

Technical Appendix

This Technical Appendix, a supplement to “Calibrated Nonparametric Scan Statistics for Pattern Detection in Graphs”, consists of:

- Appendix A, Additional Background on Nonparametric Scan Statistics
- Appendix B, Additional Details of CNSS
- Appendix C, Additional Experimental Details
- Appendix D, Case Studies

A Additional Background on Nonparametric Scan Statistics

Here we present additional background on nonparametric scan statistics (Neill and Lingwall 2007; McFowlan III, Speakman, and Neill 2013; Chen and Neill 2014).

As we describe in the main paper, the fundamental problem that NPSSs solve is to find a subset of the data \mathcal{S} , often subject to additional constraints (such as connectedness in the graph setting), and a corresponding significance level α , such that the proportion of significant p-values (at level α) in \mathcal{S} is significantly higher than expected. Or equivalently, if p-values are drawn uniformly on $[0,1]$ under the null hypothesis \mathcal{H}_0 , and under the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$ the p-values in subset \mathcal{S} are drawn with a higher than expected proportion of low (significant) p-values, we wish both to distinguish \mathcal{H}_0 from \mathcal{H}_1 , thus detecting whether a signal is present, and if so, to correctly identify the affected subset \mathcal{S} .

More precisely, NPSSs optimize an objective function $F(\mathcal{S}) = \max_{\alpha \leq \alpha_{\max}} \Phi(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S}))$, where $\Phi(\cdot)$ compares the observed number of significant p-values $N_\alpha(\mathcal{S})$ at level α to the expected number of significant p-values $\mathbb{E}[N_\alpha(\mathcal{S})] = \alpha N(\mathcal{S})$ under the null hypothesis \mathcal{H}_0 . The expectation $\mathbb{E}[N_\alpha(\mathcal{S})] = \alpha N(\mathcal{S})$ follows because, under \mathcal{H}_0 , the current data from which the p-values are generated is exchangeable with the historical data against which the current data values are ranked, leading to p-values that are asymptotically uniform on $[0,1]$ under the null. We discuss the Berk-Jones statistic $\Phi_{BJ}(\cdot)$ in detail in Appendix A.1, and other variants of NPSS in Appendix A.3.

Critically, NPSSs optimize the significance level α between 0 and some constant $\alpha_{\max} < 1$. As noted by McFowlan III, Speakman, and Neill (2013), the purpose of maximizing over a range of α values is to ensure that the statistic can reliably detect either a small number of highly significant p-values or a larger number of moderately significant p-values. If α was fixed at a high value, the statistic would have poor detection performance in the former case; if α was fixed at a low value, it would perform poorly in the latter case. However, maximization over α also presents a serious drawback for the uncalibrated scan: for large real-world graphs, NPSSs select overly large values of α , contributing to their failure to correctly identify the affected subgraph \mathcal{S} .

A number of algorithms have been proposed to optimize $F(\mathcal{S})$ over connected subgraphs, including DFGS (Speakman, McFowlan III, and Neill 2015), AdditiveScan (Speakman, Zhang, and Neill 2013),

NPHGS (Chen and Neill 2014), and ColorCoding (Cadena, Chen, and Vullikanti 2019), but none of these approaches are both exact and scalable to large graphs. Moreover, calibration of NPSSs requires us to identify the subgraph with the largest number of significant p-values N_α for *each subgraph size N and each significance level α* , rather than a single highest-scoring subgraph, thus necessitating our new (approximate) optimization algorithm described in the main paper and in Appendix B.4 below.

Once the highest scoring subgraph, $\mathcal{S}^* = \arg \max_{\mathcal{S} \in \mathbb{M}} F(\mathcal{S})$, has been identified, randomization testing can be used to compute the statistical significance of \mathcal{S}^* . To do so, a large number R of replica graphs are generated under the null hypothesis, i.e., each replica graph has the same structure as the original graph but all p-values p_i are drawn i.i.d. from $\text{Uniform}[0, 1]$. The same search procedure is used to identify the highest scoring subgraph $\mathcal{S}^{(r)}$ for each replica graph $r = 1 \dots R$, and the score $F(\mathcal{S}^*)$ is compared to the distribution of replica scores $F(\mathcal{S}^{(r)})$. To be significant at the standard 0.05 significance level, $F(\mathcal{S}^*)$ must exceed the 95th percentile of the null distribution. This standard approach corrects for *multiple testing*, in that the family-wise error rate (probability of detecting any false positive subgraphs if data is generated under the null) is bounded by the nominal level (e.g., 0.05). However, we note that it does not correct for *miscalibration* across subgraph sizes N and significance levels α , in that large, high-scoring subgraphs \mathcal{S}^* are likely to be detected even when the true subset \mathcal{S} is small or no signal is present.

Finally, as is typical in the scan statistics literature, we can perform *multiple cluster detection* by repeated single cluster detection. That is, after we detect the single highest-scoring subgraph, and test it for statistical significance, we “remove” that subgraph from the data in one of two ways, either assigning the p-value of each detected node as 1, or deleting the detected nodes from the network structure. (The former approach allows secondary clusters to overlap with the primary cluster, while the latter approach does not.) In either case, we apply the same procedure to the updated network to detect the new highest-scoring subgraph, compare the score of this cluster to the significance threshold determined by randomization testing, and repeat until no further significant clusters are present. This statistical testing approach is conservative for the secondary clusters, and several variants of the multiple cluster scan have been proposed to increase detection power for secondary clusters (Zhang, Assuncao, and Kullendorff, 2010; Li et al., 2011).

In the remainder of Appendix A, we present additional details on the fundamental modeling assumptions of the Berk-Jones nonparametric scan statistic (Appendix A.1), computation of empirical p-values (Appendix A.2), other variants of NPSS (Appendix A.3), and differences between NPSS and the Gaussian scan approach of Reyna et al. (2021) and Chitra et al. (2021) (Appendix A.4).

A.1 Fundamental modeling assumptions

In this section, we describe the fundamental modeling assumptions of nonparametric scan statistics, following Mc-

Fowlard III, Speakman, and Neill (2013), and focusing primarily on the Berk-Jones (BJ) likelihood ratio statistic. Unlike parametric scan statistics such as the Poisson and Gaussian statistics (Kulldorff 1997; Neill 2009), NPSSs do not assume that the raw data is drawn from any particular parametric distribution. Instead, the data is converted to empirical p-values by ranking the current data against a reference distribution (e.g., historical values), as described in Appendix A.2. The assumption under the null hypothesis \mathcal{H}_0 is that the current and historical data are exchangeable, and thus, ranking the current data against the historical data (and normalizing) will result in empirical p-values that are uniformly distributed on [0,1]. This also implies that, for any significance level α , the probability that a given p-value is significant ($p_i < \alpha$) is equal to α .

Under the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$, NPSSs must make some assumption about how the distribution of p-values in subset \mathcal{S} differs from the uniform distribution on [0,1]. For the BJ statistic, the assumption under $\mathcal{H}_1(\mathcal{S})$ is that there exist some α and β , where $0 < \alpha < \beta \leq 1$, such that the probability that a given p-value $p_i \in \mathcal{S}$ is significant ($p_i < \alpha$) is equal to β . This is typically framed as an assumption that the distribution of p-values p_i for $v_i \in \mathcal{S}$ is piecewise constant. More precisely, we have $\mathcal{H}_0 : p_i \sim \text{Uniform}[0, 1] \forall v_i \in \mathcal{V}$. Under $\mathcal{H}_1(\mathcal{S})$, we have $p_i \sim \text{Uniform}[0, \alpha]$ with probability β and $p_i \sim \text{Uniform}[\alpha, 1]$ with probability $1 - \beta$, $\forall v_i \in \mathcal{S}$, for $\beta > \alpha$; and $p_i \sim \text{Uniform}[0, 1] \forall v_i \in \mathcal{V} \setminus \mathcal{S}$. Here the values of both α and β are fit by maximum likelihood estimation.

The resulting generalized log-likelihood ratio scan statistic can be written as:

$$\begin{aligned} F(\mathcal{S}) &= \max_{\beta > \alpha} \log \frac{\beta^{N_\alpha(\mathcal{S})} (1 - \beta)^{N(\mathcal{S}) - N_\alpha(\mathcal{S})}}{\alpha^{N_\alpha(\mathcal{S})} (1 - \alpha)^{N(\mathcal{S}) - N_\alpha(\mathcal{S})}} \\ &= \max_{\beta > \alpha} N_\alpha(\mathcal{S}) \log \left(\frac{\beta}{\alpha} \right) + (N(\mathcal{S}) - N_\alpha(\mathcal{S})) \log \left(\frac{1 - \beta}{1 - \alpha} \right). \end{aligned}$$

Then plugging in the maximum likelihood estimate $\beta = N_\alpha(\mathcal{S})/N(\mathcal{S})$ and simplifying, we obtain:

$$F(\mathcal{S}) = \max_{\alpha} N(\mathcal{S}) \text{KL} \left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha \right),$$

where the Kullback-Liebler (KL) divergence is defined as $\text{KL}(a, b) = a \log \frac{a}{b} + (1 - a) \log \frac{1-a}{1-b}$, if $a > b$, and 0 otherwise. (Note that we use a one-sided form of KL divergence throughout, since we care only about subgraphs with a *higher* than expected proportion of significant p-values.) Thus we can see that the score $F(\mathcal{S}) = \max_{\alpha} \Phi_{BJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ is maximized over both α and \mathcal{S} , identifying a subgraph \mathcal{S} and significance level α for which \mathcal{S} has a higher than expected proportion of significant p-values at level α .

The assumption of piecewise constant p-values under the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$ is a relatively lightweight assumption, in that all significant p-values at level α are treated identically: for a given α , the precise value of each p-value does not impact the score $F(\mathcal{S})$, only whether or not that p-value is less than α . The resulting log-likelihood

ratio statistic is equivalent to a log-likelihood ratio defined in terms of the number of significant p-values at level α . That is, the null hypothesis \mathcal{H}_0 can be written as $N_\alpha(\mathcal{S}) \sim \text{Binomial}(N(\mathcal{S}), \alpha)$ for all \mathcal{S} , and the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$ can be written as $N_\alpha(\mathcal{S}) \sim \text{Binomial}(N(\mathcal{S}), \beta)$ for $\beta > \alpha$. Again, we must not only optimize over β , using the maximum likelihood estimate $\beta = \frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}$, but also the significance level α , to identify the highest-scoring subgraph \mathcal{S} .

A.2 Computation of empirical p-values

As noted in Section 3 of the main paper, empirical p-values p_i for the nonparametric scan statistic are computed for each graph node v_i using the two-stage empirical calibration process described by Chen and Neill (2014). We provide more details on this process and explain why it follows that p-values are asymptotically uniform on [0,1] under the null hypothesis \mathcal{H}_0 . Assume that node v_i has a current feature vector $\mathbf{x}_i \in \mathbb{R}^N$ and historical feature vectors $\{\mathbf{x}_i^{(1)}, \dots, \mathbf{x}_i^{(T)}\}$. Moreover, under the null hypothesis \mathcal{H}_0 , we assume that the current data is exchangeable with the historical data, i.e., \mathbf{x}_i and $\mathbf{x}_i^{(1)}, \dots, \mathbf{x}_i^{(T)}$ are all drawn from the same (unknown) distribution.

