Written Assignment 1

Vinay Patil

vpatil3@wisc.edu

1 Calculator language: Syntax

```
\begin{array}{l} {\rm digits} = \ \ "0" \ | \ "1" \ | \ "2" \ | \ "3" \ | \ "4" \ | \ "5" \ | \ "6" \ | \ "7" \ | \ "8" \ | \ "9" \ ; \\ {\rm boolconst} = \ "tt" \ | \ "ff" \ ; \\ {\rm int} = [\ "-"] \ {\rm digits} \ \{{\rm digits}\} \ ; \\ {\rm be} = {\rm boolconst} \\ | \ {\rm be} \ "\&\&" \ {\rm be} \\ | \ {\rm be} \ "\|" \ {\rm be} \\ | \ {\rm be} \ "\|" \ {\rm be} \\ | \ {\rm ae} \ "==" \ {\rm ae} \\ | \ {\rm ae} \ "<" \ {\rm ae} \ ; \\ {\rm ae} = {\rm int} \\ | \ {\rm ae} \ "\times" \ {\rm ae} \\ | \ {\rm ae} \ "\times" \ {\rm ae} \\ | \ "if" \ {\rm be} \ "then" \ {\rm ae} \ "else" \ {\rm ae} \ ; \\ \end{array}
```

2 Calculator language: Operational semantics

2.1 Values that do not step (Defining here)

$$\begin{split} & \text{b1, b2, b} = \text{``tt"} \mid \text{``ff"} \ ; \\ & i_1, \, i_2, \, \text{i} = [\text{``-"}] \text{ digits } \{ \text{digits} \} \ ; \end{split}$$

2.2 Semantics for be "&&" be

Let
$$b_1$$
, b_2 , $b \in boolconst$
Let be_1 , be_1^1 , be_2 , $be_2^1 \in be$

$$\begin{aligned} \frac{b_1\&\&b_2 = b}{(b_1``\&\&"b_2) \to b} \\ \frac{be_1 \to b_1}{(be_1``\&\&"b_2) \to (b_1``\&\&"b_2)} \\ \\ \frac{be_1 \to be_1^1}{(be_1``\&\&"b_2) \to (be_1^1``\&\&"b_2)} \end{aligned}$$

$$\begin{array}{c} be_2 \to b_2 \\ \hline (b_1 \text{``\&\&'''} be_2) \to (b_1 \text{``\&\&'''} b_2) \\ \hline be_2 \to be_2^1 \\ \hline (b_1 \text{``\&\&'''} be_2) \to (b_1 \text{``\&\&'''} be_2^1) \\ \hline be_1 \to b_1, \ be_2 \to be_2^1 \\ \hline (be_1 \text{``\&\&'''} be_2) \to (b_1 \text{``\&\&'''} be_2^1) \\ \hline be_1 \to be_1^1, \ be_2 \to b_2 \\ \hline (be_1 \text{``\&\&'''} be_2) \to (be_1^1 \text{``\&\&'''} be_2) \\ \hline be_1 \to b_1, \ be_2 \to b_2 \\ \hline (be_1 \text{``\&\&'''} be_2) \to (b_1 \text{``\&\&'''} be_2) \\ \hline be_1 \to be_1^1, \ be_2 \to be_2^1 \\ \hline (be_1 \text{``\&\&'''} be_2) \to (be_1^1 \text{``\&\&'''} be_2^1) \\ \hline \end{array}$$

2.3 Semantics for be "||" be

Let $b_1, b_2, b \in \text{boolconst}$ Let $be_1, be_1^1, be_2, be_2^1 \in \text{be}$

$$\begin{array}{c} b_1 \mid |b_2 = b \\ \hline (b_1`` \mid |"b_2) \to b \\ \\ be_1 \to b_1 \\ \hline (be_1`` \mid |"b_2) \to (b_1`` \mid |"b_2) \\ \\ \underline{be_1 \to be_1^1} \\ \hline (be_1`` \mid |"b_2) \to (be_1^1`` \mid |"b_2) \\ \\ \underline{be_2 \to b_2} \\ \hline (b_1`` \mid |"be_2) \to (b_1`` \mid |"b_2) \\ \\ \underline{be_2 \to be_2^1} \\ \hline (b_1`` \mid |"be_2) \to (b_1`` \mid |"be_2^1) \\ \\ \underline{be_1 \to b_1, \ be_2 \to be_2^1} \\ \hline (be_1`` \mid |"be_2) \to (b_1`` \mid |"be_2^1) \\ \\ \underline{be_1 \to be_1^1, \ be_2 \to be_2} \\ \hline (be_1`` \mid |"be_2) \to (be_1^1`` \mid |"b_2) \\ \\ \underline{be_1 \to be_1, \ be_2 \to be_2} \\ \hline (be_1`` \mid |"be_2) \to (b_1`` \mid |"be_2) \\ \\ \underline{be_1 \to be_1^1, \ be_2 \to be_2^1} \\ \hline (be_1`` \mid |"be_2) \to (be_1^1`` \mid |"be_2^1) \\ \\ \\ \underline{be_1 \to be_1^1, \ be_2 \to be_2^1} \\ \hline (be_1`` \mid |"be_2) \to (be_1^1`` \mid |"be_2^1) \\ \\ \hline \end{array}$$

