## ICPC Notebook

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## Data Structure and Basics

#### 1.1 Bitmask

1 2 2

```
#define isOn(S, j) (S & (1 << j))
#define setBit(S, j) (S |= (1 << j))
#define clearBit(S, j) (S &= ~(1 << j))
#define toggleBit(S, j) (S ^= (1 << j))
#define lowBit(S) (S & (-S))
#define setAll(S, n) (S = (1 << n) - 1)
#define modulo(S, N) ((S) & (N - 1)) // returns S \% N, where N is a power of 2
#define isPowerOfTwo(S) (!(S & (S - 1)))
#define nearestPowerOfTwo(S) ((int)pow(2.0, (int)((log((double)S) / log(2.0)) + 0.5)))
#define turnOffLastBit(S) ((S) & (S - 1))
#define turnOnLastZero(S) ((S) | (S + 1))
#define turnOffLastConsecutiveBits(S) ((S) & (S + 1))
#define turnOnLastConsecutiveZeroes(S) ((S) | (S - 1))
#define LSOne(S) (S & (-S))
```

#### 1.2 Sam fenwick tree

```
class FenwickTree { private:
    vector<long long int> ft;
    int n;
    FenwickTree(int _n) : n(_n) {
      ft.assign(n+1, 0);
    FenwickTree(const vector<int>& f) : n(f.size() - 1) {
      ft.assign(n+1, 0);
       for (int i = 1; i \le n; ++i) {
        ft[i] += f[i];
if(i + LSOne(i) <= n) {
           ft[i + LSOne(i)] += ft[i];
    long long int rsq(int b) {
      long long int sum = 0;
for(; b; b -= LSOne(b)) {
        sum += ft[b];
      return sum;
    long long int rsq(int a, int b) {
      return rsq(b) - rsq(a);
    void adjust(int k, long long int v) {
  for(; k <=n; k += LSOne(k)) {</pre>
         ft[k] += v;
};
```

### 1.3 Range fenwick tree

```
class FenwickTree {
private: vi ft1, ft2;
  int query(vi &ft. int b) {
    int sum = 0; for (; b; b -= LSOne(b)) sum += ft[b];
    return sum; }
  void adjust(vi &ft, int k, int v) {
    for (; k < (int)ft.size(); k += LSOne(k)) ft[k] += v; }</pre>
public:
 FenwickTree() {}
  FenwickTree(int n) { ft1.assign(n + 1, 0); ft2.assign(n+1, 0);}
  int query(int a) { return a * query(ft1, a) - query(ft2, a); }
  int query(int a, int b) { return query(b) - (a == 1 ? 0 : query(a-1)); }
  void adjust(int a, int b, int value){
    adjust(ft1, a, value);
    adjust(ft1, b+1, -value);
    adjust(ft2, a, value * (a-1));
adjust(ft2, b+1, -1 * value * b);
  int get (int n) {
    return query(n) - query(n-1); }
```

#### 1.4 **UFDS**

```
class UnionFind {
  public:
    vector<int> p, rank, setSize;
    int numSets:
    UnionFind(int N) {
     numSets = N;
      setSize.assign(N, 1);
      rank.assign(N, 0);
      p.assign(N, 0);
      for(int i = 0; i < N; ++i) {
       p[i] = i;
   int findSet(int i) {
     return (p[i] == i) ? i : (p[i] = findSet(p[i]));
    bool isSameSet(int i, int j) {
     bool x = findSet(i) == findSet(j);
    void unionSet(int i, int j) {
     if(!isSameSet(i, j)) {
       numSets--;
       int x = findSet(i);
       int y = findSet(j);
        if(rank[x] > rank[y]) {
         p[y] = x; setSize[x] += setSize[y];
         p[x] = y; setSize[y] += setSize[x];
          if(rank[x] == rank[y]) {
           rank[y]++;
    int numDisjoinSets() {
     return numSets;
    int sizeOfSet(int i) {
     return setSize[findSet(i)];
};
```

## 1.5 Segment tree

```
class SegmentTree {
                        // the segment tree is stored like a heap array
private: vi st, A;
                            // recall that vi is: typedef vector<int> vi;
  int n;
  int left (int p) { return p << 1; } // same as binary heap operations</pre>
  int right(int p) { return (p << 1) + 1; }</pre>
  void build(int p, int L, int R) {
   if (L == R)
                                          // as L == R, either one is fine
     st[p] = L;
                                                       // store the index
                                         // recursively compute the values
    else {
     build(left(p) , L
                                      (L + R) / 2);
     build(right(p), (L + R) / 2 + 1, R
     int p1 = st[left(p)], p2 = st[right(p)];
st[p] = (A[p1] <= A[p2]) ? p1 : p2;</pre>
  if (L >= i && R <= j) return st[p];</pre>
     \ensuremath{//} compute the min position in the left and right part of the interval
   if (p1 == -1) return p2; // if we try to access segment outside query
    if (p2 == -1) return p1;
                                                         // same as above
   return (A[p1] <= A[p2]) ? p1 : p2; } // as as in build routine
  int update_point(int p, int L, int R, int idx, int new_value) {
    // this update code is still preliminary, i == j
    // must be able to update range in the future!
    int i = idx, j = idx;
    // if the current interval does not intersect
    // the update interval, return this st node value!
    if (i > R | | j < L)
     return st[p];
    // if the current interval is included in the update range,
    // update that st[node]
    if (L == i && R == j) {
     A[i] = new_value; // update the underlying array
     return st[p] = L; // this index
    // compute the minimum pition in the
    // left and right part of the interval
    int p1, p2;
    p1 = update_point(left(p) , L
                                              , (L + R) / 2, idx, new_value);
    p2 = update_point(right(p), (L + R) / 2 + 1, R
                                                        , idx, new_value);
    \ensuremath{//} return the pition where the overall minimum is
    return st[p] = (A[p1] <= A[p2]) ? p1 : p2;</pre>
public:
 SegmentTree(const vi & A) {
   A = A; n = (int)A.size();
                                          // copy content for local usage
    st.assign(4 * n, 0);
                                   // create large enough vector of zeroes
   build(1, 0, n - 1);
                                                       // recursive build
  int rmq(int i, int j) { return rmq(1, 0, n - 1, i, j); } // overloading
  int update_point(int idx, int new_value) {
    return update_point(1, 0, n - 1, idx, new_value); }
```

## 1.6 Backtracking

```
/* 8 Queens Chess Problem */
#include <cstdlib>
                                                             // we use the int version of 'abs'
#include <cstdio>
#include <cstring>
using namespace std;
int row[8], TC, a, b, lineCounter;
                                                       // ok to use global variables
bool place(int r, int c) {
 for (int prev = 0; prev < c; prev++) // check previously placed queens</pre>
     if (row[prev] == r || (abs(row[prev] - r) == abs(prev - c)))
      return false;
                                     // share same row or same diagonal -> infeasible
 return true; }
void backtrack(int c) {
  if (c == 8 && row[b] == a) {
 // candidate sol, (a, b) has 1 queen
```

```
printf("%2d %d", ++lineCounter, row[0] + 1);
     for (int j = 1; j < 8; j++) printf(" %d", row[j] + 1);
    printf("\n"); }
  for (int r = 0; r < 8; r++)
                                                                                    // try all possible row
      if (place(r, c)) {
                                                 // if can place a queen at this col and row
         r \circ w [c] = r; backtrack(c + 1);
            // put this queen here and recurse
int main() {
  scanf("%d", &TC);
 while (TC--) {
    scanf ("%d %d", &a, &b); a--; b--;
    memset (row, 0, sizeof row); lineCounter = 0;
    printf ("SOLN COLUMN \n");
                                                               // switch to 0-based indexing
    printf("#
                              1 2 3 4 5 6 7 8\n\n");
    backtrack(0);
                                             // generate all possible 8! candidate solutions
      i f (TC) printf("\n");
} } // return 0;
```

#### 1.7 Sparse Table

#### 1.8 LCA

```
inline void makeP(){
   int logg = lg[n];
   for (int j = 1; j <= logg; j++) {
      for (int i = 1; i <= n; i++) {
        if [P[1][-1]! =-1] {
            P[i][j] = P[P[i][j-1]][j-1];
        }
    }
}

inline int LCA(int p, int q) {
   if (L[p] < L[q]) swap(p, q);
   int logg = lg[L[p]];
   for (int i = logg; i >= 0; i--) {
      if (L[p] - (1 << i) >= L[q]) {
        p = P[p][i] }
   }
if (p == q) return p;
   for (int i = logg; i >= 0; i--) {
      if (P[p][i] != -1 && P[p][i] != P[q][i]) {
        p = P[p][i];
        q = P[q][i];
    }
}
return P[p][0];
}
```

## 2 String

#### 2.1 KMP

```
Searches for the string w in the string s (of length k). Returns the
0-based index of the first match (k if no match is found). Algorithm
runs in O(k) time.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildTable(string& w, VI& t)
  t = VI(w.length());
 int i = 2, j = 0;

t[0] = -1; t[1] = 0;
  while(i < w.length())</pre>
    if(w[i-1] == w[j]) { t[i] = j+1; i++; j++; }
else if(j > 0) j = t[j];
    else { t[i] = 0; i++; }
int KMP(string& s, string& w)
  int m = 0, i = 0;
  VI t:
  buildTable(w, t);
  while (m+i < s.length())</pre>
    if(w[i] == s[m+i])
      if(i == w.length()) return m;
    else
      m += i-t[i];
      if(i > 0) i = t[i];
  return s.length();
```

### 2.2 Aho-Corasick

```
#include <bits/stdc++.h>
using namespace std;
// {{{ Aho Corasick
namespace AhoCorasick {
 struct Node {
    int next[26];
    int parent:
    int fromParent;
    int fail;
    int wordCount;
  int last;
 Node node[MAXNODE];
  void init() {
    last = 0;
    for (int i = 0; i < MAXNODE; i++) {
      for (int j = 0; j < 26; j++) {
  node[i].next[j] = -1;</pre>
      node[i].parent = node[i].fail = -1;
      node[i].wordCount = 0;
```

```
node[0].fail = 0;
int addString(char *s, int offset = 'a') {
 int len = strlen(s);
 int cur = 0;
 for (int i = 0; i < len; i++) { //no C++11? BibleThump</pre>
   if (node[cur].next[s[i] - offset] == -1) {
      node[cur].next[s[i] - offset] = ++last;
     node[last].parent = cur;
     node[last].fromParent = s[i] - offset;
    cur = node[cur].next[s[i] - offset];
 node[cur].wordCount++;
 return cur;
void constructFail() {
 queue<int> q;
  q.push(0);
  while (!q.empty()) {
   int u = q.front();
    g.pop();
   for (int i = 0; i < 26; i++) {
   if (node[u].next[i] != -1) {</pre>
        q.push(node[u].next[i]);
    if (u == 0) {
     continue;
    int v = node[node[u].parent].fail;
    int c = node[u].fromParent;
    while (v != 0 && node[v].next[c] == -1) {
     v = node[v].fail;
    int search(char *s, int offset) {
 int len = strlen(s);
 int cur = 0;
 int match = 0;
 for (int i = 0; i < len; i++) {</pre>
    while (cur != 0 \&\& node[cur].next[s[i] - offset] == -1) {
     assert (node[cur].fail != -1);
      cur = node[cur].fail;
    \textbf{if} \ (\texttt{node}[\texttt{cur}].\texttt{next}[\texttt{s}[\texttt{i}] \ - \ \texttt{offset}] \ != \ -1) \ \{
     cur = node[cur].next[s[i] - offset];
    assert (cur != -1);
   match += node[cur].wordCount;
 return match:
```

