SOUTH CHINA UNIVERSITY OF TECHNOLOGY

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TEMPLATE



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1 Graph Theory

1.1 Shortest Path

1.1.1 Dijkstra

```
1 typedef long long LL;
const int MAXN = ;
3 const int MAXM = ;
4 const LL DINF = ;
5 typedef pair<LL, int> P;
6 struct Edge {
7
       int to, nxt;
8
       LL w;
9
   }e[MAXM];
int head[MAXN], ecnt;
11 LL d[MAXN];
12 priority_queue<P, vector<P>, greater<P> > q;
   inline void addEdge(int x, int y, LL w) {
13
14
       e[++ecnt] = (Edge) \{y, head[x], w\}; head[x] = ecnt;
15
   void dijkstra(int st, int n) {
16
17
        for(int i = 0; i <= n; i++) d[i] = DINF;</pre>
18
       d[st] = 0;
       q.push(make_pair(0, st));
19
20
       while(!q.empty()) {
            P x = q.top(); q.pop();
21
            int u = x.second;
22
            if(d[u] != x.first) continue;
23
            for(int i = head[u], v; i; i = e[i].nxt) {
24
25
                v = e[i].to;
26
                if(d[v] > d[u] + e[i].w) {
27
                    d[v] = d[u] + e[i].w;
28
                    q.push(make_pair(d[v], v));
29
                }
30
           }
31
       }
32
```

1.1.2 SPFA

```
1
   struct Edge {
2
       int to, nxt;
        LL w;
3
   }e[MAXE];
4
   int head[MAXN], ecnt;
   LL d[MAXN];
6
   bool exist[MAXN];
7
   queue<int> q;
8
   inline void addEdge(int x, int y, LL w) {
9
       e[++ecnt] = (Edge) \{y, head[x], w\}; head[x] = ecnt;
10
11
12
   void SPFA(int st) {
13
       memset(d,0x3f,sizeof(d));
14
        d[st] = 0;
        q.push(st);
15
16
        exist[st] = 1;
17
       while(!q.empty()) {
```

```
18
            int u = q.front(); q.pop();
            exist[u] = 0;
19
20
            for(int i = head[u], v; i; i = e[i].nxt) {
                 v = e[i].to;
21
                 if(d[v] > d[u] + e[i].w) {
22
                     d[v] = d[u] + e[i].w;
23
                     //pre[v] = u;
24
                     if(!exist[v]) {
25
                         q.push(v);
26
27
                         exist[v] = 1;
                     }
28
                 }
29
30
            }
31
        }
32
   }
```

1.2 Johnson

```
1
   void johnson() {
2
       //建图中, Edge需要from, w1, w2, 去掉<math>w;
       spfa(1);
3
4
       for(int u = 1; u <= n; u++)</pre>
            for(int i = head[u]; i; i = e[i].nxt)
5
                e[i].w2 = e[i].w1 + d[e[i].from] - d[e[i].to];
6
7
       dijkstra(s,n);
  }
8
```

1.2.1 K Shortest Path (A*)

```
//可重复走同一条边 利用了反向边表示,ecnt初始化为1
    //调用 dijkstra (ed, n) 跑反向图,注意 if (i& 1) {松弛}
3
   int shrt[MAXN];
   LL A_star(int st, int ed, int k, int n) {
4
5
       if(d[st] == d[0]) return -1;
6
        if(st == ed) k++;
7
        q.push(make_pair(d[st], st));
8
        while(!q.empty()) {
           P x = q.top(); q.pop();
9
            int u = x.second;
10
           LL xd = x.first;
11
            ++shrt[u];
12
13
            if(u == ed) {
14
                if(shrt[ed] == k) return xd;
15
           for(int i = head[u], v; i; i = e[i].nxt)
16
17
           if(!(i & 1)) {
18
                v = e[i].to;
                if(shrt[v] < k)q.push(make_pair(xd-d[u]+e[i].w+d[v], v));</pre>
19
           }
20
       }
21
22
        return -1;
23
   }
```

1.2.2 K Shortest Path (Protractable Heap)

```
//可重复走同一条边
1
2 typedef double LD;
3 const int MAXN = ;
4 const int MAXM = ;
   const int MAXLT = MAXM * 20;
5
   const LD DINF = ;
6
7
   const LD eps = ;
   namespace LT{
9
        int tcnt;
10
        int ls[MAXLT], rs[MAXLT], dis[MAXLT];
11
        int to[MAXLT];
12
        LD val[MAXLT];
13
14
15
        inline int newnode(LD w, int _to) {
16
            ++tcnt;
            ls[tcnt] = rs[tcnt] = 0; dis[tcnt] = 1;
17
            val[tcnt] = w; to[tcnt] = _to;
18
19
            return tcnt;
20
        inline int copynode(int id) {
21
22
            ++tcnt;
            ls[tcnt] = ls[id]; rs[tcnt] = rs[id]; dis[tcnt] = dis[id];
23
            val[tcnt] = val[id]; to[tcnt] = to[id];
24
25
            return tcnt;
26
        void push_up(int x) {
27
28
            if(dis[ls[x]] < dis[rs[x]]) swap(ls[x], rs[x]);</pre>
29
            dis[x] = dis[rs[x]] + 1;
30
        int merge(int x, int y) {
31
            if(!x || !y) return x^y;
32
            if(val[x] - val[y] > eps) swap(x, y);
33
34
            int p = copynode(x);
35
            rs[p] = merge(rs[p],y);
36
            push_up(p);
37
            return p;
        }
38
39
   }
40
  int rt[MAXN];
41
42 typedef pair<LD, int> P;
   struct Edge {
43
        int to, nxt;
44
        LD w;
45
   }e[MAXM];
46
   int head[MAXN], ecnt;
47
   int stan, sta[MAXN], fa[MAXN];
48
   int vis[MAXN], cov[MAXM];
49
50
   LD d[MAXN];
   priority_queue<P, vector<P>, greater<P> > q;
51
   inline void addEdge(int x, int y, LD w) {
   e[++ecnt] = (Edge) {y, head[x], w}; head[x] = ecnt;
52
53
54
   void init(int n, int m) {
55
        ecnt = 1; stan = 0;
56
        for(int i = 1; i <= n; i++)</pre>
57
            head[i] = cov[i] = fa[i] = 0;
58
        for(int i = 1; i <= m; i++) {</pre>
59
            int u, v; LD w;
60
            scanf("%d%d%lf", &u,&v,&w);
```

```
62
             addEdge(u, v, w);
63
            addEdge(v, u, w);
        }
64
65
    void dijkstra(int st, int n) {
66
        for(int i = 0; i \le n; i++) {d[i] = DINF; vis[i] = 0;}
67
        while(!q.empty()) q.pop();
68
69
        d[st] = 0;q.push(make_pair(0, st));
        while(!q.empty()) {
70
            P x = q.top(); q.pop();
71
            int u = x.second;
72
            if(vis[u]) continue;
73
74
            vis[u] = 1;
             for(int i = head[u], v; i; i = e[i].nxt)
75
76
                 if(i & 1) {
                     v = e[i].to;
77
                     if(d[v]-(d[u] + e[i].w) > eps) {
78
79
                         d[v] = d[u] + e[i].w;
                         q.push(make_pair(d[v], v));
80
                     }
81
                 }
82
        }
83
    }
84
    void buildT(int u) {
85
        sta[++stan] = u;
86
        vis[u] = 1;
87
88
        for(int i = head[u], v; i; i = e[i].nxt)
89
        if(i & 1) {
90
            v = e[i].to;
            if(fabs(d[v] - (d[u] + e[i].w)) < eps && !vis[v]) {
91
                 fa[v] = u; cov[i^1] = 1; buildT(v);
92
93
        }
94
95
    void buildH(int st, int n) {
96
        buildT(st);
97
        for(int i = 2, u, v; i <= ecnt; i += 2) {
98
             if(!cov[i]) {
99
                 u = e[i^1].to; v = e[i].to;
100
                 if(fabs(d[u] - d[0]) < eps | | fabs(d[v] - d[0]) < eps) continue;
101
102
                 rt[u] = LT::merge(rt[u], LT::newnode(d[v]+e[i].w-d[u], v));
103
            }
104
        for(int i = 2, u; i <= n; i++)
105
            if(fa[u = sta[i]]) rt[u] = LT::merge(rt[u], rt[fa[u]]);
106
107
108
    //求前k短路径,其和不超过W,问最大k int ans = 0;
    //求第 k短路径,LD ans = 0;
109
    void getKth(int st, int ed, LD W) {
110
        while(!q.empty()) q.pop();
111
112
        //最短路要记入答案,注意d[st]-Weps与W-d[st]<eps
        //! if(d[st] - W > eps) return;
113
        //! else\{W-=d/st\}; ++ans;\}
114
        //注意st和ed相同
115
        //\#if(st == ed) k++;
116
        //\#if(d[st] == d[0]) \{ans = -1; return;\}
117
        //\#if(--k == 0) \{ans = d[st]; return;\}
118
119
        int u = rt[st], v;
120
121
        if(u) q.push(make_pair(LT::val[u], u));
122
        while (!q.empty()) {
```

```
u = q.top().second; LD cur = q.top().first;
123
124
             q.pop();
125
             //! if (cur + d[st] - W > eps) break;
126
             //! else \{W-= cur + d[st]; ++ans;\}
             //\#if (--k == 0) \{ans = cur + d/st\}; break;\}
127
             v = LT::ls[u];
128
129
             if (v) q.push(make_pair(cur - LT::val[u] + LT::val[v], v));
130
             v = LT::rs[u];
             if (v) q.push(make_pair(cur - LT::val[u] + LT::val[v], v));
131
             v = rt[LT::to[u]];
132
             if (v) q.push(make_pair(cur + LT::val[v], v));
133
        }
134
135
    void sol() {
136
        int n, m, st = , ed = ;
137
         //LD W; int k;
138
         init(n, m);
139
140
         dijkstra(ed, n);
141
         for(int i = 0; i \le n; i++) vis[i] = 0;
         buildH(ed, n);
142
         getKth(st, ed, W);
143
        printf("%d\n", ans);
144
145
    }
```

1.3 Network Flow

1.3.1 ISAP

```
namespace NWF {
1
2
        struct Edge{
             int to, nxt;LL f;
3
        }e[MAXM << 1];
4
5
        int S, T, tot;
        int ecnt, head[MAXN], cur[MAXN], pre[MAXN], num[MAXN], dis[MAXN];
6
7
        queue<int> q;
        void init(int _S, int _T, int _tot){
8
9
             ecnt = 1; S = \_S; T = \_T; tot = _tot;
             memset(num, 0, (tot + 1) * sizeof(int));
10
             memset(head, 0, (tot + 1) * sizeof(int));
11
12
        inline void addEdge(int u, int v, LL f) {
13
             e[++ecnt] = (Edge) \{v, head[u], f\}; head[u] = ecnt; e[++ecnt] = (Edge) \{u, head[v], 0\}; head[v] = ecnt;
14
15
16
        void bfs() {
17
             memset(dis, 0, (tot + 1) * sizeof(int));
18
             q.push(T);
19
20
             dis[T] = 1;
             while(!q.empty()) {
21
                 int u = q.front(), v; q.pop();
22
23
                 num[dis[u]]++;
                 for(int i = cur[u] = head[u]; i; i = e[i].nxt) {
24
                      if(!dis[v = e[i].to]) {
25
                          dis[v] = dis[u] + 1;
26
27
                           q.push(v);
                      }
28
29
                 }
             }
30
31
```

```
LL augment() {
32
33
             LL flow = INF;
             for(int i = S; i != T; i = e[cur[i]].to)
34
             flow = min(flow, e[cur[i]].f);
for(int i = S; i != T; i = e[cur[i]].to) {
    e[cur[i]].f -= flow;
35
36
37
38
                  e[cur[i] \land 1].f += flow;
39
             return flow;
40
41
         LL isap() {
42
             bfs();
43
             int u = S, v;
44
             LL flow = 0;
45
46
             while(dis[S] <= tot) {</pre>
                  if(u == T) {
47
                       flow += augment();
48
49
50
                  bool fg = 0;
51
                  for(int i = cur[u]; i; i = e[i].nxt) {
52
                       if(e[i].f && dis[u] > dis[v = e[i].to]) {
53
                           pre[v] = u;
54
                           cur[u] = i;
55
                           u = v;
56
57
                           fg = 1;
                           break;
58
                       }
59
60
                  if(fg) continue;
61
                  if(!--num[dis[u]]) break;
62
                  int maxDis = tot;
63
                  for(int i = head[u]; i; i = e[i].nxt) {
64
                       if(e[i].f \&\& maxDis > dis[v = e[i].to]) {
65
                           maxDis = dis[v];
66
                           cur[u] = i;
67
                       }
68
                  }
69
70
                  num[dis[u] = maxDis + 1]++;
71
                  if(u != S) u = pre[u];
72
73
             return flow;
74
        }
75
```

1.3.2 HLPP

```
namespace NWF{
1
2
         struct Edge{
3
             int to,nxt;LL f;
4
         }e[MAXM << 1];
         int S, T, tot;
int ecnt, head[MAXN], dis[MAXN], num[MAXN];
5
6
         LL sumf[MAXN];
7
8
         queue<int> q;
         list<int> dep[MAXN];
9
         void init(int _S,int _T,int _tot){
10
             ecnt = 1;S = _S;T = _T;tot = _tot;
memset(num, 0, (tot + 1) * sizeof(int));
11
12
             memset(head, 0, (tot + 1) * sizeof(int));
13
```

```
memset(sumf, 0, (tot + 1) * sizeof(LL));
14
15
16
         void addEdge(int u,int v,LL f){
              \begin{array}{l} \texttt{e[++ecnt]} = (\texttt{Edge}) \; \{\texttt{v}, \; \texttt{head[u]}, \; \texttt{f}\}; \texttt{head[u]} = \texttt{ecnt}; \\ \texttt{e[++ecnt]} = (\texttt{Edge}) \; \{\texttt{u}, \; \texttt{head[v]}, \; \texttt{0}\}; \texttt{head[v]} = \texttt{ecnt}; \\ \end{array}
17
18
19
         void bfs(){
20
              memset(dis, 0, (tot + 1) * sizeof(int));
21
              q.push(T); dis[T] = 1;
22
23
              while(!q.empty())
                   int u=q.front(), v; q.pop();
24
                   for(int i = head[u]; i; i = e[i].nxt)
25
26
                   if(!dis[v = e[i].to]){
                        dis[v] = dis[u] + 1;
27
28
                        q.push(v);
29
                   }
              }
30
31
32
         LL hlpp(){
33
              bfs();
              dis[S] = tot + 1;
34
              for(int i = 1;i <= tot; ++i)num[dis[i]]++;</pre>
35
              for(int i = tot + 1; ~i; --i)dep[i].clear();
36
37
              int maxd = dis[S];LL f;
              dep[maxd].push_back(S);sumf[S] = INF;
38
39
              for(;;){
40
                   while(maxd && dep[maxd].empty())maxd--;
41
                   if(!maxd)break;
                   int u = dep[maxd].back(), v;dep[maxd].pop_back();
42
                   int minDis = tot + 1;
43
                   for(int i = head[u]; i;i = e[i].nxt)
44
                   if(e[i].f){
45
                        if(dis[u] > dis[v = e[i].to]){
46
                             f = min(sumf[u], e[i].f);
47
                             e[i].f -= f;e[i^1].f += f;
48
                             if(sumf[u] != INF) sumf[u] -= f;
49
                             if(sumf[v] != INF) sumf[v] += f;
50
                             if(v!=S \&\& v!=T \&\& sumf[v] == f){
51
                                  maxd = max(maxd, dis[v]);
52
53
                                  dep[dis[v]].push_back(v);
54
                             if(!sumf[u])break;
55
                        }else minDis=min(minDis, dis[v] + 1);
56
57
                   if(sumf[u]){
58
                        if(!--num[dis[u]]){
59
                             for(int i = dis[u];i <= maxd;++i){</pre>
60
61
                                  while(!dep[i].empty()){
                                        --num[i];
62
                                       dis[dep[i].back()] = tot + 1;
63
64
                                       dep[i].pop_back();
65
66
                             maxd = dis[u] - 1; dis[u] = tot + 1;
67
                        }else{
68
                             dis[u] = minDis;
69
                             if(minDis > tot)continue;
70
71
                             num[minDis]++;
                             maxd = max(maxd, minDis);
72
                             dep[minDis].push_back(u);
73
74
```

1.3.3 Dinic

```
namespace NWF {
1
2
         struct Edge {
3
             int to, nxt;LL f;
         e[MAXM << 1];
4
5
         int S, T, tot;
         int ecnt, head[MAXN], cur[MAXN], dis[MAXN];
6
7
         queue<int> q;
        void init(int _S, int _T, int _tot){
8
             ecnt = 1; S = \_S; T = \_T; tot = _tot;
9
             memset(head, 0, (tot + 1) * sizeof(int));
10
11
        void addEdge(int u, int v, LL f) {
    e[++ecnt] = (Edge) {v, head[u], f}; head[u] = ecnt;
    e[++ecnt] = (Edge) {u, head[v], 0}; head[v] = ecnt;
12
13
14
15
         bool bfs() {
16
             memset(dis, 0, (tot + 1) * sizeof(int));
q.push(S); dis[S] = 1;
17
18
             while (!q.empty()) {
19
                  int u = q.front(), v; q.pop();
20
                  for (int i = cur[u] = head[u]; i ; i = e[i].nxt) {
21
                       if (e[i].f && !dis[v = e[i].to]) {
22
                           q.push(v);
23
                           dis[v] = dis[u] + 1;
24
25
                       }
26
                  }
             }
27
             return dis[T];
28
29
         LL dfs(int u, LL maxf) {
30
             if (u == T) return maxf;
31
             LL sumf = maxf;
32
33
             for (int &i = cur[u]; i; i = e[i].nxt) {
34
                  if (e[i].f && dis[e[i].to] > dis[u]) {
35
                       LL tmpf = dfs(e[i].to, min(sumf, e[i].f));
36
                       e[i].f -= tmpf; e[i ^ 1].f += tmpf;
                       sumf -= tmpf;
37
38
                       if (!sumf) return maxf;
                  }
39
             }
40
             return maxf - sumf;
41
42
         LL dinic() {
43
             LL ret = 0;
44
             while (bfs()) ret += dfs(S, INF);
45
46
             return ret;
47
         }
48
```

1.3.4 MCMF

```
namespace NWF{
1
2
        struct Edge {
3
             int to, nxt;LL f, c;
4
        e[MAXM << 1];
        int S, T, tot;
int ecnt, head[MAXN], cur[MAXN];LL dis[MAXN];
5
6
7
        bool exist[MAXN];
        queue<int> q;
8
        void init(int _S, int _T, int _tot){
   ecnt = 1; S = _S; T = _T; tot = _tot;
9
10
             memset(head, 0, (tot + 1) * sizeof(int));
11
12
        void addEdge(int u, int v, LL f, LL c) {
13
             e[++ecnt] = (Edge) \{v, head[u], f, c\}; head[u] = ecnt;
14
             e[++ecnt] = (Edge) \{u, head[v], 0, -c\}; head[v] = ecnt;
15
16
        bool spfa() {
17
             for(int i = 0;i <= tot; ++i){</pre>
18
19
                 dis[i] = INF; exist[i] = cur[i] = 0;
20
             q.push(S);dis[S] = 0;exist[S] = 1;
21
22
             while(!q.empty()) {
                 int u = q.front(), v; q.pop();exist[u] = 0;
23
                 for(int i = head[u]; i; i = e[i].nxt) {
24
25
                      if(e[i].f && dis[v = e[i].to] > dis[u] + e[i].c) {
                          dis[v] = dis[u] + e[i].c;
26
                          cur[v] = i;
27
28
                          if(!exist[v]) {
                               q.push(v);
29
30
                               exist[v] = 1;
                          }
31
                      }
32
                 }
33
             }
34
35
             return dis[T] != INF;
36
        LL mcmf() {
37
             LL cost = 0;
38
39
             while(spfa()) {
40
                 LL flow = INF;
                 for(int i = T; i != S; i = e[cur[i] ^ 1].to)
41
                      flow = min(flow, e[cur[i]].f);
42
                 for(int i = T; i != S; i = e[cur[i] ^ 1].to) {
43
                      e[cur[i]].f -= flow;
44
                      e[cur[i] \land 1].f += flow;
45
46
                 cost += flow * dis[T];
47
48
49
             return cost;
50
        }
51
    }
```

1.4 Tree Related

1.4.1 Union Set

```
int fa[MAXN], rnk[MAXN];
int Find(int x) { return x == fa[x] ? x : fa[x] = Find(fa[x]); }
bool same(int x, int y){ return Find(x) == Find(y); }
```

```
void unite(int x, int y)
4
5
   {
6
        x = Find(x);
7
        y = Find(y);
        if(x == y) return;
8
9
        if(rnk[x] < rnk[y]) {
10
            fa[x] = y;
11
        else {
12
            fa[y] = x;
13
            if(rnk[x] == rnk[y]) rnk[x]++;
14
        }
15
   }
16
```

1.4.2 Kruskal

```
1
   namespace MST{
2
        struct Edge{
3
            int u,v; LL w;
4
            bool operator < (const Edge& x) const { return w < x.w; }</pre>
5
        }e[MAXM];
6
        int ecnt, fa[MAXN];
7
        void addEdge(int u, int v, LL w) {
8
            e[++ecnt] = (Edge)\{v, u, w\}; //headp[u] = ecnt;
9
        int Find(int x) { return x == fa[x] ? x : fa[x] = Find(fa[x]); }
10
        LL kruskal(int n) {
11
            sort(e + 1, e + ecnt + 1);
12
            for(int i = 1; i <= n; i++) fa[i] = i;
13
            LL sum = 0;
14
            for (int i = 1; i <= ecnt; i++){</pre>
15
                 int fu = Find(e[i].u), fv = Find(e[i].v);
16
                 if(fu != fv){
17
                     fa[fu] = fv;
18
19
                     sum += e[i].w;
20
21
22
            return sum;
23
        }
   }
24
```

1.4.3 Prim

```
1
    namespace MST {
2
         struct Edge{
3
              int to,nxt; LL w;
         }e[MAXM];
4
         int ecnt, head[MAXN], vis[MAXN]; // pre[MAXN];
5
         LL dis[MAXN];
6
7
         void addEdge(int u, int v, LL w){
              e[++ecnt] = (Edge)\{v, head[u], w\}; head[u] = ecnt; e[++ecnt] = (Edge)\{u, head[v], w\}; head[v] = ecnt;
8
9
10
         LL Prim(int n){
11
              for (int i = 1; i \le n; i++){
12
                   //pre[i] = 0;
13
                   vis[i] = 0;
14
                   dis[i] = INF;
15
```

```
16
             vis[1] = 1;
17
             LL sum = 0;
18
             for (int i = head[1]; i; i = e[i].nxt)
19
                  dis[e[i].to] = min(dis[e[i].to],e[i].w);
20
             for (int j = 1; j < n; j++){
21
22
                  int u; LL minDis = INF;
                  for (int i = 1; i <= n; ++i)
    if (!vis[i] && dis[i] < minDis){</pre>
23
24
                           minDis = dis[i];
25
                           u = i;
26
27
                  if (minDis == INF) return -1;
28
29
                  vis[u] = 1;
30
                  sum += minDis;
                  for (int i = head[u], v; i; i = e[i].nxt)
31
                  if (!vis[v = e[i].to] && e[i].w < dis[v]){</pre>
32
33
                       //pre[u] = v;
34
                       dis[v] = e[i].w;
35
             }
36
             return sum;
37
38
        }
39
    }
```

1.4.4 Spanning Tree Calculation

关联矩阵 B:n*m 的矩阵, 其中 ek=(vi,vj), Bik 和 Bjk 一个为 1 一个为-1,第 k 列其他元素为 0 度数矩阵 D:n*n 的矩阵, 其中 i!=j 时,D[i][j]=0;i==j 时,D[i][j]=vi 的度邻接矩阵 A:n*n 的矩阵,vi,vj 有边相连,为 1,否则为 0 Kirchhoff 矩阵: B*BT=D-A 生成树个数: Kirchhoff 矩阵 n-1 阶主子式的行列式值构造 Kirchhoff 矩阵,调用 det(n)

```
LL a[MAXN][MAXN];
1
2
   void getSTC(int n, int m) {
3
        for(int i = 1; i <= n; i++) {
            for(int j = 1; j \le n; j++)
4
5
                 a[i][j] = a[j][i] = 0;
6
        for(int i = 1,u, v; i <= m; i++) {</pre>
7
            scanf("%d%d", &u,&v);
8
            if(u == v) continue;
9
            a[u][v] = --a[v][u];
10
11
        for(int i = 1; i <= n; i++) {</pre>
12
            int t = 0;
13
            for(int j = 1; j \le n; j++)
14
                 t += a[i][j];
15
16
            a[i][i] = -t;
17
        LL ans = det(n);
18
    }
19
```

1.4.5 Minimum Spanning Tree Calculation

```
typedef long long LL;
const int MAXN = ;
const int MAXM = ;
int sum,ans1,ans2=1, Mod=;
```

```
int fa1[MAXN],fa2[MAXN];
   bool vis[MAXN];
6
   struct Edge{int v,u,val;}e[MAXM];
7
   bool cmp(Edge A,Edge B){return A.val<B.val;}</pre>
8
   int getfa1(int *fa,int x){return fa[x]=fa[x]==x?x:getfa1(fa,fa[x]);}
   int getfa2(int *fa,int x){return fa[x]==x?x:getfa2(fa,fa[x]);}
10
    void dfs(int tot,int l,int r){
11
        if(tot==0){++sum; return;}
12
        for(int i=l,fx,fy;i<r;++i)</pre>
13
             if(!vis[i]){
14
15
                 vis[i]=true;
                 fx=getfa2(fa2,e[i].u);fy=getfa2(fa2,e[i].v);
16
                 if(fx!=fy){
17
                     fa2[fx]=fy;
18
19
                     dfs(tot-1,i+1,r);
20
                     fa2[fx]=fx;
21
22
                 vis[i]=false;
23
            }
24
   void sol(){
25
        int n,m;
26
        scanf("%d %d",&n,&m);
27
        for(int i=1;i<=n;++i)fa1[i]=fa2[i]=i;</pre>
28
29
        for(int i=1;i<=m;++i)</pre>
30
            scanf("%d %d %d",&e[i].v,&e[i].u,&e[i].val);
31
        std::sort(e+1,e+m+1,cmp);e[m+1].val=-1;
32
        for(int i=1, j=1, fx, fy, tot; i <= m; ++ i){</pre>
33
             for(;e[i].val==e[j].val;++j);
34
            tot=0;
             for(int k=i;k<j;++k){</pre>
35
                 fx=getfa1(fa1,e[k].u);fy=getfa1(fa1,e[k].v);
36
                 if(fx!=fy){++tot;ans1++;fa1[fx]=fy;}
37
38
39
            if(!tot)continue;
            sum=0;dfs(tot,i,j);
40
             (ans2*=sum)\%=Mod;
41
            for(int k=i;k<j;++k){</pre>
42
                 fx=getfa1(fa2,e[k].u);fy=getfa1(fa2,e[k].v);
43
44
                 if(fx!=fy)fa2[fx]=fy;
45
            }
46
        if(ans1!=n-1)puts("0");else printf("%d",ans2);
47
48
```

1.4.6 Tree Divide and Conquer

```
1
   struct Edge {
2
       int to, nxt, w;
   }e[MAXM];
3
   int head[MAXN], ecnt;
   int sz[MAXN];
5
   int d[MAXN], t[5], ans;
6
   bool vis[MAXN];
7
   inline void add_edge(int u, int v, int w) {
8
        e[++ecnt] = (Edge) \{v, head[u], w\}; head[u] = ecnt;
9
        e[++ecnt] = (Edge) \{u, head[v], w\}; head[v] = ecnt;
10
   }
11
12
   int getsz(int x, int fa) {
       sz[x] = 1;
```

```
for(int i = head[x]; i; i = e[i].nxt) {
14
            int y = e[i].to;
15
16
            if(vis[y] || y == fa) continue;
17
            sz[x] += getsz(y, x);
18
19
        return sz[x];
20
   }
21
    int getrt(int x) {
        int tot = getsz(x, 0) >> 1;
22
        while(1) {
23
24
            int u = -1;
            for(int i = head[x]; i; i = e[i].nxt) {
25
                 int y = e[i].to;
26
                 if(vis[y] || sz[y] > sz[x]) continue;
27
28
                 if(u == -1 \mid | sz[y] > sz[u]) u = y;
29
            if(\sim u \&\& sz[u] > tot) x = u;
30
31
            else break;
32
        }
33
        return x;
34
   void getdep(int x, int fa) {
35
36
        t[d[x]]++;
        for(int i = head[x]; i; i = e[i].nxt) {
37
38
            int y = e[i].to;
            if(vis[y] || y == fa) continue;
39
40
            d[y] = (d[x] + e[i].w) % 3;
41
            getdep(y, x);
        }
42
43
    int cal(int x, int v) {
44
        t[0] = t[1] = t[2] = 0;
45
        d[x] = v \% 3;
46
        getdep(x, 0);
47
        return t[0] * t[0] + t[1] * t[2] * 2;
48
   }
49
   void solve(int x) {
50
        vis[x] = 1;
51
52
        ans += cal(x, 0);
53
        for(int i = head[x]; i; i = e[i].nxt) {
54
            int y = e[i].to;
55
            if(vis[y]) continue;
56
            ans -= cal(y, e[i].w);
57
            solve(getrt(y));
58
        }
59
   int main() {
60
61
        solve(getrt(1));
62
```

1.5 LCA

1.5.1 Tree Decomposition LCA

```
int sz[MAXN], dep[MAXN], top[MAXN], fa[MAXN], son[MAXN], num[MAXN], totw;
struct Edge {
   int to, nxt;
}e[MAXN << 1];
int head[MAXN], ecnt;</pre>
```

```
inline void add_edge(int x, int y) {
7
        e[++ecnt] = (Edge) \{y, head[x]\}; head[x] = ecnt;
8
   void dfs1(int x) {
9
        sz[x] = 1; son[x] = 0;
10
        for(int i = head[x]; i; i = e[i].nxt) {
11
            int v = e[i].to;
12
13
            if(v == fa[x]) continue;
            fa[v] = x;
14
            dep[v] = dep[x] + 1;
15
            dfs1(v);
16
            SZ[X] += SZ[V];
17
            if(sz[v] > sz[son[x]]) son[x] = v;
18
19
        }
20
   }
    void dfs2(int x) {
21
        B[num[x]] = A[x];
22
23
        if(son[x]) {
24
            top[son[x]] = top[x];
25
            num[son[x]] = ++totw;
            dfs2(son[x]);
26
27
        for(int i = head[x]; i; i = e[i].nxt) {
28
29
            int v = e[i].to;
            if(v == fa[x] | v == son[x]) continue;
30
31
            top[v] = v;
32
            num[v] = ++totw;
            dfs2(v);
33
        }
34
35
   int lca(int u, int v) {
36
37
        if(u == v) return u;
        while(top[u] != top[v]) {
38
            if(dep[top[u]] > dep[top[v]]) swap(u, v);
39
            v = fa[top[v]];
40
41
        if(dep[u] > dep[v]) swap(u, v);
42
43
        return u;
   }
44
45
   inline void init() {
46
        memset(head, 0, sizeof(head)); ecnt = 0;
47
        fa[1] = 0; dep[1] = 1; top[1] = 1; num[1] = 1; totw = 1;
   }
48
   inline void pre() {
49
50
        dfs1(1); dfs2(1);
51
```

