

# Advanced R Programming - Lecture 6

## Computational complexity

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(slides based on Leif Jonsson's and Måns Magnusson's)

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# Today

Optimizing code

Performant Code

Computational complexity

Classes of problems

Big Oh notation

Determining complexity

# Questions since last time?

# Donald E. Knuth on Optimization

Programmers waste enormous amounts of time thinking about, or worrying about, the speed of noncritical parts of their programs, and these attempts at efficiency actually have a strong negative impact when debugging and maintenance are considered.

- Donald E. Knuth

# Performance

Depends on many things

1. Code
2. Complexity
3. Compiler
4. Hardware
5. Language

**If you don't measure, you don't optimize!**

# How to optimize

0. Choose optimal algorithm
1. Write code that works with accompanying test suite
2. Profile your code for bottlenecks
3. Try to eliminate the bottle necks
4. Redo 2-3 until fast enough

`proc.time()` is a basic starting tool

## Profiling (code below is deprecated, see RprofEx.R)

```
Rprof(tmp <- tempfile(),
      line.profiling = TRUE,
      memory.profiling = TRUE)
test_data <- pxweb::get_pxweb_data(
  url =
    "http://api.scb.se/OV0104/v1/doris/sv/ssd/BE/BE0101
      /BE0101A/BefolkningNy",
  dims = list(Region = c('*'),
    Civilstand = c('*'),
    Alder = c('*'),
    Kon = c('*'),
    ContentsCode = c('*'),
    Tid = as.character(1970)),
  clean = TRUE)
Rprof()
summaryRprof(tmp, lines = "show", memory = "both")
```

# Profiling

```
$by.self
```

	self.time	self.pct	total.time	total.pct	mem.total
get_pxweb_data.R#102	1.96	39.2	1.96	39.2	579.2
get_pxweb_data_internal.R#42	1.16	23.2	1.16	23.2	405.0
get_pxweb_data.R#56	0.52	10.4	0.52	10.4	31.3
get_pxweb_data.R#80	0.38	7.6	0.38	7.6	29.1
get_pxweb_data.R#82	0.32	6.4	0.32	6.4	40.7
get_pxweb_data_internal.R#48	0.26	5.2	0.26	5.2	73.2
get_pxweb_data_internal.R#74	0.26	5.2	0.26	5.2	29.8
get_pxweb_data.R#83	0.08	1.6	0.08	1.6	17.2
api_catalogue.R#75	0.02	0.4	0.02	0.4	0.0
get_pxweb_data_internal.R#44	0.02	0.4	0.02	0.4	12.6
get_pxweb_data_internal.R#71	0.02	0.4	0.02	0.4	16.0

profvis is one modern profiling package.



# Improvements

- 0. Optimal data structure and algorithm
- 1. Look for existing solutions
- 2. Do less work
- 3. Vectorise
- 0. Optimal data structure and algorithm
- 4. Parallelize
- 0. Optimal data structure and algorithm
- 5. Avoid copies

# Writing fast code

Speed is important!  
(do not forget memory)

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# Writing fast code

Speed is important!  
(do not forget memory)

Time to write code  
Time to maintain (understand) code  
Time to execute code

# Old Adage About Software

"You can have it Good, Fast, Cheap. Pick any two."

# Performance

1. Performance
2. Complexity

Complexity affects performance

# Computational complexity

Theoretical worst case  
(but what about average case?)

Big-Oh notation

Basic operations

Relationship: operations to problem size



# Types of complexity

Time complexity

Space (memory) complexity

Worst case complexity

Average case complexity

# (theoretical) Data structures

Matrix (dataframe, list)

List (**NOT** in R sense, but with pointers), FIFO, LIFO

Sets (no particular order of elements, cannot index)

Graphs (vertex, edge): vertex adjacency matrix, vertex adjacency list

# Classes of problems

**Decision problems** answer is yes or no, e.g. is  $x$  a prime number

**Optimization problems** find an object that satisfies a certain property, e.g. largest prime number smaller than  $x + 1$

**Non-algorithmic problems** cannot be solved by an algorithm, e.g. *halting problem* does a given algorithm end in finite time or fall into an infinite loop?

**Presumably nonalgorithmic problems** no algorithm is known but we do not know if non-algorithmic e.g. *Collatz problem*

```
repeat{
    if (k%%2==0){k=k/2}else{k=3*k+1}
    if (k==1){break}
}
```

Does it halt for every  $k$ ?

