

Lecture Notes

Advanced Discrete Structures

COT 4115.001 S15

2015-01-14

Recap

- Extended Euclidean Algorithm
- Solving $ax + by = d$
- Arithmetic Modulo n
- Finding $a^{-1} \pmod{n}$

Chapter 2

CLASSICAL CRYPTOSYSTEMS

Conventions

- *plaintext* will be written in lowercase
- *CIPHERTEXT* will be written in uppercase
- Letters of the alphabet are assigned numbers:

a	b	c	d	e	f	g	h	i	j
0	1	2	3	4	5	6	7	8	9
k	l	m	n	o	p	q	r	s	t
10	11	12	13	14	15	16	17	18	19
u	v	w	x	y	z				
20	21	22	23	24	25				

What to do about spaces?

- Spaces are frequent
 - They reveal information about word structure
 - If spaces encrypted, easy to identify by frequency giving the ability to crack key.



Classical Cryptosystems - Section 2.1

SHIFT CIPHERS

Shift ``Caesar'' Cipher (~45BC)

Example: Shift each letter to the right by three.

Original:

gaul is divided into three parts

By Convention: gaulisdividedintothreeparts

Shifted by 3: JDXOLVGLYLGHGLQWRWKUHHSDUWV

Mathematically:

$$x \mapsto x + \kappa \pmod{26}$$

Encryption

$$x \mapsto x - \kappa \pmod{26}$$

Decryption

Attacks on Shift Cipher

1. Ciphertext Only:

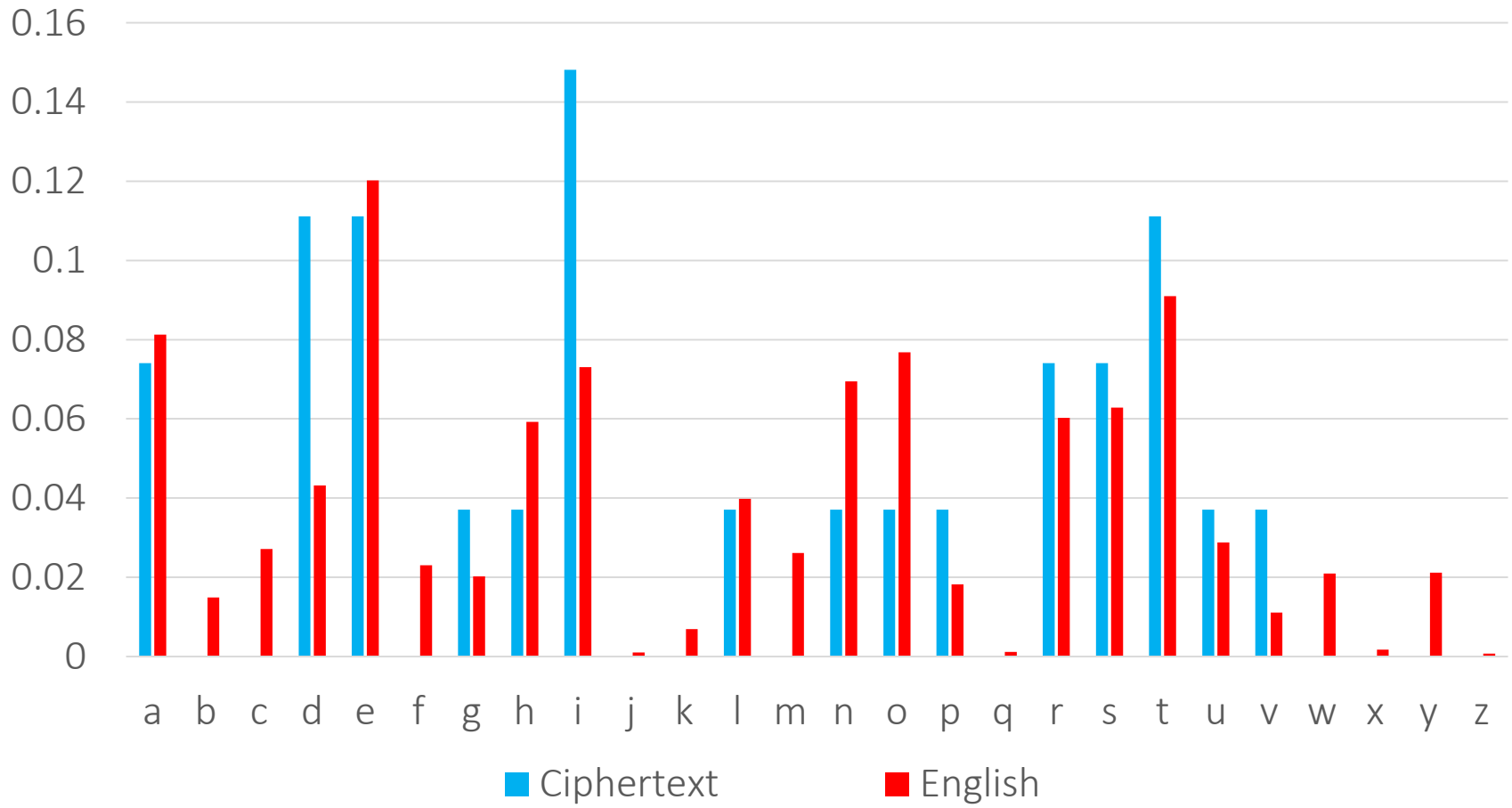
— Brute force it!

