

## Experiment No-07

**AIM:** Write a program to implement link state /Distance vector routing protocol to find suitable path for transmission.

### OBJECTIVES:

- To study TCP protocol.
- To implement socket programming.
- To use TCP socket for wired network.

### PROBLEM STATEMENT:

Write a program to implement link state /Distance vector routing protocol to find suitable path for transmission.

### SOFTWARE & HARDWARE REQUIREMENTS:

- Software: gcc compiler & wireshark.
- Hardware: Open source Linux operating system

### THEORY:

A distance-vector routing (DVR) protocol requires that a router inform its neighbors of topology changes periodically. Historically known as the old ARPANET routing algorithm (or known as Bellman-Ford algorithm).

**Bellman Ford Basics –** Each router maintains a Distance Vector table containing the distance between itself and ALL possible destination nodes. Distances, based on a chosen metric, are computed using information from the neighbors' distance vectors.

#### Distance Vector Algorithm –

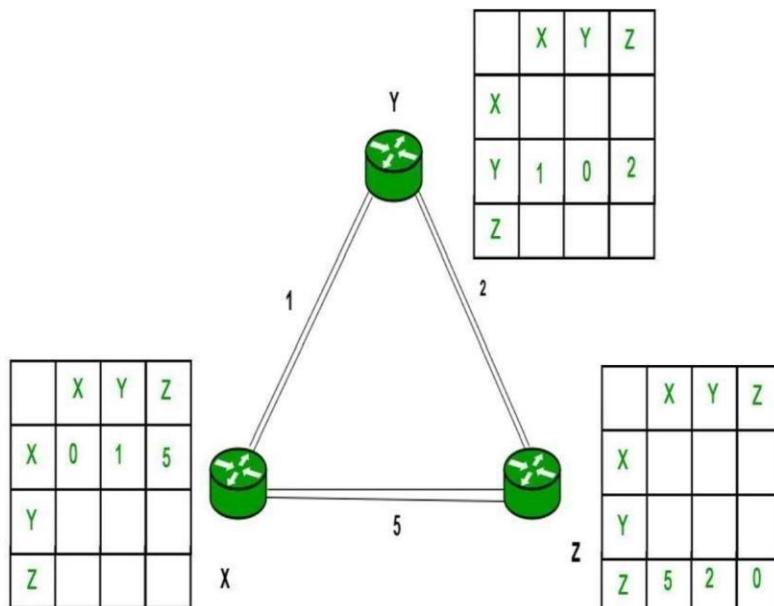
1. A router transmits its distance vector to each of its neighbors in a routing packet.

2. Each router receives and saves the most recently received distance vector from each of its neighbors.
3. A router recalculates its distance vector when:
  - It receives a distance vector from a neighbor containing different information than before.
  - It discovers that a link to a neighbor has gone down.
  - From time-to-time, each node sends its own distance vector estimate to neighbors.  $\square$

When a node  $x$  receives new DV estimate from any neighbor  $v$ , it saves  $v$ 's distance vector and it updates its own DV using B-F equation:

  - $Dx(y) = \min \{ C(x,v) + Dv(y), Dx(y) \}$  for each node  $y \in N$

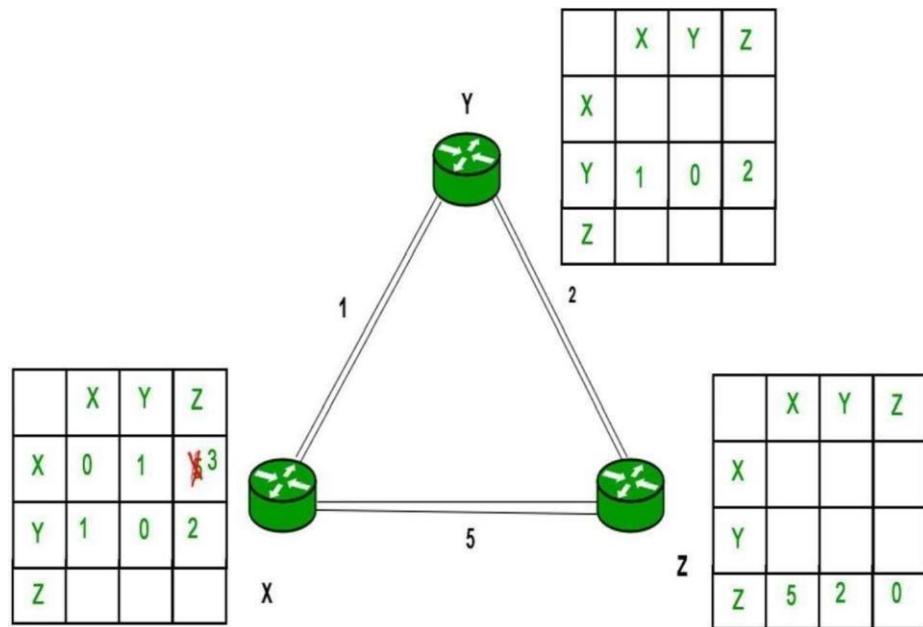
Example – Consider 3-routers X, Y and Z as shown in figure. Each router have their routing table. Every routing table will contain distance to the destination nodes.



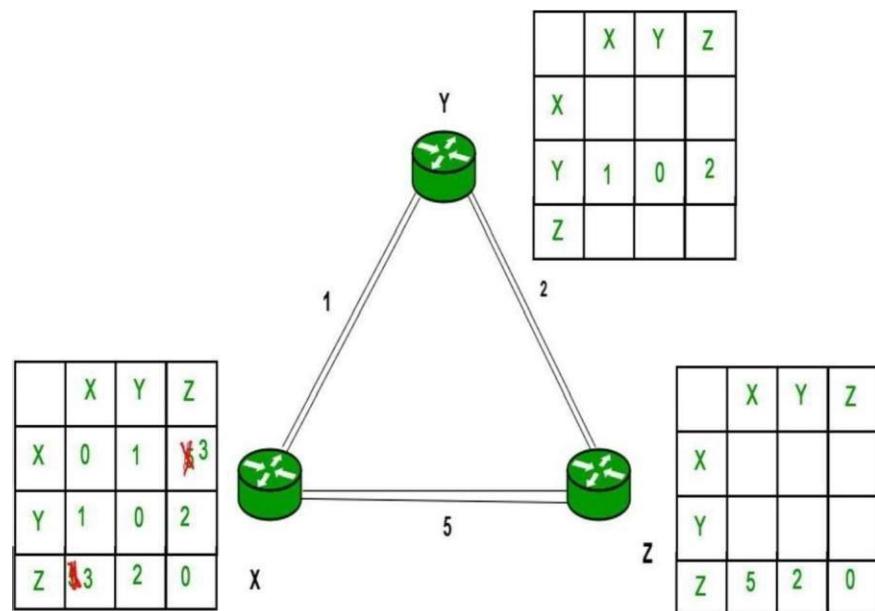
Consider router X , X will share its routing table to neighbors and neighbors will share its routing table to it to X and distance from node X to destination will be calculated using bellmen- ford equation.