2.4 Semantics for ae "==" ae

Let $i_1, i_2, i \in \text{int}$ Let $b \in \text{boolconst}$ Let $ae_1, ae_1^1, ae_2, ae_2^1 \in \text{ae}$

$$\frac{(i_1 == i_2) = b}{(i_1" == "i_2) \to b}$$

$$\frac{ae_1 \to i_1}{(ae_1" == "i_2) \to (i_1" == "i_2)}$$

$$\frac{ae_1 \to ae_1^1}{(ae_1" == "i_2) \to (ae_1^1" == "i_2)}$$

$$\frac{ae_2 \to i_2}{(i_1" == "ae_2) \to (i_1" == "i_2)}$$

$$\frac{ae_2 \to ae_2^1}{(i_1" == "ae_2) \to (i_1" == "ae_2^1)}$$

$$\frac{ae_1 \to i_1, \ ae_2 \to ae_2^1}{(ae_1" == "ae_2) \to (i_1" == "ae_2^1)}$$

$$\frac{ae_1 \to ae_1^1, \ ae_2 \to i_2}{(ae_1" == "ae_2) \to (ae_1^1" == "i_2)}$$

$$\frac{ae_1 \to ae_1^1, \ ae_2 \to i_2}{(ae_1" == "ae_2) \to (i_1" == "i_2)}$$

$$\frac{ae_1 \to ae_1^1, \ ae_2 \to ae_2^1}{(ae_1" == "ae_2) \to (i_1" == "i_2)}$$

2.5 Semantics for ae "<" ae

Let $i_1, i_2, i \in \text{int}$ Let $b \in \text{boolconst}$ Let $ae_1, ae_1^1, ae_2, ae_2^1 \in \text{ae}$

$$\frac{(i_1 < i_2) = b}{(i_1" < "i_2) \to b}$$

$$\frac{ae_1 \to i_1}{(ae_1" < "i_2) \to (i_1" < "i_2)}$$

$$\frac{ae_1 \to ae_1^1}{(ae_1" < "i_2) \to (ae_1^1" < "i_2)}$$

$$\frac{ae_2 \to i_2}{(i_1" < "ae_2) \to (i_1" < "i_2)}$$

$$\frac{ae_2 \to ae_2^1}{(i_1" < "ae_2) \to (i_1" < "ae_2^1)}$$

$$\frac{ae_1 \to i_1, \ ae_2 \to ae_2^1}{(ae_1" < "ae_2) \to (i_1" < "ae_2^1)}$$

$$\frac{ae_1 \to ae_1^1, \ ae_2 \to i_2}{(ae_1" < "ae_2) \to (ae_1^1" < "i_2)}$$

$$\frac{ae_1 \to i_1, \ ae_2 \to i_2}{(ae_1" < "ae_2) \to (i_1" < "i_2)}$$

$$\frac{ae_1 \to i_1, \ ae_2 \to i_2}{(ae_1" < "ae_2) \to (i_1" < "i_2)}$$

$$\frac{ae_1 \to ae_1^1, \ ae_2 \to ae_2^1}{(ae_1" < "ae_2) \to (ae_1^1" < "ae_2^1)}$$

