

Notes on Formal Compiler Construction with the π Framework

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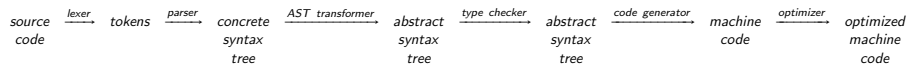
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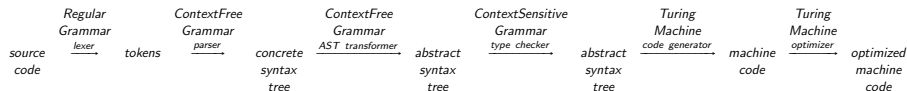
<http://github.com/ChristianoBraga/BPLC>



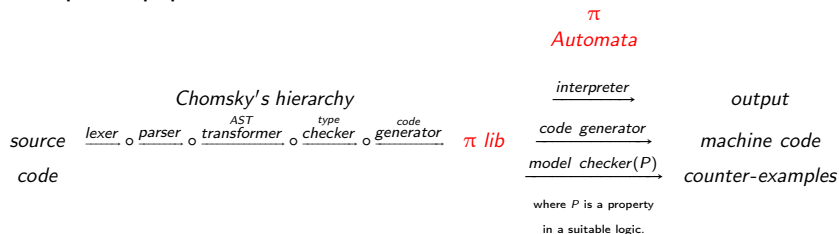
Compiler pipeline



Compiler pipeline and formal languages



Compiler pipeline with the π Framework



- π lib defines a set of constructions common to many programming languages.
- π lib constructions have a formal automata-based semantics in π automata.
- One may execute (or validate) a program in a given language by running its associated π lib program.
- π Framework: <http://github.com/ChristianoBraga/BPLC>
- Notes on Formal Compiler Construction with the π Framework: <https://github.com/ChristianoBraga/BPLC/blob/master/notes/notes.pdf>.

A calculator

We wish to compute simple arithmetic expressions such as $5 * (3 + 2)$.

A calculator: Lexer

$\langle \textit{digit} \rangle ::= [0..9]$

$\langle \textit{digits} \rangle ::= \langle \textit{digit} \rangle^+$

$\langle \textit{boolean} \rangle ::= \text{'true'} \mid \text{'false'}$

A calculator: concrete syntax

$$\langle exp \rangle ::= \langle aexp \rangle \mid \langle bexp \rangle$$
$$\langle aexp \rangle ::= \langle aexp \rangle '+' \langle term \rangle \mid \langle aexp \rangle '-' \langle term \rangle \mid \langle term \rangle$$
$$\langle term \rangle ::= \langle term \rangle '*' \langle factor \rangle \mid \langle term \rangle '/' \langle factor \rangle \mid \langle factor \rangle$$
$$\langle factor \rangle ::= '(' \langle aexp \rangle ')' \mid \langle digits \rangle$$
$$\langle bexp \rangle ::= \langle boolean \rangle \mid '\sim' \langle bexp \rangle \mid \langle bexp \rangle \langle boolop \rangle \langle bexp \rangle \\ \mid \langle aexp \rangle \langle iop \rangle \langle aexp \rangle$$
$$\langle iop \rangle ::= '=' \mid '<' \mid '>' \mid '<=' \mid '>='$$

A calculator: abstract syntax

$$\langle exp \rangle ::= \langle digits \rangle \mid \langle boolean \rangle \mid \langle exp \rangle \langle bop \rangle \langle exp \rangle$$
$$\langle bop \rangle ::= '+' \mid '-' \mid '*' \mid '|' \mid '/' \mid '=' \mid '<' \mid '>' \mid '<=' \mid '>='$$

