COP 3402 Systems Software

Compilers And Interpreters

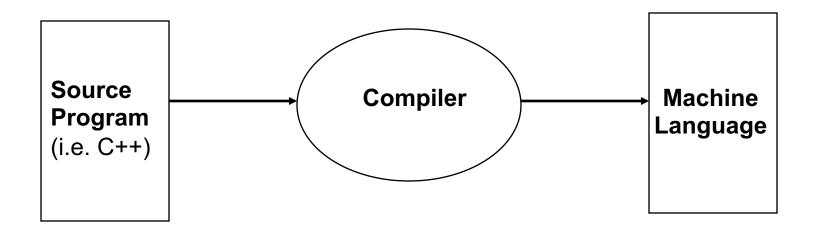
Outline

- 1. Compiler and interpreters
- 2. Compilation process
- 3. Interpreters
- 4. PL/0 Symbols (tokens)

Compilers / Interpreters

- Programming languages are notations for describing computations to people and to machines.
- Programming languages can be implemented by any of three general methods:
 - 1. Compilation
 - 2. Interpretation
 - 3. Hybrid Implementation

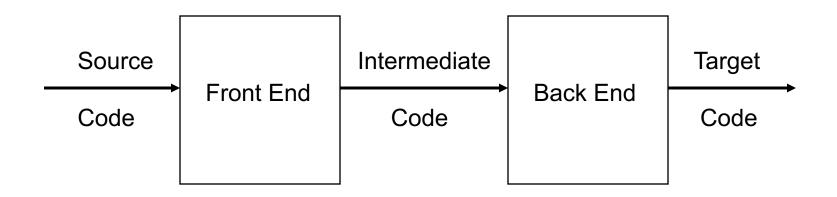
A compiler is a program that takes high level languages (i.e. Pascal, C, ML)as input, and translates it to a low-level representation (machine L\language).

