

CS 456: Quiz 7

Due on April 28, 2024

Nancy LaTorrette Section 1001

Christopher Howe

Problem

Prove that L is a context-free language using an NPDA (as explained in the lecture): $L = \{a^{2n}b^n : n \geq 1\}$

Solution

Goal

To prove that L is a context-free language using an NPDA, it must be shown that an NPDA can be created for L . If an NPDA exists that accepts the same language as defined in L , then L must be context-free language.

An NPDA that accepts L

A NPDA is defined by a tuple $M_{NPDA} = (Q, \Sigma, \Gamma, \delta, q_0, F)$ Where Q is the set of all possible states, Σ is the alphabet of input symbols, Γ is the set of stack symbols, δ is the transition function for a given input state, symbol, and stack symbol q_0 is the initial state, and F is the set of all final/accepted states.

A NPDA M_{NPDA} was designed to accept the same language as L . The general idea of how it works is it first reads all the a 's in pairs of two. When it reads the first a in the pair, it adds a b to the stack. It ignores the second a . If a single a comes in, this results in a dead configuration. If a b comes in before any a 's this also results in a dead configuration. After reading the pairs, the NPDA looks for a b . Once a b is found, the NPDA no longer can accept a 's without resulting in a dead configuration. For each b it reads in, it must find an equivalent b on the stack. Once the stack is empty it can take the lambda transition to the final state. It cannot read any more b 's at this point otherwise it will result in a dead configuration.

Below is the definition of such a NPDA which is followed by a state diagram of the NPDA.

$$M_{NPDA} = (\{q_0, q_a, q_b, q_f\}, \{a, b\}, \{\$, b\}, \delta, q_0, \{q_f\})$$

$$\delta : \left\{ \begin{array}{ll} \delta(q_0, a, \$) & \rightarrow (q_a, b\$) \\ \delta(q_0, a, b) & \rightarrow (q_a, bb) \\ \delta(q_0, b, b) & \rightarrow (q_b, \lambda) \\ \delta(q_a, a, b) & \rightarrow (q_0, b) \\ \delta(q_b, b, b) & \rightarrow (q_b, \lambda) \\ \delta(q_b, \lambda, \$) & \rightarrow (q_f, \lambda) \end{array} \right\}$$

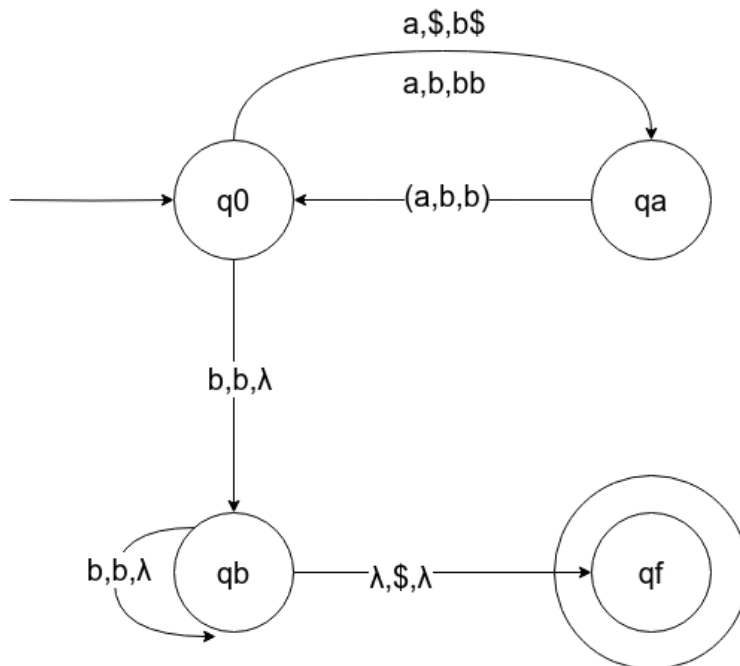


Figure 1: State Diagram of M_{NPDA} showing the various transitions used to accept L

Verification

To check that the PDA accepts only strings that are in the language, some test strings were applied. More tests were used but they are not shown here.

ba: q_0 stack:[\$] (Dead configuration)

ab: $q_0 \rightarrow q_a$ stack:[b\$] (Dead configuration)

aaaab: q_0 stack:[\$] $\rightarrow q_a$ stack:[b\$] $\rightarrow q_0$ stack:[b\$] $\rightarrow q_a$ stack:[bb\$] $\rightarrow q_0$ stack:[bb\$]
 $\rightarrow q_b$ stack:[b\$] (Dead configuration)

aab: q_0 stack:[\$] $\rightarrow q_a$ stack:[b\$] $\rightarrow q_0$ stack:[b\$] $\rightarrow q_b$ stack:[\$] $\rightarrow q_f$ stack:[] (Accepted)

Conclusion

In conclusion, the language L is a context-free language because an equivalent PDA can be constructed to accept the same language as L .