
27
28
                mn[x] = min(mn[x << 1], mn[x << 1 | 1]);
29
            }
30
            inline void pushdown(int x) {
                if(!tag[x]) return;
31
32
                mn[x << 1] += tag[x]; mn[x << 1 | 1] += tag[x];
                tag[x << 1] += tag[x]; tag[x << 1 | 1] += tag[x]; tag[x] = 0;
33
34
35
            void init(int x, int 1, int r) {
36
                mn[x] = tag[x] = 0;
37
                if (1 == r) return;
                int mid = (1 + r) >> 1;
38
                init(x << 1, 1, mid);</pre>
39
                init(x << 1 | 1, mid + 1, r);
40
41
42
            void upt(int x, int 1, int r) {
43
                if (setL <= 1 && r <= setR) {</pre>
```

```
tag[x] += setW; mn[x] += setW;
44
45
                      return;
46
                  }
47
                  pushdown(x);
48
                  int mid = (1 + r) >> 1;
49
                  if (setL <= mid) upt(x << 1, 1, mid);</pre>
50
                  if (mid < setR ) upt(x << 1 | 1, mid+1, r);</pre>
51
                  pushup(x);
52
             int que(int x, int 1, int r) {
53
                  if (1 == r) return 1;
54
55
                  pushdown(x);
                  int mid = (l+r)>>1;
56
57
                  if (!mn[x << 1]) return que(x << 1, 1, mid);</pre>
58
                  return que(x << 1 | 1, mid+1, r);</pre>
59
         } T;
60
61
         int tpmn, stmn[MAXN], tpmx, stmx[MAXN], tpk, stk[MAXN];
62
63
         int ncnt, type[MAXN<<1], L[MAXN<<1], R[MAXN<<1], M[MAXN<<1];</pre>
64
         int dep[MAXN<<1], fa[MAXN<<1][S+1], C[MAXN<<1];</pre>
65
         int id[MAXN << 1];</pre>
66
         int newnode(int _type, int _L, int _R, int _M = 0) {
             ++ncnt; type[ncnt] = _type;
67
68
             L[ncnt] = _L; R[ncnt] = _R; M[ncnt] = _M;
69
             C[ncnt] = 0;
70
             return ncnt;
71
         }
72
73
         inline bool judge(int 1, int r) {
74
             return D.Max(1, r) - D.Min(1, r) == r - 1;
75
76
         int ecnt, head[MAXN << 1];</pre>
77
         struct Edge{int to, nxt;} e[MAXN<<1];</pre>
78
79
         inline void addEdge(int x, int y) {
80
             e[++ecnt] = (Edge) \{y, head[x]\}; head[x] = ecnt;
81
             fa[y][0] = x; C[x]++;
82
         }
83
         void dfs(int u) {
84
             for(int j = 0; j < S; j++) fa[u][j+1] = fa[fa[u][j]][j];</pre>
85
             for(int i = head[u]; i; i = e[i].nxt) {
86
                  dep[e[i].to] = dep[u] + 1;
87
                  dfs(e[i].to);
88
             }
89
         }
90
         inline void init(int n) {
91
92
             ecnt = 0;
93
             for(int i = 0; i <= n; i++) head[i] = 0;</pre>
94
         void buildT(int *a, int n) {
95
96
             init(n);
97
             D.init(a, n);
98
             T.init(1, 1, n);
             tpmn = tpmx = tpk = 0;
99
             stmn[0] = stmx[0] = stk[0] = 0;
100
             for (int i = 1; i <= n; i++) {
101
                  for (;tpmn && a[i] <= a[stmn[tpmn]]; --tpmn) {</pre>
102
103
                      T.setL = stmn[tpmn - 1] + 1; T.setR = stmn[tpmn]; T.setW = a[stmn[tpmn
         ]];
```

```
104
                     T.upt(1, 1, n);
105
106
                 T.setL = stmn[tpmn] + 1; T.setR = i; T.setW = -a[i];
107
                 T.upt(1, 1, n);
108
                 stmn[++tpmn] = i;
109
110
                 for (;tpmx && a[i] >= a[stmx[tpmx]]; --tpmx) {
111
                      T.setL = stmx[tpmx - 1] + 1; T.setR = stmx[tpmx]; T.setW = -a[stmx[tpmx
        ]];
112
                     T.upt(1, 1, n);
113
                 T.setL = stmx[tpmx] + 1; T.setR = i; T.setW = a[i];
114
                 T.upt(1, 1, n);
115
116
                 stmx[++tpmx] = i;
117
118
                 int Li = T.que(1, 1, n), np = id[i] = newnode(0, i, i), nq, nw;
                 while (tpk && L[nq = stk[tpk]] >= Li) {
119
120
                     if (type[nq] && judge(M[nq], i)) {
121
                          R[nq] = i;
122
                          addEdge(nq, np);
123
                          np = nq; tpk--;
                      } else if (judge(L[nq], i)) {
124
125
                          nw = newnode(1, L[nq], i, L[np]);
                          addEdge(nw, nq); addEdge(nw, np);
126
127
                          np = nw; tpk--;
128
                     } else {
129
                          nw = newnode(0, -1, i);
130
                          addEdge(nw, np);
131
                          do {
132
                              addEdge(nw, nq);
133
                              nq = stk[--tpk];
                          } while (tpk && !judge(L[nq], i));
134
                          addEdge(nw, nq);
135
                          L[nw] = L[nq]; R[nw] = i;
136
                          np = nw; --tpk;
137
                     }
138
139
                 }
                 stk[++tpk] = np;
140
                 T.setL = 1; T.setR = i; T.setW = -1;
141
142
                 T.upt(1, 1, n);
143
             }
144
             assert(tpk == 1);
145
             dfs(stk[tpk]);
146
         void lca(int u, int v, int &aL, int &bR) {
147
             if(u == v) {
148
149
                 aL = L[u]; bR = R[v];
150
                 return;
151
             if(dep[u] > dep[v]) swap(u, v);
152
153
             for(int i = S; i >= 0; i--)
154
                 if(dep[fa[v][i]] >= dep[u]) v = fa[v][i];
155
             assert(u != v);
             for(int i = S; i >= 0; i--)
156
                 if(fa[u][i] != fa[v][i]) {
157
                     u = fa[u][i]; v = fa[v][i];
158
159
160
             if(type[fa[u][0]]) {
161
                 aL = min(L[v], L[u]);
162
                 bR = max(R[v], R[u]);
163
             }else{
```

# 3 String

## 3.1 Basics

#### 3.1.1 Hash

```
const LL p1 = 201, p2 = 301, mod1 = 1200000319, mod2 = 2147483647;
2
   struct Hash {
3
        LL a, b;
4
        void append(Hash pre, int v) {
5
            a = (pre.a * p1 + v) % mod1;
            b = (pre.b * p2 + v) \% mod2;
6
7
8
        void init(string S) {
9
            a = b = 0;
10
            for(int i = 0; i < S.size(); i++) append(*this, S[i]);</pre>
11
12
        bool operator == (const Hash &x) const {
13
           return a == x.a && b == x.b;
14
15
        bool operator < (const Hash &x) const {</pre>
16
           return a < x.a || (a == x.a && b < x.b);
17
18
   };
```

## 3.1.2 Minimum String

```
namespace minstring{
        int getmin(char *s, int sn) {
2
            int i = 0, j = 1, k = 0, t;
3
4
            while(i < sn && j < sn && k < sn) {</pre>
                 t = s[(i + k) \% sn] - s[(j + k) \% sn];
5
6
                 if(!t) k++;
7
                 else {
8
                     if(t > 0) i += k + 1; else j += k + 1;
9
                     if(i == j) j++;
10
                     k = 0;
11
12
            }
            return i < j ? i : j;</pre>
13
14
        }
15
   }
```

## 3.2 String Matching

## 3.2.1 Bitset Match

```
namespace BitMatch{
1
2
       const int S = 26;
       bitset<MAXN> bs[S], ret;
3
4
       char s[MAXN];
5
       inline int idx(char c) { return c - 'a'; }
6
       inline void init() {
7
           for(int i = 0; i < 26; i++) bs[i].reset();</pre>
8
           scanf("%s", s);
9
          int sn = strlen(s);
```

```
10
           for(int i = 0; i < sn; i++) bs[idx(s[i])].set(i);</pre>
11
        }
12
        void modify(int p, char ch) {
13
            bs[idx(s[p])].reset(p);
            s[p] = ch;
14
15
            bs[idx(s[p])].set(p);
16
17
        int match(char *t, int tn) \{//返回t串在s串中出现的次数
18
            ret = bs[idx(t[0])];
19
            for(int i = 1; i < tn; i++) {</pre>
20
                 ret <<= 1;
21
                ret &= bs[idx(t[i])];
22
23
            return ret.count();
24
25
   }
```

## 3.2.2 KMP && exKMP

```
1
   判断循环子串的充要条件: i/(i-fa[i]) > 1 && i%(i-fa[i])==0 (i是长度)且去除下面fa的优化
2
   namespace KMP {
3
        int fa[MAXN];
4
        void get_fail(char* t, int tn) {
5
            fa[0] = -1;
6
            int i = 0, j = -1;
            while(i < tn) {</pre>
7
8
                if (j == -1 || t[i] == t[j]) {
9
                    ++i; ++j;
10
                    fa[i] = t[i] != t[j] ? j : fa[j];
11
                }else{
12
                    j = fa[j];
13
14
            }
15
16
        void kmp(char* s, int sn, char* t, int tn) {
            int i = 0, j = 0;
17
            while(i < sn) {</pre>
18
                if (j == -1 || s[i] == t[j]) {
19
                    i++;j++;
20
21
                    if(j == tn) {
22
23
                }else j = fa[j];
            }
24
25
        }
26
   }
27
   namespace exKMP {
28
        int nxt[MAXN], ext[MAXN];
29
        void get_nxt(char* t, int tn) {
30
            int j = 0, mx = 0;
31
            nxt[0] = tn;
32
            for(int i = 1; i < tn; i++) {</pre>
33
                if(i >= mx || i + nxt[i - j] >= mx) {
34
                    if(i > mx) mx = i;
                    while(mx < tn && t[mx] == t[mx - i]) mx++;
35
36
                    nxt[i] = mx - i;
37
                    j = i;
38
                }else nxt[i] = nxt[i - j];
39
            }
40
        }
41
        void exkmp(char *s, int sn, char *t, int tn) {
```

```
42
            int j = 0, mx = 0;
43
             for(int i = 0; i < sn; i++) {</pre>
44
                 if(i >= mx \mid | i + nxt[i - j] >= mx) {
45
                      if(i > mx) mx = i;
46
                      while(mx < sn && mx - i < tn && s[mx] == t[mx - i]) mx++;
47
                      ext[i] = mx - i;
48
                      j = i;
                 }else ext[i] = nxt[i - j];
49
50
            }
51
        }
52
    }
```

## 3.2.3 AC Automaton

```
namespace AC {
1
2
        int ch[MAXN][sigma_size], last[MAXN];
3
        int val[MAXN], f[MAXN], sz;
        inline void init() { sz = 1; memset(ch[0], 0, sizeof(ch[0])); }
4
5
        inline int idx(char c) { return c - 'a'; }
6
        void insert(string s, int v) {
7
            int u = 0;
8
            for(int i = 0; i < s.size(); i++) {</pre>
9
                 int c = idx(s[i]);
10
                 if(!ch[u][c]) {
11
                     memset(ch[sz], 0, sizeof(ch[sz]));
12
                     val[sz] = 0;
13
                     ch[u][c] = sz++;
14
15
                u = ch[u][c];
16
17
            val[u] = v;
18
        void get_fail() {
19
20
            queue<int> q;
21
            f[0] = 0;
            for(int c = 0; c < sigma_size; c++) {</pre>
22
                 int u = ch[0][c];
23
24
                 if(u) { f[u] = 0; q.push(u); last[u] = 0; }
25
26
            while(!q.empty()) {
27
                 int r = q.front(); q.pop();
28
                 for(int c = 0; c < sigma_size; c++) {</pre>
29
                     int u = ch[r][c];
                     if(!u) { ch[r][c] = ch[f[r]][c]; continue; }
30
31
                     q.push(u);
32
                     int v = f[r];
33
                     while(v \& !ch[v][c]) v = f[v];
34
                     f[u] = ch[v][c];
35
                     last[u] = val[f[u]] ? f[u] : last[f[u]];
36
                 }
37
            }
38
39
        inline void solve(int j) {
40
            if(j) {
                 ans += val[j];
41
                 solve(last[j]);
42
43
            }
44
45
        void find(string T) {
46
            int j = 0;
```

```
for(int i = 0; i < T.size(); i++) {</pre>
47
48
                  int c = idx(T[i]);
49
                  j = ch[j][c];
50
                  if(val[j]) solve(j);
51
                  else if(last[j]) solve(last[j]);
52
             }
53
         }
54
55
    namespace AC {
56
         int root, tcnt;
57
         int ch[MAXN][sigma_size], fa[MAXN];
         inline int newnode() {
58
59
             fa[++tcnt] = 0;
60
             for(int i = 0; i < sigma_size; ++i) ch[tcnt][i] = 0;</pre>
61
             return tcnt;
62
         inline void init() {
63
64
             tcnt = -1;
65
             root = newnode();
66
         inline int idx(char c) { return c - 'a'; }
67
         void extend(char *s, int sn) {
68
69
             int cur = root;
70
             for(int i = 0, c; i < sn; i++) {</pre>
71
                  if(!ch[cur][c = idx(s[i])])
72
                      ch[cur][c] = newnode();
73
                  cur = ch[cur][c];
74
             }
75
76
         int q[MAXN], qh, qt;
77
         void get_fail() {
78
             qh = 1; qt = 0;
79
             fa[root] = 0;
             for(int c = 0, now; c < sigma_size; c++)</pre>
80
                  if((now = ch[root][c]) != 0)
81
82
                      q[++qt] = now;
83
             while(qh <= qt) {</pre>
                  int cur = q[qh++];
84
85
                  for(int c = 0, now; c < sigma_size; c++)</pre>
86
                      if((now = ch[cur][c]) != 0) {
87
                          fa[now] = ch[fa[cur]][c];
88
                          q[++qt] = now;
89
                      }else
90
                          ch[cur][c] = ch[fa[cur]][c];
91
             }
92
     //统计模板串出现次数,每个模板串只计算一次
93
94
             int \ cur = root, \ ans = 0;
             for(int \ i = 0; \ i < sn; ++i)  {
95
96
                  cur = ch[cur][idx(s[i])];
97
                  for(int j = cur; j \& cnt[j] != -1; j = fa[j])  {
98
                      ans \neq cnt[j];
                      cnt/j/ = -1;
99
100
             }
101
102
103
    }
```

## 3.3 Suffix Related

## 3.3.1 Suffix Array

```
1
   namespace SA {
2
        char s[MAXN];
        int sa[MAXN], rank[MAXN], height[MAXN];
3
4
        int t[MAXN], t2[MAXN], c[MAXN], n;
5
        void clear() { n = 0; memset(sa, 0, sizeof(sa)); }
6
        void build(int m) {
7
            int *x = t, *y = t2;
8
            for(int i = 0; i < m; i++) c[i] = 0;</pre>
9
            for(int i = 0; i < n; i++) c[x[i] = s[i]]++;
10
            for(int i = 1; i < m; i++) c[i] += c[i - 1];
11
            for(int i = n - 1; i >= 0; i--) sa[--c[x[i]]] = i;
12
            for(int k = 1; k <= n; k <<= 1) {
13
                int p = 0;
14
                for(int i = n - k; i < n; i++) y[p++] = i;
15
                for(int i = 0; i < n; i++) if(sa[i] >= k) y[p++] = sa[i] - k;
                 for(int i = 0; i < m; i++) c[i] = 0;
16
                for(int i = 0; i < n; i++) c[x[y[i]]]++;
17
                for(int i = 1; i < m; i++) c[i] += c[i - 1];</pre>
18
                for(int i = n - 1; i >= 0; i--) sa[--c[x[y[i]]]] = y[i];
19
20
                swap(x, y);
21
                p = 1; x[sa[0]] = 0;
                for(int i = 1; i < n; i++)</pre>
22
23
                     x[sa[i]] = y[sa[i - 1]] == y[sa[i]] && y[sa[i - 1] + k] == y[sa[i] + k]
        p - 1 : p + +;
24
                if(p >= n) break;
25
                m = p;
26
            }
27
        void buildHeight() {
28
29
            int k = 0;
30
            for(int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
            for(int i = 0; i < n; i++) {</pre>
31
32
                 if(k) k--;
                int j = sa[rank[i] - 1];
33
34
                while(s[i + k] == s[j + k]) k++;
35
                height[rank[i]] = k;
36
            }
37
        void init() {
38
39
            n = strlen(s) + 1;
            build('z' + 1);
40
41
            buildHeight();
42
43
    }
```

## 3.3.2 Suffix Automaton

```
namespace SAM{
   int scnt, root, last;
   int fa[MAXN<<1], len[MAXN<<1], ch[MAXN<<1][26];
   int sc[MAXN<<1], tmpl[MAXN<<1];

int newnode(int _len, int q = 0) {
   fa[++scnt] = fa[q]; len[scnt] = _len;
   sc[scnt] = 0; tmpl[scnt] = INF;</pre>
```

```
9
            for(int i = 0; i < 26; i++) ch[scnt][i] = ch[q][i];</pre>
10
            return scnt;
11
       }
       void init() {
12
13
            scnt = 0;
            root = last = newnode(0);
14
15
       void extend(int c) {
16
17
            int p = last, np = newnode(len[p] + 1);
18
            for(;p && ch[p][c] == 0; p = fa[p]) ch[p][c] = np;
19
            if(!p) fa[np] = root;
20
            else{
21
                int q = ch[p][c];
22
                if(len[p] + 1 == len[q]) fa[np] = q;
23
                else{
                    int nq = newnode(len[p] + 1, q);
24
                    fa[np] = fa[q] = nq;
25
26
                    for(; p && ch[p][c] == q; p = fa[p]) ch[p][c] = nq;
27
28
29
           last = np;
30
       int c[MAXN], rs[MAXN << 1];</pre>
31
32
       void radix_sort(int n){
33
           for(int i = 0; i <= n; i++) c[i] = 0;
34
            for(int i = 1; i <= scnt; i++) c[len[i]]++;</pre>
35
            for(int i = 1; i <= n; i++) c[i] += c[i-1];
36
           for(int i = scnt; i >= 1; i--) rs[c[len[i]]--] = i;
37
38
       void go(){
            scanf("%s",s);
39
40
            int n = strlen(s);
            for(int i = 0; i < n; ++i)</pre>
41
                extend(s[i] - 'a');
42
43
           radix_sort(n);
            //以下sc集合意义不同
44
            {//每个节点对应的位置之后有多少个不同子串
45
                for(int i = scnt; i >= 1; i--) {
46
47
                    int S = 0;
48
                    for(int j = 0; j < 26; j++)
49
                        S += sc[ ch[rs[i]][j] ];
50
                    sc[rs[i]] = S + 1;
                }
51
52
            {//right集合大小
53
                int cur = root;
54
55
                for(int i = 0; i < n; ++i) {</pre>
56
                    cur = ch[cur][s[i] - 'a'];
57
                    sc[cur]++;
58
59
                for(int i = scnt; i >= 1; --i) {
60
                    sc[ fa[rs[i]] ] += sc[rs[i]];
61
62
            //公共子串
63
            //tmpl,当前字符串:在状态cur,与模板串的最长公共后缀
64
            //minl, 多个字符串:在状态cur,与模板串的最长公共后缀
65
            //注意:在状态cur匹配成功时, cur的祖先状态与字符串的最长公共后缀
66
            for(; ~scanf("%s",s);) {
67
                int cur = root, Blen = 0;
68
69
                for(int i = 0; i <= scnt; i++)</pre>
```

```
70
                      tmpl[i] = 0;
                 n = strlen(s);
71
72
                  for(int i = 0, x; i < n; i++) {</pre>
73
                      x = s[i] - 'a';
74
                      if(ch[cur][x]) {
75
                          ++Blen;
76
                          cur = ch[cur][x];
77
                      }else{
78
                           for(; cur && ch[cur][x] == 0; cur = fa[cur]);
79
                          if(cur) {
                               Blen = len[cur] + 1;
80
                               cur = ch[cur][x];
81
82
                          }else{
83
                               cur = root; Blen = 0;
84
                          }
85
                      }
                      tmpl[cur] = max(tmpl[cur], Blen);
86
87
                  for(int i = scnt; i ; --i) {
88
89
                      if( tmpl[ fa[rs[i]] ] < tmpl[ rs[i] ])</pre>
                          tmpl[ fa[rs[i]] ] = len[ fa[rs[i]] ];
90
91
                      minl[ rs[i] ] = min(minl[ rs[i] ], tmpl[ rs[i] ]);
                 }
92
93
             }
94
95
    }
96
    namespace exSAM{
         int scnt, root;
97
98
         int fa[MAXN<<1], len[MAXN<<1], ch[MAXN<<1][26];</pre>
99
         int sc[MAXN<<1], tmpl[MAXN<<1], minl[MAXN<<1];</pre>
100
         int newnode(int _len, int q = 0) {
101
             fa[++scnt] = fa[q]; len[scnt] = _len;
102
             sc[scnt] = 0;tmpl[scnt] = 0; minl[scnt] = INF;
103
             for(int i = 0; i < 26; i++) ch[scnt][i] = ch[q][i];</pre>
104
105
             return scnt;
106
         }
         void init() {
107
108
             scnt = 0;
109
             root = newnode(0);
110
111
         int work(int p,int c){
112
             int q = ch[p][c];
             int nq = newnode(len[p] + 1, q);
113
114
             fa[q] = nq;
             for(; p \&\& ch[p][c] == q; p = fa[p]) ch[p][c] = nq;
115
             return nq;
116
117
         int extend(int p, int c) {
118
119
             if (ch[p][c]){
120
                  int q = ch[p][c];
                  if (len[p] + 1 == len[q]) return q;
121
122
                  return work(p, c);
             }
123
124
             int np = newnode(len[p] + 1);
             for(;p && ch[p][c] == 0; p = fa[p]) ch[p][c] = np;
125
126
             if (!p) fa[np] = root;
127
             else{
128
                  int q = ch[p][c];
129
                  if (len[p] + 1 == len[q]) fa[np] = q;
130
                  else fa[np] = work(p, c);
```