#### 2.3 Hash

```
// {{{ Hash
template<int pr, int MOD> class Hash {
  vector<int> p, v;
  void init() {
    p.resize(n):
    p[0] = 1;
    for (int i = 1; i < n; i++) {
      p[i] = (int)(1LL * p[i - 1] * pr % MOD);
v[i] = (int)(1LL * v[i - 1] * pr % MOD) + v[i];
      if (v[i] >= MOD) v[i] -= MOD;
public:
  Hash() {}
  Hash(string s) {
    n = (int)s.length();
    v.resize(n);
    for (int i = 0; i < n; i++) {</pre>
      v[i] = s;
    init();
  Hash(const vector<int> &v) {
    this->v = v;
```

```
n = (int)v.size();
for (auto &x : this->v) {
    x %= MOD;
    if (x < 0) x += MOD;
}
init();
}
int get(int 1, int r) {
    int res = v[r] - (l > 0 ? (int) (llL * v[l - 1] * p[r - l + 1] % MOD) : 0);
    if (res >= MOD) res -= MOD;
    if (res < 0) res += MOD;
    return res;
}
};
/// }};</pre>
```

## 2.4 SuffixArray

```
// Source: http://codeforces.com/contest/452/submission/7269543
// Efficient Suffix Array O(N*logN)
// String index from 0
// Usage:
// string s;
// SuffixArray sa(s);
// Now we can use sa.SA and sa.LCP
struct SuffixArray {
    string a;
    int N, m;
    vector<int> SA, LCP, x, y, w, c;
    SuffixArray(string _a, int m) : a(" " + _a), N(a.length()), m(m),
            SA(N), LCP(N), x(N), y(N), w(max(m, N)), c(N) {
        a[0] = 0;
        DA();
        kasaiLCP():
        #define REF(X) { rotate(X.begin(), X.begin()+1, X.end()); X.pop_back(); }
        REF(SA); REF(LCP);
        a = a.substr(1, a.size());
        for(int i = 0; i < SA.size(); ++i) --SA[i];</pre>
        #undef REF
    inline bool cmp (const int a, const int b, const int 1) { return (y[a] == y[b] && y[a+1] == y[b]
    void Sort() {
        for (int i = 0; i < m; ++i) w[i] = 0;
        for(int i = 0; i < N; ++i) ++w[x[y[i]]];</pre>
        for (int i = 0; i < m - 1; ++i) w[i + 1] += w[i];
        for (int i = N - 1; i \ge 0; --i) SA[--w[x[y[i]]]] = y[i];
    void DA() {
        for (int i = 0; i < N; ++i) x[i] = a[i], y[i] = i;
        for(int i, j = 1, p = 1; p < N; j <<= 1, m = p) {
            for (p = 0, i = N - j; i < N; i++) y[p++] = i;
            for (int k = 0; k < N; ++k) if (SA[k] >= j) y[p++] = SA[k] - j;
            for (swap(x, y), p = 1, x[SA[0]] = 0, i = 1; i < N; ++i)
                x[SA[i]] = cmp(SA[i-1], SA[i], j) ? p-1 : p++;
    void kasaiLCP() {
        for (int i = 0; i < N; i++) c[SA[i]] = i;
for (int i = 0, j, k = 0; i < N; LCP[c[i++]] = k)
            if (c[i] > 0) for (k ? k-- : 0, j = SA[c[i] - 1]; a[i + k] == a[j + k]; k++);
};
```

# 3 Graph Basic

## 3.1 Dijkstra

```
int main() {
  int V, E, s, u, v, w;
  vector < vii > AdjList;
```

```
AdjList.assign(V, vii()); // assign blank vectors of pair<int, int>s to AdjList
for (int i = 0; i < E; i++) {
  scanf ("%d %d %d", &u, &v, &w);
  AdjList[u].push_back(ii(v, w));
     directed graph
 // Dijkstra routine
vi dist(V, INF); dist[s] = 0;
                                                                    // INF = 1B to avoid overflow
priority_queue< ii, vector<ii>, greater<ii> > pq; pq.push(ii(0, s));
                                                      // ^to sort the pairs by increasing distance
while (!pq.empty()) {
                                                                                            // main
    i i front = pq.top(); pq.pop();
                                           // greedy: pick shortest unvisited vertex
   int d = front.first, u = front.second;
    i f (d > dist[u]) continue;  // this check is important, see the explanation
    f o r (int j = 0; j < (int)AdjList[u].size(); j++) {</pre>
       i i v = AdjList[u][j];
                                                                         // all outgoing edges
     from u
       i f (dist[u] + v.second < dist[v.first]) {</pre>
          dist[v.first] = dist[u] + v.second;
                                                                               // relax operation
           p q .push(ii(dist[v.first], v.first));
 }}} // note: this variant can cause duplicate items in the priority queue
for (int i = 0; i < V; i++) // index + 1 for final answer</pre>
  printf("SSSP(%d, %d) = %d\n", s, i, dist[i]);
return 0:
```

#### 3.2 Bellman Ford

```
int main() {
  int V, E, s, a, b, w;
 vector <vii>> AdjList;
 \label{eq:AdjList} AdjList.assign(V, vii()); // assign blank vectors of pair<int, int>s to AdjList for (int i = 0; i < E; i++) {s < E \ i+c} ( s < E \ i) {s < E \ i} ( %d %d", &a, &b, &w); 
    AdjList[a].push_back(ii(b, w));
   // Bellman Ford routine
  vi dist(V, INF); dist[s] = 0;
for (int i = 0; i < V - 1; i++)</pre>
                                       // relax all E edges V-1 times, overall O(VE)
     for (int u = 0; u < V; u++)
                                                                                         // these two loops =
       0 (E)
         f o r (int j = 0; j < (int)AdjList[u].size(); j++) {</pre>
              i i v = AdjList[u][j];
                                                        // we can record SP spanning here if needed
             d i s t [v.first] = min(dist[v.first], dist[u] + v.second);
  bool hasNegativeCycle = false:
  for (int u = 0; u < V; u++)
                                                                                          // one more pass to
       check
     f o r (int j = 0; j < (int)AdjList[u].size(); j++) {</pre>
         i i v = AdjList[u][j];
          i f (dist[v.first] > dist[u] + v.second)
                                                                                            // should be false
        hasNegativeCycle = true; // but if true, then negative cycle exists!
 printf("Negative Cycle Exist? %s\n", hasNegativeCycle ? "Yes" : "No");
  i f (!hasNegativeCycle)
     f o r (int i = 0; i < V; i++)</pre>
       printf("SSSP(%d, %d) = %d\n", s, i, dist[i]);
 return 0:
```

### 3.3 MCBM

```
for (int j = 0; j < (int)AdjList[1].size(); j++) {</pre>
    int r = AdjList[1][j];
     i f (match[r] == -1 || Aug(match[r])) {
       match[r] = 1; return 1;
                                                                                      // found 1
      matching
 return 0;
       // no matching
int main() {
 int V = 5, Vleft = 3;
                                                                                    // we ignore
      vertex 0
 AdjList.assign(V, vi());
AdjList[1].push_back(3); AdjList[1].push_back(4);
 AdjList[2].push_back(3);
  int MCBM = 0;
 match.assign(V, -1);
                           // V is the number of vertices in bipartite graph
  for (int 1 = 0; 1 < Vleft; 1++) { // Vleft = size of the left set
    v i s .assign(Vleft, 0);
                                                                 // reset before each recursion
    M \subset B M += Aug(1);
 printf ("Found %d matchings\n", MCBM); // the answer is 2 for Figure 4.42
 return 0:
```

# 4 Graph

#### 4.1 2-SAT

```
// VAR (x) , NOT (x)
// TwoSat(int varCount)
// bool TwoSat.solve() -> vector<int> TwoSat.value
// {{{ 2-SAT
using graph = vector<vi>;
inline int VAR(int x) { return 2 * x; }
inline int NOT(int x) { return x ^ 1; }
struct TwoSat {
 int n;
  vi value;
  graph g, grev;
  TwoSat(int n) : n(2 * n), g(2 * n) {}
  void addImplication(int x, int y) {
   g[x].emplace_back(y);
    g[NOT(y)].emplace_back(NOT(x));
  void addOr(int x, int y) {
   g[NOT(x)].emplace_back(y);
    g[NOT(y)].emplace_back(x);
  void dfs(int u, graph &g, vi &visit, vi &order) {
    visit[u] = 1;
    for (int v : g[u]) {
     if (!visit[v]) {
       dfs(v, g, visit, order);
    order.push back(u):
  bool solve() {
    grev.assign(n, vi());
    for (int i = 0; i < n; i++) {
     for (int j : g[i]) {
       grev[j].push_back(i);
    vi visited(n, false);
    vi order;
    for (int i = 0; i < n; i++) {
     if (!visited[i]) {
       dfs(i, g, visited, order);
    reverse (ALL (order));
    vi repr(n);
    fill(ALL(visited), false);
    value.assign(n, -1);
```

```
for (int u : order) {
   if (!visited[u]) {
      vi tmp;
      dfs(u, grev, visited, tmp);
      for (int v : tmp) {
        repr[v] = tmp[0];
        if (value[v] = -1) {
            value[v] = 0;
        }
      }
   }
   for (int i = 0; i < n; i++) {
      if (repr[i] == repr[NOT(i)]) {
        return false;
    }
}
   return true;
}</pre>
```

## 4.2 Biconnected Component

```
// Input graph: vector< vector<int> > a, int n
// Note: 0-indexed
// Usage: BiconnectedComponent bc; (bc.components is the list of components)
struct BiconnectedComponent {
 vector<int> low, num, s;
  vector< vector<int> > components:
  int counter:
  BiconnectedComponent() : num(n, -1), low(n, -1), counter(0) {
   for (int i = 0; i < n; i++)
     if (num[i] < 0)
       dfs(i, 1);
  void dfs(int x, int isRoot) {
   low[x] = num[x] = ++counter;
   if (a[x].empty()) {
     components.push_back(vector<int>(1, x));
     return:
   s.push_back(x);
   for (int i = 0; i < a[x].size(); i++) {</pre>
     int y = a[x][i];
     if (num[y] > -1) low[x] = min(low[x], num[y]);
       dfs(y, 0);
        low[x] = min(low[x], low[y]);
       if (isRoot || low[y] >= num[x]) {
          components.push_back(vector<int>(1, x));
          while (1) {
           int u = s.back();
           s.pop_back();
            components.back().push_back(u);
           if (u == y) break;
```

## 4.3 Bridge Articulation

```
// Assume already have undirected graph vector< vector<int> > G with V vertices
// Vertex index from 0
// Usage:
// UndirectedDfs tree;
// Then you can use tree.bridges and tree.cuts
struct UndirectedDfs {
    vector<int> low, num, parent;
    vector<bool> articulation;
    int counter, root, children;

    vector< pair<int, int> > bridges;
    vector<int> cuts;
```

```
UndirectedDfs() : low(V, 0), num(V, -1), parent(V, 0), articulation(V, false),
             counter(0), children(0)
         for (int i = 0; i < V; ++i) if (num[i] == -1) {
             root = i; children = 0;
             articulation[root] = (children > 1);
        for (int i = 0; i < V; ++i)
             if (articulation[i]) cuts.push_back(i);
private:
    void dfs(int u) {
        for(int j = num[u] = counter++;
for(int j = 0; j < G[u].size(); ++j) {
  int v = G[u][j];</pre>
             if (num[v] == -1) {
                 parent[v] = u;
                 if (u == root) children++;
                 if (low[v] >= num[u])
                     articulation[u] = true;
                 if (low[v] > num[u]) bridges.push_back(make_pair(u, v));
                 low[u] = min(low[u], low[v]);
             } else if (v != parent[u])
                 low[u] = min(low[u], num[v]);
1:
```

