irry+CSE 232 Fall 2018

# **Programming Project #10**

## **Assignment Overview**

In this assignment you will continue your practice in creating a custom data structure, BiMap. We change the specifications so that you cannot use a STL vector nor pair (as well as not use map or set or the sort algorithm). You will update your class to use **dynamic arrays** instead of STL vectors. It is due 11/28. That's a **Wednesday** because of the Thanksgiving holidays.

### **Background**

We are going to stay with the BiMap problem but update its implementation. We are going to stop using vectors and pairs and do the work ourselves with underlying dynamic arrays. We will also templatize the BiMap to work with different types.

We do this as a better way to get familiar with the underlying problem of templated classes and dynamic arrays. The specs of the BiMap should be familiar to you now, so it is a matter of changing the implementation. The focus will be on this implementation.

#### Details

As this is a templated class you can <u>only</u> write a header file to solve the problem. Thus, we will provide a proto-header file, proj10\_bimap\_starter.h, which you will fill in. In the specifications below, we will focus on what is different first for this vs. the previous BiMap, then update the rest of the already familiar set of methods.

### Class Node

Since we cannot use the STL pair, we are going to have to write our own struct called Node. This will take the place of the pair from the STL. Here is the header part of Node:

```
template<typename K, typename V>
    struct Node {
    K first;
    V second;
    Node() = default;
    Node(K,V);
    bool operator<(const Node&) const;
    bool operator==(const Node&) const;
    friend ostream& operator<<(ostream &out, const Node &n) {
        // YOUR CODE HERE
    }
};</pre>
```

First note it is doubly templated on the Key (K) and the Value (V). Conveniently we name the data members first and second to be compatible with std::pair. You are required to write in this header file:

- 2 arg constructor
- operator< Compares two Node instances based on the first values (so that, if we need to compare using something like lower\_bound, they would be ordered by key).
- operator== Similarly, two Node instances are equal if their two first values are equal (can't have duplicate keys).

• operator<< As discussed in the videos, and unlike the other definitions (which should appear below the struct), the easiest way to do operator<< for a templated class is to place the definition <u>right here</u> in the class section (not below where the other methods would go). It should print "first:second" to the ostream, whatever the values of first and second are.

### Class BiMap

The underlying data member for the storage of Node is an array (not a vector, a basic array) of type Node. Since this is not a vector we have to grow the vector when we add more Node elements than the array can presently store. The array has to grow dynamically during the course of the run. As before, the elements of the array must be in sorted order of the Nodes, using either the key or the value to order the two arrays named ordered by keys or ordered by vals, as we did before.

As this pertains to the BiMap class, where we previously had the two vectors, we now have two Node arrays ordered\_by\_keys\_ and ordered\_by\_vals\_. These are to be ordered and treated in the same way as the previous project – only now we are managing the dynamic memory ourselves. The header portion of the BiMap class should appear as follows:

```
template<typename K, typename V>
class BiMap {
private:
     Node<K,V>* ordered by keys ;
     Node<K,V>* ordered_by_vals_;
     size t last;
     size t capacity;
     Node<K,V>* find key(K key);
     Node<K,V>* find value(V value);
     void grow ();
public:
     BiMap(int sz = 2);
     BiMap(initializer list< Node<K, V> >);
     BiMap (const BiMap&);
     BiMap operator=(BiMap);
     ~BiMap();
     size t size();
     V value from key(K);
     K key from value(V);
     bool update(K, V);
     K remove val(V value);
     V remove key(K key);
     bool add(K, V);
     int compare(BiMap &);
     friend ostream& operator<<(ostream &out, const BiMap &ms) {
        // YOUR CODE HERE
     }
};
```

## **Data Members new to Project 10**

- capacity\_ the actual size of the underlying arrays. It is the number of the Node elements the arrays can hold before they must grow.
- last\_- the index of the first open (unusued) location in the arrays. When last\_ == capacity then the arrays must grow.
  - O Note that these values are the same for both arrays, and will remain so, since the data contained in each is the same, only with a different ordering

### Methods new to Project 10

- 1 arg constructor (default value for size). Creates an array of size sz (the parameter) of type Node using the new operator. The sz parameter also is the is the capacity\_ of the MapSet. last is set to 0 (the first open index).
- grow. This method doubles the capacity of the underlying array. It:
  - o create new arrays of twice the capacity
  - o copies the content of the old arrays into the new arrays
  - o update capacity\_ and the two array pointers: ordered\_by\_keys\_ and ordered by vals and then delete the old arrays
- copy constructor
- operator=
- destructor

These are all methods that we have ignored previously because we took the defaults for those methods. That worked because their operations were taken care of by the underlying STL elements. Now you must provide them yourself. You have to <u>look at the example code</u>, videos and notes to create these. In fact, **copy the example code and modify to suit your needs here** (you won't get accused of cheating for this, I promise).