We first consider the simplest case, in which each node v_i has only a single feature x_i . In this case, the empirical calibration process reduces to ranking the current feature value x_i against its historical values $x_i^{(1)}, \dots, x_i^{(T)}$ and normalizing, i.e.,

$$p_i = \frac{1 + \sum_{t=1 \dots T} \mathbf{1}\{x_i^{(t)} \geq x_i\}}{1 + T}$$

Here we assume one-sided p-values (i.e., higher values of x_i correspond to lower, more significant p-values), but two-sided p-values, or one-sided p-values where lower values of x_i are more significant, can be easily constructed as well. See McFowlard III, Speakman, and Neill (2013) for details.

Under the null hypothesis of exchangeability, it is easy to see that p_i is discrete uniform, taking on values $\frac{1}{1+T}, \frac{2}{1+T}, \dots, 1$ with equal probabilities, and converges in distribution to $\text{Uniform}[0,1]$ as the number of historical observations T becomes large. An alternative is to use p-value ranges, as proposed by McFowlard III, Speakman, and Neill (2013), which guarantee uniform (rather than asymptotically uniform) p-values under the null. We also note that an arbitrary reference set can be used in place of historical data for node v_i , in which case the assumption under \mathcal{H}_0 becomes exchangeability of the current observation with that reference set. See Chen and Neill (2014) for details.

In the more general case where the feature vectors \mathbf{x}_i and $\mathbf{x}_i^{(t)}$ have more than one feature, the two-stage empirical calibration process first ranks each feature value x_{ij} against its historical values $x_{ij}^{(t)}$, and computes a “first-stage p-value” corresponding to each feature value as above:

$$p_{ij} = \frac{1 + \sum_{t=1 \dots T} \mathbf{1}\{x_{ij}^{(t)} \geq x_{ij}\}}{1 + T}.$$

A similar “first-stage p-value” is computed for each historical value:

$$p_{ij}^{(t)} = \frac{1 + \mathbf{1}\{x_{ij} \geq x_{ij}^{(t)}\} + \sum_{t'=1 \dots T, t' \neq t} \mathbf{1}\{x_{ij}^{(t')} \geq x_{ij}^{(t)}\}}{1 + T}.$$

Next, the two-stage empirical calibration process computes the minimum (most significant) p-value for each feature vector:

$$p_{i,\min} = \min_j p_{ij},$$

$$p_{i,\min}^{(t)} = \min_j p_{ij}^{(t)}.$$

And finally, it computes the “second-stage p-value”, using the normalized rank of the minimum p-value $p_{i,\min}$ (here, lower is more significant):

$$p_i = \frac{1 + \sum_{t=1 \dots T} \mathbf{1}\{p_{i,\min}^{(t)} \leq p_{i,\min}\}}{1 + T}.$$

The exchangeability of the first-stage p-values p_{ij} and $p_{ij}^{(t)}$, and the exchangeability of their minima $p_{i,\min}$ and $p_{i,\min}^{(t)}$, follow from the exchangeability of x_{ij} and $x_{ij}^{(t)}$ under the null, and thus the second-stage p-values are asymptotically uniform on [0,1] under \mathcal{H}_0 as above. See Theorem 1 of Chen and Neill (2014) and Section 2.2 of McFowlan III, Speakman, and Neill (2013) for additional details.

The uniformity of p-values is critical since it follows that the expected proportion of significant p-values for a randomly selected connected subset under \mathcal{H}_0 is equal to the significance level α . This is the basis for the NPSS approach of comparing the observed proportion of p-values that are significant at level α to the expected proportion of significant p-values α . As we discuss in detail in the main paper and in Appendix B.1 below, this uncalibrated NPSS approach fails to adjust for the multiplicity of subgraphs, and thus we propose to calibrate α by replacing it with the expected *maximum* proportion of significant p-values $\alpha'(N, \alpha)$.

A.3 Variants of nonparametric scan

While we focused on the Berk-Jones (BJ) nonparametric scan statistic, our calibration approach can easily be applied to other NPSSs such as Higher Criticism (HC) and Kolmogorov-Smirnov (KS), since like BJ these statistics can be written as the maximum (over α values from 0 to α_{\max}) of a scaled divergence $\Phi(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ between the observed and expected proportion of significant p-values at level α . Some examples are provided below:

Berk-Jones:

$$\Phi_{BJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) = N(\mathcal{S}) \text{KL} \left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha \right) \quad (5)$$

Higher Criticism:

$$\Phi_{HC}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) = \frac{N_\alpha(\mathcal{S}) - \alpha N(\mathcal{S})}{\sqrt{N(\mathcal{S})\alpha(1-\alpha)}} \quad (6)$$

Kolmogorov-Smirnov:

$$\Phi_{KS}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) = \sqrt{N(\mathcal{S})} \cdot \left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})} - \alpha \right) \quad (7)$$

Note that in each of these cases, we use a one-sided divergence, since we only wish to detect subgraphs where the observed proportion of significant p-values $N_\alpha(\mathcal{S})/N(\mathcal{S})$ is *greater* than α . We have defined the one-sided KL divergence in Appendix A.1 above, and for the other statistics we simply set them to zero whenever $N_\alpha(\mathcal{S})/N(\mathcal{S}) \leq \alpha$.

Once the value of $\alpha'(N, \alpha) = \mathbb{E}[\max_{\mathcal{S}: |\mathcal{S}|=N} N_\alpha(\mathcal{S})/N]$ has been computed for each N and α , this value can be substituted for α in any of the above equations to obtain a calibrated nonparametric scan statistic:

Calibrated Berk-Jones:

$$\begin{aligned} & \Phi_{CBJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) \\ &= N(\mathcal{S}) \text{KL} \left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha'(N(\mathcal{S}), \alpha) \right) \end{aligned} \quad (8)$$

Calibrated Higher Criticism:

$$\begin{aligned} & \Phi_{CHC}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) \\ &= \frac{N_\alpha(\mathcal{S}) - \alpha'(N(\mathcal{S}), \alpha) N(\mathcal{S})}{\sqrt{N(\mathcal{S})\alpha'(N(\mathcal{S}), \alpha)(1 - \alpha'(N(\mathcal{S}), \alpha))}} \end{aligned} \quad (9)$$

Calibrated Kolmogorov-Smirnov:

$$\begin{aligned} & \Phi_{CKS}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) \\ &= \sqrt{N(\mathcal{S})} \cdot \left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})} - \alpha'(N(\mathcal{S}), \alpha) \right) \end{aligned} \quad (10)$$

Our future work will evaluate the impact of calibration on the detection performance of HC, KS, and other nonparametric scan statistics, as well as exploring whether the calibration approach can be adapted to other scan statistics outside the NPSS family.

We note that the HC nonparametric scan statistic, despite its interpretation as a Gaussian approximation of the BJ likelihood ratio statistic, is distinct from the Gaussian scan statistics described in the following subsection. HC does not assume that individual p-values in \mathcal{S} follow a (transformed) mean-shifted Gaussian distribution under the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$. The HC statistic (like BJ) is based only on the number of significant p-values at level α , which implicitly assumes that the pdf of the p-values is piecewise constant. Rather, it is the number of significant p-values $N_\alpha(\mathcal{S})$ that is assumed to be Gaussian, as a large sample Gaussian approximation to the Binomial distribution for $N_\alpha(\mathcal{S})$ assumed by BJ. Assuming that the number of p-values $N(\mathcal{S})$ is large, then by the Central Limit Theorem, the number of significant p-values at level α , $N_\alpha(\mathcal{S}) \sim \text{Binomial}(N(\mathcal{S}), \alpha)$, converges in distribution to $\text{Gaussian}(\alpha N(\mathcal{S}), \alpha(1-\alpha)N(\mathcal{S}))$, and the HC statistic is the z-score of $N_\alpha(\mathcal{S})$ given this Gaussian distribution, which is similar to a Wald test. Thus HC is a large-sample Gaussian approximation to BJ, regardless of whether the individual p-values are Gaussian.

A.4 Differences between NPSS and Gaussian scan

As we note in the main paper, two recent papers (Reyna et al. 2021; Chitra et al. 2021) investigate miscalibration of scan statistics in the Gaussian setting, demonstrating that the Gaussian scan statistic tends to identify subgraphs that are much larger than the true anomalous subgraph, and presenting an approach (based on Gaussian mixture modeling) that can reduce this bias. In this subsection we explain how our nonparametric scan statistic setting is fundamentally different than the Gaussian setting, resulting in a different source of bias (miscalibration of the parameter α) and thus motivating a different approach to correcting this bias, i.e., recalibration of α using $\alpha'(N, \alpha)$.

First, we note that the typical use of the Gaussian scan, assuming that the raw data follows a Gaussian distribution and computing the likelihood ratio statistic based on this assumption, differs from the nonparametric scan setting where the raw data is converted to p-values that (because of the assumption of exchangeability of current and historical observations) will be uniformly distributed on $[0,1]$ under the null hypothesis. Nonparametric scans do not rely on strong distributional assumptions (like Gaussianity) of the raw data, but rather assume that sufficient reference data (e.g., historical data) are available to convert the raw data to empirical p-values (by ranking it against the reference data and normalizing) as described in Appendix A.2 above.

However, the Gaussian scan approach of Reyna et al. (2021) and Chitra et al. (2021) differs from this typical use in that the raw data are first converted to p-values p_i by ranking and then converted to Gaussian z-scores z_i by the Gaussian probability integral transform, $z_i = \text{CDF}^{-1}(1 - p_i)$, where $\text{CDF}(\cdot)$ assumes the standard normal distribution. The null hypothesis \mathcal{H}_0 is that $z_i \sim \text{Gaussian}(0, 1)$ for all vertices v_i , while the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$ is that $z_i \sim \text{Gaussian}(\mu, 1)$ for $v_i \in \mathcal{S}$ and $z_i \sim \text{Gaussian}(0, 1)$ for $v_i \in \mathcal{V} \setminus \mathcal{S}$. Converting back to p-values, we can write equivalently that $p_i \sim \text{Uniform}[0, 1]$ under \mathcal{H}_0 , as in the nonparametric scan, and $p_i \sim 1 - \text{CDF}(z_i)$ where $z_i \sim \text{Gaussian}(\mu, 1)$ for $v_i \in \mathcal{S}$ under the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$, where the parameter μ is fit by maximum likelihood estimation.

While this framing of the Gaussian scan leads to identical null hypotheses, with p-values uniformly distributed on $[0,1]$, the NPSS alternative hypothesis $\mathcal{H}_1(\mathcal{S})$ is fundamentally different from the Gaussian setting in two ways. First, as derived in Appendix A.1 above, the nonparametric scan fits two parameters α (significance level of the subgraph) and β (fraction of significant nodes in the subgraph) by maximum likelihood estimation, while the Gaussian scan fits only the parameter μ . The additional parameter α is critical to the NPSS setting for two reasons: (1) maximizing over a range of significance levels α gives the nonparametric scan high power to detect compact signals (a small number of highly significant p-values), dispersed signals (a large number of slightly significant p-values), or anything in between; and (2) as we show, miscalibration of the estimated proportion N_α/N across different α values leads to an incorrect (overly large) choice of α , obscuring the true signal, but using the corrected $\alpha'(N, \alpha) = \mathbb{E}[\max_{\mathcal{S} \in \mathcal{M}: |\mathcal{S}|=N} N_\alpha(\mathcal{S})/N]$

in place of the uncorrected $\alpha = \mathbb{E}[N_\alpha(\mathcal{S})/N]$ solves this problem. The previous approaches for calibrating the Gaussian scan cannot solve this issue of miscalibration over α , nor is a Gaussian mixture modeling approach appropriate when p-values do not follow a transformed Gaussian under the null. On the other hand, as our experimental results show, our new approach to calibrating NPSSs is effective regardless of whether the signal is a (transformed) Gaussian or piecewise constant p-values.