1.5.2 Tarjan LCA

```
vector< pair<int,int> > G[MAXN],ask[MAXN];
   int fa[MAXN], ans[MAXN], vis[MAXN] ,dis[MAXN];
3
   int Find(int x){
       return x == fa[x] ? x : fa[x] = Find(fa[x]);
4
5
   void init(int n){
6
       memset(ans, 0,sizeof ans);
7
       memset(vis, 0,sizeof vis);
8
9
       for(int i = 0; i <= n; i++){
10
           G[i].clear();
11
           ask[i].clear();
```

```
12
       }
   }
13
   void LCA(int u){
14
        int v;
15
        fa[u] = u;
16
17
        vis[u] = true;
        for(auto it : ask[u])
18
            if(vis[v = it.first])
19
                ans[it.second] = dis[u] + dis[v] - 2 * dis[Find(it.first)];
20
21
        for(auto it : G[u])
        if(!vis[v = it.first]){
22
            dis[v] = dis[u] + it.second;
23
24
            LCA(v);
25
            fa[v] = u;
26
        }
27
   }
```

1.6 Tarjan

1.6.1 SCC

```
1
   namespace SCC{
2
        vector<int> G[MAXN];
3
        int dfs_clock, scc_cn, dfn[MAXN], low[MAXN], sccno[MAXN];
4
        stack<int> S;
        void addEdge(int u, int v) {
5
6
            G[u].push_back(v);
7
        void tarjan(int u) {
8
            dfn[u] = low[u] = ++dfs\_clock;
9
            S.push(u);
10
            for(auto v : G[u]) {
11
                 if(!dfn[v]) {
12
                     tarjan(v);
13
                     low[u] = min(low[u], low[v]);
14
                }else if(!sccno[v]) {
15
16
                     low[u] = min(low[u], dfn[v]);
17
18
            if(dfn[u] == low[u]) {
19
                 scc_cnt++;
20
21
                 for(;;) {
22
                     int v = S.top(); S.pop();
23
                     sccno[v] = scc_cnt;
24
                     if(v == u) break;
25
                }
26
            }
27
        void findSCC(int n) {
28
            for(int i = 1; i <= n; i++)
29
                 if(!dfn[i]) tarjan(i);
30
31
        void init(int n){
32
33
            dfs_clock = scc_cnt = 0;
            for(int i = 0;i <= n;++i){</pre>
34
                 dfn[i] = low[i] = sccno[i] = 0;
35
                 G[i].clear();
36
37
            }
38
```

39 **}**

1.6.2 BCC

```
namespace BCC{
1
2
        struct Edge {
3
            int to, nxt;
 4
        }e[MAXM << 1];
        int ecnt, head[MAXN];
5
        int dfs_clock, dfn[MAXN], low[MAXN];
6
7
        int is_vertex[MAXN], vbcc_cnt, vbccno[MAXN];
8
9
        vector<int> vbcc[MAXN];
10
        stack<int> vS;
11
12
        int ebcc_cnt, ebccno[MAXN];
13
        stack<int> eS;
14
15
        inline void addEdge(int u, int v) {
16
            e[++ecnt] = (Edge) \{v, head[u]\}; head[u] = ecnt;
17
            e[++ecnt] = (Edge) \{u, head[v]\}; head[v] = ecnt;
18
        inline void init(int n) {
19
            ecnt = 1;
20
21
            dfs\_clock = 0;
22
            vbcc\_cnt = 0;
23
            ebcc\_cnt = 0;
            for(int i = 1; i \le n; ++i){
24
25
                 head[i] = dfn[i] = low[i] = 0;
26
                 is_vertex[i] = 0;
27
                 vbccno[i] = 0;
28
                 ebccno[i] = 0;
29
30
            while(!vS.empty()) vS.pop();
31
        //root's edge = -1;
32
        void tarjan(int u, int edge) {
33
            dfn[u] = low[u] = ++dfs\_clock;
34
            int ch = 0;
35
            vS.push(u);
36
37
            eS.push(u);
            for(int i = head[u], v; i; i = e[i].nxt) {
38
                 if(!dfn[v = e[i].to]) {
39
                     tarjan(v, i ^ 1);
40
                     low[u] = min(low[u], low[v]);
41
                     if(low[v] >= dfn[u]) {
42
43
                         ++ch;
                         if(edge > 0 || ch > 1) is_vertex[u] = 1;
44
                         vbcc[++vbcc_cnt].clear();
45
                         vbcc[vbcc_cnt].push_back(u);
46
                         for(int x;;){
47
                              x = vS.top();vS.pop();
48
49
                              vbcc[vbcc_cnt].push_back(x);
                              vbccno[x] = vbcc_cnt;
50
                              if(x == v)break;
51
                         }
52
53
54
                     if(low[v] > dfn[u]) {
55
                                ^{\smallfrown} 1 is bridge
56
```

```
57
58
                 else if(dfn[v] < dfn[u] && i != edge)</pre>
59
                     low[u] = min(low[u], dfn[v]);
60
             if(dfn[u] == low[u]) {
61
                 ebcc_cnt++;
62
                 for(int v;;) {
63
                     v = eS.top(); eS.pop();
64
65
                     ebccno[v] = ebcc_cnt;
                     if(v == u) break;
66
67
            }
68
69
        void findBCC(int n){
70
            for(int i = 1; i <= n; i++)</pre>
71
                 if(!dfn[i]) tarjan(i, -1);
72
73
74
             //findBridge
75
             for(int u = 1; u <= n; u++) {
                 for(int i = head[u], v; i; i = e[i].nxt)
76
                 if(ebccno[u] != ebccno[v = e[i].to]) {
77
78
                     //is bridge
79
                 }
80
            }
81
        }
82
   }
```

1.7 Cactus

1.7.1 Circle-Square Tree

```
#include <bits/stdc++.h>
1
using namespace std;
3 typedef pair<int, int> P;
   const int MAXN = 2e4 + 5;
   const int S = 15;
6
   namespace Tree {
7
        struct Edge {
8
            int to, nxt, w;
        }e[MAXN << 1];</pre>
9
        int ecnt, head[MAXN];
10
        int rt, isrt[MAXN], fa[MAXN][S + 3];
11
        int sz[MAXN];
12
13
        inline void addEdge(int u, int v, int w) {
            e[++ecnt] = (Edge) \{v, head[u], w\}; head[u] = ecnt;
14
15
            fa[v][0] = u;
16
        }
17
18
   int n, m, Q;
   namespace BCC {
19
        struct Edge {
20
21
            int to, nxt, w;
        }e[MAXN << 1];</pre>
22
        int ecnt, head[MAXN];
23
        int dfs_clock, dfn[MAXN], low[MAXN];
24
25
        int is_vertex[MAXN], vbcc_cnt, vbccno[MAXN];
26
        vector<P> vbcc[MAXN];
27
        stack<P> vs;
        int tag[MAXN];
28
```

```
inline void addEdge(int u, int v, int w) {
29
            e[++ecnt] = (Edge) \{v, head[u], w\}; head[u] = ecnt;
30
31
            e[++ecnt] = (Edge) \{u, head[v], w\}; head[v] = ecnt;
32
        inline void init(int n) {
33
            ecnt = 1;
34
            dfs\_clock = 0;
35
36
            vbcc\_cnt = 0;
            for(int i = 0; i \le 2 * n; i++){
37
                 head[i] = dfn[i] = low[i] = 0;
38
                vbccno[i] = 0;
39
                tag[i] = 0;
40
41
42
            while(!vs.empty()) vs.pop();
43
        //root 's edge = -1;
44
        void tarjan(int u, int edge) {
45
            dfn[u] = low[u] = ++dfs\_clock;
46
            vs.push(P(u, e[edge ^ 1].w));
47
            for(int i = head[u], v; i; i = e[i].nxt) {
48
                 if(!dfn[v = e[i].to]) {
49
                     tarjan(v, i ^ 1);
50
                     low[u] = min(low[u], low[v]);
51
                     if(low[v] >= dfn[u]) {
52
                         if(vs.top().first == v) {
53
                              Tree::addEdge(u, v, vs.top().second);
54
55
                              vs.pop();
56
                              continue;
57
                         vbcc[++vbcc_cnt].clear();
58
                         vbcc[vbcc_cnt].push_back(P(u, 0));
59
                         Tree::isrt[u] = 1;
60
                         int &sz = Tree::sz[n + vbcc_cnt];
61
                         tag[vs.top().first] = n + vbcc_cnt;
62
                          //Tree::addEdge(u, rt, 0);
63
                         for(P x;;) {
64
                              x = vs.top(); vs.pop();
65
66
                              sz += x.second;
                              //Tree::addEdge(rt, x.first, sz);
67
68
                              vbcc[vbcc_cnt].push_back(x);
69
                              vbccno[x.first] = vbcc_cnt;
70
                              if(x.first == v) break;
                         }
71
                     }
72
73
74
                else if(dfn[v] < dfn[u] && i != edge)</pre>
75
                     low[u] = min(low[u], dfn[v]);
76
            for(int i = head[u], v; i; i = e[i].nxt) {
77
                 if(tag[v = e[i].to]) {
78
79
                     int r = tag[v]; Tree::sz[r] += e[i].w;
                     tag[v] = 0;
80
                }
81
            }
82
83
        void findBCC(int n) {
84
            for(int i = 1; i <= n; i++)</pre>
85
                 if(!dfn[i]) tarjan(i, -1);
86
87
        }
88
   namespace Tree {
```

```
int dis[MAXN], dep[MAXN], len[MAXN];
90
91
         inline void init(int n) {
             BCC::init(n);
92
93
             rt = n;
94
             ecnt = 1;
             for(int i = 0; i <= 2 * n; i++) {
95
                  head[i] = 0;
96
                  fa[i][0] = isrt[i] = dis[i] = dep[i] = len[i] = 0;
97
98
99
         void dfs(int x) {
100
             for(int i = head[x], y; i; i = e[i].nxt) {
   if(!dep[y = e[i].to]) {
101
102
                      dep[y] = dep[x] + 1;
103
104
                      dis[y] = dis[x] + e[i].w;
105
                      dfs(y);
106
                  }
             }
107
108
         void pre() {
109
             for(int k = 1; k <= BCC::vbcc_cnt; k++) {</pre>
110
                  rt++;
111
                  vector<P> &E = BCC::vbcc[k];
112
                  addEdge(E[0].first, rt, 0);
113
                  int cnt = 0;
114
                  for(int i = E.size() - 1; i >= 1; i--) {
115
116
                      cnt += E[i].second;
117
                      len[E[i].first] = cnt;
118
                      addEdge(rt, E[i].first, min(cnt, sz[rt] - cnt));
                  }
119
120
             for(int k = 1; k <= S; k++) {
121
                  for(int i = 1; i <= rt; i++) {</pre>
122
                      fa[i][k] = fa[fa[i][k - 1]][k - 1];
123
124
125
             dep[1] = 1;
126
             dfs(1);
127
128
129
         int up(int x, int d) {
130
             for(int i = S; i >= 0; i--) {
131
                  if(dep[fa[x][i]] >= d) x = fa[x][i];
             }
132
             return x;
133
134
         int lca(int u, int v) {
135
136
             if(dep[u] > dep[v]) swap(u, v);
137
             v = up(v, dep[u]);
             if(u == v) return u;
138
             for(int i = S; i >= 0; i--) {
139
                  if(fa[u][i] != fa[v][i]) {
140
                      u = fa[u][i], v = fa[v][i];
141
142
143
             return fa[u][0];
144
145
         int query(int u, int v) {
146
             int l = lca(u, v);
147
             if(l <= n) return dis[u] + dis[v] - 2 * dis[l];</pre>
148
             int x = up(u, dep[l] + 1), y = up(v, dep[l] + 1);
149
150
             int res = dis[u] - dis[x] + dis[v] - dis[y];
```

```
int tmp = abs(len[x] - len[y]);
151
             return res + min(tmp, sz[l] - tmp);
152
         }
153
154
    }
155
    int main() {
156
         ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
157
158
         using namespace Tree;
         cin >> n >> m >> Q;
159
         init(n);
160
         for(int i = 1, u, v, w; i <= m; i++) {</pre>
161
162
             cin >> u >> v >> w;
             BCC::addEdge(u, v, w);
163
164
         BCC::findBCC(n);
165
166
         pre();
167
         int u, v;
         while(Q--) {
168
169
             cin >> u >> v;
170
             cout << query(u, v) << endl;</pre>
171
         return 0;
172
    }
173
```

2 Data Structures

2.1 Basic Structures

2.1.1 RMQ

```
struct RMQ {
1
        int d[MAXN][S + 2];
2
        inline void init(int *a, int n) {
3
            for(int i = 1; i \le n; i++) d[i][0] = a[i];
4
5
            for(int k = 1; (1 << k) <= n; k++)
6
                for(int i = 1; i + (1 << k) - 1 <= n; i++)
                    d[i][k] = min(d[i][k - 1], d[i + (1 << (k - 1))][k - 1]);
7
8
9
        inline int query(int 1, int r) {
            if(l > r) swap(l, r);
10
            int k = 0;
11
12
            while((1 << (k + 1)) <= r - l + 1) k++;
            return min(d[l][k], d[r - (1 << k) + 1][k]);
13
14
15
   }rmq;
   const int MAXM = 2e5 + 5, MAXN = 3e6 + 5, S = 22;
16
   const LL INF = 1e18;
17
   #define belong(x) (x / S + 1)
18
   #define pos(x) (x % S + 1)
19
   int Log[MAXN], sz;
20
   struct RMQ {
21
        LL a[MAXN];
22
        LL d[MAXM][S + 2];
23
        LL pre[MAXM][S + 2], aft[MAXM][S + 2];
24
        inline void init(int n) {
25
            sz = n / S + 1;
26
27
            Log[0] = -1; for(int i = 1; i <= n; i++) Log[i] = Log[i / 2] + 1;
28
            for(int i = 1; i <= sz; i++) {
                pre[i][0] = aft[i][S + 1] = INF;
29
30
            }
            for(int i = 1; i <= n; i++) {</pre>
31
                pre[belong(i)][pos(i)] = min(pre[belong(i)][pos(i) - 1], a[i]);
32
33
            for(int i = n; i >= 1; i--) {
34
                aft[belong(i)][pos(i)] = min(aft[belong(i)][pos(i) + 1], a[i]);
35
36
37
            for(int i = 1; i <= sz; i++) {
38
                d[i][0] = aft[i][1];
39
            for(int k = 1; k \le S; k++)
40
                for(int i = 1; i + (1 << k) <= sz; i++)
41
                    d[i][k] = min(d[i][k - 1], d[i + (1 << (k - 1))][k - 1]);
42
43
        inline LL ask(int 1, int r) {
44
            assert(l <= r);</pre>
45
            LL res = INF;
46
            if(belong(l) == belong(r)) {
47
                for(int i = 1; i <= r; i++) res = min(res, a[i]);</pre>
48
49
50
            res = min(aft[belong(l)][pos(l)], pre[belong(r)][pos(r)]);
51
            int k = Log[belong(r) - belong(l) - 1];
52
            if(~k) {
53
```

```
res = min(res, d[belong(l) + 1][k]);
res = min(res, d[belong(r) - (1 << k)][k]);
return res;
return res;
}
return res;
}
return res;</pre>
```

2.1.2 Divide Blocks

```
int belong[MAXN], l[MAXN], r[MAXN];
   int sz, num;
   void build(int n) {
3
        sz = sqrt(n);
4
        num = n / sz; if(n % sz) num++;
5
        for(int i = 1; i <= num; i++) {</pre>
6
7
            l[i] = (i - 1) * sz + 1;
            r[i] = i * sz;
8
9
        }
10
        r[num] = n;
11
        for(int i = 1; i <= n; i++) {
12
            belong[i] = (i - 1) / sz + 1;
13
14
   }
```

2.2 Heap Structures

2.2.1 Leftist Tree

```
const int MAXN = ;
   namespace LeftistTree{
2
3
        int ls[MAXN], rs[MAXN];
        int dis[MAXN];
4
        int fg[MAXN], sfa[MAXN], rt[MAXN];//利用 rt 得到 堆根 节点
5
6
        int val[MAXN];
7
       void push_down(int x) {};
8
9
        void push_up(int x) {
10
            if(dis[ls[x]] < dis[rs[x]]) swap(ls[x], rs[x]);</pre>
11
            dis[x] = dis[rs[x]] + 1;
12
        int merge(int x, int y) {
13
            if(!x || !y) return x^y;
14
            if(val[x] > val[y] | | (val[x] == val[y] && x > y)) swap(x, y);
15
16
            push_down(x);
            rs[x] = merge(rs[x], y);
17
18
            push_up(x);
19
            return x;
20
        int getSfa(int x) {return sfa[x] == x ? x : sfa[x] = getSfa(sfa[x]);}
21
22
        int uni(int x, int y) {//返回合并后的根
            if(!x || !y) return x^y;
23
            if(fg[x] || fg[y]) return;
24
            x = getSfa(x); y = getSfa(y);
25
            if(x == y) return;
26
            int z = merge(x, y);
27
28
            return sfa[x] = sfa[y] = sfa[z] = z;
29
        }
       void uni2(int a, int b) {
```

```
31
           //val[a] \gg = 1;
           int c = merge(ls[a], rs[a]);
32
33
           ls[a] = rs[a] = dis[a] = 0;
34
           int a1 = merge(c, a);
           //val[b] >>= 1;
35
           c = merge(ls[b], rs[b]);
36
37
           ls[b] = rs[b] = dis[b] = 0;
38
           int b1 = merge(c, b);
           c = merge(a1, b1);
39
           sfa[a] = sfa[b] = sfa[c] = c;
40
           printf("%d\n", val[c])
41
42
       int pop(int x) {//返回堆顶值,也可以用于返回根
43
           if(!x || fg[x]) return -1;
44
45
           x = getSfa(x); fg[x] = 1;
           push_down(x);//在删除堆顶时要下传标记
46
           int y = merge(ls[x], rs[x]);
47
48
           sfa[x] = sfa[y] = y;
49
           return val[x];
50
       void init(int n) {
51
52
           for(int i = 1; i <= n; i++) {
53
               sfa[i] = i;
               ls[i] = rs[i] = dis[i] = fg[i] = 0;
54
           }
55
56
       }
57
   //可持久化版本见k短路
```

2.3 Sequence Structures

2.3.1 Cartesian Tree

```
struct CartesianTree{
1
2
        int rt, fa[MAXN], ls[MAXN], rs[MAXN];
3
        int top, st[MAXN];
        int cnt[MAXN];
4
        void build(LL *a,int n) {
5
6
            top = rt = 0;
            for(int i = 1; i <= n; i++) {
7
                ls[i] = rs[i] = fa[i] = 0;
8
9
                while(top && a[st[top]] > a[i]) ls[i] = st[top--];
10
                fa[i] = st[top]
                if(ls[i]) fa[ls[i]] = i;
11
                if(fa[i]) rs[fa[i]] = i; else rt = i;
12
                st[++top] = i;
13
            }
14
15
        void dfs(int x) {
16
17
            cnt[x] = 1;
            if(ls[x]) {dfs(ls[x]); cnt[x] += cnt[ls[x]];}
18
19
            if(rs[x]) {dfs(rs[x]); cnt[x] += cnt[rs[x]];}
20
        LL getAns(LL *a, int n) {
21
22
            //dfs(rt);
23
24
            return res;
25
        }
26
   }T;
```

2.3.2 TreeArray

```
1  //树状数组上二分
2  int BS(int x) {
3    int res = 0;
4    for (int i = 1 << 18; i; i >>= 1)
5        if ((res | i) <= Tn && T[res | i] <= x)
            x -= T[res |= i];
7    return res;
8 }</pre>
```

2.3.3 Segment Tree

```
#define Ls(x) (x << 1)
1
2
   #define Rs(x) (x << 1 | 1)
3
   struct Tree {
        int l, r, lazy;
5
6
        LL sum, mx;
   }tree[MAXN << 2];</pre>
7
   int A[MAXN];
8
   void push_up(int x) {
9
       tree[x].sum = tree[Ls(x)].sum + tree[Rs(x)].sum;
10
       tree[x].mx = max(tree[Ls(x)].mx, tree[Rs(x)].mx);
11
12
13
   void push_down(int x) {
14
        if(tree[x].lazy) {
            tree[Ls(x)].sum += tree[x].lazy * (tree[Ls(x)].r - tree[Ls(x)].l + 1);
15
            tree[Rs(x)].sum += tree[x].lazy * (tree[Rs(x)].r - tree[Rs(x)].l + 1);
16
            tree[Ls(x)].mx += tree[x].lazy;
17
            tree[Rs(x)].mx += tree[x].lazy;
18
            tree[Ls(x)].lazy += tree[x].lazy;
19
            tree[Rs(x)].lazy += tree[x].lazy;
20
21
            tree[x].lazy = 0;
       }
22
23
   }
   void build(int x, int L, int R) {
24
25
       tree[x].lazy = 0;
26
        tree[x].l = L; tree[x].r = R;
27
        if(L == R) {
28
            tree[x].sum = A[L];
            tree[x].mx = A[L];
29
30
31
            return;
32
33
        int mid = (L + R) \gg 1;
        build(Ls(x), L, mid);
34
35
        build(Rs(x), mid + 1, R);
36
       push_up(x);
37
   void update(int x, int L, int R, LL val) {
38
        if(tree[x].l >= L && tree[x].r <= R) {
39
            tree[x].lazy += val;
40
            tree[x].sum += val * (tree[x].r - tree[x].l + 1);
41
            tree[x].mx += val;
42
            return;
43
44
       push_down(x);
45
        int mid = (tree[x].l + tree[x].r) >> 1;
46
```

```
if(L <= mid) update(Ls(x), L, R, val);</pre>
47
48
        if(R > mid) update(Rs(x), L, R, val);
49
        push_up(x);
50
    LL query(int x, int L, int R) {
   if(tree[x].l >= L && tree[x].r <= R)</pre>
51
52
             return tree[x].sum;
53
        push_down(x);
54
        int mid = (tree[x].l + tree[x].r) >> 1;
55
        LL res = 0;
56
        if(L <= mid) res += query(Ls(x), L, R);</pre>
57
        if(R > mid) res += query(Rs(x), L, R);
58
59
        return res;
60
61
62
    LL query2(int x, int L, int R) {
        if(tree[x].l >= L && tree[x].r <= R)
63
             return tree[x].mx;
64
        push_down(x);
65
66
        int mid = (tree[x].l + tree[x].r) >> 1;
67
        LL res = -INF;
        if(L \le mid) res = max(res, query2(Ls(x), L, R));
68
69
        if(R > mid) res = max(res, query2(Rs(x), L, R));
70
        return res;
71
```

2.3.4 LiChao Tree

```
const double eps = 1e-12;
2
    namespace LiT{
        const int MLIMIT = 40000;
3
        typedef double LD;
4
        struct line{LD k,b;int l,r,id;} T[MAXN << 2];</pre>
5
6
         //inline \; LD \; calc(line \; \&a, int \; pos) \; \{return \; a.k*vec[pos]+a.b; \}
        inline LD calc(line &a,int pos) {return a.k*pos+a.b;}
7
        inline double cross(line &a,line &b) {
8
9
             if(b.k == a.k) return -1e9;
             return (double)(a.b-b.b)/(b.k-a.k);
10
11
        void build(int v, int l, int r) {
12
13
             T[v].k = 0; T[v].b = -1e18;
14
             T[v].l = 0;T[v].r = MLIMIT;
             T[v].id = 0;
15
             if(l == r)return;
16
             int mid = (l+r)>>1;
17
             build(v<<1,1,mid);</pre>
18
19
             build(v<<1|1,mid+1,r);</pre>
20
        void ins(int v,int l,int r, line k) {
21
22
             if(k.1 \ll 1 \& r \ll k.r) {
                 LD fl = calc(k, l), fr = calc(k, r);

LD gl = calc(T[v], l), gr = calc(T[v], r);
23
24
                 if(fl - gl > eps \&\& fr - gr > eps) T[v] = k;
25
                 else if(fl - gl > eps || fr - gr > eps) {
26
                      int mid = (l+r)>>1;
27
                      if(calc(k, mid) - calc(T[v], mid) > eps) swap(k, T[v]);
28
                        /if(vec[mid] - cross(k, T[v]) > eps)
29
30
                      if(mid - cross(k, T[v]) > eps)
31
                          ins(v<<1, l, mid, k); else ins(v<<1|1, mid+1, r, k);
32
```

```
33
                  return;
34
             }
35
             int mid=(l+r)>>1;
             if(k.l <= mid) ins(v<<1, l, mid, k);</pre>
36
37
             if(mid < k.r) ins(v<<1|1, mid+1, r, k);</pre>
38
39
        LD ans; int ansid;
        void que(int v, int l, int r, int x) {
   LD tmp = calc(T[v], x);
40
41
             if(tmp > ans | I (tmp == ans \&\& T[v].id < ansid)) {
42
                  ans = tmp;
43
                  ansid = T[v].id;
44
45
             if(l == r) return;
46
             int mid = (l+r)>>1;
47
48
             if(x <= mid) que(v<<1,1,mid,x);else que(v<<1|1,mid+1,r,x);
        }
49
50
    };
    //左闭右闭
```

2.3.5 Splay Tree

```
1
    namespace splay{
2
        int n, m, sz, rt;
        int val[MAXN], id[MAXN];
3
        int tr[MAXN][2], size[MAXN], fa[MAXN], rev[MAXN], s[MAXN], lazy[MAXN];
4
        void push_up(int x) {
5
6
            int l = tr[x][0], r = tr[x][1];
7
            s[x] = max(val[x], max(s[l], s[r]));
8
            size[x] = size[l] + size[r] + 1;
9
        void push_down(int x) {
10
            int l = tr[x][0], r = tr[x][1];
11
            if(lazy[x]) {
12
13
                 if(1) {
                     lazy[l] += lazy[x];
14
                     s[l] += lazy[x];
15
                     val[l] += lazy[x];
16
17
                if(r) {
18
                     lazy[r] += lazy[x];
19
20
                     s[r] += lazy[x];
                     val[r] += lazy[x];
21
22
                lazy[x] = 0;
23
24
            if(rev[x]) {
25
26
                 rev[x] = 0;
27
                 rev[l] ^= 1; rev[r] ^= 1;
28
                 swap(tr[x][0], tr[x][1]);
29
30
        void rotate(int x, int &k) {
31
            int y = fa[x];
32
            int z = fa[y];
33
            int l, r
34
            if(tr[y][0] == x) l = 0;
35
36
            else l = 1;
37
            r = 1 \wedge 1;
38
            if(y == k) k = x;
```

```
39
            else {
                 if(tr[z][0] == y) tr[z][0] = x;
40
41
                else tr[z][1] = x;
42
            fa[x] = x; fa[y] = x; fa[tr[x][r]] = y;
43
            tr[y][l] = tr[x][r]; tr[x][r] = y;
44
            push_up(y); push_up(x);
45
46
        void splay(int x, int &k) {
47
            int y, z;
48
            while(x != k) {
49
                y = fa[x];
50
                z = fa[y];
51
52
                if(y != k) {
                     if((tr[y][0] == x) \land (tr[z][0] == y)) rotate(x, k);
53
54
                     else rotate(y, k);
55
56
57
                rotate(x, k);
            }
58
59
        int find(int x, int rank) {
60
            push_down(x);
61
62
            int l = tr[x][0], r = tr[x][1];
63
64
            if(size[l] + 1 == rank) return x;
65
            else if(size[l] >= rank) return find(l, rank);
66
            else return find(r, rank - size[l] - 1);
67
        void update(int l, int r, int v) {
68
            int x = find(rt, 1), y = find(rt, r + 2);
69
            splay(x, rt); splay(y, tr[x][1]);
70
            int z = tr[y][0];
71
            lazy[z] += v;
72
            val[z] += v;
73
            s[z] += v;
74
75
        void reverse(int l, int r) {
76
77
            int x = find(rt, 1), y = find(rt, r + 2);
78
            splay(x, rt); splay(y, tr[x][1]);
79
            int z = tr[y][0];
            rev[z] ^= 1;
80
81
        void query(int l, int r) {
82
83
            int x = find(rt, 1), y = find(rt, r + 2);
            splay(x, rt); splay(y, tr[x][1]);
84
            int z = tr[y][0];
85
            printf("%d\n", s[z]);
86
87
        void build(int 1, int r, int f) {
88
            if(l > r) return;
89
            int now = id[l], last = id[f];
90
            if(l == r) {
91
                fa[now] = last; size[now] = 1;
92
                if(l < f) tr[last][0] = now;</pre>
93
                else tr[last][1] = now;
94
95
                return;
96
            int mid = (l + r) \gg 1; now = id[mid];
97
            build(l, mid - 1, mid); build(mid + 1, r, mid);
98
99
            fa[now] = last;
```

```
100
             push_up(now);
             if(mid < f) tr[last][0] = now;</pre>
101
102
             else tr[last][1] = now;
103
        void init() {
104
             s[0] = -INF;
105
             scanf("%d%d", &n, &m);
106
             for(int i = 1; i <= n + 2; i++) id[i] = ++sz;</pre>
107
108
             build(1, n + 2, 0); rt = (n + 3) >> 1;
        }
109
110
    namespace splay{
111
         //内存回收池见fhq_treap
112
        int tcnt, root;
113
114
        int sz[MAXN]
        int tr[MAXN][2], fa[MAXN];
115
        int val[MAXN];
116
         int newnode(int w) {
117
             ++tcnt;
118
119
             sz[tcnt] = 1;
             fa[tcnt] = tr[tcnt][0] = tr[tcnt][1] = 0;
120
             //val[tcnt] = w;
121
             return tcnt;
122
123
        void push_up(int v) {
124
             int l = tr[v][0], r = tr[v][1];
125
126
             sz[v] = sz[l] + 1 + sz[r];
127
128
        void push_down(int v) {
129
             if(!v) return;
130
        void init() {
131
             tcnt = 2;
132
             tr[root = fa[1] = 2][0] = 1;
133
             sz[1] = 1; sz[2] = 2;
134
             //val[1] = -INF; val[2] = INF;//权值平衡树
135
             //val[1] = val[2] = 0;//位置平衡树
136
             //1,2为哨兵节点,根据题意也可设置为n+1,n+2或1,n+1
137
138
139
         void rotate(int x) {
140
             int y = fa[x], z = fa[y];
141
             push_down(y);push_down(x);
             int lr = tr[y][1] == x;
142
             if(z) tr[z][tr[z][1]==y] = x;
143
             fa[x] = z;
144
             fa[tr[y][lr] = tr[x][lr^1]] = y;
145
             fa[tr[x][lr^1] = y] = x;
146
147
             push_up(y); push_up(x);
148
149
         void splay(int x, int k) {
             for(int y, z; (y = fa[x]) != k; rotate(x)) {
150
                 if((z = fa[y]) != k) {
151
                     if((tr[y][0] == x) \wedge (tr[z][0] == y))
152
                          rotate(x); else rotate(y);
153
154
155
             if(!k) root = x;
156
157
         int find(int x, int rank) {
158
             push_down(x);
159
             int l = tr[x][0], r = tr[x][1];
160
```

```
if(sz[l] + 1 == rank) return x;
161
162
             if(sz[l] >= rank) return find(l, rank);
163
             return find(r, rank - sz[l] - 1);
164
         int build(int 1, int r) {
165
             if(l > r) return 0;
166
            if(1 == r) {
167
                 int num;scanf("%d", &num);
168
                 return newnode(num);
169
170
            int mid = (l + r) >> 1;
171
            int ls = build(l, mid-1);
172
            int num; scanf("%d", &num);
173
174
            int v = newnode(num);
175
            int rs = build(mid+1, r);
            if(ls) fa[ls] = v;
176
            tr[v][0] = ls;
177
            if(rs) fa[rs] = v;
178
179
            tr[v][1] = rs;
180
            push_up(v);
181
             return v;
182
        void insert(int pos, ...) {
183
            int x = find(root, pos+1), y = find(root, pos+2);
184
            splay(x, 0); splay(y, x);
185
186
             //int z = newnode(w); //插入一个节点
187
             //int z = build(1, n); //插入n个节点
188
            fa[tr[y][0] = z] = y;
189
            splay(z,0);
190
        void modifyOrQuery(int l, int r, int v) {
191
            int x = find(root, 1), y = find(root, r + 2);
192
             splay(x, 0); splay(y, x);
193
             int z = tr[y][0];
194
            if(!z) return;
195
             //标记对本身无效,处理时将2点重新计算
196
197
            splay(z,0);
198
        }
         void display(int v) {
199
200
            if(!v) return;
201
            push_down(v);
            display(tr[v][0]);
202
             if(val[v]) printf("%d ", val[v]);
203
204
            display(tr[v][1]);
205
         /*int find Value (int v)  {
206
207
             int res = root;
             for(int \ cur = root; cur; \ res = cur, \ cur = tr[cur][val[cur] <= v]);
208
209
             return res;
210
211
         void insert(int w) {
212
             int y = find Value(w);
213
             int z = newnode(w);
214
             fa[tr[y][val[y] \le w] = z] = y;
215
             splay(z,0);
         }*/
216
217
         /*void split(int v) {//splay维护区间[l,r],区间分裂为[l,k-1],[k,k],[k+1,r];
             //ump查看标号是否出现,mp维护子区间左端点
218
             if(ump.find(v) == ump.end()) {
219
                 auto\ it = mp.upper\_bound(v); --it;
220
221
                 int z = it -> second;
```

```
222
                      splay(z, 0);
223
                      int pos = sz[tr[z]/0]];
                      int \ x \, = \, find \, (\, root \, , \, pos) \, , \ y \, = \, find \, (\, root \, , \, pos + rc \, [\, z] - \, l\, c \, [\, z] + 2) \, ;
224
225
                      splay(x, 0); splay(y, x);
226
                      z = tr[y][0];
                      if(lc[z] != v)  {
227
228
                           tr[z][0] = newnode(lc[z], v-1);
229
                           fa[tr[z][0]] = z;
230
                           mp/lc/z// = tr/z//0/;
231
                      if(rc[z] != v)  {
232
                           tr[z][1] = newnode(v+1, rc[z]);
233
234
                           fa / tr / z / / 1 / / = z;
235
                          mp[v+1] = tr[z][1];
236
237
                      lc[z] = rc[z] = v;
                      splay(z,0);
238
239
                     mp / v / = z;
                     ump / v / = z;
240
241
           }*/
242
243
     }
244
```