$$Dx(y) = \min \{ C(x,v) + Dv(y) \} \text{ for each node } y \in N$$

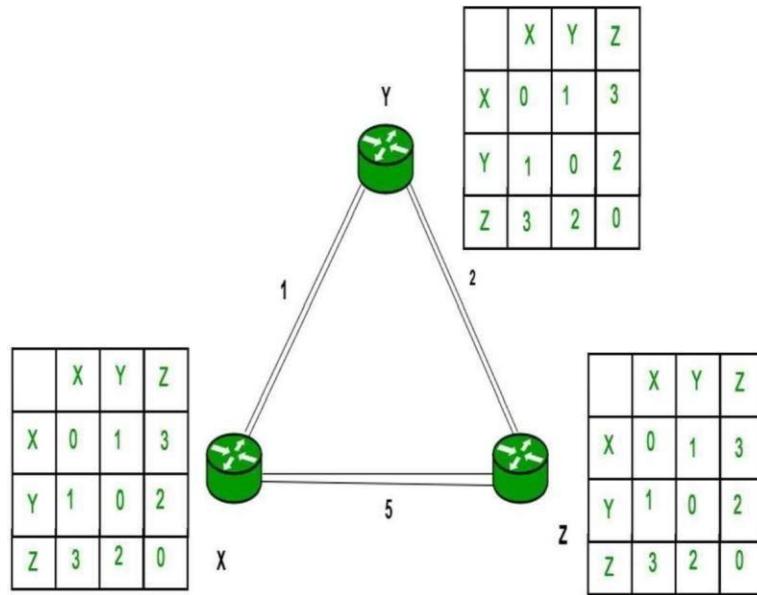
As we can see that distance will be less going from X to Z when Y is intermediate node(hop) so it will be update in routing table X.



Similarly for Z also –



Finally the routing table for all –



#### Advantages of Distance Vector routing – □

It is simpler to configure and maintain than link state routing.

#### Disadvantages of Distance Vector routing – □

It is slower to converge than link

state.

- It is at risk from the count-to-infinity problem.
- It creates more traffic than link state since a hop count change must be propagated to all routers and processed on each router. Hop count updates take place on a periodic basis, even if there are no changes in the network topology, so bandwidth-wasting broadcasts still occur.
- For larger networks, distance vector routing results in larger routing tables than link state since each router must know about all other routers. This can also lead to congestion on WAN links.

**Write a program to implement link state /Distance vector routing protocol to find suitable path for transmission.**

```
void find(int x,int y) COMPUTER NETWORKS LABORATORY { int i,j; i=x-1; j=y-1; printf("%d->",x);

if(n[i].t[j][3]!=j) { find(n[i].t[j][3]+1,y); return; } }
```

```
Enter the no of nodes required:3
Enter the dists from the nodes 1 to other node...
Pls enter 999 if there is no direct
    Enter dist to node 2=10
    Enter dist to node 3=999
Enter the dists from the nodes 2 to other node...
Pls enter 999 if there is no direct
    Enter dist to node 1=999
    Enter dist to node 3=15
Enter the dists from the nodes 3 to other node...
Pls enter 999 if there is no direct
    Enter dist to node 1=20
    Enter dist to node 2=25
The configuration of the nodes after initialization is as follows:
The routing table for node no 1 is as follows
      DESTINATION      DISTANCE      NEXT_HOP
      1              0                  1
      2              10                 2
      3          NO LINK            NO HOP
The routing table for node no 2 is as follows
      DESTINATION      DISTANCE      NEXT_HOP
      1          NO LINK            NO HOP
      2              0                  2
      3              15                 3
The routing table for node no 3 is as follows
      DESTINATION      DISTANCE      NEXT_HOP
      1              20                 1
      2              25                 2
      3              0                  3
The config of the nodes after the comp of the paths is as follows:
The routing table for node no 1 is as follows
      DESTINATION      DISTANCE      NEXT_HOP
      1              0                  1
      2              10                 2
      3              25                 3
The routing table for node no 2 is as follows
      DESTINATION      DISTANCE      NEXT_HOP
      1              35                 3
      2              0                  2
      3              15                 3
The routing table for node no 3 is as follows
      DESTINATION      DISTANCE      NEXT_HOP
      1              20                 1
      2              25                 2
      3              0                  3
Enter 0 to exit or any other key to find the shortest path:1
Enter the nodes b/w which path is to be found:1 3
The most suitable route from node 1 to 3 is as follows
--->2--->3
The length of the shortest path between node 1 & 3 is 25
Enter 0 to exit or any other key to find the shortest path:0
```

## CONCLUSION: -