WEEK3-RogerThat WP

萎了萎了,没时间做pwn了(游戏玩多了)

2.Re

2.FAKE

乱打一通,拿了一个假flag,没注意到init函数对调试器操作了一波。

```
int sub_406A11()
{
  int result; // eax
  char buf[100]; // [rsp+0h] [rbp-70h] BYREF
  int fd; // [rsp+64h] [rbp-Ch]
  char *i; // [rsp+68h] [rbp-8h]

  fd = open("/proc/self/status", 0);
  read(fd, buf, 0x64uLL);
  for ( i = buf; *i != 84 || i[1] != 114 || i[2] != 97 || i[3] != 99 || i[4] != 101 || i[5] != 114; ++i )
    ;
  result = atoi(i + 11);
  if ( !result )
    result = sub_40699B();
  return result;
}
```

这里应该是检测了是否有debugger,通过文件系统来达成,在windows中估计就是lsDebugger了,也可能有别的骚操作,暂且没见着,大胆猜测可以有多进程通信检测什么的。

```
int64 sub_40699B()
{
    int64 result; // rax
    unsigned int i; // [rsp+Ch] [rbp-4h]

mprotect(&dword_400000, 0x100000uLL, 7);
for ( i = 0; ; ++i )
{
    result = i;
    if ( i > 0x43E )
        break;
    *((_BYTE *)sub_401216 + (int)i) ^= byte_409080[i];
}
return result;
}
```

针对这个比较简单的加密函数有脚本

```
#include <idc.idc>

static main(){
    auto addr = 0x401216;
    auto i = 0;
    for(i = 0;i<0x43E;i++){
        PatchByte(addr+i,Byte(addr+i)^Byte(0x409080+i));
    }
}</pre>
```

解密后得到真正的比较函数

```
__int64 __fastcall sub_401216(__int64 a1) {
```

```
int v2[36]; // [rsp+8h] [rbp-1D0h]
  int v3[36]; // [rsp+98h] [rbp-140h]
  int v4[38]; // [rsp+128h] [rbp-B0h] BYREF
  int m; // [rsp+1C0h] [rbp-18h]
  int 1; // [rsp+1C4h] [rbp-14h]
  int k; // [rsp+1C8h] [rbp-10h]
  int j; // [rsp+1CCh] [rbp-Ch]
  int i; // [rsp+1D0h] [rbp-8h]
  unsigned int v10; // [rsp+1D4h] [rbp-4h]
  memset(v4, 0, 0x90uLL);
  v3[0] = 55030; v3[1] = 61095; v3[2] = 60151; v3[3] = 57247; v3[4] = 56780; v3[5] =
55726;
  v3[6] = 46642; v3[7] = 52931; v3[8] = 53580; v3[9] = 50437; v3[10] = 50062; v3[11]
= 44186;
  v3[12] = 44909; v3[13] = 46490; v3[14] = 46024; v3[15] = 44347; v3[16] =
43850; v3[17] = 44368;
  v3[18] = 54990; v3[19] = 61884; v3[20] = 61202; v3[21] = 58139; v3[22] =
57730; v3[23] = 54964;
  v3[24] = 48849; v3[25] = 51026; v3[26] = 49629; v3[27] = 48219; v3[28] =
47904; v3[29] = 50823;
  v3[30] = 46596; v3[31] = 50517; v3[32] = 48421; v3[33] = 46143; v3[34] =
46102; v3[35] = 46744;
  v2[0] = 104; v2[1] = 103; v2[2] = 97; v2[3] = 109; v2[4] = 101; v2[5] = 123;
  v2[6] = 64; v2[7] = 95; v2[8] = 70; v2[9] = 65; v2[10] = 75; v2[11] = 69;
  v2[12] = 95; v2[13] = 102; v2[14] = 108; v2[15] = 97; v2[16] = 103; v2[17] = 33;
  v2[18] = 45; v2[19] = 100; v2[20] = 111; v2[21] = 95; v2[22] = 89; v2[23] = 48;
  v2[24] = 117; v2[25] = 95; v2[26] = 107; v2[27] = 111; v2[28] = 110; v2[29] = 119;
  v2[30] = 95; v2[31] = 83; v2[32] = 77; v2[33] = 67; v2[34] = 63; v2[35] = 125;
  v10 = 1;
  for (i = 0; i \le 5; ++i)
    for (j = 0; j \le 5; ++j)
      for (k = 0; k \le 5; ++k)
        v4[6 * i + j] += v2[6 * k + j] * *(_DWORD *)(4LL * (6 * i + k) + a1);
    }
  for (1 = 0; 1 \le 5; ++1)
    for (m = 0; m \le 5; ++m)
      if (v4[6 * 1 + m] != v3[6 * 1 + m])
        v10 = 0;
    }
  return v10;
}
```