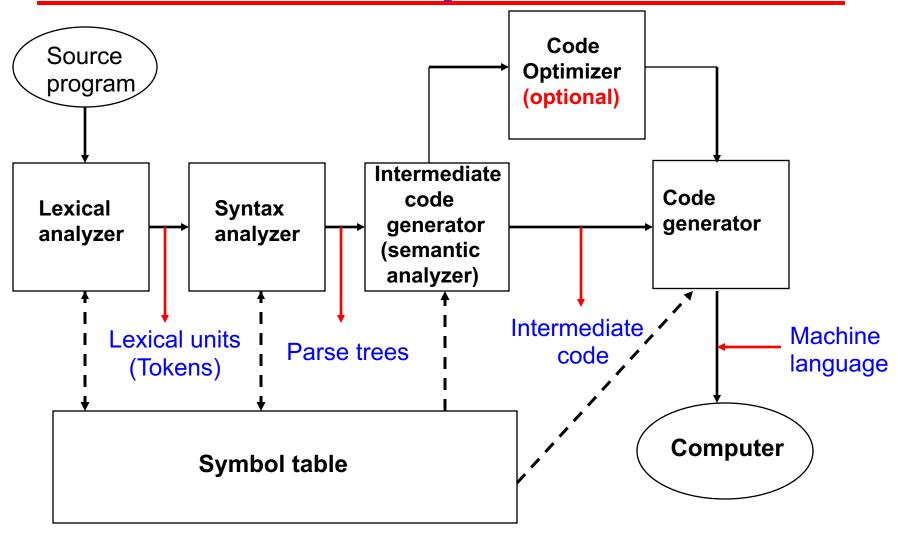


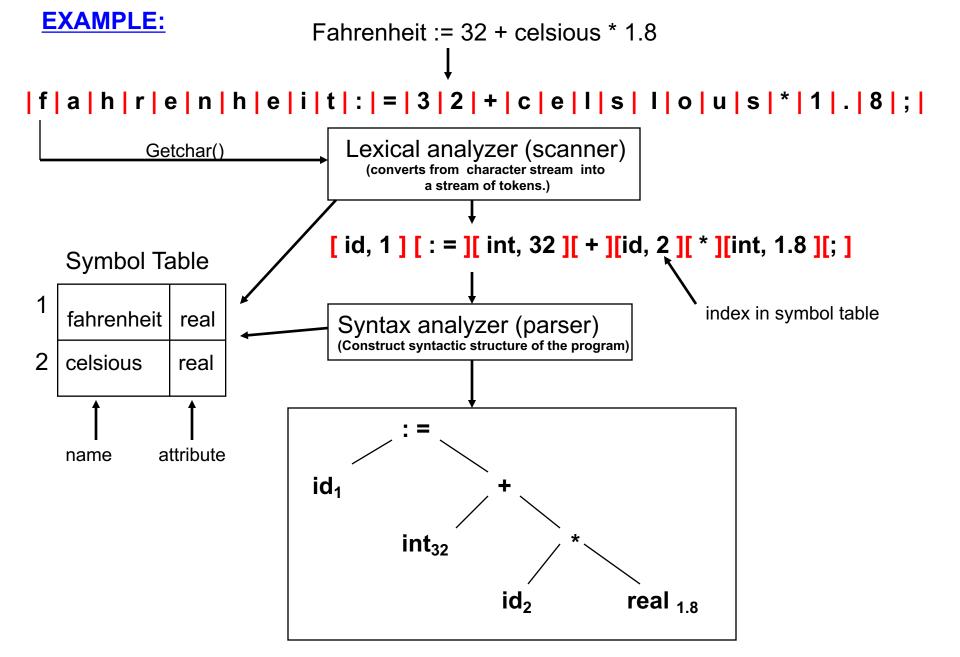
The process of compilation and program execution take place in several phases:

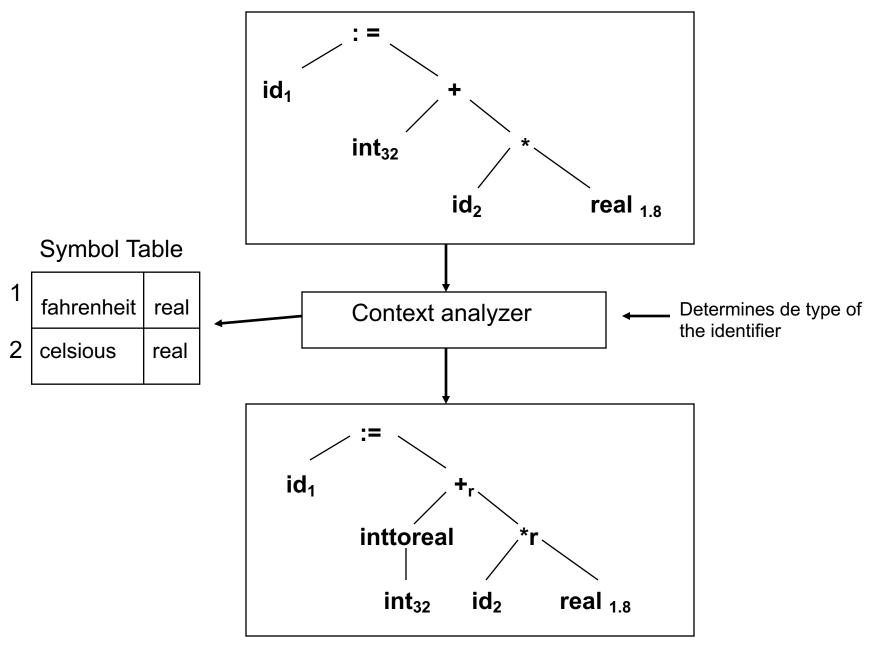
<u>Front end:</u> Scanner → Parser → Semantic Analyzer