# Classes of problems

**Non-polynomial problems** cannot be solved by an algorithm whose running time is bounded by a polynomial of its input's size  
e.g. generate all permutations of an  $n$  element set,  $n!$

**Polynomial problems** can be solved by an algorithm whose running time is bounded by a polynomial of its input's size e.g.  
sorting  $n$  elements

**Non-polynomial problems** cannot be solved by an algorithm whose running time is bounded by a polynomial

**P class** polynomial problems

# NP

**NP class** *Nondeterministic polynomial* class of problems, there exists a polynomial time procedure that verifies if something is an admissible solution, e.g. check if *graph colouring* is admissible

$$P \subset NP \quad \text{but} \quad P \stackrel{???}{=} NP$$

**NP-complete** every problem in NP can be reduced to it in polynomial time

e.g. bin packing, knapsack, longest common subsequence, chromatic number of graph,

TSP ( $\mathbb{N}$ ), multiprocessor scheduling (some)

*satisfiability (SAT)*: is there a way to assign TRUE, FALSE values so that a logical statement is TRUE?

**NP-hard**: if it can be solved in polynomial time, then  $SAT \in P$

# Big Oh

"How fast does a function grow?"

$$f(n) = O(g(n)) \quad \text{or} \quad f(n) \in O(g(n))$$

$$\exists C > 0 \quad \exists N_0 \in \mathbb{N} \quad \forall n \ni n > N_0 \quad |f(n)| \leq C * |g(n)|$$

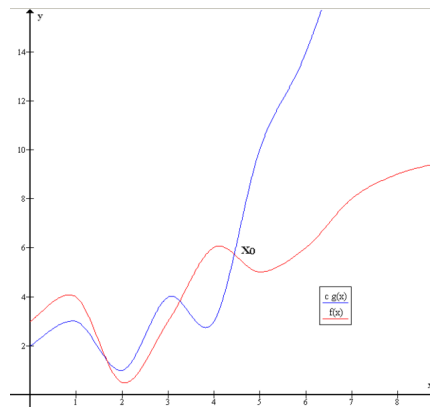
or

$$\limsup_{n \rightarrow \infty} \frac{|f(n)|}{|g(n)|} < \infty$$

$n$    number of operations

$f(n)$  does not (up to a scaling constant) grow faster than  $g(n)$

# Big Oh



[https://en.wikipedia.org/wiki/Big\\_O\\_notation](https://en.wikipedia.org/wiki/Big_O_notation)

# Big Oh

## Example

$$f(n) = n^2 + 100n + 100$$



# Big Oh

## Example

$$f(n) = n^2 + 100n + 100$$

$$f(n) = O(n^2)$$

## Other Oh

$$f = o(g) \quad \forall_{C>0} \exists_{N_0 \in \mathbb{N}} \forall_{N \ni n > N_0} |f(n)| \leq C|g(n)| \quad \lim_{n \rightarrow \infty} \frac{|f(n)|}{|g(n)|} = 0$$

$$f = O(g) \quad \exists_{C>0} \exists_{N_0 \in \mathbb{N}} \forall_{N \ni n > N_0} |f(n)| \leq C|g(n)| \quad \limsup_{n \rightarrow \infty} \frac{|f(n)|}{|g(n)|} < \infty$$

$$f = \omega(g) \quad \forall_{C>0} \exists_{N_0 \in \mathbb{N}} \forall_{N \ni n > N_0} |f(n)| \geq C|g(n)| \quad \lim_{n \rightarrow \infty} \frac{|f(n)|}{|g(n)|} = \infty$$

$$f = \Omega(g) \quad \exists_{C>0} \exists_{N_0 \in \mathbb{N}} \forall_{N \ni n > N_0} f(n) \geq Cg(n) \quad \liminf_{n \rightarrow \infty} \frac{f(n)}{g(n)} > 0$$

$$f = \Theta(g) \quad f = O(g) \text{ and } f = \Omega(g)$$

$$f \sim g \quad \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 1$$

# Complexities (the data size is a lower bound)

Big Oh	Name	Example, optimal
$O(1)$	constant	assignments, $O(1)$
$O(\log N)$	logarithmic	binary search (sorted input), $O(\log N)$
$O(N)$	linear	max., $O(N)$
$O(N \log N)$	log-linear	sorting, $O(N \log N)$
$O(N^2)$	quadratic	naive vector-matrix mult., preprocessing
$O(N^3)$	cubic	naive matrix inversion, $O(n^{2.373})$
$O(N^3)$	cubic	naive matrix-matrix mult., $O(n^{2.373})$
$O(N^c)$	polynomial	
$O(c^n)$	exponential	brute force cracking of password, ???