A calculator: π denotations I

Let D in $\langle \text{digits} \rangle$, B in $\langle \text{boolean} \rangle$ and E_1, E_2 in $\langle \text{exp} \rangle$,

$$\llbracket D \rrbracket_{\pi} = \text{Num}(D) \quad (1)$$

$$\llbracket B \rrbracket_{\pi} = \text{Boo}(B) \quad (2)$$

$$\llbracket E_1 + E_2 \rrbracket_{\pi} = \text{Sum}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (3)$$

$$\llbracket E_1 - E_2 \rrbracket_{\pi} = \text{Sub}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (4)$$

$$\llbracket E_1 * E_2 \rrbracket_{\pi} = \text{Mul}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (5)$$

$$\llbracket E_1 / E_2 \rrbracket_{\pi} = \text{Div}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (6)$$

$$\llbracket E_1 < E_2 \rrbracket_{\pi} = \text{Lt}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (7)$$

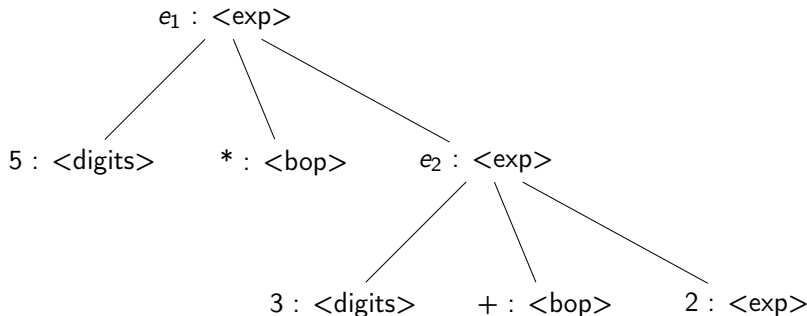
$$\llbracket E_1 \leq E_2 \rrbracket_{\pi} = \text{Le}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (8)$$

$$\llbracket E_1 > E_2 \rrbracket_{\pi} = \text{Gt}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (9)$$

$$\llbracket E_1 \geq E_2 \rrbracket_{\pi} = \text{Ge}(\llbracket E_1 \rrbracket_{\pi}, \llbracket E_2 \rrbracket_{\pi}) \quad (10)$$

A calculator: π denotations II

- π denotations are *functions* $\llbracket \cdot \rrbracket_{\pi} : AST \rightarrow \pi \text{ lib}$, where *AST* denotes the *datatype* for the abstract syntax tree and $\pi \text{ lib}$ denotes the datatype for $\pi \text{ lib}$ programs.
- Note that $\llbracket \cdot \rrbracket_{\pi}$ has *trees* as parameters, instances of *AST*. The example expression $5 * (3 + 2)$ becomes



A calculator: π denotations III

$$\llbracket 5 * (3 + 2) \rrbracket_{\pi} = \text{Mul}(\llbracket 5 \rrbracket_{\pi}, \llbracket (3 + 2) \rrbracket_{\pi}) \quad \text{by Equation 5}$$

$$\text{Mul}(\llbracket 5 \rrbracket_{\pi}, \llbracket (3 + 2) \rrbracket_{\pi}) = \text{Mul}(\text{Num}(5), \llbracket (3 + 2) \rrbracket_{\pi}) \quad \text{by Equations 1}$$

$$\text{Mul}(\text{num}(5), \llbracket (3 + 2) \rrbracket_{\pi}) = \text{Mul}(\text{Num}(5), \text{Sum}(\llbracket 3 \rrbracket_{\pi}, \llbracket 2 \rrbracket_{\pi})) \quad \text{by Equation 3}$$

$$\text{Mul}(\text{Num}(5), \text{Sum}(\llbracket 3 \rrbracket_{\pi}, \llbracket 2 \rrbracket_{\pi})) = \text{Mul}(\text{Num}(5), \text{Sum}(\text{Num}(3), \llbracket 2 \rrbracket_{\pi})) \quad \text{by Equation 1}$$

$$\text{Mul}(\text{Num}(5), \text{Sum}(\text{Num}(3), \llbracket 2 \rrbracket_{\pi})) = \text{Mul}(\text{Num}(5), \text{Sum}(\text{Num}(3), \text{Num}(2))) \quad \text{by Equation 1}$$