Back end: Code generator

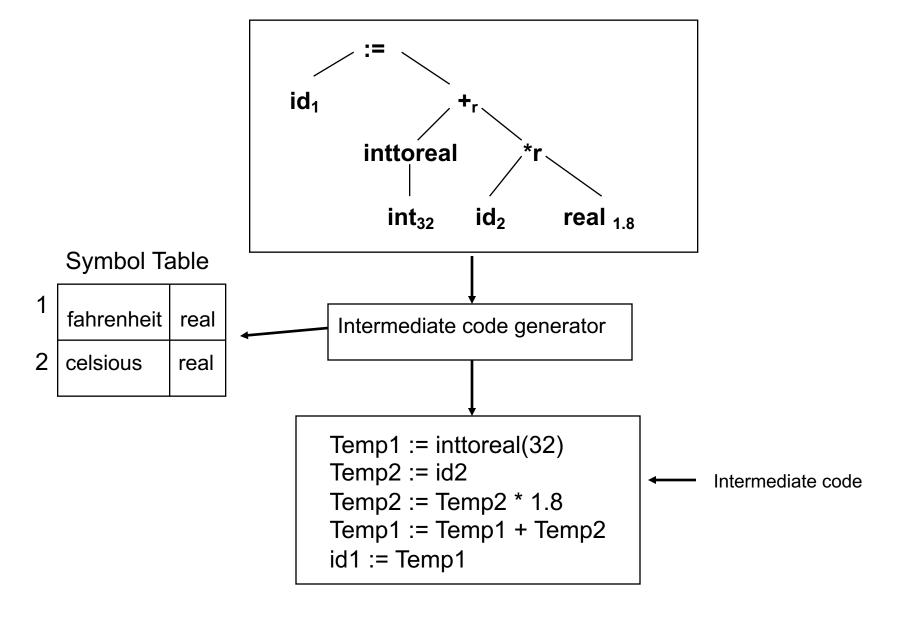


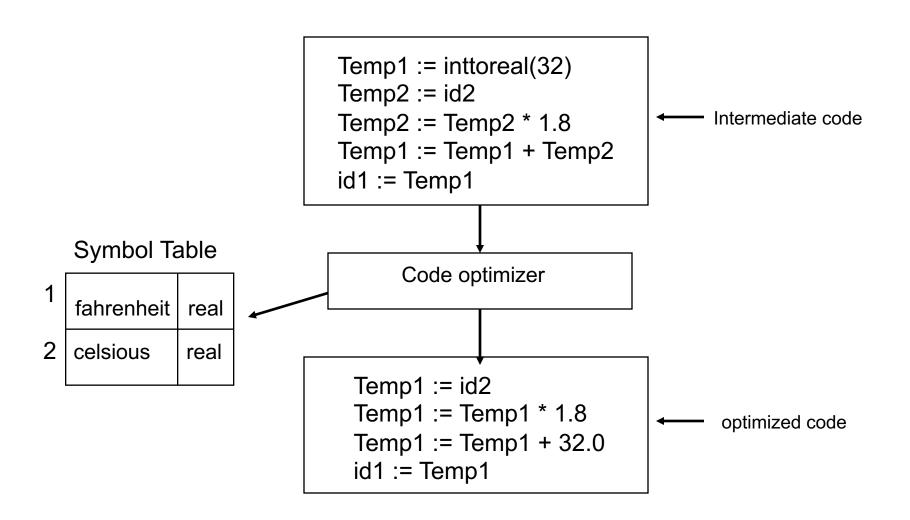


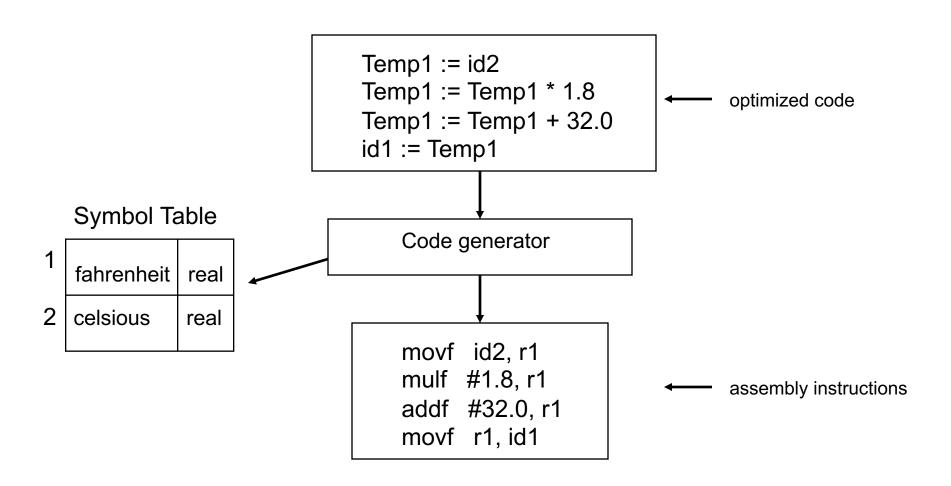




University of Central Florida







Lexical analyzer:

Gathers the characters of the source program into lexical units.

Lexical units of a program are:

identifiers

special words (reserved words)

operators

special symbols

Comments are ignored!

Syntax analyzer:

Takes lexical units from the lexical analyzer and use them to construct a hierarchical structure called **parse tree**

Parse trees represent the syntactic structure of the program.

Intermediate code:

Produces a program in a different lenguage representation:

Assembly language

Similar to assembly language

Something higher than assembly language

Note: semantic analysis is an integral part of the intermediate code generator

Optimization:

Makes programs smaller or faster or both.

Most optimization is done in the intermediate code. (i.e. tree reduction, vectorization)

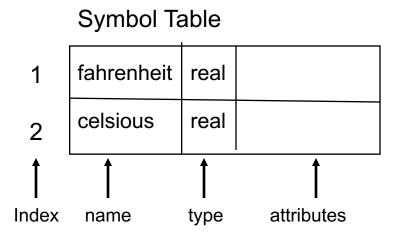
Code generator:

Translate the optimized intermediate code into machine language.