```
131
              }
132
              return np;
133
         }
         void solve() {
134
             int n; scanf("%d",&n);
135
              for(int i = 1; i <= n; i++) {</pre>
136
                  scanf("%s", s);
137
138
                  int sn = strlen(s);
139
                  int last = root;
140
                  for(int j = 0; j < sn; ++j)
                       last = extend(last, s[j] - 'a');
141
             }
142
143
    }
144
```

## 3.4 Palindrome Related

## 3.4.1 Manacher

```
namespace Manachar {
1
2
        char S[MAXN << 1];</pre>
3
        int scnt, ans;
4
        int p[MAXN << 1]; //p[i] - 1
5
        void init(char *s0, int sn0) {
6
             S[0] = '\$'; S[1] = '\#';
             for(int i = 0; i < sn0; i++) {</pre>
7
                 S[2 * i + 2] = s0[i];
8
                 S[2 * i + 3] = '\#';
9
10
             scnt = sn0 * 2 + 2;
11
             S[scnt] = \frac{1}{2};
12
13
        void manachar() {
14
             int id = 0, mx = 0;
15
             for(int i = 1; i < scnt; i++) {</pre>
16
17
                 p[i] = mx > i ? min(p[2 * id - i], mx - i) : 1;
                 while(S[i + p[i]] == S[i - p[i]]) p[i]++;
18
19
                 if(i + p[i] > mx) {
20
                      mx = i + p[i];
21
                      id = i;
22
                 }
23
             }
24
        }
25
```

## 3.4.2 Palindromic Automaton

```
1
   namespace PAM {
2
        int scnt, S[MAXN];
        int pcnt, last, len[MAXN], fail[MAXN], ch[MAXN][26];
3
        int\ cnt[MAXN]; //节点i表示的本质不同的串的个数(调用count())
4
       int num[MAXN]; //以节点 i表示的最长回文串的最右端点为回文串结尾的回文串个数int newnode(int _len) {
    len[pcnt] = _len;
5
6
7
            cnt[pcnt] = num[pcnt] = 0;
8
            for(int i = 0; i < 26; i++) ch[pcnt][i] = 0;</pre>
9
10
            return pcnt++;
11
```

```
inline void init() {
12
13
            S[scnt = 0] = -1;
            pcnt = 0;newnode(0);newnode(-1);
14
15
            fail[0] = 1;last = 0;
16
        int getfail(int x) {
17
18
            while(S[scnt - len[x] - 1] != S[scnt]) x = fail[x];
19
            return x;
20
21
        void extend(int c) {
22
            S[++scnt] = c;
23
            int cur = getfail(last);
24
            if(!ch[cur][c]) {
25
                int now = newnode(len[cur] + 2);
                fail[now] = ch[getfail(fail[cur])][c];
26
27
                ch[cur][c] = now;
28
                num[now] = num[fail[now]] + 1;
29
30
            last = ch[cur][c];
31
            cnt[last]++;
32
        void count() {
33
            for(int i = pcnt - 1; i >= 0; i--) cnt[fail[i]] += cnt[i];
34
35
36
   };
```

# 3.5 Substring Automaton

当字符集过大时使用主席树维护 ch

# 4 Math

# 4.1 Algebra

### 4.1.1 FFT

```
//不预处理精度
      const double pi = acos(-1.0);
       const int MAXN = 300003;
        struct comp {
 4
 5
                 double x, y;
                 comp operator + (const comp& a) const { return (comp) {x + a.x, y + a.y}; }
 6
 7
                 comp operator - (const comp& a) const { return (comp) {x - a.x, y - a.y}; }
                 comp operator * (const comp& a) const { return (comp) \{x * a.x - y * a.y, x * a.y + a.y 
 8
                y * a.x}; }
 9
       };
10
       int rev[MAXN], T;
        comp tmp;
11
        void fft(comp *a, int r) {
12
                 if(r == -1) for(int i = 0; i < T; i++) a[i] = a[i] * a[i];</pre>
13
14
                 for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
15
                 for(int i = 2, mid = 1; i \leftarrow T; mid = i, i \leftarrow 1) {
16
                          comp step = (comp) {cos(pi / mid), r * sin(pi / mid)};
17
                          for(int j = 0; j < T; j += i) {
                                   comp cur = (comp) {1, 0};
18
                                   for(int k = j; k < j + mid; k++, cur = cur * step) {
19
                                           tmp = a[k + mid] * cur;
20
                                           a[k + mid] = a[k] - tmp;
21
22
                                           a[k] = a[k] + tmp;
23
                                   }
24
                         }
25
26
                 if(r == -1) for(int i = 0; i < T; i++) a[i].y = (int)(a[i].y / T / 2 + 0.5);
27
28
       comp A[MAXN];
29
       void init(int n) {
30
                 //A[0] = (comp) \{0, 0\};
31
                 for(T = 1; T <= n; T <<= 1);
32
                 for(int i = 1; i < T; i++) {</pre>
33
                          if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
34
                          else rev[i] = rev[i >> 1] >> 1;
35
                          //A[i] = (comp) \{0, 0\};
36
                 }
37
        //预处理精度
38
39
        int rev[MAXN], T;
        comp Sin[MAXN], tmp;
40
        void fft(comp *a, int r) {
41
42
                 if(r == -1) {
                          for(int i = 0; i < (T >> 1); i++) Sin[i].y = -Sin[i].y;
43
44
                          for(int i = 0; i < T; i++) a[i] = a[i] * a[i];</pre>
45
                 for(int i = 1; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
46
                 for(int i = 2, mid = 1, s = (T >> 1); i <= T; mid = i, i <<= 1, s >>= 1) {
47
48
                          for(int j = 0; j < T; j += i) {</pre>
49
                                   for(int k = j, cur = 0; k < j + mid; k++, cur += s) {
                                           tmp = a[k + mid] * Sin[cur];
50
51
                                           a[k + mid] = a[k] - tmp;
52
                                           a[k] = a[k] + tmp;
```

```
53
54
            }
55
        if(r == -1) for(int i = 0; i < T; i++) a[i].y = (int)(a[i].y / T / 2 + 0.5);
56
57
58
   comp A[MAXN];
    void init(int n) {
59
60
        for(T = 1; T <= n; T <<= 1);
61
        for(int i = 0; i < T; i++) {</pre>
            if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
62
            else rev[i] = rev[i >> 1] >> 1;
63
64
            //A[i] = (comp) \{0, 0\};
65
        for(int i = 0; i < (T >> 1); i++) {
66
            Sin[i] = (comp) \{cos(2 * pi * i / T), sin(2 * pi * i / T)\};
67
68
69
   }
70
   int main() {
        scanf("%d%d", &n, &m);
71
72
        init(n + m);
        for(int i = 0; i <= n; i++) scanf("%lf", &A[i].x);</pre>
73
74
        for(int i = 0; i <= m; i++) scanf("%lf", &A[i].y);</pre>
75
        fft(A, 1);
76
        fft(A, -1);
77
        for(int i = 0; i <= n + m; i++) printf("%d%c", (int)(A[i].y), i == n + m ? '\n': '
        ');
78
        return 0;
79
   }
```

## 4.1.2 NTT

**4.**常用NTT模数:

以下模数的共同g=3189

$p=r\times 2^k+1$	k	g
104857601	22	3
167772161	25	3
469762049	26	3
950009857	21	7
998244353	23	3
1004535809	21	3
2013265921	27	31
2281701377	27	3
3221225473	30	5

```
const int MAXN = 300005, G = 3, mod = 998244353; //or (479LL<<21) + 1
2
   int rev[MAXN], T;
3
   LL qpow(LL x, LL y) {
       LL res = 1;
4
       while(y) {
5
6
           if(y & 1) res = res * x % mod;
7
           x = x * x % mod;
           y >>= 1;
8
9
10
       return res;
11
12 LL A[MAXN], B[MAXN];
```

```
13
   void ntt(LL *a, int r) {
        if(r == -1) for(int i = 0; i < T; i++) A[i] = A[i] * B[i] % mod;
14
        for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
15
16
        for(int i = 2, mid = 1; i <= T; mid = i, i <<= 1) {
17
            LL gn = qpow(G, (mod - 1) / i);
            if(r == -1) gn = qpow(gn, mod - 2);
18
            for(int j = 0; j < T; j += i) {</pre>
19
                 LL cur = 1, tmp;
20
21
                 for(int k = j; k < j + mid; k++, cur = cur * gn % mod) {</pre>
                     tmp = a[k + mid] * cur % mod;
22
                     a[k + mid] = ((a[k] - tmp) \% mod + mod) \% mod;
23
                     a[k] = (a[k] + tmp) \% mod;
24
25
                 }
26
            }
27
28
        if(r == -1) {
            LL inv = qpow(T, mod - 2);
29
30
            for(int i = 0; i < T; i++) a[i] = a[i] * inv % mod;</pre>
31
32
    void init(int n) {
33
        for(T = 1; T <= n; T <<= 1);
34
35
        for(int i = 0; i < T; i++) {</pre>
36
            if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
37
            else rev[i] = rev[i >> 1] >> 1;
38
        }
39
   }
```

## 4.1.3 MTT

```
#include <bits/stdc++.h>
   using namespace std;
   typedef long long LL;
   const double pi = acos(-1.0);
5
   const int MAXN = 300003;
6
   struct comp {
7
        double x, y;
        comp operator + (const comp& a) const { return (comp) {x + a.x, y + a.y}; }
8
        comp operator - (const comp& a) const { return (comp) {x - a.x, y - a.y}; }
9
       comp operator * (const comp& a) const { return (comp) {x * a.x - y * a.y, x * a.y +
10
       y * a.x}; }
11
   };
   #define conj(a) ((comp){a.x, -a.y})
12
13
   int rev[MAXN], T;
   comp Sin[MAXN], tmp;
14
   void fft(comp *a, int r) {
15
16
        for(int i = 1; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
17
        for(int i = 2, mid = 1, s = (T >> 1); i <= T; mid = i, i <<= 1, s >>= 1) {
            for(int j = 0; j < T; j += i) {</pre>
18
19
                for(int k = j, cur = 0; k < j + mid; k++, cur += s) {</pre>
20
                    tmp = a[k + mid] * Sin[cur];
21
                    a[k + mid] = a[k] - tmp;
22
                    a[k] = a[k] + tmp;
23
                }
            }
24
25
        }
   }
26
27
   void init(int n) {
28
        for(T = 1; T <= n; T <<= 1);
        for(int i = 0; i < T; i++) {</pre>
```

```
30
            if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
31
            else rev[i] = rev[i >> 1] >> 1;
32
        for(int i = 0; i < (T >> 1); i++) {
33
            Sin[i] = (comp) \{cos(2 * pi * i / T), sin(2 * pi * i / T)\};
34
35
36
37
   int n, m, mod;
    void mtt(int *x, int *y) {
38
        for(int i = 0; i < T; i++) (x[i] += mod) %= mod, (y[i] += mod) %= mod;
39
40
        static comp a[MAXN], b[MAXN];
        static comp dfta[MAXN], dftb[MAXN], dftc[MAXN], dftd[MAXN];
41
        for(int i = 0; i < T; i++) {</pre>
42
43
            a[i] = \{x[i] \& 0x7fff, x[i] >> 15\};
44
            b[i] = {y[i] \& 0x7fff, y[i] >> 15};
45
        fft(a, 1); fft(b, 1);
46
        for(int i = 0; i < T; i++) {</pre>
47
            int j = (T - i) & (T - 1);
48
49
            static comp da, db, dc, dd;
            da = (a[i] + conj(a[j])) * (comp){0.5, 0};
50
            db = (a[i] - conj(a[j])) * (comp){0, -0.5};
51
            dc = (b[i] + conj(b[j])) * (comp){0.5, 0};
52
            dd = (b[i] - conj(b[j])) * (comp){0, -0.5};
53
            dfta[j] = da * dc;
54
55
            dftb[j] = da * dd;
56
            dftc[j] = db * dc;
57
            dftd[j] = db * dd;
58
        for(int i = 0; i < T; i++) {</pre>
59
            a[i] = dfta[i] + dftb[i] * (comp) {0, 1};
60
            b[i] = dftc[i] + dftd[i] * (comp) {0, 1};
61
62
        //for(int \ i = 0; \ i < (T >> 1); \ i++) \ Sin[i].y = -Sin[i].y;
63
        fft(a, -1); fft(b, -1);
64
        for(int i = 0; i < T; i++) {</pre>
65
            static int da, db, dc, dd;
66
            da = (LL)(a[i].x / T + 0.5) \% mod;
67
            db = (LL)(a[i].y / T + 0.5) \% mod;
68
69
            dc = (LL)(b[i].x / T + 0.5) \% mod;
70
            dd = (LL)(b[i].y / T + 0.5) \% mod;
71
            x[i] = ((da + ((LL)(db + dc) << 15) + ((LL)dd << 30)) % mod + mod) % mod;
72
        }
73
   }
74
   int main() {
        static int a[MAXN], b[MAXN];
75
76
        scanf("%d%d%d", &n, &m, &mod);
        for(int i = 0; i <= n; i++) scanf("%d", a + i);</pre>
77
78
        for(int i = 0; i <= m; i++) scanf("%d", b + i);</pre>
79
        init(n + m);
80
        mtt(a, b);
        for(int i = 0; i <= n + m; i++) printf("%d%c", a[i], i == n + m? '\n': '');
81
82
        return 0;
83
```

## 4.1.4 FWT

```
void FWT(LL *a,int n) {
   for(int i = 2;i <= n; i <<= 1) {
    for(int j = 0; j < n; j += i) {</pre>
```

```
4
                for(int d = 0, w = i >> 1; d < w; d++){
5
                     LL u = a[j + d], v = a[j + d + w];
6
                     //xor: a[j + d] = u + v, a[j + d + w] = u - v;
7
                     //and: a[j + d] = u + v;
8
                     //or : a/j + d + w/ = u + v;
9
                }
10
            }
11
        }
12
    void UFWT(LL *a, int n) {
13
        for(int i = 2; i <= n; i <<= 1) {
14
            for(int j = 0; j < n; j += i) {</pre>
15
                for(int d = 0, w = i >> 1; d < w; d++) {
16
17
                     LL u = a[j + d], v = a[j + d + w];
                     //xor: a[j + d] = (u + v) / 2, a[j + d + w] = (u - v) / 2;
18
                     //and: a[j + d] = u - v;
19
                     //or : a/j + d + w/ = v - u;
20
21
                }
22
            }
23
        }
24
   }
25
   void solve(int n) {
26
        FWT(a, n); FWT(b, n);
27
        for(int i = 0; i < n; i++) a[i] = a[i] * b[i];</pre>
28
        UFWT(a, n);
29
   }
```

## 4.1.5 FFT Divide and Conquer

$$f_i = \sum_{j=1}^{i-1} f_j \cdot g_{i-j}$$

```
1 #include <bits/stdc++.h>
2 using namespace std;
   typedef long long LL;
3
   const int MAXN = 300005, G = 3, mod = 998244353;
4
   namespace NTT {
5
6
        LL A[MAXN], B[MAXN]
        int rev[MAXN], T;
7
8
        LL qpow(LL x, LL y) {
9
            LL res = 1;
10
            while(y) {
                if(y & 1) res = res * x % mod;
11
                x = x * x % mod;
12
13
                y >>= 1;
            }
14
            return res;
15
16
        void ntt(LL *a, int r) {
17
18
            for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
19
            for(int i = 2, mid = 1; i <= T; mid = i, i <<= 1) {
20
                LL gn = qpow(G, (mod - 1) / i);
21
                if(r == -1) gn = qpow(gn, mod - 2);
22
                for(int j = 0; j < T; j += i) {
23
                    LL cur = 1, tmp;
24
                    for(int k = j; k < j + mid; k++, cur = cur * gn % mod) {
                        tmp = a[k + mid] * cur % mod;
25
26
                        a[k + mid] = ((a[k] - tmp) \% mod + mod) \% mod;
```

```
27
                          a[k] = (a[k] + tmp) \% mod;
28
                     }
29
                 }
30
            }
            if(r == -1) {
31
                 LL inv = qpow(T, mod - 2);
32
33
                 for(int i = 0; i < T; i++) a[i] = a[i] * inv % mod;</pre>
34
             }
35
        void init(int n) {
36
            for(T = 1; T <= n; T <<= 1);
37
             for(int i = 0; i < T; i++) {</pre>
38
                 if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
39
40
                 else rev[i] = rev[i >> 1] >> 1;
41
            }
42
        }
43
   }
   LL f[MAXN], g[MAXN];
44
   using namespace NTT;
45
46
   void solve(int 1, int r) {
47
        if(1 == r) return;
        int mid = (1 + r) >> 1;
48
49
        solve(l, mid);
50
        init(r - 1);
51
        for(int i = 0; i < T; i++) A[i] = B[i] = 0;
52
        for(int i = 0; i <= mid - 1; i++) A[i] = f[i + 1];
53
        for(int i = 0; i <= r - 1; i++) B[i] = g[i];</pre>
54
        ntt(A, 1); ntt(B, 1);
55
        for(int i = 0; i < T; i++) A[i] = A[i] * B[i] % mod;
        ntt(A, -1);
56
        for(int i = mid + 1; i <= r; i++) f[i] =(f[i] + A[i - 1]) % mod;</pre>
57
58
        solve(mid + 1, r);
   }
59
   int main() {
60
        int n; scanf("%d", &n);
61
        for(int i = 1; i < n; i++) scanf("%lld", g + i);</pre>
62
        f[0] = 1;
63
64
        solve(0, n - 1);
        for(int i = 0; i < n; i++) printf("%lld%c", f[i], i == n - 1 ? ' \setminus n' : ' ');
65
66
67
```

# 4.1.6 Linear Basis

```
//dynamic
1
2
   const int D = 60;
3
   struct Basis {
4
        vector<int> ind;
5
        vector<LL> base;
6
        Basis() {
7
            ind.resize(D, -1);
8
            base.resize(D);
9
10
        bool update(LL x, int id) {
            for(int i = 0; i < D; i++) if(~ind[i] && x >> i & 1) {
11
                x ^= base[i];
12
13
            if(!x) return 1;
14
15
            int pos = __builtin_ctzll(x);
16
            ind[pos] = id;
```

```
17
             base[pos] = x;
18
             return 0;
19
        }
20
    };
21
    //array
22
    int Gauss(int n, int m) {
23
        int num = 1;
24
        for(int x = 1; x <= n && x <= m; x++) {
25
             int t = 0;
26
             for(int j = x; j \le m; j++) if(g[j][x]) { t = j; break; }
27
             if(t) {
28
                 swap(g[x], g[t]);
29
                 for(int i = x + 1; i <= n; i++) {</pre>
                     if(g[i][x]) {
30
31
                          for(int k = 1; k \le m; k++) g[i][k] ^= g[x][k];
32
33
                 }
34
                 num++;
35
            }
36
37
        return --num;
38
   }
39
    //long long
40
    int Gauss() {
41
        int num = 1;
42
        for(int k = 61; k >= 0; k--) {
43
             int t = 0;
44
             for(int j = num; j \leftarrow cnt; j++) if((A[j] >> k) & 1) { t = j; break; }
45
             if(t) {
46
                 swap(A[t], A[num]);
                 for(int j = num + 1; j <= cnt; j++) if((A[j] >> k) & 1) A[j] ^= A[num];
47
48
49
             }
50
        }
51
        return --num;
52
   }
53
    //det
   LL det(int n){
54
55
        LL ret = 1;
56
        for(int i = 1; i < n; i++){</pre>
57
             for(int j = i + 1; j < n; j++)</pre>
58
                 while(a[j][i]){
                     LL t = (LL)a[i][i] / a[j][i];
59
60
                     for(int k = i; k < n; k++) {</pre>
                          a[i][k] = (a[i][k]-a[j][k] * t);
61
62
                          swap(a[i][k], a[j][k]);
63
                     }
64
                     ret = -ret;
65
66
             if(a[i][i] == 0)return 0;
67
             ret = ret * a[i][i];
68
        if(ret<0)ret = -ret;</pre>
69
70
        return ret;
71
    }
```

## 4.1.7 Polynomial

Inverse:

$$A_x * B'_x \equiv 1 \qquad (\text{mod } x^{\frac{n}{2}})$$

$$A_x * B_x \equiv 1 \qquad (\text{mod } x^n)$$

$$B_x \equiv 2 \cdot B'_x - A_x \cdot {B'}_x^2 \qquad (\text{mod } x^n)$$

Division:

$$A_r(x) = x^n A(\frac{1}{x}) \Rightarrow A_r(x)[i] = A(x)[n-i]$$

$$A(x) = B(x) * Q(x) + R(x)$$

$$Q_r(x) = A_r(x) * B_r^{-1}(x)$$

$$R(x) = A(x) - B(x) * Q(x)$$

```
1 //NTT模数
2 #include <bits/stdc++.h>
3 using namespace std;
4 typedef long long LL;
   const int MAXN = 300005, G = 3, mod = 998244353;
6
   namespace NTT {
7
        int rev[MAXN], T;
        LL qpow(LL x, LL y) {
8
            LL res = 1;
9
10
            while(y) {
                if(y & 1) res = res * x % mod;
11
                x = x * x % mod;
12
13
                y >>= 1;
            }
14
15
            return res;
16
17
        void ntt(LL *a, int r) {
            for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
18
            for(int i = 2, mid = 1; i <= T; mid = i, i <<= 1) {
19
20
                 LL gn = qpow(G, (mod - 1) / i);
21
                if(r == -1) gn = qpow(gn, mod - 2);
22
                for(int j = 0; j < T; j += i) {</pre>
                     LL cur = 1, tmp;
23
                     for(int k = j; k < j + mid; k++, cur = cur * gn % mod) {
24
25
                         tmp = a[k + mid] * cur % mod;
26
                         a[k + mid] = ((a[k] - tmp) \% mod + mod) \% mod;
                         a[k] = (a[k] + tmp) \% mod;
27
                     }
28
29
                }
30
            if(r == -1) {
31
                 LL inv = qpow(T, mod - 2);
32
                for(int i = 0; i < T; i++) a[i] = a[i] * inv % mod;</pre>
33
34
35
        void init(int n) {
36
37
            for(T = 1; T <= n; T <<= 1);
            for(int i = 0; i < T; i++) {</pre>
38
                 if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
39
                else rev[i] = rev[i >> 1] >> 1;
40
```