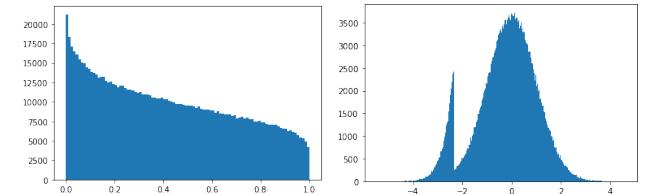


Figure 4: Examples of the difference in signal shape between piecewise constant and transformed Gaussian p-values. Left panel: histogram of p-values corresponding to a transformed Gaussian signal ($\mu = 0.3$). Right panel: histogram of z-scores corresponding to a piecewise-constant p-value signal.

A second fundamental difference is that, even for a given value of α , the nonparametric scan assumes a qualitatively different signal shape, i.e., piecewise constant rather than transformed Gaussian p-values, under the alternative hypothesis $\mathcal{H}_1(\mathcal{S})$. As shown in the left panel of Figure 4, if we plot the histogram of p-values corresponding to a transformed mean-shifted Gaussian (with $\mu = 0.3$ in our example), we see that they are decreasing on the entire interval $[0,1]$, as opposed to the NPSS assumption of piecewise constant p-values. Conversely, if we plot the histogram of z-scores corresponding to the alternative hypothesis for the BJ statistic (with 10% of p-values significant at $\alpha = .01$ in our example), as in the right panel of Figure 4, we see that the distribution is neither Gaussian nor a mixture of Gaussians, as the mixture component with low p-values is heavily left-skewed, with a sharp cutoff of the right tail at the z-score corresponding to α .

While a performance comparison of nonparametric and (transformed) Gaussian scans is beyond the scope of this paper, we note that the differing alternative hypotheses have important implications as to what types of signals can be detected. One might expect the transformed Gaussian scan to have somewhat higher power if its modeling assumptions are correct and the signals are in fact Gaussian. However, certain signals (such as cases where the p-value distribution is symmetric around $p = 0.5$) would not be detectable in the Gaussian setting, while the nonparametric scans can detect signals as long as there exist some (α, β) , where $\beta > \alpha$, such that $\Pr(p < \alpha) = \beta$.

In summary, the nonparametric scan statistics that we consider here differ fundamentally from the (transformed) Gaussian scans considered by Reyna et al. (2021) and Chitra et al. (2021), both in their assumptions about the true signal (distribution of p-values under \mathcal{H}_1) and in their maximization over a critical parameter, the significance level α .

These differences motivate both our new empirical studies to understand and quantify the miscalibration of α for the (uncalibrated) nonparametric scan statistics, as well as our new approach to correctly calibrate α .

B Additional Details of CNSS

B.1 Correctness of the calibration approach

As we discuss in detail in the main paper, the primary source of miscalibration for NPSSs is the discrepancy between α and $\alpha'(N, \alpha)$. That is, for a given significance level α , the expected proportion of significant nodes (at level α) for a randomly selected subgraph under \mathcal{H}_0 is equal to α . However, for a non-random subgraph that is selected by maximizing the number of significant nodes, $\arg \max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} N_\alpha(\mathcal{S})$, the expected proportion of significant nodes $\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}$ under \mathcal{H}_0 , called $\alpha'(N, \alpha)$, is much greater than α , as shown in Figure 1. Previous, uncalibrated NPSS approaches do not account for this discrepancy between α' and α , causing them to incorrectly detect large, high-scoring subgraphs even when no signal is present, or equivalently, these incorrectly detected subgraphs will obscure a true signal with lower score. This motivates our development of *calibrated* NPSS score functions that compare the observed proportion of significant nodes $\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}$ to $\alpha'(N, \alpha)$ rather than α . Moreover, we observe that the amount of discrepancy between α' and α varies not only with N and α , but also with the graph structure (including its size and sparsity), thus motivating our decision to empirically calibrate $\alpha'(N, \alpha)$ based on randomization testing.

In the discussion below, we consider the correctness of using $\alpha'(N, \alpha) = \mathbb{E} \left[\max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \frac{N_\alpha(\mathcal{S})}{N} \right]$ in place of the expectation $\alpha = \mathbb{E} \left[\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})} \right]$ in NPSS. For example, for the Berk-Jones score function $\Phi_{BJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) = N(\mathcal{S}) \text{KL} \left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha \right)$, we instead define the calibrated Berk-Jones score function, $\Phi_{CBJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) = N(\mathcal{S}) \text{KL} \left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha'(N(\mathcal{S}), \alpha) \right)$.

We now present a more formal argument for the correctness of calibration, followed by a numeric example showing how the uncalibrated Berk-Jones statistic fails and why our calibration approach corrects this issue.

First, we note that the objective function $\max_{\mathcal{S} \in \mathbb{M}} F(\mathcal{S})$ can be written as $\max_{N \in \{1, \dots, |\mathcal{V}|\}, \alpha \leq \alpha_{\max}} g(N, \alpha)$, where

$$g(N, \alpha) = \max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \Phi(\alpha, N_\alpha(\mathcal{S}), N)$$

is the maximum score over all subgraphs of size N at significance level α . Since $\Phi(\alpha, N_\alpha, N)$ is an increasing function of N_α , we can rewrite $g(N, \alpha) = \Phi(\alpha, \max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} N_\alpha(\mathcal{S}), N)$. Now, we would like $g(N, \alpha)$ to satisfy two intuitive properties if the null hypothesis \mathcal{H}_0 is true: (i) $g(N, \alpha)$ should be small for all N and α , and (ii) $g(N, \alpha)$ should be similar in magnitude across all values of N and α . The first property makes it possible to differentiate \mathcal{H}_0 from $\mathcal{H}_1(\mathcal{S})$, since large values of $g(N, \alpha)$

would obscure the true signal \mathcal{S} . The second property provides similar detection power across different subgraph sizes N and significance levels α .

However, the uncalibrated Berk-Jones statistic does not satisfy either of these properties. For brevity, let $h(N, \alpha)$ denote the maximum proportion of significant p-values in a subgraph of size N ,

$$h(N, \alpha) = \max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \frac{N_\alpha(\mathcal{S})}{N}.$$

Then under \mathcal{H}_0 , we have $g(N, \alpha) = N \text{KL}(h(N, \alpha), \alpha)$. Replacing $h(N, \alpha)$ with its expectation under the null, $\alpha'(N, \alpha)$, we obtain $g(N, \alpha) \approx N \text{KL}(\alpha'(N, \alpha), \alpha)$. As shown in Figure 1 and the numeric example below, the discrepancy between α' and α , and the resulting values of $g(N, \alpha)$, are large, obscuring the true signal when one is present. Moreover, $g(N, \alpha)$ is much larger for high values of the significance level α and subgraph size N , leading to the detection of overly large, incorrect subgraphs.

On the other hand, for the calibrated Berk-Jones statistic, we have:

$$\begin{aligned} g(N, \alpha) &= N \text{KL} \left(\max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \frac{N_\alpha(\mathcal{S})}{N}, \alpha'(N, \alpha) \right) \\ &= N \text{KL}(h(N, \alpha), \mathbb{E}[h(N, \alpha)]), \end{aligned}$$

where the expectation assumes that \mathcal{H}_0 is true, and is computed by averaging $h(N, \alpha) = \max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \frac{N_\alpha(\mathcal{S})}{N}$ over a large number of instantiations of the graph \mathbb{G} with p-values drawn from the null distribution, $p_i \sim \text{Uniform}[0, 1] \forall v_i \in \mathcal{V}$. Thus if the null hypothesis \mathcal{H}_0 is true, the value of $h(N, \alpha)$ for the real data is drawn from the same distribution as the null data, and then compared (using one-sided KL divergence) to the expectation of that distribution $\alpha'(N, \alpha)$ in order to compute the score $g(N, \alpha)$. From this, it is clear that the scores $g(N, \alpha)$ will be close to zero under \mathcal{H}_0 , diverging from zero only when the value of $h(N, \alpha)$ happens by chance to be greater than its expectation.

To more precisely quantify the impact of the variance of $h(N, \alpha)$ on the score $g(N, \alpha)$, assuming that the null hypothesis \mathcal{H}_0 is true, we use a second-order Taylor expansion of the KL divergence to obtain:

$$\mathbb{E}[g(N, \alpha)] \approx \frac{N \text{Var}[h(N, \alpha)]}{2\alpha'(N, \alpha)(1 - \alpha'(N, \alpha))},$$

where the variance is taken over instantiations of the graph \mathbb{G} with p-values drawn under the null distribution. We observe empirically that, for large N , the variance of $h(N, \alpha)$ under the null is approximated well by $\frac{\alpha'(N, \alpha)}{N}$, giving us $\mathbb{E}[g(N, \alpha)] \approx \frac{1}{2}(1 - \alpha'(N, \alpha))^{-1}$, which is small and slowly decreases with N . For small N , we observe empirically that $\text{Var}[h(N, \alpha)]$ is much smaller than $\frac{\alpha'(N, \alpha)}{N}$. For example, $\text{Var}[h(N, \alpha)] = 0$ when $\alpha'(N, \alpha) = 1$. As a result, we observe scores $g(N, \alpha)$ that are close to zero (typically peaking in the low single digits) for all N and α , as illustrated in Figure 5c. This observation has two important implications: first, since the null scores are much lower than for the uncalibrated BJ statistic, a true signal $\mathcal{H}_1(\mathcal{S})$ can be more easily

detected and the true subgraph \mathcal{S} more accurately identified. Second, we no longer observe the biases toward large N and α which led the uncalibrated BJ statistic to detect large, incorrect subgraphs.

Thus the calibrated BJ statistic corrects for the multiplicity of subgraphs of a given size N , by comparing the observed maximum value of $N_\alpha(\mathcal{S})$ across all size- N subgraphs to the expectation of that maximum value under the null. In doing so, it calibrates the statistic across all values of N and α , giving similar values of $g(N, \alpha)$ under \mathcal{H}_0 . However, we note that calibration alone does not correct for the multiple testing resulting from maximization of $g(N, \alpha)$ over all $N \in \{1 \dots |\mathcal{V}|\}$ and $\alpha \leq \alpha_{\max}$. We must still apply the standard randomization testing approach described in Appendix A to compare the maximum calibrated score $\max_{\mathcal{S} \in \mathbb{M}} F(\mathcal{S})$ to the distribution of the maximum calibrated score under \mathcal{H}_0 , thus bounding the overall false positive rate.

