(2.1) The Bisection Method MATH 4701 Numerical Analysis

Intermediate Value Theorem

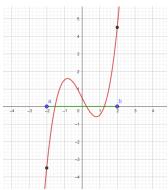
Intermediate value theorem (special case): If $f(x):[a,b]\to\mathbb{R}$ is continuous and f(a)f(b)<0 (f(a) and f(b) have opposite signs) then for some $a<\alpha< b$ we have $f(\alpha)=0$.



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While the theorem seems to only have theoretical implications on existence of a solution without giving any information on how that solution can be found, one can in fact approximate a solution f(x) = 0 using **bisection method** which is based on the intermediate value theorem.

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Examples

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- For $F(x) = (x+1)x^3(x-1)(x-2)^2(x-3)(x-5)$ if we start with $x_0 = -2$ and $x_1 = 7$ which of the solutions of F(x) = 0 will the bisection method find.