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Exercise 1

```
<program> → begin <statement-list> end  
<statement-list> → <statement>; <statement-list>  
<statement-list> → <statement>  
<statement> → id = <expression>  
<expression> → <factor> + <factor>  
<expression> → <factor> - <factor>  
<expression> → <factor>  
<factor> → id | num
```

① Eliminating Left recursion

The grammar does not contain left recursion as none of the grammar fragments are in the form $A \rightarrow A\alpha \mid \beta$.

② Left factoring

After left factoring, you get the following grammar

```
<program> → begin <statement-list> end  
<statement-list> → <statement> <statement-list'>  
<statement-list'> → ; <statement-list> |  $\epsilon$   
<statement> → id = <expression>  
<expression> → <factor> <expression'>  
<expression'> → + <factor> | - <factor> |  $\epsilon$   
<factor> → id | num
```

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Exercise 3

① $\begin{aligned} \langle \text{program} \rangle &\rightarrow \text{begin } \langle \text{statement_list} \rangle \text{ end} \\ \langle \text{statement_list} \rangle &\rightarrow \langle \text{statement} \rangle \langle \text{statement_list}' \rangle \\ \langle \text{statement_list}' \rangle &\rightarrow ; \langle \text{statement_list} \rangle \mid \epsilon \\ \langle \text{statement} \rangle &\rightarrow \text{id} = \langle \text{expression} \rangle \\ \langle \text{expression} \rangle &\rightarrow \langle \text{factor} \rangle \langle \text{expression}' \rangle \\ \langle \text{expression}' \rangle &\rightarrow + \langle \text{factor} \rangle \mid - \langle \text{factor} \rangle \mid \epsilon \\ \langle \text{factor} \rangle &\rightarrow \text{id} \mid \text{num} \end{aligned}$

- * $\text{FIRST}(\text{program}) = \{ \text{begin} \}$
- * $\text{FIRST}(\text{statement_list}) = \{ \text{FIRST}(\text{statement}) \} = \{ \text{id}, = \}$
- * $\text{FIRST}(\text{statement_list}') = \{ ;, \epsilon \}$ Need to ask whether this should be one terminal or two
- * $\text{FIRST}(\text{statement}) = \{ \text{id}, = \}$ ←
- * $\text{FIRST}(\text{expression}) = \{ \text{FIRST}(\text{factor}) \} = \{ \text{id}, \text{num} \}$
- * $\text{FIRST}(\text{expression}') = \{ +, -, \epsilon \}$
- * $\text{FIRST}(\text{factor}) = \{ \text{id}, \text{num} \}$

- * $\text{FOLLOW}(\text{program}) = \{ \$ \}$
- * $\text{FOLLOW}(\text{statement_list}) = \{ \text{end}, \text{FOLLOW}(\text{statement_list}') \}$
 $= \{ \text{end} \}$ Asked the prof about what to do in this situation. Waiting for a response.
- * $\text{FOLLOW}(\text{statement_list}') = \{ \text{FOLLOW}(\text{statement_list}) \}$
 $= \{ \text{end} \}$
- * $\text{FOLLOW}(\text{statement}) = \{ (\text{FIRST}(\text{statement_list}') - \epsilon) \cup \text{FOLLOW}(\text{statement_list}) \}$
 $= \{ ;, \text{end} \}$ ←

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- * $\text{FOLLOW}(\text{expression}) = \{\text{FOLLOW}(\text{statement})\}$
 $= \{\text{;}, \text{end}\}$ ← remember to redo if wrong
- * $\text{FOLLOW}(\text{expression}') = \{\text{FOLLOW}(\text{expression})\}$
 $= \{\text{;}, \text{end}\}$
- * $\text{FOLLOW}(\text{factor}) = \{(\text{FIRST}(\text{expression}') - \epsilon) \cup$
 $\quad \text{FOLLOW}(\text{expression}), \text{FOLLOW}(\text{expression}')\}$
 $= \{+, -, ;, \text{end}\}$ ← Redo

Exercise 3.2

	begin	end	;	id	num	=	+	-	\$
program		-	-	-	-	-	-	-	-
	<program> → begin <statement_list> end								
statement_list	-	-	-	<statement_list> → <statement> <statement_list'>	-	<statement_list> → <statement> <statement_list'>	-	-	-
statement_list'	-	<statement_list'> → ε	<statement_list'> → ; <statement_list>	-	-	-	-	-	-
statement	-	-	-	<statement> id = <expression>	-	<statement> id = <expression>	-	-	-
expression	-	-	-	<expression> → <factor> <expression'>	<expression> → <factor> <expression'>	-	-	-	-
expression '	-	<expression> → ε	<expression> → ε	-	-	-	<expression> →	<expression> →	-

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							+<factor>	-<factor>	
factor	-	-	-	<factor> → id	<factor> → num	-	-	-	-