```
41
             }
42
        }
43
    }
44
    namespace poly {
    using namespace NTT;
45
         void mul(LL *a, LL *b, LL *c, int n, int m) {
46
47
             init(n + m);
             static LL x[MAXN], y[MAXN];
48
49
             for(int i = 0; i < T; i++) {
50
                 x[i] = i <= n ? a[i] : 0;
51
                 y[i] = i <= m ? b[i] : 0;
             }
52
53
             ntt(x, 1); ntt(y, 1);
54
             for(int i = 0; i < T; i++) c[i] = x[i] * y[i] % mod;</pre>
55
             ntt(c, -1);
56
         void poly_inv(LL *a, LL *b, int n) {
57
58
             if(n == 1) {
                 b[0] = qpow(a[0], mod - 2);
59
                 return;
60
61
             static LL c[MAXN], d[MAXN];
62
63
             memset(c, 0, n * sizeof(LL));
64
             poly_inv(a, c, n >> 1);
             for(int i = 0; i < n; i++) {</pre>
65
66
                 d[i] = a[i];
67
             }
             init(n);
68
69
             ntt(c, 1); ntt(d, 1);
             for(int i = 0; i < T; i++) b[i] = c[i] * (2 + mod - d[i] * c[i] % mod) % mod;
70
71
             ntt(b, -1);
             for(int i = n; i < T; i++) b[i] = 0;</pre>
72
73
         void inv(LL *a, LL *b, int n) {//A must be different from B
74
75
             init(n);
76
             poly_inv(a, b, T);
77
         }
         //A_x = B_x * Q_x + R_x;
78
79
         void div(LL *A, LL *B, LL *Q, LL *R, int n, int m) {
80
             static LL f[MAXN], g[MAXN], inv_g[MAXN];
81
             for(int i = 0; i <= n; i++) f[n - i] = A[i];
82
             for(int i = 0; i <= m; i++) g[m - i] = B[i];</pre>
83
             //inv(G_r)
             for(int i = n - m + 1; i \leftarrow m; i++) g[i] = 0;
84
85
             inv(g, inv_g, n - m);
86
             //Q
87
             mul(f, inv_g, f, n, n - m);
88
             for(int i = 0; i <= n - m; i++) Q[i] = f[n - m - i];
89
90
             mul(Q, B, f, n - m, m);
91
             for(int i = 0; i < m; i++) R[i] = (A[i] + mod - f[i]) % mod;
92
93
    LL A[MAXN], B[MAXN];
94
95
    LL Q[MAXN], R[MAXN];
96
    int main() {
97
         ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
98
         int n, m; cin >> n >> m;
99
         for(int i = 0; i <= n; i++) cin >> A[i];
100
         for(int i = 0; i <= m; i++) cin >> B[i];
101
         poly::div(A, B, Q, R, n, m);
```

```
102
        for(int i = 0; i <= n - m; i++) cout << Q[i] << " "; cout << endl;</pre>
         for(int i = 0; i < m; i++) cout << R[i] << " "; cout << endl;</pre>
103
104
         return 0;
105
    //非NTT模数求逆
106
107
    #include <bits/stdc++.h>
108
    using namespace std;
    typedef long long LL;
109
110
    const double pi = acos(-1.0);
    const int MAXN = 300003;
111
112
    struct comp {
         double x, y;
113
         comp operator + (const comp& a) const { return (comp) {x + a.x, y + a.y}; }
114
         comp operator - (const comp& a) const { return (comp) {x - a.x, y - a.y}; }
115
        comp operator * (const comp& a) const { return (comp) {x * a.x - y * a.y, x * a.y +
116
        y * a.x}; }
117
    };
    #define conj(a) ((comp){a.x, -a.y})
118
    int rev[MAXN], T;
120 comp Sin[MAXN], tmp;
121
    void fft(comp *a, int r) {
         for(int i = 1; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
122
         for(int i = 2, mid = 1, s = (T >> 1); i <= T; mid = i, i <<= 1, s >>= 1) {
123
             for(int j = 0; j < T; j += i) {</pre>
124
                 for(int k = j, cur = 0; k < j + mid; k++, cur += s) {</pre>
125
126
                      tmp = a[k + mid] * Sin[cur];
127
                      a[k + mid] = a[k] - tmp;
128
                      a[k] = a[k] + tmp;
129
                 }
130
             }
131
132
    void init(int n) {
133
         //for(T = 1; T \le n; T \le 1);
134
135
         T = n << 1;
         for(int i = 0; i < T; i++) {</pre>
136
137
             if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
             else rev[i] = rev[i >> 1] >> 1;
138
139
140
         for(int i = 0; i < (T >> 1); i++) {
141
             Sin[i] = (comp) \{cos(2 * pi * i / T), sin(2 * pi * i / T)\};
142
143
    }
144
    const int mod = 1e9 + 7;
    void mtt(int *x, int *y) {
145
146
         for(int i = 0; i < T; i++) (x[i] += mod) %= mod, (y[i] += mod) %= mod;
         static comp a[MAXN], b[MAXN];
147
         static comp dfta[MAXN], dftb[MAXN], dftc[MAXN], dftd[MAXN];
148
149
         for(int i = 0; i < T; i++) {</pre>
             a[i] = \{x[i] \& 0x7fff, x[i] >> 15\};
150
151
             b[i] = {y[i] \& 0x7fff, y[i] >> 15};
152
153
         fft(a, 1); fft(b, 1);
         for(int i = 0; i < T; i++) {</pre>
154
             int j = (T - i) & (T - 1);
155
             static comp da, db, dc, dd;
156
             da = (a[i] + conj(a[j])) * (comp){0.5, 0};
157
             db = (a[i] - conj(a[j])) * (comp){0, -0.5};
158
             dc = (b[i] + conj(b[j])) * (comp){0.5, 0};
159
             dd = (b[i] - conj(b[j])) * (comp){0, -0.5};
160
161
             dfta[j] = da * dc;
```

```
162
             dftb[j] = da * dd;
163
             dftc[j] = db * dc;
164
             dftd[j] = db * dd;
165
166
         for(int i = 0; i < T; i++) {</pre>
             a[i] = dfta[i] + dftb[i] * (comp) {0, 1};
167
168
             b[i] = dftc[i] + dftd[i] * (comp) {0, 1};
169
170
         //for(int \ i = 0; \ i < (T >> 1); \ i++) \ Sin[i].y = -Sin[i].y;
171
         fft(a, -1); fft(b, -1);
         for(int i = 0; i < T; i++) {</pre>
172
             static int da, db, dc, dd;
173
             da = (LL)(a[i].x / T + 0.5) \% mod;
174
             db = (LL)(a[i].y / T + 0.5) \% mod;
175
             dc = (LL)(b[i].x / T + 0.5) \% mod;
176
             dd = (LL)(b[i].y / T + 0.5) \% mod;
177
             x[i] = ((da + ((LL)(db + dc) << 15) + ((LL)dd << 30)) % mod + mod) % mod;
178
         }
179
180
    }
181
    LL qpow(LL x, LL y) {
182
         LL res = 1;
183
         while(y) {
             if(y & 1) res = res * x % mod;
184
             x = x * x % mod;
185
186
             y >>= 1;
187
188
         return res;
189
190
    void poly inv(int *a, int *b, int n) {
191
         if(n == 1) {
192
             b[0] = qpow(a[0], mod - 2);
193
             return;
194
         static int c[MAXN], cc[MAXN], d[MAXN];
195
         memset(c, 0, n * sizeof(int)); memset(cc, 0, n * sizeof(int)); memset(d, 0, n *
196
         sizeof(int));
197
         poly_inv(a, c, n >> 1);
198
         for(int i = 0; i < n; i++) cc[i] = c[i], d[i] = a[i];</pre>
199
         init(n);
200
         mtt(cc, c);
201
         mtt(d, cc);
202
         for(int i = 0; i < T; i++) b[i] = (2 * c[i] % mod + mod - d[i]) % mod;
203
         for(int i = 0; i < n; i++) b[n + i] = 0;
204
    }
205
    int A[MAXN], B[MAXN];
    int main() {
206
         ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
207
208
         int n; cin >> n; n--;
         for(T = 1; T <= n; T <<= 1);
209
210
         for(int i = 0; i <= n; i++) cin >> A[i];
         poly_inv(A, B, T);
211
         for(int i = 0; i <= n; i++) cout << B[i] << " ";
212
213
         cout << endl;</pre>
214
         return 0;
215
    }
216
    //LOJ挑战多项式(备用)
217
218 #include <bits/stdc++.h>
219 #define rep(i, a, b) for (int i = (a); i \leftarrow (b); i++)
220 #define per(i, a, b) for (int i = (a); i >= (b); i--)
221 #define REP(i, n) for (int i = (0); i < (n); i++)
```

```
222 #define fi first
223 #define se second
224
    #define mp make_pair
225
    #define pb push_back
226
    using namespace std;
    typedef unsigned long long ull;
227
228
    typedef pair<int, int> pii;
229
    typedef vector<int> poly;
    typedef long long 11;
230
231
     struct ano {
232
         char a[1 << 25], *s;</pre>
233
         char b[1 << 25], *t;</pre>
234
         ano() : s(a), t(b) { a[fread(a, 1, sizeof a, stdin)] = 0; }
235
         ~ano() { fwrite(b, 1, t - b, stdout); }
236
         operator int() {
237
             int x = 0;
             while (*s < 48) ++s;
238
239
             while (*s > 32) x = x * 10 + *s++ - 48;
240
             return x;
241
         11 in() {
242
243
             11 x = 0;
             while (*s < 48) ++s;
244
             while (*s > 32) x = x * 10 + *s++ - 48;
245
246
             return x;
247
248
         void out(int x, char e = ' ') {
249
             static char c[12];
250
             char *i = c;
251
             if (!x)
                  *t++ = 48;
252
253
             else {
                 while (x) {
254
255
                     int y = x / 10;
256
                      *i++ = x - y * 10 + 48, x = y;
257
258
                 while (i != c) *t++ = *--i;
259
260
             *t++ = e;
261
262 } buf;
263 const int mod = 998244353;
264 namespace Poly {
265
    const int N = (1 << 20) + 5, g = 3;
266
    inline int power(int x, int p) {
267
         int res = 1;
         for (; p; p >>= 1, x = (11)x * x % mod)
268
269
             if (p & 1)
270
                 res = (11)res * x % mod;
271
         return res;
272
273
    inline int fix(const int x) { return x >= mod ? x - mod : x; }
     void dft(poly &A, int n) {
274
         static ull W[N << 1], *H[30], *las = W, mx = 0;</pre>
275
276
         for (; mx < n; mx++) {
277
             H[mx] = las;
278
             ull w = 1, wn = power(g, (mod - 1) >> (mx + 1));
             REP(i, 1 << n) *las++ = w, w = w * wn % mod;
279
280
281
         if (A.size() != (1 << n))</pre>
282
             A.resize(1 << n);
```

```
283
         static ull a[N];
284
         for (int i = 0, j = 0; i < (1 << n); ++i) {
285
             a[i] = A[j];
286
             for (int k = 1 << (n - 1); (j ^= k) < k; k >>= 1)
287
288
         for (int k = 0, d = 1; k < n; k++, d <<= 1)
289
             for (int i = 0; i < (1 << n); i += (d << 1)) {
290
291
                 ull *1 = a + i, *r = a + i + d, *w = H[k], t;
                 for (int j = 0; j < d; j++, l++, r++) {
292
                      t = (*r) * (*w++) % mod;
293
                      *r = *1 + mod - t, *1 += t;
294
295
296
297
         REP(i, 1 << n) A[i] = a[i] \% mod;
298
    }
    void idft(poly &a, int n) {
299
300
         a.resize(1 << n), reverse(a.begin() + 1, a.end());</pre>
301
         dft(a, n);
302
         int inv = power(1 << n, mod - 2);</pre>
303
         REP(i, 1 << n) a[i] = (ll)a[i] * inv % mod;
304
    }
305
    poly FIX(poly a) {
306
         while (!a.empty() && !a.back()) a.pop_back();
         return a;
307
308
    }
309
    poly add(poly a, poly b, int op = 0) {
310
         a.resize(max(a.size(), b.size()));
311
         REP(i, b.size()) a[i] = fix(op ? a[i] + mod - b[i] : a[i] + b[i]);
312
         return FIX(a);
313
314
    poly mul(poly a, poly b, int t = 1) {
         if (t == 1 && a.size() + b.size() <= 24) {</pre>
315
316
             poly c(a.size() + b.size(), 0);
             REP(i, a.size()) REP(j, b.size()) c[i + j] = (c[i + j] + (ll)a[i] * b[j]) % mod;
317
             return FIX(c);
318
319
         }
         int n = 1, aim = a.size() * t + b.size();
320
321
         while ((1 << n) <= aim) n++;</pre>
322
         dft(a, n), dft(b, n);
323
         if (t == 1)
324
             REP(i, 1 << n) a[i] = (ll)a[i] * b[i] % mod;
325
326
             REP(i, 1 << n) a[i] = (ll)a[i] * a[i] % mod * b[i] % mod;
327
         idft(a, n), a.resize(aim);
328
         return FIX(a);
329
    }
    poly mul(poly a, int b) {
330
331
         REP(i, a.size()) a[i] = (ll)a[i] * b % mod;
332
         return FIX(a);
333
334
335
    poly inv(poly a, int n) { // a[0] != 0
336
         a.resize(n);
337
         poly b;
         if (n == 1) {
338
339
             b.pb(power(a[0], mod - 2));
340
             return b;
341
         b = inv(a, n + 1 >> 1);
342
         b = add(mul(b, 2), mul(b, a, 2), 1);
343
```

```
344
    return b.resize(n), b;
345
    }
346
347
    poly Der(poly a) {
348
         REP(i, a.size() - 1) a[i] = (ll)(i + 1) * a[i + 1] % mod;
349
         return a.pop_back(), a;
350
    }
351
    poly Int(poly a) {
352
         static int inv[N];
         inv[1] = 1;
353
354
         a.pb(0);
         rep(i, 2, a.size()) inv[i] = (11)(mod - mod / i) * inv[mod % i] % mod;
355
         per(i, a.size() - 1, 1) a[i] = (ll)a[i - 1] * inv[i] % mod;
356
357
         return a[0] = 0, a;
358
    }
359
    poly Ln(poly a, int n) { // a[0] = 1
360
        a = mul(Der(a), inv(a, n)), a.resize(n - 1);
361
         return FIX(Int(a));
362
363
    poly Exp(poly a, int n) { // a[0] = 0
364
         a.resize(n);
365
         poly b, one(1, 1);
366
         if (n == 1)
367
            return one;
368
         b = Exp(a, n + 1 >> 1);
369
         b = mul(b, add(add(a, Ln(b, n), 1), one));
370
        return b.resize(n), b;
371
    }
372
373
    poly Div(poly a, poly b) {
374
         poly c;
375
         int n = a.size() - 1, m = b.size() - 1;
         if (n < m)
376
377
             return c;
378
         reverse(a.begin(), a.end());
379
         a.resize(n - m + 1);
380
        reverse(b.begin(), b.end());
381
         b.resize(n - m + 1);
382
         c = mul(a, inv(b, n - m + 1));
383
         c.resize(n - m + 1);
384
        return reverse(c.begin(), c.end()), c;
385
    }
    poly Mod(poly a, poly b) { return FIX(add(a, mul(Div(a, b), b), 1)); }
386
387
    inline int chk(int x) { return power(x, (mod - 1) / 2) == 1; }
388
    inline int R() { return rand() % mod; }
389
    inline pii mul(pii a, pii b, int w) {
         return mp(((11)a.fi * b.fi + (11)a.se * b.se % mod * w) % mod, ((11)a.fi * b.se + (
390
        11)a.se * b.fi) % mod);
391
    inline int Sqrt(int x) {
392
393
         if (!chk(x))
394
             return -1;
395
         int a = R();
         while (chk(((11)a * a - x + mod) % mod)) a = R();
396
397
         int w = ((11)a * a - x + mod) % mod, p = (mod + 1) / 2;
         pii res = mp(1, 0), t = mp(a, 1);
398
         for (; p; p >>= 1, t = mul(t, t, w))
399
             if (p & 1)
400
                 res = mul(res, t, w);
401
         assert(!res.se);
402
        return min(res.fi, mod - res.fi);
403
```

```
404
405
    poly Sqrt(poly a, int n) {
406
         if (n == 1) {
407
             poly b(1, Sqrt(a[0]));
408
             return b;
409
410
         a.resize(n);
         poly b = Sqrt(a, n + 1 >> 1);
411
412
         b = mul(add(b, mul(a, inv(b, n))), (mod + 1) / 2);
413
         return b.resize(n), b;
414
     poly fastpow(poly a, 11 k, int n) {
415
         a.resize(n), a = FIX(a);
416
         if (!a.size())
417
418
             return a;
419
         int st = 0, base = 0;
420
         while (!a[st]) ++st;
421
         if (st * k >= n)
422
             return a.resize(0), a;
423
         REP(i, a.size() - st) a[i] = a[i + st];
424
         if (st)
425
             a.resize(a.size() - st);
426
         base = a[0];
427
         11 inv = power(base, mod - 2);
428
         REP(i, a.size()) a[i] = a[i] * inv % mod;
429
         a = FIX(Exp(mul(Ln(a, n), k % mod), n));
430
431
         if (st) {
             reverse(a.begin(), a.end());
432
433
             a.resize(a.size() + st * k);
434
             reverse(a.begin(), a.end());
435
             a.resize(n), a = FIX(a);
436
         base = power(base, k);
437
         REP(i, a.size()) a[i] = (ll)a[i] * base % mod;
438
         return FIX(a);
439
    }
440
       // namespace Poly
441
    using namespace Poly;
442
443
    int main() {
444
         int n = buf + 1, K = buf;
445
         poly a(n, 0), b;
446
         REP(i, n) a[i] = buf;
447
         b = add(a, Exp(Int(inv(Sqrt(a, n), n)), n), 1);
         b[0] = (b[0] + 2 - a[0] + mod) \% mod;
448
449
         a = Ln(b, n);
450
         a[0]++;
         b = Der(fastpow(a, K, n));
451
452
         b.resize(n - 1);
         REP(i, n - 1) buf.out(b[i]);
453
454
         return 0;
455
```

# 4.1.8 Lagrange Polynomial

$$L(x) = \sum_{i=0}^{n} y_i \prod_{j=0, j \neq i}^{n} \frac{x - x_j}{x_i - x_j}$$

```
1
   //O(n^2)
2
   #include <bits/stdc++.h>
3 using namespace std;
   typedef long long LL;
4
5
   typedef pair<int, int> P;
    const int MAXN = 3005, mod = 998244353;
7
    int exgcd(int a, int b, int &x, int &y) {
8
        int d = a;
9
        if(b != 0) {
            d = exgcd(b, a % b, y, x);
10
11
            y -= (a / b) * x;
12
13
        else {
14
            x = 1; y = 0;
15
16
        return d;
17
    }
   int inv(int a) {
18
19
        int x, y;
20
        exgcd(a, mod, x, y);
21
        return (x % mod + mod) % mod;
22
   }
23
    struct Lagrange {
24
        int n, a[MAXN][2];
25
        void init() {
26
            for(int i = 0; i <= n; i++) a[i][0] = a[i][1] = 0;
27
            n = 0;
            a[0][1] = 1;
28
29
30
        int query(int x, int q = 0) {
31
            int res = 0;
            for(int i = n; i >= 0; i--) res = ((LL)res * x + a[i][q]) % mod;
32
33
            return res;
34
        void update(int x, int y) {
35
36
            a[n][0] = 0;
            int v = (LL)(y - query(x) + mod) % mod * inv(query(x, 1)) % mod;
37
            for(int i = 0; i <= n; i++) a[i][0] = (a[i][0] + (LL)a[i][1] * v) % mod;
38
39
            a[++n][1] = 0;
40
            for(int i = n; i; i--) a[i][1] = (a[i - 1][1] + (LL)a[i][1] * (mod - x)) % mod;
41
            a[0][1] = (LL)a[0][1] * (mod - x) % mod;
42
        }
43
   }p;
44
   int main() {
        ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
45
        int Q;
46
        cin >> Q;
47
48
        int op, x, y;
49
        p.n = 0;
50
        p.init();
51
        while(Q--) {
52
            cin >> op >> x;
53
            if(op == 1) {
                 cin >> y;
54
55
                 p.update(x, y);
56
57
            else cout << p.query(x) << endl;</pre>
58
        }
59
        return 0;
60
   //f(x) = y, \quad \stackrel{\cdot}{x} f(z_i)
```

```
62 #include <bits/stdc++.h>
63
    using namespace std;
64
    #define ll long long
    int n, m;
65
    vector<int> x, y, z;
66
67
    namespace Poly{
         const int P = 998244353;
68
69
         vector<int> ans;// for Evaluate()
70
         vector<vector<int>> p;// for Evaluate() & Interpolate()
71
         inline int Pow(11 x, int y=P-2){ // x^y
72
             int ans=1;
             for(; y; y>>=1, x=x*x%P) if(y&1) ans=ans*x%P;
73
74
             return ans;
75
76
         inline int Ge(int x){ int n=1; while(n<=x) n<<=1; return n;}</pre>
77
         inline int Mod(int x){ return x<P?x:x-P;}</pre>
78
         inline void NTT(vector<int> &f, int g, int n){
79
             f.resize(n);
             for(int i=0, j=0; i<n; ++i){</pre>
80
81
                  if(i>j) swap(f[i], f[j]);
82
                 for(int k=n>>1; (j^=k)<k; k>>=1);
83
             }
84
             vector<int> w(n>>1);
85
             for(int i=1; i<n; i<<=1){</pre>
                  for(int j=w[0]=1, w0=(g==1?Pow(3, (P-1)/i/2):Pow(Pow(3, (P-1)/i/2))); j<i;
86
        ++j) w[j]=(ll)w[j-1]*w0%P;
87
                 for(int j=0; j<n; j+=i<<1){</pre>
88
                      for(int k=j; k<j+i; ++k){</pre>
89
                          int t=(11)f[k+i]*w[k-j]%P;
90
                          f[k+i]=Mod(f[k]-t+P);
91
                          f[k]=Mod(f[k]+t);
92
                      }
                 }
93
94
             if(g==-1) for(int i=0, I=Pow(n); i< n; ++i) f[i]=(11)f[i]*I%P;
95
96
97
         inline vector<int> Add(const vector<int> &f, const vector<int> &g){
98
             vector<int> ans=f;
             for(unsigned i=0; i<f.size(); ++i) (ans[i]+=g[i])%=P;</pre>
99
100
             return ans;
101
102
         inline vector<int> Mul(const vector<int> &f, const vector<int> &g)\{//f*g
103
             vector<int> F=f, G=g;
104
             int p=Ge(f.size()+g.size()-2);
105
             NTT(F, 1, p), NTT(G, 1, p);
             for(int i=0; i<p; ++i) F[i]=(l1)F[i]*G[i]%P;</pre>
106
107
             NTT(F, -1, p);
108
             return F.resize(f.size()+g.size()-1), F;
109
         inline vector<int> PolyInv(const vector<int> &f, int n=-1){// 1/f
110
111
             if(n==-1) n=f.size();
112
             vector<int> ans;
113
             if(n==1) return ans.push_back(Pow(f[0])), ans;
             ans=PolyInv(f, (n+1)/2);
114
             vector<int> tmp(&f[0], &f[0]+n);
115
             int p=Ge(n*2-2);
116
             NTT(tmp, 1, p), NTT(ans, 1, p);
117
             for(int i=0; i<p; ++i) ans[i]=(2-(11)ans[i]*tmp[i]%P+P)*ans[i]%P;</pre>
118
119
             NTT(ans, -1, p);
120
             return ans.resize(n), ans;
121
```

