Analysis Report

jacobi_gpu1(int, double, int, double*, double*, double*)

Duration	1.67398 ms (1,673,976 ns)
Grid Size	[1,1,1]
Block Size	[1,1,1]
Registers/Thread	40
Shared Memory/Block	0 B
Shared Memory Executed	0 B
Shared Memory Bank Size	4 B

[0] Tesla V100-PCIE-32GB

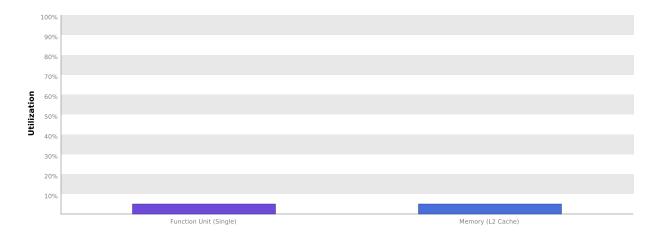
GPU-ced0b34e-8a1e-4f06-af19-2cde27755034
7.0
1024
2048
48 KiB
96 KiB
65536
65536
[2147483647, 65535, 65535]
[1024, 1024, 64]
64
32
28.262 TeraFLOP/s
14.131 TeraFLOP/s
7.066 TeraFLOP/s
80
1.38 GHz
true
4
32
898.048 GB/s
31.719 GiB
64 KiB
6 MiB
7
3
8 Gbit/s
16

1. Compute, Bandwidth, or Latency Bound

The first step in analyzing an individual kernel is to determine if the performance of the kernel is bounded by computation, memory bandwidth, or instruction/memory latency. The results below indicate that the performance of kernel "jacobi_gpu1" is most likely limited by instruction and memory latency. You should first examine the information in the "Instruction And Memory Latency" section to determine how it is limiting performance.

1.1. Kernel Performance Is Bound By Instruction And Memory Latency

This kernel exhibits low compute throughput and memory bandwidth utilization relative to the peak performance of "Tesla V100-PCIE-32GB". These utilization levels indicate that the performance of the kernel is most likely limited by the latency of arithmetic or memory operations. Achieved compute throughput and/or memory bandwidth below 60% of peak typically indicates latency issues.



2. Instruction and Memory Latency

Instruction and memory latency limit the performance of a kernel when the GPU does not have enough work to keep busy. The results below indicate that the GPU does not have enough work because the kernel does not execute enough blocks.

2.1. Grid Size Too Small To Hide Compute And Memory Latency

The kernel does not execute enough blocks to hide memory and operation latency. Typically the kernel grid size must be large enough to fill the GPU with multiple "waves" of blocks. Based on theoretical occupancy, device "Tesla V100-PCIE-32GB" can simultaneously execute 32 blocks on each of the 80 SMs, so the kernel may need to execute a multiple of 2560 blocks to hide the compute and memory latency. If the kernel is executing concurrently with other kernels then fewer blocks will be required because the kernel is sharing the SMs with those kernels.

Optimization: Increase the number of blocks executed by the kernel.

2.2. GPU Utilization May Be Limited By Block Size

Theoretical occupancy is less than 100% but is large enough that increasing occupancy may not improve performance. You can attempt the following optimization to increase the number of warps on each SM but it may not lead to increased performance.

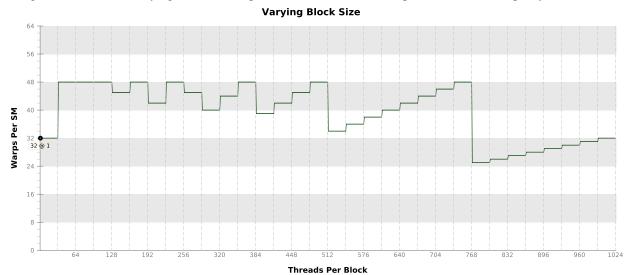
The kernel has a block size of 1 thread. This block size is likely preventing the kernel from fully utilizing the GPU. Device "Tesla V100-PCIE-32GB" can simultaneously execute up to 32 blocks on each SM. Because each block uses 1 warp to execute the block's 1 thread, the kernel is using only 32 warps on each SM. Chart "Varying Block Size" below shows how changing the block size will change the number of warps that can execute on each SM.

Optimization: Increase the number of threads in each block to increase the number of warps that can execute on each SM.