— Frequency analysis

0: gaulisdividedintothreeparts
1: hbvmjtejwjefejoupuisffqbsut
2: icwnkufkxkfgfkpvqvjtggrectvu
3: jdxolvgllylghglqwrwkuhhsduwv
4: keypmwhmzmhihmrxxslviitevwx
5: lfzqnxinanijinsytymwjjufwyx
6: mgaroyjobojkjotzuznxkkvgxzy
7: nhbspzkpcpklkpuavaoyllwhyaz
8: oictqalqdqlmlqvbwbpzmmxizba
9: pjdurbmrmrmnrmwxcxcqannyjacb
10: qkevscnsfsnonsxdydrboozkbdc
11: rlfwtdotgtopotyezescppalced
12: smgxuepuhupqpuzfaftdqgbmdfe

13: tnhyvfqvivqrqvagbguerrcnegf
14: uoizwgrwjwrsrwbhchvfssdofhg
15: vpjaxhsxkxstscidiwgttepgih
16: wqkbyitylytutydjejxhuufqhji
17: xrlczjuzmzuvuzekfkyivvgrikj
18: ysmdakvanavwvaflglzjwwhsjlk
19: ztneblwbobwxwbgmhmakxxitkml
20: auofcmxcpcxyxchninblyyjulnm
21: bvpgdnydqdyzydiojocmzzkvmon
22: cwqheozerezazejpkpdnaalwnpo
23: dxrifpafsfabafkqlqeobbmxoqp
24: eysjgqbggtgbcbglrmpccnyprq
25: fztkhrchuhcdchmsnsgqddozqsr

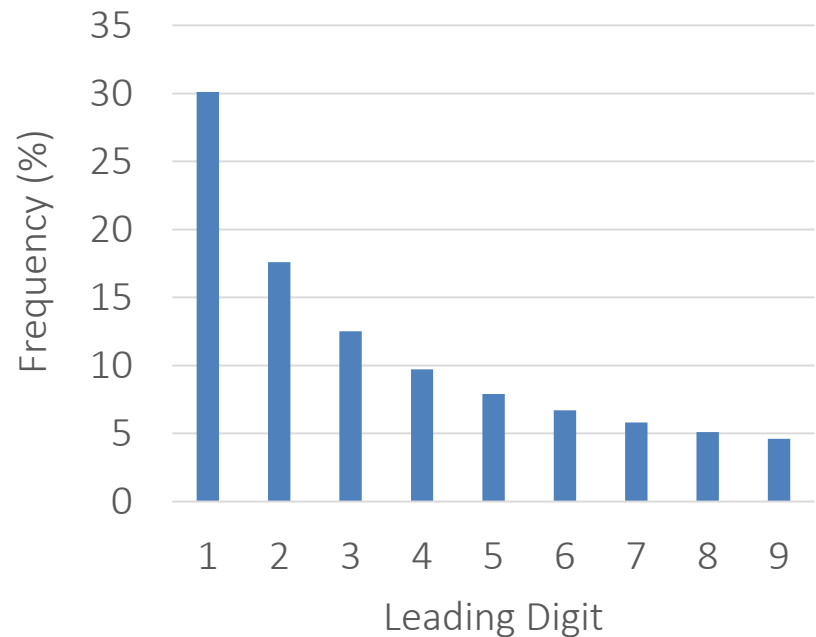
English Language (Frequency Analysis)



Benford's Law

- **Benford's Law**, also called the **First-Digit Law**, refers to the frequency distribution in many (but not all) real-life sources of data.