#### 4.4 Dinic Max Flow

```
#include <bits/stdc++.h>
using namespace std;
//Add Edge : MaxFlow.addEdge(int from, int to, F capacity)
//Solve : MaxFlow.calcMaxFlow(int source, int sink)
// {{{ Dinic Max Flow<F=int>
template<typename F = int>
struct DinicMaxFlow {
 struct Edge {
    int to, rev;
    F flow, cap;
  };
  const F INF_FLOW = numeric_limits<F>::max() / 2;
  const int INF_DIST = 1e9;
  int source, sink;
  vector<vector<Edge>> graph;
  vector<int> dist, ptr;
 DinicMaxFlow(int N) : graph(N), dist(N), ptr(N) {}
  void addEdge(int from, int to, F cap) {
  if (from == to) {
      return;
    Edge e1 { to, SZ(graph[to]), 0, cap };
    Edge e2 { from, SZ(graph[from]), 0, 0 };
    graph[from].push_back(e1);
    graph[to].push_back(e2);
  bool buildLevelGraph() {
    fill(begin(dist), end(dist), INF_DIST);
    queue<int> q;
    dist[source] = 0;
    a.push (source):
    while (!q.empty()) {
      int u = q.front();
      g.pop();
      for (const auto &e : graph[u]) {
        if (e.flow < e.cap && dist[e.to] > dist[u] + 1) {
    dist[e.to] = dist[u] + 1;
          q.push(e.to);
    return dist[sink] != INF_DIST;
  F dfs(int u, F bottleneck) {
  if (u == sink || bottleneck == 0) {
      return bottleneck;
    for (int &i = ptr[u]; i < SZ(graph[u]); i++) {</pre>
      auto &e = graph[u][i];
```

```
auto &b = graph[e.to][e.rev];
      if (dist[e.to] != dist[u] + 1 || e.flow >= e.cap) {
      F pushed = dfs(e.to, min(e.cap - e.flow, bottleneck));
      if (pushed > 0) {
       e.flow += pushed;
        b.flow -= pushed;
        return pushed;
    return 0;
  F calcMaxFlow(const int source, const int sink) {
    this->source = source;
    this->sink = sink;
    F mf = 0;
    while (buildLevelGraph()) {
     fill(begin(ptr), end(ptr), 0);
      F pushed;
      while ((pushed = dfs(source, INF_FLOW))) {
       mf += pushed;
    return mf;
};
```

#### 4.5 Directed MST

```
#include "../../template.h"
const int maxe = 100111:
const int maxy = 100;
// Index from 0
// Running time O(E*V)
namespace chuliu {
    struct Cost:
    vector<Cost> costlist;
    struct Cost {
       int id, val, used, a, b, pos;
        Cost() { val = -1; used = 0; }
        Cost(int _id, int _val, bool temp) {
            a = b = -1; id = _id; val = _val; used = 0;
            pos = costlist.size(); costlist.push_back(*this);
        Cost(int _a, int _b) {
            a = _a; b = _b; id = -1; val = costlist[a].val - costlist[b].val;
            used = 0; pos = costlist.size(); costlist.push_back(*this);
        void push() {
           if (id == -1) {
               costlist[a].used += used;
               costlist[b] used -= used;
    };
    struct Edge {
        int u, v;
        Cost cost;
        Edge() {}
        Edge(int id, int _u, int _v, int c) {
           u = _u; v = _v; cost = Cost(id, c, 0);
    } edge[maxe];
    int n, m, root, pre[maxv], node[maxv], vis[maxv], best[maxv];
    void init(int _n) {
        n = _n; m = 0;
        costlist.clear();
    void add(int id, int u, int v, int c) {
        edge[m++] = Edge(id, u, v, c);
    int mst(int root) {
       int ret = 0;
        while (true) {
            REP(i, n) best[i] = -1;
```

```
REP(e, m) {
                int u = edge[e].u, v = edge[e].v;
                if ((best[v] == -1 || edge[e].cost.val < costlist[best[v]].val) && u != v) {</pre>
                    best[v] = edge[e].cost.pos;
            REP(i, n) if (i != root && best[i] == -1) return -1;
            int cntnode = 0;
            memset(node, -1, sizeof node); memset(vis, -1, sizeof vis);
            REP(i, n) if (i != root)
                ret += costlist[best[i]].val;
                costlist[best[i]].used++;
                while (vis[v] != i \&\& node[v] == -1 \&\& v != root) {
                    v = pre[v];
                if (v != root && node[v] == -1) {
                    for (int u = pre[v]; u != v; u = pre[u]) node[u] = cntnode;
                    node[v] = cntnode++;
            if (cntnode == 0) break;
            REP(i, n) if (node[i] == -1) node[i] = cntnode++;
            REP(e, m) {
                int v = edge[e].v;
                edge[e].u = node[edge[e].u];
                edge[e].v = node[edge[e].v];
                if (edge[e].u != edge[e].v) edge[e].cost = Cost(edge[e].cost.pos, best[v]);
            n = cntnode:
            root = node[root];
        return ret;
    vector<int> trace() {
        vector<int> ret;
        FORD(i, costlist.size()-1,0) costlist[i].push();
        REP(i, costlist.size()) {
            Cost cost = costlist[i];
            if (cost.id != -1 && cost.used > 0) ret.push_back(cost.id);
        return ret;
int main() {
```

# 4.6 Eulerian Cycle

```
// directed!
// careful with empty graph!
// {{{ enlerianTour
template<typename T>
bool eulerianTour(const vector<vector<T>> &g, vi &tour) {
  int n = SZ(q);
  int edgeCount = 0;
  vi inDeg(n), outDeg(n);
  for (int i = 0; i < n; i++) {
    outDeg[i] = SZ(g[i]);
    for (auto &it : q[i]) {
      inDeg[it.fi]++;
      edgeCount++;
  int out = -1, in = -1;
for (int i = 0; i < n; i++) {</pre>
    if (inDeg[i] == outDeg[i] + 1) {
      debug printf("finish vertex %d\n", i);
if (out == -1) {
        out = i;
      } else {
        return false;
```

```
} else if (outDeg[i] == inDeg[i] + 1) {
    debug printf("start vertex %d\n", i);
    if (in == -1) {
     in = i;
    } else {
      return false;
  } else if (outDeg[i] != inDeg[i]) {
    return false;
// either zero or both
if ((in != -1) ^ (out != -1)) {
 return false;
if (in == -1) {
  for (int i = 0; i < n; i++) {
   if (outDeg[i] != 0) {
      in = out = i;
     break;
  // empty graph?
  if (in == -1) {
    in = out = 0;
tour.clear();
stack<int> st;
vi visit(n);
vi ptr(n):
st.push(in);
while (!st.empty()) {
  int u = st.top();
  visit[u] = true;
  if (ptr[u] == SZ(g[u])) {
   tour.push_back(u);
    st.pop();
  } else {
   st.push(g[u][ptr[u]++].fi);
reverse (ALL(tour));
return SZ(tour) == edgeCount + 1;
```

## 4.7 Edmonds Blossom

// General matching on graph

```
const int maxv = 1000:
const int maxe = 50000;
// Index from 1
// Directed
struct EdmondsLawler {
    int n, E, start, finish, newRoot, qsize, adj[maxe], next[maxe], last[maxv], mat[maxv], que[maxv],
          dad[maxv], root[maxv];
    bool inque[maxv], inpath[maxv], inblossom[maxv];
    void init(int _n) {
        n = _n; E = 0;
        for (int x=1; x<=n; ++x) { last[x] = -1; mat[x] = 0; }</pre>
    void add(int u, int v) {
   adj[E] = v; next[E] = last[u]; last[u] = E++;
    int lca(int u, int v) {
        for(int x=1; x<=n; ++x) inpath[x] = false;</pre>
        while (true) {
            u = root[u];
             inpath[u] = true;
            if (u == start) break;
            u = dad[mat[u]];
        while (true) {
             v = root[v];
            if (inpath[v]) break;
            v = dad[mat[v]];
        return v;
    void trace(int u) {
        while (root[u] != newRoot) {
            int v = mat[u];
            inblossom[root[u]] = true;
```

```
inblossom[root[v]] = true;
             if (root[u] != newRoot) dad[u] = v;
    void blossom(int u, int v) {
         for(int x=1; x<=n; ++x) inblossom[x] = false;</pre>
         newRoot = lca(u, v);
         trace(u); trace(v);
         if (root[u] != newRoot) dad[u] = v;
         if (root[v] != newRoot) dad[v] = u;
         for(int x=1; x<=n; ++x) if (inblossom[root[x]]) {</pre>
             root[x] = newRoot;
             if (!inque[x]) {
                  inque[x] = true;
                  que[qsize++] = x;
    bool bfs() {
         for(int x=1; x<=n; ++x) {</pre>
             inque[x] = false;
dad[x] = 0;
root[x] = x;
         qsize = 0;
         que[qsize++] = start;
         inque[start] = true;
         for(int i=0; i<qsize; ++i) {</pre>
             int u = que[i];
             for (int e = last[u]; e != -1; e = next[e]) {
                  int v = adj[e];
                  if (root[v] != root[u] && v != mat[u]) {
   if (v == start || (mat[v] > 0 && dad[mat[v]] > 0)) blossom(u, v);
   else if (dad[v] == 0) {
                           dad[v] = u;
                           if (mat[v] > 0) que[qsize++] = mat[v];
                           else {
                               finish = v;
                               return true;
         return false:
    void enlarge() {
         int u = finish;
while (u > 0) {
             int v = dad[u], x = mat[v];
             mat[v] = u;
             mat[u] = v;
             u = x:
         for(int x=1; x<=n; ++x) if (mat[x] == 0) {</pre>
             if (bfs()) enlarge();
         int ret = 0;
         for(int x=1; x<=n; ++x) if (mat[x] > x) ++ret;
         return ret:
| edmonds:
```

## 4.8 Hungarian Assignment

```
const int MN = 3210;
const int inf = 1000111000;

// Vertex: 1 --> nx, 1 --> ny
// cost >= 0
// fx[x] + fy[y] <= cost[x][y] for all x, y
// Min cost matching

struct Hungary {
   int nx, ny, cost[MN][MN], fx[MN], fy[MN], matx[MN], which[MN], dist[MN];
   bool used[MN];</pre>
```

```
void init(int _nx, int _ny) {
        nx = _nx; ny = _ny;
        memset (fx, 0, sizeof fx);
        memset(fy, 0, sizeof fy);
        memset (used, false, sizeof used);
        memset(matx, 0, sizeof matx);
        for(int i=0; i<=nx; ++i) for(int j=0; j<=ny; ++j) cost[i][j] = inf;</pre>
    void add(int x, int y, int c) { cost[x][y] = min(cost[x][y],c); }
    int mincost() {
        for (int x=1; x<=nx; ++x) {</pre>
             int y0 = 0; matx[0] = x;
             for(int y=0; y<=ny; ++y) { dist[y] = inf + 1; used[y] = false; }</pre>
                 used[y0] = true;
                 int x\bar{0} = matx[y0], delta = inf + 1, y1;
                 for(int y=1; y<=ny; ++y) if (!used[y]) {</pre>
                     int curdist = cost[x0][y] - fx[x0] - fy[y];
if (curdist < dist[y]) {</pre>
                          dist[y] = curdist;
                          which[y] = y0;
                     if (dist[y] < delta) {</pre>
                         delta = dist[y];
                         y1 = y;
                 for(int y=0; y<=ny; ++y) if (used[y]) {</pre>
                     fx[matx[y]] += delta;
                     fy[y] -= delta;
                 } else dist[y] -= delta;
                 y0 = y1;
             } while (matx[y0] != 0);
                 int y1 = which[y0];
                 matx[y0] = matx[y1];
             } while (y0);
        return -fy[0]; // n u lu n
                                            m b o c b ghp
        int ret = 0;
        for (int y=1; y<=ny; ++y) {</pre>
            int x = matx[y];
             if (cost[x][y] < inf) ret += cost[x][y];</pre>
        return ret;
} hungary;
```