A calculator: executing π lib with π automata

A π automaton is a 5-tuple $\mathcal{A} = (G, Q, \delta, q_0, F)$, where G is a context-free grammar, Q is the set of states, q_0 is the initial state, $F \subseteq Q$ is the set of final states and

$$\delta : L(G)^* \times L(G)^* \times Store \rightarrow Q,$$

where $L(G)$ is the language generated by G and $Store$ represents the memory. (Elements in a set S^* are represented by terms $[s_1, s_2, \dots, s_n].$)

$$\begin{aligned} \delta([Mul(Num(5), Sum(Num(3), Num(2))), \emptyset, \emptyset) &= \delta([Num(5), Sum(Num(3), Num(2)), \#MUL], \emptyset, \emptyset) \\ \delta([Num(5), Sum(Num(3), Num(2)), \#MUL], \emptyset, \emptyset) &= \delta([Sum(Num(3), Num(2)), \#MUL], [Num(5)], \emptyset) \\ \delta([Sum(Num(3), Num(2)), \#MUL], [Num(5)], \emptyset) &= \delta([Num(3), Num(2), \#SUM, \#MUL], [Num(5)], \emptyset) \\ \delta([Num(3), Num(2), \#SUM, \#MUL], [Num(5)], \emptyset) &= \delta([Num(2), \#SUM, \#MUL], [Num(3), Num(5)], \emptyset) \\ \delta([Num(2), \#SUM, \#MUL], [Num(3), Num(5)], \emptyset) &= \delta([\#SUM, \#MUL], [Num(2), Num(3), Num(5)], \emptyset) \\ \delta([\#SUM, \#MUL], [Num(2), Num(3), Num(5)], \emptyset) &= \delta([\#MUL], [Num(5), Num(5)], \emptyset) \\ \delta([\#MUL], [Num(5), Num(5)], \emptyset) &= \delta(\emptyset, [Num(25)], \emptyset) \\ \delta(\emptyset, [Num(25)], \emptyset) &= Num(25) \end{aligned}$$

π lib expressions

$\langle \text{Statement} \rangle ::= \langle \text{Exp} \rangle$

$\langle \text{Exp} \rangle ::= \langle \text{ArithExp} \rangle \mid \langle \text{BoolExp} \rangle$

$\langle \text{ArithExp} \rangle ::= \text{'Num'}(\langle \text{digits} \rangle) \mid \text{'Sum'}(\langle \text{Exp} \rangle, \langle \text{Exp} \rangle) \mid$
 $\text{'Sub'}(\langle \text{Exp} \rangle, \langle \text{Exp} \rangle) \mid \text{'Mul'}(\langle \text{Exp} \rangle, \langle \text{Exp} \rangle)$

$\langle \text{BoolExp} \rangle ::= \text{'Eq'}(\langle \text{Exp} \rangle, \langle \text{Exp} \rangle) \mid \text{'Not'}(\langle \text{Exp} \rangle)$

π automata semantics for π lib expressions

- Recall that $\delta : L(G)^* \times L(G)^* \times Store \rightarrow Q$, and let where $N, N_i \in \mathbb{N}$, $C, V \in L(G)^*$, $S \in Store$,

$$\delta(Num(N) :: C, V, S) = \delta(C, Num(N) :: V, S) \quad (11)$$

$$\delta(Sum(E_1, E_2) :: C, V, S) = \delta([E_1, E_2, \#SUM] :: C, V, S) \quad (12)$$

$$\delta(\#SUM :: C, [Num(N_1), Num(N_2)] :: V, S) = \delta(C, Num(N_1 + N_2) :: V, S) \quad (13)$$

...

$$\delta(Not(E) :: C, V, S) = \delta([E, \#NOT] :: C, V, S) \quad (14)$$

$$\delta(\#NOT :: C, [Boo(true)] :: V, S) = \delta(C, [Boo(false)] :: V, S) \quad (15)$$

$$\delta(\#NOT :: C, [Boo(false)] :: V, S) = \delta(C, [Boo(true)] :: V, S) \quad (16)$$

- Notation $h :: ls$ denotes the concatenation of element h with the list ls . The same notation is used for appending two lists.
- C represents the *control* stack. V represents the *value* stack. S denotes the memory store.
- $\delta(\emptyset, V, S)$ denotes an *accepting state*.