Leading Digit	Occurring Frequency
1	30.1%
2	17.6%
3	12.5%
4	9.7%
5	7.9%
6	6.7%
7	5.8%
8	5.1%
9	4.6%



Attacks on Shift Cipher

2. Known plaintext:

- If you know one letter, you can decipher the rest

Example:

$$d (= 23) \mapsto N (= 13), \quad \kappa \equiv 13 - 23 \equiv -10 \equiv 16 \pmod{26}$$

3. Chosen plaintext:

- Encrypt the letter a and compute the difference.

4. Chosen ciphertext:

- Decrypt the letter A and compute the difference.

Classical Cryptosystems - Section 2.2

AFFINE CIPHERS

Affine Ciphers

Shift Cipher:

$$x \mapsto x + \kappa \pmod{26}$$

for some $\kappa \in \mathbb{Z}$.

Affine Cipher:

$$x \mapsto \alpha x + \beta \pmod{26}$$

for some $\alpha, \beta \in \mathbb{Z}$ with $\gcd(\alpha, 26) = 1$.

Affine Ciphers

Example: $\alpha = 3, \beta = 5$ (Check: $\gcd(3, 26) = 1$ ✓)

	m	a	t	h	i	s	f	u	n	
	12	0	19	7	8	18	5	20	13	
x3	10	0	5	21	24	2	15	8	13	(mod 26)
+5	15	5	10	0	3	7	20	13	18	
	P	F	K	A	D	H	U	N	S	

$$3x + 5 \equiv X \pmod{26}$$

$$x \equiv 3^{-1}(X - 5) \pmod{26}$$

$$3^{-1} \equiv 9 \pmod{26}$$

$$x \equiv 9X + 7 \pmod{26}$$

Affine Ciphers

Example:

$$x \equiv 9X + 7 \pmod{26}$$

	m	a	t	h	i	s	f	u	n	
	12	0	19	7	8	18	5	20	13	
x3	10	0	5	21	24	2	15	8	13	(mod 26)
+5	15	5	10	0	3	7	20	13	18	
	P	F	K	A	D	H	U	N	S	
	15	5	10	0	3	7	20	13	18	
x9	5	19	12	0	1	11	24	13	6	(mod 26)
+7	12	0	19	7	8	18	5	20	13	
	m	a	t	h	i	s	f	u	n	

Encryption Key: $(\alpha, \beta) = (3, 5)$ Decryption Key: $(\alpha', \beta') = (9, 7)$

Affine Ciphers

What can happen if $\gcd(\alpha, 26) > 1$?

Example:

$$x \mapsto 13x + 4$$

$$\gcd(13, 26) = 13$$

input \mapsto ERRER

alter \mapsto ERRER

How many possible pairs (α, β) with $\gcd(\alpha, 26) = 1$?

- $2 \mid \gcd(2k, 26)$ leaves 13 odd numbers
- $13 \mid \gcd(13k, 26)$ leaves 12 possibilities for α
- $12 \times 26 = \mathbf{312}$ choices for the key (α, β)

Attacks on Affine Ciphers

1. Ciphertext only:

- Brute force all 312 keys!
- Frequency Analysis

2. Known plaintext:

- Knowing two letters allows you to find the key:

$$X \equiv \alpha x + \beta \pmod{26} \quad \text{and} \quad Y \equiv \alpha y + \beta \pmod{26}$$

$$\alpha (x - y) \equiv X - Y \pmod{26}$$

If $\gcd(x - y, 26) = d > 1$, then $d \mid X - Y$ first divide both $x - y$ and $X - Y$ by d . Then:

$$\alpha \equiv (X - Y)(x - y)^{-1} \pmod{26}$$

$$\beta \equiv X - \alpha x \pmod{26}$$

Attacks on Affine Ciphers

Example:

$$\begin{array}{ll} c \mapsto \mathbb{Z} & \text{and} \quad e \mapsto \mathbb{M} \\ 2 \mapsto 25 & 15 \mapsto 12 \end{array}$$