#### 4.9 Min Cost Max Flow

```
// Min Cost Max Flow - SPFA
// Index from 0
// Lots of double comparison --> likely to fail for double
// Example:
// MinCostFlow mcf(n);
// mcf.addEdge(1, 2, 3, 4);
// cout << mcf.minCostFlow() << endl;</pre>
template<class Flow=int, class Cost=int>
struct MinCostFlow {
    const Flow INF_FLOW = 1000111000;
    const Cost INF_COST = 1000111000111000LL;
    int n, t, S, T;
    Flow totalFlow;
    Cost totalCost;
    vector<int> last, visited;
    vector<Cost> dis;
    struct Edge {
        int to;
        Flow cap:
        Cost cost;
        int next:
        Edge (int to, Flow cap, Cost cost, int next) :
                to(to), cap(cap), cost(cost), next(next) {}
    vector<Edge> edges;
```

```
 \texttt{MinCostFlow(int n)} : \texttt{n(n), t(0), totalFlow(0), totalCost(0), last(n, -1), visited(n, 0), dis(n, 0) } 
         edges.clear();
    int addEdge(int from, int to, Flow cap, Cost cost) {
         edges.push_back(Edge(to, cap, cost, last[from]));
         last[from] = t++;
         edges.push_back(Edge(from, 0, -cost, last[to]));
         last[to] = t++;
         return t - 2;
    pair<Flow, Cost> minCostFlow(int _S, int _T) {
         S = \_S; T = \_T;
         SPFA();
         while (1) {
             while (1) {
                  REP(i,n) visited[i] = 0;
                  if (!findFlow(S, INF_FLOW)) break;
             if (!modifyLabel()) break;
         return make_pair(totalFlow, totalCost);
private:
    void SPFA() {
         REP(i,n) dis[i] = INF_COST;
         priority_queue< pair<Cost, int> > Q;
         Q.push (make_pair(dis[S]=0, S));
         while (!Q.empty()) {
             int x = Q.top().second;
             Cost d = -Q.top().first;
             Q.pop();
              // For double: dis[x] > d + EPS
             if (dis[x] != d) continue;
             for(int it = last[x]; it >= 0; it = edges[it].next)
                  if (edges[it].cap > 0 && dis[edges[it].to] > d + edges[it].cost)
                      Q.push(make_pair(-(dis[edges[it].to] = d + edges[it].cost), edges[it].to));
         REP(i,n) dis[i] = dis[T] - dis[i];
    Flow findFlow(int x, Flow flow) {
         if (x == T) {
             totalCost += dis[S] * flow;
             totalFlow += flow;
             return flow;
         visited[x] = 1;
         Flow now = flow;
         for(int it = last[x]; it >= 0; it = edges[it].next)
             // For double: fabs(dis[edges[it].to] + edges[it].cost - dis[x]) < EPS
if (edges[it].cap && !visited[edges[it].to] && dis[edges[it].to] + edges[it].cost == dis[x</pre>
                  Flow tmp = findFlow(edges[it].to, min(now, edges[it].cap));
                  edges[it].cap -= tmp;
edges[it ^ 1].cap += tmp;
                  now -= tmp;
                  if (!now) break;
         return flow - now;
    bool modifyLabel() {
         Cost d = INF_COST;
         REP(i,n) if (visited[i])
             for(int it = last[i]; it >= 0; it = edges[it].next)
if (edges[it].cap && !visited[edges[it].to])
    d = min(d, dis[edges[it].to] + edges[it].cost - dis[i]);
         // For double: if (d > INF_COST / 10)
if (d == INF_COST) return false;
                                                         INF\ COST = 1e20
         REP(i,n) if (visited[i])
         return true;
};
```

## 4.10 Strongly Connected Component

```
// Assume that already have directed graph vector< vector<int> > G with V vertices // Index from 0 // Usage: // DirectedDfs tree; // Now you can use tree.scc
```

```
struct DirectedDfs {
    vector<int> num, low, current, S;
    int counter;
    vector< vector<int> > scc;
    DirectedDfs(): num(V, -1), low(V, 0), current(V, 0), counter(0) {
        REP(i, V) if (num[i] == -1) dfs(i);
    void dfs(int u) {
        low[u] = num[u] = counter++;
        S.push_back(u);
        current[u] = 1;
        REP(j, G[u].size()) {
            int v = G[u][j];
if (num[v] == -1) dfs(v);
            if (current[v]) low[u] = min(low[u], low[v]);
        if (low[u] == num[u]) {
            scc.push_back(vector<int>());
            while (1) {
                int v = S.back(); S.pop_back(); current[v] = 0;
                scc.back().push_back(v);
                if (u == v) break;
};
```

#### 4.11 LRFlow

```
#include "Dinic.cpp"
//Construct : LRFlow(int N, int originalSource, int originalSink)
//Add Edge : LRFlow.addEdge(int from, int to, F lowerBound, F upperBound) []
//Solve: LRFlow.calcLRFlow()
// {{{ LR FLow<F=int, MF=DinicMaxFlow<F>>
template<typename F = int, typename MF = DinicMaxFlow<F>>
struct LRFlow {
  MF mf:
  F lowerBoundSum = 0;
  int n, realSource, realSink;
  LRFlow(int n, int src, int snk) : mf(n + 2), n(n + 2) {
  realSource = this->n - 2;
  realSink = this->n - 1;
    mf.addEdge(snk, src, mf.INF_FLOW);
  void addEdge(int from, int to, F lowerBound, F upperBound) {
    assert(lowerBound <= upperBound);</pre>
    lowerBoundSum += lowerBound;
    mf.addEdge(from, to, upperBound - lowerBound);
    mf.addEdge(realSource, to, lowerBound);
mf.addEdge(from, realSink, lowerBound);
  pair<F, bool> calcLRFlow() {
    F maxFlow = mf.calcMaxFlow(realSource, realSink);
    return make_pair(maxFlow, maxFlow == lowerBoundSum);
};
```

## 4.12 Hopcroft Karp

```
#include <algorithm>
#include <iostream>
using namespace std;

const int MAXN1 = 50000;
const int MAXN2 = 50000;
const int MAXM = 150000;
int n1, n2, edges, last[MAXN1], prev[MAXM], head[MAXM];
int matching[MAXN2], dist[MAXN1], Q[MAXN1];
bool used[MAXN1], vis[MAXN1];

void init(int _n1, int _n2) {
    n1 = _n1;
    n2 = _n2;
    edges = 0;
```

```
fill(last, last + n1, -1);
void addEdge(int u, int v) {
    head[edges] = v;
prev[edges] = last[u];
    last[u] = edges++;
void bfs() {
    fill(dist, dist + n1, -1);
    int sizeQ = 0;
    for (int u = 0; u < n1; ++u) {
        if (!used[u]) {
            Q[sizeQ++] = u;
            dist[u] = 0;
    for (int i = 0; i < sizeQ; i++) {</pre>
        int u1 = Q[i];
        for (int e = last[u1]; e >= 0; e = prev[e]) {
            int u2 = matching[head[e]];
            if (u2 >= 0 && dist[u2] < 0) {
                dist[u2] = dist[u1] + 1;
Q[sizeQ++] = u2;
bool dfs(int u1) {
    vis[u1] = true;
    for (int e = last[u1]; e >= 0; e = prev[e]) {
        int v = head[e];
        int u2 = matching[v];
        if (u2 < 0 \mid \mid !vis[u2] \&\& dist[u2] == dist[u1] + 1 \&\& dfs(u2)) {
            matching[v] = u1;
            used[u1] = true;
            return true;
    return false:
int maxMatching() {
    fill(used, used + n1, false);
    fill (matching, matching + n2, -1);
    for (int res = 0;;) {
        bfs();
        fill(vis, vis + n1, false);
        int f = 0;
        for (int u = 0; u < n1; ++u)
            if (!used[u] && dfs(u))
                ++f;
        if (!f)
           return res:
        res += f:
int main() {
    init(2, 2);
    addEdge(0, 0);
    addEdge(0, 1);
    addEdge(1, 1);
    cout << (2 == maxMatching()) << endl;
```

## 4.13 Heavy Light Decomposition

```
template <class T, int V>
class HeavyLight {
  int parent[V], heavy[V], depth[V];
  int root[V], treePos[V];
SegmentTree<T> tree;

template <class G>
  int dfs(const G& graph, int v) {
   int size = 1, maxSubtree = 0;
   for (int u : graph[V]) if (u != parent[V]) {
     parent[u] = V;
   depth[u] = depth[V] + 1;
   int subtree = dfs(graph, u);
   if (subtree > maxSubtree) heavy[V] = u, maxSubtree = subtree;
   size += subtree;
```

```
return size;
  template <class BinaryOperation>
  void processPath(int u, int v, BinaryOperation op) {
    for (; root[u] != root[v]; v = parent[root[v]]) {
      if (depth[root[u]] > depth[root[v]]) swap(u, v);
      op(treePos[root[v]], treePos[v] + 1);
    if (depth[u] > depth[v]) swap(u, v);
    op(treePos[u], treePos[v] + 1);
public:
  template <class G>
  void init (const G& graph)
    int n = graph.size();
    fill_n(heavy, n, -1);
    parent[0] = -1;
    depth[0] = 0;
    dfs(graph, 0);
    for (int i = 0, currentPos = 0; i < n; ++i)</pre>
     if (parent[i] == -1 || heavy[parent[i]] != i)
       for (int j = i; j != -1; j = heavy[j]) {
  root[j] = i;
          treePos[j] = currentPos++;
    tree.init(n):
  void set(int v, const T& value) {
    tree.set(treePos[v], value);
  void modifyPath(int u, int v, const T& value) {
   processPath(u, v, [this, &value](int 1, int r) { tree.modify(1, r, value); });
  T queryPath(int u, int v) {
    T res = T();
    processPath(u, v, [this, &res](int 1, int r) { res.add(tree.query(1, r)); });
    return res;
};
```

## 4.14 Gomory-Hu

```
// tree which represents all-pair min-cut
// min-cut of u-v = minimum edge in path from u to v in tree
// combine the following function with usual max flow routine
// also because we need to reset the flow graph during each max flow, // don't forget to add variable
    in struct to record initial capacity
void buildGomoryHu() {
    for(int i = 2; i <= n; i++)
        parent[i] = 1;
    for(int i = 2; i <= n; i++) {
        int minCut = maxFlow(i,parent[i]);
        tree[parent[i]].pb((i,minCut));
        // the one below is not necessary if we want to make a rooted tree tree[i].pb({parent[i],minCut});
    for(int j = i + 1; j <= n; j++) {
        if(dist[j] != -1 && parent[j] == parent[i]) {
            parent[j] = i;
        }
    }
}</pre>
```