Finally, we note that correcting the miscalibration of NPSSs does not require an exact solution to the optimization problem, i.e., maximizing $N_\alpha(\mathcal{S})$ over subgraphs of size N . Under \mathcal{H}_0 , the current and null data are exchangeable, so for a given α , $\tilde{h}(N, \alpha) \approx \max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \frac{N_\alpha(\mathcal{S})}{N}$ will be distributed identically for the current and null data, as long as the *same* approximation algorithm is used for the current and null data. That is, we compare the observed value of the (approximate) maximum proportion of significant p-values, $\tilde{h}(N, \alpha)$, to the expectation of $\tilde{h}(N, \alpha)$ under \mathcal{H}_0 , $\tilde{\alpha}'(N, \alpha)$, so $g(N, \alpha) \approx N \text{KL}(\tilde{h}(N, \alpha), \tilde{\alpha}'(N, \alpha))$ remains well-calibrated. The downside of a using an approximate rather than exact search is some potential loss of detection power and accuracy under $\mathcal{H}_1(\mathcal{S})$, but our experiments demonstrate that the approximate algorithm achieves high detection performance across five large real-world datasets, outperforming the uncalibrated scan and baseline methods by a wide margin.

Motivating Example. As a concrete example of how the uncalibrated BJ statistic fails, and how calibration solves this problem, let us consider a single instantiation of the `WikiVote` graph ($|\mathcal{V}| = 7,066$) generated under $\mathcal{H}_1(\mathcal{S})$. The true subgraph \mathcal{S} was generated using a random walk on the graph structure, with $|\mathcal{S}| = 100$ and a relatively strong signal injected, such that 75% of the p-values in \mathcal{S} are significant at $\alpha = .01$. This corresponds to $q = 75$ for the piecewise constant p-value simulations in Appendix C.5 below. In this case, the uncalibrated BJ score of the true subgraph at the significance threshold $\alpha = 0.01$ can be computed as $N(\mathcal{S}) \text{KL}\left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha\right) = 100 \text{KL}(0.75, 0.01) \approx 289$. Thus the true subset \mathcal{S} (into which the signal has been injected) has a high BJ score, corresponding to the true significance threshold $\alpha = 0.01$.

However, another, much larger subset has an even higher score, corresponding to a high significance threshold α . More precisely, at the highest α value considered ($\alpha = .09$), the uncalibrated scan picks out a subgraph containing nearly all of the approximately 700 significant p-values in the graph, plus some additional nodes needed

to connect them, resulting in a 900 node subgraph with 74.4% significant p-values. If this observed proportion of 74.4% is compared to the 9% of significant p-values that one would expect to see in a random subgraph at $\alpha = .09$, it would have an extremely high BJ score of $N(\mathcal{S}) \text{KL}\left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha\right) = 900 \text{KL}(0.744, 0.09) \approx 1100$, and as a result, the uncalibrated scan incorrectly identifies this subgraph instead of the true subgraph. However, under \mathcal{H}_0 , for a graph of this size and structure, one would *expect* to see some subgraph of this size $N = 900$ with about $\alpha'(900, 0.09) = 69.9\%$ significant p-values. Then the calibrated score of that incorrect subgraph (comparing the observed 74.4% to the expected 69.9%) would be close to zero: $N(\mathcal{S}) \text{KL}\left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha'(N(\mathcal{S}), \alpha)\right) = 900 \text{KL}(0.744, 0.699) = 4.47$, allowing a subgraph closer to the true subgraph (with $N = 202$ and 73.3% of p-values significant at $\alpha = .01$, as compared to $\alpha'(202, 0.01) = 0.347$, for a score of 62.26) to be found instead. Comparing this subgraph to the true subgraph using the metrics in Appendix C.3 below, we compute precision = 0.72, recall = 0.69, and F-score = 0.70, while the subgraph found by the uncalibrated scan had slightly higher recall (0.75) but much lower precision (0.08) and F-score (0.15). This is why calibration using $\alpha'(N, \alpha)$ in place of α improves detection performance: it prevents the scan from detecting large, incorrect subgraphs whose log-likelihood ratio score exceeds the score of the true subgraph.

For additional clarity, we show examples of score computation for the calibrated and uncalibrated BJ statistics in Figure 5 and 6. Figure 5 is for an instantiation of the `WikiVote` graph with no injected signal, such that all p-values are uniform on [0,1], while Figure 6 is for the instantiation of the `WikiVote` graph with injected piecewise-constant signal discussed above. For illustration, we consider only $\alpha = .01$ (red lines) and $\alpha = .09$ (blue lines). In each graph, the left panel shows the observed value of $h(N, \alpha) = \max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \frac{N_\alpha(\mathcal{S})}{N}$ as a function of N , shown as a dashed line, as compared to the expected value $\alpha'(N, \alpha) = \mathbb{E}\left[\max_{\mathcal{S} \in \mathbb{M}: |\mathcal{S}|=N} \frac{N_\alpha(\mathcal{S})}{N}\right]$ under the null hypothesis \mathcal{H}_0 , shown as a solid line. As expected, when there is no signal (Figure 5), the observed maximum matches α' almost exactly (very slight differences are visible when zooming in on the graph), demonstrating that α' is correctly calibrated, and the resulting calibrated BJ score (right panel) is close to zero. On the other hand, we note that the observed maximum is much larger than α , and thus the resulting uncalibrated BJ score (center panel) is very high. When the signal is present (Figure 6), we see clear differences between the observed maximum and $\alpha'(N, \alpha)$ (left panel), resulting in a high calibrated score (right panel) that is maximized at the true α value, $\alpha = .01$, for a subgraph that closely matches the true subgraph as described above. On the other hand, for the uncalibrated scan, the score of the true subgraph at the true α value is exceeded by the score of the much larger subgraph described above (center panel), at the incorrect $\alpha = .09$.

In summary, this example, along with the discussion

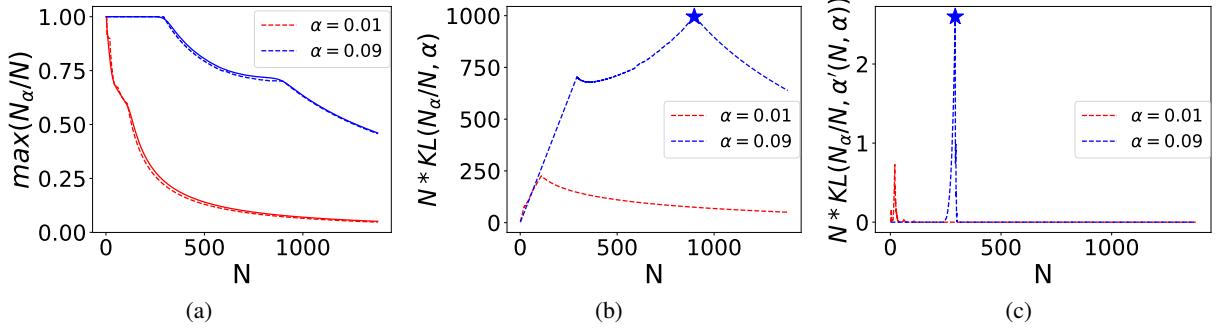


Figure 5: Example of calibrated versus uncalibrated score computation for the `WikiVote` graph with p-values generated under \mathcal{H}_0 . Red lines: $\alpha = .01$. Blue lines: $\alpha = .09$. (a) Observed maximum value of N_α/N (dashed line) and expected maximum value $\alpha'(N, \alpha)$ (solid line). (b) BJ score of the uncalibrated scan statistic, comparing the observed N_α/N to the expected value α . (c) BJ score of the calibrated scan statistic, comparing the observed N_α/N to the expected maximum value $\alpha'(N, \alpha)$.

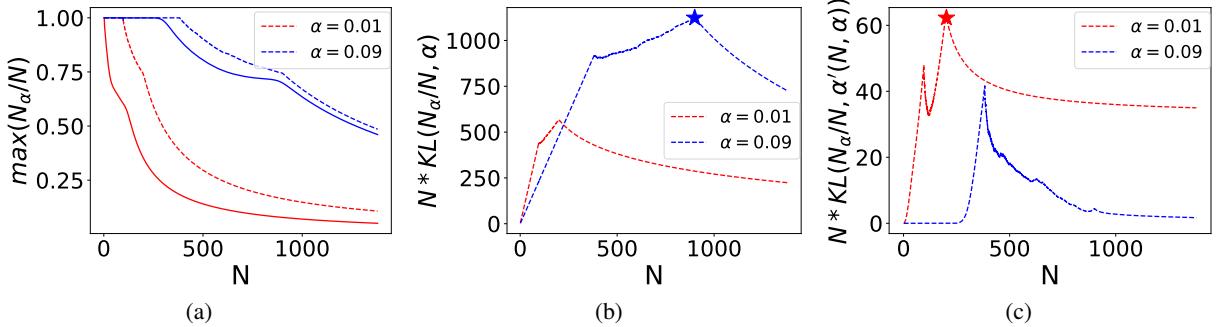


Figure 6: Example of calibrated versus uncalibrated score computation for the `WikiVote` graph with piecewise constant p-values generated under $\mathcal{H}_1(\mathcal{S})$. Red lines: $\alpha = .01$. Blue lines: $\alpha = .09$. (a) Observed maximum value of N_α/N (dashed line) and expected maximum value $\alpha'(N, \alpha)$ (solid line). (b) BJ score of the uncalibrated scan statistic, comparing the observed N_α/N to the expected value α . (c) BJ score of the calibrated scan statistic, comparing the observed N_α/N to the expected maximum value $\alpha'(N, \alpha)$.

above, demonstrates how miscalibration can harm the detection performance of uncalibrated NPSSs, resulting in a very large, incorrect detected subgraph, at an incorrect α value, that swamps the true signal. However, when the uncalibrated α value is replaced with the calibrated $\alpha'(N, \alpha)$, the miscalibration issue is solved, leading to substantially improved detection performance for the calibrated scan.

B.2 Proofs of Theorems 1-2

Theorem 1. For each $c \in \{1, \dots, |\mathcal{V}|\}$, let k_c be the largest ext-degree of a connected subgraph of size c . Then for any $N \in \{1, \dots, |\mathcal{V}|\}$ such that $c \leq N \leq c + k_c$, a lower bound for $\mathbb{E}[\max_{\mathcal{S} \in \mathbb{M}, |\mathcal{S}|=N} N_\alpha(\mathcal{S})]$ is: $c\alpha + \min(k_c\alpha, N - c)$.

Proof. We consider the size- N subgraph consisting of all c nodes from the subgraph and $N - c$ nodes from the neighbors. There are two cases. For $c \leq N \leq c + k_c\alpha$, we choose only significant nodes from the neighbors, so the expected number of significant nodes for a given N is $c\alpha + (N - c)$. For $c + k_c\alpha \leq N \leq c + k_c$, all significant neighbors are included, and thus the expected number of significant nodes for the given N is $c\alpha + k_c\alpha$. Thus we have $\mathbb{E}[N_\alpha] = c\alpha + \min(k_c\alpha, N - c)$. \square

Theorem 2. For an Erdos-Renyi $(|\mathcal{V}|, p)$ graph with average degree $\langle k \rangle = (|\mathcal{V}| - 1)p$, with high probability,

$$\alpha' \geq \min \left(1, \frac{\alpha|\mathcal{V}|}{N} \left(1 - \exp \left(-\langle k \rangle \frac{N}{|\mathcal{V}|} \right) \right) \right).$$

Proof. To find a lower bound on α' given N and α , let $Z = \frac{|\mathcal{V}|}{N}(1 - \exp(-\langle k \rangle \frac{N}{|\mathcal{V}|}))$. We note that $Z \leq \langle k \rangle$, since $1 - \exp(-x) \leq x$ for all non-negative x . Now there are two cases.