Quicksort:  $O(N^2)$  worst case, but  $O(N \log N)$  on average

# Determine complexity

```
statement 1  
statement 2  
...  
statement c
```

$O(1)$

# Determine complexity

```
if (a)
    statement a
else
    statement b
```

$\max(O(a), O(b))$

# Determine complexity

```
for(i in 1:N)  
  statement i
```

$O(n)$

# Determine complexity

```
for(i in 1:N)
  for (j in 1:M)    O ?
    statement i,j
```

# Determine complexity

```
for(i in 1:N)
  for (j in 1:M)    O(N * M)
    statement i,j
```



# Determine complexity

```
for(i in 1:N)  
  g(i)
```

$$g(n) = O(n^2)$$

$$O(n^3)$$

# Sorting

naïve sorting:  $O(n^2)$

merge sort:  $O(n \log n)$  but large number of copies

“merge sorted lists of two into four, then those and so on”

`sort()`

quicksort: average (uniform)  $O(n \log n)$ , worst  $O(n^2)$ , low overhead

radix sort:  $O(n \cdot k)$ , sorts numbers on  $k$  digits, by using the digits

shell sort:  $O(n^{4/3})$  sorts in-place by swapping elements

## Analysis of recursive algorithms (mergesort)

```
mergesort<-function(L){
  ## assume  $n=2^k$ 
  n<-length(L)
  if (n==1){return(L)}
  else{
    L1<-mergesort(L[1:(n/2)])
    L2<-mergesort(L[(n/2+1):n])
    ## merge is done in  $O(n)$  time
    return(merge(L1,L2))
  }
}
```

$$T(n) \leq \begin{cases} c_1 & n = 1 \\ 2T(n/2) + c_2n & n > 1 \end{cases}$$

## Analysis of recursive algorithms (Master Theorem)

A function  $f$  is multiplicative if  $f(xy) = f(x)f(y)$

Let  $a, b, c > 0$ ,  $k \in \mathbb{N}$  and  $d(n)$  be a multiplicative function. Then the solution to the recurrence equation

$$T(n) = \begin{cases} \Theta(1) & n = 1 \\ aT(n/b) + d(n) & n = b^k \end{cases}$$

is

$$T(n) = \Theta(a^k) + \sum_{j=0}^{k-1} a^j d(b^{k-j})$$

with asymptotic behaviour

$$T(n) = \begin{cases} \Theta(n^{\log_a d(b)}) & a < d(b) \\ \Theta(n^{\log_b a} \log n) & a = d(b) \\ \Theta(n^{\log_b a}) & a > d(b) \end{cases}$$

## Analysis of recursive algorithms (mergesort)

$c_n n$  is not multiplicative so take  $T(n) = c_2 \tilde{T}(n)$ , then

$$\tilde{T}(1) = T(1)/c_2 = c_1/c_2 = c$$

$$T(n) = 2T(n/2) + c_2 n \text{ becomes } c_2 \tilde{T}(n) = 2c_2 \tilde{T}(n/2) + c_2 n$$

Consider

$$U(n) = \begin{cases} c & n = 1 \\ 2U(n/2) + n & n > 1 \end{cases}$$

$n$  is multiplicative and using the Master Theorem we obtain

$$U(n) = \Theta(n \log n) \text{ and hence } U(n) \geq T(n) = O(n \log n).$$

Actually  $T(n) = \Theta(n \log n)$ .

## Approximate algorithms

If we cannot solve a hard problem let us approximate its solution.  
Let  $S_{opt}$  be the optimal solution and  $S_{approx}$  the approximate one

*k-absolute approximate* algorithm if  $|S_{opt} - S_{approx}| \leq k$

*k-(relative) approximate* algorithm if  $s \leq k$ , where

$$s = \max(S_{opt}/S_{approx}, S_{approx} - S_{opt})$$

**LAB:** knapsack problem

The End... for today.  
Questions?  
See you next time!