## 4.15 Online-Bridge

```
// O(N + M)
int psetBridge[N], psetComp[N]; // UF super vertice, UF tree int par[N];
bool seen[N];
int size[N];
int bridge;
int n;
void reset() {
  for(int i = 0 ; i <= n ; i++) {
    psetBridge[i] = psetComp[i] = i;
    par[i] = -1;
    seen[i] = 0;
    size[i] = 1;</pre>
```

```
bridge = 0;
int findsB(int x) {
 if(x == -1) return -1;
  return x == psetBridge[x] ? x : psetBridge[x] = findsB(psetBridge[x]);
int findsC(int x) {
  x = findsB(x);
  return x == psetComp[x] ? x : psetComp[x] = findsC(psetComp[x]);
void makeRoot(int x) { // make super vertice which contains x root of its tree
  x = findsB(x);
  int root = x;
  int chld = -1;
  while (x != -1) {
   int papa = findsB(par[x]);
   par[x] = chld;
    psetComp[x] = root;
    chld = x;
    x = papa;
  size[root] = size[chld];
void mergePath(int u,int v) { // remove bridge between u and v in tree vector<int> vu,vv;
  int lca = -1;
  while(1) {
   if(u != -1) {
     u = findsB(u):
      vu.push_back(u);
      if(seen[u]) {
        lca = u:
        break;
      seen[u] = 1;
      u = par[u];
    if(v != -1) {
      v = findsB(v);
      vv.push_back(v);
      if(seen[v]) {
        lca = v;
        break:
      seen[v] = 1;
      v = par[v];
  for (int i = 0; i < 2; i++) {
     vector<int> &proc = i ? vv : vu;
    for(int x : proc) {
      psetBridge[x] = lca;
      if(x == lca) {
        break:
      bridge--:
    for(int x : proc)
      seen[x] = 0;
void addEdge(int u,int v) {
 u = findsB(u);
  v = findsB(v);
  if(u != v) {
    int uu = findsC(u);
    int vv = findsC(v);
    if(uu != vv) {
      bridge++;
      if(size[uu] > size[vv]) {
        swap(u,v);
        swap (vv, uu);
      makeRoot(u);
      par[u] = psetComp[u] = v;
      size[vv] += size[u];
    } else
      mergePath(u,v);
```

## 5 Math Basic

## 5.1 Big Integer

import java.io.\*;

```
import java.util.*;
import java.util.Scanner;
import java.math.BigInteger;
class BigInteger2 {
 public static void main(String[] args) throws Exception {
   BufferedReader br = new BufferedReader(new InputStreamReader(System.in));
   PrintWriter pw = new PrintWriter(new BufferedWriter(new OutputStreamWriter(System.out)));
    while (true) {
       String N = br.readLine();
         i f (N == null) break;
      BigInteger BN = new BigInteger(N);
BigInteger BRN = new BigInteger(
            n e w StringBuffer(BN.toString()).reverse().toString());
         p w .printf("%s is ", N);
i f (!BN.isProbablePrime(10)) // 10 is enough
             p w .printf("not prime.\n";
         else if (!BN.equals(BRN) && BRN.isProbablePrime(10))
            p w .printf("emirp.\n")
             p w .printf("prime.\n")
     p w .close();
import java.util.Scanner; // Scanner class is inside package java.util
import java.math.BigInteger; // BigInteger class is inside package java.math
class Main { /* UVa 10925 - Krakovia, 0.732s in Java */
 public static void main(String[] args) {
    Scanner sc = new Scanner(System.in);
     int caseNo = 1;
    while (true) {
         int N = sc.nextInt(), F = sc.nextInt(); // N bills, F friends
         i f (N == 0 && F == 0) break;
      BigInteger sum = BigInteger ZERO; // BigInteger has this constant ZERO
         f o r (int i = 0; i < N; i++) { // sum the N large bills
BigInteger V = sc.nextBigInteger(); // for reading next BigInteger!</pre>
             s u m = sum.add(V); // this is BigInteger addition
       System.out.println("Bill #" + (caseNo++) + " costs " +
            s u m + ": each friend should pay " + sum.divide(BigInteger.valueOf(F)));
       System.out.println(); // the line above is BigInteger division
                                                   // divide the large sum to F friends
} } }
```

#### 5.2 Prime

```
bitset<10000010> bs; // 10^7 should be enough for most cases
vi primes;
             // compact list of primes in form of vector<int>
// first part
void sieve(ll upperbound) {
                                                // create list of primes in [0..upperbound]
 _sieve_size = upperbound + 1;
                                                                        // add 1 to include upperbound
  bs.set();
        // set all bits to 1
  bs[0] = bs[1] = 0;
       except index 0 and 1
  for (11 i = 2; i <= _sieve_size; i++) if (bs[i]) {
        // cross out multiples of i starting from i * i!
      for (ll j = i * i; j \leftarrow sieve\_size; j \leftarrow i) bs[j] = 0;
    primes.push_back(i); // also add this vector containing list of primes
                                                                                             // call this
      method in main method
bool isPrime(11 N) {
                                                        // a good enough deterministic prime tester
  if (N < _sieve_size) return bs[N];</pre>
                                                                             // O(1) for small primes
  for (int i = 0; i < (int)primes.size(); i++)</pre>
      i f (N % primes[i] == 0) return false;
 return true:
                                                         // it takes longer time if N is a large prime!
                                               // note: only work for N <= (last prime in vi "primes")^2
// second part
vi primeFactors(11 N) { // remember: vi is vector of integers, 11 is long long
 vi factors; // vi 'primes' (generated by sieve) is optional ll PF_idx = 0, PF = primes[PF_idx]; // using PF = 2, 3, 4, ..., is also ok while (N != 1 && (PF + PF <= N)) { // stop at sqrt(N), but N can get smaller
    while (N % PF == 0) { N /= PF; factors.push_back(PF); } // remove this PF
     P F = primes[++PF_idx];
                                                                                             // only
       consider primes!
  if (N != 1) factors.push back(N);
                                               // special case if N is actually a prime
 return factors;
                                    // if pf exceeds 32-bit integer, you have to change vi
```

```
// third part
11 numPF(11 N) {
  11 PF_idx = 0, PF = primes[PF_idx], ans = 0;
 while (N != 1 && (PF * PF <= N)) {
    while (N % PF == 0) { N /= PF; ans++; }
     P F = primes[++PF_idx];
  if (N != 1) ans++;
 return ans;
ll numDiffPF(ll N) {
 11 PF_idx = 0, PF = primes[PF_idx], ans = 0;
while (N != 1 && (PF * PF <= N)) {</pre>
     i f (N % PF == 0) ans++;
                                                                                       // count this pf
       only once
    while (N % PF == 0) N /= PF;
     P F = primes[++PF_idx];
  if (N != 1) ans++;
 return ans;
11 sumPF(11 N) {
 11 PF_idx = 0, PF = primes[PF_idx], ans = 0;
 while (N != 1 && (PF * PF <= N)) {
   while (N % PF == 0) { N /= PF; ans += PF; }
     P F = primes[++PF_idx];
  if (N != 1) ans += N;
 return ans;
ll numDiv(ll N) {
 11 PF_idx = 0, PF = primes[PF_idx], ans = 1;
while (N != 1 && (PF * PF <= N)) {</pre>
                                                                           // start from ans = 1
     1 1 power = 0;
       count the power
    while (N % PF == 0) { N /= PF; power++; }
    ans *= (power + 1);
                                                                                         // according to
       the formula
     P F = primes[++PF_idx];
  if (N != 1) ans *= 2;
                                                // (last factor has pow = 1, we add 1 to it)
 return ans;
ll sumDiv(ll N) {
 11 PF_idx = 0, PF = primes[PF_idx], ans = 1;
                                                                          // start from ans = 1
 while (N != 1 && (PF * PF <= N)) {
     11 power = 0;
    while (N % PF == 0) { N /= PF; power++; }
    ans *= ((11)pow((double)PF, power + 1.0) - 1) / (PF - 1);
                                                                                    // formula
     P F = primes[++PF_idx];
  if (N != 1) ans *= ((11)pow((double)N, 2.0) - 1) / (N - 1);
                                                                                // last one
 return ans:
ll EulerPhi(ll N) {
  11 PF_idx = 0, PF = primes[PF_idx], ans = N;
                                                                           // start from ans = N
 while (N != 1 && (PF * PF <= N)) {
   if (N % PF == 0) ans -= ans / PF;
   while (N % PF == 0) N /= PF;
                                                                          // only count unique factor
     P F = primes[++PF_idx];
  if (N != 1) ans -= ans / N:
                                                                                  // last factor
 return ans;
```

## 5.3 Pollard Rho

```
#include <cstdio>
using namespace std;
#define abs_val(a) (((a)>=0)?(a):-(a))
typedef long long ll;

ll mulmod(ll a, ll b, ll c) { // returns (a * b) % c, and minimize overflow
    ll x = 0, y = a % c;
    while (b > 0) {
        i f (b % 2 == 1) x = (x + y) % c;
        y = (y * 2) % c;
        b /= 2;
```

```
return x % c;
11 gcd(ll a, ll b) { return !b ? a : gcd(b, a % b); }
                                                                   // standard gcd
11 pollard_rho(ll n) {
 int i = 0, k = 2;
11 x = 3, y = 3;
                                                  // random seed = 3, other values possible
 while (1) {
      i ++;
      x = (mulmod(x, x, n) + n - 1) % n;
                                                                       // generating function
     1 1 d = gcd(abs_val(y - x), n);
i f (d != 1 && d != n) return d;
                                                                                   // the key insight
                                                       // found one non-trivial factor
     if (i == k) y = x, k *= 2;
int main() {
  11 n = 2063512844981574047LL;
                                       // we assume that n is not a large prime
                                       // break n into two non trivial factors
  11 ans = pollard_rho(n);
  if (ans > n / ans) ans = n / ans;
                                                   // make ans the smaller factor
 printf("%1ld %1ld\n", ans, n / ans); // should be: 1112041493 1855607779
} // return 0:
```

## 6 Math

## 6.1 Number theory

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std:
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
       return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
       while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b) {
       return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
               b >>= 1:
        return ret:
// returns q = qcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
       int xx = y = 0;
        int yy = x = 1;
        while (b) {
               int q = a / b;
               int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
               t = yy; yy = y - q*yy; y = t;
        return a:
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
```

```
int x, y;
        VI ret,
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                 x = mod(x*(b / g), n);
                 for (int i = 0; i < g; i++)
                         ret.push_back(mod(x + i*(n / g), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
if (g > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = 1 cm (m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make_pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {
    ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);</pre>
                 if (ret.second == -1) break;
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
        if (!a && !b)
                 if (c) return false;
                 x = 0; y = 0;
                 return true;
        if (!a)
                 if (c % b) return false:
                 x = 0; y = c / b;
                 return true;
        if (!b)
                 if (c % a) return false;
                 x = c / a; y = 0;
                 return true;
        int g = gcd(a, b);
        if (c % g) return false;
        x = c / g * mod_inverse(a / g, b / g);

y = (c - a*x) / b;
        return true:
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;
        // expected: 2 -2 1
        int x, y;
int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
        // expected: 95 451
        VI sols = modular_linear_equation_solver(14, 30, 100);
        for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
        cout << endl:
        // expected: 8
        cout << mod_inverse(8, 9) << endl;</pre>
        // expected: 23 105
        PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
        cout << ret.first << " " << ret.second << endl;</pre>
```

```
ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;

// expected: 5 -15
if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;
return 0;</pre>
```

### 6.2 Chinese Remainder Theorem

#### 6.3 Fast Fourier Transform