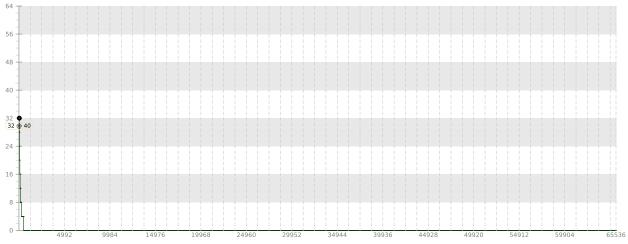
Variable	Achieved	Theoretical	Device Limit	Grid Size: [1,1,1] (1 block) Block Size: [1,1,1] (1 thread)
Occupancy Per SM				
Active Blocks		32	32	0 3 6 9 12 15 18 21 24 27 30 32
Active Warps	1	32	64	0 7 14 21 28 35 42 49 56 664
Active Threads		1024	2048	0 256 512 768 1024 1280 1536 1792 2048
Occupancy	1.6%	50%	100%	
Warps				
Threads/Block		1	1024	0 128 256 384 512 640 768 896 1024
Warps/Block		1	32	0 3 6 9 12 15 18 21 24 27 30 32
Block Limit		64	32	0 3 6 9 12 15 18 21 24 27 30 32
Registers				
Registers/Thread		40	65536	0 8192 16384 24576 32768 40960 49152 57344 65536
Registers/Block		1280	65536	0 16k 32k 48k 64k
Block Limit		48	32	0 3 6 9 12 15 18 21 24 27 30 32
Shared Memory				
Shared Memory/Block		0	98304	0 32k 64k 96k
Block Limit		0	32	0 3 6 9 12 15 18 21 24 27 30 32

2.3. Occupancy Charts

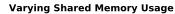
The following charts show how varying different components of the kernel will impact theoretical occupancy.

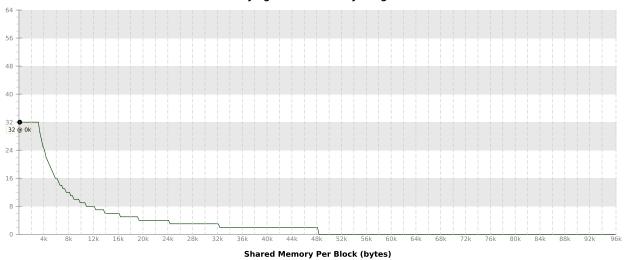






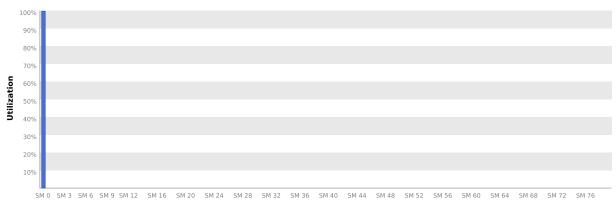
Registers Per Thread





2.4. Multiprocessor Utilization

The kernel's blocks are distributed across the GPU's multiprocessors for execution. Depending on the number of blocks and the execution duration of each block some multiprocessors may be more highly utilized than others during execution of the kernel. The following chart shows the utilization of each multiprocessor during execution of the kernel.



Multiprocessor

3. Compute Resources

GPU compute resources limit the performance of a kernel when those resources are insufficient or poorly utilized. Compute resources are used most efficiently when all threads in a warp have the same branching and predication behavior. The results below indicate that a significant fraction of the available compute performance is being wasted because branch and predication behavior is differing for threads within a warp.

3.1. Low Warp Execution Efficiency

Warp execution efficiency is the average percentage of active threads in each executed warp. Increasing warp execution efficiency will increase utilization of the GPU's compute resources. The kernel's maximum warp execution efficiency is 3.1% because the number of threads per block is not a multiple of the warp size.

Optimization: Reduce the amount of intra-warp divergence and predication in the kernel.

3.2. Function Unit Utilization

Different types of instructions are executed on different function units within each SM. Performance can be limited if a function unit is over-used by the instructions executed by the kernel. The following results show that the kernel's performance is not limited by overuse of any function unit.

Load/Store - Load and store instructions for shared and constant memory.

Texture - Load and store instructions for local, global, and texture memory.