The symbol table:

Serve as a database for the compilation process.

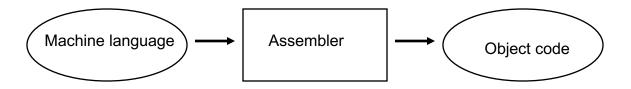
Contents type and attribute information of each user-defined name in the program.



Machine language

A program in its machine language form needs in general

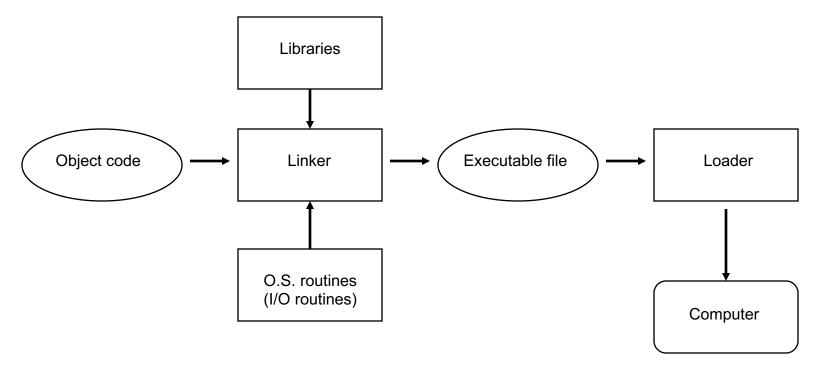
- -- To be translated to object code for execution
- -- To translate the program from its machine language form (assembly language) into object code, an assembler is required.
- -- An assembler is a program that translate machine code into object code



Machine language

To run a program in its object code form, it needs in general

- -- some other code (libraries)
- -- programs from the O.S. (i.e. input/output routines)



Interpreters

Programs are interpreted (executed) by another program called the interpreter.

Advantages: Easy implementation of many source-level debugging operations, because all run-time errors operations refer to source-level units.

Disadvantages: 10 to 100 times slower because statements are interpreted each time the statement is executed.