```
122
        inline void PolyDiv(const vector<int> &a, const vector<int> &b, vector<int> &d,
        vector<int> &r){//a=d*b+r}
123
             if(b.size()>a.size()) return (void)(d.clear(), r=a);
124
125
             vector<int> A=a, B=b, iB;
126
             int n=a.size(), m=b.size();
127
             reverse(A.begin(), A.end()), reverse(B.begin(), B.end());
128
             B.resize(n-m+1), iB=PolyInv(B, n-m+1);
129
             d=Mul(A, iB);
             d.resize(n-m+1), reverse(d.begin(), d.end());
130
131
132
             r=Mul(b, d);
             for(int i=0; i<m-1; ++i) r[i]=(P+a[i]-r[i])%P;</pre>
133
134
             r.resize(m-1);
135
136
         inline vector<int> Derivative(const vector<int> &a){ // a '
137
             vector<int> ans;
138
             ans.resize(a.size()-1);
             for(unsigned i=1; i<a.size(); ++i) ans[i-1]=(ll)a[i]*i%P;</pre>
139
140
141
         void Evaluate_Interpolate_Init(int 1, int r, int t, const vector<int> &a){
142
             if(l==r) return p[t].clear(), p[t].push_back(P-a[l]), p[t].push_back(1);
143
             int mid=(l+r)/2, k=t<<1;</pre>
144
145
             Evaluate_Interpolate_Init(1, mid, k, a), Evaluate_Interpolate_Init(mid+1, r, k
        |1, a);
146
             p[t]=Mul(p[k], p[k|1]);
147
         }
148
         inline void Evaluate(int 1, int r, int t, const vector<int> &f, const vector<int> &a
        ){
149
             if(r-1+1<=512){
150
                 for(int i=1; i<=r; ++i){</pre>
                      int x=0, j=f.size(), a1=a[i], a2=(ll)a[i]*a[i]%P, a3=(ll)a[i]*a2%P, a4=(
151
        ll)a[i]*a3%P, a5=(ll)a[i]*a4%P, a6=(ll)a[i]*a5%P, a7=(ll)a[i]*a6%P, a8=(ll)a[i]*a7%P
                     while(j>=8)
152
                     x=((11)x*a8+(11)f[j-1]*a7+(11)f[j-2]*a6+(11)f[j-3]*a5+(11)f[j-4]*a4+(11)
153
        f[j-5]*a3+(l1)f[j-6]*a2+(l1)f[j-7]*a1+f[j-8])%P, j-=8;
                     while(j--) x=((11)x*a[i]+f[j])%P;
154
155
                     ans.push_back(x);
156
                 }
157
                 return;
158
             }
159
             vector<int> tmp;
160
             PolyDiv(f, p[t], tmp, tmp);
161
             Evaluate(1, (1+r)/2, t<<1, tmp, a), Evaluate((1+r)/2+1, r, t<<1|1, tmp, a);
162
163
        inline vector<int> Evaluate(const vector<int> &f, const vector<int> &a, int flag=-1)
        \{// f(a_i)
             if(flag==-1) p.resize(a.size()<<2), Evaluate Interpolate Init(0, a.size()-1, 1,</pre>
164
        a);
165
             ans.clear(), Evaluate(0, a.size()-1, 1, f, a);
166
             return ans;
167
         vector<int> Interpolate(int 1, int r, int t, const vector<int> &x, const vector<int>
168
         &f){
169
             if(1==r){
                 vector<int> ans;
170
171
                 return ans.push_back(f[1]), ans;
172
             int mid=(l+r)/2, k=t<<1;</pre>
173
```

```
174
             return Add(Mul(Interpolate(l, mid, k, x, f), p[k|1]), Mul(Interpolate(mid+1, r,
        k|1, x, f), p[k]);
175
        }
176
         inline vector<int> Interpolate(const vector<int> &x, const vector<int> &y)\{//(x_i,
        y_{\perp}i)
177
             int n=x.size();
178
             p.resize(n<<2), Evaluate_Interpolate_Init(0, n-1, 1, x);</pre>
179
             vector<int> f=Evaluate(Derivative(p[1]), x, 0);
180
             for(int i=0; i<n; ++i) f[i]=(ll)y[i]*Pow(f[i])%P;</pre>
181
             return Interpolate(0, n-1, 1, x, f);
182
    }
183
    using namespace Poly;
184
185
    int main() {
186
         cin >> n; x.resize(n), y.resize(n);
         for(int i=0; i<n; ++i) cin >> x[i], cin >> y[i];
187
188
         cin >> m, z.resize(m);
         for(int i=0; i<m; ++i) cin >> z[i];
189
         x=Evaluate(Interpolate(x, y), z);
190
         for(int i:x) cout << i << " ";</pre>
191
192
         return 0;
193
```

#### 4.1.9 BM Alogrithm

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 #define rep(i,a,n) for (int i=a;i<n;i++)</pre>
   #define per(i,a,n) for (int i=n-1;i>=a;i--)
   #define pb push_back
6
   #define mp make_pair
7
   #define all(x) (x).begin(),(x).end()
   #define fi first
9 #define se second
#define SZ(x) ((int)(x).size())
11 typedef vector<int> VI;
12 typedef long long 11;
13 typedef pair<int,int> PII;
14 const 11 mod=1000000007;
  ll powmod(ll a,ll b) {ll res=1;a%=mod; assert(b>=0); for(;b;b>>=1){if(b&1)res=res*a%mod;
       a=a*a%mod;}return res;}
16
   // head
17
   namespace linear_seq {
18
        const int N=10010;
19
        11 res[N],base[N],_c[N],_md[N];
20
21
        vector<int> Md;
22
        void mul(ll *a,ll *b,int k) {
            rep(i,0,k+k) _c[i]=0;
23
24
            rep(i,0,k) if (a[i]) rep(j,0,k) _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;
25
            for (int i=k+k-1;i>=k;i--) if (_c[i])
26
                rep(j,0,SZ(Md))
                                 _c[i-k+Md[j]]=(_c[i-k+Md[j]]-_c[i]*_md[Md[j]])%mod;
27
            rep(i,0,k) a[i]=_c[i];
28
        int solve(ll n,VI a,VI b) { // a 系数 b 初值 b[n+1]=a[0]*b[n]+...
29
              printf("%d \mid n", SZ(b));
30
31
            11 ans=0,pnt=0;
32
            int k=SZ(a);
33
            assert(SZ(a)==SZ(b));
34
            rep(i,0,k) _md[k-1-i]=-a[i];_md[k]=1;
```

```
35
            Md.clear();
36
             rep(i,0,k) if (_md[i]!=0) Md.push_back(i);
37
             rep(i,0,k) res[i]=base[i]=0;
38
             res[0]=1;
39
             while ((111<<pnt)<=n) pnt++;</pre>
40
             for (int p=pnt;p>=0;p--) {
                 mul(res,res,k);
41
42
                 if ((n>>p)&1) {
43
                     for (int i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
44
                     rep(j,0,SZ(Md)) res[Md[j]]=(res[Md[j]]-res[k]*_md[Md[j]])%mod;
45
                 }
             }
46
             rep(i,0,k) ans=(ans+res[i]*b[i])%mod;
47
48
             if (ans<0) ans+=mod;</pre>
49
             return ans;
50
        VI BM(VI s) {
51
52
            VI C(1,1), B(1,1);
53
             int L=0, m=1, b=1;
54
             rep(n,0,SZ(s)) {
55
                 11 d=0;
                 rep(i,0,L+1) d=(d+(l1)C[i]*s[n-i])%mod;
56
57
                 if (d==0) ++m;
                 else if (2*L<=n) {</pre>
58
59
                     VI T=C;
                     11 c=mod-d*powmod(b,mod-2)%mod;
60
61
                     while (SZ(C) < SZ(B) + m) C.pb(0);
62
                     rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
63
                     L=n+1-L; B=T; b=d; m=1;
64
                 } else {
                     11 c=mod-d*powmod(b,mod-2)%mod;
65
66
                     while (SZ(C) < SZ(B) + m) C.pb(0);
67
                     rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
68
                     ++m;
69
                 }
70
             }
71
             return C;
72
        }
73
        int gao(VI a,ll n) {
74
            VI c=BM(a);
75
             c.erase(c.begin());
76
             rep(i,0,SZ(c)) c[i]=(mod-c[i])%mod;
77
             return solve(n,c,VI(a.begin(),a.begin()+SZ(c)));
78
        }
79
    };
80
    int main() {
81
        while (~scanf("%d",&n)) {
82
83
             vector<int>v;
             v.push_back(1);
84
85
            v.push_back(2);
86
            v.push_back(4);
87
            v.push_back(7);
88
            v.push_back(13);
89
            v.push_back(24);
90
             //VI\{1,2,4,7,13,24\}
             printf("%d\n",linear_seq::gao(v,n-1));
91
92
        }
93
    }
```

# 4.2 Math Theory

#### **4.2.1** Inverse

```
//O(logn) 求n的 逆元
   const int mod = 1e6 + 3;
   int exgcd(int a, int b, int &x, int &y) {
4
        int d = a;
5
        if(b != 0) {
            d = exgcd(b, a % b, y, x);
6
7
            y -= (a / b) * x;
8
9
        else {
10
           x = 1; y = 0;
11
12
       return d;
13
   }
14
   int inverse(int a) {
15
       int x, y;
16
        exgcd(a, mod, x, y);
        return (x % mod + mod) % mod;
17
18
   int inverse(int a) { return qpow(a, mod - 2); }
19
20
   //O(n)求 1\sim n的 逆 元
21 int inv[MAXN];
   void init() {
22
23
        inv[0] = inv[1] = 1;
       for(int i = 2; i < MAXN; i++) inv[i] = (long long)(mod - mod / i) * inv[mod % i] %</pre>
24
25
   }
```

## 4.2.2 Lucas

```
//mod很小可以预处理逆元的情况
1
2
   void init() {
       fac[0] = 1;
3
        for(int i = 1; i < mod; i++) fac[i] = (long long)fac[i - 1] * i % mod;</pre>
4
5
        inv[0] = inv[1] = 1;
6
       for(int i = 2; i < mod; i++) inv[i] = (long long)(mod - mod / i) * inv[mod % i] %</pre>
       mod;
7
       for(int i = 1; i < mod; i++) inv[i] = (long long)inv[i] * inv[i - 1] % mod;</pre>
8
9
   int C(int a, int b) {
       if(b > a) return 0;
10
       if(a < mod) return (long long)fac[a] * inv[b] % mod * inv[a - b] % mod;</pre>
11
        return (long long)C(a / mod, b / mod) * C(a % mod, b % mod) % mod;
12
13
   }
   //mod过大不能预处理逆元的情况
14
   LL qpow(LL x, LL y) {
15
       LL res = 1;
16
17
        while(y) {
18
           if(y \& 1) res = res * x % mod;
19
           x = x * x % mod;
20
           y >>= 1;
21
       }
22
       return res;
23
   }
24
   LL C(LL a, LL b) {
25
   if(b > a) return 0;
```

```
26
       if(b > a - b) b = a - b;
27
        LL s1 = 1, s2 = 1;
28
        for(LL i = 0; i < b; i++) {
29
            s1 = s1 * (a - i) % mod;
            s2 = s2 * (i + 1) % mod;
30
31
        }
32
        return s1 * qpow(s2, mod - 2) % mod;
33
34
   LL lucas(LL a, LL b) {
        if(a < mod) return C(a, b);</pre>
35
        return lucas(a / mod, b / mod) * C(a % mod, b % mod);
36
37
```

#### 4.2.3 CRT && exCRT

 $x \equiv a_i \pmod{m_i}$ 

```
1
   namespace CRT {
2
        LL m[MAXN], a[MAXN];
3
        LL exgcd(LL _a, LL _b, LL &x, LL &y) {
4
            if(!_b) {
                x = 1; y = 0;
5
                return _a;
6
7
            LL d = exgcd(_b, _a % _b, y, x);
8
            y -= (_a / _b) * x;
9
            return d;
10
11
        LL crt(int n) {
12
            LL M = 1, tmp, res = 0, x, y;
13
            for(int i = 1; i <= n; i++) M *= m[i];</pre>
14
            for(int i = 1; i <= n; i++) {</pre>
15
16
                tmp = M / m[i];
17
                exgcd(tmp, m[i], x, y);
18
                x = (x + m[i]) % m[i];
19
                res = (a[i] * x % M * tmp % M + res) % M;
20
            }
21
            return res;
22
        }
23
   }
24
   namespace EXCRT {
25
        LL m[MAXN], a[MAXN];
26
        LL exgcd(LL _a, LL _b, LL &x, LL &y) {
27
            if(!_b) {
28
                x = 1; y = 0;
                return _a;
29
30
31
            LL d = exgcd(_b, _a % _b, y, x);
            y -= (_a / _b) * x;
32
33
            return d;
34
        LL excrt(int n) {
35
            LL M = m[1], A = a[1], x, y, d, tmp;
36
37
            for(int i = 2; i <= n; i++) {</pre>
38
                d = exgcd(M, m[i], x, y);
39
                if((A - a[i]) % d) return -1; //No solution
                tmp = M / d; M *= m[i] / d;
40
                y = (A - a[i]) / d % M * y % M;
41
```

```
42
                y = (y + tmp) \% tmp;
43
                A = (m[i] \% M * y \% M + a[i]) \% M;
44
                A = (A + M) \% M;
45
            }
46
            return A;
47
        LL inv(LL _a, LL _b) {
48
            LL x, y;
49
50
            exgcd(_a, _b, x, y);
            return (x % _b + _b) % _b;
51
52
        LL excrt(int n) {
53
            LL M = m[1], A = a[1], x, y, d, c, tmp;
54
            for(int i = 2; i <= n; i++) {</pre>
55
56
                d = exgcd(M, m[i], x, y);
57
                c = a[i] - A;
                if(c % d) return -1;
58
                c = (c \% m[i] + m[i]) \% m[i];
59
60
                M /= d; m[i] /= d;
61
                c = c / d * inv(M % m[i], m[i]) % m[i];
62
                tmp = M;
                M *= m[i] * d;
63
                A = (c * tmp % M * d % M + A) % M;
64
65
            }
66
            return A;
67
        // = a[i] 一开始就是负数时,转成正数a[i] = (a[i] \% m[i] + m[i]) \% m[i];
68
69
```

### 4.2.4 BSGS

```
const int MOD = 76543;
   int hs[MOD + 5], head[MOD + 5], nxt[MOD + 5], id[MOD + 5], ecnt;
2
3
    void insert(int x, int y) {
        int k = x \% MOD;
4
5
        hs[ecnt] = x, id[ecnt] = y, nxt[ecnt] = head[k], head[k] = ecnt++;
6
    }
7
    int find(int x) {
8
        int k = x \% MOD;
9
        for(int i = head[k]; i; i = nxt[i])
10
            if(hs[i] == x)
11
                return id[i];
12
        return -1;
13
   }
    int BSGS(int a, int b, int c){
14
        memset(head, 0, sizeof head); ecnt = 1;
15
16
        if(b == 1) return 0;
17
        int m = sqrt(c * 1.0), j;
18
        LL x = 1, p = 1;
        for(int i = 0; i < m; i++, p = p * a % c)</pre>
19
20
            insert(p * b % c, i);
        for(LL i = m; ;i += m){
21
22
            if((j = find(x = x * p % c)) != -1) return i - j;
23
            if(i > c) break;
24
25
        return -1;
26
   }
```

# 4.2.5 Miller-Rabin && PollardRho

```
LL ksc(LL a, LL n, LL mod){
1
2
        LL ret=0;
3
        for(;n;n>>=1){
4
            if(n&1){ret+=a;if(ret>=mod)ret-=mod;}
5
            a<<=1; if(a>=mod)a-=mod;
6
7
        return ret;
   }
8
   LL ksm(LL a,LL n,LL mod){
9
10
        LL ret = 1;
11
        for(;n;n>>=1){
            if(n&1)ret=ksc(ret,a,mod);
12
13
            a=ksc(a,a,mod);
14
        }
15
        return ret;
16
   }
17
    int millerRabin(LL n){
18
        if(n<2 || (n!=2 && !(n&1)))return 0;
19
        LL d=n-1; for(;!(d&1);d>>=1);
20
        for(int i=0;i<20;++i){</pre>
21
             LL a=rand()%(n-1)+1;
22
             LL t=d,m=ksm(a,d,n);
23
             for(;t!=n-1 && m!=1 && m!=n-1;m=ksc(m,m,n),t<<=1);
24
             if(m!=n-1 && !(t&1)) return 0;
25
        }
26
        return 1;
27
28
   LL cnt, fact[100];
29
   LL gcd(LL a, LL b) { return !b?a:gcd(b,a%b);}
30
   LL pollardRho(LL n, int a){
31
        LL x=rand()%n, y=x, d=1, k=0, i=1;
        while(d==1){
32
33
            ++k;
            x=ksc(x,x,n)+a; if(x>=n)x-=n;
34
35
            d=gcd(x>y?x-y:y-x,n);
36
            if(k==i){y=x;i<<=1;}</pre>
37
        if(d==n)return pollardRho(n,a+1);
38
39
        return d;
40
   }
41
    void findfac(LL n){
42
        if(millerRabin(n)){fact[++cnt]=n;return;}
43
        LL p=pollardRho(n,rand()%(n-1)+1);
44
        findfac(p);
45
        findfac(n/p);
46
   }
```

# **4.2.6** $\varphi(n)$

```
int phi(int x) {
   int res = x;
   for(int i = 2; i * i <= x; i++) {
      if(x % i == 0) {
        res = res / i * (i - 1);
        while(x % i == 0) x /= i;
    }
}</pre>
```

```
9    if(x > 1) res = res / x * (x - 1);
10    return res;
11 }
```

### 4.2.7 Euler Sieve

```
1
   int prime[MAXN], cnt, phi[MAXN], mu[MAXN];
   bool isp[MAXN];
3
                         //最小质因子最高次幂
   int min_pow[MAXN];
                         //1+p+p^2+\ldots+p^k
5
   int min_sum[MAXN];
   int div_sum[MAXN];
6
                         //约数和
   int min_index[MAXN]; //最小质因子的指数
8
                          //约数个数
   int div_num[MAXN];
9
   void Euler(int n) {
10
        mu[1] = phi[1] = div_num[1] = div_sum[1] = 1;
11
        for(int i = 2; i <= n; i++) {</pre>
12
13
            if(!isp[i]) {
                prime[++cnt] = min_pow[i] = i;
14
15
                phi[i] = i - 1;
16
                mu[i] = -1;
17
                min_index[i] = 1; div_num[i] = 2;
18
                div_sum[i] = min_sum[i] = i + 1;
19
            }
            for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
20
21
                isp[i * prime[j]] = 1;
22
                if(i % prime[j] == 0) {
                    phi[i * prime[j]] = phi[i] * prime[j];
23
24
                    mu[i * prime[j]] = 0;
25
26
                    min_index[i * prime[j]] = min_index[i] + 1;
                    div_num[i * prime[j]] = div_num[i] / (min_index[i] + 1) * (min_index[i *
27
        prime[j]] + 1);
28
                    min_sum[i * prime[j]] = min_sum[i] + min_pow[i] * prime[j];
29
                    div_sum[i * prime[j]] = div_sum[i] / min_sum[i] * min_sum[i * prime[j]];
30
                    min_pow[i * prime[j]] = min_pow[i] * prime[j];
31
32
                    break:
33
                phi[i * prime[j]] = phi[i] * (prime[j] - 1);
34
35
                mu[i * prime[j]] = -mu[i];
36
                div_num[i * prime[j]] = div_num[i] << 1;</pre>
37
                min_index[i * prime[j]] = 1;
38
39
                div_sum[i * prime[j]] = div_sum[i] * (prime[j] + 1);
40
41
                min_pow[i * prime[j]] = prime[j];
42
                min_sum[i * prime[j]] = prime[j] + 1;
43
            }
44
        }
45
   }
```

### 4.2.8 DuJiao Sieve

$$\sum_{i=1}^{n} \varphi(i)$$

```
1 vector<int> prime;
2
   int phi[MAXN], P[MAXN];
3 bool isp[MAXN];
4
   unordered_map<LL, int> mp;
5
   void Euler(int n) {
6
        phi[1] = 1;
7
        for(int i = 2; i <= n; i++) {</pre>
8
            if(!isp[i]) {
9
                 prime.push_back(i);
10
                 phi[i] = i - 1;
11
            for(auto x : prime) {
12
                 if(i * x > n) break;
13
                 isp[i * x] = 1;
14
                 if(i % x == 0) {
15
                     phi[i * x] = phi[i] * x;
16
                     break;
17
18
19
                 phi[i * x] = phi[i] * (x - 1);
20
            }
21
22
        for(int i = 1; i \le n; i++) P[i] = (P[i - 1] + phi[i]) % mod;
23
   }
24
   LL cal(LL n) {
25
        if(n < MAXN) return P[n];</pre>
26
        if(mp.count(n)) return mp[n];
27
        LL res = 0;
28
        for(LL i = 2, last; i <= n; i = last + 1) {</pre>
            last = n / (n / i);
29
30
            res += (last - i + 1) % mod * cal(n / i) % mod;
31
            res %= mod;
32
        mp[n] = ((\_int128)n * (n + 1) / 2 % mod + mod - res) % mod;
33
34
        return mp[n];
35
```

$$\sum_{i=1}^{n} \mu(i)$$

```
1
    LL cal(LL n) {
2
        if(n < MAXN) return M[n];</pre>
3
        if(mp.count(n)) return mp[n];
4
        LL res = 0;
5
        for(LL i = 2, last; i <= n; i = last + 1) {</pre>
6
            last = n / (n / i);
7
             res += (last - i + 1) * cal(n / i);
8
9
        mp[n] = 1 - res;
10
        return 1 - res;
11
   }
```