The destructor is tested in that memory leaks are tested (Test 35). Other than that these are tested implicity in the code but YOU should test them yourself

- K remove val(V value)
- V remove key(K key)

Since we no longer have the same type for the keys and values, we have to break up the remove operator into two separate functions, to either remove the Node with that value from the arrays, and return the associated key, or remove the Node with that key and return the associated value. If the key or value is not present, a default value or key should be returned, respectively

## **Methods modified for Project 10**

- Node<K,V>\* find key(K key)
- Node<K,V>\* find value(V value)
- These both serve the same functionality as in the previous project, but return a Node<K, V>\* instead of an iterator. You can still use lower\_bound if you do pointer arithmetic, though the begin() and end() functions won't work as the array size is not fixed at compile time. The pointer is either the first Node in the array that is equal to (by key/value) or greater than the key/value, or nullptr (the last two meaning that the key isn't in the arrays). It must be private because the arrays are private and we cannot return a pointer to private data.

To make lower\_bound work, you are either going to have to write a function or a lambda that compares a Node<K, V> and a K or a V

These functions are <u>not tested</u> in the Mimir test set but necessary everywhere. However, it is essentially the use of lower bound. It is a good to isolate this usage however for future projects.

- BiMap(initializer\_list< Node<K,V>>): Create an array of size initializer\_list.size(). Take each Node and place in the vectors. The initializer\_list does not have to be in sorted order but the vector should be after you add all the elements. (Hint: write add first and use it here)
- size t size(): returns the size of the BiMap (number of Nodes) as an unsigned int
- K key\_from\_value (V value): member function. Return the key associated with the value. If the value does not exist, then return a default K
- V value\_from\_key(K key): member function. Return the value associated with the key. If the key does not exist, then return a default V
- bool update (K, V): if the key is in the BiMap, and the value is not, update the key-value Node to the parameter value. Return true. If the key is not in BiMap, do nothing and return false.
- bool add (K, V): if the key or the value are in the BiMap, do nothing and return false. Otherwise create a Node with the argument values and insert the new Node into the vectors, in sorted order, and return true.
- bool compare (BiMap&): compare the two BiMaps lexicographically, that is element by element using the string-key of the Nodes as comparison values. If you compare two Nodes, then the comparison is based on the .first of each Node. The first difference that occurs determines the compare result. If the calling BiMap is greater, return 1. If the argument BiMap is greater, return -1. If all of the comparable Nodes are equal but one BiMap is bigger (has more Nodes), then the longer determines the return value (1 if the first is longer, -1 if the second).
- friend ostream& operator<< (ostream&, BiMap&). Returns the ostream after writing the BiMap to the ostream. The formatting should have each Node colon (':') separated, and each Node comma + space separated (', '), with no trailing comma.
  - o E.g., Ann: ABCD, Bob: EFGH, Charlie: IJKL

#### **Requirements**

We provide proj10\_bimap\_starter.h as a skeleton that you must fill in. You submit the completed proj10 bimap.h to Mimir

We will test your files using Mimir, as always.

### **Deliverables**

proj10/proj10 bimap.h

- 1. Remember to include your section, the date, project number and comments.
- 2. Please be sure to use the specified directory and file name.

### **Assignment Notes**

### lower bound

You should get how lower\_bound works now. The difference here is now you have no iterators to work with. You must instead use pointer arithmetic.

```
lower_bound(ary_, ary_+last_, value_to_search_for)
```

The return value is a pointer (not an iterator) to the either the element in the container that meets the criteria, or the value of the last element in the range searched (in this case, nullptr)

#### That means that either:

- the value to search for is already in the container and the iterator points to it.
- value to search for is not in the container. Not in the container means:
  - o the iterator points to a value "just greater" than the value to search for
  - o the iterator points to ayr +last

## array insert or erase

Bad news, there is no insert nor erase. You can write these separately or put them in add and remove (respectively). To make these work, you are going to have to remake the arrays.

- for insert: make a new array that has: copy of the old array up to to insert point + new element + the old array after the insert point.
  - o if you made an array with new, better delete it or you will leak!
- for erase: make a new array that has: copy of old array up to the remove point, then copy after the old array to the end. In other words, skip the element being erased in the copy.

#### add

The critical method is add. Get that right first and them much of the rest is easy. For example, the initializer list constructor can then use add to put elements into the vector at the correct location (in sorted order).

#### sort

As before, no use of sort allowed. If you use sort in a test case you will get 0 for that test case. Do a combination of lower bound and vector insert to get an element where it needs to be in a vector.