```
// {{{ FFT Generic
template<typename T>
struct FFTTrait {
  static T getRoot(int, bool) { assert(false); }
template<>
struct FFTTrait<complex<double>> {
  static complex<double> getRoot(int len, bool rev) {
    double ang = 2 * M_PI / len;
    return rev ? complex<double>{cos(ang), sin(ang)} : complex<double>{cos(-ang), sin(-ang)};
};
template<typename T = complex<double>, typename Trait = FFTTrait<T>>
  const static int MAXLEN = 1 << 20;</pre>
  static void reorder(vector<T> &data)
    int n = SZ(data);
    assert(n > 0 && (n & (n - 1)) == 0);
int bitCount = __builtin_ctz(n);
for (int i = 0; i < n; i++) {</pre>
      int j = 0;
      for (int b = 0; b < bitCount; b++) {
        if (i & (1 << b)) {
           j |= 1 << (bitCount - 1 - b);
      if (i < j) {
        swap(data[i], data[j]);
  static void fix(vector<T> &data) {
    int n = 1:
    while (n < SZ(data)) n \neq 2;
    n \star = 2:
    data.resize(n);
  static void fft(vector<T> &data, bool rev) {
    reorder (data);
    int n = SZ(data);
    for (int len = 2; len <= n; len <<= 1) {</pre>
      static T roots[MAXLEN];
      int len2 = len >> 1;
      T w = Trait::getRoot(len, rev);
       roots[0] = T(1);
       for (int i = 1; i < len2; i++) {</pre>
        roots[i] = roots[i - 1] * w; // binary? can be precomputed
      for (int i = 0; i < n; i += len) {
        for (int j = 0; j < len2; j++) {
   T p = data[i + j];</pre>
           T q = roots[j] * data[i + j + len2];
```

## 6.4 Sqrt Mod

```
// Jacobi Symbol (m/n), m, n 0 and n is odd // (m/n)==1 x^2 == m (mod n) solvable, -1 unsolvable #define NEGFOW(e) ((e) % 2 ? -1 : 1)
int jacobi(int a, int m) {
   if (a == 0) return m == 1 ? 1 : 0;
if (a % 2) return NEGPOW((a-1)*(m-1)/4)*jacobi(m%a, a);
   else return NEGPOW((m*m-1)/8)*jacobi(a/2, m);
int invMod(int a, int m) {
   int x, y;
   if (extgcd(a, m, x, y) == 1) return (x + m) % m;
   else
                                       return 0: // unsolvable
// No solution when: n(p-1)/2 = -1 \mod p
int sqrtMod(int n, int p) { //find x: x2 = n (mod p) p is prime
   int S, Q, W, i, m = invMod(n, p);

for (Q = p - 1, S = 0; Q % 2 == 0; Q /= 2, ++S);

do { W = rand() % p; } while (W == 0 || jacobi(W, p) != -1);
   for (int R = powMod(n, (Q+1)/2, p), V = powMod(W, Q, p); ;) {
       int z = R * R * m % p;
       for (i = 0; i < S && z % p != 1; z *= z, ++i);
       if (i == 0) return R;
       R = (R * powMod(V, 1 << (S-i-1), p)) % p;
int powMod (int a, int b, int p) {
         int res = 1;
         while (b)
                   if (b & 1)
                             res = int (res * 111 * a % p), --b;
                    else
                             a = int (a * 111 * a % p), b >>= 1;
         return res;
```

## 6.5 Brent, Miller Rabin

```
#include "../Prime/RabinMiller.h"
long long mul(long long a, long long b, long long mod) {
    if (b == 0) return 0;
    if (b == 1) return a % mod;
    long long mid = mul(a, b >> 1, mod);
    mid = (mid + mid) % mod;
    if (b & 1) return (mid + a) % mod;
    else return mid:
long long brent(long long n) {
    if (n == 1) return 1;
    if (!(n & 1)) return 2;
    if (!(n % 3)) return 3;
    const int p[3] = \{1, 3, 5\};
    long long y, q, x, ys, g, my = 3;
    int i, j, k, m, r, c;
    for (i = 0; i < my; ++i) {</pre>
       y = 1; r = 1; q = 1; m = 111; c = p[i];
            x = y; k = 0;
            for (j = 1; j \le r; ++j) y = (mul(y, y, n) + c) % n;
            do {
                for (j = 1; j \le min(m, r-k); ++j) {
                   y = (mul(y, y, n) + c) % n;
```

# 7 Geometry

## 7.1 Rectangle in Rectangle

```
// Checks if rectangle of sides x, y fits inside one of sides X, Y
// Not tested with doubles but should work fine :)
  Code as written rejects rectangles that just touch.
bool rect_in_rect(int X, int Y, int x, int y) {
   if (Y > X) swap(Y, X);
   if (y > x) swap(y, x);
   double diagonal = sqrt(double(X)*X + double(Y)*Y);
   if (x < X && y < Y) return true;
   else if (y >= Y || x >= diagonal) return false;
   else {
      double w, theta, tMin = PI/4, tMax = PI/2;
      while (tMax - tMin > EPS) {
        theta = (tMax + tMin)/2.0;
         w = (Y-x*cos(theta))/sin(theta);
         if (w < 0 \mid \mid x * sin(theta) + w * cos(theta) < X) tMin = theta;
        else tMax = theta:
     return (w > v);
```

## 7.2 Geometry Basic

```
#define EPS 1e-6
double DEG_to_RAD(double d) { return d * M_PI / 180.0; }
double RAD_to_DEG(double r) { return r * 180.0 / M_PI; }
inline int cmp(double a, double b) {
    return (a < b - EPS) ? -1 : ((a > b + EPS) ? 1 : 0);
struct Point {
    Point (double x = 0.0, double y = 0.0) : x(x), y(y) {}
    Point operator + (Point a) { return Point(x+a.x, y+a.y); }
    Point operator - (Point a) { return Point(x-a.x, y-a.y); }
    Point operator * (double k) { return Point(x*k, y*k); }
    Point operator / (double k) { return Point (x/k, y/k); }
    double operator * (Point a) { return x*a.x + y*a.y; } // dot product
    double operator % (Point a) { return x*a.y - y*a.x; } // cross product
    int cmp(Point q) const { if (int t = ::cmp(x,q.x)) return t; return ::cmp(y,q.y); }
    \#define Comp(x) bool operator x (Point q) const { return cmp(q) x 0; }
    Comp (>) Comp (<) Comp (==) Comp (>=) Comp (<=) Comp (!=)
    #undef Comp
    Point conj() { return Point(x, -y); }
    double norm() { return x*x + y*y; }
    // Note: There are 2 ways for implementing len(): 
// 1. \sqrt{1 - x} fast, but inaccurate (produce some values that are of order \sqrt{2}) 
// 2. \sqrt{2 - x} ow, but much more accurate
    double len() { return sqrt(norm()); }
```

```
Point rotate(double alpha) {
        double cosa = cos(alpha), sina = sin(alpha);
        return Point(x * cosa - y * sina, x * sina + y * cosa);
};
int ccw(Point a, Point b, Point c) {
    return cmp((b-a)%(c-a),0);
double angle (Point a, Point o, Point b) { // angle AOB
    a = a - o; b = b - o;
    return acos((a * b) / sqrt(a.norm() * b.norm()));
// Distance from p to Line ab (closest Point --> c)
double distToLine (Point p, Point a, Point b, Point &c) {
    Point ap = p - a, ab = b - a;
    double u = (ap * ab) / ab.norm();
    c = a + (ab * u);
    return (p-c).len();
// Distance from p to segment ab (closest Point --> c)
double distToLineSegment (Point p, Point a, Point b, Point &c) {
    Point ap = p - a, ab = b - a;
    double u = (ap * ab) / ab.norm();
    if (u < 0.0) {
        c = Point(a.x. a.v);
        return (p - a).len();
    if (u > 1.0) {
        c = Point(b.x, b.y);
        return (p - b).len();
    return distToLine(p, a, b, c);
struct Line {
    double a, b, c;
    Point A, B; // Added for polygon intersect line. Do not rely on assumption that these are valid
    Line(double a, double b, double c) : a(a), b(b), c(c) {}
    Line(Point A, Point B) : A(A), B(B) {
        a = B.v - A.v;
        b = A.x - B.x;
        c = - (a * A.x + b * A.y);
    Line(Point P, double m)
        a = -m; b = 1;
        c = -((a * P.x) + (b * P.y));
    double f(Point A)
        return a*A.x + b*A.v + c:
1:
bool areParallel(Line 11, Line 12) {
    return cmp(11.a*12.b, 11.b*12.a) == 0;
bool areSame(Line 11, Line 12) {
    return areParallel(11 ,12) && cmp(11.c*12.a, 12.c*11.a) == 0;
bool areIntersect(Line 11, Line 12, Point &p) {
    if (areParallel(11, 12)) return false;
    double dx = 11.b*12.c - 12.b*11.c;
   double dy = 11.c*12.a - 12.c*11.a;
double d = 11.a*12.b - 12.a*11.b;
    p = Point(dx/d, dy/d);
    return true;
void closestPoint(Line 1, Point p, Point &ans) {
    if (fabs(1.b) < EPS) {
        ans.x = -(1.c) / 1.a; ans.y = p.y;
        return;
    if (fabs(1.a) < EPS) {
        ans.x = p.x; ans.y = -(1.c) / 1.b;
        return:
    Line perp(l.b, -l.a, - (l.b*p.x - l.a*p.y));
    areIntersect(1, perp, ans);
void reflectionPoint(Line 1, Point p, Point &ans) {
    closestPoint(l, p, b);
```

ans = p + (b - p) \* 2;

## Geometry Circle

```
struct Circle : Point {
    double r;
    Circle(double x = 0, double y = 0, double r = 0) : Point(x, y), r(r) {}
    Circle(Point p, double r) : Point(p), r(r) {}
    bool contains(Point p) { return (*this - p).len() <= r + EPS; }</pre>
// Find common tangents to 2 circles
// Helper method
void tangents(Point c, double r1, double r2, vector<Line> & ans) {
    double r = r2 - r1;
    double z = sqr(c.x) + sqr(c.y);
    double d = z - sqr(r);
    if (d < -EPS) return;
    d = sqrt(fabs(d));
    Line l((c.x * r + c.y * d) / z,
             (c.y * r - c.x * d) / z,
            r1);
    ans.push_back(1);
// Actual method: returns vector containing all common tangents
vector<Line> tangents(Circle a, Circle b) {
  vector<Line> ans; ans.clear();
    for (int i=-1; i<=1; i+=2)
        for (int j=-1; j<=1; j+=2)
    tangents(b-a, a.r*i, b.r*j, ans);</pre>
    for(int i = 0; i < ans.size(); ++i)</pre>
        ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
    vector<Line> ret;
for(int i = 0; i < (int) ans.size(); ++i) {</pre>
        bool ok = true;
        for (int j = 0; j < i; ++j)
             if (areSame(ret[j], ans[i])) {
                 ok = false:
                 break:
        if (ok) ret.push_back(ans[i]);
    return ret:
// Circle & line intersection
vector<Point> intersection(Line 1, Circle cir) {
    double r = cir.r, a = 1.a, b = 1.b, c = 1.c + 1.a*cir.x + 1.b*cir.y;
    vector<Point> res;
    double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if (c*c > r*r*(a*a+b*b)+EPS) return res;
    else if (fabs(c*c - r*r*(a*a+b*b)) < EPS) {
         res.push_back(Point(x0, y0) + Point(cir.x, cir.y));
        return res;
    else {
        double d = r*r - c*c/(a*a+b*b);
        double mult = sqrt (d / (a*a+b*b));
        double ax, ay, bx, by;
        ax = x0 + b * mult;
        bx = x0 - b * mult;
        ay = y0 - a * mult;
        by = y0 + a * mult;
        res.push_back(Point(ax, ay) + Point(cir.x, cir.y));
        res.push_back(Point(bx, by) + Point(cir.x, cir.y));
        return res:
double commonCircleArea(Circle c1, Circle c2) { //return the common area of two circle
    double d = hypot(c1.x-c2.x, c1.y-c2.y), area;
    if (c1.r+c2.r <= d) area = 0;</pre>
    else if (c2.r+d <= c1.r) area = (c1.r * c1.r - c2.r * c2.r) * M_PI;</pre>
    else if (c1.r+d <= c2.r) area = (c2.r * c2.r - c1.r * c1.r) * M_PI;</pre>
        double p1 = c2.r * c2.r * acos((d*d + c2.r*c2.r - c1.r*c1.r) / (2*d*c2.r));
        double p2 = c1.r * c1.r * acos((d*d + c1.r*c1.r - c2.r*c2.r) / (2*d*c1.r));
        double p3 = 0.5 * sqrt((-d+c2.r+c1.r) * (d+c2.r-c1.r) * (d+c1.r-c2.r) * (d+c1.r+c2.r));
        area = p1 + p2 - p3;
    return area;
```