# 4.2.9 Min\_25 Sieve

思路为把结果分为 i 为质数的和,i 为合数的和,i=1 的和 g(n,j) 表示从 1 累加到 n 的 f(i),其中的 i 满足要么 i 自己是质数,要么 i 的最小质因子大于第 j 个质数 要求  $f(p)(p\ is\ prime)$  可被多项式表示, $f(p^k)$  可快速计算 分别计算多项式  $x^0, x^1, x^2$  的 g(n,j) 和 h(n,j) h 由欧拉筛计算,递归起点 g(n,0)

公式

$$\sum_{i} f(i)$$
 
$$g(n,j) = \begin{cases} g(n,j-1) & p_j^2 > n \\ \\ g(n,j-1) - f(p_j)[g(\frac{n}{p_j},j-1) - g(p_j-1,j-1)] & p_j^2 \le n \end{cases}$$

其中

$$g(p_j - 1, j - 1) = \sum_{i=1}^{j-1} f(p_i)$$

可用 h 表示

S(n,j) 表示从 1 累加到 n 的 f(i),同样的, i 满足要么 i 自己是质数, 要么 i 的最小质因子大于第 j 个质数和 g 顺序相反, 最后答案为 S(n,1),递推式为

$$S(n,j) = g(n,|P|) - \sum_{i=1}^{j} f(p_i) + \sum_{k \ge j} \sum_{\substack{n^{e+1} < n}} (f(p_k^e)S(\frac{n}{p_k^e}, k+1) + f(p_k^{e+1}))$$

前一部分  $g(n, |P|) - \sum_{i=1}^{j} f(p_i)$  是相应的质数部分和。例题: 欧拉函数前缀和

$$\sum_{i=1}^{n} \varphi(i)$$

 $g_{k,n}$  and  $h_{k,n}$  Count

$$\sum_{i=1}^{n} i^k$$

```
#include <bits/stdc++.h>
2 using namespace std;
3 typedef long long LL;
4 const int MAXN = 1e6 + 5, mod = 1e9 + 7;
5 const int inv2 = (mod + 1) / 2, inv6 = (mod + 1) / 6;
6 int prime[MAXN], isp[MAXN], cnt;
   LL g[3][MAXN << 1], h[3][MAXN << 1];
8 LL w[MAXN << 1];</pre>
9 int id1[MAXN], id2[MAXN];
inline int MOD(LL x) { return x >= mod ? x - mod : x; }
   //inline int MOD(LL x) \{ return x \% mod; \}
11
   inline int add(LL x, LL y) { return MOD(MOD(x) + MOD(y)); }
   void Euler(int n) {
        for(int i = 2; i <= n; i++) {</pre>
14
15
            if(!isp[i]) {
                prime[++cnt] = i;
16
17
                h[0][cnt] = h[0][cnt - 1] + 1;
                h[1][cnt] = add(h[1][cnt - 1], i);
18
                h[2][cnt] = add(h[2][cnt - 1], (LL)i * i % mod);
19
20
            for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
21
                isp[i * prime[j]] = 1;
22
23
                if(i % prime[j] == 0) {
24
                    break;
25
                }
26
```

```
}
27
28
   }
29
   LL n;
30
   int sz, m;
31
    inline int id(LL x) {
        return x <= sz ? id1[x] : id2[n / x];</pre>
32
33
    //f(p \hat{k})
34
    inline int f(int p, LL pk) {
35
        return pk / p * (p - 1) % mod;
36
37
    LL S(LL x, int y) {
38
39
        if(x \leftarrow 1 \mid | prime[y] > x) return 0;
        //G(x) - H(j-1) (first part)
40
41
        LL res = add(add(g[1][id(x)], mod - g[0][id(x)]), mod - add(h[1][y - 1], mod - h[0][id(x)])
        y - 1]));
        for(int j = y, k = 1; j <= cnt && (LL)prime[j] * prime[j] <= x; j++, k = 1) {
42
            for(LL pk = prime[j]; pk * prime[j] <= x; pk *= prime[j], k++) {
43
                 res = add(res, S(x / pk, j + 1) * f(prime[j], pk) % mod + f(prime[j], pk *
44
        prime[j]));
45
46
        return res;
47
48
    }
49
    int main() {
50
        ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
51
        cin >> n;
52
        sz = sqrt(n);
53
        Euler(sz);
        for(LL i = 1, last, t; i <= n; i = last + 1) {</pre>
54
            last = n / (n / i);
55
            w[++m] = n / i, t = n / i \% mod;
56
            w[m] \leftarrow sz ? id1[w[m]] = m : id2[last] = m;
57
            g[0][m] = MOD(t + mod - 1);
58
            g[1][m] = add(t * (t + 1) % mod * inv2 % mod, mod - 1);
59
            g[2][m] = add((2 * t + 1) % mod * t * (t + 1) % mod * inv6 % mod, mod - 1);
60
61
        for(int j = 1; j <= cnt; j++) {</pre>
62
            for(int i = 1; i <= m && (LL)prime[j] * prime[j] <= w[i]; i++) {</pre>
63
64
                 g[0][i] = MOD(g[0][i] + mod - (g[0][id(w[i] / prime[j])] - h[0][j - 1]));
65
                 g[1][i] = MOD(g[1][i] + mod - ((LL)prime[j]) * MOD(g[1][id(w[i] / prime[j]))]
        + mod - h[1][j - 1]) % mod));
                 g[2][i] = MOD(g[2][i] + mod - ((LL)prime[j] * prime[j] % mod * MOD(g[2][id(w)])
66
        [i] / prime[j])] + mod - h[2][j - 1]) % mod));
67
            }
68
        }
        //S(n, 1) + F(1);
69
        LL ans = MOD(S(n, 1) + 1);
70
        cout << ans << endl;</pre>
71
72
        return 0;
73
```

$$\sum_{i=1}^{n} \mu(i)$$

```
#include <bits/stdc++.h>
using namespace std;

typedef long long LL;
const int MAXN = 1e6 + 5;
int prime[MAXN], isp[MAXN], cnt;
LL g[3][MAXN << 1], h[3][MAXN << 1];</pre>
```

```
LL w[MAXN << 1];
8
    int id1[MAXN], id2[MAXN];
9
    void Euler(int n) {
        for(int i = 2; i <= n; i++) {</pre>
10
11
             if(!isp[i]) {
12
                 prime[++cnt] = i;
13
                 h[0][cnt] = h[0][cnt - 1] + 1;
14
             for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
15
                 isp[i * prime[j]] = 1;
16
                 if(i % prime[j] == 0) {
17
18
                     break;
19
20
            }
21
        }
22
   }
23 LL a, b;
24 LL n;
   int sz, m;
26
   inline int id(LL x) {
27
        return x <= sz ? id1[x] : id2[n / x];</pre>
28
   }
29
    //f(p \hat{k})
   inline int f(int p, int k) {
30
31
        return k == 1 ? -1 : 0;
32
33
    LL S(LL x, int y) {
        if(x \leftarrow 1 || prime[y] \rightarrow x) return 0;
34
35
        //g(x) - h(j-1)
36
        LL res = -g[0][id(x)] + h[0][y - 1];
        for(int j = y, k = 1; j <= cnt && (LL)prime[j] * prime[j] <= x; j++, k = 1) {
37
            for(LL pk = prime[j]; pk * prime[j] <= x; pk *= prime[j], k++) {</pre>
38
                 res += S(x / pk, j + 1) * f(prime[j], k) + f(prime[j], k + 1);
39
            }
40
        }
41
        return res;
42
43
   }
   LL cal(LL x) {
44
45
        n = x;
46
        m = 0;
47
        sz = sqrt(n);
48
        for(LL i = 1, last, t; i <= n; i = last + 1) {
49
             last = n / (n / i);
50
            w[++m] = n / i, t = n / i;
51
            w[m] \le sz ? id1[w[m]] = m : id2[last] = m;
52
            g[0][m] = t - 1;
53
54
        for(int j = 1; j <= cnt; j++) {</pre>
             for(int i = 1; i <= m && (LL)prime[j] * prime[j] <= w[i]; i++) {</pre>
55
56
                 g[0][i] = g[0][i] - (g[0][id(w[i] / prime[j])] - h[0][j - 1]);
57
58
59
        return S(x, 1) + 1;
60
61
    int main() {
        ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
62
        cin >> a >> b;
63
64
        Euler(sqrt(b));
65
        //S(n, 1) + F(1);
66
        cout << cal(b) - cal(a - 1) << endl;</pre>
```

```
68 return 0;
69 }
```

#### 4.2.10 Möbius Inversion

$$\sum_{i=1}^{n} \sum_{j=1}^{m} lcm(i,j) \pmod{p}$$

```
int mu[MAXN], prime[MAXN], sum[MAXN], cnt;
1
2
    bool isp[MAXN];
    void getmu(int n) {
3
4
        mu[1] = 1;
5
        for(int i = 2; i <= n; i++) {
6
            if(!isp[i]) {
7
                 mu[i] = -1;
                 prime[++cnt] = i;
8
9
10
            for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
                 isp[i * prime[j]] = 1;
11
                 if(i % prime[j] == 0) {
12
                     mu[i * prime[j]] = 0;
13
14
                     break;
15
16
                 mu[i * prime[j]] = -mu[i];
17
            }
        }
18
19
   }
20
   ll n, m, ans;
21
   ll query(ll x, ll y) { return (x * (x + 1) / 2 % mod) * (y * (y + 1) / 2 % mod) % mod; }
22
   11 F(11 x, 11 y) {
23
        11 \text{ res} = 0, \text{ last};
24
        for(ll i = 1; i \leftarrow min(x, y); i = last + 1) {
25
            last = min(x / (x / i), y / (y / i));
26
            res = (res + (sum[last] - sum[i - 1]) * query(x / i, y / i) % mod) % mod;
27
        }
28
        return res;
   }
29
30
   int main() {
31
        cin>>n>>m;
32
        getmu(min(n, m));
        for(11 i = 1; i \le min(n, m); i++) sum[i] = (sum[i - 1] + (i * i * mu[i]) % mod) %
33
        mod;
34
        ll last;
35
        for(ll d = 1; d <= min(n, m); d = last + 1) {</pre>
36
            last = min(n / (n / d), m / (m / d));
37
            ans = (ans + (last - d + 1) * (d + last) / 2 % mod * F(n / d, m / d) % mod) %
        mod;
38
        }
39
        ans = (ans + mod) \% mod;
40
        cout<<ans<<endl;</pre>
41
        return 0;
42
```

# 4.2.11 Primitive Root

简易版如果模 m 有原根,则有  $\varphi(\varphi(m))$  个原根当 d 遍历  $\varphi(m)$  的简化剩余系时, $g^d$  遍历完 m 的全部原根

```
const int MAXN = ;
1
    LL qpow(LL a, LL b, LL mod) {
2
3
        LL res = 1;
4
        for (;b; b >>= 1) {
5
             if(b&1) res = res * a % mod;
6
            a = a * a % mod;
7
        }
8
        return res;
9
10
    LL phi(LL x) {
11
        LL res = x;
        for(int i = 2; i <= x / i; i++) {</pre>
12
            if(x % i == 0) {
13
                 res = res / i * (i - 1);
14
15
                 while(x % i == 0) x /= i;
16
            }
17
        if(x > 1) res = res / x * (x - 1);
18
19
        return res;
20
21
    int has_primitive_root(int p) {
22
        if (p == 4) return 1;
        if (p % 2 == 0) p /= 2;
23
24
        if (p % 2 == 0) return false;
25
        for (int i = 2; i <= p/i; i++)
26
        if (p % i == 0) {
27
            while(p % i == 0) p /= i;
28
            return p == 1 ? i : 0;
29
        }
30
        return p;
31
32
    //int indg[MAXN];
33
    int get_g_init(LL p) {
        //p : 2 or 4 or (p**n) or (2 * p**n); else return -1;
34
        if(p == 2) {puts("1");return 1;}
35
        if(p == 4) {puts("3"); return 3;}
36
        if(!has_primitive_root(p)) {puts("-1"); return -1;}
37
38
        vector<int> p_fact;
39
        LL p_{phi} = phi(p), tp = p_{phi};
40
        for(int i = 2, in = sqrt(p) + 0.5; i \le in; i++)
41
        if(tp % i == 0) {
42
             p_fact.push_back(p_phi / i);
43
            while(tp % i == 0) tp /= i;
44
45
        if(tp != 1) p_fact.push_back(p_phi / tp);
46
        int g = 1;
47
        for (bool fg = 0; !fg ;) {
            fg = 1; g++;
48
             if(qpow(g, p_phi, p) != 1) {fg = 0; continue;}
49
             for(auto it : p_fact)
50
51
                 if(qpow(g, it, p) == 1) {fg = 0; break;}
52
        //for(int \ i = 0, \ tg = 1; \ i < p\_phi; \ i++, \ tg = tg * g \% p) \ indg[tg] = i;
53
54
        /*vector < int > fac;
55
        for(int \ d = 1, \ tg = g; \ d < p\_phi; \ d++, \ tg = tg * g \% p)
             if(\underline{\phantom{a}}gcd((LL)d, p\_phi) == 1) fac.push\_back(tg);
56
57
        sort (fac.begin (), fac.end ());
        int \ ed = fac.back();
58
        for (auto it : fac)
59
60
             printf("\%d\%c", it, it! = ed?' ':' \setminus n');*/
61
        return g;
```

62 }

# 5 Geometry

# 5.1 Commonly Definition and Functions

### 5.1.1 Const and Functions

```
namespace CG{
1
2
        #define Point Vector
3
        const double pi=acos(-1.0);
4
        const double inf=1e100;
5
        const double eps=1e-9;
6
        template <typename T> inline T Abs(T x){return x>0?x:-x;}
7
        template <typename T> inline bool operator == (T x,T y){return Abs(x-y)<eps;}</pre>
8
        int sgn(double x){
9
            if (Abs(x)<eps) return 0;</pre>
10
            if (x>0) return 1;
            else return -1;
11
12
        }
13
   }
```

#### 5.1.2 Point Definition

```
1
   namespace CG{
2
       struct Point{
3
            double x,y;
4
            Point(double x=0, double y=0):x(x),y(y){}
5
6
       Vector operator + (const Vector a,const Vector b){return Vector(a.x+b.x,a.y+b.y);}
7
       Vector operator - (const Vector a,const Vector b){return Vector(a.x-b.x,a.y-b.y);}
       Vector operator * (const Vector a,const double k){return Vector(a.x*k,a.y*k);}
8
       Vector operator / (const Vector a,const double k){return Vector(a.x/k,a.y/k);}
9
10
       bool operator < (const Vector a,const Vector b) {return a.x==b.x?a.y<b.y:a.x<b.x;}</pre>
11
       bool operator == (const Vector a,const Vector b) {return a.x==b.x && a.y==b.y;}
12
       double Dot(const Vector a,const Vector b){return a.x*b.x+a.y*b.y;}
       double Cross(const Vector a,const Vector b){return a.x*b.y-a.y*b.x;}
13
       double mult_Cross(const Vector a,const Vector b,const Vector c){return (a.x-c.x)*(b.
14
       y-c.y)-(b.x-c.x)*(a.y-c.y);}
15
       double mult_Dot(const Vector a,const Vector b,const Vector c){return (a.x-c.x)*(b.x-
       c.x)+(a.y-c.y)*(b.y-c.y);}
       double Norm(const Vector a){return sqrt(Dot(a,a));}
16
       double Angle(const Vector a,const Vector b){return acos(Dot(a,b)/Norm(a)/Norm(b));}
17
18
       Vector Rotate(const Vector a, const double theta){return Vector(a.x*cos(theta)-a.y*
       sin(theta),a.x*sin(theta)+a.y*cos(theta));}
19
       bool ToLeftTest(const Vector a,const Vector b){return Cross(a,b)<0;}</pre>
       double DisPP(const Vector a,const Vector b){return sqrt((a.x-b.x)*(a.x-b.x)+(a.y-b.y
20
       )*(a.y-b.y));}
21
   }
```

# 5.1.3 Line Definition

```
namespace CG{
struct Line{
    point p0,v,p1;
    double t,theta;
    Line(Point _p0=0,Point _v=0,double _t=1):p0(_p0),v(_v),t(_t){p1=p0+v*t; theta=
    atan2(v.y,v.x);}
```

```
6
                              // Line(Point \_p0=0, Point \_v=0, double \_t=1): p0(\_p0), p1(\_v)\{v=(p1-p0)/t; theta=0, double \_t=1): p0(\_v)\{v=(p1-p0)/t; theta=0, double \_t=1): p0(\_v)\{v=(p1-p0)/t; theta=0, d
                   atan2(v.y,v.x);}
  7
                   };
 8
                    bool operator < (const Line n,const Line m) {return n.theta<m.theta;}</pre>
 9
                    Point GetIntersection(const Line n,const Line m){return n.p0+n.v*Cross(m.v,(n.p0-m.
                   p0))/Cross(n.v,m.v);}
                    bool OnLine(const Vector a,const Line 1){return Cross(1.p0-a,1.p1-a)==0;}
10
                    bool OnSegment(const Point a,const Line 1){return sgn(Cross(1.p0-a,1.p1-a))==0 &&
11
                   sgn(Dot(1.p0-a,1.p1-a))<0;}
                    double DisPL(const Point a,const Line 1){return Abs(Cross(1.p1-1.p0,a-1.p0)/Norm(1.
12
                   p1-l.p0));}
                    double DisPS(const Point a,const Line 1){
13
                              if (1.p0==1.p1) return Norm(a-1.p0);
14
                              Vector v1=l.p1-l.p0, v2=a-l.p0, v3=a-l.p1;
15
                              if (sgn(Dot(v1,v2))<0) return Norm(v2);</pre>
16
17
                              if (sgn(Dot(v1,v3))>0) return Norm(v3);
18
                              return DisPL(a,1);
19
20
                    Point GetProjection(const Point a, const Line 1){
21
                              Vector v=1.p1-l.p0;
22
                              return 1.p0+v*(Dot(v,a-1.p0)/Dot(v,v));
23
24
                    bool SegmentIntersection(const Line n,const Line m,bool p){
25
                              double c1=Cross(n.p1-n.p0,m.p1-m.p0);
                              double c2=Cross(n.p1-n.p0,m.p1-n.p0);
26
27
                              double c3=Cross(m.p1-m.p0,n.p0-m.p0);
                              double c4=Cross(m.p1-m.p0,n.p1-m.p0);
28
29
                              if (p){
30
                                          if (!sgn(c1) || !sgn(c2) || !sgn(c3) || !sgn(c4)){
31
                                                    return OnSegment(n.p0,m) || OnSegment(n.p1,m) || OnSegment(m.p0,n) ||
                   OnSegment(m.p0,m);
32
33
                                         }
                              }
34
35
                              return (sgn(c1)*sgn(c2)<0 && sgn(c3)*sgn(c4)<0);</pre>
36
                    }
37
         }
```

### 5.1.4 Get Area

```
namespace CG{
double GetArea(Point *p,int n){
         double area=Cross(p[n],p[1]);
         for (int i=2;i<=n;i++) area+=0.5*Cross(p[i-1],p[i]);
         return Abs(area);
}
</pre>
```

# 5.1.5 Get Circumference

## 5.1.6 Anticlockwise Sort

```
1
   namespace CG{
        \\p为一个凸包,只是不知其点集是否为逆时针
2
3
        void clockwise_sort(Point *p,int n){
4
            for(int i=0;i<n-2;i++){</pre>
5
                double tmp = mult_Cross(p[i+1],p[i+2],p[i]);
                if(tmp>0) return;
6
7
                else if(tmp<0){</pre>
8
                    reverse(p,p+n);
9
                    return;
10
                }
11
           }
12
       }
13
```

# 5.2 Convex Hull

## 5.2.1 Get Convex Hull

```
namespace CG{
1
2
        Point p[MAXN],s[MAXN]; // both based from 0
        int ConvexHull(Point *p,int n,Point *s){
3
4
            sort(p,p+n,cmp); //x从小到大,y从小到大;
5
            int m=0;
            for (int i=0;i<n;i++){</pre>
6
7
                for (;m>=2 && Cross(s[m-1]-s[m-2],p[i]-s[m-1])<=0;m--);</pre>
8
                s[m++]=p[i];
9
            }
10
            int k=m;
            for (int i=n-2;i;i--){
11
                for (;m>=k+1 && Cross(s[m-1]-s[m-2],p[i]-s[m-1])<=0;m--);</pre>
12
13
                s[m++]=p[i];
            }
14
15
            return m-1;
16
        }
17
   }
```

# 5.2.2 Point in Convex Hull

```
1
   namespace CG{
2
        bool PointInConvexHull(Point A){
3
            int l=1,r=tot-2,mid;
            while(l<=r){</pre>
4
5
                 mid=(l+r)>>1;
6
                 double a1=Cross(p[mid]-p[0],A-p[0]);
7
                 double a2=Cross(p[mid+1]-p[0],A-p[0]);
8
                 if(a1>=0 && a2<=0){
9
                     if(Cross(p[mid+1]-p[mid],A-p[mid])>=0) return true;
10
                     return false;
11
                 else if(a1<0) r=mid-1;</pre>
12
13
                 else l=mid+1;
14
15
            return false;
16
        }
17
   }
```

