

In the computer world, use restricted resource you have to generate maximum benefit is what we always want to pursue.

For now, suppose you are a dominator of m 0s and n 1s respectively. On the other hand, there is an array with strings consisting of only 0s and 1s.

Now your task is to find the maximum number of strings that you can form with given m 0s and n 1s. Each 0 and 1 can be used at most **once**.

Note:

1. The given numbers of 0s and 1s will both not exceed 100
2. The size of given string array won't exceed 600.

Example 1:

Input: Array = {"10", "0001", "111001", "1", "0"}, m = 5, n = 3

Output: 4

Explanation: This are totally 4 strings can be formed by the using of 5 0s and 3 1s, which are "10", "0001", "1", "0"

Example 2:

Input: Array = {"10", "0", "1"}, m = 1, n = 1

Output: 2

Explanation: You could form "10", but then you'd have nothing left. Better form "0" and "1".

Intution :

This problem wants us to **select maximum** no of strings **given some number of ones and zeroes** .

So this is a typical 0/1 Knapsack problem

But the twist is that it has 2 constraints

approach :

for every string :

starting from 0 number of **ones** and 0 number of **zeroes**

To m number of **zeroes** and n number **ones**

find the number of strings possible with this much resource

using recursive formula :

$dp[i][j][k] = \max(dp[i-1][j][k], 1 + dp[i-1][j-z][k-o])$

//If you pay attention the recursive formula is exactly same as 0/1 knapsack except one more additional constraint where :

i <--- tracks string

j <--- tracks number of zeroes for i th string

k <--- tracks number of ones for j th string

z <--- number of zeroes in string i

o <--- number of ones in string i

if(o > j || z > k) then $dp[i][j][k] = dp[i-1][j][k]$ // implying resources are not enough

CODE :

class Solution {

public:

int findMaxForm(vector<string>& strs, int m, int n) {

if(strs.size() == 0) return 0;

int dp[strs.size()+1][m+1][n+1];

for(int s = 0; s <= strs.size(); s++)

for(int i=0; i<=m; i++)

for(int j=0; j<=n; j++){

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        if(s == 0 || (j==0 && i==0)) dp[s][i][j] = 0;
        else{
            int ct0 = 0, ct1 = 0;
            for(char a : strs[s-1])
            {
                if(a == '0')
                    ct0++;
                else
                    ct1++;
            }
            if(ct0 > i || ct1 > j) dp[s][i][j] = dp[s-1][i][j];
            else
                dp[s][i][j] = max(1 + dp[s-1][i-ct0][j-ct1], dp[s-1][i][j]);
        }
    }
    return dp[strs.size()][m][n];
}
};

```