

CS323 Project1 Report

Haotian Liu@SUSTech 11613015

Project contents: Lexical Analyzer & Syntax analyzer for SUSTech programming language (SPL)

1. Write Regular Expression using flex to recognize tokens
2. Write Context-Free Grammar (CFG) using bison to specify the SPL's grammar

Data Structure: The tree-structured AST with siblings linked together.

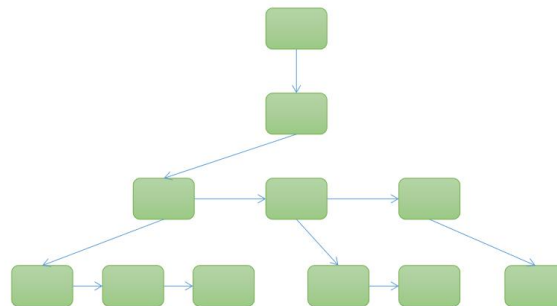
Each node in the tree is defined as (see [/ast.c](#))

```
struct ast {  
    int lineno;  
    char *name;  
    char *value;  
    struct ast *sub_ast;  
    struct ast *next_sibling;  
};
```

The [lineno](#) corresponds to the line number, the [name](#) to the token name, all of the token's values are stored as the string directly in [value](#). The [sub_ast](#) points to the child node lower one level while [next_sibling](#) points to the node at the same level.

In this project, all of the tokens and grammar units are defined as ast.

The AST nodes are linked as below:



Function Implementations (see [/ast.c](#))

```
struct ast *new_node (char *name, char*value, int lineno);
```

When a token is recognized by the regular expression defined in flex file (see [/lex.l](#)), the function [new_node](#) will be invoked to generate a new ast with the token's name, value and the line number of it.

```
struct ast *new_ast (char *name, int num, ...);
```

When a grammar in SPL is matched the CFG defined in bison file (see [/syntax.y](#)), the function [new_ast](#) will be invoked to connect the nodes and generate a new AST. The [name](#) represents the name grammar unit, the [num](#) is the number of children, and also the variable number of the variable argument list, where variables will be parsed as ast's.

```
void preorder (struct ast *ast, int level);
```

Traverse the AST in pre-order way with function [preorder](#), print the expected result.

Optional Features

- ✓ Support single-line comments.

Realization: when recognizes the “//” pattern, using input() function to read characters from the input buffer without handling. That is, discard all characters behind “//” in the line by the following code:

```
//" { char c; while((c=input()) != '\n'); }
```

- ✓ Support multi-line comments and detect nested multi-line comments.

Realization: when recognizes the “/*” pattern, read and discard the characters until meet “*/”.

The code is as below:

```
/*" {  
    char c;  
    while (1) {  
        if ((c=input()) == '*') {  
            if ((c=input()) == '/') {  
                break;  
            }  
            else {unput(c);}  
        }  
    }  
    if ((c=input()) == '/') {  
        if ((c=input()) == '*') {  
            Raise an syntax error;  
        }  
        else {unput(c);}  
    }  
}  
    }  
    /*/" {  
        Raise an syntax error;  
    }  
}
```

I directly raise the syntax error in **lex.l** file for such implementation is simple and the additional syntax unit can be omitted.

If “/*” or “*/” are contained in pair of double-quotes, they will be recognized as string, whose matching length is longer.

- ✓ Support hexadecimal representation and detect illegal form of hex-int.

Define the regular expression below

```
HEX_INT [-+]?0[xX]([0-9a-f]{0,8})
```

```
HEX_INT_WRONG [-+]?0[xX][0-9a-zA-Z]*
```

And processing HEX_INT before HEX_INT_WRONG, when a correct hexadecimal is recognized, although both of the pattern can match it, the former one will be returned.

When a wrong hex-int or hex-int more than 32bits detected, the latter one will be returned.

- ✓ Support lex-form characters, and detect illegal form like ‘\xt0’

```
CHAR ('[a-zA-Z0-9]')|('\x[0-9a-fA-F]{2}')
```

```
CHAR_WRONG ('\x[0-9a-zA-Z]*')
```

Similar as hexadecimal representation detection written above.