而假的函数却也像真的一样,要是解析不出来数据段嘛,那我倒是会先去想SMC,这点看真的是非常巧妙!!! 后来仔细想了一下,应该是encrypt(f1(x)) == f2(2),只要满足能转换f1和f2,任何加密方式都可以。而异或恰恰是比较简单方便的加密方式。

基本是同一穿代码, 异或以下 (加密以下) 能性成一种看似很像真代码的代码。

最后经过一些探索,发现是一个线性的加密,矩阵乘法。因此得到解密脚本

```
import numpy as np
v2 = [0]*36
v3 = [0]*36
v3[0] = 55030; v3[1] = 61095; v3[2] = 60151; v3[3] = 57247; v3[4] = 56780; v3[5] =
55726;
v3[6] = 46642; v3[7] = 52931; v3[8] = 53580; v3[9] = 50437; v3[10] = 50062; v3[11] =
v3[12] = 44909; v3[13] = 46490; v3[14] = 46024; v3[15] = 44347; v3[16] =
43850; v3[17] = 44368;
v3[18] = 54990; v3[19] = 61884; v3[20] = 61202; v3[21] = 58139; v3[22] =
57730; v3[23] = 54964;
v3[24] = 48849; v3[25] = 51026; v3[26] = 49629; v3[27] = 48219; v3[28] =
47904; v3[29] = 50823;
v3[30] = 46596; v3[31] = 50517; v3[32] = 48421; v3[33] = 46143; v3[34] =
46102;v3[35] = 46744;
v2[0] = 104; v2[1] = 103; v2[2] = 97; v2[3] = 109; v2[4] = 101; v2[5] = 123;
v2[6] = 64; v2[7] = 95; v2[8] = 70; v2[9] = 65; v2[10] = 75; v2[11] = 69;
v2[12] = 95; v2[13] = 102; v2[14] = 108; v2[15] = 97; v2[16] = 103; v2[17] = 33;
v2[18] = 45; v2[19] = 100; v2[20] = 111; v2[21] = 95; v2[22] = 89; v2[23] = 48;
v2[24] = 117; v2[25] = 95; v2[26] = 107; v2[27] = 111; v2[28] = 110; v2[29] = 119;
v2[30] = 95; v2[31] = 83; v2[32] = 77; v2[33] = 67; v2[34] = 63; v2[35] = 125;
v2_{matrix} = []
for i in range(6):
    temp_list = []
    for j in range(6):
        temp_list.append(v2[6*j+i])
    v2_matrix.append(temp_list)
v3_{matrix} = []
for i in range(6):
    temp_list = []
    for j in range(6):
        temp_list.append(v3[6*i+j])
    v3_matrix.append(temp_list)
# print(v2_matrix)
# print(v3_matrix)
v2_matrix = np.array(v2_matrix)
v3_matrix_compute = []
for i in range(6):
    v3_matrix_compute.append(np.array(v3_matrix[i]))
# print(v2_matrix)
# print(np.linalg.inv(v2_matrix))
# print(v3_matrix_compute[0])
for i in range(6):
    v3_matrix_compute[i] = np.array([v3_matrix_compute[i]]).T
    #print(v3_matrix_compute[i])
    result = np.matmul(np.linalg.inv(v2_matrix),v3_matrix_compute[i])
    for i in np.array(result):
        for j in i:
            #print(j,end = '')
            print(chr(round(j)),end = '')
```

