

Problem solving 1

- There are many problems that use sorting as like a pre-step.
- This is where you would need to pick a sorting algorithm based on various trade offs.

- Problems we will look at:
 - sum to target problem
 - Intersection of two sorted arrays

sum to target problem

- Input:
 - An array of numbers
 - Target number
- Output:
 - Return any pair of numbers from the input that add up to the target number.
- Example 1:
 - Input: (2, 7, 16, 25) and target = 27
 - Output: 2, 25
- Example 2:
 - Input: (2, 7, 16, 25) and target = 34
 - Output: empty set

sum to target problem

- What are our options for a solution?
- Brute Force?

```
for each number num1 in input
    num2 = target – num1
    Linear search for num2 in input.
    if found
        return (num1, num2)

return empty set
```

Time complexity: ?

Space: ?

What about a better option?

sum to target problem

- How about if we sort the input?
 - What will that cost?

sort input

for each number num1 in input

num2 = target – num1

Binary search for num2 in input.

if found

return (num1, num2)

return empty set

Time complexity: ?

Space: ?

Any other option after having sorted the array?

sum to target problem

Once we sort the array, we could use a 2 indices (pointers).

Traverse sorted array from both ends, each time summing the two numbers and comparing with target.

Input: 7 9 11 14 16 17 21 26

Target: 25

leftElement = 7	rightElement = 26	sum = 33
leftElement = 7	rightElement = 21	sum = 28
leftElement = 7	rightElement = 17	sum = 24
leftElement = 9	rightElement = 17	sum = 26
leftElement = 9	rightElement = 16	sum = 25 Found a solution (9, 16)
Answer found, we are done. Can continue if more pairs r needed. If so, ++ leftIndex or -- rightIndex.		
leftElement = 9	rightElement = 14	sum = 23
leftElement = 11	rightElement = 14	sum = 25 Found a solution (11, 14)
leftElement = 11	rightElement = 11	leftIndex == rightIndex, we scanned the entire input array

sum to target problem

- Once we sort the array, we could use 2 indices (pointers).
 - Traverse sorted array from both ends, each time summing the two numbers and comparing with target.
- Input: 7 9 11 14 16 17 21 26 Target: 25

sort input

leftIndex = 0;

rightIndex = n-1;

// n is size of input

while leftIndex < rightIndex

 sum = input [leftIndex] + input [rightIndex]

 if sum == target

 return (input [leftIndex], input [rightIndex])

 else if (sum > target)

 -- rightIndex

// sum is bigger, so we want smaller numbers

// we move right index towards the left

 else

 ++ leftIndex

// we want bigger numbers

return empty set

Time complexity: Sorting + looking for solution : ?

Space: ?

sum to target problem

- Any other options ?

sum to target problem

- Using a hash table
- Would the following work, or do u see any issues?

HashTable HT;

for each number num1 in input

HT.Insert(num1)

for each number num1 in input

num2 = target - num1

if HT.Contains (num2)

return (num1, num2)

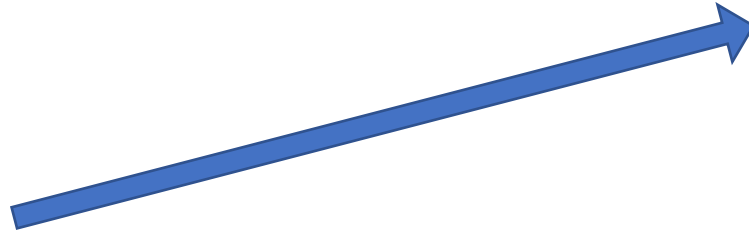
return empty set

sum to target problem

- Lets look at the bug here:
 - Consider this input: 7 2 1 8 25 3 Target: 16

HashTable HT;
for each number num1 in input
 HT.Insert(num1)

for each number num1 in input
 num2 = target - num1
 if HT.Contains (num2)
 return (num1, num2)
return empty set



HT
2
25
7
1
3
8



num1	num2	num2 In HT?
7	9	no
2	14	no
1	15	no
8	8	yes

sum to target problem

- We will fix the bug:
 - input: 7 2 1 8 25 3 Target: 16

HashTable HT;

```
for each number num1 in input
    num2 = target - num1
    if HT.Contains( num2 )
        return (num1, num2)
    else
        HT.Insert( num1 )
return empty set
```

Time complexity: ?

Space: ?

num1	num2	num2 In HT?	HT
7	9	no	7
2	14	no	7,2
1	15	no	7,2,1
8	8	no	7,2,1,8
25	-9	no	7,2,1,8,25
3	13	No	7,2,1,8,25,3

sum to target problem (summary)

Option	Time	Space
Brute force	$O(N^2)$	$O(1)$
Sort, then binary search	$O(N \log N) + O(N \log N)$ $= O(N \log N)$	$O(1)$ if in-place sort.
Using 2 pointers	$O(N \log N) + O(N)$ $= O(N \log N)$	$O(1)$ if in-place sort.
Hash table	$O(N)$	$O(N)$

LAB 1 of 2 (sum to target problem)

- Write a function **GetPairSum** (input, target)
- Input:
 - An array of numbers: 7 2 1 8 25 3
 - Target number: 16
- Output:
 - Return any pair of numbers from the input that add up to the target number, if no such pair exists, return an empty pair.
 - Example 1:
 - Input: (2, 7, 16, 25) and target = 27
 - Output: 2, 25
 - Example 2:
 - Input: (2, 7, 16, 25) and target = 34
 - Output: empty set

Intersection of two sorted arrays

- Lets say we need to find the intersection of two sorted arrays.
- Many ways to do this.