2.6 Semantics for ae "+" ae

Let $i_1, i_2, i \in \text{int}$ Let $ae_1, ae_1^1, ae_2, ae_2^1 \in \text{ae}$

$$\begin{split} \frac{i_1+i_2=i}{(i_1``+"i_2)\to i} \\ \frac{ae_1\to i_1}{(ae_1``+"i_2)\to (i_1``+"i_2)} \\ \frac{ae_1\to ae_1^1}{(ae_1``+"i_2)\to (ae_1^1``+"i_2)} \\ \frac{ae_2\to i_2}{(i_1``+"ae_2)\to (i_1``+"i_2)} \\ \frac{ae_2\to ae_2^1}{(i_1``+"ae_2)\to (i_1``+"ae_2^1)} \\ \frac{ae_1\to i_1,\ ae_2\to ae_2^1}{(ae_1``+"ae_2)\to (i_1``+"ae_2^1)} \\ \frac{ae_1\to i_1,\ ae_2\to ae_2^1}{(ae_1``+"ae_2)\to (ae_1^1``+"i_2)} \\ \frac{ae_1\to ae_1^1,\ ae_2\to i_2}{(ae_1``+"ae_2)\to (i_1``+"i_2)} \\ \frac{ae_1\to i_1,\ ae_2\to i_2}{(ae_1``+"ae_2)\to (i_1``+"i_2)} \\ \frac{ae_1\to ae_1^1,\ ae_2\to ae_2^1}{(ae_1``+"ae_2)\to (ae_1^1``+"ae_2^1)} \\ \end{split}$$

2.7 Semantics for ae "x" ae

Let $i_1, i_2, i \in \text{int}$ Let $ae_1, ae_1^1, ae_2, ae_2^1 \in \text{ae}$

$$\begin{aligned} \frac{i_1 \times i_2 = i}{(i_1 " \times " i_2) \to i} \\ \frac{ae_1 \to i_1}{(ae_1 " \times " i_2) \to (i_1 " \times " i_2)} \\ \frac{ae_1 \to ae_1^1}{(ae_1 " \times " i_2) \to (ae_1^1 " \times " i_2)} \\ \frac{ae_2 \to i_2}{(i_1 " \times " ae_2) \to (i_1 " \times " i_2)} \\ \frac{ae_2 \to ae_2^1}{(i_1 " \times " ae_2) \to (i_1 " \times " ae_2^1)} \\ \frac{ae_1 \to i_1, \ ae_2 \to ae_2^1}{(ae_1 " \times " ae_2) \to (i_1 " \times " ae_2^1)} \\ \frac{ae_1 \to i_1, \ ae_2 \to ae_2^1}{(ae_1 " \times " ae_2) \to (ae_1^1 " \times " i_2)} \\ \frac{ae_1 \to ae_1^1, \ ae_2 \to i_2}{(ae_1 " \times " ae_2) \to (i_1 " \times " i_2)} \\ \frac{ae_1 \to ae_1^1, \ ae_2 \to ae_2^1}{(ae_1 " \times " ae_2) \to (i_1 " \times " i_2)} \\ \frac{ae_1 \to ae_1^1, \ ae_2 \to ae_2^1}{(ae_1 " \times " ae_2) \to (ae_1^1 " \times " ae_2^1)} \end{aligned}$$

2.8 Semantics for if be then ae else ae: (Lazy-If)

Let $i_1, i_2, i \in \text{int}$ Let $ae_1, ae_1^1, ae_2, ae_2^1 \in \text{ae}$ Let $b_1, b_2, b \in \text{boolconst}$ Let $be_1, be_1^1, be_2, be_2^1 \in \text{be}$

$$\begin{array}{c} i1\\ \hline if "tt" then \ i_1 \ else \ ae_2 \rightarrow i_1 \\ \hline \\ i2\\ \hline if "ff" then \ ae_1 \ else \ i_2 \rightarrow i_2 \\ \hline \\ ae_1 \rightarrow i_1\\ \hline if "tt" then \ ae_1 \ else \ ae_2 \rightarrow if "tt" then \ i_1 \ else \ ae_2\\ \hline \\ ae_1 \rightarrow ae_1^1\\ \hline if "tt" then \ ae_1 \ else \ ae_2 \rightarrow if "tt" then \ ae_1^1 \ else \ ae_2\\ \hline \\ if "ff" then \ ae_1 \ else \ ae_2 \rightarrow if "ff" then \ ae_1 \ else \ i_2\\ \hline \end{array}$$

$$\begin{array}{c} ae_2 \rightarrow ae_2^1 \\ \hline if \ "ff" \ then \ ae_1 \ else \ ae_2 \rightarrow if \ "ff" \ then \ ae_1 \ else \ ae_2^1 \\ \hline \\ \underline{be_1 \rightarrow b_1} \\ \hline if \ be_1 \ then \ ae_1 \ else \ ae_2 \rightarrow if \ b1 \ then \ ae_1 \ else \ ae_2 \\ \hline \\ \underline{be_1 \rightarrow be_1^1} \\ \hline if \ be_1 \ then \ ae_1 \ else \ ae_2 \rightarrow if \ be_1^1 \ then \ ae_1 \ else \ ae_2 \\ \end{array}$$