Case 1: If $\alpha \geq \frac{1}{Z}$, then the fraction of significant nodes α is also greater than $\frac{1}{\langle k \rangle}$. Thus there exists w.h.p. a giant cluster of size $|\mathcal{V}|P_\infty$, where $P_\infty = \alpha(1 - \exp(-\langle k \rangle P_\infty)) \geq \frac{1}{Z}(1 - \exp(-\langle k \rangle P_\infty))$, consisting entirely of significant nodes. Now we can see that $P_\infty \geq \frac{N}{|\mathcal{V}|}$, since $\frac{N}{|\mathcal{V}|} = \frac{1}{Z}(1 - \exp(-\langle k \rangle \frac{N}{|\mathcal{V}|}))$, and thus there exists a cluster of size N consisting entirely of significant nodes, i.e., $\alpha' = 1$ for the given N and α .

Case 2: If $\alpha < \frac{1}{Z}$, then mark all of the significant nodes and fraction $\frac{1/Z - \alpha}{1 - \alpha}$ of non-significant nodes, so that the proportion of marked nodes is $1/Z$, and the probability that

a marked node is significant is $Z\alpha$. Since the fraction of marked nodes $\frac{1}{Z} > \frac{1}{\langle k \rangle}$, there exists w.h.p. a giant cluster of size $|\mathcal{V}|P_\infty$, where $P_\infty = \frac{1}{Z}(1 - \exp(-\langle k \rangle P_\infty)) = \frac{N}{|\mathcal{V}|}$, consisting entirely of marked nodes. Thus there exists a cluster of size N such that the fraction of significant nodes in that cluster is $Z\alpha = \frac{\alpha|\mathcal{V}|}{N}(1 - \exp(-\langle k \rangle \frac{N}{|\mathcal{V}|}))$, i.e., $\alpha' \geq \frac{\alpha|\mathcal{V}|}{N}(1 - \exp(-\langle k \rangle \frac{N}{|\mathcal{V}|}))$ for the given N and α . Combining these two cases, we obtain the lower bound on α' , for all α and N . \square

B.3 Algorithm 1

As described in Section 4, Algorithm 1 searches for the highest-scoring connected subgraph, $\arg \max_{\mathcal{S}} F(\mathcal{S}) = \arg \max_{\mathcal{S}} \Phi_{CBJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S}))$. To do so, it steps over a range of significance thresholds $\alpha \in \mathcal{L}$, calling Algorithm 2 for each α value, and collecting the subgraph \mathcal{S} with largest $N_\alpha(\mathcal{S})$ for each $N \in \{1, \dots, |\mathcal{V}|\}$ and each $\alpha \in \mathcal{L}$. These subgraphs are then scored with the calibrated Berk-Jones statistic Φ_{CBJ} , and the subgraph with the highest score is returned.

Algorithm 1: Anomalous Subgraph Detection via Calibrated Non-parametric Scan Statistics

Input: Graph $\mathbb{G} = (\mathcal{V}, \mathcal{E})$, \mathbf{X} which contains all historical feature observations of each node, and a list of α denoted as \mathcal{L} , i.e., $\mathcal{L} = [0.001, \dots, 0.009, 0.01, \dots, 0.09]$;

Output: A set of nodes $\mathcal{S} \subseteq \mathcal{V}$ that form a connected subgraph of \mathbb{G} ;

- 1 Compute empirical p-values \mathbf{p} for all nodes based on \mathbf{X} , following the approach described in Chen and Neill (2014).
- 2 **for** $\alpha \in \mathcal{L}$ **do**
- 3 Apply Algorithm 2 on graph \mathbb{G} with empirical p-values \mathbf{p} and significance threshold α , to collect a list of candidate subgraphs \mathcal{S} and corresponding $N_\alpha(\mathcal{S})$ and $N(\mathcal{S})$.
- 4 **for** each collected triplet $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ **do**
- 5 Compute the calibrated BJ scan statistic:
- 6
$$\Phi_{CBJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S})) = N(\mathcal{S}) \times \text{KL}\left(\frac{N_\alpha(\mathcal{S})}{N(\mathcal{S})}, \alpha'(N(\mathcal{S}), \alpha)\right) \quad (11)$$

where

$$\alpha'(N, \alpha) = \frac{\mathbb{E} [\max_{\mathcal{S}:|\mathcal{S}|=N} N_\alpha(\mathcal{S})]}{N} \quad (12)$$
- 7 **end**
- 8 **end**
- 9 **return** \mathcal{S} with the highest $\Phi_{CBJ}(\alpha, N_\alpha(\mathcal{S}), N(\mathcal{S}))$

At line 1, we use the two-stage empirical calibration procedure described in (Chen and Neill 2014) to convert the observed node features into a single p-value for each node

based on the historical node features. Note that the computation of these p-values is not the focus of the present work; thus, in our simulated experiments on the five datasets, we simulate the p-values directly (as discussed in Section 5) rather than simulating the node features and computing p-values from them. This is not only faster, but provides a more natural way to measure the strength of the injected signal.

At line 2, we iterate over a list of significance thresholds α . In our experiments, we used significance thresholds $\alpha \in \{0.001, \dots, 0.009, 0.01, \dots, 0.09\}$.

At line 5, the algorithm assumes that we have pre-computed the $\alpha'(N, \alpha)$ values for each N and α under consideration. To do so, there are two options. First, we could use the randomization tests that apply Algorithm 2 on K replicas of datasets under \mathcal{H}_0 to collect K number of N_α values for each $N \in \{1, \dots, |\mathcal{V}|\}$ and $\alpha \in \mathcal{L}$. Then we use the averaged N_α to compute the $\alpha'(N, \alpha)$. We could also replace the randomization tests with the lower bounds of $\alpha'(N, \alpha)$ as discussed in Section 4.2, which we describe as the method CNSS+LowerBound.

In addition, we can also apply the core-tree decomposition and tree compression steps described in Section 4.3 and Appendix B.5 to obtain a compressed core \mathbb{C} and corresponding p-values \mathbf{p} , and then apply the algorithm on \mathbb{C} and \mathbf{p} to speed up the search. The core-tree decomposition can be done just once for a given graph (between lines 1 and 2), while the tree compression is done separately for each α value (between lines 2 and 3). This method is called CNSS+CoreTree.

B.4 Algorithm 2

As described in Section 4.1, for a given graph \mathbb{G} with corresponding empirical p-values \mathbf{p} for each node, and a given significance threshold α , Algorithm 2 searches for the most significant subgraph $\max_{\mathcal{S} \in \mathbb{M}, |\mathcal{S}|=N} N_\alpha(\mathcal{S})$ for each $N \in \{1, \dots, |\mathcal{V}|\}$. The algorithm is a greedy merging approach, which enables it to scale to large graph sizes but has the drawback of not guaranteeing that the subgraph with maximum N_α will be found.

At line 1: After we merge adjacent significant nodes, the merged nodes have significance ratio equal to 1. Those merged nodes \mathcal{S} could be viewed as candidate detected subgraphs, and will be merged further to create larger candidate subgraphs.

At line 2: We maintain the ordered list \mathcal{Z} throughout the algorithm, where \mathcal{Z} is sorted by significance ratio first, highest to lowest. If two items in \mathcal{Z} have the same significance ratio, we sort them based on the merged node size $N(\mathcal{S})$, highest to lowest.

At line 4: Initially, all nodes in the list have the same significance ratio of 1, so the merged node representing the largest subgraph of significant nodes will be the initial root \mathcal{S}^T .

At line 8: We evaluate three merge options on the root \mathcal{S}^T as follows:

1. If there exists a neighbor n of \mathcal{S}^T which contains some or all significant p-values, merge n into \mathcal{S}^T .
2. If there exists a non-significant neighbor n of \mathcal{S}^T which is also adjacent to at least one other significant node,

Algorithm 2: Greedily Merge Graph and Estimate
 $\max N_\alpha$ for $N \in \{1, \dots, |\mathcal{V}|\}$.

Input: graph $\mathbb{G} = (\mathcal{V}, \mathcal{E})$, $\mathbf{p} \in [0, 1]^{|\mathcal{V}|}$, and
 $\alpha \in [0, 1]$.
Output: a set of $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ triplets

- 1 Merge all adjacent significant nodes, and get a list of merged nodes \mathcal{S} , denoted as \mathcal{Z} .
- 2 Sort the list \mathcal{Z} by significance ratio $N_\alpha(\mathcal{S})/N(\mathcal{S})$, from highest to lowest, and breaking ties using larger size $N(\mathcal{S})$.
- 3 Initialize an empty list \mathcal{P} to store $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ triplets.
- 4 Get the merged node \mathcal{S}^T with the highest significance ratio in \mathcal{Z} .
- 5 Add $(\mathcal{S}^T, N_\alpha(\mathcal{S}^T), N(\mathcal{S}^T))$ of the merged node \mathcal{S}^T to \mathcal{P} .
- 6 **while** $\text{length}(\mathcal{Z}) > 1$ **do**
- 7 Select the merged node \mathcal{S}^T with the highest significance ratio in \mathcal{Z} as root.
- 8 Select and apply the best merge option among the three options described below, to merge another node into \mathcal{S}^T .
- 9 **if** the best merge option is option 1 **then**
- 10 Add $(\mathcal{S}^T, N_\alpha(\mathcal{S}^T), N(\mathcal{S}^T))$ of the merged node \mathcal{S}^T to \mathcal{P} .
- 11 **end**
- 12 **end**
- 13 Sort the list \mathcal{P} by $N(\mathcal{S})$ from highest to lowest.
- 14 $\text{previous_ratio} \leftarrow N_\alpha(\mathcal{S})/N(\mathcal{S})$ of the first element of \mathcal{P} .
- 15 **for** each successive element $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ in list \mathcal{P} **do**
- 16 $\text{current_ratio} \leftarrow N_\alpha(\mathcal{S})/N(\mathcal{S})$.
- 17 **if** $\text{current_ratio} > \text{previous_ratio}$ **then**
- 18 $\text{previous_ratio} \leftarrow \text{current_ratio}$
- 19 **else**
- 20 Delete element $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ from \mathcal{P} .
- 21 **end**
- 22 **end**
- 23 **return** \mathcal{P}

merge n into \mathcal{S}^T .

3. Merge the highest-degree non-significant neighbor n into \mathcal{S}^T .

Then we select and apply the option that leads to the highest significance ratio for the merged node among these three options. If they result in same significance ratio, we use the priority order $1 > 2 > 3$.

At line 9: We also collect the $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ triplets of the merged nodes after each merge of option 1, which produces a list of $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ triplets where $N(\mathcal{S}) \in \{1, \dots, |\mathcal{V}|\}$. Only the largest N_α value for each N must be kept. Note that we do not need to record the triplets formed after option 2 or option 3 since no significant p-values have been added to \mathcal{S}^T .

The purpose of lines 13 - 22 is to remove sub-optimal $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ from \mathcal{P} , as a subgraph with smaller significance ratio $N_\alpha(\mathcal{S})/N(\mathcal{S})$ and smaller size $N(\mathcal{S})$ is guaranteed to have lower score.

We note that, as written, the algorithm only returns $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ for a subset of N values, $\mathcal{Q} \subseteq \{1, \dots, |\mathcal{V}|\}$. For the original graph, these are the only values of $N(\mathcal{S})$ for which the corresponding subgraph \mathcal{S} could have optimal score $F(\mathcal{S})$. For the replica graphs used in randomization testing, we apply linear interpolation to estimate the N_α for the remaining values of N , that is, $N \in \{1, \dots, |\mathcal{V}|\} \setminus \mathcal{Q}$. Also, when we apply Algorithm 2 under the null hypothesis for randomization tests, we only record the $(N_\alpha(\mathcal{S}), N(\mathcal{S}))$ pairs, without recording \mathcal{S} , to save memory space.