- Use case?

- ?

- **Brute force**

for each element a in array 1

search a in array 2

if found

add to result collection

- Time complexity ?
- Space complexity ?

Intersection of two sorted arrays

- But of course, we wont use brute force
- What are some other options?

Intersection of two sorted arrays

- What are some other options?
 - Traverse one and do binary search on the other.

```
for each element a in array1           // or if array2 , then use array1 below for binary search
    if ( BinarySearch ( array2, a ) )
        add to result collection
```

Assume one array is size m , other is size n

- Time complexity: ?
- Space complexity ?

Would you have a preference on which array to traverse and which one to run a binary search on?

Intersection of two sorted arrays

- Another option : Using 2 indices (one in each array). Traverse both arrays, looking for common elements.

array1: 1 4 7 11

array2: 5 7 9 11

index1 = 0

index2 = 0

while index1 < array1_Length and index2 < array2_Length

 if array1 [index1] < array2 [index2]

 ++ index1

 else if array1 [index1] > array2 [index2]

 ++ index2

 else

 add to result collection

 ++ index1 and ++ index2

What's the space complexity of pseudocode above?

What's the time complexity of pseudocode above?

Intersection of two sorted arrays

```
while index1 < array1_Length and index2 < array2_Length  
    if array1[ index1 ] < array2 [ index2 ]  
        ++ index1  
    else if array1[ index1 ] > array2 [ index2 ]  
        ++ index2  
    else  
        add to result collection  
        ++ index1 and ++ index2
```

Worst case: Traverse both arrays fully $\rightarrow O (m + n)$

Best case: Have to traverse only one array $\rightarrow O (m)$ or $O (n)$, depending on which array is traversed.

Lets see what sort of input would lead to the worst and best case.

Intersection of two sorted arrays

- Worst case: $O(m + n)$
 - Traverse both arrays fully... last element in both arrays is the same.

array1:	1	4	7	11
array2:	5	7	9	11

- Best case: $O(m)$ or $O(n)$
 - You end up traversing only one of the arrays
- | | | | | |
|---------|----|----|----|----|
| array1: | 1 | 4 | 7 | 11 |
| array2: | 15 | 17 | 19 | 25 |
- Which array would you traverse fully for the inputs above?

Intersection of two sorted arrays

- Question for u all:
 - Which one do u think will give u better runtime for finding the intersection of two sorted arrays
 - A. Traverse one and run a binary search on the other?
 - $O(m \log n)$
 - B. Using 2 indices : Linear traversal of both arrays, looking for common elements?
 - $O(m + n)$
 - C. Cannot tell without more information.
 - D. Both approaches, binary search and 2 indices, are about the same

Intersection of two sorted arrays

Lets take some input sizes and see how the math works out.

Let $m = 1000,000$ and
 $n = 2000,000$

- Traverse one and run a **binary search** on the other?
 - $O(m \log n)$
 - $1000,000 * \log(2000,000)$
 - $1000,000 * 21$
 - **21,000,0000**
- **Linear traversal of both arrays**, looking for common elements?
 - $O(m + n)$
 - $1000,000 + 2000,000$
 - **3,000,000**
- Which one gives u better performance?

Intersection of two sorted arrays

Let $m = 1,000,000$

$n = 20,000,000,000$

- Traverse one and run a **binary search** on the other?
 - $O(m \log n)$
 - $1,000,000 * \log(20,000,000,000)$
 - Now, $\log(2 * 10,000,000,000) = \log(2) + \log(10,000,000,000)$
 - $= \log(2) + \log(10 * 1,000,000,000)$
 - $= \log(2) + \log(10) + \log(1,000,000,000)$
 - approx $= 1 + 3 + 30 = 34$
 - So, $1,000,000 * \log(20,000,000,000) = 1,000,000 * 34 = 34,000,000$
- **Linear traversal of both arrays**, looking for common elements?
 - $O(m + n)$
 - $1,000,000 + 20,000,000,000$
 - approx $20,000,000,000$ (can ignore 1,000,000)
- Which one gives u better performance in this scenario?

Intersection of two sorted arrays

Could we use a hash table to find the intersection?

If yes, what would be:

Time complexity ?

Space complexity?

Pseudocode on next slide

Intersection of two sorted arrays

Using a hash table

for each element a in array1 $\leftarrow O(m)$

 Insert a in HT

for each element a in array2 $\leftarrow O(n)$

 if HT.Contains(a) $\leftarrow O(1)$

 add to result collection

Time complexity : $O(m + n)$

Space : $O(m)$

LAB 2 of 2 (intersection of 2 sorted arrays)

- Input:
 - Two sorted arrays (or unsorted, and call sort in initialization step).
- Output:
 - Collection of elements that form the intersection of two given arrays
- Write this function using binary search
 - ***FindIntersection_BinarySearch***(array1, array2)