2.9 Semantics for if be then ae else ae: (Eager-If)

Let $i_1, i_2, i \in \text{int}$ Let $ae_1, ae_1^1, ae_2, ae_2^1 \in \text{ae}$ Let $b_1, b_2, b \in \text{boolconst}$ Let $be_1, be_1^1, be_2, be_2^1 \in \text{be}$

$$i1$$

$$if "tt" then i_1 else i_2 \rightarrow i_1$$

$$i2$$

$$if "ff" then i_1 else i_2 \rightarrow i_2$$

$$be_1 \rightarrow b_1$$

$$if be_1 then i_1 else i_2 \rightarrow if b_1 then i_1 else i_2$$

$$be_1 \rightarrow be_1^1$$

$$if be_1 then i_1 else i_2 \rightarrow if be_1^1 then i_1 else i_2$$

$$ae_1 \rightarrow i_1$$

$$if be_1 then ae_1 else i_2 \rightarrow if be_1 then i_1 else i_2$$

$$ae_2 \rightarrow i_2$$

$$if be_1 then i_1 else ae_2 \rightarrow if be_1 then i_1 else i_2$$

$$ae_1 \rightarrow i_1, ae_2 \rightarrow ae_2^1$$

$$if be_1 then ae_1 else ae_2 \rightarrow if be_1 then i_1 else ae_2^1$$

$$if be_1 then ae_1 else ae_2 \rightarrow if be_1 then ae_1^1 else i_2$$

$$ae_1 \rightarrow ae_1^1, ae_2 \rightarrow ae_2^1$$

$$if be_1 then ae_1 else ae_2 \rightarrow if be_1 then ae_1^1 else i_2$$

$$ae_1 \rightarrow ae_1^1, ae_2 \rightarrow ae_2^1$$

$$if be_1 then ae_1 else ae_2 \rightarrow if be_1 then ae_1^1 else ae_2^1$$

- 1. Lazy-If is preferred over Eager-If as it ends up computing less to arrive at the same result when compared to Eager-If as in Eager-If all the expressions in body are evaluated irrespective of the guard condition.
 - 2. Also in the case of Eager-If if the arguments doesn't terminate the function call doesn't terminate.

3 Lambda calculus

3.1

```
(\lambda f. \ \lambda x. \ f \ x) \ (\lambda x. \ x+1) \ 5
\rightarrow (\lambda x. \ (\lambda x. \ x+1) \ x) \ 5
\rightarrow (\lambda x. \ x+1) \ 5
\rightarrow 6
```

3.2

$$\begin{array}{l} (\lambda f. \ \lambda x. \ \lambda y. \ f \ y \ x) \ (\lambda x. \ \lambda y. \ x - y) \ 5 \ 3 \\ \rightarrow & ((\lambda x. \ \lambda y. \ (\lambda x. \ \lambda y. \ x - y) \ y \ 5) \ 3) \\ \rightarrow & ((\lambda y. \ (\lambda x. \ \lambda y. \ x - y) \ y \ 5) \ 3) \\ \rightarrow & ((\lambda x. \ \lambda y. \ x - y) \ 3 \ 5) \\ \rightarrow & ((\lambda y. \ 3 - y) \ 5) \\ \rightarrow & (3 - 5) \\ \rightarrow & -2 \end{array}$$

3.3

```
(\lambda x. \ x \ x) \ (\lambda x. \ x \ x)
\rightarrow (\lambda x. \ x \ x) \ (\lambda x. \ x \ x)
\rightarrow (\lambda x. \ x \ x) \ (\lambda x. \ x \ x)
\rightarrow .
\rightarrow .
\rightarrow .
\rightarrow .
\rightarrow . (Neverstops! - Infinitely recursive)
```

4 Recursion

```
fibonacci = fix f. \lambda n. if n < 2 then 1 else (f(n-1) + f(n-2))

Test n=3:

\rightarrow [\lambda n. if \ n < 2 \ then \ 1 \ else \ ((fix \ f...)(n-1) + (fix \ f...)(n-2)] \ 3

\rightarrow if \ 3 < 2 \ then \ 1 \ else \ ((fix \ f...)(3-1) + (fix \ f...)(3-2))

\rightarrow ((fix \ f...)(2) + (fix \ f...)(1))

\rightarrow ((if \ 1 < 2 \ then \ 1 \ else \ ...) + (if \ 0 < 2 \ then \ 1 \ else \ ...) + (if \ 1 < 2 \ then \ 1 \ else \ ...)))

\rightarrow 1 \ + 1 \ + 1
```