2.3.6 Scapegoat Tree

```
struct ScapegoatTree{
1
        int Tsn; queue<int> q;
2
        int val[MAXM], ext[MAXM];
3
        int sz[MAXM], tsz[MAXM];
4
        int fa[MAXM], tr[MAXM][2];
5
        int root;
6
7
        double alp;
        void init() {
8
9
            root = 0;
10
            alp = 0.7;
11
        int newnode(int x) {
12
            if(q.empty()) q.push(++Tsn);
13
            int tcnt = q.front(); q.pop();
14
            val[tcnt] = x; ext[tcnt] = 1;
15
            fa[tcnt] = tr[tcnt][0] = tr[tcnt][1] = 0;
16
            sz[tcnt] = tsz[tcnt] = 0;
17
            return tcnt;
18
19
20
        void push_up(int v) {
21
            sz[v] = ext[v];tsz[v] = 1;
            if(tr[v][0]) {
22
                sz[v] += sz[tr[v][0]];
23
24
                tsz[v] += tsz[tr[v][0]];
25
            if(tr[v][1]) {
26
                sz[v] += sz[tr[v][1]];
27
28
                tsz[v] += tsz[tr[v][1]];
            }
29
30
        bool isBad(int v) {
31
32
            return (double(tsz[ tr[v][0] ]) > double(tsz[v]) * alp) ||
33
                 (double(tsz[ tr[v][1] ]) > double(tsz[v]) * alp) ||
34
                (sz[v] * 2 < tsz[v]);
```

```
35
36
        vector<int> vec;
37
        void rRecycle(int v) {
            if(tr[v][0]) rRecycle(tr[v][0]);
38
            if(ext[v]) vec.push_back(v); else q.push(v);
39
40
            if(tr[v][1]) rRecycle(tr[v][1]);
41
        int rBuild(int l, int r) {
42
            int mid = (l + r) \gg 1, v = vec[mid];
43
            tr[v][0] = (1 \le mid-1) ? rBuild(1, mid - 1) : 0;
44
            if(tr[v][0]) fa[tr[v][0]] = v;
45
            tr[v][1] = (mid+1 \ll r) ? rBuild(mid + 1, r) : 0;
46
            if(tr[v][1]) fa[tr[v][1]] = v;
47
            push_up(v);
48
49
            return v;
50
        void rebuild(int x) {
51
            int v = 0;
52
53
            for(;x; x= fa[x]) {
54
                 push_up(x);
                 if(isBad(x)) v = x;
55
56
            if(v && isBad(v)){
57
                vec.clear();
58
59
                 int u = fa[v], lr = tr[u][1] == v;
60
                 rRecycle(v);
61
                 if(vec.size()) v = rBuild(0, vec.size() - 1); else v = 0;
62
                if(u == 0) fa[root = v] = 0;
63
                 else{
                     tr[u][lr] = v;
64
65
                     if(v) fa[v] = u;
                }
66
            }
67
68
        void ins(int x) {
69
70
            int p = root, q = root;
            for(;p && val[p] != x; q = p, p = tr[p][val[p] < x]);
71
72
            if(!q) {
73
                 p = root = newnode(x);
74
            }else if(p) {
75
                ext[p]++;
76
            }else{
                 fa[p = tr[q][val[q] < x] = newnode(x)] = q;
77
78
            rebuild(p);
79
80
81
        void del(int x) {
82
            int p = root;
            for(;p && val[p] != x; p = tr[p][ val[p] < x]);</pre>
83
            if(p && ext[p]){
84
85
                 --ext[p];
                 rebuild(p);
86
            }
87
88
89
        int get_rank(int x) {
            int ret = 0;
90
91
            for(int p = root;p;) {
92
                 if(val[p] < x) {
                     ret += sz[tr[p][0]] + ext[p];
93
94
                     p = tr[p][1];
95
                 }else p = tr[p][0];
```

```
96
97
             return ret + 1;
98
99
         int get_Kth(int p, int k) {
             if(sz[tr[p][0]] >= k) return get_Kth(tr[p][0] ,k);
100
101
             k -= sz[tr[p][0]];
             if(ext[p] >= k) return val[p];
102
103
             k -= ext[p];
             return get_Kth(tr[p][1], k);
104
105
         int pre(int x) {
106
             int id = get_rank(x);
107
             return get_Kth(root, id - 1);
108
109
110
         int nxt(int x) {
             int id = get_rank(x + 1);
111
             return get_Kth(root ,id);
112
113
         void display(int v) {
114
             if(tr[v][0]) display(tr[v][0]);
115
             cerr<<val[v]<<" "
116
             if(tr[v][1]) display(tr[v][1]);
117
         }
118
    }T;
119
```

2.3.7 FHQ Treap

```
namespace fhq_treap{
1
2
       int Tsz; queue<int> q; //内存回收池
3
       int tcnt, root;
       //int rt[MAXN]; //可持久化时使用rt,维护版本号(int &root), 空间开大
4
       int sz[MAXN], rnd[MAXN];
5
       int tr[MAXN][2];
6
       //int fa [MAXN];//维护fa时除了在pushup更新v节点左右孩子父节点信息还要在split和merge结
7
       東时维护root的fa信息,fa[root]=0;
       int val[MAXN], rev[MAXN];
8
9
       void init() {
           srand(time(0));
10
           Tsz = tcnt = root = 0;
11
12
       int newnode(int v) {
13
           if(q.empty()) q.push(++Tsz);
14
           tcnt = q.front(); q.pop();
15
           sz[tcnt] = 1;
16
           rnd[tcnt] = rand();
17
           tr[tcnt][0] = tr[tcnt][1] = 0;
18
19
           //val[tcnt] = v;
20
           return tcnt;
21
22
       /*int copynode(int id) {
23
           //++tcnt;//获取一个新的节点编号
24
           sz/tcnt/ = sz/id/;
           rnd[tcnt] = rnd[id];
25
           tr[tent][0] = tr[id][0];
26
           tr[tcnt][1] = tr[id][1];
27
           //val[tcnt] = val[id];
28
29
           return tent;
30
       }*/
31
       void push_up(int v) {
32
           int l = tr[v][0], r = tr[v][1];
```

```
sz[v] = sz[l] + 1 + sz[r];
33
34
        }
35
       void push_down(int v) {
36
            if(!v) return;
            int l = tr[v][0], r = tr[v][1];
37
38
            //if(l) ;
39
            //if(r) ;
40
            //swap时候注意交换 tr[v][0]和 tr[v][1], 而不是 l和 r;
41
        /* //可持久化在push_down是要新建节点,否者历史版本有可能被多次下传,以rev为例
42
        void push_down(int v) {
43
            if(!v \mid / | !rev[v]) return;
44
            int \ \mathcal{E}l = tr[v][0], \ \mathcal{E}r = tr[v][1];
45
46
            if(l) {
47
                l = copynode(l);
                rev[l]^{=1};
48
                swap(tr[l][0], tr[l][1]);
49
50
            };
            if(r) {
51
52
                r = copynode(r);
53
                rev[r]^=1;
                swap(tr[r][0], tr[r][1]);
54
            };
55
56
            rev[v] = 0;
57
58
        void split(int v,int k,int &x,int &y) {
59
            if(!v) {x=y=0;return;}
60
            push_down(v);
            //v = copynode(v); //可持久化时复制节点
61
62
            /*if(k > sz[tr[v][0]]) {
63
                split(tr[v]/1], k-sz[tr[v]/0]/-1, tr[v]/1], y);
64
            }else{
65
66
                split(tr[v][0], k, x, tr[v][0]);
67
68
            if(val[v] \leftarrow k) {
69
70
                X = V;
                split(tr[v][1], k, tr[v][1], y);
71
72
            }else{
73
74
                split(tr[v][0], k, x, tr[v][0]);
            }
75
            push_up(v);
76
77
        int merge(int x, int y) \{//x堆所有值均小于y堆
78
            if(!x || !y) return x|y;
79
            push_down(x); push_down(y);
80
81
            if(rnd[x]<rnd[y]){</pre>
                 //x = copynode(x); //可持久化时复制节点,可不写
82
                tr[x][1] = merge(tr[x][1],y);
83
84
                push_up(x);
85
                return x;
            }else{
86
87
                 //y = copynode(y); //可持久化时复制节点,可不写
                tr[y][0] = merge(x,tr[y][0]);
88
89
                push_up(y);
                return y;
90
            }
91
92
        void insert(int k) {
93
```

```
int x,y;
94
95
             split(root,k,x,y);
             root = merge(merge(x,newnode(k)),y);
96
97
98
         void recycle(int v) {//回收一颗 treap上所有节点
99
             if(!v) return;
             q.push(v);
100
             recycle(tr[v][0]); recycle(tr[v][1]);
101
102
         void erase(int k) {
103
104
             int x,y,z;
             split(root,k,x,y);
105
106
             split(x,k-1,x,z);
             z = merge(tr[z][0],tr[z][1]);
107
108
             root = merge(x,merge(z,y));
109
         void krank(int k) {
110
             int x,y;
111
             split(root,k-1,x,y);
112
             printf("%d\n",sz[x]+1);
113
             root = merge(x,y);
114
115
         int find(int v,int k) {
116
             if(sz[tr[v][0]]==k-1) return val[v];
117
             if(sz[tr[v][0]]>=k) return find(tr[v][0],k);
118
119
             return find(tr[v][1],k-sz[tr[v][0]]-1);
120
121
         void pre(int k) {
             int x,y;
122
123
             split(root,k-1,x,y);
             printf("%d\n",find(x,sz[x]));
124
125
             root=merge(x,y);
126
         void nxt(int k){
127
             int x,y;
128
             split(root,k,x,y);
129
             printf("%d\n",find(y,1));
130
             root=merge(x,y);
131
132
133
         void reverse(int l,int r){
134
             int x,y,z;
135
             split(root, r, x, y);
136
             split(x, l-1, x, z);
             //rev[z] \stackrel{\sim}{=} 1;标记对本身无效,处理时将z点重新计算
137
             root = merge(merge(x,z),y);
138
139
140
         /*int getRank(int S) {
141
             int res = sz[tr[S][0]]+1;
             for(;fa[S];S = fa[S])
142
143
                  if(tr[fa[S]]/1] == S) res += sz[tr[fa[S]]/0] + 1;
144
            return res;
145
         }*/
         void display(int v) {
146
             if(!v) return;
147
             push_down(v);
148
             display(tr[v][0]);
149
             printf("%d ",val[v]);
150
             display(tr[v][1]);
151
152
         }
153
    }
```

154 //一种可持久化平衡树的替代(非强制在线),由历史版本向当前版本连边,在dfs遍历中利用权值树状数组,普通平衡树等获取答案

2.4 Persistent Data Structures

2.4.1 Chairman Tree

```
struct Node {
        int 1, r;
2
        LL sum;
3
   }t[MAXN * 40];
4
5
   int cnt, n;
   int rt[MAXN];
   void update(int pre, int &x, int l, int r, int v) {
7
       x = ++cnt; t[x] = t[pre]; t[x].sum++;
9
        if(l == r) return;
        int mid = (l + r) >> 1;
10
        if(v <= mid) update(t[pre].1, t[x].1, l, mid, v);</pre>
11
       else update(t[pre].r, t[x].r, mid + 1, r, v);
12
13
   int query(int x, int y, int l, int r, int v) {
14
        if(l == r) return l;
15
16
        int mid = (l + r) >> 1;
17
        int sum = t[t[y].l].sum - t[t[x].l].sum;
18
        if(sum >= v) return query(t[x].1, t[y].1, 1, mid, v);
19
        else return query(t[x].r, t[y].r, mid + 1, r, v - sum);
20
```

2.4.2 Unite Chairman Tree

```
//Q x到 y路 径 第 k大
   //L link(x, y)
2
3 #include <bits/stdc++.h>
4 using namespace std;
5 typedef int LL;
6 const int MAXN = 8e4 + 5;
   const int S = 18;
7
   struct Node {
8
9
       int 1, r;
10
       LL sum;
11 }t[MAXN * 800]; //2 * log^2(n)
12 int n, m, Q;
13 int cnt;
   int rt[MAXN], sz[MAXN];
14
   void update(int pre, int &x, int l, int r, int v) {
15
16
       x = ++cnt; t[x] = t[pre]; t[x].sum++;
17
       if(l == r) return;
       int mid = (l + r) >> 1;
18
       if(v <= mid) update(t[pre].l, t[x].l, l, mid, v);</pre>
19
20
       else update(t[pre].r, t[x].r, mid + 1, r, v);
21
   int query(int x, int y, int z, int w, int l, int r, int v) {
22
       if(l == r) return l;
23
        int mid = (l + r) >> 1;
24
25
        int sum = t[t[x].l].sum + t[t[y].l].sum - t[t[z].l].sum - t[t[w].l].sum;
26
       if(sum >= v) return query(t[x].l, t[y].l, t[z].l, t[w].l, l, mid, v);
27
       return query(t[x].r, t[y].r, t[z].r, t[w].r, mid + 1, r, v - sum);
   }
```

```
int fa[MAXN][S + 3], dep[MAXN];
30
   int val[MAXN];
   vector<int> G[MAXN];
31
32
   inline void addEdge(int x, int y) {
33
        G[x].push_back(y);
        G[y].push_back(x);
34
35
   inline void upd(int x) {
36
        update(rt[fa[x][0]], rt[x], 1, n, val[x]);
for(int i = 1; i <= S; i++) fa[x][i] = fa[fa[x][i - 1]][i - 1];</pre>
37
38
39
   inline void Go(int &x, int step) {
40
        for(int i = S; i >= 0; i--) if(step >> i \& 1) x = fa[x][i];
41
42
43
   int lca(int x, int y) {
        if(dep[x] < dep[y]) swap(x, y);
44
        Go(x, dep[x] - dep[y]);
45
        if(x == y) return x;
46
        for(int i = S; i >= 0; i--) if(fa[x][i] != fa[y][i]) {
47
48
            x = fa[x][i], y = fa[y][i];
49
        return fa[x][0];
50
51
   int get_rt(int x) {
52
        for(int i = S; i \ge 0; i--) if(fa[x][i]) x = fa[x][i];
53
54
        return x;
55
56
   void dfs(int x, int f) {
57
        fa[x][0] = f;
        dep[x] = dep[f] + 1;
58
59
        upd(x);
        for(auto y : G[x]) {
60
             if(y == f) continue;
61
            dfs(y, x);
62
        }
63
   }
64
   void unite(int x, int y) {
65
        int rx = get_rt(x), ry = get_rt(y);
66
        if(sz[rx] > sz[ry]) swap(x, y), swap(rx, ry);
67
68
        addEdge(x, y);
69
        dfs(x, y);
70
        sz[ry] += cnt - rt[x] + 1;
   }
71
   void init() {
72
        cnt = 0;
73
        for(int i = 1; i <= n; i++) {</pre>
74
75
             rt[i] = sz[i] = 0;
            G[i].clear();
76
            dep[i] = 0;
77
        }
78
79
    int main() {
80
        ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
81
        int T; cin >> T;
82
        while(T--) {
83
            cin >> n >> m >> Q;
84
            init();
85
            vector<int> b;
86
87
            map<int, int> mp;
            for(int i = 1; i <= n; i++) {</pre>
88
                 cin >> val[i];
```

```
b.push_back(val[i]);
90
91
             }
92
             sort(b.begin(), b.end());
             b.erase(unique(b.begin(), b.end());
93
94
             for(int i = 1, tmp; i <= n; i++) {</pre>
                  tmp = val[i];
95
                  val[i] = lower_bound(b.begin(), b.end(), val[i]) - b.begin() + 1;
96
97
                  mp[val[i]] = tmp;
98
              for(int i = 1, u, v; i <= m; i++) {
99
                  cin >> u >> v;
100
                  addEdge(u, v);
101
102
103
              for(int i = 1; i <= n; i++) if(!dep[i]) {</pre>
104
                  dep[i] = 1;
                  dfs(i, 0);
105
                  sz[i] = cnt - rt[i] + 1;
106
107
108
             char s[3]; int x, y, z, k, ans = 0;
109
             while(Q--) {
110
                  cin >> s >> x >> y;
                  x \stackrel{\wedge}{=} ans; y \stackrel{\wedge}{=} ans;
111
                  if(s[0] == Q') {
112
                       cin >> k; k ^= ans;
113
                       z = lca(x, y);
114
115
                       ans = query(rt[x], rt[y], rt[z], rt[fa[z][0]], 1, n, k);
116
                       ans = mp[ans];
                       cout << ans << endl;
117
118
                  }
                  else {
119
                       unite(x, y);
120
121
             }
122
123
         return 0;
124
125
```

2.4.3 Persistent Trie

```
//区间异或最值查询
2
   const int N=5e4+10;
3 int t[N];
   int ch[N*32][2],val[N*32];
4
   int cnt;
5
   void init(){
6
7
       mem(ch,0);
8
       mem(val,0);
9
        cnt=1;
   }
10
11
   int add(int root,int x){
12
        int newroot=cnt++,ret=newroot;
        for(int i=30;i>=0;i--){
13
            ch[newroot][0]=ch[root][0];
14
            ch[newroot][1]=ch[root][1];
15
            int now=(x>>i)&1;
16
            root=ch[root][now];
17
18
19
            ch[newroot][now]=cnt++;
            newroot=ch[newroot][now];
20
21
            val[newroot]=val[root]+1;
```

```
22
23
24
        return ret;
25
   int query(int lt,int rt,int x){
26
        int ans=0;
27
28
        for(int i=30;i>=0;i--){
29
             int now=(x>i)&1;
             if(val[ch[rt][now^1]]-val[ch[lt][now^1]]){
30
                 ans l = (1 << i);
31
                 rt=ch[rt][now^1];
32
                 lt=ch[lt][now^1];
33
34
                 } else{
                 rt=ch[rt][now];
35
36
                 lt=ch[lt][now];
37
            }
        }
38
39
        return ans;
40
```

2.4.4 SGT in BBST

```
#include <bits/stdc++.h>
1
using namespace std;
   const int MAXN = 1e5;
3
   const int MAXM = 2e7;
4
   const int LM = 0;
5
   const int RM = 70005;
6
7
    namespace T{
8
        int Tsz; queue<int> q;
        int ls[MAXM], rs[MAXM], val[MAXM];
9
        int Tan, Ta[105], Tbn, Tb[105];
10
        int newnode() {
11
            if(q.empty()) q.push(++Tsz);
12
            int x = q.front(); q.pop();
13
            ls[x] = rs[x] = val[x] = 0;
14
            return x;
15
16
        void insert(int &x, int l, int r, int k, int f) {
17
            if(!x) x = newnode();
18
19
            val[x] += f;
20
            if(l == r) return;
            int mid = (l + r) >> 1;
21
            if(k \ll mid) {
22
                insert(ls[x], l, mid, k,f);
23
24
            }else{
25
                insert(rs[x], mid+1, r, k,f);
26
            }
27
28
        int query(int 1, int r, int y) {
29
            if(l == r) return l;
            int mid = (l + r) \gg 1;
30
            int sum = 0;
31
            for(int i = 0; i <= Tan; i++) sum -= val[ls[Ta[i]]];</pre>
32
            for(int i = 0; i <= Tbn; i++) sum += val[ls[Tb[i]]];</pre>
33
34
            if(y \le sum) {
35
36
                 for(int i = 0; i <= Tan; i++) Ta[i] = ls[Ta[i]];</pre>
37
                 for(int i = 0; i <= Tbn; i++) Tb[i] = ls[Tb[i]];</pre>
38
                 return query(l, mid, y);
```

```
39
            for(int i = 0; i <= Tan; i++) Ta[i] = rs[Ta[i]];</pre>
40
            for(int i = 0; i <= Tbn; i++) Tb[i] = rs[Tb[i]];</pre>
41
42
            return query(mid+1, r, y - sum);
43
        void recycle(int v) {
44
            if(!v) return;
45
46
            q.push(v);
            recycle(ls[v]);
47
48
            recycle(rs[v]);
        }
49
   };
50
    namespace TT{
51
        int Tsn;
52
        int val[MAXN], Trt[MAXN];
53
54
        int sz[MAXN];
        int fa[MAXN], tr[MAXN][2];
55
56
57
        int root;
58
        double alp;
59
        int newnode(int x) {
60
            int tcnt = ++Tsn;
61
            val[tcnt] = x; Trt[tcnt] = 0;
62
            fa[tcnt] = tr[tcnt][0] = tr[tcnt][1] = 0;
63
64
            sz[tcnt] = 0;
65
            return tcnt;
66
67
        bool isBad(int v) {
            return (double(sz[ tr[v][0] ]) > double(sz[v]) * alp) ||
68
                 (double(sz[ tr[v][1] ]) > double(sz[v]) * alp);
69
70
        vector<int> vec;
71
        void rRecycle(int v) {
72
            if(tr[v][0]) rRecycle(tr[v][0]);
73
            vec.push_back(v);
74
            T::recycle(Trt[v]);
75
            if(tr[v][1]) rRecycle(tr[v][1]);
76
77
78
        int rBuild(int 1, int r) {
79
            int mid = (l + r) \gg 1, v = vec[mid];
80
            Trt[v] = 0; sz[v] = r - l + 1;
            for(int i = 1; i <= mid; i++) {</pre>
81
82
                 T::insert(Trt[v], LM, RM, val[vec[i]], 1);
83
            tr[v][0] = (l \ll mid-1) ? rBuild(l, mid - 1) : 0;
84
            if(tr[v][0]) fa[tr[v][0]] = v;
85
            tr[v][1] = (mid+1 \le r) ? rBuild(mid + 1, r) : 0;
86
            if(tr[v][1]) fa[tr[v][1]] = v;
87
            return v;
88
89
        void rebuild(int v) {
90
91
            if(isBad(v)) {
                 vec.clear()
92
93
                 int u = fa[v], lr = tr[u][1] == v;
94
                 rRecycle(v)
                 if(vec.size()) v = rBuild(0, vec.size() - 1); else v = 0;
95
                if(u == 0) fa[root = v] = 0;
96
                 else{
97
                     tr[u][lr] = v;
98
99
                     if(v) fa[v] = u;
```

```
100
             }
101
102
103
         int find(int x, int k) {
             int l = tr[x][0], r = tr[x][1];
104
             if(sz[l] + 1 == k) return x;
105
             if(sz[l] >= k) return find(l, k);
106
107
             return find(r, k - sz[l] - 1);
108
         void ins(int x, int y) {
109
             int v = find(root, x);
110
             int p = tr[v][1], q = v;
111
             if(p) {
112
                 for(;p; q = p, p = tr[p][0]);
113
114
                 fa[p = tr[q][0] = newnode(y)] = q;
115
             }else {
                 fa[p = tr[q][1] = newnode(y)] = q;
116
117
             int fg = 0;
118
             T::insert(Trt[p], LM, RM, y, 1);
119
120
             sz[p] = 1;
             for(;fa[p]; p = fa[p]) {
121
                 if(tr[fa[p]][0] == p)
122
                     T::insert(Trt[fa[p]], LM, RM, y, 1);
123
                 sz[fa[p]]++;
124
125
                 if(isBad(fa[p])) fg = fa[p];
126
127
             rebuild(fg);
128
         void upd(int x, int y) {
129
             int p = find(root, x+1);
130
             int ty = val[p]; val[p] = y;
131
             T::insert(Trt[p], LM, RM, ty, -1);
132
             T::insert(Trt[p], LM, RM, y, 1);
133
             for(;fa[p]; p = fa[p]) {
134
                 if(tr[fa[p]][0] == p) {
135
                     T::insert(Trt[fa[p]], LM, RM, ty, -1);
136
                     T::insert(Trt[fa[p]], LM, RM, y, 1);
137
                 }
138
             }
139
140
141
         int que(int x, int y,int z) {
             x = find(root, x);
142
             T::Ta[T::Tan = 0] = Trt[x];
143
             for(;fa[x]; x = fa[x])
144
                 if(tr[fa[x]][1] == x) {
145
                     T::Ta[++T::Tan] = Trt[fa[x]];
146
147
             y = find(root, y+1);
148
             T::Tb[T::Tbn = 0] = Trt[y];
149
             for(;fa[y]; y = fa[y]) {
150
                 if(tr[fa[y]][1] == y)
151
                     T::Tb[++T::Tbn] = Trt[fa[y]];
152
153
154
             return T::query(LM, RM, z);
155
156
         void init(int n) {
157
             alp = 0.7;
158
             vec.clear();
159
             vec.push_back(newnode(RM));
160
```

```
for(int i = 1, a; i <= n; i++) {</pre>
161
                  scanf("%d", &a);
162
                  vec.push_back(newnode(a));
163
164
             root = rBuild(0,vec.size() - 1);
165
             fa[root] = 0;
166
167
168
         void display(int v) {
             if(tr[v][0]) display(tr[v][0]);
169
             cerr<<val[v]<<" ";
170
             if(tr[v][1]) display(tr[v][1]);
171
172
    };
173
```

2.5 Tree Structures

2.5.1 dsu on tree

```
const int MAXN = 1e5 + 7;
1
   vector<int> G[MAXN];
2
   int bgison, dfs_clock, sz[MAXN], st[MAXN], bt[MAXN], et[MAXN];
3
   int fg[MAXN], col[MAXN];
4
   long long ans[MAXN];
5
6
   void dfs1(int u, int fa) {
        sz[u] = 1;
7
8
        st[bt[u] = ++dfs\_clock] = u;
9
        for(auto v : G[u])
        if(v != fa) {
10
            dfs1(v, u);
11
            sz[u] += sz[v];
12
13
        et[u] = dfs_clock;
14
   }
15
   int maxx = 0;
16
    void dfs2(int u, int fa, int keep) {
17
        int mx = -1, bigson = -1;
18
19
        for(auto &v : G[u])
20
        if(v != fa) {
21
            if(sz[v] > mx)
22
                 mx = sz[v], bigson = v;
23
        for(auto &v : G[u])
24
25
        if(v != fa && v != bigson)
26
             dfs2(v,u,0);
        if(bigson != -1) {
27
            dfs2(bigson, u, 1);
28
29
            ans[u] = ans[bigson];
30
             for(int &v : G[u])
                 if(v != fa && v != bigson)
31
                 for(int i = bt[v]; i <= et[v]; i++) {</pre>
32
                      ++fg[col[st[i]]]
33
                     if(fg[ col[st[i]] ] > maxx) maxx=fg[col[st[i]]], ans[u] = 0;
if(fg[ col[st[i]] ] == maxx) ans[u] += col[st[i]];
34
35
36
37
        ++fq[col[u]];
38
        if(fg[col[u]] > maxx) maxx = fg[col[u]], ans[u] = 0;
39
        if(fg[col[u]] == maxx) ans[u] += col[u];
40
41
        if(keep == 0) {
```