B.5 Core-Tree Decomposition

We adopt the implementation of core-tree decomposition from (Maehara et al. 2014). After we decompose the whole graph into the core part $\mathbb{C} = (\mathcal{V}_C, \mathcal{E}_C)$ and tree part $\mathbb{T} = (\mathcal{V}_T, \mathcal{E}_T)$, we utilize an additional tree-compression step before applying Algorithm 2 on \mathbb{C} .

Tree compression merges the significant nodes in each single tree into an adjacent core node. We could conceivably optimize over each single tree to identify and merge the highest scoring sub-tree, but this would be time-consuming given the large number of trees resulting from the core-tree decomposition. Instead, we use breadth-first tree search to find and merge each significant sub-tree that is adjacent to the core. If a significant tree node is adjacent to multiple core nodes, then we merge this tree node into the most significant adjacent core node. In order to achieve this, we first sort the core nodes \mathcal{V}_C by p-value (lowest to highest), and also remove all non-significant tree nodes from the graph (this may disconnect some of the significant tree nodes, which are removed from the graph as well). We then iteratively select the most significant core node in the sorted \mathcal{V}_C as the root of breadth-first tree search until all remaining tree nodes \mathcal{V}_T are explored.

The time complexity of core-tree decomposition is $\mathcal{O}(d|\mathcal{V}| + |\mathcal{E}|)$ where d denotes a user specified tree width, and the time complexity of tree compression is mainly on the sequence of breadth-first tree search, which has $\mathcal{O}(|\mathcal{V}_T| + |\mathcal{E}_T|)$, as well as the sorting of core nodes, which has $\mathcal{O}(|\mathcal{V}_C| \log |\mathcal{V}_C|)$.

The impact of core-tree decomposition and tree compression is to substantially reduce the effective graph size down to the size of the core, while keeping many of the significant p-values in the trees. However, we note that there is a potential trade-off to this computationally efficient approach. Any significant tree node that is not adjacent to a core node, and is not connected to the core by a path consisting only of other significant nodes, will not be merged into the core and therefore will not be part of the detected subgraph returned by CNSS+CoreTree. In practice, however, we find that the loss of accuracy from this approach is minimal, while the speedup in runtime is substantial.

B.6 Time Complexity Analysis of CNSS

Since CNSS Algorithm 1 applies Algorithm 2 for each significance threshold α under consideration, we focus on the time complexity of Algorithm 2 first. Algorithm 2 includes three main steps:

- Step 1 (line 1): merge adjacent significant nodes in the graph;
- Step 2 (lines 2-12): greedily merge the whole graph according to the three options as described in Appendix B.4 and record $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ throughout the merge process;
- Step 3 (lines 13-22): filter out recorded $(\mathcal{S}, N_\alpha(\mathcal{S}), N(\mathcal{S}))$ with suboptimal $N_\alpha(\mathcal{S})$ for the corresponding $N(\mathcal{S})$.

For step 1, the algorithm must iterate over all edges and record all significant and non-significant neighbors for each node. Therefore, the time complexity of step 1 is $\mathcal{O}(|\mathcal{V}| + |\mathcal{E}|)$. The time complexity of step 2 is mainly based on sorting and searching for the best merge option, which iterates over the root node’s neighbors. The sorting takes $\mathcal{O}(|\mathcal{V}| \log |\mathcal{V}|)$. For the first option of merge, we randomly merge one significant neighbor. For the second and third options, they need to iterate over all neighbors of current root node. Hence, the time complexity would be $\mathcal{O}(k|\mathcal{V}|)$ where k denotes the largest degree of a node in the network. The overall time complexity of step 2 is $\mathcal{O}(k|\mathcal{V}| + |\mathcal{V}| \log |\mathcal{V}|)$. For step 3, the time complexity is mainly on the sorting, thus time complexity is $\mathcal{O}(|\mathcal{V}| \log |\mathcal{V}|)$. Therefore, the overall time complexity of the Algorithm 2 is $\mathcal{O}(k|\mathcal{V}| + |\mathcal{V}| \log |\mathcal{V}|)$.

For CNSS Algorithm 1, it applies Algorithm 2 for each significance threshold $\alpha \in \mathcal{L}$. For obtaining $\alpha'(N, \alpha)$ for a given N and α , it requires us to apply Algorithm 2 on K replicas of datasets under the null hypothesis, unless the lower bound method is used in place of randomization testing. Therefore, the time complexity is $\mathcal{O}(K|\mathcal{L}|(k|\mathcal{V}| + |\mathcal{V}| \log |\mathcal{V}|))$. As we note below, these K replicas can be run in parallel, or alternatively, the lower bound approach avoids the need for randomization testing; in either case, the time needed to compute $\alpha'(N, \alpha)$ can be reduced by a factor of K .

B.7 Implementation Details and Reproducibility

We performed all experiments on Linux servers with the same hardware configuration (64-bit machines with Intel(R)Xeon(R) CPU E5-2680 v4 @ 2.40GHz and 251GB RAM). We implemented the CNSS in Python, and the code is accessible via the following link: <https://bit.ly/2QTVdzM>.

We set the random seed of each run under the alternative hypothesis and null hypothesis as the run index, and we used significance thresholds $\alpha \in [0.001, 0.002, \dots, 0.009, 0.01, \dots, 0.09]$.

The implementation details of each baseline method used in our evaluation are provided in Appendix C.2.

C Additional Experimental Details

C.1 Datasets

Five real-world networks were obtained from the Stanford Network Analysis Project (SNAP)^{3,4}, including 1) Twitter: a social network in where every node is a user and every edge represents a relation of follower and followee, where we do not consider the edge direction and thus treat the dataset as an undirected graph; 2) DBLP: a co-authorship network where every author is represented by a node, and two authors are connected if they publish at least one paper together; 3) SlashDot: a technology-related news social network in where every node is an user and every link represents the friendship between two users; 4) CondMat: Arxiv COND-MAT (Condensed Matter Physics) collaboration network is from the e-print arXiv and covers scientific collaborations between authors on papers submitted to the Condensed Matter category. If an author i co-authored a paper with author j , the graph contains a undirected edge between i and j . If the paper is co-authored by k authors, the co-authorship graph contains a completely connected subgraph on these k nodes; and 5) WikiVote: the network contains all the Wikipedia voting data from the inception of Wikipedia until January 2008. Nodes in the network represent Wikipedia users and a directed edge from node i to node j represents that user i voted on user j . In our experiments, we treat it as an undirected graph. The descriptive statistics of all datasets are described in Table 2.

We have also tried different sizes of true subgraphs (up to 5% of $|\mathcal{V}|$), and we only report one of them since they have consistent relative performance of methods as our reported results in the paper.

C.2 Details of Comparison Methods

We compare our proposed algorithm with six state of the art methods for event detection and anomalous subgraph detection. These six methods are commonly used as baselines for detection of anomalous subgraphs, and include:

- Linear Time Subset Scanning (LTSS) (Neill 2012) is an efficient event detection algorithm in massive data sets, where the event detection problem could be viewed as a problem of finding the subset which maximizes some score function. For score functions satisfying the LTSS property (e.g., the Berk-Jones scan statistic), the subset of data records which maximizes $F(\mathcal{S})$ can be found by ordering the records according to some “priority” function and searching over groups consisting of the top- k highest priority records, requiring a linear rather than exponential number of subsets to be evaluated (Neill 2012). The time complexity of LTSS is $O(|\mathcal{V}| \log |\mathcal{V}|)$. However, LTSS does not enforce the graph connectivity constraint and may produce a disconnected subset of graph nodes. Thus we use the largest connected component in the detected subset of nodes as the detected subgraph \mathcal{S} . We obtained the code from the authors, and

³<https://snap.stanford.edu/data/>

⁴License information for these datasets is found at: <https://snap.stanford.edu/snap/license.html>

Table 2: Descriptive Statistics of Real-World Networks.

Dataset	Vertices $ \mathcal{V} $	Edges $ \mathcal{E} $	Density	Core Vertices $ \mathcal{V}_C $	Core Density	True Nodes $ \mathcal{S} $
WikiVote	7,066	100,736	0.00403	1,823	0.0425	100
CondMat	21,363	91,286	0.0004	2,513	0.00487	200
Twitter	81,309	1,342,296	0.000406	17,337	0.0041	1,000
SlashDot	82,168	504,230	0.000149	10,599	0.0046	1,000
DBLP	317,080	1,049,866	0.0000208	22,354	0.00054	1,000

we use the BJ scan statistic as the objective to maximize. We iterate over the list of significance thresholds $\alpha \in \{0.001, \dots, 0.009, 0.01, \dots, 0.09\}$ to find the maximum BJ score.

- EventTree (Rozenshtein et al. 2014): is an event detection algorithm, which defines an event to be a connected subgraph of nodes in the network that are close to each other and have high activity levels. Unlike the other methods considered here, the objective function of EventTree is not a log-likelihood ratio statistic. Rather, the objective is to maximize $Q(\mathcal{S}) = \lambda W(\mathcal{S}) - D(\mathcal{S})$, where $W(\mathcal{S}) = \sum_{v \in \mathcal{S}} w(v)$ measures the total node weight value of a subgraph \mathcal{S} , $D(\mathcal{S}) = \frac{1}{2} \sum_{u \in \mathcal{S}} \sum_{v \in \mathcal{S}} d(u, v)$ measures the distance value of subgraph \mathcal{S} , and λ is a normalization coefficient. We optimized the code provided by the authors with a more efficient PCST solver⁵ that reduces the time complexity from $\mathcal{O}(|\mathcal{V}|^2 \log |\mathcal{V}|)$ to $\mathcal{O}(|\mathcal{E}| \log |\mathcal{V}|)$ but has identical detection performance. Since we use the simulated p-values for other methods, we use the reciprocal of the p-value as the node weight for EventTree. The parameter λ also controls the granularity of the detected event, and we tuned it over the list $[0.001, \dots, 0.009, 0.01, \dots, 0.09]$.
- ColorCoding (Cadena, Chen, and Vullikanti 2019): is an unified framework for optimizing a large class of parametric and non-parametric scan statistics for networks with connectivity constraints. It is the only baseline method other than DFGS that provides a rigorous solution guarantee. The time complexity of ColorCoding is $\mathcal{O}(2^k \cdot e^k |\mathcal{E}| \log(\frac{|\mathcal{V}|}{\epsilon}))$ for a $(1 - \epsilon)$ approximate solution, where k is the effective solution size (Cadena, Chen, and Vullikanti 2019). It is extremely expensive when k is large. However, it provides additional, heuristic preprocessing steps to reduce the graph size such that, empirically, $k < 10$ is sufficient to find good solutions. We used the code provided by the authors and implemented the approximation refinement based on suggestions from the authors. We tuned the parameter refinement coefficient β over $[0.1, \dots, 0.9]$ for different signal strengths, and ended up with $\beta = 0.9$ for $\mu = 5$, $\beta = 0.8$ for $\mu = 4$, $\beta = 0.7$ for $\mu = 3$, $\beta = 0.6$ for $\mu = 2$, and $\beta = 0.5$ for $\mu = 1.5$ which achieve the best performance. In addition,

we set $k = 5$ with 300 iterations as suggested by the authors.