π lib expressions in Python I

<https://github.com/ChristianBraga/BPLC/blob/master/python/pi.ipynb>

```
1 class Statement:
2     def __init__(self, *args):
3         self.opr =args
4     def __str__(self):
5         ret =str(self.__class__.__name__)+ "("
6         for o in self.opr:
7             ret +=str(o)
8         ret +=")"
9         return ret
10 class Exp(Statement): pass
11 class ArithExp(Exp): pass
```

π lib expressions in Python II

```
1 class Num(ArithExp):
2     def __init__(self, f):
3         assert(isinstance(f, int))
4         ArithExp.__init__(self, f)
5 class Sum(ArithExp):
6     def __init__(self, e1, e2):
7         assert(isinstance(e1, Exp) and isinstance(e2, Exp))
8         ArithExp.__init__(self, e1, e2)
9 ...
```


π lib expressions in Python III

```
1 class BoolExp(Exp): pass
2 class Eq(BoolExp):
3     def __init__(self, e1, e2):
4         assert(isinstance(e1, Exp) and isinstance(e2, Exp))
5         BoolExp.__init__(self, e1, e2)
6 ...
```

π lib expressions in Python IV

```
1 exp = Sum(Num(1), Mul(Num(2), Num(4)))  
2 print(exp)  
3  
4 Sum(Num(1)Mul(Num(2)Num(4)))
```

π lib expressions in Python V

```

1 exp2 =Mul(2, 1)
2 -----
3
4
5
6
7
8
9
10
11
12
13
14

```

AssertionError Traceback (most recent call last)

<ipython-input-7-00fd40a79a54> in <module>()

---->1 exp2 =Mul(2, 1)

<ipython-input-5-42a82e58862f> in __init__(self, e1, e2)

```

28 class Mul(ArithExp):
29     def __init__(self, e1, e2):
30     assert(isinstance(e1, Exp) and isinstance(e2, Exp))
31     ArithExp.__init__(self, e1, e2)
32     class BoolExp(Exp): pass

```

AssertionError:

π automaton for π lib expressions I

```
1  ## Expressions
2  class ValueStack(list): pass
3  class ControlStack(list): pass
4  class ExpKW:
5      SUM = "#SUM"
6      SUB = "#SUB"
7      MUL = "#MUL"
8      EQ = "#EQ"
9      NOT = "#NOT"
```

π automaton for π lib expressions II

```
1 class ExpPiAut(dict):
2     def __init__(self):
3         self["val"] = ValueStack()
4         self["cnt"] = ControlStack()
5     def __evalSum(self, e):
6         e1 = e.opr[0]
7         e2 = e.opr[1]
8         self.pushCnt(ExpKW.SUM)
9         self.pushCnt(e1)
10        self.pushCnt(e2)
11    def pushCnt(self, e):
12        cnt = self.cnt()
13        cnt.append(e)
14    ...
```

π automaton for π lib expressions III

```
1 ea =ExpPiAut()  
2 print(exp)  
3 ea.pushCnt(exp)  
4 while not ea.emptyCnt():  
5     ea.eval()  
6     print(ea)
```

π automaton for π lib expressions IV

```

1 Sum(Num(1)Mul(Num(2)Num(4)))
2 {'val': [], 'cnt': ['#SUM', <__main__.Num object at 0x111851470>, <
    __main__.Mul object at 0x1118516d8>]}
3 {'val': [], 'cnt': ['#SUM', <__main__.Num object at 0x111851470>, '#MUL',
    , <__main__.Num object at
    0x111851630>, <__main__.Num object
    at 0x1118516a0>]}
4 {'val': [4], 'cnt': ['#SUM', <__main__.Num object at 0x111851470>, '#MUL',
    , <__main__.Num object at
    0x111851630>]}
5 {'val': [4, 2], 'cnt': ['#SUM', <__main__.Num object at 0x111851470>, '#MUL']}
6 {'val': [8], 'cnt': ['#SUM', <__main__.Num object at 0x111851470>]}
7 {'val': [8, 1], 'cnt': ['#SUM']}
8 {'val': [9], 'cnt': []}

```