2.5.2 Vitural Tree

```
const int MAXN = ;
   const int MAXM = ;
   const LL INF =
3
   const int S = 19;
4
   int ecnt, head[MAXN];
5
   struct Edge{int to, nxt; LL w;} e[MAXM], ve[MAXM];
   inline void addEdgeT(int x, int y, LL w) {
7
8
        e[++ecnt] = (Edge) \{y, head[x], w\}; head[x] = ecnt;
9
   }
   int dep[MAXN], dfn_time, dfn[MAXN], fa[MAXN][S+1];
10
   LL dis[MAXN][S+1];
12
   void dfs(int v,int _fa) {
13
        dfn[v] = ++dfn_time;
14
        dep[v] = dep[_fa] + 1;
15
        fa[v][0] = _fa;
        for(int i = 1; i <= S; i++) {
16
            fa[v][i] = fa[fa[v][i-1]][i-1];
17
18
            dis[v][i] = min(dis[v][i-1], dis[fa[v][i-1]][i-1]);
19
        for(int i = head[v], u; i; i = e[i].nxt)
20
21
        if((u = e[i].to) != _fa) {
22
            dis[u][0] = e[i].w;
23
            dfs(u, v);
        }
24
25
   int getLca(int u, int v) {
26
        if(dep[u] < dep[v]) swap(u, v);</pre>
27
        for(int i = S; i >= 0; i--)
28
            if(dep[fa[u][i]] >= dep[v]) u = fa[u][i];
29
        if(u == v) return u;
30
        for(int i = S; i >= 0; i--)
31
            if(fa[u][i] != fa[v][i])
32
33
                u = fa[u][i], v = fa[v][i];
        return fa[u][0];
34
35
   LL getDis(int u, int v) {
36
        if(dep[u] < dep[v]) swap(u, v);</pre>
37
38
        LL res = INF;
        for(int i = S; i >= 0; i--)
39
40
            if(dep[fa[u][i]] >= dep[v]) {
                res = min(res, dis[u][i]);
41
                u = fa[u][i];
42
43
        return res;
44
45
   namespace VituralTree{
46
        int hn, h[MAXN];
47
        int vecnt, vhead[MAXN];
48
        Edge ve[MAXM];
49
50
        int top, st[MAXN], cln, cl[MAXN];
51
        int fgn, fg[MAXN];//利用fg[i]==fgn判断是否为当前有效点
        void addEdgeVT(int x, int y, LL w) {
```

```
ve[++vecnt] = (Edge) \{y, vhead[x], w\}; vhead[x] = vecnt;
53
54
55
        inline void link(int u, int v) {
56
            LL w = getDis(u, v);
            addEdgeVT(u, v, w);
57
            addEdgeVT(v, u, w);
58
59
60
        inline bool cmp(int a, int b) {return dfn[a] < dfn[b];}</pre>
61
        void build() {
62
            ++fgn;
            for(int i = 1; i <= hn; i++) fg[h[i]] = fgn;</pre>
63
            sort(h + 1, h + hn + 1, cmp);
64
            cl[cln = 1] = st[top = 1] = 1;
65
            for(int i = 1; i <= hn; i++) {
66
                int rem = getLca(st[top], h[i]);
67
                if(rem == st[top]) {
68
                    if(rem != st[top]) cl[++cln] = st[++top] = h[i];
69
70
71
                while(top > 1 && dep[st[top - 1]] >= dep[rem]) {
72
73
                    link(st[top - 1], st[top]); top--;
74
                if(dep[st[top]] > dep[rem]) {
75
76
                    link(rem, st[top]), top --;
77
78
                if(rem != st[top]) cl[++cln] = st[++top] = rem;
79
                if(h[i] != st[top]) cl[++cln] = st[++ top] = h[i];
80
81
            while(top > 1) {
                link(st[top - 1], st[top]), top --;
82
83
84
        void clear() {
85
            vecnt = 0;
86
            for(;cln; --cln) vhead[cl[cln]] = 0;
87
88
        void sol() {
89
90
            build();
91
            //注意一号节点可能在虚树外面
92
            clear();
93
        }
94
   }
```

2.5.3 Tree Decomposition

```
int sz[MAXN], dep[MAXN], top[MAXN], fa[MAXN], son[MAXN], num[MAXN], totw;
1
   struct Edge {
2
3
       int to, nxt;
   }e[MAXN << 1];
4
   int head[MAXN], ecnt;
5
   int n, m, Q;
   #define Ls(x) (x << 1)
7
   #define Rs(x) (x << 1 | 1)
   struct Tree {
9
        int l, r, lazy;
10
       LL sum, mx;
11
   }tree[MAXN << 2]</pre>
12
int A[MAXN], B[MAXN];
14 void push_up(int x) {
       tree[x].sum = tree[Ls(x)].sum + tree[Rs(x)].sum;
```

```
tree[x].mx = max(tree[Ls(x)].mx, tree[Rs(x)].mx);
16
17
   void push_down(int x) {
18
        if(tree[x].lazy) {
19
            tree[Ls(x)].sum += tree[x].lazy * (tree[Ls(x)].r - tree[Ls(x)].l + 1);
20
            tree[Rs(x)].sum += tree[x].lazy * (tree[Rs(x)].r - tree[Rs(x)].l + 1);
21
            tree[Ls(x)].mx += tree[x].lazy;
22
23
            tree[Rs(x)].mx += tree[x].lazy;
            tree[Ls(x)].lazy += tree[x].lazy;
24
            tree[Rs(x)].lazy += tree[x].lazy;
25
26
            tree[x].lazy = 0;
        }
27
28
29
   void build(int x, int L, int R) {
30
        tree[x].lazy = 0;
31
        tree[x].l = L; tree[x].r = R;
        if(L == R) {
32
            tree[x].sum = B[L];
33
            tree[x].mx = B[L];
34
35
36
        int mid = (L + R) \gg 1;
37
        build(Ls(x), L, mid);
38
39
        build(Rs(x), mid + 1, R);
40
        push_up(x);
41
42
   void update(int x, int L, int R, LL val) {
        if(tree[x].l >= L && tree[x].r <= R) {</pre>
43
44
            tree[x].lazy += val;
            tree[x].sum += val * (tree[x].r - tree[x].l + 1);
45
46
            tree[x].mx += val;
47
            return;
        }
48
        push_down(x);
49
        int mid = (tree[x].l + tree[x].r) >> 1;
50
        if(L <= mid) update(Ls(x), L, R, val);</pre>
51
52
        if(R > mid) update(Rs(x), L, R, val);
        push_up(x);
53
54
55
   LL query(int x, int L, int R) {
56
        if(tree[x].l >= L && tree[x].r <= R)
57
            return tree[x].sum;
58
        push_down(x);
59
        int mid = (tree[x].l + tree[x].r) >> 1;
60
        LL res = 0;
61
        if(L \leftarrow mid) res += query(Ls(x), L, R);
62
        if(R > mid) res += query(Rs(x), L, R);
        return res;
63
64
   LL query2(int x, int L, int R) {
   if(tree[x].l >= L && tree[x].r <= R)</pre>
65
66
67
            return tree[x].mx;
        push_down(x);
68
        int mid = (tree[x].l + tree[x].r) >> 1;
69
        LL res = -INF;
70
        if(L \le mid) res = max(res, query2(Ls(x), L, R));
71
72
        if(R > mid) res = max(res, query2(Rs(x), L, R));
73
        return res;
74
   inline void add_edge(int x, int y) {
75
        e[++ecnt] = (Edge) \{y, head[x]\}; head[x] = ecnt;
```

```
77
    void dfs1(int x) {
78
         sz[x] = 1; son[x] = 0;
79
80
         for(int i = head[x]; i; i = e[i].nxt) {
81
             int v = e[i].to;
             if(v == fa[x]) continue;
82
             fa[v] = x;
83
             dep[v] = dep[x] + 1;
84
85
             dfs1(v);
86
             sz[x] += sz[v];
             if(sz[v] > sz[son[x]]) son[x] = v;
87
        }
88
89
    void dfs2(int x) {
90
91
         B[num[x]] = A[x];
         if(son[x]) {
92
             top[son[x]] = top[x];
93
94
             num[son[x]] = ++totw;
95
             dfs2(son[x]);
96
         for(int i = head[x]; i; i = e[i].nxt) {
97
             int v = e[i].to;
98
             if(v == fa[x] | | v == son[x]) continue;
99
             top[v] = v;
100
101
             num[v] = ++totw;
102
             dfs2(v);
103
         }
104
105
    void up(int a, int b, int c) {
         int f1 = top[a], f2 = top[b];
106
        while(f1 != f2) {
107
             if(dep[f1] < dep[f2]) { swap(a, b); swap(f1, f2); }</pre>
108
             update(1, num[f1], num[a], c);
109
             a = fa[f1];
110
             f1 = top[a];
111
112
         if(dep[a] > dep[b]) swap(a, b);
113
        update(1, num[a], num[b], c);
114
115
116
    int qsum(int a, int b) {
117
         if(a == b) return query(1, num[a], num[a]);
118
         int f1 = top[a], f2 = top[b];
         int res = 0;
119
        while(f1 != f2) {
120
             if(dep[f1] < dep[f2]) { swap(a, b); swap(f1, f2); }</pre>
121
122
             res += query(1, num[f1], num[a]);
             a = fa[f1];
123
             f1 = top[a];
124
125
126
         if(dep[a] > dep[b]) swap(a, b);
127
         res += query(1, num[a], num[b]);
         return res;
128
129
    int qmax(int a, int b) {
130
         if(a == b) return query2(1, num[a], num[a]);
131
         int f1 = top[a], f2 = top[b];
132
         int res = -10000000000;
133
        while(f1 != f2) {
134
             if(dep[f1] < dep[f2]) { swap(a, b); swap(f1, f2); }
135
             res = max(res, query2(1, num[f1], num[a]));
136
137
             a = fa[f1];
```

```
138
            f1 = top[a];
139
        if(dep[a] > dep[b]) swap(a, b);
140
141
        res = max(res, query2(1, num[a], num[b]));
142
        return res;
143
    inline void init() {
144
        memset(head, 0, sizeof(head)); ecnt = 0;
145
         fa[1] = 0; dep[1] = 1; top[1] = 1; num[1] = 1; totw = 1;
146
147
    inline void pre() {
148
        dfs1(1); dfs2(1); build(1, 1, totw);
149
150
```

2.5.4 Link-Cut Tree

```
1
    namespace LCT {
2
        int fa[MAXN], rev[MAXN], tr[MAXN][2];
3
        int s[MAXN], val[MAXN];
4
        void push_up(int x) {
5
            int l = tr[x][0], r = tr[x][1];
6
            s[x] = s[1] + s[r] + val[x];
 7
8
        void Rev(int x) {
9
            rev[x] = 1; swap(tr[x][0], tr[x][1]);
10
        void push_down(int x) {
11
            if(!rev[x]) return;
12
13
            int l = tr[x][0], r = tr[x][1];
            rev[x] = 0;
14
            if(l) Rev(l); if(r) Rev(r);
15
16
        bool isroot(int x) {
17
            return tr[fa[x]][0] != x && tr[fa[x]][1] != x;
18
19
        void pre(int x) {
20
            if(!isroot(x)) pre(fa[x]);
21
            push_down(x);
22
23
24
        void rotate(int x) {
25
            int y = fa[x]; int z = fa[y];
            int l = tr[y][1] == x;
26
            int r = 1 \wedge 1;
27
            if(!isroot(y)) tr[z][tr[z][1] == y] = x;
28
            fa[x] = z; fa[y] = x; fa[tr[x][r]] = y;
29
30
            tr[y][l] = tr[x][r]; tr[x][r] = y;
31
            push_up(y);
32
        void splay(int x) {
33
            pre(x);
34
35
            int y, z;
            while(!isroot(x)) {
36
                y = fa[x]; z = fa[y];
37
                if(!isroot(y)) {
38
                     if((tr[z][0] == y) == (tr[y][0] == x))rotate(y);
39
                     else rotate(x);
40
41
42
                rotate(x);
43
44
            push_up(x);
```

```
45
        void access(int x) {
46
47
            int y = 0;
            while(x) {
48
49
                splay(x); tr[x][1] = y;
                push_up(x);
50
51
                y = x; x = fa[x];
            }
52
53
        void makeroot(int x) {
54
            access(x); splay(x); Rev(x);
55
56
        void lnk(int x, int y) {
57
            makeroot(x); fa[x] = y;
58
59
        void cut(int x, int y) {
60
            makeroot(x); access(y); splay(y);
61
62
            tr[y][0] = fa[x] = 0; push_up(y);
63
        void update(int x, int y) {
64
            makeroot(x); val[x] = y; push_up(x);
65
66
        int query(int x, int y) {
67
            makeroot(x); access(y); splay(y);
68
            return s[y];
69
70
        bool check(int x, int y) {
71
            int tmp = y;
makeroot(x); access(y); splay(x);
72
73
            while(!isroot(y)) y = fa[y];
74
            splay(tmp);
75
76
            return x == y;
        }
77
   }
78
```

3 String

3.1 Basics

3.1.1 Hash

```
const LL p1 = 201, p2 = 301, mod1 = 12000000319, mod2 = 2147483647;
2
   struct Hash {
3
        LL a, b;
        void append(Hash pre, int v) {
4
5
            a = (pre.a * p1 + v) \% mod1;
            b = (pre.b * p2 + v) \% mod2;
6
7
        void init(string S) {
8
9
            a = b = 0;
            for(int i = 0; i < S.size(); i++) append(*this, S[i]);</pre>
10
11
        bool operator == (const Hash &x) const {
12
13
            return a == x.a \&\& b == x.b;
14
15
        bool operator < (const Hash &x) const {</pre>
16
            return a < x.a | | (a == x.a \&\& b < x.b);
17
        }
18
   };
```

3.1.2 Minimum String

```
namespace minstring{
1
        int getmin(char *s, int sn) {
2
            int i = 0, j = 1, k = 0, t;
3
            while(i < sn && j < sn && k < sn) \{
4
5
                t = s[(i + k) % sn] - s[(j + k) % sn];
6
                if(!t) k++;
7
                 else {
8
                     if(t > 0) i += k + 1; else j += k + 1;
9
                     if(i == j) j++;
                     k = 0;
10
11
12
            }
13
            return i < j ? i : j;</pre>
14
        }
15
   }
```

3.2 String Matching

3.2.1 Bitset Match

```
namespace BitMatch{
1
2
       const int S = 26;
       bitset<MAXN> bs[S], ret;
3
4
       char s[MAXN];
5
       inline int idx(char c) { return c - 'a'; }
6
       inline void init() {
           for(int i = 0; i < 26; i++) bs[i].reset();</pre>
7
           scanf("%s", s);
8
9
           int sn = strlen(s);
```

```
for(int i = 0; i < sn; i++) bs[idx(s[i])].set(i);
10
11
12
       void modify(int p, char ch) {
13
            bs[idx(s[p])].reset(p);
            s[p] = ch;
14
            bs[idx(s[p])].set(p);
15
16
17
        int match(char *t, int tn) {//返回 t 串在 s 串中出现的次数
            ret = bs[idx(t[0])];
18
            for(int i = 1; i < tn; i++) {
19
                ret <<= 1;
20
                ret &= bs[idx(t[i])];
21
22
23
            return ret.count();
24
        }
25
   }
```

3.2.2 KMP && exKMP

```
1
   namespace KMP {
2
        int fa[MAXN];
        void get_fail(char* t, int tn) {
3
4
            fa[0] = -1;
            int i = 0, j = -1;
5
6
            while(i < tn) {</pre>
                 if (j == -1 || t[i] == t[j]) {
7
                     ++i; ++j;
8
9
                     fa[i] = t[i] != t[j] ? j : fa[j];
10
                 }else{
11
                     j = fa[j];
12
            }
13
14
        void kmp(char* s, int sn, char* t, int tn) {
15
            int i = 0, j = 0;
16
            while(i < sn) {</pre>
17
                 if (j == -1 || s[i] == t[j]) {
18
                     i++;j++;
19
20
                     if(j == tn) {
21
                     }
22
                 }else j = fa[j];
            }
23
        }
24
25
   }
   namespace exKMP {
26
27
        int nxt[MAXN], ext[MAXN];
28
        void get_nxt(char* t, int tn) {
29
            int j = 0, mx = 0;
            nxt[0] = tn;
30
            for(int i = 1; i < tn; i++) {</pre>
31
                 if(i \ge mx \mid | i + nxt[i - j] \ge mx) {
32
33
                     if(i > mx) mx = i;
                     while(mx < tn && t[mx] == t[mx - i]) mx++;
34
                     nxt[i] = mx - i;
35
36
                     j = i;
                 }else nxt[i] = nxt[i - j];
37
            }
38
39
40
        void exkmp(char *s, int sn, char *t, int tn) {
41
            int j = 0, mx = 0;
```

```
for(int i = 0; i < sn; i++) {</pre>
42
                 if(i >= mx || i + nxt[i - j] >= mx) {
43
44
                     if(i > mx) mx = i;
                     while(mx < sn && mx - i < tn && s[mx] == t[mx - i]) mx++;
45
46
                     ext[i] = mx - i;
47
                     j = i;
                 }else ext[i] = nxt[i - j];
48
49
            }
        }
50
   }
51
```

3.2.3 AC Automaton

```
namespace AC {
1
        int ch[MAXN][sigma_size], last[MAXN];
2
3
        int val[MAXN], f[MAXN], sz;
4
        inline void init() { sz = 1; memset(ch[0], 0, sizeof(ch[0])); }
        inline int idx(char c) { return c - 'a'; }
5
6
        void insert(string s, int v) {
7
            int u = 0;
8
            for(int i = 0; i < s.size(); i++) {</pre>
9
                 int c = idx(s[i]);
10
                 if(!ch[u][c]) {
                     memset(ch[sz], 0, sizeof(ch[sz]));
11
                     val[sz] = 0;
12
13
                     ch[u][c] = sz++;
14
                 u = ch[u][c];
15
16
            val[u] = v;
17
18
        void get_fail() {
19
            queue<int> q;
20
            f[0] = 0;
21
            for(int c = 0; c < sigma_size; c++) {</pre>
22
23
                 int u = ch[0][c];
                 if(u) { f[u] = 0; q.push(u); last[u] = 0; }
24
25
            while(!q.empty()) {
26
27
                 int r = q.front(); q.pop();
28
                 for(int c = 0; c < sigma_size; c++) {</pre>
                     int u = ch[r][c];
29
                     if(!u) { ch[r][c] = ch[f[r]][c]; continue; }
30
                     q.push(u);
31
                     int v = f[r];
32
                     while(v \&\& !ch[v][c]) v = f[v];
33
34
                     f[u] = ch[v][c];
                     last[u] = val[f[u]] ? f[u] : last[f[u]];
35
36
                 }
            }
37
38
39
        inline void solve(int j) {
40
            if(j) {
                 ans += val[j];
41
                 solve(last[j]);
42
            }
43
        }
44
45
        void find(string T) {
46
            int j = 0;
            for(int i = 0; i < T.size(); i++) {</pre>
47
```

```
int c = idx(T[i]);
48
49
                  j = ch[j][c];
50
                  if(val[j]) solve(j);
51
                  else if(last[j]) solve(last[j]);
52
             }
53
         }
54
55
     namespace AC {
         int root, tcnt;
56
         int ch[MAXN][sigma_size], fa[MAXN];
57
         inline int newnode() {
58
             fa[++tcnt] = 0;
59
60
              for(int i = 0; i < sigma_size; ++i) ch[tcnt][i] = 0;</pre>
61
              return tcnt;
62
         inline void init() {
63
             tcnt = -1;
64
65
              root = newnode();
66
         inline int idx(char c) { return c - 'a'; }
67
         void extend(char *s, int sn) {
68
             int cur = root;
69
              for(int i = 0, c; i < sn; i++) {
70
                  if(!ch[cur][c = idx(s[i])])
71
72
                       ch[cur][c] = newnode();
73
                  cur = ch[cur][c];
             }
74
75
76
         int q[MAXN], qh, qt;
77
         void get_fail() {
             qh = 1; qt = 0;
78
              fa[root] = 0;
79
              for(int c = 0, now; c < sigma_size; c++)</pre>
80
                  if((now = ch[root][c]) != 0)
81
                       q[++qt] = now;
82
             while(qh <= qt) {</pre>
83
                  int cur = q[qh++];
84
                  for(int c = 0, now; c < sigma_size; c++)</pre>
85
86
                       if((now = ch[cur][c]) != 0) {
87
                           fa[now] = ch[fa[cur]][c];
88
                           q[++qt] = now;
89
                       }else
                           ch[cur][c] = ch[fa[cur]][c];
90
91
             }
92
93
     //统计模板串出现次数,每个模板串只计算一次
94
              int\ cur\ =\ root\,,\ ans\ =\ 0\,;
95
              for(int \ i = 0; \ i < sn; ++i)
                  cur \; = \; ch \, [\, cur \, ] \, [\, idx \, (\, s \, [\, i \, ]) \, ] \, ;
96
97
                  for(int j = cur; j \& cnt[j] != -1; j = fa[j]) {
98
                       ans \neq cnt[j];
                       cnt/j/=-1;
99
100
              }
101
102
    }
103
```

3.3 Suffix Related

3.3.1 Suffix Array

```
1
    namespace SA {
2
        char s[MAXN];
        int sa[MAXN], rank[MAXN], height[MAXN];
3
4
        int t[MAXN], t2[MAXN], c[MAXN], n;
5
        void clear() { n = 0; memset(sa, 0, sizeof(sa)); }
        void build(int m) {
6
            int *x = t, *y = t2;
7
            for(int i = 0; i < m; i++) c[i] = 0;
8
9
             for(int i = 0; i < n; i++) c[x[i] = s[i]]++;
             for(int i = 1; i < m; i++) c[i] += c[i - 1];
10
             for(int i = n - 1; i \ge 0; i--) sa[--c[x[i]]] = i;
11
             for(int k = 1; k <= n; k <<= 1) {
12
13
                 int p = 0;
14
                 for(int i = n - k; i < n; i++) y[p++] = i;
                 for(int i = 0; i < n; i++) if(sa[i] >= k) y[p++] = sa[i] - k;
for(int i = 0; i < m; i++) c[i] = 0;</pre>
15
16
                 for(int i = 0; i < n; i++) c[x[y[i]]]++;
17
                 for(int i = 1; i < m; i++) c[i] += c[i - 1];
18
                 for(int i = n - 1; i \ge 0; i--) sa[--c[x[y[i]]]] = y[i];
19
20
                 swap(x, y);
                 p = 1; x[sa[0]] = 0;
21
                 for(int i = 1; i < n; i++)</pre>
22
                     x[sa[i]] = y[sa[i - 1]] == y[sa[i]] & y[sa[i - 1] + k] == y[sa[i] + k]
23
        ? p - 1 : p++;
24
                 if(p >= n) break;
25
                 m = p;
26
            }
27
        void buildHeight() {
28
29
            int k = 0;
30
            for(int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
31
            for(int i = 0; i < n; i++) {</pre>
32
                 if(k) k--;
                 int j = sa[rank[i] - 1];
33
34
                 while(s[i + k] == s[j + k]) k++;
35
                 height[rank[i]] = k;
            }
36
37
        void init() {
38
39
            n = strlen(s) + 1;
            build(z' + 1);
40
            buildHeight();
41
        }
42
43
   }
```

3.3.2 Suffix Automaton

```
namespace SAM{
   int scnt, root, last;
   int fa[MAXN<<1], len[MAXN<<1][26];
   int sc[MAXN<<1], tmpl[MAXN<<1];

int newnode(int _len, int q = 0) {
   fa[++scnt] = fa[q]; len[scnt] = _len;
   sc[scnt] = 0; tmpl[scnt] = INF;</pre>
```

```
for(int i = 0; i < 26; i++) ch[scnt][i] = ch[q][i];</pre>
9
10
           return scnt;
11
12
       void init() {
13
           scnt = 0;
           root = last = newnode(0);
14
15
       void extend(int c) {
16
           int p = last, np = newnode(len[p] + 1);
17
           for(;p \&\& ch[p][c] == 0; p = fa[p]) ch[p][c] = np;
18
19
           if(!p) fa[np] = root;
           else{
20
21
               int q = ch[p][c];
               if(len[p] + 1 == len[q]) fa[np] = q;
22
23
               else{
24
                    int nq = newnode(len[p] + 1, q);
                   fa[np] = fa[q] = nq;
25
26
                   for(; p && ch[p][c] == q; p = fa[p]) ch[p][c] = nq;
27
28
           last = np;
29
30
       int c[MAXN], rs[MAXN << 1];</pre>
31
       void radix_sort(int n){
32
33
           for(int i = 0; i \le n; i++) c[i] = 0;
34
           for(int i = 1; i <= scnt; i++) c[len[i]]++;</pre>
35
           for(int i = 1; i <= n; i++) c[i] += c[i-1];
36
           for(int i = scnt; i >= 1; i--) rs[c[len[i]]--] = i;
37
       void go(){
38
           scanf("%s",s);
39
           int n = strlen(s);
40
           for(int i = 0; i < n; ++i)
41
               extend(s[i] - 'a');
42
           radix_sort(n);
43
            //以下sc集合意义不同
44
           {//每个节点对应的位置之后有多少个不同子串
45
                for(int i = scnt; i >= 1; i--) {
46
47
                    int S = 0;
48
                    for(int j = 0; j < 26; j++)
49
                        S += sc[ ch[rs[i]][j] ];
50
                    sc[rs[i]] = S + 1;
               }
51
52
           {//right集合大小
53
               int cur = root;
54
                for(int i = 0; i < n; ++i) {</pre>
55
                   cur = ch[cur][s[i] - 'a'];
56
57
                   sc[cur]++;
58
               for(int i = scnt; i >= 1; --i) {
59
                    sc[ fa[rs[i]] ] += sc[rs[i]];
60
61
           }
62
           //公共子串
63
           //tmpl,当前字符串:在状态cur,与模板串的最长公共后缀
64
           //minl, 多个字符串:在状态cur,与模板串的最长公共后缀
65
            //注意:在状态cur匹配成功时,cur的祖先状态与字符串的最长公共后缀
66
           for(; ~scanf("%s",s);) {
67
               int cur = root, Blen = 0;
68
                for(int i = 0; i <= scnt; i++)
69
```

```
70
                      tmpl[i] = 0;
                 n = strlen(s);
71
                 for(int i = 0, x; i < n; i++) {
72
                      X = S[i] - a';
73
                      if(ch[cur][x]) {
74
                          ++Blen;
75
                          cur = ch[cur][x];
76
77
                      }else{
                          for(; cur \&\& ch[cur][x] == 0; cur = fa[cur]);
78
                          if(cur) {
79
                              Blen = len[cur] + 1;
80
                              cur = ch[cur][x];
81
                          }else{
82
83
                              cur = root; Blen = 0;
84
                          }
85
                      tmpl[cur] = max(tmpl[cur], Blen);
86
87
                 for(int i = scnt; i ; --i) {
88
                      if( tmpl[ fa[rs[i]] ] < tmpl[ rs[i] ])</pre>
89
                          tmpl[ fa[rs[i]] ] = len[ fa[rs[i]] ];
90
                      minl[ rs[i] ] = min(minl[ rs[i] ], tmpl[ rs[i] ]);
91
                 }
92
             }
93
94
95
96
    namespace exSAM{
97
         int scnt, root;
         int fa[MAXN<<1], len[MAXN<<1], ch[MAXN<<1][26];</pre>
98
         int sc[MAXN<<1], tmpl[MAXN<<1], minl[MAXN<<1];</pre>
99
100
         int newnode(int _len, int q = 0) {
101
             fa[++scnt] = fa[q]; len[scnt] = _len;
102
             sc[scnt] = 0;tmpl[scnt] = 0; minl[scnt] = INF;
103
             for(int i = 0; i < 26; i++) ch[scnt][i] = ch[q][i];
104
105
             return scnt;
106
         }
         void init() {
107
             scnt = 0;
108
109
             root = newnode(0);
110
         int work(int p,int c){
111
112
             int q = ch[p][c];
             int nq = newnode(len[p] + 1, q);
113
             fa[q] = nq;
114
             for(; p && ch[p][c] == q; p = fa[p]) ch[p][c] = nq;
115
116
             return nq;
117
         int extend(int p, int c) {
118
             if (ch[p][c]){
119
120
                  int q = ch[p][c];
                 if (len[p] + 1 == len[q]) return q;
121
                 return work(p, c);
122
123
124
             int np = newnode(len[p] + 1);
             for(;p \&\& ch[p][c] == 0; p = fa[p]) ch[p][c] = np;
125
             if (!p) fa[np] = root;
126
             else{
127
                 int q = ch[p][c];
128
                 if (len[p] + 1 == len[q]) fa[np] = q;
129
                 else fa[np] = work(p, c);
130
```

```
131
              return np;
132
133
134
          void solve() {
              int n; scanf("%d",&n);
135
              for(int i = 1; i <= n; i++) {
    scanf("%s", s);</pre>
136
137
138
                    int sn = strlen(s);
                    int last = root;
139
                    for(int j = 0; j < sn; ++j)
140
                        last = extend(last, s[j] - 'a');
141
              }
142
          }
143
144
     }
```

3.4 Palindrome Related

3.4.1 Manacher

```
namespace Manachar {
 1
 2
          char S[MAXN << 1];</pre>
3
          int scnt, ans;
          int p[MAXN << 1]; //p[i] - 1
void init(char *s0, int sn0) {
   S[0] = '$'; S[1] = '#';</pre>
 4
 5
 6
               for(int i = 0; i < sn0; i++) {
    S[2 * i + 2] = s0[i];
 7
 8
                     S[2 * i + 3] = \frac{\pi}{3};
9
10
               scnt = sn0 * 2 + 2;
11
               S[scnt] = \frac{1}{2};
12
13
          void manachar() {
14
               int id = 0, mx = 0;
15
               for(int i = 1; i < scnt; i++) {</pre>
16
                     p[i] = mx > i ? min(p[2 * id - i], mx - i) : 1;
17
                     while(S[i + p[i]] == S[i - p[i]]) p[i]++;
18
19
                     if(i + p[i] > mx) {
20
                          mx = i + p[i];
                          id = i;
21
22
                     }
23
               }
24
          }
25
    }
```

3.4.2 Palindromic Automaton

```
1
    namespace PAM {
         int scnt, S[MAXN];
2
         int pcnt, last, len[MAXN], fail[MAXN], ch[MAXN][26];
3
        int cnt[MAXN]; //节点i表示的本质不同的串的个数(调用count()) int num[MAXN]; //以节点i表示的最长回文串的最右端点为回文串结尾的回文串个数int newnode(int _len) {
4
5
6
             len[pcnt] = _len;
7
             cnt[pcnt] = num[pcnt] = 0;
8
9
             for(int i = 0; i < 26; i++) ch[pcnt][i] = 0;</pre>
10
             return pcnt++;
11
         }
```

```
inline void init() {
12
13
            S[scnt = 0] = -1;
            pcnt = 0;newnode(0);newnode(-1);
14
            fail[0] = 1; last = 0;
15
16
        int getfail(int x) {
17
            while(S[scnt - len[x] - 1] != S[scnt]) x = fail[x];
18
19
            return x;
20
        void extend(int c) {
21
22
            S[++scnt] = c;
            int cur = getfail(last);
23
            if(!ch[cur][c]) {
24
                int now = newnode(len[cur] + 2);
25
                fail[now] = ch[getfail(fail[cur])][c];
26
27
                ch[cur][c] = now;
28
                num[now] = num[fail[now]] + 1;
29
            last = ch[cur][c];
30
            cnt[last]++;
31
32
       void count() {
33
34
            for(int i = pcnt - 1; i >= 0; i--) cnt[fail[i]] += cnt[i];
35
36
   };
```

3.5 Substring Automaton

```
for(int j = 0; j < 26; j++)
    ch[n][j] = ch[n+1][j] = n + 1;

for(int i = n; i >= 1; i--) {
    for(int j = 0; j < 26; j++)
        ch[i-1][j] = ch[i][j];
    ch[i-1][s[i]-'a'] = i;
}</pre>
```