$$25 \equiv 2\alpha + \beta \pmod{26} \quad \text{and} \quad 12 \equiv 15\alpha + \beta \pmod{26}$$

$$13 \equiv -13\alpha \pmod{26}$$

- No multiplicative inverse, $(-13)^{-1}$, modulo n .
- Note that: $\gcd(-13, 26) = 13$.
- Divide both sides by 13 and solve:

$$1 \equiv -1 \cdot \alpha \pmod{26} \quad \text{and} \quad (-1)^{-1} \equiv -1 \pmod{26}$$

$$\alpha \equiv -1 \cdot -1 \equiv 1 \pmod{26}$$

$$\beta \equiv 25 - 2 \cdot 1 \equiv 23 \pmod{26}$$

- Check:

$$25 \equiv 2 \cdot 1 + 23 \pmod{26} \quad \text{and} \quad 12 \equiv 15 \cdot 1 + 23 \pmod{26} \quad \checkmark$$

Attacks on Affine Ciphers

Chosen plaintext:

- Choose ab ($=01$) as the plaintext. This gives:

$$X \equiv \alpha \cdot 0 + \beta = \beta \pmod{26} \implies \beta \equiv X \pmod{26}$$

$$Y \equiv \alpha \cdot 1 + \beta \pmod{26} \implies \alpha \equiv Y - \beta \equiv Y - X \pmod{26}$$

$$\text{Encryption: } (Y - X, X) \qquad \text{Decryption: } ((Y - X)^{-1}, -(Y - X)^{-1}X)$$

Chosen ciphertext:

- Choose AB ($=01$) as the ciphertext. This gives:

$$0 \equiv \alpha \cdot x + \beta \pmod{26}$$

$$1 \equiv \alpha \cdot y + \beta \equiv \alpha (y - x) \pmod{26} \implies \alpha \equiv (y - x)^{-1} \pmod{26}$$

$$\beta \equiv -\alpha \cdot x \equiv -(y - x)^{-1}x \pmod{26}$$

$$\text{Encryption: } ((y - x)^{-1}, -(y - x)^{-1}x) \qquad \text{Decryption: } (y - x, x)$$

Classical Cryptosystems - Section 2.3

THE VIGENÈRE CIPHER

Vigenère Cipher Encryption

1. Choose a random key:

`key: bacon`

2. Match message to key repeats:

`plaintext: whatdoesthefoxsay`

`key: baconbaconbaconba`

3. Add message and key to get ciphertext

Vigenère Cipher En(De)ryption

w	h	a	t	d	o	e	s	t	h	e	f	o	x	s	a	y
22	7	0	19	3	14	4	18	19	7	4	5	14	23	18	0	24
+																
b	a	c	o	n	b	a	c	o	n	b	a	c	o	n	b	a
1	0	2	14	13	1	0	2	14	13	1	0	2	14	13	1	0
<hr/>																
23	7	2	7	16	15	4	20	7	20	5	5	16	11	5	1	24
X	H	C	H	Q	P	E	U	H	U	F	F	Q	L	F	B	Y

To decrypt, subtract key in the say fashion.

Attacks on the Vigenère Cipher

- Plaintext attack
 - If enough characters are known, the key is obtained.
- Chosen plaintext
 - Choose the plaintext: `aaaa...` . Ciphertext is the key
- Chosen ciphertext
 - Choose the ciphertext `AAAA...` . Plaintext is negative of the key

Attacks on the Vigenère Cipher

- Ciphertext only

- Does frequency analysis work?

- In the example, e can become either $F, E, G, S,$ or R
 - Letter frequencies become “averaged”

- Can we guess the length of the key?

- Try displacing the message and counting the coincidences, when a letter matches the displaced letter

- If we can guess the key length, perhaps we can recover the letter frequency and recover the key

Classical Cryptosystems - Section 2.3.1

FINDING KEY LENGTH

Attacks on the Vigenère Cipher

- Example:

plain: whatdoesthefoxsay

cipher: XHCHQPEUHUFFQLFBY

d=1	XHCHQPEUHU F QLFBY	1
d=2	X H CHQPE U HUFFQLFBY	2
d=3	XHCHQPEUHUF F QLFBY	1
d=4	XHCHQPEUHU F QLFBY	1
d=5	XHC H QPEUHUFFQLFBY	1

Message is too short!

English Language (Frequency Analysis)