## 5.3 Minkowski Sum

```
1
                     namespace CG{
    2
                                              void Minkowski(Point *C1,int n,Point *C2,int m){
    3
                                                                      for(int i=1;i<=n;i++) s1[i]=C1[i]-C1[i-1];</pre>
    4
                                                                      for(int i=1;i<=m;i++) s2[i]=C2[i]-C2[i-1];</pre>
    5
                                                                      A[tot=1]=C1[1]+C2[1];
    6
                                                                      int p1=1,p2=1;
     7
                                                                     while (p1 \le n \& p2 \le m) ++tot, A[tot] = A[tot-1] + (s1[p1] *s2[p2] >= 0?s1[p1++]:s2[p2] + (s1[p1] *s2[p2] + (s1[p1] *s
                                           ++]);
    8
                                                                      while (p1<=n) ++tot,A[tot]=A[tot-1]+s1[p1++];</pre>
    9
                                                                     while (p2<=m) ++tot,A[tot]=A[tot-1]+s2[p2++];</pre>
10
                                                                     tot=ConvexHull(A,tot);
11
                                              }
12
                     }
```

# 5.4 Rotating Calipers

#### 5.4.1 The Diameter of Convex Hull

```
1
   namespace CG{
2
        double RotatingCalipers(Point *p,int n){
3
            double dis=0;
4
            for(int i=0,j=2;i<n;++i){</pre>
5
                 while (abs(Cross(p[i+1]-p[i],p[j]-p[i]))<abs(Cross(p[i+1]-p[i],p[j+1]-p[i]))</pre>
        ) j=(j+1)%n;
6
                 dis=max(dis,max(DisPP(p[j],p[i]),DisPP(p[j],p[i+1])));
7
            }
8
            return dis;
9
        }
10
   }
```

#### 5.4.2 The Min Distance Bewteen two Convex Hull

```
1
   namespace CG{
2
      ///点c到线段ab的最短距离
3
      double GetDist(Point a, Point b, Point c){
          if(dis(a,b)<esp) return dis(b,c); ///a,b是同一个点
4
5
          if(mult_Dot(b,c,a)<-esp) return dis(a,c); ///投影
6
          if(mult_Dot(a,c,b)<-esp) return dis(b,c);</pre>
7
          return fabs(mult_Cross(b,c,a)/dis(a,b));
8
9
      }
      ///求一条线段ab的两端点到另外一条线段bc的距离,反过来一样,共4种情况
10
      double MinDist(Point a, Point b, Point c, Point d){
11
12
          return min(min(GetDist(a,b,c),GetDist(a,b,d)),min(GetDist(c,d,a),GetDist(c,d,b))
      );
13
      }
14
      double RotatingCalipers(Point *p,int n,Point *q,int m){
15
          int yminP = 0,ymaxQ=0;
          for(int i=1;i< n;i++){ ///找到点集p组成的凸包的左下角
16
             if(p[i].y < p[yminP].y||(p[i].y = p[yminP].y)&&(p[i].x < p[yminP].x)) yminP = i;
17
18
19
          for(int i=1;i<m;i++){ ///找到点集q组成的凸包的右上角
20
             21
22
          double ans = DisPP(p[yminP],q[ymaxQ]); //距离(yminP,ymaxQ)维护为当前最小值。
```

```
23
            for(int i=0;i<n;i++){</pre>
24
                 double tmp;
25
                 while(tmp=(mult_Cross(q[ymaxQ+1],p[yminP],p[yminP+1])-mult_Cross(q[ymaxQ],p[
        yminP],p[yminP+1]))>esp)
                     ymaxQ = (ymaxQ+1)%m;
26
                 if(tmp<-esp) ans = min(ans,GetDist(p[yminP],p[yminP+1],q[ymaxQ]));</pre>
27
                 else ans=min(ans,MinDist(p[yminP],p[yminP+1],q[ymaxQ],q[ymaxQ+1]));
28
29
                 yminP = (yminP+1)%n;
30
            }
31
            return ans;
32
        }
33
   }
```

## 5.5 Half Plane Intersection

```
1
   namespace CG{
2
        void HalfPlaneIntersection(Line 1[],int n){
3
            deque <Point> p;
4
            sort(l+1,l+1+n);
5
            deque <Line> q;
6
            q.push_back(l[1]);
7
            for (int i=2;i<=n;i++){</pre>
8
                for (;!p.empty() && !ToLeftTest(p.back()-l[i].p0,l[i].v);q.pop_back(),p.
       pop_back());
9
                for (;!p.empty() && !ToLeftTest(p.front()-1[i].p0,1[i].v);q.pop_front(),p.
       pop front());
10
                if (sgn(Cross(l[i].v,q.back().v))==0)
                     if (ToLeftTest(l[i].p0-q.back().p0),q.back().v){
11
12
                         q.pop back();
13
                         if (!p.empty()) p.pop_back();
14
15
                if (!q.empty()) p.push_back(GetIntersection(q.back(),1[i]));
16
                q.push_back(1[i]);
17
            }
            for (;!p.empty() && !ToLeftTest(p.back()-q.front().p0,q.front().v);q.pop_back(),
18
       p.pop_back());
            p.push_back(GetIntersection(q.back(),q.front()));
19
            if (p.size() < 3) printf("0\n");</pre>
20
21
            else{
                cerr << "!" << endl;</pre>
22
                double area = 0.5 * Cross(p.back(), p.front());
23
24
                Point last = p.front();
                for (p.pop_front(); !p.empty(); last = p.front(), p.pop_front())
25
                     area += 0.5 * Cross(last, p.front());
26
27
                printf("%.8lf\n", fabs(area));
28
            }
29
        }
30
   }
```

#### 5.6 Min Circle Cover

```
namespace CG{
Point GetCircleCenter(const Point a,const Point b,const Point c){
Point p=(a+b)/2.0,q=(a+c)/2.0;
Vector v=Rotate(b-a,pi/2.0),w=Rotate(c-a,pi/2.0);
if (sgn(Norm(Cross(v,w)))==0){
    if (sgn(Norm(a-b)+Norm(b-c)-Norm(a-c))==0) return (a+c)/2;
    if (sgn(Norm(b-a)+Norm(a-c)-Norm(b-c))==0) return (b+c)/2;
```

```
8
                 if (sgn(Norm(a-c)+Norm(c-b)-Norm(a-b))==0) return (a+c)/2;
9
             }
10
             return GetIntersection(Line(p,v),Line(q,w));
11
12
        void MinCircleCover(Point p[],int n){
13
             random_shuffle(p+1,p+1+n);
             Point c=p[1];
14
15
             double r=0;
16
             for (int i=2;i<=n;i++)</pre>
17
                 if (sgn(Norm(c-p[i])-r)>0){
18
                     c=p[i],r=0;
                     for (int j=1;j<i;j++)</pre>
19
20
                          if (sgn(Norm(c-p[j])-r)>0){
21
                              c=(p[i]+p[j])/2.0;
22
                              r=Norm(c-p[i]);
23
                              for (int k=1;k<j;k++)</pre>
24
                                   if (sgn(Norm(c-p[k])-r)>0){
                                       c=GetCircleCenter(p[i],p[j],p[k]);
25
26
                                       r=Norm(c-p[i]);
27
                                   }
28
                          }
29
             printf("%.10f\n%.10f %.10f",r,c.x,c.y);
30
31
        }
32
    }
```

## 5.7 Circle Union Area

```
//k次覆盖
1
   //圆并去重后s[0]
3 typedef pair<double, int> P;
4 const double pi = acos(-1.0);
5 const int MAXN = 10003;
6 P arc[MAXN << 1];</pre>
7
   int acnt, cnt;
8
   double s[1003];
9
   bool del[1003];
10
   void add(double st, double en) {
11
        if(st < -pi) {
12
            add(st + 2 * pi, pi);
13
            add(-pi, en);
14
            return;
15
16
        if(en > pi) {
17
            add(st, pi);
            add(-pi, en - 2 * pi);
18
19
            return;
20
21
        arc[++acnt] = P(st, 1);
22
        arc[++acnt] = P(en, -1);
23
   }
24
   double F(double x) {
25
        return (x - \sin(x)) / 2;
26
   }
27
   struct Node {
28
        int x, y, r;
29
        Node(int _x = 0, int _y = 0, int _r = 0):x(_x), y(_y), r(_r) {}
30
        bool operator == (const Node& t) {
31
           return x == t.x && y == t.y && r == t.r;
```

```
32
33
        inline void read() {
34
            scanf("%d%d%d", &x, &y, &r);
35
    }a[1003];
36
37
    int main() {
        int n;
38
        scanf("%d", &n);
39
        for(int i = 1; i <= n; i++) a[i].read();</pre>
40
41
        //去重
42
        int nn = 0;
43
        for(int \ i = 1; \ i <= n; \ i++)  {
44
             bool\ same = 0;
45
46
             for(int \ j = 1; \ j < i; \ j++) \ \{
                 if(a/i) == a/j) {
47
                     same = 1; break;
48
49
50
51
             if(!same) \ a/++nn/ = a/i/;
52
53
        n = nn;
        //去包含
54
        for(int \ i = 1; \ i \ll n; \ i++)
55
56
             for(int j = 1; j \le n; j++) if(i != j) {
57
                 if(hypot(a[i].x - a[j].x, a[i].y - a[j].y) < (double)(a[i].r - a[j].r)) \ del[i].y = (double)(a[i].r - a[j].r)
        j = 1;
58
59
60
        nn = 0;
61
        for(int \ i = 1; \ i \le n; \ i++) \ if(!del[i]) 
62
            a/++nn/ = a/i/;
63
64
        n = nn;
65
        for(int i = 1; i <= n; i++) {</pre>
66
67
             acnt = 0;
68
             for(int j = 1; j <= n; j++) if(i != j) {</pre>
                 int dis = (a[i].x - a[j].x) * (a[i].x - a[j].x) + (a[i].y - a[j].y) * (a[i].
69
        y - a[j].y);
70
                 if(a[j].r > a[i].r && dis <= (a[j].r - a[i].r) * (a[j].r - a[i].r)) add(-pi,
         pi);
71
                 else if(dis > (a[i].r - a[j].r) * (a[i].r - a[j].r) && dis < (a[i].r + a[j].
        r) * (a[i].r + a[j].r)){
72
                     double c = sqrt(dis);
                     double angle = acos((a[i].r * a[i].r + c * c - a[j].r * a[j].r) / (2 * a
73
        [i].r * c));
74
                     double k = atan2(a[j].y - a[i].y, a[j].x - a[i].x);
75
                     add(k - angle, k + angle);
76
                 }
77
             }
78
             arc[++acnt] = P(pi, -1);
79
             sort(arc + 1, arc + acnt + 1);
             cnt = 0;
80
81
             double last = -pi;
            for(int j = 1; j <= acnt; j++) {</pre>
82
                 s[cnt] += F(arc[j].first - last) * a[i].r * a[i].r; //扇形 - 三角形
83
84
                 double xa = a[i].x + a[i].r * cos(last);
                 double ya = a[i].y + a[i].r * sin(last);
85
86
                 last = arc[j].first;
87
                 double xb = a[i].x + a[i].r * cos(last);
```

```
88
                double yb = a[i].y + a[i].r * sin(last);
89
                s[cnt] += (xa * yb - xb * ya) / 2; //到圆心的三角形面积
90
                cnt += arc[j].second;
91
            }
92
        //printf("\%.3f | n", s[0]);
93
        for (int i = 0; i < n; i++) {</pre>
94
95
            printf("[%d] = %.3f\n", i + 1, s[i] - s[i + 1]);
96
97
        return 0;
98
```

## 5.8 Simpson Integrate

```
double Simpson(double 1,double r){
1
       return (r-1)*(F(1)+4*F((1+r)/2)+F(r))/6;
2
3
4
   double Integrate(double 1,double r,double S){
5
        double mid=(1+r)/2;
6
        double A=Simpson(1,mid);
7
        double B=Simpson(mid,r);
8
        if(A+B-S<eps)return S;</pre>
9
        return Integrate(1,mid,A)+Integrate(mid,r,B);
10
   }
```

#### 5.9 Closest Point

```
vector <Point> P;
1
2
    DB CP(int 1, int r) {
3
        if (1 == r) return DB_INF;
4
        if (1 + 1 == r) return DisPP(P[1], P[r]);
5
        int mid = (1 + r) >> 1;
6
        DB d = min(CP(1, mid), CP(mid + 1, r));
7
        vector <Point> tmp;
8
        for (int i = 1; i <= r; i++)
9
            if (fabs(P[mid].x - P[i].x) < d) tmp.push_back(P[i]);</pre>
10
        sort(tmp.begin(), tmp.end(), cmpy);
        for (int i = 0; i < tmp.size(); i++)</pre>
11
            for (int j = i + 1; j < tmp.size() && tmp[j].y - tmp[i].y < d; j++)
12
                d = min(d, DisPP(tmp[i], tmp[j]));
13
14
        return d;
    }
15
```

## 5.10 K-D Tree

```
#include <iostream>
#include <algorithm>
#include <stack>
#include <math.h>
using namespace std;
/*function of this program: build a 2d tree using the input trainingdata
the input is exm_set which contains a list of tuples (x,y)
the output is a 2d tree pointer*/

struct data
```

```
12
  {
13
        double x = 0;
14
        double y = 0;
15
   };
16
17
   struct Tnode
18
   {
19
        struct data dom elt;
20
        int split;
        struct Tnode * left;
21
        struct Tnode * right;
22
23
   };
24
25
   bool cmp1(data a, data b){
26
        return a.x < b.x;</pre>
27
   }
28
29
   bool cmp2(data a, data b){
30
        return a.y < b.y;</pre>
31
   }
32
   bool equal(data a, data b){
33
34
        if (a.x == b.x && a.y == b.y)
35
36
            return true;
37
        }
38
        else{
39
            return false;
40
        }
41
    }
42
    void ChooseSplit(data exm_set[], int size, int &split, data &SplitChoice{
43
        /*compute the variance on every dimension. Set split as the dismension that have the
44
        biggest
45
         variance. Then choose the instance which is the median on this split dimension.*/
        /*compute variance on the x,y dimension. DX=EX^2-(EX)^2*/
46
        double tmp1,tmp2;
47
        tmp1 = tmp2 = 0;
48
49
        for (int i = 0; i < size; ++i)</pre>
50
51
            tmp1 += 1.0 / (double)size * exm_set[i].x * exm_set[i].x;
52
            tmp2 += 1.0 / (double)size * exm_set[i].x;
53
        double v1 = tmp1 - tmp2 * tmp2; //compute variance on the xdimension
54
55
        tmp1 = tmp2 = 0;
56
57
        for (int i = 0; i < size; ++i)</pre>
58
            tmp1 += 1.0 / (double)size * exm_set[i].y * exm_set[i].y;
59
60
            tmp2 += 1.0 / (double)size * exm_set[i].y;
61
        double v2 = tmp1 - tmp2 * tmp2; //compute variance on the ydimension
62
63
        split = v1 > v2 ? 0:1; //set the split dimension
64
65
        if (split == 0)
66
67
        {
            sort(exm_set,exm_set + size, cmp1);
68
69
        }
70
        else{
71
            sort(exm_set,exm_set + size, cmp2);
```

```
72
73
74
         //set the split point value
75
         SplitChoice.x = exm_set[size / 2].x;
76
         SplitChoice.y = exm_set[size / 2].y;
77
78
    }
79
80
    Tnode* build_kdtree(data exm_set[], int size, Tnode* T){
81
         // call \ function \ Choose Split \ to \ choose \ the \ split \ dimension \ and \ split point
82
         if (size == 0){
             return NULL;
83
         }
84
         else{
85
86
             int split;
87
             data dom elt;
             ChooseSplit(exm_set, size, split, dom_elt);
88
             data exm set right [100];
89
             data exm set left [100];
90
91
             int sizeleft ,sizeright;
92
             sizeleft = sizeright = 0;
93
             if (split == 0)
94
95
                  for (int i = 0; i < size; ++i)</pre>
96
97
                  {
98
                      if (!equal(exm_set[i],dom_elt) && exm_set[i].x <=dom_elt.x)</pre>
99
100
                      {
101
                           exm_set_left[sizeleft].x = exm_set[i].x;
102
                           exm_set_left[sizeleft].y = exm_set[i].y;
103
                           sizeleft++;
                      }
104
                      else if (!equal(exm_set[i],dom_elt) && exm_set[i].x >dom_elt.x)
105
106
                           exm_set_right[sizeright].x = exm_set[i].x;
107
108
                           exm_set_right[sizeright].y = exm_set[i].y;
109
                           sizeright++;
                      }
110
111
                  }
112
             }
113
             else{
114
                  for (int i = 0; i < size; ++i)</pre>
115
116
                      if (!equal(exm_set[i],dom_elt) && exm_set[i].y <=dom_elt.y)</pre>
117
118
                      {
                           exm_set_left[sizeleft].x = exm_set[i].x;
119
120
                           exm_set_left[sizeleft].y = exm_set[i].y;
121
                           sizeleft++;
122
                      }
                      else if (!equal(exm_set[i],dom_elt) && exm_set[i].y >dom_elt.y)
123
124
                      {
                           exm_set_right[sizeright].x = exm_set[i].x;
125
126
                           exm_set_right[sizeright].y = exm_set[i].y;
127
                           sizeright++;
128
                      }
129
                  }
130
             T = new Tnode;
131
132
             T->dom_elt.x = dom_elt.x;
```

```
133
            T->dom_elt.y = dom_elt.y;
134
            T->split = split;
            T->left = build_kdtree(exm_set_left, sizeleft, T->left);
135
136
            T->right = build_kdtree(exm_set_right, sizeright, T->right);
137
            return T;
138
139
        }
140
141
142
143
    double Distance(data a, data b){
        double tmp = (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y);
144
145
        return sqrt(tmp);
146
    }
147
148
    void searchNearest(Tnode * Kd, data target, data &nearestpoint, double &distance){
149
150
151
        //1. 如果Kd是空的,则设dist为无穷大返回
152
        //2. 向下搜索直到叶子结点
153
154
        stack<Tnode*> search_path;
155
156
        Tnode* pSearch = Kd;
157
        data nearest;
158
        double dist;
159
160
        while(pSearch != NULL)
161
        {
162
             //pSearch加入到search_path中;
163
             search_path.push(pSearch);
164
            if (pSearch->split == 0)
165
166
                 if(target.x <= pSearch->dom_elt.x) /* 如果小于就进入左子树 */
167
168
                 {
169
                     pSearch = pSearch->left;
170
                }
                else
171
172
                {
173
                     pSearch = pSearch->right;
174
                 }
175
            }
             else{
176
                 if(target.y <= pSearch->dom_elt.y) /* 如果小于就进入左子树 */
177
178
                 {
179
                     pSearch = pSearch->left;
180
                 }
181
                else
182
                 {
183
                     pSearch = pSearch->right;
184
185
            }
186
187
        //取出 search_path最后一个赋给 nearest
        nearest.x = search_path.top()->dom_elt.x;
188
        nearest.y = search_path.top()->dom_elt.y;
189
190
        search_path.pop();
191
192
193
        dist = Distance(nearest, target);
```

```
194
        //3. 回溯搜索路径
195
196
        Tnode* pBack;
197
198
        while(search path.size() != 0)
199
        {
            //取出search path最后一个结点赋给pBack
200
201
            pBack = search path.top();
202
           search_path.pop();
203
           if(pBack->left == NULL && pBack->right == NULL) /* 如果pBack为子结点 */
204
205
            {
206
207
208
               if( Distance(nearest, target) > Distance(pBack->dom_elt,target) )
209
210
                   nearest = pBack->dom elt;
                   dist = Distance(pBack->dom elt, target);
211
212
               }
213
214
           }
215
216
           else
217
218
            {
219
220
               int s = pBack->split;
221
               if (s == 0)
222
               {
                   if( fabs(pBack->dom_elt.x - target.x) < dist) /* 如果 target为中心的圆
223
        (球或超球),半径为dist的圆与分割超平面交, 那么就要跳到另一边的子空间去搜索 */
224
                       if( Distance(nearest, target) > Distanc(pBack->dom_elt, target) )
225
226
227
                           nearest = pBack->dom_elt;
                           dist = Distance(pBack->dom_elt, target);
228
229
                       }
                       if(target.x <= pBack->dom_elt.x) /* 如果 target位 pBack的 左子空间, 那
230
        么就要跳到右子空间去搜索 */
231
                           pSearch = pBack->right;
232
                       else
233
                           pSearch = pBack->left; /* 如果 target位于 pBack的子空间, 那么就要
       跳到左子空间去搜索 */
                       if(pSearch != NULL)
234
235
                           //pSearch加入到search_path中
236
                           search_path.push(pSearch);
237
                   }
238
               }
239
               else {
240
                   if( fabs(pBack->dom_elt.y - target.y) < dist) /* 如果 target为中心的圆
        (球或超球),半径为dist的圆与分割超平面交, 那么就要跳到另一边的子空间去搜索 */
                   {
241
242
                       if( Distance(nearest, target) > Distanc(pBack->dom_elt, target) )
243
                       {
244
                           nearest = pBack->dom_elt;
                           dist = Distance(pBack->dom_elt, target);
245
246
                       if(target.y <= pBack->dom_elt.y) /* 如果 target位 pBack的 左子空间, 那
247
        么就要跳到右子空间去搜索 */
                           pSearch = pBack->right;
248
249
```

```
pSearch = pBack->left; /* 如果target位于pBack的子空间, 那么就要
250
         跳到左子空间去搜索 */
                          if(pSearch != NULL)
251
                             // pSearch加入到search_path中
252
253
                              search_path.push(pSearch);
254
                      }
255
                 }
256
257
             }
258
259
         nearestpoint.x = nearest.x;
260
261
         nearestpoint.y = nearest.y;
262
         distance = dist;
263
264
    }
265
266
    int main(){
267
         data exm_set[100]; //assume the max training set size is 100
268
         double x,y;
269
         int id = 0;
         cout<<"Please input the training data in the form x y. One instanceper line. Enter
270
         -1 -1 to stop."<<endl;</pre>
         while (cin>>x>>y){
271
272
             if(x == -1)
273
             {
274
                 break;
275
             }
276
             else{
277
                 exm_set[id].x = x;
278
                 exm_set[id].y = y;
279
                 id++;
             }
280
         }
281
         struct Tnode * root = NULL;
282
         root = build_kdtree(exm_set, id, root);
283
284
         data nearestpoint;
285
286
         double distance;
287
         data target;
288
         cout <<"Enter search point"<<endl;</pre>
289
         while (cin>>target.x>>target.y)
290
291
             searchNearest(root, target, nearestpoint, distance);
             cout<<"The nearest distance is "<<distance<<",and the nearestpoint is "<<</pre>
292
         nearestpoint.x<<","<<nearestpoint.y<<endl;</pre>
293
             cout <<"Enter search point"<<endl;</pre>
294
295
         }
296
```