4 Math

4.1 Algebra

4.1.1 FFT

```
//不预处理精度
   const double pi = acos(-1.0);
   const int MAXN = 300003;
4
   struct comp {
5
        double x, y;
        comp operator + (const comp& a) const { return (comp) \{x + a.x, y + a.y\}; }
6
        comp operator - (const comp& a) const { return (comp) {x - a.x, y - a.y}; }
7
        comp operator * (const comp& a) const { return (comp) {x * a.x - y * a.y, x * a.y +
8
       y * a.x; }
9
   };
   int rev[MAXN], T;
10
   comp tmp;
11
   void fft(comp *a, int r) {
12
13
        if(r == -1) for(int i = 0; i < T; i++) a[i] = a[i] * a[i];
        for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
14
        for(int i = 2, mid = 1; i <= T; mid = i, i <<= 1) {
15
16
            comp step = (comp) \{ cos(pi / mid), r * sin(pi / mid) \};
            for(int j = 0; j < T; j += i) {
17
                comp cur = (comp) \{1, 0\};
18
                for(int k = j; k < j + mid; k++, cur = cur * step) {
19
                    tmp = a[k + mid] * cur;
20
                    a[k + mid] = a[k] - tmp;
21
22
                    a[k] = a[k] + tmp;
23
                }
            }
24
25
26
        if(r == -1) for(int i = 0; i < T; i++) a[i].y = (int)(a[i].y / T / 2 + 0.5);
27
28
   comp A[MAXN];
   void init(int n) {
29
        for(T = 1; T \le n; T \le 1);
30
31
        for(int i = 1; i < T; i++) {
            if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
32
33
            else rev[i] = rev[i >> 1] >> 1;
34
            //A[i] = (comp) \{0, 0\};
35
36
    //预处理精度
37
   int rev[MAXN], T;
38
   comp Sin[MAXN], tmp;
void fft(comp *a, int r) {
39
40
        if(r == -1) {
41
            for(int i = 0; i < (T >> 1); i++) Sin[i].y = -Sin[i].y;
42
            for(int i = 0; i < T; i++) a[i] = a[i] * a[i];</pre>
43
44
        for(int i = 1; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
45
        for(int i = 2, mid = 1, s = (T >> 1); i <= T; mid = i, i <<= 1, s >>= 1) {
46
            for(int j = 0; j < T; j += i) {
47
                for(int k = j, cur = 0; k < j + mid; k++, cur += s) {
48
                    tmp = a[k + mid] * Sin[cur];
49
                    a[k + mid] = a[k] - tmp;
50
                    a[k] = a[k] + tmp;
51
52
```

```
53
             }
54
        if(r == -1) for(int i = 0; i < T; i++) a[i].y = (int)(a[i].y / T / 2 + 0.5);
55
56
   comp A[MAXN];
57
    void init(int n) {
58
59
        for(T = 1; T \le n; T \le 1);
        for(int i = 0; i < T; i++) {
60
             if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
61
             else rev[i] = rev[i >> 1] >> 1;
62
             //A[i] = (comp) \{0, 0\};
63
64
        for(int i = 0; i < (T >> 1); i++) {
    Sin[i] = (comp) {cos(2 * pi * i / T), sin(2 * pi * i / T)};
65
66
        }
67
68
   }
   int main() {
69
        scanf("%d%d", &n, &m);
70
        init(n + m);
71
        for(int i = 0; i <= n; i++) scanf("%lf", &A[i].x);</pre>
72
        for(int i = 0; i <= m; i++) scanf("%lf", &A[i].y);</pre>
73
        fft(A, 1);
74
        fft(A, -1);
75
        for(int i = 0; i \le n + m; i++) printf("%d%c", (int)(A[i].y), i == n + m? '\n': '
76
        ');
        return 0;
77
78
    }
```

4.1.2 NTT

4.常用NTT模数:

以下模数的共同g=3189

$p=r\times 2^k+1$	k	g
104857601	22	3
167772161	25	3
469762049	26	3
950009857	21	7
998244353	23	3
1004535809	21	3
2013265921	27	31
2281701377	27	3
3221225473	30	5

```
const int MAXN = 300005, G = 3, mod = 998244353; //or (479LL << 21) + 1
   int rev[MAXN], T;
2
   LL qpow(LL x, LL y) {
3
       LL res = 1;
4
5
       while(y) {
            if(y \& 1) res = res * x % mod;
6
            x = x * x % mod;
7
8
            y >>= 1;
9
10
       return res;
11
   LL A[MAXN], B[MAXN];
12
   void ntt(LL *a, int r) {
```

```
if(r == -1) for(int i = 0; i < T; i++) A[i] = A[i] * B[i] % mod;
14
         for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
15
         for(int i = 2, mid = 1; i <= T; mid = i, i <<= 1) {
    LL gn = qpow(G, (mod - 1) / i);</pre>
16
17
              if(r == -1) gn = qpow(gn, mod - 2);
18
              for(int j = 0; j < T; j += i) {
19
                  LL cur = 1, tmp;
20
                  for(int k = j; k < j + mid; k++, cur = cur * gn % mod) {
    tmp = a[k + mid] * cur % mod;</pre>
21
22
                       a[k + mid] = ((a[k] - tmp) \% mod + mod) \% mod;
23
                       a[k] = (a[k] + tmp) \% mod;
24
                  }
25
26
             }
27
28
         if(r == -1) {
              LL inv = qpow(T, mod - 2);
29
              for(int i = 0; i < T; i++) a[i] = a[i] * inv % mod;
30
31
32
    void init(int n) {
33
         for(T = 1; T <= n; T <<= 1);
34
         for(int i = 0; i < T; i++) {</pre>
35
             if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
36
             else rev[i] = rev[i >> 1] >> 1;
37
38
39
    }
```

4.1.3 MTT

```
#include <bits/stdc++.h>
   using namespace std;
   typedef long long LL;
   const double pi = acos(-1.0);
   const int MAXN = 300003;
6
   struct comp {
7
        double x, y;
        comp operator + (const comp& a) const { return (comp) \{x + a.x, y + a.y\}; }
8
       comp operator - (const comp& a) const { return (comp) {x - a.x, y - a.y}; }
9
       comp operator * (const comp& a) const { return (comp) {x * a.x - y * a.y, x * a.y +
10
       y * a.x; }
11
   #define conj(a) ((comp)\{a.x, -a.y\})
12
   int rev[MAXN], T;
13
   comp Sin[MAXN], tmp;
14
   void fft(comp *a, int r) {
15
        for(int i = 1; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
16
        for(int i = 2, mid = 1, s = (T >> 1); i <= T; mid = i, i <<= 1, s >>= 1) {
17
            for(int j = 0; j < T; j += i) {
18
                for(int k = j, cur = 0; k < j + mid; k++, cur += s) {</pre>
19
                    tmp = a[k + mid] * Sin[cur];
20
                    a[k + mid] = a[k] - tmp;
21
22
                    a[k] = a[k] + tmp;
                }
23
            }
24
       }
25
26
   void init(int n) {
27
28
        for(T = 1; T <= n; T <<= 1);
29
        for(int i = 0; i < T; i++) {</pre>
30
            if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
```

```
else rev[i] = rev[i >> 1] >> 1;
31
32
        for(int i = 0; i < (T >> 1); i++) {
   Sin[i] = (comp) {cos(2 * pi * i / T), sin(2 * pi * i / T)};
33
34
35
36
37
    int n, m, mod;
    void mtt(int *x, int *y) {
38
         for(int i = 0; i < T; i++) (x[i] += mod) %= mod, (y[i] += mod) %= mod;
39
         static comp a[MAXN], b[MAXN];
40
         static comp dfta[MAXN], dftb[MAXN], dftc[MAXN], dftd[MAXN];
41
         for(int i = 0; i < T; i++) {
42
             a[i] = \{x[i] \& 0x7fff, x[i] >> 15\};
43
             b[i] = {y[i] \& 0x7fff, y[i] >> 15};
44
45
         fft(a, 1); fft(b, 1);
46
         for(int i = 0; i < T; i++) {</pre>
47
             int j = (T - i) & (T - 1);
48
             static comp da, db, dc, dd;
49
             da = (a[i] + conj(a[j])) * (comp){0.5, 0};
50
             db = (a[i] - conj(a[j])) * (comp){0, -0.5};
51
             dc = (b[i] + conj(b[j])) * (comp){0.5, 0};
52
             dd = (b[i] - conj(b[j])) * (comp){0, -0.5};
53
             dfta[j] = da * dc;
54
             dftb[j] = da * dd;
55
56
             dftc[j] = db * dc;
57
             dftd[j] = db * dd;
58
         for(int i = 0; i < T; i++) {</pre>
59
             a[i] = dfta[i] + dftb[i] * (comp) {0, 1};
60
             b[i] = dftc[i] + dftd[i] * (comp) {0, 1};
61
62
         //for(int \ i = 0; \ i < (T >> 1); \ i++) \ Sin[i].y = -Sin[i].y;
63
         fft(a, -1); fft(b, -1);
64
         for(int i = 0; i < T; i++) {</pre>
65
             static int da, db, dc, dd;
66
             da = (LL)(a[i].x / T + 0.5) \% mod;
67
             db = (LL)(a[i].y / T + 0.5) \% mod;
68
             dc = (LL)(b[i].x / T + 0.5) \% mod;
69
             dd = (LL)(b[i].y / T + 0.5) \% mod;
70
71
             x[i] = ((da + ((LL)(db + dc) << 15) + ((LL)(dd << 30)) % mod + mod) % mod;
        }
72
73
   }
    int main() {
74
         static int a[MAXN], b[MAXN];
75
        scanf("%d%d%d", &n, &m, &mod);
for(int i = 0; i <= n; i++) scanf("%d", a + i);
for(int i = 0; i <= m; i++) scanf("%d", b + i);</pre>
76
77
78
         init(n + m);
79
80
        mtt(a, b);
         for(int i = 0; i <= n + m; i++) printf("%d%c", a[i], i == n + m ? ^{\prime}\n' : '');
81
82
         return 0;
83
```

4.1.4 FWT

```
void FWT(LL *a,int n) {
for(int i = 2; i <= n; i <<= 1) {
    for(int j = 0; j < n; j += i) {
    for(int d = 0, w = i >> 1; d < w; d++){</pre>
```

```
LL u = a[j + d], v = a[j + d + w];
5
6
                      //xor: a[j + d] = u + v, a[j + d + w] = u - v;
                      //and: a[j + d] = u + v;
7
8
                      //or : a[j + d + w] = u + v;
9
                 }
            }
10
        }
11
12
    void UFWT(LL *a, int n) {
13
        for(int i = 2; i <= n; i <<= 1) {
14
             for(int j = 0; j < n; j += i) {
   for(int d = 0, w = i >> 1; d < w; d++) {</pre>
15
16
                      LL u = a[j + d], v = a[j + d + w];
17
                      //xor: a[j + d] = (u + v) / 2, a[j + d + w] = (u - v) / 2;
18
                      //and: a[j + d] = u - v;
19
                      //or : a[j + d + w] = v - u;
20
21
                 }
22
            }
23
        }
24
    void solve(int n) {
25
        FWT(a, n); FWT(b, n);
26
        for(int i = 0; i < n; i++) a[i] = a[i] * b[i];
27
28
        UFWT(a, n);
29
    }
```

4.1.5 FFT Divide and Conquer

$$f_i = \sum_{j=1}^{i-1} f_j \cdot g_{i-j}$$

```
#include <bits/stdc++.h>
1
   using namespace std;
2
   typedef long long LL:
3
   const int MAXN = 300005, G = 3, mod = 998244353;
4
   namespace NTT {
5
6
        LL A[MAXN], B[MAXN]
        int rev[MAXN], T;
7
8
        LL qpow(LL x, LL y) {
9
            LL res = 1;
10
            while(y) {
                if(y \& 1) res = res * x % mod;
11
                x = x * x % mod;
12
13
                y >>= 1;
            }
14
            return res;
15
16
        void ntt(LL *a, int r) {
17
            for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
18
19
            for(int i = 2, mid = 1; i <= T; mid = i, i <<= 1) {
20
                LL gn = qpow(G, (mod - 1) / i);
21
                if(r == -1) gn = qpow(gn, mod - 2);
                for(int j = 0; j < T; j += i) {
22
23
                    LL cur = 1, tmp;
                    for(int k = j; k < j + mid; k++, cur = cur * gn % mod) {
24
                         tmp = a[k + mid] * cur % mod;
25
                        a[k + mid] = ((a[k] - tmp) \% mod + mod) \% mod;
26
27
                        a[k] = (a[k] + tmp) \% mod;
```

```
28
29
30
            if(r == -1) {
31
                 LL inv = qpow(T, mod - 2);
32
                 for(int i = 0; i < T; i++) a[i] = a[i] * inv % mod;
33
34
35
        void init(int n) {
36
            for(T = 1; T <= n; T <<= 1);</pre>
37
            for(int i = 0; i < T; i++) {</pre>
38
                 if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
39
                 else rev[i] = rev[i >> 1] >> 1;
40
41
            }
42
        }
43
   LL f[MAXN], g[MAXN];
44
   using namespace NTT;
45
   void solve(int 1, int r) {
46
47
        if(l == r) return;
        int mid = (l + r) >> 1;
48
49
        solve(l, mid);
50
        init(r - l);
        for(int i = 0; i < T; i++) A[i] = B[i] = 0;
51
52
        for(int i = 0; i <= mid - l; i++) A[i] = f[i + l];
        for(int i = 0; i <= r - l; i++) B[i] = g[i];
53
54
        ntt(A, 1); ntt(B, 1);
55
        for(int i = 0; i < T; i++) A[i] = A[i] * B[i] % mod;</pre>
        ntt(A, -1);
56
        for(int i = mid + 1; i \le r; i++) f[i] = (f[i] + A[i - l]) % mod;
57
58
        solve(mid + 1, r);
59
   int main() {
60
        int n; scanf("%d", &n);
61
        for(int i = 1; i < n; i++) scanf("%lld", g + i);</pre>
62
        f[0] = 1;
63
64
        solve(0, n - 1);
        for(int i = 0; i < n; i++) printf("%lld%c", f[i], i == n - 1 ? '\n' : '');</pre>
65
66
        return 0;
67
```

4.1.6 Linear Basis

```
1
   //dynamic
2
   const int D = 60;
3
   struct Basis {
4
       vector<int> ind;
5
        vector<LL> base;
6
        Basis() {
7
            ind.resize(D, -1);
8
            base.resize(D);
9
        bool update(LL x, int id) {
10
            for(int i = 0; i < D; i++) if(\simind[i] && x >> i & 1) {
11
                x \sim base[i];
12
13
            if(!x) return 1;
14
            int pos = __builtin_ctzll(x);
15
16
            ind[pos] = id;
17
            base[pos] = x;
```

```
18
            return 0;
        }
19
20
   };
21
    //array
   int Gauss(int n, int m) {
22
        int num = 1;
23
24
        for(int x = 1; x <= n && x <= m; x++) {
25
            int t = 0;
            for(int j = x; j <= m; j++) if(g[j][x]) { t = j; break; }</pre>
26
27
            if(t) {
                 swap(g[x], g[t]);
28
                 for(int i = x + 1; i \le n; i++) {
29
30
                     if(g[i][x]) {
31
                         for(int k = 1; k \le m; k++) g[i][k] ^= g[x][k];
32
                     }
33
                 }
                 num++;
34
35
            }
36
        }
37
        return --num;
38
39
    //long long
   int Gauss() {
40
        int num = 1;
41
42
        for(int k = 61; k >= 0; k--) {
43
            int t = 0;
44
            for(int j = num; j \le cnt; j++) if((A[j] >> k) & 1) { t = j; break; }
45
            if(t) {
46
                 swap(A[t], A[num]);
                 for(int j = num + 1; j \le cnt; j++) if((A[j] >> k) & 1) A[j] ^- A[num];
47
48
                 num++;
            }
49
        }
50
        return --num;
51
   }
52
53
    //det
   LL det(int n){
54
        LL ret = 1;
55
56
        for(int i = 1; i < n; i++){</pre>
            for(int j = i + 1; j < n; j++)
57
58
                 while(a[j][i]){
59
                     LL t = (LL)a[i][i] / a[j][i];
                     for(int k = i; k < n; k++) {
60
                         a[i][k] = (a[i][k]-a[j][k] * t);
61
62
                         swap(a[i][k], a[j][k]);
                     }
63
64
                     ret = -ret;
65
            if(a[i][i] == 0)return 0;
66
            ret = ret * a[i][i];
67
68
        if(ret<0)ret = -ret;</pre>
69
70
        return ret;
   }
71
```

4.1.7 Polynomial

Inverse:

$$A_x * B'_x \equiv 1 \qquad (\text{mod } x^{\frac{n}{2}})$$

$$A_x * B_x \equiv 1 \qquad (\text{mod } x^n)$$

$$B_x \equiv 2 \cdot B'_x - A_x \cdot B'^2_x \qquad (\text{mod } x^n)$$

Division:

$$A_r(x) = x^n A(\frac{1}{x}) \Rightarrow A_r(x)[i] = A(x)[n-i]$$

$$A(x) = B(x) * Q(x) + R(x)$$

$$Q_r(x) = A_r(x) * B_r^{-1}(x)$$

$$R(x) = A(x) - B(x) * Q(x)$$

```
1 //NTT模数
2 #include <bits/stdc++.h>
3 using namespace std;
   typedef long long LL;
const int MAXN = 300005, G = 3, mod = 998244353;
   namespace NTT {
6
        int rev[MAXN], T;
LL qpow(LL x, LL y) {
7
8
             LL res = 1;
9
10
             while(y) {
                  if(y \& 1) res = res * x % mod;
11
                 x = x * x % mod;
12
13
                 y >>= 1;
14
15
             return res;
16
        void ntt(LL *a, int r) {
17
             for(int i = 0; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
18
             for(int i = 2, mid = 1; i <= T; mid = i, i <<= 1) {
19
                  LL gn = qpow(G, (mod - 1) / i);
20
                  if(r == -1) gn = qpow(gn, mod - 2);
21
                  for(int j = 0; j < T; j += i) {
22
                      LL cur = 1, tmp;
23
                      for(int k = j; k < j + mid; k++, cur = cur * gn % mod) {
    tmp = a[k + mid] * cur % mod;</pre>
24
25
                           a[k + mid] = ((a[k] - tmp) \% mod + mod) \% mod;
26
                           a[k] = (a[k] + tmp) \% mod;
27
                      }
28
                 }
29
30
             if(r == -1) {
31
                  LL inv = qpow(T, mod - 2);
32
                  for(int i = 0; i < T; i++) a[i] = a[i] * inv % mod;
33
34
35
        void init(int n) {
36
             for(T = 1; T <= n; T <<= 1);</pre>
37
             for(int i = 0; i < T; i++) {
38
                  if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
39
                  else rev[i] = rev[i >> 1] >> 1;
40
```

```
}
41
42
43
    namespace poly {
44
    using namespace NTT;
45
        void mul(LL *a, LL *b, LL *c, int n, int m) {
46
47
             init(n + m);
             static LL x[MAXN], y[MAXN];
48
49
             for(int i = 0; i < T; i++)
                 x[i] = i \le n ? a[i] : 0;
50
                 y[i] = i \le m ? b[i] : 0;
51
             }
52
             ntt(x, 1); ntt(y, 1);
53
             for(int i = 0; i < T; i++) c[i] = x[i] * y[i] % mod;
54
55
             ntt(c, -1);
56
         void poly_inv(LL *a, LL *b, int n) {
57
             if(n == 1) {
58
                 b[0] = qpow(a[0], mod - 2);
59
60
                 return;
61
             static LL c[MAXN], d[MAXN];
62
             memset(c, 0, n * sizeof(LL));
63
             poly_inv(a, c, n >> 1);
64
             for(int i = 0; i < n; i++) {</pre>
65
66
                 d[i] = a[i];
67
             init(n);
68
             ntt(c, 1); ntt(d, 1);
69
             for(int i = 0; i < T; i++) b[i] = c[i] * (2 + mod - d[i] * c[i] % mod) % mod;
70
71
             ntt(b, -1);
             for(int i = n; i < T; i++) b[i] = 0;
72
73
        void inv(LL *a, LL *b, int n) {//A must be different from B
74
75
             init(n);
76
             poly_inv(a, b, T);
77
        }
         //A_x = B_x * Q_x + R_x;
78
79
         void div(LL *A, LL *B, LL *Q, LL *R, int n, int m) {
             static LL f[MAXN], g[MAXN], inv_g[MAXN];
80
81
             for(int i = 0; i <= n; i++) f[n - i] = A[i];
82
             for(int i = 0; i <= m; i++) g[m - i] = B[i];
83
             //inv(G_r)
             for(int i = n - m + 1; i \le m; i++) g[i] = 0;
84
             inv(g, inv_g, n - m);
85
             //Q
86
87
             mul(f, inv_g, f, n, n - m);
             for(int i = 0; i \le n - m; i++) Q[i] = f[n - m - i];
88
89
90
             mul(Q, B, f, n - m, m);
             for(int i = 0; i < m; i++) R[i] = (A[i] + mod - f[i]) % mod;
91
        }
92
93
    LL A[MAXN], B[MAXN];
94
    LL Q[MAXN], R[MAXN];
95
    int main() {
96
         ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
97
98
         int n, m; cin >> n >> m;
         for(int i = 0; i <= n; i++) cin >> A[i];
99
         for(int i = 0; i <= m; i++) cin >> B[i];
100
101
        poly::div(A, B, Q, R, n, m);
```

```
for(int i = 0; i <= n - m; i++) cout << Q[i] << " "; cout << endl;
102
         for(int i = 0; i < m; i++) cout << R[i] << " "; cout << endl;</pre>
103
         return 0;
104
105
106
     //非NTT模数求逆
    #include <bits/stdc++.h>
107
    using namespace std;
108
     typedef long long LL;
109
    const double pi = acos(-1.0);
110
     const int MAXN = 300003;
111
112
     struct comp {
113
         double x, y;
         comp operator + (const comp& a) const { return (comp) \{x + a.x, y + a.y\}; }
114
         comp operator - (const comp& a) const { return (comp) {x - a.x, y - a.y}; }
115
         comp operator * (const comp& a) const { return (comp) {x * a.x - y * a.y, x * a.y +
116
        y * a.x; }
117
     };
    #define conj(a) ((comp)\{a.x, -a.y\})
118
    int rev[MAXN], T;
119
    comp Sin[MAXN], tmp;
120
    void fft(comp *a, int r) {
121
         for(int i = 1; i < T; i++) if(rev[i] > i) swap(a[rev[i]], a[i]);
122
         for(int i = 2, mid = 1, s = (T >> 1); i <= T; mid = i, i <<= 1, s >>= 1) {
123
             for(int j = 0; j < T; j += i) {
124
                  for(int k = j, cur = 0; k < j + mid; k++, cur += s) {
125
                      tmp = a[k + mid] * Sin[cur];
126
127
                      a[k + mid] = a[k] - tmp;
128
                      a[k] = a[k] + tmp;
129
                 }
             }
130
         }
131
132
    void init(int n) {
133
         //for(T = 1; T \le n; T \le 1);
134
         T = n << 1;
135
         for(int i = 0; i < T; i++) {
136
             if(i & 1) rev[i] = (rev[i >> 1] >> 1) ^ (T >> 1);
137
             else rev[i] = rev[i >> 1] >> 1;
138
139
140
         for(int i = 0; i < (T >> 1); i++) {
             Sin[i] = (comp) \{cos(2 * pi * i / T), sin(2 * pi * i / T)\};
141
142
143
    const int mod = 1e9 + 7;
144
    void mtt(int *x, int *y) {
145
         for(int i = 0; i < T; i++) (x[i] += mod) %= mod, (y[i] += mod) %= mod;
146
147
         static comp a[MAXN], b[MAXN];
         static comp dfta[MAXN], dftb[MAXN], dftc[MAXN], dftd[MAXN];
148
         for(int i = 0; i < T; i++) {
    a[i] = {x[i] & 0x7fff, x[i] >> 15};
    b[i] = {y[i] & 0x7fff, y[i] >> 15};
149
150
151
152
         fft(a, 1); fft(b, 1);
153
         for(int i = 0; i < T; i++) {
154
             int j = (T - i) & (T - 1);
155
             static comp da, db, dc, dd;
156
             da = (a[i] + conj(a[j])) * (comp){0.5, 0};
157
             db = (a[i] - conj(a[j])) * (comp){0, -0.5};
158
             dc = (b[i] + conj(b[j])) * (comp){0.5, 0};
159
             dd = (b[i] - conj(b[j])) * (comp){0, -0.5};
160
             dfta[j] = da * dc;
161
```

```
162
             dftb[j] = da * dd;
             dftc[j] = db * dc;
163
             dftd[j] = db * dd;
164
165
         for(int i = 0; i < T; i++) {
    a[i] = dfta[i] + dftb[i] * (comp) {0, 1};
166
167
             b[i] = dftc[i] + dftd[i] * (comp) {0, 1};
168
169
170
          //for(int \ i = 0; \ i < (T >> 1); \ i++) \ Sin[i].y = -Sin[i].y;
         fft(a, -1); fft(b, -1);
171
         for(int i = 0; i < T; i++) {</pre>
172
              static int da, db, dc, dd;
173
             da = (LL)(a[i].x / T + 0.5) \% mod;
174
             db = (LL)(a[i].y / T + 0.5) \% mod;
175
             dc = (LL)(b[i].x / T + 0.5) \% mod;
176
             dd = (LL)(b[i].y / T + 0.5) \% mod;
177
             x[i] = ((da + ((LL)(db + dc) << 15) + ((LL)(dd << 30)) % mod + mod) % mod;
178
         }
179
180
    LL qpow(LL x, LL y) {
181
182
         LL res = 1;
         while(y) {
183
             if(y & 1) res = res * x % mod;
184
             x = x * x % mod;
185
186
             y >>= 1;
187
188
         return res;
189
190
     void poly_inv(int *a, int *b, int n) {
191
         if(n == 1) {
             b[0] = qpow(a[0], mod - 2);
192
193
             return;
194
         static int c[MAXN], cc[MAXN], d[MAXN];
195
         memset(c, 0, n * sizeof(int)); memset(cc, 0, n * sizeof(int)); memset(d, 0, n *
196
         sizeof(int));
197
         poly_inv(a, c, n >> 1);
         for(int i = 0; i < n; i++) cc[i] = c[i], d[i] = a[i];
198
         init(n);
199
200
         mtt(cc, c);
201
         mtt(d, cc);
         for(int i = 0; i < T; i++) b[i] = (2 * c[i] % mod + mod - d[i]) % mod;
202
         for(int i = 0; i < n; i++) b[n + i] = 0;
203
204
205
    int A[MAXN], B[MAXN];
    int main() {
206
         ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
207
208
         int n; cin >> n; n--;
         for(T = 1; T <= n; T <<= 1);</pre>
209
         for(int i = 0; i <= n; i++) cin >> A[i];
poly_inv(A, B, T);
210
211
         for(int i = 0; i <= n; i++) cout << B[i] << " ";
212
213
         cout << endl;
214
         return 0;
215
    }
216
     //LOJ挑战多项式(备用)
217
    #include <bits/stdc++.h>
218
219 #define rep(i, a, b) for (int i = (a); i <= (b); i++)
    #define per(i, a, b) for (int i = (a); i >= (b); i--)
    #define REP(i, n) for (int i = (0); i < (n); i++)
```

```
#define fi first
222
223
    #define se second
    #define mp make_pair
224
225
    #define pb push_back
226
    using namespace std;
    typedef unsigned long long ull;
227
228
    typedef pair<int, int> pii;
229
    typedef vector<int> poly;
    typedef long long ll;
230
231
     struct ano {
         char a[1 << 25], *s;
char b[1 << 25], *t;</pre>
232
233
234
         ano(): s(a), t(b) { a[fread(a, 1, size of a, stdin)] = 0; }
235
         ~ano() { fwrite(b, 1, t - b, stdout); }
         operator int() {
236
237
             int x = 0;
             while (*s < 48) ++s;
238
             while (*s > 32) x = x * 10 + *s++ - 48;
239
240
             return x;
241
         ll in() {
242
             11 \times = 0;
243
             while (*s < 48) ++s;
244
             while (*s > 32) x = x * 10 + *s++ - 48;
245
             return x;
246
247
248
         void out(int x, char e = ' ') {
249
             static char c[12];
             char *i = c;
250
             if (!x)
251
                  *t++ = 48;
252
             else {
253
                  while (x) {
254
                      int y = x / 10;
255
                      *i++ = x - y * 10 + 48, x = y;
256
257
                  while (i != c) *t++ = *--i;
258
259
             *t++ = e;
260
261
262
    } buf;
263
    const int mod = 998244353;
264
    namespace Poly {
    const int N = (1 \ll 20) + 5, g = 3;
265
    inline int power(int x, int p) {
266
267
         int res = 1;
268
         for (; p; p >>= 1, x = (ll)x * x % mod)
269
             if (p & 1)
270
                  res = (ll)res * x % mod;
271
         return res;
272
    inline int fix(const int x) { return x >= mod ? x - mod : x; }
273
     void dft(poly &A, int n) {
274
         static ull W[N << 1], *H[30], *las = W, mx = 0;</pre>
275
         for (; mx < n; mx++) {
276
277
             H[mx] = las;
278
             ull w = 1, wn = power(g, (mod - 1) >> (mx + 1));
             REP(i, 1 << n) *las++ = w, w = w * wn % mod;
279
280
         if (A.size() != (1 << n))</pre>
281
282
             A.resize(1 \ll n);
```

```
static ull a[N];
283
         for (int i = 0, j = 0; i < (1 << n); ++i) {
284
285
             a[i] = A[j];
             for (int k = 1 \ll (n - 1); (j ^{=} k) < k; k >>= 1)
286
287
288
         for (int k = 0, d = 1; k < n; k++, d <<= 1)
289
             for (int i = 0; i < (1 << n); i += (d << 1)) {
    ull *l = a + i, *r = a + i + d, *w = H[k], t;
290
291
                  for (int j = 0; j < d; j++, l++, r++) {
292
                      t = (*r) * (*w++) % mod;
293
                      *r = *l + mod - t, *l += t;
294
295
296
         REP(i, 1 << n) A[i] = a[i] \% mod;
297
298
    void idft(poly &a, int n) {
299
         a.resize(1 \ll n), reverse(a.begin() + 1, a.end());
300
301
         dft(a, n);
302
         int inv = power(1 << n, mod - 2);</pre>
         REP(i, 1 \ll n) a[i] = (ll)a[i] * inv % mod;
303
304
305
    poly FIX(poly a) {
         while (!a.empty() && !a.back()) a.pop_back();
306
         return a;
307
308
309
    poly add(poly a, poly b, int op = 0) {
310
         a.resize(max(a.size(), b.size()));
311
         REP(i, b.size()) a[i] = fix(op ? a[i] + mod - b[i] : a[i] + b[i]);
312
         return FIX(a);
313
    poly mul(poly a, poly b, int t = 1) {
314
         if (t == 1 && a.size() + b.size() <= 24) {
315
             poly c(a.size() + b.size(), 0)
316
             REP(i, a.size()) REP(j, b.size()) c[i + j] = (c[i + j] + (ll)a[i] * b[j]) % mod;
317
             return FIX(c);
318
319
         }
         int n = 1, aim = a.size() * t + b.size();
320
321
         while ((1 << n) <= aim) n++;
322
         dft(a, n), dft(b, n);
323
         if (t == 1)
             REP(i, 1 << n) a[i] = (ll)a[i] * b[i] % mod;
324
325
             REP(i, 1 << n) a[i] = (ll)a[i] * a[i] % mod * b[i] % mod;
326
327
         idft(a, n), a.resize(aim);
         return FIX(a);
328
329
    poly mul(poly a, int b) {
330
         REP(i, a.size()) a[i] = (ll)a[i] * b % mod;
331
332
         return FIX(a);
333
334
     poly inv(poly a, int n) { // a[0] != 0
335
336
         a.resize(n);
         poly b;
337
         if (n == 1) {
338
             b.pb(power(a[0], mod - 2));
339
340
             return b;
341
         b = inv(a, n + 1 >> 1);
342
         b = add(mul(b, 2), mul(b, a, 2), 1);
343
```

```
return b.resize(n), b;
344
345
    }
346
    poly Der(poly a) {
347
         REP(i, a.size() - 1) a[i] = (ll)(i + 1) * a[i + 1] % mod;
348
349
         return a.pop_back(), a;
350
    }
351
    poly Int(poly a) {
352
         static int inv[N];
         inv[1] = 1;
353
        a.pb(0);
354
        rep(i, 2, a.size()) inv[i] = (ll)(mod - mod / i) * inv[mod % i] % mod;
355
        per(i, a.size() - 1, 1) a[i] = (ll)a[i - 1] * inv[i] % mod;
356
357
         return a[0] = 0, a;
358
359
    poly Ln(poly a, int n) { // a[0] = 1
        a = mul(Der(a), inv(a, n)), a.resize(n - 1);
360
         return FIX(Int(a));
361
362
363
    poly Exp(poly a, int n) { // a[0] = 0
364
        a.resize(n);
365
         poly b, one(1, 1);
366
         if (n == 1)
             return one;
367
        b = Exp(a, n + 1 >> 1);
368
369
        b = mul(b, add(add(a, Ln(b, n), 1), one));
370
        return b.resize(n), b;
371
    }
372
    poly Div(poly a, poly b) {
373
374
         poly c;
         int n = a.size() - 1, m = b.size() - 1;
375
        if (n < m)
376
377
             return c;
        reverse(a.begin(), a.end());
378
        a.resize(n - m + 1);
379
        reverse(b.begin(), b.end());
380
        b.resize(n - m + 1);
381
382
        c = mul(a, inv(b, n - m + 1));
383
         c.resize(n - m + 1);
384
        return reverse(c.begin(), c.end()), c;
385
    poly Mod(poly a, poly b) { return FIX(add(a, mul(Div(a, b), b), 1)); }
386
    inline int chk(int x) { return power(x, (mod - 1) / 2) == 1; }
387
    inline int R() { return rand() % mod; }
388
389
    inline pii mul(pii a, pii b, int w) {
         return mp(((ll)a.fi * b.fi + (ll)a.se * b.se % mod * w) % mod, ((ll)a.fi * b.se + (
390
        ll)a.se * b.fi) % mod);
391
    inline int Sqrt(int x) {
392
393
         if (!chk(x))
394
             return -1;
         int a = R()
395
        while (chk(((11)a * a - x + mod) % mod)) a = R();
396
        int w = ((11)a * a - x + mod) % mod, p = (mod + 1) / 2;
397
        pii res = mp(1, 0), t = mp(a, 1);
398
         for (; p; p >>= 1, t = mul(t, t, w))
399
             if (p & 1)
400
                 res = mul(res, t, w);
401
        assert(!res.se);
402
403
        return min(res.fi, mod - res.fi);
```

```
404
    poly Sqrt(poly a, int n) {
405
         if (n == 1) {
406
407
             poly b(1, Sqrt(a[0]));
408
             return b;
409
         a.resize(n);
410
411
         poly b = Sqrt(a, n + 1 >> 1);
         b = mul(add(b, mul(a, inv(b, n))), (mod + 1) / 2);
412
413
         return b.resize(n), b;
414
    poly fastpow(poly a, ll k, int n) {
415
         a.resize(n), a = FIX(a);
416
         if (!a.size())
417
             return a;
418
         int st = 0, base = 0;
419
         while (!a[st]) ++st;
420
421
         if (st * k >= n)
422
             return a.resize(0), a;
         REP(i, a.size() - st) a[i] = a[i + st];
423
         if (st)
424
             a.resize(a.size() - st);
425
         base = a[0];
426
         11 inv = power(base, mod - 2);
427
         REP(i, a.size()) a[i] = a[i] * inv % mod;
428
429
         a = FIX(Exp(mul(Ln(a, n), k \% mod), n));
430
         if (st) {
431
432
             reverse(a.begin(), a.end());
433
             a.resize(a.size() + st * k)
434
             reverse(a.begin(), a.end());
435
             a.resize(n), a = FIX(a);
436
         base = power(base, k);
437
         REP(i, a.size()) a[i] = (ll)a[i] * base % mod;
438
         return FIX(a);
439
    }
440
        // namespace Poly
441
    using namespace Poly;
442
443
    int main() {
         int n = buf + 1, K = buf;
444
445
         poly a(n, 0), b;
         REP(i, n) a[i] = buf;
446
         b = add(a, Exp(Int(inv(Sqrt(a, n), n)), n), 1);
447
         b[0] = (b[0] + 2 - a[0] + mod) \% mod;
448
         a = Ln(b, n);
449
450
         a[0]++;
         b = Der(fastpow(a, K, n));
451
         b.resize(n - 1);
452
453
         REP(i, n - 1) buf.out(b[i]);
454
         return 0;
455
```