对于脚本的一点mark,希望下次能快点写出来,踩了大坑,不要拿4血,真的淦。不过,最重要的应该还是知识点,其实这种操作复杂化一点就变成了加壳,那些脱壳的操作,正是这个进化版本,可能还会用到什么动态库,在seh上设置一个加壳函数,抛出异常再脱壳等等(仅猜测)

```
#初始化矩阵
matrix = np.array(<list>)
#矩阵乘法
result = np.matmul()
#矩阵点乘
result = np.dot()
#一维矩阵转置
matrix = np.array([matrix]).T
#矩阵求逆
matrix_inv = np.linalg.inv(matrix_inv)
```

再mark一波idc https://blog.csdn.net/jazrynwong/article/details/84875699

其实看文档也可以,不过还是中文舒服点<u>https://www.hex-rays.com/products/ida/support/idadoc/16</u>2.shtml

3.helloRe3

跟踪StringOutDebug,也就是debugger函数就会发现这个东西,可以看见



框框中对一个全局变量进行了改变,可见,这里是关键,但是发现我一旦把这地方变成1,它马上就变回07.

当我们跟踪内存时候发现这个&(&addr)[72*i][4*v16+1]其实在内存中代表着对应字符对应的一个数字&(&addr)[72*i][4*v16]则代表内存中的一个指向对应字符的指针

直接在这个地方下硬件写入断点,就能的到被改变的原因了,原来是一个线程搞的鬼。不过不下断点也能猜测到是别的线程下的手脚。不过仔细观察一下线程的一些特征也能找到一些不对劲的地方。



除了主线程以外,还有一个线程在用户代码上花了很多时间,调试过后会发现这个线程的时间比主线程还长,由此可见应该去分析一下那个线程的入口,不出所料。



看见那个地方为0,那么就会在循环里面不断重复。一旦变成1,他就开始了!!!



在右边的方框中先把每个地方取反然后再进入左边,经过动态调试可以发现第一个函数是加密,静态分析后会发现第二个是比较两个是否相等。跟进第一个函数后发现,里面有两个函数,第一个生成一个key,第二个是加密,而key以明文256位存在内存中,直接dump下来就可以了。很奇怪的是为什么我dump了好几次,但是有两个结果??第二天的时候dump下来的key才能用,很奇怪。可能是我之前dump错了地方,搞得我想半天以为是反调试,比较这个代码