```
// Given 2 circles: (0, 0, r) and (c2.x, c2.y, c2.r)
// Intersections are intersections with line Ax + By + C where
// A = -2 * c2.x
// C = c2.x^2 + c2.y^2 + r^2 - c2.r^2
```

## 7.4 Geometry Polygon

```
typedef vector< Point > Polygon;
// Convex Hull: returns minimum number of vertices. Should work when N <= 1 and when all points
// are collinear
struct comp hull {
    Point pivot:
    bool operator() (Point q, Point w) {
        Point Q = q - pivot, W = w - pivot;

double R = Q % W;
         if (cmp(R,0)) return R < 0;</pre>
        return cmp (Q*Q, W*W) < 0;
Polygon convex_hull(Polygon p) { // minimum vertices
    int j = 0, k, n = p.size();
    Polygon r(n);
    if (!n) return r;
    comp_hull comp;
    comp.pivot = *min_element(p.begin(),p.end());
    sort(p.begin(),p.end(),comp);
for(int i = 0; i < n; ++i) {</pre>
        while (j > 1 && ccw(r[j-1],r[j-2],p[i]) <= 0) j--;
        r[j++] = p[i];
    r.resize(j);
    if (r.size() >= 2 && r.back() == r.front()) r.pop_back();
    return r:
// Area, perimeter, centroid
double signed_area(Polygon p) {
    double area = 0;
    for(int i = 0; i < p.size(); i++) {</pre>
        int j = (i+1) % p.size();
        area += p[i].x*p[j].y - p[j].x*p[i].y;
    return area / 2.0;
double area(const Polygon &p) {
    return fabs(signed_area(p));
Point centroid(Polygon p) {
    Point c(0,0);
double scale = 6.0 * signed_area(p);
for (int i = 0; i < p.size(); i++) {
   int j = (i+1) % p.size();</pre>
        c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
    return c / scale;
double perimeter(Polygon P) {
    double res = 0;
    for(int i = 0; i < P.size(); ++i) {</pre>
        int j = (i + 1) % P.size();
        res += (P[i] - P[j]).len();
    return res;
.
// Is convex: checks if polygon is convex. Assume there are no 3 collinear points
bool is convex(const Polygon &P) {
    int sz = (int) P.size();
    if (sz <= 2) return false;</pre>
    int isLeft = ccw(P[0], P[1], P[2]);
    for (int i = 1; i < sz; i++)
        if (ccw(P[i], P[(i+1) % sz], P[(i+2) % sz]) * isLeft < 0)</pre>
    return true;
// Inside polygon: O(N). Works with any polygon
// Does not work when point is on edge. Should check separately
bool in_polygon(const Polygon &P, Point pt) {
    if ((int)P.size() == 0) return false;
    double sum = 0;
for (int i = 0; i < (int)P.size(); i++) {</pre>
        Point Pj = P[(i+1) % P.size()];
        if (ccw(pt, P[i], Pj) > 0)
             sum += angle(P[i], pt, Pj);
         else sum -= angle(P[i], pt, Pj);
```

```
return fabs(fabs(sum) - 2*M_PI) < EPS;
// Check point in convex polygon, O(logN)
// Source: http://codeforces.com/contest/166/submission/1392387
// On edge --> false
\texttt{\#define} \ \ \texttt{Det} \ (\texttt{a},\texttt{b},\texttt{c}) \quad (\texttt{(double)} \ (\texttt{b}.\texttt{x-a}.\texttt{x}) * (\texttt{double}) \ (\texttt{c}.\texttt{y-a}.\texttt{y}) - (\texttt{double}) \ (\texttt{b}.\texttt{y-a}.\texttt{y}) * (\texttt{c}.\texttt{x-a}.\texttt{x}))
bool in_convex(vector<Point>& 1, Point p){
    int a = 1, b = 1.size()-1, c;
    if (\text{Det}(1[0], 1[a], 1[b]) > 0) swap(a,b); 
// Allow on edge --> if (\text{Det}... > 0 \mid | \text{Det} ... < 0) if (\text{Det}(1[0], 1[a], p) >= 0 \mid | \text{Det}(1[0], 1[b], p) <= 0) return false;
    while (abs(a-b) > 1) {
          c = (a+b)/2:
         if (Det(1[0], 1[c], p) > 0) b = c; else a = c;
     // Alow on edge --> return Det... <= 0
    return Det(1[a], 1[b], p) < 0;</pre>
// Cut a polygon with a line. Returns one half.
// To return the other half, reverse the direction of Line 1 (by negating 1.a, 1.b)
// The line must be formed using 2 points
Polygon polygon_cut(Polygon P, Line 1) {
    Polygon O:
    for(int i = 0; i < P.size(); ++i) {
         Point A = P[i], B = (i == P.size()-1) ? P[0] : P[i+1];
         if (ccw(1.A, 1.B, A) != -1) Q.push_back(A);
         if (ccw(1.A, 1.B, A) *ccw(1.A, 1.B, B) < 0) {
              Point p; areIntersect(Line(A, B), 1, p);
              Q.push_back(p);
     return Q;
// Find intersection of 2 polygons
// Helper method
bool intersect_1pt(Point a, Point b,
    Point c, Point d, Point &r) {
double D = (b - a) % (d - c);
    if (cmp(D, 0) == 0) return false;
    double t = ((c - a) % (d - c)) / D;
    double s = -((a - c) % (b - a)) / D;
    r = a + (b - a) * t;
    return cmp(t, 0) > 0 && cmp(t, 1) < 0 && cmp(s, 0) > 0 && cmp(s, 1) < 0;
Polygon convex_intersect(Polygon P, Polygon Q) {
    const int n = P.size(), m = Q.size();
    int a = 0, b = 0, aa = 0, ba = 0;
    enum { Pin, Qin, Unknown } in = Unknown;
    Polygon R;
    do {
         int a1 = (a+n-1) % n, b1 = (b+m-1) % m;
         double C = (P[a] - P[a1]) % (Q[b] - Q[b1]);
double A = (P[a1] - Q[b]) % (P[a] - Q[b]);
         double B = (Q[b1] - P[a]) % (Q[b] - P[a]);
         if (intersect_1pt(P[a1], P[a], Q[b1], Q[b], r)) {
              if (in == Unknown) aa = ba = 0;
              R.push_back( r );
              in = B > 0 ? Pin : A > 0 ? Qin : in;
         if (C == 0 && B == 0 && A == 0) {
              if (in == Pin) { b = (b + 1) % m; ++ba; }
              else
                                \{ a = (a + 1) % m; ++aa; \}
         } else if (C >= 0) {
             } else {
              if (B > 0) { if (in == Qin) R.push back(Q[b]); b = (b+1)%m; ++ba; }
                            { if (in == Pin) R.push_back(P[a]); a = (a+1)%n; ++aa; }
     while ( (aa < n || ba < m) && aa < 2*n && ba < 2*m );</pre>
    if (in == Unknown) {
         if (in_convex(Q, P[0])) return P;
         if (in_convex(P, Q[0])) return Q;
    return R;
// Find the diameter of polygon.
// Rotating callipers
double convex_diameter(Polygon pt) {
    const int n = pt.size();
    int is = 0, js = 0;
for (int i = 1; i < n; ++i) {
         if (pt[i].y > pt[is].y) is = i;
         if (pt[i].y < pt[js].y) js = i;</pre>
```

```
double maxd = (pt[is]-pt[js]).norm();
    int i, maxi, j, maxj;
    i = maxi = is;
     j = maxj = js;
        int jj = j+1; if (jj == n) jj = 0;
        if ((pt[i] - pt[jj]).norm() > (pt[i] - pt[j]).norm()) j = (j+1) % n;
        else i = (i+1) % n;
        if ((pt[i]-pt[j]).norm() > maxd) {
            maxd = (pt[i]-pt[j]).norm();
            maxi = i; maxj = j;
    } while (i != is || j != js);
    return maxd; /* farthest pair is (maxi, maxj). */
// Closest pair
// Source: e-maxx.ru
#define upd_ans(x, y) {}
#define MAXN 100
double mindist = 1e20; // will be the result
void rec(int 1, int r, Point a[]) {
    if (r - 1 \le 3) {
        for (int i=1; i<=r; ++i)
            for (int j=i+1; j<=r; ++j)</pre>
                   upd_ans(a[i], a[j]);
        sort (a+1, a+r+1); // compare by y
        return:
    int m = (1 + r) >> 1;
    int midx = a[m].x;
    rec(1, m, a), rec(m+1, r, a);
    static Point t[MAXN];
    merge(a+1, a+m+1, a+m+1, a+r+1, t); // compare by y
    copy(t, t+r-1+1, a+1);
    int tsz = 0;
    for (int i=1; i<=r; ++i)
       if (fabs(a[i].x - midx) < mindist) {
    for (int j=tsz-1; j>=0 && a[i].y - t[j].y < mindist; --j)</pre>
                upd_ans(a[i], t[j]);
            t[tsz++] = a[i];
// Pick theorem
// Given non-intersecting polygon.
// I = number of integer points strictly Inside
// B = number of points on sides of polygon
//S = I + B/2 - 1
```

## 7.5 Smallest Enclosing Circle

```
// Smallest enclosing circle:
// Given N points. Find the smallest circle enclosing these points.
// Amortized complexity: O(N)
struct SmallestEnclosingCircle {
   Circle getCircle(vector<Point> points) {
       assert(!points.empty());
        random_shuffle(points.begin(), points.end());
        Circle c(points[0], 0);
       int n = points.size();
       for (int i = 1; i < n; i++)</pre>
            if ((points[i] - c).len() > c.r + EPS)
                c = Circle(points[i], 0);
               for (int j = 0; j < i; j++)
                   if ((points[j] - c).len() > c.r + EPS)
                        c = Circle((points[i] + points[j]) / 2, (points[i] - points[j]).len() / 2);
                        for (int k = 0; k < j; k++)
                            if ((points[k] - c).len() > c.r + EPS)
                                c = getCircumcircle(points[i], points[j], points[k]);
       return c:
    Circle getCircumcircle(Point a, Point b, Point c) {
       double d = 2.0 * (a.x * (b.y - c.y) + b.x * (c.y - a.y) + c.x * (a.y - b.y));
       assert (fabs(d) > EPS);
```

```
double x = (a.norm() * (b.y - c.y) + b.norm() * (c.y - a.y) + c.norm() * (a.y - b.y)) / d;
double y = (a.norm() * (c.x - b.x) + b.norm() * (a.x - c.x) + c.norm() * (b.x - a.x)) / d;
Point p(x, y);
return Circle(p, (p - a).len());
};
```

### 7.6 Geometry Complex