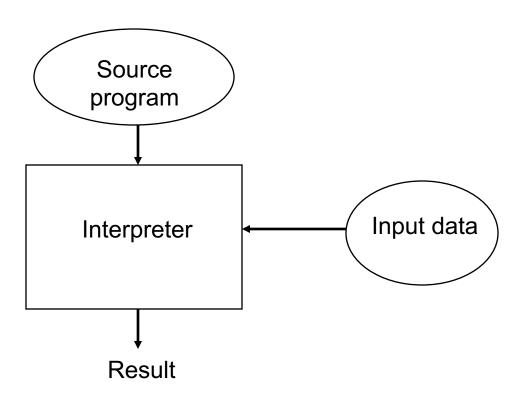
Background:

Early sixties → APL, SNOBOL, Lisp.

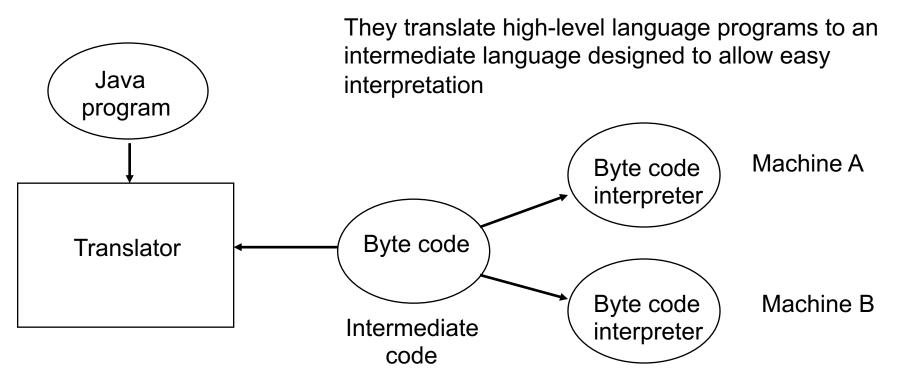
By the 80s \rightarrow rarely used.

Recent years → Significant comeback (some Web scripting languages: JavaScritp, php)

Interpreters



Hybrid implementation systems



Example: PERL and initial implementations of Java

Interpreters

Just-In-Time (JIT) implementation

Programs are translated to an intermediate language.

During execution, it compiles intermediate language methods into machine code when they are called.

The machine code version is kept for subsequent calls.

.NET and Java programs are implemented with JIT system.

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i,x,y,z,q,r;
procedure mult;
 var a, b;
begin
 a := x; b := y; z := 0;
 while b > 0 do
  begin
    if odd x then z := z+a;
    a := 2*a:
    b := b/2;
  end
end;
begin
x := m;
 y := n;
 call mult;
end.
```

As in any language, in PL/0 we need to identify what is the vocabulary and what are the valid names and special symbols that we accept as valid:

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i,x,y,z,q,r;
procedure mult;
 var a, b;
begin
 a := x; b := y; z := 0;
 while b > 0 do
  begin
    if odd x then z := z+a;
    a := 2*a:
    b := b/2;
  end
end:
begin
x := m;
 y := n;
 call mult;
end.
```

As in any language, in PL/0 we need to identify what is the vocabulary and what are the valid names and special symbols that we accept as valid:

For instance, in the on the example we notice that there are many reserved words (keywords)

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i,x,y,z,q,r;
procedure mult;
 var a, b;
begin
 a := x; b := y; z := 0;
 while b > 0 do
  begin
    if odd x then z := z + a;
    a := 2 * a:
    b := b / 2;
  end
end:
begin
x := m;
 y := n;
 call mult;
end.
```

Also there are some operators and special symbols:

```
a) Operators (+, -, *, /, <, =, >, <=, <>, >=, :=)
```

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i, x, y, z, q, r;
procedure mult;
 var a, b:
begin
 a := x; b := y; z := 0;
 while b > 0 do
  begin
    if odd x then z := z + a;
    a := 2 * a;
    b := b / 2;
  end
end:
begin
x := m;
 y := n;
 call mult;
end.
```