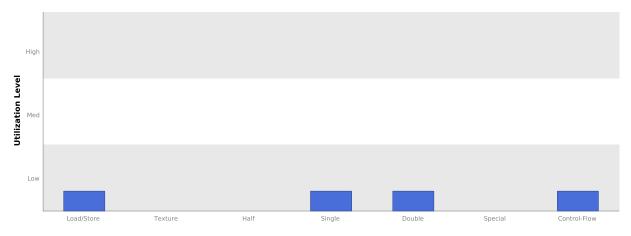
Half - Half-precision floating-point arithmetic instructions.

Single - Single-precision integer and floating-point arithmetic instructions.

Double - Double-precision floating-point arithmetic instructions.

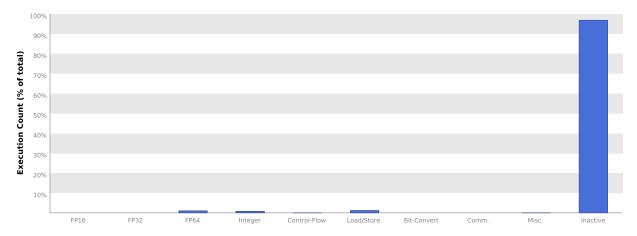
Special - Special arithmetic instructions such as sin, cos, popc, etc.

Control-Flow - Direct and indirect branches, jumps, and calls.



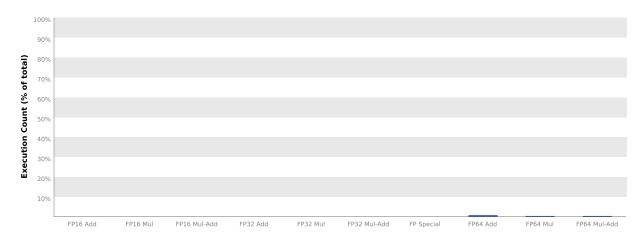
3.3. Instruction Execution Counts

The following chart shows the mix of instructions executed by the kernel. The instructions are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing instructions in that class. The "Inactive" result shows the thread executions that did not execute any instruction because the thread was predicated or inactive due to divergence.



3.4. Floating-Point Operation Counts

The following chart shows the mix of floating-point operations executed by the kernel. The operations are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing operations in that class. The results do not sum to 100% because non-floating-point operations executed by the kernel are not shown in this chart.

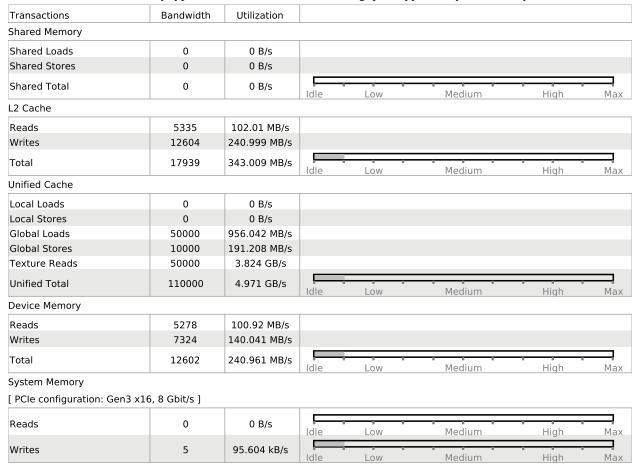


4. Memory Bandwidth

Memory bandwidth limits the performance of a kernel when one or more memories in the GPU cannot provide data at the rate requested by the kernel.

4.1. Memory Bandwidth And Utilization

The following table shows the memory bandwidth used by this kernel for the various types of memory on the device. The table also shows the utilization of each memory type relative to the maximum throughput supported by the memory.



4.2. Memory Statistics

The following chart shows a summary view of the memory hierarchy of the CUDA programming model. The green nodes in the diagram depict logical memory space whereas blue nodes depicts actual hardware unit on the chip. For the various caches the reported percentage number states the cache hit rate; that is the ratio of requests that could be served with data locally available to the cache over all requests made.

The links between the nodes in the diagram depict the data paths between the SMs to the memory spaces into the memory system. Different metrics are shown per data path. The data paths from the SMs to the memory spaces report the total number of memory instructions executed, it includes both read and write operations. The data path between memory spaces and "Unified Cache" or "Shared Memory" reports the total amount of memory requests made (read or write). All other data paths report the total amount of transferred memory in bytes.