- Non-parametric Heterogeneous Graph Scan (NPHGS) (Chen and Neill 2014): optimizes the original Berk-Jones nonparametric scan statistic over connected subgraphs, using a greedy growth heuristic. The nonparametric scan statistics are free of distributional assumptions and can be applied to anomalous connected subgraph detection in heterogeneous graph data, in contrast to traditional parametric scan statistics (e.g., the Kulldorff statistic). The time complexity of NPHGS is $\mathcal{O}(|\mathcal{V}|^2 \log |\mathcal{V}|)$. We used the code provided by the authors, and chose the BJ scan statistic as the objective score to optimize with parameter $\alpha \in \{0.001, \dots, 0.009, 0.01, \dots, 0.09\}$.
- Additive Graph Scan (AdditiveScan) (Speakman, Zhang, and Neill 2013): was proposed as an efficient heuristic alternative to DFGS which can be used to identify the high-scoring (most positive) connected subsets in a given graph structure with real-valued weights at each node. This method stems from two facts that 1) additive functions satisfy the LTSS property, which could leverage DFGS, however, 2) computation time of DFGS is exponential in the graph size. The time complexity of AdditiveScan is $\mathcal{O}(|\mathcal{V}|^2 \sqrt{|\mathcal{V}|})$. We adopted the implementation by the authors, and chose the BJ scan statistic as the objective score to optimize with parameter $\alpha \in \{0.001, \dots, 0.009, 0.01, \dots, 0.09\}$.
- Depth First Graph Scan (DFGS) (Speakman, McFowland III, and Neill 2015): is a graph scan method that is guaranteed to find the exact solution, however, the worst case complexity of DFGS is exponential in the neighborhood size k . If no pruning was performed, DFGS would evaluate all connected subsets, requiring $\mathcal{O}(2^k)$ run time; however, it is able to rule out many connected subsets as provably suboptimal, reducing complexity to $\mathcal{O}(q^k)$ for some constant $1 < q < 2$, where q is dependent on the proportion of subsets that are pruned. We adopted the implementation by the authors directly with suggested hyperparameters, and chose the BJ scan statistic as the objective score to optimize with parameter $\alpha \in \{0.001, \dots, 0.009, 0.01, \dots, 0.09\}$.

C.3 Experimental Setup and Evaluation Metrics

For each of the five real-world graph structures enumerated above, we simulate 200 runs of p-values generated under the null hypothesis and 50 runs of p-values generated under each of five different alternative hypotheses with $\mu \in [1.5, 2, 3, 4, 5]$. We run all algorithms on these 2250

⁵Hegde, Chinmay, Piotr Indyk, and Ludwig Schmidt. “A fast, adaptive variant of the Goemans-Williamson scheme for the prize-collecting Steiner tree problem.” Workshop of the 11th DIMACS Implementation Challenge. Vol. 2, 2014.

graphs and record the detected subgraphs and their corresponding scores, and then performance evaluations in terms of detection power, precision, recall, and F-score are conducted. Let \mathcal{R} be the ground truth nodes in the anomalous subgraph and let \mathcal{S} be the detected subgraph.

- Detection power measures the ability of a method to distinguish between graphs with or without an affected subgraph. It is computed based on the following steps: 1) compute BJ score for each detected subgraph; 2) for each alternative run, we conduct a hypothesis test with significance level $\alpha = 0.05$ by setting p-value as the proportion of null runs that have higher BJ score than the alternative run; 3) compute the proportion of hypothesis tests (for each method, for each real-world graph, for each signal strength μ) that reject the null hypothesis.
- Precision is the ratio of the number of detected true positive nodes divided by the total number of detected nodes, that is,

$$\text{Precision} = \frac{|\mathcal{R} \cap \mathcal{S}|}{|\mathcal{S}|}. \quad (13)$$

- Recall is the ratio of the number of detected true positive nodes divided by the number of nodes in true subgraph, that is,

$$\text{Recall} = \frac{|\mathcal{R} \cap \mathcal{S}|}{|\mathcal{R}|}. \quad (14)$$

- F-score is the harmonic mean of precision and recall, that is,

$$\text{F-score} = \frac{2 \cdot \text{Precision} \cdot \text{Recall}}{\text{Precision} + \text{Recall}}. \quad (15)$$

We note that precision, recall, and F-score are each averaged over the 50 graphs created for a given signal strength μ .

C.4 Additional Experimental Results

We now present a detailed comparison of the CNSS and baseline methods with respect to run time, detection power, and detection performance (precision, recall, and F-measure). Note that, for the larger datasets, we do not report the run time and performance of some baseline methods due to extremely long clock run time (over 2 weeks on 250 CPUs) to finish all experiments over 50 replicas of datasets under the alternative hypothesis.

Run time. As shown in Table 3, not including the pre-processing time needed to compute the distribution of $\alpha'(N, \alpha)$ for a given graph structure, our method has competitive run time that is faster than NPHGS, AdditiveScan, and DFGS, which is aligned with the time complexity analysis. We also observe substantial speedups (ranging from 2.6x for WikiVote to 28x for CondMat) for CNSS+CoreTree as compared to CNSS without core-tree decomposition. These timing results are not impacted by the approach used to compute $\alpha'(N, \alpha)$, i.e., calibration by randomization tests versus calibration by lower bounds versus no calibration.

However, as shown in Table 4, the total pre-processing time needed to perform calibration by randomization testing

is large because the same search must be performed on a large number K of replica datasets generated under the null hypothesis \mathcal{H}_0 , thus multiplying the run time by K . We used $K = 200$ for our experiments, thus substantially increasing run time. However, we note that these null runs are entirely independent and thus can be easily parallelized. Moreover, this preprocessing step must only be performed once for a given graph structure, and can be reused if the graph structure remains constant, for example, for daily monitoring of disease cases with a graph structure defined by zip code adjacency. Additionally, we observe that the use of core-tree decomposition resulted in similar speedups, ranging from 2.4x for WikiVote to 28x on CondMat, on the replica datasets. Most critically, the use of lower bounds in place of randomization testing eliminates the need to generate and search over the large number of replica datasets, resulting in huge savings in preprocessing time. These speedups ranged from 300x to 2000x as compared to randomization testing with core-tree decomposition, assuming $K = 200$ and no parallelization for the randomization tests.

Detection power. Table 5 compares the detection power (i.e., the proportion of signals detected, at a fixed false positive rate of 0.05) for CNSS and baseline methods across the five real-world datasets and varying signal strengths. Note that several of the slower methods were not run for the largest graphs due to excessive run times needed to perform detection on the null graphs.

We observe that all three of the calibrated CNSS methods (CNSS, CNSS+CoreTree, and CNSS+LowerBound) achieve perfect detection power across all of the real-world graphs and signal strengths considered. In contrast, the uncalibrated CNSS methods and the baseline methods have substantially reduced detection power for low signal strengths, particularly on the smaller graphs (WikiVote and CondMat) which also had a smaller number of true nodes generated under the alternative hypothesis \mathcal{H}_1 . Interestingly, even the uncalibrated CNSS methods outperformed the baseline methods with respect to detection power; among the baseline methods, EventTree and ColorCoding tended to outperform LTSS, NPHGS, and AdditiveScan.

Detection performance. The detection performances are shown in Table 6. The bold number indicates that method is significantly better than other methods. Overall, our proposed CNSS outperforms baseline methods under different signal strengths μ on the various network structures. Specifically, the calibrated BJ score helps to precisely pinpoint the true affected subgraph as the strength of signal increases. The use of core-tree decomposition and lower bounds do not have substantial effects on detection performance for these five real-world datasets, while significantly reducing run time. On the other hand, the baseline methods do not have consistent performance over different values of μ with different network structures. LTSS has the worst detection performance, because it does not enforce the graph connectivity constraints and thus picks out a subset of disconnected, individually anomalous nodes that do not accurately reflect the true affected subgraph. EventTree has relatively good performance on Twitter and SlashDot datasets when the event signal is not strong. However, when event signal

Table 3: Run Time on All Datasets. The run time of our method is reported for 18 different α values using a single processor, which could be parallelized to speed up the run time. We implement all discussed algorithms and perform the experiments on Linux servers with the same hardware configuration (64-bit machines with Intel(R)Xeon(R) CPU E5-2680 v4 @ 2.40GHz and 251GB memory). Note that the run times for CNSS and CNSS+CoreTree do not include the pre-processing time required to estimate $\alpha'(N, \alpha)$, and thus are independent of the calibration approach (randomization test versus lower bounds versus no calibration).

Methods	WikiVote	CondMat	Twitter	SlashDot	DBLP
	Run Time (sec.)				
LTSS	21	24	619	243	1425
EventTree	23	25	179	186	1019
ColorCoding	5220	8295	66690	29790	124956
NPHGS	8912	52046	998624	496587	\times
AdditiveScan	17950	123100	\times	\times	\times
DFGS	22791	\times	\times	\times	\times
CNSS	1771	43325	489624	447800	\times
CNSS+CoreTree	685	1544	128812	45208	185053

Table 4: Preprocessing Time Comparison for the Computation of $\alpha'(N, \alpha)$. The run time of randomization testing is the run time of a single CNSS run under \mathcal{H}_0 times the number of replica datasets K , and we used $K = 200$ runs for our experiments.

Methods	WikiVote	CondMat	Twitter	SlashDot	DBLP
	Run Time (sec.)	Run Time (sec.)	Run Time (sec.)	Run Time (sec.)	Run Time (sec.)
RandomizationTest	$1602 \times K$	$28341 \times K$	$299349 \times K$	$375999 \times K$	\times
RandomizationTest+CoreTree	$660 \times K$	$1026 \times K$	$107192 \times K$	$40124 \times K$	$147086 \times K$
LowerBounds	59	504	16094	9073	87832

is strong, the subgraph detected by EventTree includes many noisy nodes around true nodes to reach a higher score, thus dramatically harming precision; this pattern is also observed for the other baseline methods, while our proposed calibrated scan approach converges to both high precision and high recall as the signal strength increases. DFGS has relatively good performance when the graph is small, such as CondMat data with a strong event signal, but does not perform well on WikiVote and quickly becomes computationally infeasible for the larger graphs. ColorCoding, AdditiveScan, and NPHGS have similar performance: all of them suffer from the similar issue as EventTree in that, when the event signal is strong, they include many individually significant nodes that are not part of the true affected subgraph, thus reducing precision and F-score.

In addition, we show the average performance (F-score) over various signal strengths and network structures in Figure 7. We can see that CNSS, CNSS+CoreTree, and CNSS+LowerBound have much better average performance than all baselines, while CNSS without calibration (CNSS+NoCalib) performs poorly.

These results demonstrate the advantage of the novel calibration approach proposed here for precisely identifying the true affected subgraph, as well as the utility of our core-tree decomposition and lower bound approaches for enabling scalability and computational feasibility for large graphs.