Also there are some operators and special symbols:

a) Operators (+, -, *, /, <, =, >, <=, <>, >=, :=)

b) Special symbols

```
b) Special symbols
     ( , ),[,],,,.,;;
```

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i, x, y, z, q, r;
procedure mult;
 var a, b:
begin
 a := x; b := y; z := 0;
 while b > 0 do
  begin
    if odd x then z := z + a;
    a := 2 * a;
    b := b / 2:
  end
end:
begin
x := m;
 y := n;
 call mult;
end.
```

```
There are also:
Numerals such as : 5, 0, 85, 2, 346, . . .
```

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i, x, y, z, q, r;
procedure mult;
 var a, b:
begin
 a := x; b := y; z := 0;
 while b > 0 do
  begin
    if odd x then z := z + a:
    a := 2 * a:
    b := b / 2;
  end
end:
begin
x := m;
y := n;
 call mult;
end.
```

```
There are also:
Numerals such as: 5, 0, 85, 2, 346, ...

And names (identifiers):
A letter
or a letter followed by more letters
or a letter followed by more letters or digits.
```

Examples: x, m, celsious, mult, intel486

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i,x,y,z,q,r;
                                      In addition there are also:
procedure mult;
                                      Comments:
 var a, b;
                                                  /* in C */
begin
 a := x; b := y; z := 0;
                                                  (* in Pascal *)
 while b > 0 do
  begin
                                      Separators:
   if odd x then z := z+a;
                                           white spaces
   a := 2*a; <
   b := b/2;
                                           invisible characters like: tab "\t"
  end
                                                                          new line "\n"
end:
begin
                                           Example: \t a :=
x := m;
y := n;
 call mult;
end.
```

Given the following program written in PL/0:

```
const m = 7, n = 85;
var i,x,y,z,q,r;
procedure mult;
 var a, b;
begin
 a := x; b := y; z := 0;
 while b > 0 do
  begin
    if odd x then z := z+a:
    a := 2*a:
    b := b/2;
  end
end;
begin
x := m;
y := n;
 call mult;
end.
```

Every language has an alphabet (a finite set of characters)

```
PL/0 alphabet { a, b, c, d, e , f, g, h, i, j, k, l , m ,n, o, p q, r, s, t, u, v, w, x, y, z, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, , +, -, *, /, <, =, >, :, . , , , ; }
```

Using concatenation (joining two or more characters) we obtain a string of symbols.

A language L, is simply <u>any</u> set of strings over a fixed alphabet.