4.1.8 Lagrange Polynomial

$$L(x) = \sum_{i=0}^{n} y_i \prod_{j=0, j \neq i}^{n} \frac{x - x_j}{x_i - x_j}$$

```
//O(n^2)
1
2
   #include <bits/stdc++.h>
3 using namespace std;
   typedef long long LL;
4
   typedef pair<int, int> P;
const int MAXN = 3005, mod = 998244353;
5
    int exgcd(int a, int b, int &x, int &y) {
7
8
        int d = a;
        if(b != 0) {
9
            d = exgcd(b, a \% b, y, x);
10
            y = (a / b) * x;
11
12
        else {
13
            x = 1; y = 0;
14
15
16
        return d;
17
   int inv(int a) {
18
        int x, y;
19
20
        exgcd(a, mod, x, y);
        return (x % mod + mod) % mod;
21
22
23
   struct Lagrange {
24
        int n, a[MAXN][2];
25
        void init() {
26
            for(int i = 0; i <= n; i++) a[i][0] = a[i][1] = 0;
27
            n = 0;
            a[0][1] = 1;
28
29
        int query(int x, int q = 0) {
30
31
            int res = 0;
            for(int i = n; i >= 0; i--) res = ((LL)res * x + a[i][q]) % mod;
32
            return res;
33
34
        void update(int x, int y) {
35
            a[n][0] = 0;
36
            int v = (LL)(y - query(x) + mod) % mod * inv(query(x, 1)) % mod;
37
            for(int i = 0; i \le n; i++) a[i][0] = (a[i][0] + (LL)a[i][1] * v) % mod;
38
39
            a[++n][1] = 0;
40
            for(int i = n; i; i--) a[i][1] = (a[i - 1][1] + (LL)a[i][1] * (mod - x)) % mod;
41
            a[0][1] = (LL)a[0][1] * (mod - x) % mod;
42
        }
43
   }p;
44
   int main() {
        ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
45
46
        int Q;
47
        cin >> Q;
        int op, x, y;
48
        p.n = 0;
49
50
        p.init();
51
        while(Q--) {
52
            cin >> op >> x;
53
            if(op == 1) {
                 cin >> y;
54
55
                 p.update(x, y);
56
            else cout << p.query(x) << endl;</pre>
57
58
59
        return 0;
60
   //f(x) = y, \sharp f(z_i)
```

```
62 #include <bits/stdc++.h>
63
    using namespace std;
64 #define ll long long
    int n, m;
65
    vector<int> x, y, z;
66
67
    namespace Poly{
         const int P = 998244353;
68
69
         vector<int> ans;// for Evaluate()
70
         vector<vector<int>>> p;// for Evaluate() & Interpolate()
         inline int Pow(ll x, int y=P-2){ // x^y
71
72
             int ans=1;
             for(; y; y>>=1, x=x*x%P) if(y&1) ans=ans*x%P;
73
74
             return ans;
75
76
         inline int Ge(int x){ int n=1; while(n<=x) n<<=1; return n;}</pre>
         inline int Mod(int x){ return x<P?x:x-P;}</pre>
77
         inline void NTT(vector<int> &f, int g, int n){
78
79
             f.resize(n);
             for(int i=0, j=0; i<n; ++i){
80
81
                 if(i>j) swap(f[i], f[j]);
                 for(int k=n>1; (j^k=k)<k; k>=1);
82
83
             vector<int> w(n>>1);
84
             for(int i=1; i<n; i<<=1){
85
                 for(int j=w[0]=1, w0=(g==1?Pow(3, (P-1)/i/2):Pow(Pow(3, (P-1)/i/2))); j<i;
86
        ++j) w[j]=(ll)w[j-1]*w0%P;
87
                 for(int j=0; j<n; j+=i<<1){</pre>
88
                      for(int k=j; k<j+i; ++k){</pre>
89
                          int t=(ll)f[k+i]*w[k-j]%P;
90
                          f[k+i]=Mod(f[k]-t+P);
91
                          f[k]=Mod(f[k]+t);
                     }
92
                 }
93
94
95
             if(g==-1) for(int i=0, I=Pow(n); i<n; ++i) f[i]=(ll)f[i]*I%P;
96
97
         inline vector<int> Add(const vector<int> &f, const vector<int> &g){
98
             vector<int> ans=f;
             for(unsigned i=0; i<f.size(); ++i) (ans[i]+=g[i])%=P;</pre>
99
100
             return ans;
101
102
         inline vector<int> Mul(const vector<int> &f, const vector<int> &g){// f*g
             vector<int> F=f, G=g;
103
             int p=Ge(f.size()+g.size()-2);
104
             NTT(F, 1, p), NTT(G, 1, p);
105
             for(int i=0; i<p; ++i) F[i]=(ll)F[i]*G[i]%P;</pre>
106
             NTT(F, -1, p);
107
             return F.resize(f.size()+g.size()-1), F;
108
109
         inline vector<int> PolyInv(const vector<int> &f, int n=-1){// 1/f
110
111
             if(n==-1) n=f.size();
             vector<int> ans;
112
             if(n==1) return ans.push_back(Pow(f[0])), ans;
113
             ans=PolyInv(f, (n+1)/2);
114
             vector<int> tmp(&f[0], &f[0]+n);
115
             int p=Ge(n*2-2);
116
             NTT(tmp, 1, p), NTT(ans, 1, p);
117
             for(int i=0; i<p; ++i) ans[i]=(2-(11)ans[i]*tmp[i]%P+P)*ans[i]%P;</pre>
118
             NTT(ans, -1, p);
119
             return ans.resize(n), ans;
120
121
```

```
inline void PolyDiv(const vector<int> &a, const vector<int> &b, vector<int> &d,
122
        vector<int> &r){//a=d*b+r}
             if(b.size()>a.size()) return (void)(d.clear(), r=a);
123
124
125
             vector<int> A=a, B=b, iB;
             int n=a.size(), m=b.size();
126
             reverse(A.begin(), A.end()), reverse(B.begin(), B.end());
127
             B.resize(n-m+1), iB=PolyInv(B, n-m+1);
128
129
             d=Mul(A, iB);
             d.resize(n-m+1), reverse(d.begin(), d.end());
130
131
             r=Mul(b, d);
132
             for(int i=0; i<m-1; ++i) r[i]=(P+a[i]-r[i])%P;</pre>
133
134
             r.resize(m-1);
135
         inline vector<int> Derivative(const vector<int> &a){ // a '
136
137
             vector<int> ans;
             ans.resize(a.size()-1);
138
             for(unsigned i=1; i<a.size(); ++i) ans[i-1]=(ll)a[i]*i%P;</pre>
139
140
141
        void Evaluate_Interpolate_Init(int l, int r, int t, const vector<int> &a){
142
             if(l==r) return p[t].clear(), p[t].push_back(P-a[l]), p[t].push_back(1);
143
             int mid=(l+r)/2, k=t<<1;</pre>
144
             Evaluate_Interpolate_Init(l, mid, k, a), Evaluate_Interpolate_Init(mid+1, r, k
145
        |1, a);
146
             p[t]=Mul(p[k], p[k|1]);
147
        }
         inline void Evaluate(int l, int r, int t, const vector<int> &f, const vector<int> &a
148
        ){
             if(r-l+1<=512){
149
                 for(int i=1; i<=r; ++i){</pre>
150
                     int x=0, j=f.size(), a1=a[i], a2=(l1)a[i]*a[i]%P, a3=(l1)a[i]*a2%P, a4=(
151
        ll)a[i]*a3%P, a5=(ll)a[i]*a4%P, a6=(ll)a[i]*a5%P, a7=(ll)a[i]*a6%P, a8=(ll)a[i]*a7%P
                     while(j > = 8)
152
                     x=((ll)x*a8+(ll)f[j-1]*a7+(ll)f[j-2]*a6+(ll)f[j-3]*a5+(ll)f[j-4]*a4+(ll)
153
        f[j-5]*a3+(11)f[j-6]*a2+(11)f[j-7]*a1+f[j-8])%P, j-=8;
                     while(j--) x=((ll)x*a[i]+f[j])%P;
154
155
                     ans.push_back(x);
156
                 }
157
                 return;
158
             }
             vector<int> tmp;
159
             PolyDiv(f, p[t], tmp, tmp);
160
             Evaluate((1, (1+r)/2, t<<1, tmp, a), Evaluate((1+r)/2+1, r, t<<1/1, tmp, a);
161
162
        inline vector<int> Evaluate(const vector<int> &f, const vector<int> &a, int flag=-1)
163
        \{// f(a_i)
             if(flag==-1) p.resize(a.size()<<2), Evaluate_Interpolate_Init(0, a.size()-1, 1,</pre>
164
        a);
             ans.clear(), Evaluate(0, a.size()-1, 1, f, a);
165
166
             return ans;
167
         vector<int> Interpolate(int l, int r, int t, const vector<int> &x, const vector<int>
168
         &f){
169
             if(l==r){
                 vector<int> ans;
170
                 return ans.push_back(f[l]), ans;
171
172
             int mid=(l+r)/2, k=t<<1;
173
```

```
return Add(Mul(Interpolate(l, mid, k, x, f), p[k|1]), Mul(Interpolate(mid+1, r,
174
         k|1, x, f), p[k]);
175
         inline vector<int> Interpolate(const vector<int> &x, const vector<int> &y)\{// (x_i, x_i)\}
176
         y_{\perp}i)
177
              int n=x.size();
              p.resize(n<<2), Evaluate_Interpolate_Init(0, n-1, 1, x);</pre>
178
              vector<int> f=Evaluate(Derivative(p[1]), x, 0);
179
              for(int i=0; i<n; ++i) f[i]=(l1)y[i]*Pow(f[i])%P;</pre>
180
181
              return Interpolate(0, n-1, 1, x, f);
         }
182
     }
183
    using namespace Poly;
184
    int main() {
185
         cin >> n; x.resize(n), y.resize(n);
186
         for(int i=0; i<n; ++i) cin >> x[i], cin >> y[i];
187
         cin >> m, z.resize(m);
188
         for(int i=0; i<m; ++i) cin >> z[i];
189
         x=Evaluate(Interpolate(x, y), z);
for(int i:x) cout << i << " ";</pre>
190
191
192
         return 0;
193
```

4.1.9 BM Alogrithm

```
1 #include<bits/stdc++.h>
using namespace std;
3 #define rep(i,a,n) for (int i=a;i<n;i++)</pre>
   #define per(i,a,n) for (int i=n-1;i>=a;i--)
   #define pb push_back
6
   #define mp make_pair
   #define all(x) (x).begin(),(x).end()
7
   #define fi first
9
   #define se second
10 #define SZ(x) ((int)(x).size())
   typedef vector<int> VI;
11
12 typedef long long ll;
13 typedef pair<int,int> PII;
14 const ll mod=1000000007;
   ll powmod(ll a, ll b) {ll res=1; a\%=mod; assert(b>=0); for(;b;b>>=1){if(b&1)res=res*a\%mod;
       a=a*a%mod;}return res;}
16
   // head
   namespace linear_seq {
17
        const int N=10010;
18
       11 res[N],base[N],_c[N],_md[N];
19
20
21
       vector<int> Md;
       void mul(ll *a,ll *b,int k) {
22
23
            rep(i,0,k+k) _c[i]=0;
            rep(i,0,k) if (a[i]) rep(j,0,k) _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;
24
25
            for (int i=k+k-1;i>=k;i--) if (_c[i])
26
                rep(j,0,SZ(Md)) _c[i-k+Md[j]]=(_c[i-k+Md[j]]-_c[i]*_md[Md[j]])%mod;
27
            rep(i,0,k) a[i]=_c[i];
28
        int solve(ll n,VI a,VI b) { // a 系数 b 初值 b[n+1]=a[0]*b[n]+...
29
              printf("%d \mid n", SZ(b));
30
            ll ans=0,pnt=0;
31
32
            int k=SZ(a);
33
            assert(SZ(a)==SZ(b));
34
            rep(i,0,k) _md[k-1-i]=-a[i];_md[k]=1;
```

```
35
            Md.clear();
36
             rep(i,0,k) if (_md[i]!=0) Md.push_back(i);
37
             rep(i,0,k) res[i]=base[i]=0;
             res[0]=1;
38
            while ((1ll<<pnt)<=n) pnt++;</pre>
39
             for (int p=pnt;p>=0;p--) {
40
                 mul(res,res,k);
41
42
                 if ((n>>p)&1) {
                     for (int i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
43
                     rep(j,0,SZ(Md)) res[Md[j]]=(res[Md[j]]-res[k]*_md[Md[j]])%mod;
44
                 }
45
46
            rep(i,0,k) ans=(ans+res[i]*b[i])%mod;
47
            if (ans<0) ans+=mod;</pre>
48
49
            return ans;
50
        VI BM(VI s) {
51
            VI C(1,1),B(1,1);
52
            int L=0, m=1, b=1;
53
54
             rep(n,0,SZ(s)) {
55
                 ll d=0;
                 rep(i,0,L+1) d=(d+(ll)C[i]*s[n-i])%mod;
56
                 if (d==0) ++m;
57
                 else if (2*L <= n) {
58
59
                     VI T=C;
60
                     11 c=mod-d*powmod(b,mod-2)%mod;
61
                     while (SZ(C) < SZ(B) + m) C.pb(0);
62
                     rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
63
                     L=n+1-L; B=T; b=d; m=1;
64
                 } else {
                     11 c=mod-d*powmod(b,mod-2)%mod;
65
                     while (SZ(C)<SZ(B)+m) C.pb(0);</pre>
66
                     rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
67
68
                     ++m;
69
                 }
70
            }
71
            return C;
72
73
        int gao(VI a,ll n) {
74
            VI c=BM(a);
75
            c.erase(c.begin());
             rep(i,0,SZ(c)) c[i]=(mod-c[i])%mod;
76
             return solve(n,c,VI(a.begin(),a.begin()+SZ(c)));
77
        }
78
79
    };
80
81
    int main() {
        while (~scanf("%d",&n)) {
82
            vector<int>v;
83
            v.push_back(1);
84
85
            v.push_back(2);
86
            v.push_back(4);
            v.push_back(7);
87
            v.push_back(13);
88
89
            v.push_back(24);
90
             //VI\{1,2,4,7,13,24\}
            printf("%d\n",linear_seq::gao(v,n-1));
91
92
        }
93
    }
```

4.2 Math Theory

4.2.1 Inverse

```
//O(logn) 求n的 逆元
   const int mod = 1e6 + 3;
   int exgcd(int a, int b, int &x, int &y) {
4
        int d = a;
5
        if(b != 0) {
            d = exgcd(b, a \% b, y, x);
6
            y = (a / b) * x;
7
8
9
        else {
10
            x = 1; y = 0;
11
12
       return d;
13
14
   int inverse(int a) {
        int x, y;
15
16
        exgcd(a, mod, x, y);
        return (x % mod + mod) % mod;
17
18
   int inverse(int a) { return qpow(a, mod - 2); }
19
   //O(n)求1~n的逆元
20
   int inv[MAXN];
21
   void init() {
22
        inv[0] = inv[1] = 1;
23
        for(int i = 2; i < MAXN; i++) inv[i] = (long long)(mod - mod / i) * <math>inv[mod \% i] \%
24
25
   }
```

4.2.2 Lucas

```
//mod很小可以预处理逆元的情况
1
   void init() {
2
        fac[0] = 1;
3
        for(int i = 1; i < mod; i++) fac[i] = (long long)fac[i - 1] * i % mod;
4
5
        inv[0] = inv[1] = 1;
       for(int i = 2; i < mod; i++) inv[i] = (long long)(mod - mod / i) * <math>inv[mod \% i] \%
6
7
       for(int i = 1; i < mod; i++) inv[i] = (long long)inv[i] * inv[i - 1] % mod;</pre>
8
   int C(int a, int b) {
9
       if(b > a) return 0;
10
       if(a < mod) return (long long)fac[a] * inv[b] % mod * inv[a - b] % mod;</pre>
11
        return (long long)C(a / mod, b / mod) * C(a % mod, b % mod) % mod;
12
   }
13
    //mod过大不能预处理逆元的情况
14
   LL qpow(LL x, LL y) {
15
       LL res = 1;
16
17
       while(y) {
18
           if(y \& 1) res = res * x % mod;
           x = x * x % mod;
19
20
           y >>= 1;
       }
21
22
       return res;
23
   LL C(LL a, LL b) {
24
   if(b > a) return 0;
25
```

```
if(b > a - b) b = a - b;
26
27
         LL s1 = 1, s2 = 1;
         for(LL i = 0; i < b; i++) {
    s1 = s1 * (a - i) % mod;
28
29
             s2 = s2 * (i + 1) % mod;
30
31
         }
32
        return s1 * qpow(s2, mod - 2) % mod;
33
    LL lucas(LL a, LL b) {
34
        if(a < mod) return C(a, b);</pre>
35
         return lucas(a / mod, b / mod) * C(a % mod, b % mod);
36
37
```

4.2.3 CRT && exCRT

 $x \equiv a_i \pmod{m_i}$

```
namespace CRT {
1
2
        LL m[MAXN], a[MAXN];
        LL exgcd(LL _a, LL _b, LL &x, LL &y) {
    if(!_b) {
3
4
5
                 x = 1; y = 0;
                 return _a;
6
7
            LL d = exgcd(_b, _a % _b, y, x);
8
            y = (a / b) * x;
9
            return d;
10
11
        LL crt(int n) {
12
            LL M = 1, tmp, res = 0, x, y;
13
             for(int i = 1; i <= n; i++) M *= m[i];</pre>
14
             for(int i = 1; i <= n; i++) {
15
16
                 tmp = M / m[i];
17
                 exgcd(tmp, m[i], x, y);
                 x = (x + m[i]) \% m[i];

res = (a[i] * x % M * tmp % M + res) % M;
18
19
20
            }
21
            return res;
22
        }
23
    namespace EXCRT {
24
        LL m[MAXN], a[MAXN];
25
26
        LL exgcd(LL _a, LL _b, LL &x, LL &y) {
             if(!_b) {
27
28
                 x = 1; y = 0;
29
                 return _a;
30
31
            LL d = exgcd(_b, _a % _b, y, x);
            y = (_a / _b) * x;
32
33
            return d;
34
        LL excrt(int n) {
35
            LL M = m[1], A = a[1], x, y, d, tmp;
36
             for(int i = 2; i <= n; i++) {</pre>
37
                 d = exgcd(M, m[i], x, y);
38
                 if((A - a[i]) % d) return -1; //No solution
39
                 tmp = M / d; M *= m[i] / d;
40
                 y = (A - a[i]) / d % M * y % M;
41
```

```
y = (y + tmp) \% tmp;
42
                 A = (m[i] \% M * y \% M + a[i]) \% M;
43
                 A = (A + M) \% M;
44
45
            return A;
46
47
        }
            LL inv(LL _a, LL _b) {
48
            LL x, y;
exgcd(_a, _b, x, y);
49
50
             return (x % _b + _b) % _b;
51
52
        LL excrt(int n) {
53
            LL M = m[1], A = a[1], x, y, d, c, tmp;
54
             for(int i = 2; i <= n; i++) {</pre>
55
56
                 d = exgcd(M, m[i], x, y);
57
                 c = a[i] - A;
                 if(c % d) return -1;
58
                 c = (c \% m[i] + m[i]) \% m[i];
59
                 M /= d; m[i] /= d;
60
                 c = c / d * inv(M % m[i], m[i]) % m[i];
61
62
                 tmp = M;
                 M *= m[i] * d;
63
                 A = (c^* tmp \% M * d \% M + A) \% M;
64
65
66
            return A;
67
        }
68
   }
```

4.2.4 BSGS

```
const int MOD = 76543;
   int hs[MOD + 5], head[MOD + 5], nxt[MOD + 5], id[MOD + 5], ecnt;
   void insert(int x, int y) {
        int k = x \% MOD;
4
       hs[ecnt] = x, id[ecnt] = y, nxt[ecnt] = head[k], head[k] = ecnt++;
5
   }
6
   int find(int x) {
7
        int k = x \% MOD;
8
9
        for(int i = head[k]; i; i = nxt[i])
10
            if(hs[i] == x)
                return id[i];
11
12
        return -1;
13
   }
14
   int BSGS(int a, int b, int c){
15
       memset(head, 0, sizeof head); ecnt = 1;
16
        if(b == 1) return 0;
        int m = sqrt(c * 1.0), j;
17
18
        LL x = 1, p = 1;
        for(int i = 0; i < m; i++, p = p * a % c)
19
            insert(p * b % c, i);
20
        for(LL i = m; ;i += m){
21
            if((j = find(x = x * p % c)) != -1) return i - j;
22
23
            if(i > c) break;
24
       }
25
       return -1;
26
```

4.2.5 Miller-Rabin && PollardRho

```
1
   LL ksc(LL a,LL n,LL mod){
2
        LL ret=0;
3
        for(;n;n>>=1){
             if(n&1){ret+=a;if(ret>=mod)ret-=mod;}
4
5
            a \le 1; if(a \ge mod)a = mod;
6
        }
7
        return ret;
8
    LL ksm(LL a, LL n, LL mod){
9
10
        LL ret = 1;
        for(;n;n>>=1){
11
             if(n&1)ret=ksc(ret,a,mod);
12
            a=ksc(a,a,mod);
13
14
        }
15
        return ret;
16
    int millerRabin(LL n){
17
        if(n<2 || (n!=2 && !(n&1)))return 0;
18
19
        LL d=n-1; for(;!(d&1); d>>=1);
20
        for(int i=0;i<20;++i){
            LL a=rand()\%(n-1)+1;
21
            LL t=d, m=ksm(a,d,n);
22
            for(;t!=n-1 && m!=1 && m!=n-1;m=ksc(m,m,n),t<<=1);</pre>
23
24
            if(m!=n-1 && !(t&1)) return 0;
25
        }
26
        return 1;
27
   LL cnt,fact[100];
28
    LL gcd(LL a,LL b){return !b?a:gcd(b,a%b);}
29
30
   LL pollardRho(LL n, int a){
        LL x=rand()%n, y=x, d=1, k=0, i=1;
31
32
        while(d==1){
33
            ++k;
34
            x=ksc(x,x,n)+a;if(x>=n)x-=n;
35
            d=gcd(x>y?x-y:y-x,n);
36
            if(k==i){y=x;i<<=1;}
37
        if(d==n)return pollardRho(n,a+1);
38
39
        return d;
40
   }
41
    void findfac(LL n){
        if(millerRabin(n)){fact[++cnt]=n; return;}
42
        LL p=pollardRho(n,rand()%(n-1)+1);
43
        findfac(p);
44
45
        findfac(n/p);
46
    4.2.6 \varphi(n)
 1
    int phi(int x) {
        int res = x;
2
        for(int i = 2; i * i <= x; i++) {</pre>
3
            if(x \% i == 0) {
4
                 res = res / i * (i - 1);
5
                 while(x % i == 0) x /= i;
6
7
            }
```

8

10

11 }

if(x > 1) res = res / x * (x - 1);

return res;

4.2.7 Euler Sieve

```
int prime[MAXN], cnt, phi[MAXN], mu[MAXN];
1
   bool isp[MAXN];
3
   int min_pow[MAXN];
                         //最小质因子最高次幂
4
   int min_sum[MAXN];
5
                         //1+p+p^2+...+p^k
   int div_sum[MAXN];
6
                         //约数和
   int min_index[MAXN]; //最小质因子的指数
8
                         //约数个数
   int div_num[MAXN];
9
   void Euler(int n) {
10
        mu[1] = phi[1] = div_num[1] = div_sum[1] = 1;
11
        for(int i = 2; i <= n; i++) {
12
13
            if(!isp[i]) {
                prime[++cnt] = min_pow[i] = i;
14
15
                phi[i] = i - 1;
                mu[i] = -1;
16
17
                min_index[i] = 1; div_num[i] = 2;
                div_sum[i] = min_sum[i] = i + 1;
18
19
20
            for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
                isp[i * prime[j]] = 1;
21
22
                if(i % prime[j] == 0) {
23
                    phi[i * prime[j]] = phi[i] * prime[j];
                    mu[\bar{i} * prime[\bar{j}]] = 0;
24
25
                    min_index[i * prime[j]] = min_index[i] + 1;
26
                    div_num[i * prime[j]] = div_num[i] / (min_index[i] + 1) * (min_index[i *
27
        prime[j]] + 1);
28
                    min_sum[i * prime[j]] = min_sum[i] + min_pow[i] * prime[j];
29
                    div_sum[i * prime[j]] = div_sum[i] / min_sum[i] * min_sum[i * prime[j]];
30
                    min_pow[i * prime[j]] = min_pow[i] * prime[j];
31
32
                    break;
33
                phi[i * prime[j]] = phi[i] * (prime[j] - 1);
34
                mu[i * prime[j]] = -mu[i];
35
36
                div_num[i * prime[j]] = div_num[i] << 1;</pre>
37
                min_index[i * prime[j]] = 1;
38
39
40
                div_sum[i * prime[j]] = div_sum[i] * (prime[j] + 1);
                min_pow[i * prime[j]] = prime[j];
41
42
                min_sum[i * prime[j]] = prime[j] + 1;
43
           }
       }
44
45
   }
```

4.2.8 DuJiao Sieve

$$\sum_{i=1}^{n} \varphi(i)$$

```
vector<int> prime;
int phi[MAXN], P[MAXN];
bool isp[MAXN];
unordered_map<LL, int> mp;
```

```
void Euler(int n) {
5
6
        phi[1] = 1;
7
        for(int i = 2; i <= n; i++) {</pre>
8
            if(!isp[i]) {
9
                prime.push_back(i);
10
                 phi[i] = i - 1;
11
            for(auto x : prime) {
12
                 if(i * x > n) break;
13
                 isp[i * x] = 1;
14
                 if(i % x == 0) {
15
                     phi[i * x] = phi[i] * x;
16
                     break;
17
18
                phi[i * x] = phi[i] * (x - 1);
19
            }
20
21
        for(int i = 1; i <= n; i++) P[i] = (P[i - 1] + phi[i]) % mod;
22
23
   LL cal(LL n) {
24
        if(n < MAXN) return P[n];</pre>
25
        if(mp.count(n)) return mp[n];
26
        LL res = 0;
27
        for(LL i = 2, last; i <= n; i = last + 1) {
28
29
            last = n / (n / i);
30
            res += (last - i + 1) % mod * cal(n / i) % mod;
31
            res %= mod;
32
        mp[n] = ((_int128)n * (n + 1) / 2 % mod + mod - res) % mod;
33
34
        return mp[n];
   }
35
```

