# Relational Algebra Tutorial

**SWEN304/SWEN439** 

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**Engineering and Computer Science** 





- Unary Operation: Select, Project, Rename
- Binary Operation: Join, Cartesian Product, Outer Join, Union, Interaction, Difference
- Relational algebra exercises

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## **Unary Operations**

- Project:  $\pi_{AL}(N)$ 
  - Example:  $\pi_{LName, FName}(Student)$
- Select:  $\sigma_c(N)$ 
  - Example:  $\sigma_{FName = 'Susan'}$  (Student)
- Rename:  $\delta_{AI \rightarrow BI,...,Ak \rightarrow Bk}(N)$ 
  - Example:  $\delta_{\text{FName} \rightarrow \text{FirstName}, \text{LName} \rightarrow \text{LastName}}$ (Student)



## **Binary Operations**

- Union:  $N_1 \cup N_2$ 
  - Example:  $\pi_{StudId}(Student) \cup \pi_{StudId}(Grades)$
- Interaction:  $N_1 \cap N_2$ 
  - Example:  $\pi_{StudId}(Student) \cap \pi_{StudId}(Grades)$
- Difference:  $N_1$   $N_2$ 
  - Example:  $\pi_{StudId}(Student) \pi_{StudId}(Grades)$



### Binary Operation: Join Operations

- Join operation joins two relations by merging those tuples from two relations that satisfy a given condition
  - The condition is defined on attributes belonging to relations to be joined
- Equijoin, natural join operations



## **Equijoin Operation**

• Notation:  $N = N_1 \bowtie_C N_2$ where  $JC = jc_1 \wedge ... \wedge jc_n$  $jc_i \equiv A = B, A \in R_1, B \in R_2,$ 

For example,

Student ⋈ StudId = StudId Grades

In SQL:

SELECT \*

FROM Student s, Grades g WHERE s.StudId = g.StudId;



# Equijoin

Equijoin:  $N_1 \bowtie_{N1.B=N2.B} N_2$ 

$N_{I}$		
А	В	
1	2	
3	3	
4	4	

$$N1.B = N2.B$$

	В	C
2	2	7
	4	9
	$\omega$	0

 $N_2$ 

A	В	В	С
1	2	2	7
4	4	4	9

Natural Join :  $N = N_1 * N_2$ 

$N_1$		_	$N_{2}$	2		Ì	V	
Α	В	*	В	С		A	В	C
1	2	,,	2	7	=	1	2	7
3	3		4	9		4	4	9
4	4		$\omega$	0				



## **Cartesian Product**

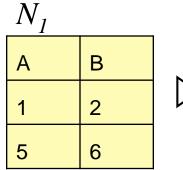
$N_1$			$N_2$		
Α	В		В	C	
1	2	×	2	7	=
3	3		4	9	
4	4		$\omega$	0	
			-		

$N_1 \times N_2$				
Α	В	В	C	
1	2	2	7	
1	2	4	9	
1	2	$\omega$	0	
3	3	2	7	
3	3	4	9	
3	3	$\omega$	0	
4	4	2	7	
4	4	4	9	
4	4	$\omega$	0	



#### **Outer Join**

#### Right Outer Join



$$N_2$$
 $C$ 
 $D$ 
 $2$ 
 $7$ 
 $2$ 
 $9$ 
 $\omega$ 
 $7$ 

$$N_1 \bowtie_{\mathsf{B}=\mathsf{C}} N_2$$

Α	В	С	D
1	2	2	7
1	2	2	9
ω	ω	ω	7

#### Full Outer Join

$$\begin{array}{c|cccc}
N_I \\
\hline
A & B \\
1 & 2 \\
\hline
5 & 6 \\
\end{array}$$

$N_1 \supset I_{B=C} N_2$			
А	В	С	D
1	2	2	7
1	2	2	9
5	6	ω	ω
ω	ω	ω	7



#### **Summary or Relational Operations**

- SELECT  $\sigma_c(N)$  : choose rows
- PROJECT  $\pi_{A1,...,Ak}(N)$ : choose columns
- RENAME  $\delta_{AI \to BI,...,Ak \to Bk}(N)$ : rename attributes
- JOIN: combine tables
  - Natural Join  $N_1 * N_2$  or
  - Equi-Join  $N_1 \bowtie_{A1=B1,...,Ak=Bk} N_2$
- CARTESIAN PRODUCT ( x ): combine tables
- Set operations
  - UNION ( ∪ ),
  - INTERSECTION ( ∩ ),
  - DIFFERENCE (or MINUS, )
- Additional Relational Operations
  - OUTER JOINS



# A Sample Relational Database

#### **Student**

LName	FName	StudId	Major
Smith	Susan	131313	Comp
Bond	James	007007	Math
Smith	Susan	555555	Comp
Cecil	John	010101	Math

#### **Course**

CName	CourId	Points	Dept
DB Sys	C302	15	Comp
SofEng	C301	15	Comp
DisMat	M214	22	Math
Pr&Sys	C201	22	Comp

#### **Grades**

StudId	CourId	Grade
007007	C302	A+
555555	C302	ω
007007	C301	Α
007007	M214	A+
131313	C201	B-
555555	C201	С
131313	C302	ω
007007	C201	Α
010101	C201	ω

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#### **Exercises**

- Suppose we are given the university database instance as in slide 9. Write queries in relational algebra for the following queries
  - 1. Find all students with their ID who got at least one 'A+'
  - 2. Find students with their ID, FName, who have enrolled in C302
  - 3. Find students with their IDs who have enrolled in 'C201' but not 'C302'
  - 4. Find students who have enrolled in both 'M214' and 'C302'
  - 5. Find students who have neither enrolled in 'M214' nor in 'C302'
  - 6. Find students who major in 'Math' and got 'A+' in at least one course offered by computer science department