# 6 Conclusion

### 6.1 Math

#### 6.1.1 Euler's Theorem

$$a^{b} \equiv \begin{cases} a^{b\%\varphi(p)} & \gcd(a,p) = 1 \\ a^{b} & \gcd(a,p) \neq 1, b < \varphi(p) \\ a^{b\%\varphi(p) + \varphi(p)} & \gcd(a,p) \neq 1, b \geq \varphi(p) \end{cases} \pmod{p}$$

#### 6.1.2 Möbius Inversion

Dirichlet Convolution is  $(f \times g)(N) = \sum_{d|N} f(d) * g(\frac{N}{d})$ 

Theorem:

$$\begin{cases} f = g \times 1 \\ g = f \times \mu \end{cases}$$

$$\begin{cases} id(n) = \sum_{d|n} \varphi(d) \\ e(n) = \sum_{d|n} \mu(d) \end{cases}$$
 (1)

$$\begin{cases}
\sum_{i}^{n} \sum_{j}^{m} \gcd(i,j) = \sum_{d}^{\max(n,m)} \varphi(d) * \lfloor \frac{n}{d} \rfloor \lfloor \frac{m}{d} \rfloor \\
\sum_{i}^{n} \sum_{j}^{m} e(\gcd(i,j)) = \sum_{d}^{\min(n,m)} \mu(d) * \lfloor \frac{n}{d} \rfloor \lfloor \frac{m}{d} \rfloor \\
\sum_{i=1}^{n} |\mu(i)| = \sum_{i=1}^{\lfloor \sqrt{n} \rfloor} \mu(i) * \lfloor \frac{n}{i*i} \rfloor
\end{cases} \tag{2}$$

$$\begin{cases} sum(x,y) = \sum_{i}^{x} \sum_{j}^{y} i * j = \frac{x * (x+1)}{2} * \frac{y * (y+1)}{2} \\ F(x,y) = \sum_{i=1}^{\min(x,y)} i^{2} * \mu(i) * sum(\lfloor \frac{x}{i} \rfloor, \lfloor \frac{y}{i} \rfloor) \\ \sum_{i}^{n} \sum_{j}^{m} lcm(i,j) = \sum_{i=1}^{\min(n,m)} d * F(\lfloor \frac{n}{i} \rfloor, \lfloor \frac{y}{i} \rfloor) \end{cases}$$

$$(3)$$

## 6.1.3 Sieve Tips

$$\varphi(nm) = \varphi(n) \cdot \varphi(m) \cdot \frac{\gcd(n,m)}{\varphi(\gcd(n,m))} \tag{4}$$

$$\varphi(n) = \sum_{i=1}^{n} [(n,i) = 1] \cdot i = \frac{n * \varphi(n) + [n=1]}{2}$$
(5)

$$\begin{cases}
id = \varphi \times 1 \\
\frac{n \cdot (n+1)}{2} = \sum_{i=1}^{n} i = \sum_{i=1}^{n} \sum_{d|i} \varphi(d) = \sum_{\frac{i}{d}=1}^{n} \sum_{d=1}^{\lfloor \frac{n}{i} \rfloor} \varphi(d) = \sum_{i=1}^{n} \phi(\lfloor \frac{n}{i} \rfloor)
\end{cases}$$
(6)

$$\begin{cases} e = \mu \times 1 \\ 1 = \sum_{i=1}^{n} [i = 1] = \sum_{i=1}^{n} \sum_{d|i} \mu(d) = \sum_{i=1}^{n} \sum_{d=1}^{\lfloor \frac{n}{i} \rfloor} \mu(d) = \sum_{i=1}^{n} M(\lfloor \frac{n}{i} \rfloor) \end{cases}$$
 (7)

$$\begin{cases}
id^{2} = (id \cdot \varphi) \times id \\
\phi'(n) = \sum_{i=1}^{n} i \cdot \varphi(i) \\
\frac{n \cdot (n+1) \cdot (2n+1)}{6} = \sum_{i=1}^{n} i^{2} = \sum_{i=1}^{n} \sum_{d|i} d \cdot \varphi(d) \cdot \frac{i}{d} = \sum_{\frac{i}{d}=1}^{n} \frac{i}{d} \sum_{d=1}^{\frac{n}{d}} d \cdot \varphi(d) = \sum_{i=1}^{n} i \cdot \phi'(\lfloor \frac{n}{i} \rfloor)
\end{cases}$$
(8)

#### 6.1.4 Newton's method

Ιff

$$f^{'}(x) \neq 0 \&\& f^{''}(x) continuous$$

Then

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

#### 6.1.5 Cantor Expansion

康托展开是一个全排列到一个自然数的双射,常用于构建哈希表时的空间压缩。康托展开的实质是计算当前排列在所有由小到大全排列中的顺序,因此是可逆的。以下称第 x 个全排列是都是指由小到大的顺序。

$$X = a_n(n-1)! + a_{n-1}(n-2)! + \cdots + a_1 \cdot 0!$$

其中  $a_i$  为排列的第 i 个元素值

例如,357412968展开为 98884。因为 X=2\*8!+3\*7!+4\*6!+2\*5!+0\*4!+0\*3!+2\*2!+0\*1!+0\*0!=98884. 解释:

排列的第一位是 3, 比 3 小的数有两个, 以这样的数开始的排列有 8! 个, 因此第一项为 2\*8!

排列的第二位是 5,比 5 小的数有 1、2、3、4,由于 3 已经出现,因此共有 3 个比 5 小的数,这样的排列有 7! 个,因此第二项为 3\*7!

以此类推,直至 0\*0!

用途:

显然,n 位(0 n-1)全排列后,其康托展开唯一且最大约为 n!,因此可以由更小的空间来储存这些排列。由公式可将 X 逆推出唯一的一个排列。

逆运算:

如 n=5,x=96 时:

首先用 96-1 得到 95, 说明 x 之前有 95 个排列.(将此数本身减去 1)

用 95 去除 4! 得到 3 余 23, 说明有 3 个数比第 1 位小, 所以第一位是 4.

用 23 去除 3! 得到 3 余 5, 说明有 3 个数比第 2 位小, 所以是 4, 但是 4 已出现过, 因此是 5.

用 5 去除 2! 得到 2 余 1, 类似地, 这一位是 3.

用 1 去除 1! 得到 1 余 0, 这一位是 2.

最后一位只能是 1.

所以这个数是 45321.

按以上方法可以得出通用的算法。

**6.1.6**  $\sum_{i=1}^{n} i^k$ 

$$\sum_{i}^{n} i = \frac{n(n+1)}{2}$$

$$\sum_{i}^{n} i^{2} = \frac{n(n+1)(2n+1)}{6}$$

$$\sum_{i}^{n} i^{3} = \left[\frac{n(n+1)}{2}\right]^{2}$$

# 差分和组合数: $O(k^2)$

考虑数列  $\{1_k, 2_k, \cdots, i_k\}$ 

相邻两项做差之后,得到的数列的每项应该是一个 k-1 次关于 i 的多项式。 再次相邻两项做差之后,得到的数列的每项应该是一个 k-2 次关于 i 的多项式。 如此进行 k 次,得到的数列的每项应该是一个 0 次关于 i 的多项式,即常数数列。 再次相邻两项做差之后,一定会得到一个全是 0 的数列。

假设经过 i 次差分数列之后的数列第一项为  $r_i$  那么答案就是

$$\sum_{i=0}^{k} r_i \binom{n}{i+1}$$

举例来说,考虑数列

 $1,8,27,64,125,216,\dots$ 

求差分可得

7,19,37,61,91,...

12,18,24,30,...

6,6,6,...

0,0,...

每个数列的第一项为1,7,12,6。所以最终的答案即为

$$\binom{n}{1} + 7\binom{n}{2} + 12\binom{n}{3} + 6\binom{n}{4} = \frac{n^2(n+1)^2}{4}$$

## 拉格朗日插值:O(k)

### 6.1.7 Generating Function

## 生成函数

普通型

$$A(x) = \sum_{i=0}^{\infty} a_i x^i$$

已知  $\{a_i\}$ ,  $\{b_i\}$  的生成函数分别是 A(x), B(x)。

那么  $\{a_i \pm b_i\}$  的生成函数是  $A(x) \pm B(x)$ 。

值得注意的是数列  $\{a_i\}$ ,  $\{b_i\}$  的卷积的生成函数,恰好是 C(x) = A(x)B(x)。

$$c_i = \sum_j a_j b_{i-j}$$

## 关键公式

$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + \dots = \sum_{i} x^i$$
$$\frac{1}{1-x^2} = 1 + x^2 + x^4 + x^6 + \dots = \sum_{i} x^{2i}$$

推广的二项式定理

$$(1+x)^n = \binom{n}{0}x^0 + \binom{n}{1}x^1 + \binom{n}{2}x^2 + \dots = \sum_i \binom{n}{i}x^i$$

其中有  $\beta$  函数和  $\Gamma$  函数

$$\begin{split} \Gamma(s) &= \int_0^{+\infty} x^{s-1} e^{-x} dx \\ B(p,q) &= \binom{p+q}{p} = \int_0^1 x^{p-1} (1-x)^{q-1} dx = \frac{\Gamma(p)\Gamma(q)}{\Gamma(p+q)} \end{split}$$

分别有递推公式

$$\begin{split} \Gamma(s+1) &= s\Gamma(s) \\ B(p,q) &= B(q,p) \\ &= \frac{q-1}{p+q-1}B(p,q-1) \\ &= \frac{p-1}{p+q-1}B(p-1,q) \\ &= \frac{(p-1)(q-1)}{(p+q-1)(p+q-2)}B(p-1,q-1) \end{split}$$

## 例子

## Fibonacci 生成函数

构造一个与  $f_n$  有关得幂级数

$$F(x) = f_0 + f_1 x + f_2 x^2 + \dots + f_n x^n + \dots$$

亦即

$$F(x) = xF(x) + x^{2}F(x) + f_{0} + (f_{1} - f_{0})x$$

带入 
$$f_0 = 0, f_1 = 1$$

$$F(x) = \frac{x}{1 - x - x^2}$$

如果想继续计算通项, 我们需要解方程得到

$$\frac{x}{1-x-x^2} = \frac{a}{1-\alpha x} + \frac{b}{1-\beta x}$$

解出

$$\begin{cases} \alpha = \frac{1 + \sqrt{5}}{2} \\ \beta = \frac{1 - \sqrt{5}}{2} \end{cases}$$

带入方程有

$$a = \frac{1}{\sqrt{5}}, b = -\frac{1}{\sqrt{5}}$$

则

$$F(x) = \frac{a}{1 - \alpha x} + \frac{b}{1 - \beta x}$$
  
=  $a(1 + \alpha x + \alpha^2 x^2 + \dots) + b(1 + \beta x + \beta^2 x^2 + \dots)$ 

蕴含通项为

$$f_n = a\alpha^n + b\beta^n = \frac{(\frac{1+\sqrt{5}}{2})^n - (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$$

## Catalan 数

$$c_n = \sum_{i=1}^{n-1} c_i c_{n-i}$$

生成函数

$$C(x) = c_1 x + c_2 x^2 + \cdots$$

亦即

$$C(x) = xC(x)^2 + 1$$

解得

$$C(x) = \frac{1 - \sqrt{1 - 4x}}{2x}$$

利用推广的二项式定理展开  $\sqrt{1-4x}$  可得通项

$$c_n = \frac{1}{n} \binom{2n-2}{n-1}$$

# 指数型生成函数

$$A(x) = \sum_{i=0}^{\infty} \frac{a_i x^i}{i!}$$

已知  $\{a_i\}$ ,  $\{b_i\}$  的生成函数分别是 A(x), B(x)。

那么  $\{a_i \pm b_i\}$  的生成函数是  $A(x) \pm B(x)$ 。

值得注意的是数列  $\{a_i\}$ ,  $\{b_i\}$  的卷积的生成函数,恰好是 C(x) = A(x)B(x)。

$$c_i = \sum_{j} \binom{i}{j} a_j b_{i-j}$$

关键公式

$$e^{x} = 1 + x + \frac{1}{2}x^{2} + \frac{1}{6}x^{3} + \dots = \sum_{i} \frac{x^{i}}{i!}$$

$$e^{-x} = 1 - x + \frac{1}{2}x^{2} - \frac{1}{6}x^{3} + \dots = \sum_{i} \frac{(-x)^{i}}{i!}$$

$$\frac{e^{x} + x^{-x}}{2} = 1 + \frac{1}{2}x^{2} + \frac{1}{24}x^{4} + \dots = \sum_{i} \frac{x^{2i}}{(2i)!}$$

#### 6.1.8 Polya

设 G 是 p 个对象的一个置换群,用 k 种颜色去染这 p 个对象,若一种染色方案在群 G 的作用下变为另一种方案,则这两个方案当作是同一种方案,这样的不同染色方案数为:

$$\frac{1}{|G|} \times \sum (k^{C(f)}), f \in G$$

C(f) 为循环节,|G| 表示群的置换方法数 对于有 n 个位置的手镯,有 n 种旋转置换和 n 种翻转置换 对于旋转置换:

$$C(f_i) = gcd(n, i)$$

i 表示一次转过 i 颗宝石, i = 0 时 c = n;

对于翻转置换:

如果 n 为偶数: 则有  $\frac{n}{2}$  个置换  $C(f) = \frac{n}{2}$ ,有  $\frac{n}{2}$  个置换  $C(f) = \frac{n}{2} + 1$ 

如果 n 为奇数:  $C(f) = \frac{n}{2} + 1$ 

#### 6.1.9 FWT

$$\begin{cases} C_k = \sum_{i \oplus j = k} A_i * B_j \\ DWT(A)_i = \sum_{j}^n A_j * f_{i,j} \\ DWT(C)_i = DWT(A)_i * DWT(B)_i \\ f_{i,j} \cdot f_{i,k} = f_{i,j \oplus k} \\ f_{i,j} = [i \ and \ j == i] \\ f_{i,j} = [i \ and \ j == j] \\ f_{i,j} = [i \ and \ j == j] \end{cases}$$
 (and)
$$f_{i,j} = [i \ and \ j == j]$$
 (or)
$$f_{i,j} = (-1)^{|i| \ and \ j|}$$
 (xor)

## 6.2 Geometry

#### 6.2.1 The Number of Ingeter Point on a Circle

Set r = const is the radius of the circle.

$$r^2 = p_1^{a_1} + p_2^{a_2} + \dots + p_m^{a_m} = \sum_{i=1}^m p_i^{a_i}$$

Define

$$\chi(n) = \begin{cases} 1 & n\%4 = 1 \\ -1 & n\%4 = 3 \\ 0 & n\%2 = 0 \end{cases}$$

By the way,  $\chi(n)$  is a multiplicative function.

Define

$$\Gamma(p_i, a_i) = \sum_{j=0}^{a_i} \chi(p_i^j) = \begin{cases} 1 & p_i = 2 & || & (p_i\%4 = 3 & \&\& & a_i\%2 = 0) \\ 0 & p_i\%4 = 3 & \&\& & a_i\%2 = 1 \\ a_i + 1 & p_i\%4 = 1 \end{cases}$$

Define cnt is the number of integer point on circle

$$cnt(r) = 4 \prod_{i=1}^{m} \sum_{j=0}^{a_i} \chi(p_i^j) = 4 \prod_{i=1}^{m} \Gamma(p_i, a_i) = 4 \sum_{k|r^2} \chi(k)$$

Define CNT is the number of integer point in circle

$$CNT(r) = 1 + \sum_{i=1}^{r^2} cnt(i) = 1 + \sum_{i=1}^{r^2} \lfloor \frac{r^2}{i} \rfloor \chi(i)$$

## 6.3 Josephus

#### **6.3.1** J(n,m): The Last Surviving Person

based-0, m-1 was the first be killed.

$$J_{0}(n,m) = \begin{cases} 0 & n = 1; \\ (J_{0}(n-1,m) + m)\%n & 1 < n < m; \\ \lfloor \frac{k((f(n',k) - nmodk)modn')}{k-1} \rfloor & where \ n' = n - \lfloor \frac{n}{k} \rfloor & m <= n \end{cases}$$

$$J_k(n,m) = (J_0(n,m) + k)\%n$$

```
int J0(int n, int m) {
    if (m == 1) return n - 1;
    int ans = 0;
    for (int i = 2; i <= n; ) {
        if (ans + m >= i) {
            ans = (ans + m) % i;
            i++;
            continue;
    }
}
```

```
9
10
            int step = (i - 1 - ans - 1) / (m - 1);
11
            if (i + step > n) {
12
                ans += (n - (i - 1)) * m;
13
                break;
14
            }
15
            i += step; ans += step * m;
16
        }
17
        return ans;
18
```

## **6.3.2** aJ(n,1,K): Survival Time of K-th Person

based-1, 1 was the first be killed.

```
int aJ1(int size, int be, int goal){
    if (be > size) be = (be - 1) % size + 1;
    if (goal % m == be % m) return (goal - be) / m + 1;
    return (size - be) / m + 1 + get(size - (size - be) / m - 1, (be + ((size - be) / m + 1) * m) - size, goal > be ? goal - (goal - be) / m - 1 : goal);
}
{
    \\ size : the size of current group;
\\ be : the start person of current step;
\\ goal : index of the asked person.
```

## 7 Others

## 7.1 Offline Algorithm

### 7.1.1 CDQ Divide and Conquer

```
struct Node {
1
2
        int x, y, z, ans;
3
        Node() {}
4
        Node(int _x, int _y, int _z):x(_x), y(_y), z(_z) {}
5
        bool operator < (const Node &b) const {</pre>
 6
            if(y == b.y) {
 7
                 if(z == b.z) return x < b.x;
8
                 return z < b.z;</pre>
9
10
            return y < b.y;</pre>
11
   }A[MAXN], B[MAXN], C[MAXN];
12
   int bit[MAXN];
13
   void add(int k, int v) {
14
15
        for(; k <= m; k += k & -k) bit[k] = max(bit[k], v);</pre>
16
17
   void clear(int k) {
        for(; k \le m; k += k \& -k) bit[k] = 0;
18
19
20
   int sum(int k) {
21
        int res = 0;
        for(; k; k -= k & -k) res = max(res, bit[k]);
22
23
        return res;
24
   }
25
   void solve(int 1, int r) {
26
        if(1 == r) {
27
            B[1] = A[1];
28
             return;
29
30
        int mid = (1 + r) >> 1;
31
        solve(l, mid);
32
        for(int i = mid + 1; i <= r; i++) B[i] = A[i];</pre>
33
        //sort(B + l, B + mid + 1);
34
        sort(B + mid + 1, B + r + 1);
        int L = 1;
35
        for(int R = mid + 1; R <= r; R++) {</pre>
36
37
            while(L <= mid && B[L].y < B[R].y) add(B[L].z, B[L].ans), L++;
38
            A[B[R].x].ans = max(A[B[R].x].ans, sum(B[R].z - 1) + 1);
39
            B[R].ans = A[B[R].x].ans;
40
        for(int i = 1; i <= L; i++) clear(B[i].z);</pre>
41
42
        solve(mid + 1, r);
43
        L = 1;
        int p = 1, q = mid + 1;
44
45
        while(p <= mid || q <= r) {
            if(q > r | | (p \le mid \&\& B[p].y \le B[q].y)) C[L++] = B[p++];
46
47
            else C[L++] = B[q++];
48
49
        for(int i = 1; i <= r; i++) B[i] = C[i];</pre>
50
```

## 7.1.2 Mo's Algorithm

```
struct Node{
1
2
        int 1, r, t, id;
3
        bool operator < (const Node& a) const {</pre>
4
             if(1 /sz == a.1 / sz) {
5
                 if(r == a.r) return t < a.t;</pre>
6
                 return r < a.r;</pre>
7
             return 1 / sz < a.1 / sz;</pre>
8
        }
9
   }q[MAXN];
10
   void solve() {
11
12
        while (t < q[i].t) addTime(t++, 1);</pre>
13
        while (t > q[i].t) addTime(--t, -1);
14
        while(L < q[i].1) add(L++, -1);
15
        while(L > q[i].1) add(--L, 1);
16
        while(R < q[i].r) add(++R, 1);</pre>
17
        while(R > q[i].r) add(R--, -1);
18
```