 $\sum_{i=1}^{n} \mu(i)$

```
LL cal(LL n) {
1
        if(n < MAXN) return M[n];</pre>
2
3
        if(mp.count(n)) return mp[n];
4
        LL res = 0;
        for(LL i = 2, last; i <= n; i = last + 1) {</pre>
5
            last = n / (n / i);
6
             res += (last - i + 1) * cal(n / i);
7
8
9
        mp[n] = 1 - res;
10
        return 1 - res;
11
   }
```

4.2.9 Min_25 Sieve

$$\sum_{i=1}^{n} \varphi(i)$$

$$g_{k,n} \text{ and } h_{k,n} \text{ Count}$$

$$\sum_{i=1}^{n} i^{k}$$

```
#include <bits/stdc++.h>
1
2
   using namespace std;
3 typedef long long LL;
   const int MAXN = 1e6 + 5, mod = 1e9 + 7;
   const int inv2 = (mod + 1) / 2, inv6 = (mod + 1) / 6;
   int prime[MAXN], isp[MAXN], cnt;
   LL g[3][MAXN << 1], h[3][MAXN << 1];
   LL w[MAXN << 1];
   int id1[MAXN], id2[MAXN];
9
   inline int MOD(LL x) { return x >= mod ? x - mod : x; }
10
   //inline int MOD(LL\ x) { return x\ \%\ mod; } inline int add(LL x, LL y) { return MOD(MOD(x) + MOD(y)); }
11
12
   void Euler(int n) {
13
        for(int i = 2; i <= n; i++) {
14
15
             if(!isp[i]) {
16
                 prime[++cnt] = i;
                 h[0][cnt] = h[0][cnt - 1] + 1;
17
                 h[1][cnt] = add(h[1][cnt - 1], i);
18
                 h[2][cnt] = add(h[2][cnt - 1], (LL)i * i % mod);
19
20
             for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
21
                 isp[i * prime[j]] = 1;
22
                 if(i % prime[j] == 0) {
23
                     break:
24
25
26
            }
27
        }
28
29
   LL n;
30
   int sz, m;
   inline int id(LL x) {
31
        return x <= sz ? id1[x] : id2[n / x];</pre>
32
   }
33
    //f(p \hat{k})
34
   inline int f(int p, LL pk) {
35
        return pk / p * (p - 1) % mod;
36
37
   LL S(LL x, int y) {
38
39
        if(x <= 1 || prime[y] > x) return 0;
40
        //G(x) - H(j-1)
41
        LL res = add(add(g[1][id(x)], mod - g[0][id(x)]), mod - add(h[1][y - 1], mod - h[0][
       y - 1]));
        for(int j = y, k = 1; j \leftarrow cnt \&\& (LL)prime[j] * prime[j] \leftarrow x; j++, k = 1) {
42
             for(LL pk = prime[j]; pk * prime[j] <= x; pk *= prime[j], k++) {</pre>
43
                 res = add(res, S(x / pk, j + 1) * f(prime[j], pk) % mod + f(prime[j], pk *
44
        prime[j]));
45
            }
46
        }
47
        return res;
48
49
    int main() {
        ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
50
51
        cin >> n;
        sz = sqrt(n);
52
        Euler(sz);
53
        for(LL i = 1, last, t; i <= n; i = last + 1) {
54
            last = n / (n / i);
55
            w[++m] = n / i, t = n / i \% mod;
56
            w[m] \le sz ? id1[w[m]] = m : id2[last] = m;
57
            q[0][m] = MOD(t + mod - 1);
58
            g[1][m] = add(t * (t + 1) % mod * inv2 % mod, mod - 1);
```

```
g[2][m] = add((2 * t + 1) % mod * t * (t + 1) % mod * inv6 % mod, mod - 1);
60
61
62
           for(int j = 1; j <= cnt; j++) {</pre>
                 for(int i = 1; i <= m && (LL)prime[j] * prime[j] <= w[i]; i++) {</pre>
63
                       g[0][i] = MOD(g[0][i] + mod - (g[0][id(w[i] / prime[j])] - h[0][j - 1]));
g[1][i] = MOD(g[1][i] + mod - ((LL)prime[j] * MOD(g[1][id(w[i] / prime[j])]
64
65
           + \  \, \text{mod} \, - \  \, \text{h[1][j-1]) \% mod));} \\ g[2][i] = \  \, \text{MOD}(g[2][i] \, + \, \text{mod} \, - \, ((LL)prime[j] \ * \, prime[j] \ \% \, \, \text{mod} \, * \, \text{MOD}(g[2][id(w)]) ) ] 
66
           [i] / prime[j])] + mod - h[2][j - 1]) % mod));
67
68
            //S(n, 1) + F(1);
69
           LL ans = MOD(S(n, 1) + 1);
70
71
           cout << ans << endl;
72
           return 0;
73
    }
```

4.2.10 Möbius Inversion

$$\sum_{i=1}^{n} \sum_{j=1}^{m} lcm(i,j) \pmod{p}$$

```
int mu[MAXN], prime[MAXN], sum[MAXN], cnt;
1
   bool isp[MAXN];
2
   void getmu(int n) {
3
        mu[1] = 1;
4
        for(int i = 2; i <= n; i++) {
5
6
            if(!isp[i]) {
7
                mu[i] = -1;
8
                 prime[++cnt] = i;
9
10
            for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
                 isp[i * prime[j]] = 1;
11
                 if(i % prime[j] == 0) {
12
                     mu[i * prime[j]] = 0;
13
                     break;
14
15
                mu[i * prime[j]] = -mu[i];
16
            }
17
        }
18
19
20
   ll n, m, ans;
   ll query(ll x, ll y) { return (x * (x + 1) / 2 % mod) * (y * (y + 1) / 2 % mod) % mod; }
21
22
   11 F(11 x, 11 y) 
23
        ll res = 0, last;
24
        for(ll i = 1; i \le min(x, y); i = last + 1) {
            last = min(x / (x / i), y / (y / i));
25
            res = (res + (sum[last] - sum[i - 1]) * query(x / i, y / i) % mod) % mod;
26
27
28
        return res;
29
   int main() {
30
31
        cin>>n>>m;
32
        getmu(min(n, m));
        for(ll \ i = 1; \ i \le min(n, m); \ i++) \ sum[i] = (sum[i - 1] + (i * i * mu[i]) \% \ mod) \%
33
        mod;
        ll last;
34
        for(ll d = 1; d <= min(n, m); d = last + 1) {</pre>
35
```

```
last = min(n / (n / d), m / (m / d));
ans = (ans + (last - d + 1) * (d + last) / 2 % mod * F(n / d, m / d) % mod);
mod;
last = min(n / (n / d), m / (m / d));
ans = (ans + (last - d + 1) * (d + last) / 2 % mod * F(n / d, m / d) % mod) %
mod;
ans = (ans + mod) % mod;
cout < ans < endl;
return 0;
}</pre>
```

5 Geometry

5.1 Commonly Definition and Functions

5.1.1 Const and Functions

```
namespace CG{
1
        #define Point Vector
2
3
        const double pi=acos(-1.0);
        const double inf=1e100;
4
5
        const double eps=1e-9;
        template <typename T> inline T Abs(T x){return x>0?x:-x;}
6
        template <typename T> inline bool operator == (T x, T y){return Abs(x-y)<eps;}
7
8
        int sqn(double x){
9
            if (Abs(x)<eps) return 0;</pre>
            if (x>0) return 1;
10
            else return -1;
11
12
        }
13
   }
```

5.1.2 Point Definition

```
1
   namespace CG{
2
       struct Point{
3
            double x,y;
4
           Point(double x=0, double y=0):x(x),y(y){}
5
6
       Vector operator + (const Vector a,const Vector b){return Vector(a.x+b.x,a.y+b.y);}
       Vector operator - (const Vector a,const Vector b){return Vector(a.x-b.x,a.y-b.y);}
7
       Vector operator * (const Vector a,const double k){return Vector(a.x*k,a.y*k);}
8
       Vector operator / (const Vector a,const double k){return Vector(a.x/k,a.y/k);}
9
10
       bool operator < (const Vector a,const Vector b) {return a.x==b.x?a.y<b.y:a.x<b.x;}</pre>
11
       bool operator == (const Vector a,const Vector b) {return a.x==b.x && a.y==b.y;}
12
       double Dot(const Vector a,const Vector b){return a.x*b.x+a.y*b.y;}
       double Cross(const Vector a,const Vector b){return a.x*b.y-a.y*b.x;}
13
       double mult_Cross(const Vector a,const Vector b,const Vector c){return (a.x-c.x)*(b.
14
       y-c.y)-(b.x-c.x)*(a.y-c.y);}
       double mult_Dot(const Vector a,const Vector b,const Vector c){return (a.x-c.x)*(b.x-
15
       c.x)+(a.y-c.y)*(b.y-c.y);}
       double Norm(const Vector a){return sqrt(Dot(a,a));}
16
       double Angle(const Vector a,const Vector b){return acos(Dot(a,b)/Norm(a)/Norm(b));}
17
18
       Vector Rotate(const Vector a, const double theta){return Vector(a.x*cos(theta)-a.y*
       sin(theta),a.x*sin(theta)+a.y*cos(theta));}
       bool ToLeftTest(const Vector a,const Vector b){return Cross(a,b)<0;}</pre>
19
       double DisPP(const Vector a,const Vector b){return sqrt((a.x-b.x)*(a.x-b.x)+(a.y-b.y
20
       )*(a.y-b.y));}
21
   }
```

5.1.3 Line Definition

```
6
            // Line(Point \_p0=0, Point \_v=0, double \_t=1): p0(\_p0), p1(\_v)\{v=(p1-p0)/t; theta=1\}
       atan2(v.y,v.x);
7
       };
8
       bool operator < (const Line n,const Line m) {return n.theta<m.theta;}</pre>
       Point GetIntersection(const Line n,const Line m){return n.p0+n.v*Cross(m.v,(n.p0-m.
9
       p0))/Cross(n.v,m.v);}
       bool OnLine(const Vector a, const Line 1){return Cross(1.p0-a,1.p1-a)==0;}
10
       bool OnSegment(const Point a,const Line 1){return sgn(Cross(l.p0-a,l.p1-a))==0 &&
11
       sgn(Dot(l.p0-a,l.p1-a))<0;}
       double DisPL(const Point a,const Line 1){return Abs(Cross(1.p1-1.p0,a-1.p0)/Norm(1.
12
       p1-l.p0));}
       double DisPS(const Point a,const Line 1){
13
            if (l.p0==l.p1) return Norm(a-l.p0);
14
            Vector v1=l.p1-l.p0,v2=a-l.p0,v3=a-l.p1;
15
            if (sgn(Dot(v1,v2))<0) return Norm(v2);</pre>
16
            if (sgn(Dot(v1,v3))>0) return Norm(v3);
17
            return DisPL(a,1);
18
19
20
        Point GetProjection(const Point a, const Line 1){
21
            Vector v=l.p1-l.p0;
            return 1.p0+v*(Dot(v,a-1.p0)/Dot(v,v));
22
23
        bool SegmentIntersection(const Line n,const Line m,bool p){
24
            double c1=Cross(n.p1-n.p0,m.p1-m.p0);
25
            double c2=Cross(n.p1-n.p0,m.p1-n.p0);
26
27
            double c3=Cross(m.p1-m.p0,n.p0-m.p0);
28
            double c4=Cross(m.p1-m.p0,n.p1-m.p0);
            if (p){
29
                if (!sgn(c1) || !sgn(c2) || !sgn(c3) || !sgn(c4)){
30
                     return OnSegment(n.p0,m) | OnSegment(n.p1,m) | OnSegment(m.p0,n) |
31
       OnSegment(m.p0,m);
32
                }
33
            }
34
35
            return (sgn(c1)*sgn(c2)<0 && sgn(c3)*sgn(c4)<0);</pre>
36
       }
37
   }
```

5.1.4 Get Area

```
namespace CG{
    double GetArea(Point *p,int n){
        double area=Cross(p[n],p[1]);
        for (int i=2;i<=n;i++) area+=0.5*Cross(p[i-1],p[i]);
        return Abs(area);
    }
}</pre>
```

5.1.5 Get Circumference

5.1.6 Anticlockwise Sort

```
1
   namespace CG{
2
       \\p为一个凸包,只是不知其点集是否为逆时针
       void clockwise_sort(Point *p,int n){
3
            for(int i=0;i<n-2;i++){</pre>
4
                double tmp = mult_Cross(p[i+1],p[i+2],p[i]);
5
                if(tmp>0) return;
6
                else if(tmp<0){</pre>
7
8
                    reverse(p,p+n);
9
                    return;
10
                }
11
           }
12
       }
13
   }
```

5.2 Convex Hull

5.2.1 Get Convex Hull

```
namespace CG{
1
2
        Point p[MAXN],s[MAXN];
        int ConvexHull(Point *p,int n,Point *s){
3
4
            sort(p,p+n,cmp); //x从小到大,y从小到大;
5
            int m=0;
            for (int i=0;i<n;i++){</pre>
6
                 for (;m>=2 && Cross(s[m-1]-s[m-2],p[i]-s[m-1])<=0;m--);</pre>
7
8
                s[++m]=p[i];
            }
9
            int k=m;
10
            for (int i=n-2;i;i--){
11
                 for (;m>=k+1 && Cross(s[m-1]-s[m-2],p[i]-s[m-1])<=0;m--);</pre>
12
                s[++m]=p[i];
13
14
15
            return m-1;
16
        }
17
   }
```

5.2.2 Point in Convex Hull

```
namespace CG{
1
2
        bool PointInConvexHull(Point A){
            int l=1,r=tot-2,mid;
3
            while(l<=r){</pre>
4
5
                 mid=(l+r)>>1;
                 double a1=Cross(p[mid]-p[0],A-p[0]);
6
7
                 double a2=Cross(p[mid+1]-p[0],A-p[0]);
                 if(a1>=0 \&\& a2<=0){
8
                     if(Cross(p[mid+1]-p[mid],A-p[mid])>=0) return true;
9
10
                     return false;
11
                 else if(a1<0) r=mid-1;</pre>
12
13
                 else l=mid+1;
14
            return false;
15
16
        }
17
   }
```

5.3 Minkowski Sum

```
1
    namespace CG{
2
        void Minkowski(Point *C1,int n,Point *C2,int m){
            for(int i=1;i<=n;i++) s1[i]=C1[i]-C1[i-1];</pre>
3
            for(int i=1;i<=m;i++) s2[i]=C2[i]-C2[i-1];</pre>
4
5
            A[tot=1]=C1[1]+C2[1];
6
            int p1=1,p2=1;
            while (p1<=n && p2<=m) ++tot,A[tot]=A[tot-1]+(s1[p1]*s2[p2]>=0?s1[p1++]:s2[p2
7
        ++]);
8
            while (p1<=n) ++tot,A[tot]=A[tot-1]+s1[p1++];</pre>
9
            while (p2<=m) ++tot,A[tot]=A[tot-1]+s2[p2++];</pre>
10
            tot=ConvexHull(A,tot);
11
        }
12
    }
```

5.4 Rotating Calipers

5.4.1 The Diameter of Convex Hull

```
namespace CG{
1
2
        double RotatingCalipers(Point *p,int n){
3
            double dis=0;
4
            for(int i=0, j=2; i<n;++i){</pre>
                 while (abs(Cross(p[i+1]-p[i],p[j]-p[i]))<abs(Cross(p[i+1]-p[i],p[j+1]-p[i]))</pre>
5
        ) j=(j+1)%n;
                 dis=max(dis,max(DisPP(p[j],p[i]),DisPP(p[j],p[i+1])));
6
7
8
            return dis;
9
        }
10
   }
```

5.4.2 The Min Distance Bewteen two Convex Hull

```
1
   namespace CG{
2
       ///点c到线段ab的最短距离
      double GetDist(Point a,Point b,Point c){
3
          if(dis(a,b) < esp) return dis(b,c); ///a,b是同一个点
4
          if(mult_Dot(b,c,a)<-esp) return dis(a,c); ///投影
5
6
          if(mult_Dot(a,c,b)<-esp) return dis(b,c);</pre>
7
          return fabs(mult_Cross(b,c,a)/dis(a,b));
8
      }
9
       ///求一条线段ab的两端点到另外一条线段bc的距离,反过来一样,共4种情况
10
      double MinDist(Point a, Point b, Point c, Point d){
11
12
          return min(min(GetDist(a,b,c),GetDist(a,b,d)),min(GetDist(c,d,a),GetDist(c,d,b))
      );
      }
13
      double RotatingCalipers(Point *p,int n,Point *q,int m){
14
          int yminP = 0,ymaxQ=0;
15
          for(int i=1; i< n; i++){ ///找到点集p组成的凸包的左下角
16
              17
18
19
          for(int i=1;i<m;i++){ ///找到点集q组成的凸包的右上角
              if(q[i].y>q[ymaxQ].y||(q[i].y==q[ymaxQ].y)&&(q[i].x>q[ymaxQ].x))              ymaxQ = i;
20
21
          double ans = DisPP(p[yminP],q[ymaxQ]); ///距离(yminP,ymaxQ)维护为当前最小值。
22
```

```
for(int i=0;i<n;i++){</pre>
23
                 double tmp;
24
                 while(tmp=(mult_Cross(q[ymaxQ+1],p[yminP],p[yminP+1])-mult_Cross(q[ymaxQ],p[
25
       yminP],p[yminP+1]))>esp)
                     ymaxQ = (ymaxQ+1)%m;
26
                 if(tmp<-esp) ans = min(ans,GetDist(p[yminP],p[yminP+1],q[ymax0]));</pre>
27
                 else ans=min(ans,MinDist(p[yminP],p[yminP+1],q[ymaxQ],q[ymaxQ+1]));
28
29
                yminP = (yminP+1)%n;
30
            return ans;
31
        }
32
   }
33
```

5.5 Half Plane Intersection

```
1
   namespace CG{
        void HalfPlaneIntersection(Line 1[],int n){
2
3
            deque <Point> p;
4
            sort(l+1,l+1+n);
5
            deque <Line> q;
            q.push_back(l[1]);
6
            for (int i=2;i<=n;i++){</pre>
7
                for (;!p.empty() && !ToLeftTest(p.back()-l[i].p0,l[i].v);q.pop_back(),p.
8
       pop_back());
                for (;!p.empty() && !ToLeftTest(p.front()-l[i].p0,l[i].v);q.pop_front(),p.
9
       pop_front());
                if (sqn(Cross(l[i].v,q.back().v))==0)
10
                    if (ToLeftTest(l[i].p0-q.back().p0),q.back().v){
11
12
                        q.pop_back();
13
                        if (!p.empty()) p.pop_back();
14
                if (!q.empty()) p.push_back(GetIntersection(q.back(),l[i]));
15
                q.push_back(l[i]);
16
17
            for (;!p.empty() && !ToLeftTest(p.back()-q.front().p0,q.front().v);q.pop_back(),
18
       p.pop_back());
            p.push_back(GetIntersection(q.back(),q.front()));
19
            if (p.size() < 3) printf("0\n");
20
21
                cerr << "!" << endl;
22
                double area = 0.5 * Cross(p.back(), p.front());
23
                Point last = p.front();
24
                for (p.pop_front(); !p.empty(); last = p.front(), p.pop_front())
25
                    area += 0.5 * Cross(last, p.front());
26
                printf("%.8lf\n", fabs(area));
27
            }
28
29
       }
30
   }
```

5.6 Min Circle Cover

```
namespace CG{
    Point GetCircleCenter(const Point a,const Point b,const Point c){
    Point p=(a+b)/2.0,q=(a+c)/2.0;
    Vector v=Rotate(b-a,pi/2.0),w=Rotate(c-a,pi/2.0);
    if (sgn(Norm(Cross(v,w)))==0){
        if (sgn(Norm(a-b)+Norm(b-c)-Norm(a-c))==0) return (a+c)/2;
        if (sgn(Norm(b-a)+Norm(a-c)-Norm(b-c))==0) return (b+c)/2;
}
```

```
if (sgn(Norm(a-c)+Norm(c-b)-Norm(a-b))==0) return (a+c)/2;
8
9
            }
10
            return GetIntersection(Line(p,v),Line(q,w));
11
        void MinCircleCover(Point p[],int n){
12
             random_shuffle(p+1,p+1+n);
13
            Point c=p[1];
14
             double r=0;
15
             for (int i=2;i<=n;i++)</pre>
16
                 if (sgn(Norm(c-p[i])-r)>0){
17
                      c=p[i], r=0;
18
                      for (int j=1;j<i;j++)</pre>
19
20
                          if (sgn(Norm(c-p[j])-r)>0){
                              c=(p[i]+p[j])/2.0;
21
                              r=Norm(c-p[i]);
22
                              for (int k=1;k<j;k++)</pre>
23
                                   if (sgn(Norm(c-p[k])-r)>0){
24
25
                                       c=GetCircleCenter(p[i],p[j],p[k]);
26
                                       r=Norm(c-p[i]);
                                   }
27
                          }
28
29
            printf("%.10f\n%.10f %.10f",r,c.x,c.y);
30
31
        }
32
    }
```

5.7 Circle Union Area

```
//k次覆盖
1
   //圆并去重后s[0]
3 typedef pair<double, int> P;
4 const double pi = acos(-1.0);
5 const int MAXN = 10003;
6 P arc[MAXN << 1];</pre>
7
   int acnt, cnt;
8
   double s[1003];
   bool del[1003];
9
   void add(double st, double en) {
10
        if(st < -pi) {
   add(st + 2 * pi, pi);</pre>
11
12
13
            add(-pi, en);
            return;
14
15
        if(en > pi) {
16
            add(st, pi);
17
            add(-pi, en - 2 * pi);
18
19
            return;
20
        arc[++acnt] = P(st, 1);
21
22
        arc[++acnt] = P(en, -1);
23
   double F(double x) {
24
25
        return (x - \sin(x)) / 2;
26
   }
27
   struct Node {
        int x, y, r;
28
        Node(int _x = 0, int _y = 0, int _r = 0):x(_x), y(_y), r(_r) {}
29
        bool operator == (const Node& t) {
30
            return x == t.x & y == t.y & r == t.r;
31
```

```
32
        inline void read() {
33
34
            scanf("%d%d%d", &x, &y, &r);
35
   }a[1003];
36
37
    int main() {
        int n;
scanf("%d", &n);
38
39
        for(int i = 1; i <= n; i++) a[i].read();</pre>
40
41
        //去重
42
        int nn = 0;
43
        for(int \ i = 1; \ i <= n; \ i++) 
44
            bool\ same = 0;
45
46
            for(int \ j = 1; \ j < i; \ j++) \ \{
                 if(a[i] == a[j]) {
47
                     same = 1; break;
48
49
50
51
            if(!same) \ a/++nn/ = a/i/;
52
53
        n = nn;
        //去包含
54
        for(int \ i = 1; \ i \ll n; \ i++)
55
            for(int j = 1; j \le n; j++) if(i != j) 
56
57
                 if(hypot(a[i].x - a[j].x, a[i].y - a[j].y) < (double)(a[i].r - a[j].r)) \ del[i]
        j = 1;
58
59
60
        nn = 0;
61
        for(int \ i = 1; \ i \le n; \ i++) \ if(!del[i]) 
62
            a/++nn/ = a/i/;
63
64
        n = nn;
65
        for(int i = 1; i <= n; i++) {
66
67
            acnt = 0;
            for(int j = 1; j <= n; j++) if(i != j) {
68
                 int dis = (a[i].x - a[j].x) * (a[i].x - a[j].x) + (a[i].y - a[j].y) * (a[i].
69
       y - a[j].y);
70
                 if(a[j].r > a[i].r & dis <= (a[j].r - a[i].r) * (a[j].r - a[i].r)) add(-pi,
        pi);
                 else if(dis > (a[i].r - a[j].r) * (a[i].r - a[j].r) && dis < (a[i].r + a[j].
71
        r) * (a[i].r + a[j].r)){
72
                     double c = sqrt(dis);
                     double angle = a\cos((a[i].r * a[i].r + c * c - a[j].r * a[j].r) / (2 * a
73
        [i].r * c));
74
                     double k = atan2(a[j].y - a[i].y, a[j].x - a[i].x);
                     add(k - angle, k + angle);
75
76
77
            arc[++acnt] = P(pi, -1);
78
            sort(arc + 1, arc + acnt + 1);
79
            cnt = 0;
80
            double last = -pi;
81
            for(int j = 1; j <= acnt; j++) {</pre>
82
                 s[cnt] += F(arc[j].first - last) * a[i].r * a[i].r; //扇形 - 三角形
83
                 double xa = a[i].x + a[i].r * cos(last);
84
                 double ya = a[i].y + a[i].r * sin(last);
85
86
                last = arc[j].first;
                 double xb = a[i].x + a[i].r * cos(last);
87
```

```
double yb = a[i].y + a[i].r * sin(last);
88
                  s[cnt] += (xa * yb - xb * ya) / 2; //到圆心的三角形面积
89
90
                  cnt += arc[j].second;
91
             }
92
        }
         //printf("\%.3f | n", s[0]);
93
         for (int i = 0; i < n; i++) {
    printf("[%d] = %.3f\n", i + 1, s[i] - s[i + 1]);</pre>
94
95
96
97
         return 0;
    }
98
```

5.8 Simpson Integrate

```
double Simpson(double l,double r){
1
2
       return (r-1)*(F(1)+4*F((1+r)/2)+F(r))/6;
3
4
   double Integrate(double l,double r,double S){
       double mid=(l+r)/2;
5
6
       double A=Simpson(l,mid);
       double B=Simpson(mid,r);
7
8
       if(A+B-S<eps)return S;</pre>
       return Integrate(l,mid,A)+Integrate(mid,r,B);
9
   }
10
```

6 Conclusion

6.1 Game Theory

6.1.1 Bash's Game

Bash's Game 巴什博弈

有一堆个数为 n 的石子, 游戏双方依次从中拿取, 满足:

1. 每次至少取 1 个, 最多取 m 个.

最后取光者得胜。

结论: n = t(m+1) + r, 必败态:r = 0;

巴什博弈变种:

取一个指定集合的石头个数

取到最后一个石子输, n = t(m + 1) + r, r = 1;

6.1.2 Wythoff's Game

Wythoff's Game (威佐夫博弈)

有两堆分别为 (an, bn) 的石子, 游戏双方依次从中拿取, 满足:

1. 从任意一堆中取任意个 > 1。2. 从两堆中取同样多个。最后取完者胜.

结论: 对于任意的局势 (a, b)(a < b), 必败点为 (b-a)*(sqrt(5)+1)/2=a.

6.1.3 Fibonacci's Game / Zeckendorf's theory

Fibonacci's Game (斐波那契博弈)

有一堆个数为 n 的石子, 游戏双方轮流取石子, 满足:

- 1. 先手不能在第一次把所有的石子取完;
- 2. 之后每次可以取的石子数介于 1 到对手刚取的石子数的 2 倍之间(包含 1 和对手刚取的石子数的 2 倍)。 结论: 必败点是斐波那契数

齐肯多夫定理: 任何正整数可以表示为若干个不连续的 Fibonacci 数之和

6.1.4 Nim's Game / Anti-Nim's Game / K-Nim's Game / Anti-K-Nim's Game

Nim's Game (尼姆博弈)

石子的个数可以等价成某个游戏的 SG 函数。

有 n 堆石子, 游戏双方依次从中拿取, 满足:

1. 规定每次只能从一堆中取若干根,可将一堆全取走,但不可不取.

最后取完者为胜。

结论:

T态: 所有火柴数异或和为 0

S 态: 所有火柴数异或和不为 0

必胜态:S

有 n 堆石子, 游戏双方依次从中拿取, 满足:

1. 规定每次只能从一堆中取若干根,可将一堆全取走,但不可不取.

最后取完者为败。

结论:

S0 态: 即仅有奇数个孤单堆

T0 态: 即仅有偶数个孤单堆

S1 态: 异或和大于 0, 且有 1 个充裕堆

T1 态: 不存在

S2 态: 异或和大于 0, 且有多个充裕堆

T2 态: 异或和等于 0, 且有多个充裕堆

必胜态:T0,S1,S2

必败态:S0,T2

有 n 堆石子, 游戏双方依次从中拿取, 满足:

1. 规定每次只能至多 k 堆中取若干根, 可将 k 堆全取走, 但不可不取.

最后取完者为胜。

结论:

对于每一堆,把它石子的个数用二进制表示

必败态: 对所有的石子堆, 如果在任何一个二进制位上 1 的个数总是 k+1 的整数倍

有 n 堆石子, 游戏双方依次从中拿取, 满足:

1. 规定每次只能至多 k 堆中取若干根, 可将 k 堆全取走, 但不可不取

最后取完者为败。

结论:

- 1. 对于每一堆, 把它石子的个数用二进制表示
- 2. 所有的堆(非零堆,下同)全是 1,此时如果 1 堆个数模 k+1 的结果是 1 则必败,否则必胜(我们可以通过 拿走 0 到 k 个堆来随意调整当前状态模的结果,然后再将所有大于 1 的堆降到 1 就行了)
- 3. 有多于 k 个堆的个数大于 1。必胜

6.1.5 阶梯博弈

有 n 个阶梯呈升序排列,每个阶梯上有若干个石子,游戏双方轮流取石子,满足:

1. 将一个阶梯上的石子移任意个(>0)到前一个台阶。

当没有可行操作时(所有石子都被移动到了地面,即第0号台阶)输。

结论:

奇数号台阶的 Nim 游戏

变种 1: 树上, 每个石子只能往父亲节点移动.

变种 2:

游戏双方在一个 1*N 的格子内挪动棋子,刚开始在若干个位置上有棋子,每个位置至多一个棋子

每一个选手可以进行的操作时选择一个棋子并把它向左方移动,当然不能越过其它的棋子,也不能超出边界。 谁不能移动谁就输了。求谁会赢?

结论:

将棋子位置按升序排列, 然后从后往前两两绑成一对, 如果个数是奇数, 那么将第一个和边界外绑定.