最后解密脚本:

```
#include <stdio.h>
#include <stdint.h>

void exchange_a_byte(uint8_t * a,uint8_t * b){
    char v = *a;
    *a = *b;
```

```
*b = v;
}
int main() {
           uint8_t string_0[256] =
{0x90,0x1D,0x7B,0xA2,0x53,0xA5,0x73,0xB6,0xCC,0x32,0x20,0x21,0x28,0x7A,0xE4,0x97
,0xEF,0x41,0x66,0x04,0xF3,0x24,0xA9,0xBB,0x44,0x3E,0xCD,0xA4,0xD5,0x7D,0xFE,0x4B
,0x3A,0x3C,0x8F,0x08,0xB8,0x23,0x2B,0x54,0x6D,0x25,0x47,0xBD,0x4D,0xDC,0xDA,0x2A
,0xBF,0x75,0xB1,0xC3,0x05,0xD2,0x17,0x3F,0x30,0x83,0xFD,0x96,0xE0,0xAD,0xCF,0x57
,0x0E,0x85,0x18,0xB4,0x52,0x86,0x6C,0xE1,0x8C,0x77,0x38,0x8E,0xB0,0x40,0xDF,0x7E
,0xD4,0x78,0xC7,0xF2,0x5C,0x5D,0x39,0x67,0x62,0x27,0x79,0x2E,0x58,0x82,0x2F,0x01
,0x69,0x59,0xAC,0x22,0x76,0x8D,0xDD,0xC1,0xB2,0xE9,0xFA,0x2C,0xE2,0xC4,0x1A,0x0D
, 0 \times C8, 0 \times 09, 0 \times D7, 0 \times 46, 0 \times 02, 0 \times 0C, 0 \times E7, 0 \times A8, 0 \times B5, 0 \times 64, 0 \times 5A, 0 \times A6, 0 \times 1E, 0 \times 5F, 0 \times 6B, 0 \times C6, 0 \times 1E, 0 
,0xAA,0x91,0xB9,0xD8,0xD6,0x48,0x49,0xA7,0x11,0x56,0xFB,0xEC,0x8B,0x1B,0xE3,0x45
,0xAB,0xF6,0x4A,0xDB,0xB7,0x81,0x10,0x42,0x71,0x0F,0x89,0x14,0xD3,0x1F,0x50,0xEE
,0x60,0x33,0x4F,0x7C,0x98,0xAE,0x9A,0xEA,0x8A,0xC5,0x9C,0xF8,0xBA,0xD9,0x1C,0x31
,0xB3,0x68,0x5B,0x80,0x03,0xF9,0x19,0x92,0x3B,0x29,0x3D,0x5E,0xDE,0x51,0xED,0x70
,0x7F,0x0A,0x36,0xD1,0x87,0xAF,0x35,0x16,0x9E,0xC9,0xEB,0xF5,0x37,0xF1,0x61,0x43
,0xE5,0x95,0xA3,0xE8,0x34,0x9D,0xBC,0xCA,0x65,0x13,0xA1,0xF0,0xF7,0xCE,0x94,0x55
,0x6F,0x99,0x84,0x26,0x88,0x00,0xFC,0xD0,0xCB,0x9F,0x93,0xC2,0x6A,0x9B,0x4E,0x2D
,0x74,0xE6,0xBE,0x15,0x72,0xFF,0xF4,0x07,0x4C,0x12,0x6E,0x06,0xC0,0x63,0xA0,0x0B
};
           uint8_t result[20] = \{0x4D, 0xAF, 0x27, 0xAD,
                                                                        0xE1,0xEC,0x6D,0xDA,
                                                                        0xF0,0x31,0x5E,0x9A,
                                                                        0x9E,0x29,0xFA,0xBE,
                                                                        0x6B,0x08,0xC8,0x49;
           uint8_t flag[20] = \{0\};
           int v6 = 0, v7 = 0;
           for(int i = 0; i < 20; i++){
                      v7 = (v7+1)\%256;
                      v6 = (string_0[v7]+v6)%256;
                      exchange_a_byte(string_0+v7,string_0+v6);
                      flag[i] = result[i] \wedge *(uint8_t*)(string_0 + (*(uint8_t*)(string_0+v6)+*)
(uint8_t*)(string_0+v7))%256);
           }
//
                for(int i = 0; i < 256; i++){
                           if(i%16 == 0 && i!= 0)
//
//
                                      printf("\n");
//
                           printf("%x ",string_0[i]);
//
```

```
// putchar('\n');
// for(int i = 0;i<20;i++){
// printf("%x ",flag[i]);
// }

printf("\n");
for(int i = 0;i<20;i++){
 printf("%02d ",flag[i]^0xff);
}

//HGAME{6-K4K.4R+3C4T}</pre>
```

在这个题目中也有很多可以讨论的地方,比如在TLS函数中可以做手脚,检测int断点等等,对于dump那块内存最稳妥的办法是修改程序流,似乎可以用dll注入的方法,不过那个徒手操作确实比较复杂到现在我也掌握的不好(老是算错地址。。。)在WndProc中我们也可以看到其中对输入的一些操作,但是我终究不清楚它的WM COMMOND跑哪去了,那些按钮究竟在哪实现的,这点很奇怪。



我还找到了另外的一个switch分支**sub_1401CC620**,但是很迷惑,只是经过动态调试知道了它在按钮按下的时候走的default路线。在其中还有对xmm0,xmm1这些寄存器的操作,应该是经过优化后得到的结果。

mark: rep stosb就是从EDI所指的内存开始,将连续的ECX个字节写成AL的内容,多用于清零等,每次遇到这个都会忘,然后重新查,就像movsx,movss这些指令一样。

还有一点,程序中很容易产生[addr] = ??? 也就是地址写入的时候出异常,这种时候x64dbg可以用shift+f9忽略异常,也可以在选项中直接把对应异常的代码忽略掉(建议这样做)

4.Crypto

1.rsa

和之前一样,,,这算是白嫖flag嘛,,,都有点不好意思了

python3 ./RsaCtfTool.py -n 153618988782356965 6907198648267160926305328986853052888831678087029867946180609 -e 65537 --uncipher 792718762802605

不过那个网站也分出了素数