```
Alphabet
                          Languages
                        {0,10,100,1000,100000...}
{0,1}
                        {0,1,00,11,000,111,...}
                       {abc,aabbcc,aaabbbccc,...}
\{a,b,c\}
\{A, \ldots, Z\}
                       {TEE,FORE,BALL,...}
                       {FOR,WHILE,GOTO,...}
{A,...,Z,a,...,z,0,...9,}
                       { All legal PASCAL progs}
 +,-,....<>,...}
                       { All grammatically correct
                         English sentences }
```

Special Languages: Ø - EMPTY LANGUAGE

 \in - contains \in string only

The purpose of the lexical analyzer (scanner) is to decompose the source program into its elementary symbols or tokens:

- 1. Read input characters of the source program.
- 2. Group them into lexemes (a lexeme is a sequence of characters that matches the pattern for a token).
- 3. Produce a token for each lexeme

A lexeme (lowest level syntactic unit) is a sequence of characters in the source program

Scan Input
Remove WS, NL, ...
Identify Tokens
Create Symbol Table
Insert Tokens into ST
Generate Errors
Send Tokens to Parser

A lexeme (lowest level syntactic unit) is a sequence of characters in the source program

ASCII Character Set

X

The ordinal number of a character ch is computed from its coordinates (X,Y) in the table as:

$$ord(ch) = 16 * X + Y$$

Example:

	0	1	2	3	4	5	6	7
0	NUL	DLE	SP	0	@	P	•	р
1	SOH	DC1	!	1	A	Q	a	q
2	STX	DC2	**	2	В	R	b	r
3	ETX	DC3	#	3	C	S	c	s
4	EOT	DC4	\$	4	D	T	d	t
5	ENQ	NAK	%	5	E	U	e	u
6	ACK	SYN	&	6	F	V	f	v
7	BEL	ЕТВ	•	7	G	W	g	W
8	BS	CAN	(8	Н	X	h	X
9	НТ	EM)	9	I	Y	i	y
10(A)	LF	SUB	*	:	J	Z	j	Z
11(B)	VT	ESC	+	;	K	[k	{
12(C)	FF	FS	,	<	L	\	l	
13(D)	CR	GS	1	=	M]	m	}
14(E)	so	RS	٠	>	N	۸	n	~
15(F)	SI	US	/	?	0	_	0	DEL

ASCII character table

Dec	Hex	ASCII	Dec	Hex	ASCII	Dec	Hex	ASCII
0	00	NUL (null)	16	10	DLE (data link escape)	32	20	SP (space)
1	01	SOH (start of heading)	17	11	DC1 (device control 1)	33	21	!
2	02	STX (start of text)	18	12	DC2 (device control 2)	34	22	"
3	03	ETX (end of text)	19	13	DC3 (device control 3)	35	23	#
4	04	EOT (end of transmission)	20	14	DC4 (device control 4)	36	24	\$
5	05	ENQ (enquiry)	21	15	NAK (negative acknowledge)	37	25	%
6	06	ACK (acknowledge)	22	16	SYN (synchronous idle)	38	26	&
7	07	BEL (bell)	23	17	ETB (end of transmission block)	39	27	•
8	08	BS (backspace)	24	18	CAN (cancel)	40	28	(
9	09	HT (horizontal tab)	25	19	EM (end of medium)	41	29)
10	0A	LF (line feed)	26	1A	SUB (substitute)	42	2A	*
11	0B	VT (vertical tab)	27	1B	ESC (escape)	43	2B	+
12	0C	FF (form feed)	28	1C	FS (file separator)	44	2C	,
13	0D	CR (carriage return)	29	1D	GS (group separator)	45	2D	-
14	0E	SO (shift out)	30	1E	RS (record separator)	46	2E	
15	0F	SI (shift in)	31	1F	US (unit separator)	47	2F	/

ASCII character table

Dec	Hex	ASCII	Dec	Hex	ASCII	Dec	Hex	ASCII
48	30	0	64	40	@	80	50	P
49	31	1	65	41	A	81	51	Q
50	32	2	66	42	В	82	52	R
51	33	3	67	43	C	83	53	S
52	34	4	68	44	D	84	54	T
53	35	5	69	45	E	85	55	U
54	36	6	70	46	F	86	56	V
55	37	7	71	47	G	87	57	W
56	38	8	72	48	Н	88	58	X
57	39	9	73	49	I	89	59	Y
58	3A	:	74	4A	J	90	5A	Z
59	3B	;	75	4B	K	91	5B	[
60	3C	<	76	4C	L	92	5C	\
61	3D	=	77	4D	M	93	5D]
62	3E	>	78	4 E	N	94	5E	^
63	3F	?	79	4F	О	95	5F	_

ASCII character table

Dec	Hex	ASCII	Dec	Hex	ASCII
96	60	•	112	70	p
97	61	a	113	71	q
98	62	b	114	72	r
99	63	c	115	73	S
100	64	d	116	74	t
101	65	e	117	75	u
102	66	f	118	76	V
103	67	g	119	77	W
104	68	h	120	78	X
105	69	i	121	79	y
106	6A	j	122	7A	Z
107	6B	k	123	7B	{
108	6C	1	124	7C	
109	6D	m	125	7D	}
110	6E	n	126	7E	~
111	6F	0	127	7F	DEL