一对棋子的前一个和前一对棋子的后一个之间有多少个空位置对最终的结果是没有影响的。

于是我们只需要考虑同一对的两个棋子之间有多少空位,将同一对棋子间的空位视为石子,做 nim 游戏两对棋子间的空格数当奇数位石子,其他当偶数位石子,石子相右边移动

变种 3:

山上有 n 个人,每个人给出距离山顶的距离,给出其中一个人为 king,每次能挑选一个人向上移动,不能越过其他人,最后将 king 移动到山顶者获胜。问获胜者。

结论:

只要把 King 当作普通人一样处理即可。除了两种特殊情况:

- 1. 当 King 是第一个人时, Alice 直接胜
- 2. 当 King 是第二个人且一共有奇数个人时,第一堆的大小需要减 1。

6.1.6 Multi-Nim

有 n 堆石子, 游戏双方依次从中拿取, 满足:

- 1. 任意一堆石子中拿任意多个石子 (不能不拿)
- 2. 把一堆数量不少于 2 石子分为两堆不为空的石子 最后取完者为胜。

结论:

操作一与普通的 Nim 游戏等价

操作二实际上是将一个游戏分解为两个游戏,根据 SG 定理,我们可以通过异或运算把两个游戏连接到一起, 作为一个后继状态

$$SG(x) \equiv \begin{cases} x - 1 & (x \mod 4 = 0) \\ x & (x \mod 4 = 1 \text{ or } 2) \\ x + 1 & (x \mod 4 = 3) \end{cases}$$
 (1)

Multi-SG 游戏规定,在符合拓扑原则的前提下,一个单一游戏的后继可以为多个单一游戏。

Multi-SG 其他规则与 SG 游戏相同。

注意在这里要分清楚后继与多个单一游戏

对于一个状态来说,不同的划分方法会产生多个不同的后继,而在一个后继中可能含有多个独立的游戏

一个后继状态的 SG 值即为后继状态中独立游戏的异或和

该状态的 SG 值即为后继状态的 SG 值中未出现过的最小值

6.1.7 Every-SG

给定一张无向图,上面有一些棋子,两个顶尖聪明的人在做游戏,每人每次必须将可以移动的棋子进行移动, 不能移动的人

因为两个人都顶尖聪明,因此当一个人知道某一个游戏一定会输的话,它一定会尽力缩短游戏的时间,当它知 道某一个游戏一定会赢的话,一定会尽力延长游戏的时间

对于还没有结束的单一游戏,游戏者必须对该游戏进行一步决策;

其他规则与普通 SG 游戏相同

Every-SG 游戏与普通 SG 游戏最大的不同就是它多了一维时间 对于 SG 值为 0 的点,我们需要知道最少需要多少步才能走到结束 对于 SG 值不为 0 的点,我们需要知道最多需要多少步结束 这样我们用 step 变量来记录这个步数

$$step(x) \equiv \begin{cases} 0 & u \\ maxstep(v) & sg(u)! = 0vusg(v) = 0 \\ minstep(v) & sg(u) = 0vu \end{cases}$$
 (2)

6.1.8 树的删边游戏

给出一个有 N 个点的树,有一个点作为树的根节点。游戏者轮流从树中删去边,删去一条边后,不与根节点相连的部分将被移走。谁无法移动谁输。

结论:

Colon Principle: 对于树上的某一个点, ta 的分支可以转化成以这个点为根的一根竹子, 这个竹子的长度就是 ta 各个分支的边的数量的异或和

叶子节点的 SG 值为 0; 中间节点的 SG 值为它的所有子节点的 SG 值加 1 后的异或和。

6.1.9 Chomp's theory?

取一个无关紧要的位置, 如果对方必胜, 则学习其策略, 我方必胜.

6.1.10 Other's theory?

有 n 堆石子, 游戏双方依次从中拿取, 满足:

1. 规定每次能从任意多堆中取 1 根, 不可不取.

最后取完者为胜。

结论: 如果全是偶数, 先手必败, 否者先手必胜

一个无相联通图,有一个点作为图的根。

游戏者轮流从图中删去边,删去一条边后,不与根节点相连的部分将被移走。

谁无路可走谁输。

结论:

Fusion Principle: 环上的点可以融合,且不改变图的 SG 值,我们可以把一个带有奇数边的环等价成只有一个端点的一条边而偶数边的环等价于一个点

6.1.11 SG Theory

```
memset(mex, 0, sizeof mex);
for (int i = 1; i < maxN;++i) {
    for (int j = 1; j <= n;++j) {
        if (a[j] <= i)
            mex[SG[i - a[j]]] = i;
        for (int k = 1; i - k - a[j] > 0;++k)
            mex[SG[k] ^ SG[i - k - a[j]]] = 1;
}
```

6.1.12 SJ Theory

反公平游戏 Anti-SG Game

DAG 上没有出度的点为胜利状态,其它定义与一般游戏相同。

现在的问题是解决多个反公平游戏的合并。

SJ 定理说明: 先手必胜, 当且仅当以下两个条件同时成立或同时不成立:

- 1. 合并的 SG 值为 0;
- 2. 所有游戏的 SG 值不超过 1。

6.1.13 Surreal Number Theory

6.2 Math Theory

6.2.1 Euler's Theorem

$$a^b \equiv \begin{cases} a^{b\%\varphi(p)} & \gcd(a,p) = 1 \\ a^b & \gcd(a,p) \neq 1, b < \varphi(p) \\ a^{b\%\varphi(p) + \varphi(p)} & \gcd(a,p) \neq 1, b \geq \varphi(p) \end{cases} \pmod{p}$$

6.2.2 Möbius Inversion

Dirichlet Convolution is $(f \times g)(N) = \sum_{d|N} f(d) * g(\frac{N}{d})$

Theorem:

$$\begin{cases} f = g \times 1 \\ g = f \times \mu \end{cases}$$

$$\begin{cases} id(n) = \sum_{d|n} \varphi(d) \\ e(n) = \sum_{d|n} \mu(d) \end{cases}$$
(3)

$$\begin{cases}
\sum_{i}^{n} \sum_{j}^{m} gcd(i,j) = \sum_{d}^{\max(n,m)} \varphi(d) * \lfloor \frac{n}{d} \rfloor \lfloor \frac{m}{d} \rfloor \\
\sum_{i}^{n} \sum_{j}^{m} e(gcd(i,j)) = \sum_{d}^{\min(n,m)} \mu(d) * \lfloor \frac{n}{d} \rfloor \lfloor \frac{m}{d} \rfloor \\
\sum_{i=1}^{n} |\mu(i)| = \sum_{i=1}^{\lfloor \sqrt{n} \rfloor} \mu(i) * \lfloor \frac{n}{i*i} \rfloor
\end{cases} \tag{4}$$

$$\begin{cases}
sum(x,y) = \sum_{i}^{x} \sum_{j}^{y} i * j = \frac{x * (x+1)}{2} * \frac{y * (y+1)}{2} \\
F(x,y) = \sum_{i=1}^{\min(x,y)} i^{2} * \mu(i) * sum(\lfloor \frac{x}{i} \rfloor, \lfloor \frac{y}{i} \rfloor) \\
\sum_{i}^{n} \sum_{j}^{m} lcm(i,j) = \sum_{i=1}^{\min(n,m)} d * F(\lfloor \frac{n}{i} \rfloor, \lfloor \frac{y}{i} \rfloor)
\end{cases} (5)$$

6.2.3 Sieve Tips

$$\varphi(nm) = \varphi(n) \cdot \varphi(m) \cdot \frac{\gcd(n,m)}{\varphi(\gcd(n,m))} \tag{6}$$

$$\varphi(n) = \sum_{i=1}^{n} [(n,i) = 1] \cdot i = \frac{n * \varphi(n) + [n=1]}{2}$$
(7)

$$\begin{cases}
id = \varphi \times 1 \\
\frac{n \cdot (n+1)}{2} = \sum_{i=1}^{n} i = \sum_{i=1}^{n} \sum_{d|i} \varphi(d) = \sum_{\frac{i}{d}=1}^{n} \sum_{d=1}^{\lfloor \frac{n}{d} \rfloor} \varphi(d) = \sum_{i=1}^{n} \phi(\lfloor \frac{n}{i} \rfloor)
\end{cases}$$
(8)

$$\begin{cases}
e = \mu \times 1 \\
1 = \sum_{i=1}^{n} [i = 1] = \sum_{i=1}^{n} \sum_{d|i} \mu(d) = \sum_{i=1}^{n} \sum_{d=1}^{\lfloor \frac{n}{i} \rfloor} \mu(d) = \sum_{i=1}^{n} M(\lfloor \frac{n}{i} \rfloor)
\end{cases}$$
(9)

$$\begin{cases}
id^{2} = (id \cdot \varphi) \times id \\
\phi'(n) = \sum_{i=1}^{n} i \cdot \varphi(i) \\
\frac{n \cdot (n+1) \cdot (2n+1)}{6} = \sum_{i=1}^{n} i^{2} = \sum_{i=1}^{n} \sum_{d \mid i} d \cdot \varphi(d) \cdot \frac{i}{d} = \sum_{\frac{i}{d} = 1}^{n} \frac{i}{d} \sum_{d=1}^{\frac{n}{d}} d \cdot \varphi(d) = \sum_{i=1}^{n} i \cdot \phi'(\lfloor \frac{n}{i} \rfloor)
\end{cases}$$
(10)

6.2.4 Newton's method

Iff

 $f'(x) \neq 0 \&\& f''(x) continuous$

Then

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

6.2.5 Cantor Expansion

康托展开是一个全排列到一个自然数的双射,常用于构建哈希表时的空间压缩。康托展开的实质是计算当前排列在所有由小到大全排列中的顺序,因此是可逆的。以下称第 x 个全排列是都是指由小到大的顺序。

$$X = a_n(n-1)! + a_{n-1}(n-2)! + \cdots + a_1 \cdot 0!$$

其中 a_i 为排列的第 i 个元素值

例如,357412968展开为98884。因为 X=2*8!+3*7!+4*6!+2*5!+0*4!+0*3!+2*2!+0*1!+0*0!=98884.

解释:

排列的第一位是 3, 比 3 小的数有两个, 以这样的数开始的排列有 8! 个, 因此第一项为 2*8!

排列的第二位是 5,比 5 小的数有 1、2、3、4,由于 3 已经出现,因此共有 3 个比 5 小的数,这样的排列有 7! 个,因此第二项为 3*7!

以此类推,直至0*0!

用途:

显然, n 位 (0 n-1) 全排列后, 其康托展开唯一且最大约为 n!, 因此可以由更小的空间来储存这些排列。由公式可将 X 逆推出唯一的一个排列。

逆运算:

如 n=5,x=96 时:

首先用 96-1 得到 95, 说明 x 之前有 95 个排列.(将此数本身减去 1)

用 95 去除 4! 得到 3 余 23, 说明有 3 个数比第 1 位小, 所以第一位是 4.

用 23 去除 3! 得到 3 余 5, 说明有 3 个数比第 2 位小, 所以是 4, 但是 4 已出现过, 因此是 5.

用 5 去除 2! 得到 2 余 1, 类似地, 这一位是 3.

用 1 去除 1! 得到 1 余 0, 这一位是 2.

最后一位只能是 1.

所以这个数是 45321.

按以上方法可以得出通用的算法。

6.3 Convolution

6.3.1 FWT

$$\begin{cases} C_k = \sum_{i \oplus j = k} A_i * B_j \\ DWT(A)_i = \sum_{j}^n A_j * f_{i,j} \\ DWT(C)_i = DWT(A)_i * DWT(B)_i \\ f_{i,j} \cdot f_{i,k} = f_{i,j \oplus k} \\ f_{i,j} = [i \text{ and } j == i] \\ f_{i,j} = [i \text{ and } j == j] \\ f_{i,j} = [i \text{ and } j == j] \end{cases}$$
 (and)
$$f_{i,j} = [i \text{ and } j == j]$$
 (or)
$$f_{i,j} = (-1)^{|i|} \text{ and } j = (-1$$

6.4 Geometry

6.4.1 The Number of Ingeter Point on a Circle

Set r = const is the radius of the circle.

$$r^2 = p_1^{a_1} + p_2^{a_2} + \dots + p_m^{a_m} = \sum_{i=1}^m p_i^{a_i}$$

Define

$$\chi(n) = \begin{cases} 1 & n\%4 = 1 \\ -1 & n\%4 = 3 \\ 0 & n\%2 = 0 \end{cases}$$

By the way, $\chi(n)$ is a multiplicative function.

Define

$$\Gamma(p_i, a_i) = \sum_{j=0}^{a_i} \chi(p_i^j) = \begin{cases} 1 & p_i = 2 & || & (p_i\%4 = 3 & \&\& & a_i\%2 = 0) \\ 0 & p_i\%4 = 3 & \&\& & a_i\%2 = 1 \\ a_i + 1 & p_i\%4 = 1 \end{cases}$$

Define cnt is the number of integer point on circle

$$cnt(r) = 4 \prod_{i=1}^{m} \sum_{j=0}^{a_i} \chi(p_i^j) = 4 \prod_{i=1}^{m} \Gamma(p_i, a_i) = 4 \sum_{k|r^2} \chi(k)$$

Define CNT is the number of integer point in circle

$$CNT(r) = 1 + \sum_{i=1}^{r^2} cnt(i) = 1 + \sum_{i=1}^{r^2} \lfloor \frac{r^2}{i} \rfloor \chi(i)$$

7 Others

7.1 Offline Algorithm

7.1.1 CDQ Divide and Conquer

```
struct Node {
2
        int x, y, z, ans;
        Node() {}
3
4
        Node(int _x, int _y, int _z):x(_x), y(_y), z(_z) {}
5
        bool operator < (const Node &b) const {</pre>
6
            if(y == b.y) {
7
                 if(z == b.z) return x < b.x;
8
                 return z < b.z;
9
10
            return y < b.y;</pre>
11
   }A[MAXN], B[MAXN], C[MAXN];
12
   int bit[MAXN];
13
   void add(int k, int v) {
14
15
        for(; k \le m; k += k \& -k) bit[k] = max(bit[k], v);
16
17
   void clear(int k) {
        for(; k \le m; k += k \& -k) bit[k] = 0;
18
19
   int sum(int k) {
20
        int res = 0;
21
        for(; k; k \rightarrow k - k) res = max(res, bit[k]);
22
23
        return res;
   }
24
25
   void solve(int l, int r) {
26
        if(l == r) {
27
            B[l] = A[l];
28
             return;
29
30
        int mid = (l + r) >> 1;
31
        solve(l, mid);
        for(int i = mid + 1; i <= r; i++) B[i] = A[i];</pre>
32
33
        //sort(B + l, B + mid + 1);
34
        sort(B + mid + 1, B + r + 1);
        int L = 1;
35
        for(int R = mid + 1; R <= r; R++) {</pre>
36
            while(L \leftarrow mid && B[L].y \leftarrow B[R].y) add(B[L].z, B[L].ans), L++;
37
38
            A[B[R].x].ans = max(A[B[R].x].ans, sum(B[R].z - 1) + 1);
            B[R].ans = A[B[R].x].ans;
39
40
        for(int i = l; i <= L; i++) clear(B[i].z);</pre>
41
        solve(mid + 1, r);
42
        L = 1;
43
        int p = 1, q = mid + 1;
44
        while(p <= mid || q <= r) {</pre>
45
            if(q > r | l (p \le mid \&\& B[p].y \le B[q].y)) C[L++] = B[p++];
46
            else C[L++] = B[q++];
47
48
        for(int i = 1; i <= r; i++) B[i] = C[i];</pre>
49
50
```

7.1.2 Mo's Algorithm

```
struct Node{
1
2
        int l, r, t, id;
        bool operator < (const Node& a) const {</pre>
3
4
            if(l /sz == a.l / sz) {
                 if(r == a.r) return t < a.t;</pre>
5
6
                 return r < a.r;</pre>
7
            return l / sz < a.l / sz;</pre>
8
9
   }q[MAXN];
10
   void solve() {
11
        while (t < q[i].t) addTime(t++, 1);
12
13
        while (t > q[i].t) addTime(--t, -1);
14
        while(L < q[i].l) add(L++, -1);
        while(L > q[i].l) add(--L, 1);
15
        while(R < q[i].r) add(++R, 1);
16
        while(R > q[i].r) add(R--, -1);
17
18
```

7.1.3 Mo's Algorithm On Tree

```
struct Edge {
1
2
        int to, nxt;
3
    }e[MAXN << 1];
    int head[MAXN], ecnt;
    int stack[MAXN], top, belong[MAXN], cnt, sz;
    struct Node {
         int l, r, id, ti;
7
        bool operator < (const Node &x) const {</pre>
8
             return belong[1] < belong[x.1] || (belong[1] == belong[x.1] && belong[r] <</pre>
9
        belong[x.r]) || (belong[l] == belong[x.l] && belong[r] == belong[x.r] && ti < x.ti);
10
    }q[MAXN];
11
    struct Node2 {
12
         int l, r, ti;
13
   }qq[MAXN];
14
   int n, m, Q, Q0, Q1;
15
   int V[MAXN], W[MAXN], C[MAXN];
   int fa[MAXN][S + 3], dep[MAXN];
17
   long long ans[MAXN], tans;
18
   int vis[MAXN], cur[MAXN];
19
   long long sum[MAXN];
20
21
   int l, r, tm;
22
    inline int read() {
        int x = 0; char ch = getchar(); bool fg = 0;
while(ch < '0' || ch > '9') { if(ch == '-') fg = 1; ch = getchar(); }
while(ch >= '0' && ch <= '9') { x = x * 10 + ch - '0'; ch = getchar(); }</pre>
23
24
25
         return fg ? -x : x;
26
27
    inline void add_edge(int u, int v) {
28
         e[++ecnt] = (Edge) \{v, head[u]\}; head[u] = ecnt;
29
         e[++ecnt] = (Edge) \{u, head[v]\}; head[v] = ecnt;
30
31
    void dfs(int u, int f) {
32
         fa[u][0] = f;
33
         dep[u] = dep[f] + 1;
34
         int bot = top;
```

```
36
        for(int i = head[u]; i; i = e[i].nxt) {
37
            int v = e[i].to;
            if(v == f) continue;
38
            dfs(v, u);
39
            if(top - bot >= sz) {
40
41
                while(top != bot) belong[stack[top--]] = cnt;
42
43
44
        stack[++top] = u;
45
46
    void G(int &u, int step) {
47
        for(int i = 0; i < S; i++) if((1 << i) & step) u = fa[u][i];
48
49
50
   int lca(int u, int v) {
        if(dep[u] > dep[v]) swap(u, v);
51
        G(v, dep[v] - dep[u]);
52
        if(u == v) return u;
53
        for(int i = S; i >= 0; i--) if(fa[u][i] != fa[v][i]) {
54
55
            u = fa[u][i]; v = fa[v][i];
56
        return fa[u][0];
57
58
   inline void modify(int u) {
59
        tans -= V[C[u]] * sum[cur[C[u]]];
60
61
        cur[C[u]] += vis[u];
        vis[u] = -vis[u];
62
63
        tans += V[C[u]] * sum[cur[C[u]]];
64
    inline void update(int u, int v) {
65
        if(u == v) return;
66
        if(dep[u] > dep[v]) swap(u, v);
67
        while(dep[v] > dep[u]) {
68
            modify(v);
69
70
            v = fa[v][0];
71
        while(u != v) {
72
            modify(u); modify(v);
73
74
            u = fa[u][0]; v = fa[v][0];
75
        }
76
   }
77
   inline void upd(int t) {
        if(vis[qq[t].l] == -1) {
78
79
            modify(qq[t].1);
80
            swap(C[qq[t].1], qq[t].r);
81
            modify(qq[t].l);
82
        else swap(C[qq[t].1], qq[t].r);
83
84
    inline void moveto(int u, int v) {
85
86
        update(l, u); update(r, v);
87
        l = u; r = v;
88
   int main() {
89
        n = read(); m = read(); Q = read();
90
        sz = (int)pow(n, 2.0 / 3.0);
91
        for(int i = 1; i <= m; i++) V[i] = read();</pre>
92
        for(int i = 1; i <= n; i++) W[i] = read();</pre>
93
        for(int i = 1, u, v; i < n; i++) {
94
            u = read(); v = read();
95
            add_edge(u, v);
```

```
97
         for(int i = 1; i <= n; i++) {</pre>
98
              C[i] = read();
99
100
              vis[i] = 1;
              sum[i] = sum[i - 1] + W[i];
101
102
         for(int i = 1, tp; i <= Q; i++) {
103
104
              tp = read();
              if(tp) {
105
                  ++Q1:
106
                  q[Q1].l = read(); q[Q1].r = read();
107
                  q[Q1].id = Q1;
108
                  q[Q1].ti = i;
109
110
              }
111
              else {
                  ++00;
112
                  qq[Q0].l = read(); qq[Q0].r = read();
113
                  qq[Q0].ti = i;
114
              }
115
116
         dfs(1, 0);
117
         while(top) belong[stack[top--]] = cnt;
118
         sort(q + 1, q + Q1 + 1);
119
         for(int k = 1; k <= S; k++) {
120
121
              for(int i = 1; i <= n; i++) {</pre>
122
                  fa[i][k] = fa[fa[i][k - 1]][k - 1];
123
124
         for(int i = 1; i <= Q1; i++) {
    if(belong[q[i].l] > belong[q[i].r]) swap(q[i].l, q[i].r);
125
126
127
              moveto(q[i].l, q[i].r);
              int lc = lca(l, r);
128
              modify(lc);
129
              while(qq[tm + 1].ti < q[i].ti && tm < Q0) upd(++tm);</pre>
130
              while(qq[tm].ti > q[i].ti) upd(tm--);
131
              ans[q[i].id] = tans;
132
133
              modify(lc);
         }
134
135
         for(int i = 1; i \leftarrow Q1; i++) printf("%lld\n", ans[i]);
136
         return 0;
137
```

7.2 Randomized Algorithm

7.2.1 Simulated Annealing

```
void solve() {
1
2
       while(T > eps) {
3
           double alpha = ((rand() % 30001) / 15000.0) * pi;
           double theta = ((rand() % 10001) / 10000.0) * pi;
4
           tmp.x = cur.x + T * sin(theta) * cos(alpha);
5
           tmp.y = cur.y + T * sin(theta) * sin(alpha);
6
           tmp.z = cur.z + T * cos(theta);
7
           tmp.dis = cal(tmp);
8
            if(tmp.dis < cur.dis || (tmp.dis * 0.999 < cur.dis && (rand() & 7) == 7)) cur =
9
       tmp;
10
            //if(exp((cur.d - tmp.d) / T) > ((double)rand() / RAND_MAX)) cur = tmp;
11
           T *= 0.999;
12
```

```
13 } 14 }
```

7.3 Other Method

7.3.1 Enumerate Subset

7.3.2 Enumerate $\lfloor \frac{n}{d} \rfloor \lfloor \frac{m}{d} \rfloor$

```
int cal(int n, int m) {
1
2
        if(n > m) swap(n, m);
3
        int res = 0, last;
        for(int i = 1; i <= n; i = last + 1) {</pre>
4
             last = min(n / (n / i), m / (m / i));
res += (n / i) * (m / i) * (sum(last) - sum(i - 1));
5
6
7
        }
8
        return res;
9
```

7.3.3 Find Primitive Root Modulo N

```
for i in range(1,mod):
    if 3 ** i % mod == 1:
        if i == mod - 1:
            print("yes")
            break
        print("no")
```

8 Samples

8.1 vimrc

```
1  set cindent
2  set number
3  set mouse=a
4  set tabstop=4
5  set shiftwidth=4
6  syntax on
7  inoremap { {}<left>
8  map <F9> :w<CR> :! g++ % -0 %< -Wall --std=c++14 -g && ./%< <CR>
9  "ios::sync_with_stdio(0); cin.tie(0); cout.precision(6); cout << fixed;</pre>
```

8.2 Check

```
while true; do
./data > in
./tmp < in > out
./std < in > ans
diff out ans
if [ $? -ne 0 ]; then exit; fi
echo Passed
done
```

8.3 FastIO

```
//普通情况
2
   namespace IO {
        const int MB = 1048576;
3
        const int RMAX = 16 * MB;
4
        const int WMAX = 16 * MB;
5
       #define getchar() *(rp++)
6
       #define putchar(x) (*(wp++) = (x))
7
8
        char rb[RMAX], *rp = rb, wb[WMAX], *wp = wb;
9
        inline void init() {
            fread(rb, sizeof(char), RMAX, stdin);
10
11
        template <class _T> inline void read(_T &_a) {
12
            _a = 0; bool _f = 0; int _c = getchar();
13
            while (_c < '0' | | _c > '9') _f | = _c = '-', _c = getchar();
14
            while (_c >= '0' \& _c <= '9') _a = _a * 10 + (_c ^ '0'), _c = getchar();
15
16
            _a = _f ? -_a : _a;
17
        template <class _T> inline void write(_T _a) {
18
19
            static char buf[20], *top = buf;
            if (_a) {
20
                while (_a) {
21
                     _{T} tm = _{a} / 10;
22
                    *(++top) = char(_a - tm * 10) | '0';
23
24
                    _a = tm;
25
26
                while (top != buf) putchar(*(top--));
27
28
            else putchar('0');
29
```

```
void output() {
30
            fwrite(wb, sizeof(char), wp - wb, stdout);
31
32
33
34
    //EOF结尾+分块读入
   #define likely(x) __builtin_expect(!!(x), 1)
35
   #define unlikely(x) __builtin_expect(!!(x), 0)
36
37
    namespace IO {
        const int MB = 1048576;
38
        const int RMAX = 4 * MB;
39
        const int WMAX = 4 * MB;
40
        unsigned long long filesize;
41
        #define putchar(x) (*(wp++) = (x))
42
        char rb[RMAX], wb[WMAX], *wp = wb;
43
44
        int rp = 0;
        inline void init() {
45
            filesize = fread(rb, sizeof(char), RMAX, stdin);
46
47
            wp = wb;
48
49
        void output() {
50
            fwrite(wb, sizeof(char), wp - wb, stdout);
51
52
        inline char getCHAR(){
53
            if(unlikely(rp == filesize)){
54
55
                fwrite(wb, sizeof(char), wp - wb, stdout);
56
                init();
57
                if(unlikely(filesize == 0)) {
58
                     //cerr << 1.0 * (clock() - st) / CLOCKS_PER_SEC << endl;
59
                     exit(0);
                }
60
            }
61
            return rb[rp++];
62
63
        template <class _T> inline void read(_T &_a) {
64
            _a = 0; static bool _f = 0; static int _c;
65
            _f = 0; _c = getCHAR();
66
            while (_c < '0' | | _c > '9') _f | = _c == '-', _c = getCHAR();
67
            while (_c >= '0' \&\& _c <= '9') _a = _a * 10 + (_c ^ '0'), _c = getCHAR();
68
69
            _a = _f ? -_a : _a;
70
71
        template <class _T> inline void write(_T _a) {
72
            static char buf[20], *top = buf;
            if (_a) {
73
                while (_a) {
74
                     _{T} tm = _{a} / 10;
75
                     *(++top) = char(_a - tm * 10) | '0';
76
77
                     _a = tm;
78
79
                while (top != buf) putchar(*(top--));
80
            else putchar('0');
81
82
            putchar(' \setminus n');
        }
83
   }
84
```

8.4 Java BigNum

```
1 import java.math.*;
```

```
import java.util.*;
2
   import java.lang.*;
3
4
5
   public class Main{
6
       public static void main(String []args){}
7
8
   //IO
   Scanner in = new Scanner(System.in);
9
10
   while(in.hasNext()){} //EOF
11
   public static void main(String argv[]) throws IOException{}
12
   StreamTokenizer cin = new StreamTokenizer(new BufferedReader(new InputStreamReader(
13
       System.in)));
14
   PrintWriter cout = new PrintWriter(new OutputStreamWriter(System.out));
   while(cin.nextToken() != StreamTokenizer.TT_EOF) ;//EOF
15
   cin.nextToken();int n = (int)cin.nval;String s = cin.sval;
16
   cout.println( Type );cout.flush();
17
   cin.ordinaryChar('/');
18
19
20 BufferedReader br = new BufferedReader(new InputStreamReader(System.in));
21
   br.ready()//EOF
   while ((valueString=bf.readLine())!=null);
22
23 br.close();
24
   //true\ fast-IO
25
   static class InputReader {
26
       public BufferedReader reader;
27
       public StringTokenizer tokenizer;
28
29
       public InputReader(InputStream stream) {
           reader = new BufferedReader(new InputStreamReader(stream), 32768);
30
           tokenizer = null;
31
       }
32
33
       public String next() {
34
           while (tokenizer == null || !tokenizer.hasMoreTokens()) {
35
36
               try {
                   tokenizer = new StringTokenizer(reader.readLine());
37
                 catch (IOException e) {
38
                    throw new RuntimeException(e);
39
40
41
           }
42
           return tokenizer.nextToken();
       }
43
44
       public int nextInt() {
45
           return Integer.parseInt(next());
46
47
48
49
   //类 Number
50
51
   //double Value ()
52
   //int Value()
53
   //long Value()
   //shortValue()
54
   //类 BigDecimal
55
   //ROUND_CEILING 接近正无穷大的舍入模式。
56
   //ROUND_FLOOR 接近负无穷大的舍入模式。
57
   //ROUND_DOWN 接近零的舍入模式
58
   //ROUND HALF UP 四舍五入 >=0.5向上舍入
59
   //ROUND_HALF_DOWN 四舍五入 >0.5向上舍入
60
   //BigDecimal(BigInteger\ val)
```

```
62
    //BigDecimal(BigInteger\ unscaledVal\,,\ int\ scale)
63
    //BigDecimal(char[] in, int offset, int len, MathContext mc)
64
   //BigDecimal(double val, MathContext mc)不建议
    //BigDecimal(int val, MathContext mc)
65
66
    //BigDecimal(long\ val\ ,\ MathContext\ mc)
67
    //BigDecimal(String val, MathContext mc)
    //abs()
68
    //add(BigDecimal augend, MathContext mc)
69
70
    //compareTo(BigDecimal\ val)
    //divide\ (BigDecimal\ divisor\ , MathContext\ mc)
71
    //divide \, To Integral Value \, (Big Decimal \ divisor \,, \ Math Context \ mc)
72
73
    //max(BigDecimal\ val)
    //min(BigDecimal\ val)
74
75
    //multiply(BigDecimal multiplicand, MathContext mc)
76
    //negate() 其值为 (-this), 其标度为 this.scale()
77
    //pow(int n)
    //remainder(BigDecimal divisor) 返回其值为 (this % divisor) 的 BigDecimal
78
    //round(MathContext mc) 返回根据 MathContext 设置进行舍入后的 BigDecimal。
79
    //caleByPowerOfTen(int n) 返回其数值等于 (this * 10^n) 的 BigDecimal。
80
81
    //subtract(BigDecimal subtrahend, MathContext mc)
82
   //setScale (int newScale, RoundingMode roundingMode)
83
    //toString()
    //ulp()返回此 BigDecimal 的 ulp (最后一位的单位) 的大小
84
    //String \ s = b. strip Trailing Zeros(). to Plain String(); 让 big decimal 不用科学计数法显示
85
86
    //类 BigInteger
87
    //parseInt
88
    //BigInteger\ zero = BigInteger.valueOf(0);
89
    //BigInteger \ a = in.nextBigInteger();
90
    //abs()
    //and(BigInteger val) 返回其值为 (this & val)
91
    //or(BigInteger val) 返回其值为 (this | val)
92
93
    //andNot(BigInteger val) 返回其值为 (this & ~val)
    //compareTo(BigInteger val)
94
    //add(BigInteger\ val)
95
96
    //divide(BigInteger\ val)
    //BigInteger[] divideAndRemainder(BigInteger val) 返回包含 (this / val) 后跟 (this %
97
        val) 的两个 BigInteger 的数组。
    //equals(Object x)
98
    //gcd(BigInteger val)
99
100
    //isProbablePrime(int certainty) e.g. a.isProbablePrime(4)
101
    //max(BigInteger\ val)\ min(BigInteger\ val)
102
    //mod(BigInteger m)
    //modInverse(BigInteger m) 返回其值为 (this ~-1 mod m)
103
104
    //modPow(BigInteger exponent, BigInteger m) 返回其值为 (this exponent mod m)
    //multiply(BigInteger\ val)
105
    //not() 返回其值为 (~this)
106
107
    //shiftLeft(int n) 返回其值为 (this << n)
108
    //shiftRight(int n) 返回其值为 (this >> n)
109
    //toString()
110
    //valueOf(long val)
111
    //xor(BigInteger val) 返回其值为 (this ^ val)
112
    //other
113
    //Arrays.sort(array);
```

$8.5 ext{ pb_ds}$

```
1 //可持久化平衡树,不过时间和空间都不太行
2 #include <ext/rope>
3 using namespace __gnu_cxx;
```