```
// add : a + b
// scalar multi : r * a
// dot product : (conj(a) * b).x
// cross product : (conj(a) * b).y
// squared distance : norm(a - b)
// euclidean distance : abs(a - b)
// angle of elevation : arg(b-a)
// slope of line (a, b) : tan(arg(b-a))
// polar to cartesian : polar(r, theta)
// rotation about the origin : a * polar(1.0, theta) 
// rotation about pivot p : (a - p) * polar(1.0, theta) + p 
// angle ABC: abs(remainder(arg(a - b) - arg(c - b), 2.0 * M_PI)) distance: [-PI, PI]
// project p onto vector v : v * dot(p, v) / norm(v)
// project p onto line(a, b) : a + (b - a) * dot(p - a, b - a) / norm(b - a)
// reflect p across line (a, b) : a + conj((p - a) / (b - a)) * (b - a)
// relect p across line (a, b) : a + conj((p - a) / (b - a)) * (b - a)
// intersection of line(a, b) and (p, q) :
point intersection(point a, point b, point c) {
  double c1 = cross(p - a, b - a), c2 = cross(q - a, b - a);
  return (c1 * q - c2 * p) / (c1 - c2); //undefined if parallel
// read:
template<class T>
istream &operator>>(istream &is, complex<T> &p) {
  T value:
  is >> value:
  p.real(value);
  is >> value;
  p.imag(value);
  return is;
```

## 7.7 Latitude Longitude

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
#include <iostream>
#include <cmath>
using namespace std;
struct 11
  double r, lat, lon;
};
struct rect
  double x, y, z;
};
11 convert (rect& P)
 11 Q;
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
  Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
rect convert(11& Q)
  P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
P.z = Q.r*sin(Q.lat*M_PI/180);
  return P;
```

```
int main()
{
  rect A;
  11 B;

  A.x = -1.0; A.y = 2.0; A.z = -3.0;

  B = convert(A);
  cout << B.r << " " << B.lat << " " << B.lon << endl;

  A = convert(B);
  cout << A.x << " " << A.y << " " << A.z << endl;
}</pre>
```

## 7.8 Delaunay

```
#include <iostream>
#include <cstdio>
#include <algorithm>
#include <vector>
#include <complex>
#include <list>
#include <array>
#include <set>
using namespace std;
typedef long double ld;
typedef complex<ld> Pnt;
typedef array<int, 2> Edge;
typedef array<int, 3> Tri;
typedef pair<Pnt, ld> Circ;
const 1d eps = 1e-6;
int n;
Circ compute_circumcircle(Pnt a, Pnt b, Pnt c) {
 b = a;
 c -= a:
  auto center = (b*norm(c)-c*norm(b)) / (b*conj(c)-c*conj(b));
 return {center+a, norm(center)};
vector<Tri> delaunay(vector<Pnt> pnts) {
 pnts.push_back({-oo, -oo});
  pnts.push_back({oo, -oo});
  pnts.push_back({oo, oo});
  pnts.push_back({-oo, oo});
  list<pair<Tri, Circ>> tcs;
  tcs.push_back({{n, n+1, n+2}, {0, 2*00*00}});
  tcs.push_back({{n+2, n+3, n}, {0, 2*00*00}});
  for (int i = 0; i < n; i++) {
    set Edge> edges;
for (auto it = begin(tcs); it != end(tcs); ) {
  auto circ = it->second;
      ld d = norm(pnts[i] - circ.first) - circ.second;
      if (d < -eps) {
        auto tri = it->first;
        auto j = *prev(end(tri));
        for (int k: tri) {
          auto ite = edges.find({k, j});
          if (ite == end(edges)) {
            edges.insert({j, k});
          } else {
            edges.erase(ite);
          j = k;
        it = tcs.erase(it);
      } else {
        it++:
    int j0 = (*begin(edges))[0];
    for (int j = j0; ; ) {
      int k = (*edges.lower_bound({j, 0}))[1];
      Circ circ = compute_circumcircle(pnts[i], pnts[j], pnts[k]);
      tcs.push_back({{i, j, k}, circ});
      if (j == j0) break;
  vector<Tri> tris;
  for (auto& tc: tcs) {
    auto& tri = tc.first;
    bool flag = false;
    for (int j: tri) {
      flag = flag || j >= n;
```

```
}
if (flag) continue;
tris.push_back(tri);
}
return tris;
```

#### 7.9 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
    INPUT: a vector of input points, unordered.

OUTPUT: a vector of points in the convex hull, counterclockwise, starting
               with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
  BEGIN CUT
#include <map>
// END CUT
using namespace std;
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
  PT() {}
  PT(T x, T y) : x(x), y(y) {}
  bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
  bool operator==(const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE_REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y-b.y)*(c.y-b.y) <= 0);
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {
   if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();</pre>
    dn.push back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
int main() {
 int t;
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {
```

```
int n;
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    map<PT,int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h);
    double len = 0;
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
    printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {</pre>
     if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

### 8 Miscellaneaous

#### 8.1 Calendar Java

```
/*
Constructors:
GregorianCalendar()
GregorianCalendar(int year, int month, int dayOfMonth)
GregorianCalendar(int year, int month, int dayOfMonth, int hour, int minute)
Methods:
add(int field, int amount)
get(int field, int value)
Fields:
GregorianCalendar.YEAR
GregorianCalendar.MONTH
GregorianCalendar.MONTH
GregorianCalendar.DAY_OF_MONTH
GregorianCalendar.DAY_OF_MEEK
GregorianCalendar.DAY_OF_MEEK
GregorianCalendar.DAY_OF_MEEK
GregorianCalendar.DAY_OF_MEEK
```

## 8.2 Dynamic Convex Hull DP (Max)

```
const 11 is_query = -(1LL<<62);</pre>
struct Line {
    mutable function < const Line * () > succ;
    bool operator<(const Line& rhs) const
        if (rhs.b != is_query) return m < rhs.m;</pre>
        const Line* s = succ();
        if (!s) return 0:
        11 \times = rhs.m:
        return b - s->b < (s->m - m) * x;
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum
    bool bad(iterator y) {
        auto z = next(y);
        if (y == begin())
            if (z == end()) return 0;
            return y->m == z->m && y->b <= z->b;
        if (z == end()) return y->m == x->m && y->b <= x->b;
        return (x->b - y->b)*(z->m - y->m) >= (y->b - z->b)*(y->m - x->m);
    void insert_line(l1 m, l1 b) {
        auto y = insert({ m, b });
y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
        if (bad(y)) { erase(y); return; }
        while (next(y) != end() && bad(next(y))) erase(next(y));
```

```
while (y != begin() && bad(prev(y))) erase(prev(y));
}
ll eval(ll x) {
   auto l = *lower_bound((Line) { x, is_query });
   return l.m * x + l.b;
};
```

## 8.3 Floyd Cycle Finding

```
pii floyd(int x0) {
  int tortoise = f(x0), hare = f(f(x0));
  while (tortoise != hare) {
    tortoise = f(tortoise);
    hare = f(f(hare));
  int mu = 0; hare = 0;
  while (tortoise != hare) {
    tortoise = f(tortoise);
   hare = f(hare);
   mu++;
  int lambda = 1;
  hare = f(tortoise);
  while (tortoise != hare) {
   hare = f(hare);
    lambda++;
  return mp (mu, lambda);
```

### 8.4 Josephus

```
/* case k=2 -> if n=2^m+k and 0 <= k < 2^m, then f(n)=2k+1 general case (0-based) -> f(n, k)=(f(n-1, k)+k) \mod n, f(1, k)=0
```

## 8.5 Knight Shortest Path

```
// knight shortest path
int f(int x1, int y1, int x2, int y2)
{
  int dx=abs(x2-x1);
  int dy=abs(y2-y1);
  int lb=(dx+1)/2;
  lb>?=(dy+1)/2;
  lb>?=(dx+dy+2)/3;
  while ((lb%2)!=(dx+dy)%2) lb++;
  if (abs(dx)==1 && dy==0) return 3;
  if (abs(dx)==1 && dx==0) return 3;
  if (abs(dx)==2 && abs(dy)==2) return 4;
  return lb;
}
```

## 8.6 DP Knuth

```
FOR(k = A[i][j-1]..A[i+1][j])
             update(dp[i][j])
// OPTCUT
#include "../template.h"
const int MN = 2011;
int a[MN], dp[MN][MN], C[MN][MN], A[MN][MN];
int main() {
    while (cin >> n) {
         FOR(i,1,n) {
    cin >> a[i];
             a[i] += a[i-1];
         FOR(i,1,n) FOR(j,i,n) C[i][j] = a[j] - a[i-1];
         FOR(i,1,n) dp[i][i] = 0, A[i][i] = i;
         FOR(len,1,n-1)
              FOR(i,1,n-len) {
                  int j = i + len;
dp[i][j] = 2000111000;
                  FOR(k,A[i][j-1],A[i+1][j]) {
                       int cur = dp[i][k-1] + dp[k][j] + C[i][j];
if (cur < dp[i][j]) {
    dp[i][j] = cur;</pre>
                            A[i][j] = k;
         cout << dp[1][n] << endl;
```

## 8.7 C++ Template

```
#ifdef __DEBUG
#define debug if (true)
#define _GLIBCXX_DEBUG
#else
#define debug if (false)
#endif
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
#define pb push_back
#define fi first
#define se second
#define ALL(a) begin(a), end(a)
#define SZ(a) ((int)(a).size())
typedef long long 11;
typedef pair<int, int> pii;
typedef vector<int> vi;
int main() {
  ios::sync_with_stdio(false);
  cin.tie(nullptr);
  return 0;
```

# 8.8 Java Template

```
import java.util.*;
import java.io.*;

class Main {
   FastScanner in;
   PrintWriter out;

public void solve() throws IOException {
   }

public void run() {
   try {
```

```
in = new FastScanner(System.in);
    out = new PrintWriter(System.out);
    solve();
    out.close();
  } catch (IOException e) {
   e.printStackTrace();
class FastScanner {
 BufferedReader br;
  StringTokenizer st;
  FastScanner(InputStream is) {
   br = new BufferedReader(new InputStreamReader(is));
  FastScanner(File f) {
     br = new BufferedReader(new FileReader(f));
    } catch (FileNotFoundException e) {
      e.printStackTrace();
  String next() {
   while (st == null || !st.hasMoreTokens()) {
     try {
        st = new StringTokenizer(br.readLine());
     } catch (IOException e) {
       e.printStackTrace();
    return st.nextToken();
 int nextInt() {
   return Integer.parseInt(next());
public static void main(String[] arg) {
 new Main().run();
```

### 8.9 pika

## 9 SHK

## 9.1 SHK C++

```
#include <bits/stdc++.h>
using namespace std;
const double PI = acos(-1);
#define INF IE9
#define endl '\n'
#define pb push_back
// A lot of typedefs
typedef long long ll;
```

```
// struct
   struct mystruct {
       int counter;
      };
  //custom hashing
   struct custom_hash {
      inline std::size_t operator()(const std::pair<int,int> & v) const {
            return v.first*31+v.second*7;
      };
  // pq/set custom comparator, will get reversed (at least in pq)
   class mycomp {
      public:
       bool operator() (mystruct a, mystruct b) {
         return a.counter > b.counter;
  // sort custom comparator
    bool customcompare(mystruct a, mystruct b) {
      return a.counter > b.counter;
int main ()
  ios::sync_with_stdio(false);
 cin.tie(nullptr);
 return 0;
```

## 9.2 SHK Python

```
import sys
class SamInput (object):
 def __init__(self):
    self.inp = []
     for i in sys.stdin:
         i = i.replace("\n", "")
         j = list(i.split())
       self.inp.append(j)
 def readln(self):
     i f (len(self.inp) == 0) :
      return False
    else:
      return str.join(" ", self.inp.pop(0))
  def read(self):
     i f (len(self.inp) == 0):
      return False
    while (len(self.inp[0]) == 0):
       self.inp.pop(0)
        i f (len(self.inp) == 0):
         return False
   return self.inp[0].pop(0)
```

### 9.3 SHK Shortcuts

```
Comment highlighted code
Ctrl + Shift + C
Uncomment highlighted code
Ctrl + Shift + X

Indent block.
Tab
Dedent block.
Shift + Tab

Line cut.
Ctrl + L
Line copy.
Ctrl + Shift + T

g++ x.cpp -std=c++11 -o x
x.exe < test.txt
```