#### 7.1.3 Mo's Algorithm On Tree

```
struct Edge {
1
2
       int to, nxt;
3
   }e[MAXN << 1];
   int head[MAXN], ecnt;
   int stack[MAXN], top, belong[MAXN], cnt, sz;
6
    struct Node {
        int l, r, id, ti;
7
        bool operator < (const Node &x) const {</pre>
8
            return belong[1] < belong[x.1] || (belong[1] == belong[x.1] && belong[r] <</pre>
9
        belong[x.r]) || (belong[1] == belong[x.1] && belong[r] == belong[x.r] && ti < x.ti);
10
   }q[MAXN];
11
   struct Node2 {
12
13
        int 1, r, ti;
14 }qq[MAXN];
15 int n, m, Q, Q0, Q1;
16 int V[MAXN], W[MAXN], C[MAXN];
int fa[MAXN][S + 3], dep[MAXN];
18 long long ans[MAXN], tans;
19 int vis[MAXN], cur[MAXN];
20 long long sum[MAXN];
21 int 1, r, tm;
   inline int read() {
22
        int x = 0; char ch = getchar(); bool fg = 0;
while(ch < '0' || ch > '9') { if(ch == '-') fg = 1; ch = getchar(); }
23
24
        while(ch >= 0 && ch <= 9) { x = x * 10 + ch - 0; ch = getchar(); }
25
26
        return fg ? -x : x;
27
   inline void add_edge(int u, int v) {
28
29
        e[++ecnt] = (Edge) {v, head[u]}; head[u] = ecnt;
        e[++ecnt] = (Edge) \{u, head[v]\}; head[v] = ecnt;
30
31
   void dfs(int u, int f) {
32
33
        fa[u][0] = f;
34
        dep[u] = dep[f] + 1;
        int bot = top;
```

```
36
        for(int i = head[u]; i; i = e[i].nxt) {
37
            int v = e[i].to;
38
            if(v == f) continue;
39
            dfs(v, u);
40
            if(top - bot >= sz) {
41
                 cnt++:
                 while(top != bot) belong[stack[top--]] = cnt;
42
43
            }
44
45
        stack[++top] = u;
46
    void G(int &u, int step) {
47
        for(int i = 0; i < S; i++) if((1 << i) & step) u = fa[u][i];
48
49
    }
50
   int lca(int u, int v) {
51
        if(dep[u] > dep[v]) swap(u, v);
        G(v, dep[v] - dep[u]);
52
53
        if(u == v) return u;
54
        for(int i = S; i >= 0; i--) if(fa[u][i] != fa[v][i]) {
55
            u = fa[u][i]; v = fa[v][i];
56
        return fa[u][0];
57
58
   }
59
    inline void modify(int u) {
60
        tans -= V[C[u]] * sum[cur[C[u]]];
61
        cur[C[u]] += vis[u];
62
        vis[u] = -vis[u];
63
        tans += V[C[u]] * sum[cur[C[u]]];
64
65
    inline void update(int u, int v) {
66
        if(u == v) return;
67
        if(dep[u] > dep[v]) swap(u, v);
        while(dep[v] > dep[u]) {
68
69
            modify(v);
70
            v = fa[v][0];
71
        while(u != v) {
72
            modify(u); modify(v);
73
74
            u = fa[u][0]; v = fa[v][0];
75
        }
76
   }
77
    inline void upd(int t) {
78
        if(vis[qq[t].1] == -1) {
79
            modify(qq[t].1);
80
            swap(C[qq[t].1], qq[t].r);
81
            modify(qq[t].1);
82
        else swap(C[qq[t].1], qq[t].r);
83
84
    inline void moveto(int u, int v) {
85
86
        update(1, u); update(r, v);
87
        1 = u; r = v;
88
    int main() {
89
90
        n = read(); m = read(); Q = read();
91
        sz = (int)pow(n, 2.0 / 3.0);
92
        for(int i = 1; i <= m; i++) V[i] = read();</pre>
93
        for(int i = 1; i <= n; i++) W[i] = read();</pre>
94
        for(int i = 1, u, v; i < n; i++) {</pre>
            u = read(); v = read();
95
96
            add_edge(u, v);
```

```
97
98
         for(int i = 1; i <= n; i++) {</pre>
99
             C[i] = read();
100
             vis[i] = 1;
101
             sum[i] = sum[i - 1] + W[i];
102
         for(int i = 1, tp; i <= Q; i++) {</pre>
103
             tp = read();
104
105
              if(tp) {
106
                  ++Q1;
107
                  q[Q1].l = read(); q[Q1].r = read();
                  q[Q1].id = Q1;
108
                  q[Q1].ti = i;
109
110
             }
             else {
111
                  ++00;
112
                  qq[Q0].1 = read(); qq[Q0].r = read();
113
114
                  qq[Q0].ti = i;
115
             }
116
         dfs(1, 0);
117
         while(top) belong[stack[top--]] = cnt;
118
119
         sort(q + 1, q + Q1 + 1);
         for(int k = 1; k <= S; k++) {</pre>
120
121
             for(int i = 1; i <= n; i++) {</pre>
122
                  fa[i][k] = fa[fa[i][k - 1]][k - 1];
123
              }
124
125
         for(int i = 1; i <= Q1; i++) {
126
             if(belong[q[i].1] > belong[q[i].r]) swap(q[i].1, q[i].r);
127
             moveto(q[i].l, q[i].r);
             int lc = lca(1, r);
128
129
             modify(lc);
             while(qq[tm + 1].ti < q[i].ti && tm < Q0) upd(++tm);</pre>
130
             while(qq[tm].ti > q[i].ti) upd(tm--);
131
132
             ans[q[i].id] = tans;
133
             modify(lc);
134
         }
135
         for(int i = 1; i <= Q1; i++) printf("%lld\n", ans[i]);</pre>
136
         return 0;
137
```

## 7.2 Randomized Algorithm

## 7.2.1 Simulated Annealing

```
void solve() {
1
2
       while(T > eps) {
            double alpha = ((rand() % 30001) / 15000.0) * pi;
3
           double theta = ((rand() % 10001) / 10000.0) * pi;
4
           tmp.x = cur.x + T * sin(theta) * cos(alpha);
5
           tmp.y = cur.y + T * sin(theta) * sin(alpha);
6
7
           tmp.z = cur.z + T * cos(theta);
           tmp.dis = cal(tmp);
8
9
            if(tmp.dis < cur.dis || (tmp.dis * 0.999 < cur.dis && (rand() & 7) == 7)) cur =
       tmp;
10
            //if(exp((cur.d - tmp.d) / T) > ((double)rand() / RAND_MAX)) cur = tmp;
11
           T *= 0.999;
12
```

```
13 }
14 }
```

## 7.3 Other Method

### 7.3.1 Enumerate Subset

## **7.3.2** Enumerate $\lfloor \frac{n}{d} \rfloor \lfloor \frac{m}{d} \rfloor$

```
1
   int cal(int n, int m) {
2
        if(n > m) swap(n, m);
3
        int res = 0, last;
4
        for(int i = 1; i <= n; i = last + 1) {</pre>
             last = min(n / (n / i), m / (m / i));
res += (n / i) * (m / i) * (sum(last) - sum(i - 1));
5
6
7
        }
8
        return res;
9
```

### 7.3.3 Find Primitive Root Modulo N

```
for i in range(1,mod):
    if 3 ** i % mod == 1:
        if i == mod - 1:
            print("yes")
            break
        print("no")
```

# 8 Samples

### 8.1 vimrc

```
set cindent
2 set number
3 set mouse=a
4 set tabstop=4
5 set shiftwidth=4
6 syntax on
7 inoremap { {}<left>
8 map \langle F9 \rangle :w\langle CR \rangle :! g++ % -o %\langle -Wall --std=c++14 -g && ./% \langle -CR \rangle
   "ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;
1
   set nocompatible
2 source $VIMRUNTIME/vimrc_example.vim
   source $VIMRUNTIME/mswin.vim
3
   se cin nu mouse=a ts=4 sw=4 ww=b,s,<,>,[,]
4
5
    syntax on
6
    inoremap { {}<left>
    map <F9> :w<CR> :! g++ % -o %< -Wall --std=c++14 -g && ./%< <CR>
    func! AddTitle()
        call append(0,"// Cease\ to\ struggle\ and\ you\ cease\ to\ live") call append(1,"// Created\ by\ hjj")
9
10
        call append(2,"#include <bits/stdc++.h>")
11
        call append(3,"using namespace std;")
12
        call append(4,"")
13
        call append(5,"int main() {")
14
        call append(6,"")
call append(7,"
15
16
                              return 0;")
        call append(8,"}")
17
   endfunc
18
   map <F8> :call AddTitle()<CR>
```

## 8.2 Check

Linux

```
while true; do
   ./data > in
   ./tmp < in > out
   ./std < in > ans
   diff out ans
   if [ $? -ne 0 ]; then exit; fi
   echo Passed
done
```

windows

```
1  @echo off
2  :loop
3     rand.exe > data.in
4     std.exe < data.in > std.out
5     my.exe < data.in > my.out
6     fc my.out std.out
7     if not errorlevel 1 goto loop
8     pause
9     goto loop
```

#### 8.3 Random

```
1 mt19937_64 mt(chrono::steady_clock::now().time_since_epoch().count());
2 shuffle(per.begin(), per.end(), mt);
3 //random_shuffle(per.begin(), per.end());
```

#### 8.4 FastIO

```
//普通情况
1
2
   namespace IO {
        const int MB = 1048576;
3
        const int RMAX = 16 * MB;
4
        const int WMAX = 16 * MB;
5
        #define getchar() *(rp++)
6
7
        #define putchar(x) (*(wp++) = (x))
8
        char rb[RMAX], *rp = rb, wb[WMAX], *wp = wb;
9
        inline void init() {
10
            fread(rb, sizeof(char), RMAX, stdin);
11
12
        template <class _T> inline void read(_T &_a) {
13
            _a = 0; bool _f = 0; int _c = getchar();
            while (_c < '0' || _c > '9') _f |= _c == '-', _c = getchar();
14
            while (_c >= '0' \&\& _c <= '9') _a = _a * 10 + (_c ^ '0'), _c = getchar();
15
16
            _a = _f ? -_a : _a;
17
18
        template <class _T> inline void write(_T _a) {
19
            static char buf[20], *top = buf;
20
            if (_a) {
21
                while (_a) {
                    _T tm = _a / 10;
*(++top) = char(_a - tm * 10) | '0';
22
23
24
25
26
                while (top != buf) putchar(*(top--));
27
28
            else putchar('0');
29
        }
        void output() {
30
            fwrite(wb, sizeof(char), wp - wb, stdout);
31
32
33 }
   //EOF结尾+分块读入
34
35 #define likely(x) __builtin_expect(!!(x), 1)
   #define unlikely(x) __builtin_expect(!!(x), 0)
36
37
   namespace IO {
        const int MB = 1048576;
38
        const int RMAX = 4 * MB;
39
40
        const int WMAX = 4 * MB;
41
        unsigned long long filesize;
42
        #define putchar(x) (*(wp++) = (x))
43
        char rb[RMAX], wb[WMAX], *wp = wb;
44
        int rp = 0;
45
        inline void init() {
            filesize = fread(rb, sizeof(char), RMAX, stdin);
46
47
            rp = 0;
            wp = wb;
48
49
50
        void output() {
51
            fwrite(wb, sizeof(char), wp - wb, stdout);
```

```
52
53
        inline char getCHAR(){
54
            if(unlikely(rp == filesize)){
55
                fwrite(wb, sizeof(char), wp - wb, stdout);
                init();
56
                if(unlikely(filesize == 0)) {
57
                     //cerr \ll 1.0 * (clock() - st) / CLOCKS_PER_SEC \ll endl;
58
59
                    exit(0);
60
                }
61
            }
62
            return rb[rp++];
63
64
        template <class _T> inline void read(_T &_a) {
            _a = 0; static bool _f = 0; static int _c;
65
66
             while (_c < '0' | _c > '9') _f | _c = '-', _c = getCHAR();
67
            while (_c >= '0' \&\& _c <= '9') _a = _a * 10 + (_c ^ '0'), _c = getCHAR();
68
69
            _a = _f ? -_a : _a;
70
71
        template <class _T> inline void write(_T _a) {
72
            static char buf[20], *top = buf;
            if (_a) {
73
74
                while (_a) {
75
                     _{T} tm = _{a} / 10;
76
                    *(++top) = char(_a - tm * 10) | '0';
77
                    _a = tm;
78
79
                while (top != buf) putchar(*(top--));
80
            }
81
            else putchar('0');
82
            putchar(' \setminus n');
83
        }
84
   }
```

## 8.5 Java BigNum

```
1
   import java.math.*;
2
   import java.util.*;
3
   import java.lang.*;
4
5
   public class Main{
6
       public static void main(String []args){}
7
8
   //IO
9
   Scanner in = new Scanner(System.in);
10
   while(in.hasNext()){} //EOF
11
   public static void main(String argv[]) throws IOException{}
12
   StreamTokenizer cin = new StreamTokenizer(new BufferedReader(new InputStreamReader(
13
       System.in)));
14 PrintWriter cout = new PrintWriter(new OutputStreamWriter(System.out));
15 while(cin.nextToken() != StreamTokenizer.TT EOF) ;//EOF
16 cin.nextToken();int n = (int)cin.nval;String s = cin.sval;
17 cout.println( Type );cout.flush();
18 cin.ordinaryChar('/');
19
20 BufferedReader br = new BufferedReader(new InputStreamReader(System.in));
21 br.ready()//EOF
22 while ((valueString=bf.readLine())!=null);
```

```
23
   br.close();
24
   //true\ fast-IO
25
   static class InputReader {
26
       public BufferedReader reader;
27
       public StringTokenizer tokenizer;
28
29
       public InputReader(InputStream stream) {
30
            reader = new BufferedReader(new InputStreamReader(stream), 32768);
31
            tokenizer = null;
32
33
       public String next() {
34
           while (tokenizer == null || !tokenizer.hasMoreTokens()) {
35
36
37
                    tokenizer = new StringTokenizer(reader.readLine());
38
                } catch (IOException e) {
39
                    throw new RuntimeException(e);
40
41
            }
42
            return tokenizer.nextToken();
43
44
       public int nextInt() {
45
46
           return Integer.parseInt(next());
47
       }
48
49
   //类 Number
50
51
   //double Value ()
52
   //intValue()
53
   //long Value()
54
   //shortValue()
   //类 BigDecimal
55
   //ROUND_CEILING 接近正无穷大的舍入模式。
56
57
   //ROUND_FLOOR 接近负无穷大的舍入模式。
   //ROUND_DOWN 接近零的舍入模式
58
   //ROUND HALF UP 四舍五入 >=0.5向上舍入
59
   //ROUND HALF DOWN 四舍五入 >0.5向上舍入
60
   //BigDecimal(BigInteger\ val)
61
62
   //BigDecimal(BigInteger\ unscaledVal,\ int\ scale)
63
   //BigDecimal(char] in, int offset, int len, MathContext mc)
64
   //BigDecimal(double val, MathContext mc)不建议
65
   //BigDecimal(int\ val,\ MathContext\ mc)
66
   //BigDecimal(long\ val\ ,\ MathContext\ mc)
67
   //BigDecimal(String\ val,\ MathContext\ mc)
   //abs()
68
   //add(BigDecimal\ augend,\ MathContext\ mc)
69
70
   //compareTo(BigDecimal\ val)
   //divide\ (BigDecimal\ divisor\ , MathContext\ mc)
71
72
   //divideToIntegralValue(BigDecimal divisor, MathContext mc)
73
   //max(BigDecimal\ val)
74
   //min(BigDecimal\ val)
   //multiply (BigDecimal multiplicand, MathContext mc)
75
               其值为 (-this), 其标度为 this.scale()
76
   //negate()
77
    //pow(int n)
   //remainder(BigDecimal divisor) 返回其值为 (this % divisor) 的 BigDecimal
78
   //round(MathContext mc) 返回根据 MathContext 设置进行舍入后的 BiqDecimal。
79
   //caleByPowerOfTen(int n) 返回其数值等于 (this * 10^n) 的 BigDecimal。
80
81
   //subtract(BigDecimal subtrahend, MathContext mc)
82
   //setScale(int\ newScale, RoundingMode\ roundingMode)
   //toString()
```

```
//ulp()返回此 BigDecimal 的 ulp (最后一位的单位) 的大小
85
    //String s = b.stripTrailingZeros().toPlainString();让bigdecimal不用科学计数法显示
86
   //类 BigInteger
   //parseInt
87
    //BigInteger\ zero = BigInteger.valueOf(0);
88
89
    //BigInteger \ a = in.nextBigInteger();
    //abs()
90
    //and(BigInteger val) 返回其值为 (this & val)
91
92
    //or(BigInteger val) 返回其值为 (this | val)
    //andNot(BigInteger val) 返回其值为 (this & ~val)
93
94
    //compareTo(BigInteger\ val)
95
    //add(BigInteger\ val)
    //divide(BigInteger val)
96
97
    //BigInteger[] divideAndRemainder(BigInteger val) 返回包含 (this / val) 后跟 (this %
        val) 的两个 BigInteger 的数组。
98
    //equals(Object x)
99
    //gcd(BigInteger\ val)
    //isProbablePrime(int certainty) e.g. a.isProbablePrime(4)
100
   //max(BigInteger val) min(BigInteger val)
101
102
   //mod(BigInteger m)
   //modInverse(BigInteger m) 返回其值为 (this ~-1 mod m)
103
   //modPow(BigInteger exponent, BigInteger m) 返回其值为 (this exponent mod m)
104
105
   //multiply(BigInteger\ val)
   //not() 返回其值为 (~this)
106
    //shiftLeft(int n) 返回其值为 (this << n)
107
108
    //shiftRight(int n) 返回其值为 (this >> n)
109
    //toString()
110
    //valueOf(long\ val)
    //xor(BigInteger val) 返回其值为 (this ^ val)
111
112
113
    //Arrays.sort(array);
```

### $8.6 ext{ pb\_ds}$

```
//P.S.: 无脑正确使用pb_ds代替 std:: set/map/priority_queue不会变慢
2
  //可持久化平衡树,不过时间和空间都不太行
3 #include <ext/rope>
4 using namespace __gnu_cxx;
5 int a[1000];
6 rope<int> x;
  rope < int > x(a,a + n);
7
8
   rope<int> a(x);
9
  x->at(10);x[10];
                     // 在末尾添加x
10
   x->push back(x)
                     // 在pos插入x
   x->insert(pos,x)
11
                     // 从 pos 开始 删除 x 个
12
   x->erase(pos,x)
                     // 从pos开始换成x
13 x->replace(pos,x)
                     // 提取pos开始x个
14
   x->substr(pos,x)
15
16
   //不支持低级操作(如交换左右子树)
17
18 #include <ext/pb ds/assoc container.hpp>
19 #include <ext/pb_ds/tree_policy.hpp>
20 using namespace __gnu_pbds;
21 定义一颗红黑树
22 tree<int,null_type,less<int>,rb_tree_tag,tree_order_statistics_node_update>t, other;
23 int 关键字类型
24 null_type
   无映射(低版本g++为null_mapped_type)(无映射为类似set,有映射类似map)
```

```
26 less<int>从小到大排序
27
   rb_tree_tag 红黑树 (splay_tree_tag)
   tree_order_statistics_node_update结点更新(统计子树size,可自写),不写不支持order_of_key
       以及find_by_order
29
   插入:t.insert();
   删除:t.erase();
30
   比x小的个数:t.order_of_key(x);
31
   第x+1值:t.find by order(x);
33
   前驱:t.lower_bound();
   后继:t.upper_bound();
34
   合并:t.join(other); (other和*this值域不能相交)
35
   分裂:t.split(x, other); (清空other, 将t中比x小的元素移至other)
36
37
   //自定义节点更新
38
39
   template <class Node_CItr , class Node_Itr , class Cmp_Fn , class _Alloc >
40
   struct my node update {
41
       virtual Node_CItr node_begin () const = 0;
       virtual Node CItr node end() const = 0;
42
       typedef char metadata type; //节点上记录的额外信息的类型
43
44
       //以上为固定格式
45
       //operator()的功能是将节点it的信息更新为其左右孩子的信息之和, 传入的end_it表示空节点
46
       //对Node_Itr可以做的事情有:用get_l_child,get_r_child获取左右孩子,用两个星号获取节
47
       点信息,用get_metadata获取节点额外信息
       inline void operator()(Node_Itr it, Node_CItr end_it) {
48
49
           Node_Itr l = it.get_l_child(), r = it.get_r_child();
50
           int left = 0, right = 0;
51
           if(1 != end_it) left = 1.get_metadata();
52
           if(r != end it) right = r.get metadata();
53
           const_cast<metadata_type &>(it.get_metadata()) = left + right + (*it)->second;
54
           //it 是 node_Itr, 取 * 后 变 为 iterator, 再 取 -> second 变 成 mapped_value
55
       inline int prefix_sum(int x) {
56
           int ans = 0;
57
           Node_CItr it = node_begin ();
58
           while(it != node_end()) {
59
60
               Node_CItr l = it.get_l_child(), r = it.get_r_child();
61
               if(Cmp_Fn()(x, (*it)->first)) it = 1;
62
               else {
63
                  ans += (*it)->second;
64
                  if(1 != node_end()) ans += 1.get_metadata();
65
                  it = r;
66
               }
67
           }
68
           return ans;
69
70
       inline int interval_sum(int 1, int r) {
71
           return prefix_sum(r) - prefix_sum(l - 1);
72
       }
73
   };
74
75
   tree<int, char, less<int>, rb_tree_tag, my_node_update> T;//map
76
       T[2] = a'; T[3] = b'; T[4] = 1;
77
       cout << (char)T.interval_sum(3, 4) << endl;//c</pre>
78
79
       return 0;
80
   }
   //堆
81
82 #include <ext/pb_ds/priority_queue.hpp>
83 using namespace gnu pbds;
   __gnu_pbds::priority_queue<int, std::less<int>, __gnu_pbds::pairing_heap_tag> q;
```

```
85
86 template <typename Value_Type ,
87 typename Cmp_Fn = std::less<Value_Type>,
88 typename Tag = pairing_heap_tag ,
89 typename Allocator = std::allocator<char> >
90 class priority_queue
   Tag可以是binary_heap_tag(二叉堆)binomial_heap_tag(二项堆)rc_binomial_heap_tag
       pairing heap tag (配对堆) thin heap tag
   用 begin()和 end()获取迭代器从而遍历
92
   删除单个元素 void erase(point_iterator)
93
   更改一个元素的值 void modify(point_iterator, const_reference)
94
   合并 void join(priority_queue &other), 把other合并到*this, 并把other清空
95
96 push()会返回迭代器
97
   五种操作: push、pop、modify、erase、join
98 • pairing_heap_tag: push和joinO(1), 其余均摊O(logn)
99 • binary heap tag: 只支持push和pop, 均为均摊O(logn)
100 • binomial_heap_tag: push为均摊O(1),其余为O(logn)
101 • rc binomial heap tag: push为O(1), 其余为O(logn)
102 • thin heap tag: push为O(1), 不支持join, 其余为O(logn); 但是如果只有increase key, modify
       均摊0(1)
103 • 不支持不是不能用, 而是用起来很慢
104 经过实践检测得到的结论:
105 • Dijkstra算法中应用pairing_heap_tag,速度与手写数据结构相当。
106 • binary_heap_tag在绝大多数情况下优于std::priority_queue
   • pairing_heap_tag 在 绝 大 多 数 情 况 优 于 binomial_heap_tag 和 rc_binomial_heap_tag
107
108
   • 只有push, pop和join操作时, binary_heap_tag速度较快
109
   • 有modify操作时,可以考虑thin_heap_tag或者pairing_heap_tag,或手写数据结构。
110
111
   //hash\_table
112
   #include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/hash_policy.hpp>
114 using namespace __gnu_pbds;
   __gnu_pbds::cc_hash_table <Key, Mapped> mp; //使用链地址法解决哈希冲突
115
     _gnu_pbds::gp_hash_table <Key, Mapped> mp; //使用探测法解决哈希冲突
116
117 //用法和map一样
```