C.5 Results with piecewise constant p-values

Though the nonparametric scan statistics (calibrated or uncalibrated) do not make any assumptions of Gaussianity, we used Gaussian signals in our simulation experiments to show

that the calibrated NPSS formulation can achieve high detection performance even when the injected signal does not necessarily obey our specific modeling assumptions. Gaussian mean-shift signals are simple, have a natural way of measuring signal strength, and are frequently used in the literature, e.g., by Reyna et al. (2021) and Chitra et al. (2021). However, one might ask what happens when the Berk-Jones NPSS modeling assumptions are precisely correct, and the resulting p-values are piecewise constant under $\mathcal{H}_1(\mathcal{S})$. Does the uncalibrated Berk-Jones statistic still fail to detect these signals due to miscalibration, and does the calibrated BJ statistic still outperform competing approaches by a wide margin?

To explore these questions, we performed additional simulations using the WikiVote and CondMat datasets, comparing the calibrated scans (CNSS, CNSS+CoreTree, and CNSS+LowerBounds) with the uncalibrated scan (CNSS+NoCalib) and the various baselines (LTSS, EventTree, NPHGS, AdditiveScan, DFGS, and ColorCoding) with respect to detection power, precision, recall, and F-score. As in the main evaluation, we simulated the true subgraph \mathcal{S} using a random walk with size roughly $0.01|\mathcal{V}|$ (see Table 2), and reported the average performance over 50 runs of simulations of true subgraphs and p-values, for each signal strength, on each network structure. However, for these runs, we assumed piecewise constant p-values, where each p-value $p_i \in \mathcal{S}$ is drawn from Uniform[0, 0.01] with probability $q \cdot 0.01$, and from Uniform[0.01, 1] with probability $1 - q \cdot 0.01$. Signal strengths $q \in \{10, 25, 50, 75, 100\}$ were used for these simulations, and p-values outside subset S were drawn from

Table 5: Detection Power Comparison on Five Real World Datasets with Gaussian Signals.

Methods	WikiVote ($\mu = 1.5$)	WikiVote ($\mu = 2$)	WikiVote ($\mu = 3$)	WikiVote ($\mu = 4$)	WikiVote ($\mu = 5$)
	Detection Power	Detection Power	Detection Power	Detection Power	Detection Power
LTSS	0.28	0.4	0.96	1.0	1.0
EventTree	0.6	0.94	1.0	1.0	1.0
ColorCoding	0.8	1.0	1.0	1.0	1.0
NPHGS	0.0	0.0	0.0	0.0	0.0
AdditiveScan	0.0	0.0	0.0	0.0	0.0
CNSS+NoCalib	0.94	1.0	1.0	1.0	1.0
CNSS+CoreTree+NoCalib	0.86	0.98	0.98	1.0	1.0
CNSS+CoreTree	1.0	1.0	1.0	1.0	1.0
CNSS+LowerBound	1.0	1.0	1.0	1.0	1.0
CNSS	1.0	1.0	1.0	1.0	1.0

Methods	CondMat ($\mu = 1.5$)	CondMat ($\mu = 2$)	CondMat ($\mu = 3$)	CondMat ($\mu = 4$)	CondMat ($\mu = 5$)
	Detection Power	Detection Power	Detection Power	Detection Power	Detection Power
LTSS	0.3	0.68	1.0	1.0	1.0
EventTree	0.66	0.94	1.0	1.0	1.0
ColorCoding	0.0	0.8	1.0	1.0	1.0
NPHGS	0.0	0.1	0.72	0.96	1.0
AdditiveScan	0.32	0.36	0.36	0.36	0.38
CNSS+NoCalib	0.96	1.0	1.0	1.0	1.0
CNSS+CoreTree+NoCalib	0.86	1.0	1.0	1.0	1.0
CNSS+CoreTree	1.0	1.0	1.0	1.0	1.0
CNSS+LowerBound	1.0	1.0	1.0	1.0	1.0
CNSS	1.0	1.0	1.0	1.0	1.0

Methods	Twitter ($\mu = 1.5$)	Twitter ($\mu = 2$)	Twitter ($\mu = 3$)	Twitter ($\mu = 4$)	Twitter ($\mu = 5$)
	Detection Power	Detection Power	Detection Power	Detection Power	Detection Power
LTSS	0.0	0.0	0.0	1.0	1.0
EventTree	1.0	1.0	1.0	1.0	1.0
ColorCoding	1.0	1.0	1.0	1.0	1.0
NPHGS	0.0	0.0	0.0	0.02	0.34
CNSS+NoCalib	1.0	1.0	1.0	1.0	1.0
CNSS+CoreTree+NoCalib	1.0	1.0	1.0	1.0	1.0
CNSS+CoreTree	1.0	1.0	1.0	1.0	1.0
CNSS+LowerBound	1.0	1.0	1.0	1.0	1.0
CNSS	1.0	1.0	1.0	1.0	1.0

Methods	SlashDot ($\mu = 1.5$)	SlashDot ($\mu = 2$)	SlashDot ($\mu = 3$)	SlashDot ($\mu = 4$)	SlashDot ($\mu = 5$)
	Detection Power	Detection Power	Detection Power	Detection Power	Detection Power
LTSS	1.0	1.0	1.0	1.0	1.0
EventTree	1.0	1.0	1.0	1.0	1.0
ColorCoding	1.0	1.0	1.0	1.0	1.0
NPHGS	0.0	0.0	0.0	0.04	0.38
CNSS+NoCalib.	1.0	1.0	1.0	1.0	1.0
CNSS+CoreTree+NoCalib	1.0	1.0	1.0	1.0	1.0
CNSS+CoreTree	1.0	1.0	1.0	1.0	1.0
CNSS+LowerBound	1.0	1.0	1.0	1.0	1.0
CNSS	1.0	1.0	1.0	1.0	1.0

Methods	DBLP ($\mu = 1.5$)	DBLP ($\mu = 2$)	DBLP ($\mu = 3$)	DBLP ($\mu = 4$)	DBLP ($\mu = 5$)
	Detection Power	Detection Power	Detection Power	Detection Power	Detection Power
LTSS	1.0	1.0	1.0	1.0	1.0
EventTree	0.98	1.0	1.0	1.0	1.0
ColorCoding	1.0	1.0	1.0	1.0	1.0
CNSS+CoreTree	1.0	1.0	1.0	1.0	1.0

Table 6: Detection Performance Results (Average Precision, Recall, and F-score) on Five Real World Datasets with Gaussian Signals. The bold number indicates that method has a significantly higher F-score than the other methods. Statistical significance is computed using paired t-tests ($p < 0.05$).

Methods	WikiVote ($\mu = 1.5$)			WikiVote ($\mu = 2$)			WikiVote ($\mu = 3$)			WikiVote ($\mu = 4$)			WikiVote ($\mu = 5$)		
	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score
LTSS	0.023	0.144	0.039	0.024	0.154	0.042	0.022	0.145	0.039	0.024	0.157	0.042	0.024	0.155	0.041
EventTree	0.035	0.041	0.037	0.041	0.060	0.049	0.036	0.075	0.049	0.034	0.083	0.048	0.026	0.244	0.047
ColorCoding	0.074	0.664	0.132	0.111	0.801	0.195	0.145	0.953	0.252	0.174	0.997	0.297	0.244	0.920	0.376
NPHGS	0.144	0.547	0.227	0.185	0.751	0.297	0.216	0.946	0.351	0.225	0.997	0.367	0.346	0.949	0.432
AdditiveScan	0.143	0.547	0.227	0.185	0.751	0.296	0.216	0.946	0.351	0.225	0.997	0.367	0.226	1.000	0.368
DFGS	0.154	0.523	0.235	0.204	0.703	0.311	0.206	0.804	0.325	0.219	0.836	0.343	0.203	0.800	0.321
CNSS+NoCalib	0.068	0.628	0.123	0.087	0.802	0.156	0.102	0.959	0.184	0.106	0.998	0.191	0.106	1.000	0.192
CNSS+CoreTree+NoCalib	0.073	0.652	0.132	0.091	0.813	0.163	0.105	0.960	0.189	0.109	0.998	0.197	0.110	1.000	0.198
CNSS+CoreTree	0.233	0.400	0.265	0.285	0.641	0.373	0.752	0.578	0.601	0.923	0.803	0.858	0.968	0.965	0.966
CNSS+LowerBound	0.213	0.233	0.218	0.360	0.379	0.361	0.737	0.567	0.630	0.891	0.810	0.847	0.951	0.967	0.958
CNSS	0.232	0.401	0.257	0.289	0.645	0.372	0.706	0.604	0.583	0.921	0.803	0.858	0.965	0.965	0.965
Methods	CondMat ($\mu = 1.5$)			CondMat ($\mu = 2$)			CondMat ($\mu = 3$)			CondMat ($\mu = 4$)			CondMat ($\mu = 5$)		
	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score
LTSS	0.017	0.083	0.029	0.018	0.090	0.030	0.018	0.094	0.031	0.019	0.096	0.032	0.017	0.088	0.029
EventTree	0.014	0.204	0.027	0.015	0.216	0.028	0.014	0.208	0.026	0.014	0.209	0.026	0.014	0.214	0.027
ColorCoding	0.352	0.074	0.094	0.169	0.564	0.228	0.254	0.899	0.379	0.255	0.894	0.388	0.313	1.000	0.470
NPHGS	0.203	0.416	0.272	0.245	0.626	0.351	0.304	0.939	0.459	0.315	0.995	0.478	0.345	0.999	0.501
AdditiveScan	0.202	0.478	0.283	0.247	0.675	0.361	0.286	0.945	0.439	0.306	0.995	0.467	0.304	1.000	0.467
DFGS	0.459	0.078	0.132	0.616	0.193	0.290	0.819	0.656	0.725	0.861	0.937	0.897	0.866	0.995	0.926
CNSS+NoCalib	0.045	0.629	0.084	0.056	0.779	0.104	0.068	0.957	0.126	0.070	0.997	0.132	0.070	1.000	0.132
CNSS+CoreTree+NoCalib	0.057	0.620	0.105	0.071	0.779	0.130	0.086	0.958	0.158	0.089	0.995	0.163	0.089	0.999	0.164
CNSS+CoreTree	0.235	0.326	0.248	0.334	0.484	0.383	0.533	0.654	0.560	0.883	0.782	0.824	0.932	0.890	0.905
CNSS+LowerBound	0.078	0.451	0.132	0.115	0.609	0.193	0.267	0.825	0.394	0.883	0.833	0.857	0.909	0.975	0.941
CNSS	0.258	0.361	0.278	0.340	0.563	0.408	0.735	0.685	0.687	0.916	0.821	0.866	0.985	0.974	0.979
Methods	Twitter ($\mu = 1.5$)			Twitter ($\mu = 2$)			Twitter ($\mu = 3$)			Twitter ($\mu = 4$)			Twitter ($\mu = 5$)		
	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score
LTSS	0.075	0.726	0.136	0.084	0.829	0.152	0.096	0.974	0.175	0.098	0.998	0.178	0.098	0.999	0.179
EventTree	0.222	0.283	0.249	0.318	0.449	0.371	0.454	0.790	0.577	0.502	0.962	0.660	0.510	0.997	0.675
ColorCoding	0.060	0.661	0.109	0.084	0.783	0.153	0.114	0.957	0.204	0.136	0.997	0.239	0.153	1.000	0.265
NPHGS	0.110	0.553	0.183	0.136	0.737	0.230	0.166	0.951	0.282	0.172	0.996	0.294	0.172	1.000	0.293
CNSS+NoCalib	0.063	0.601	0.114	0.079	0.764	0.144	0.097	0.954	0.177	0.101	0.996	0.184	0.102	0.999	0.185
CNSS+CoreTree+NoCalib	0.071	0.617	0.127	0.088	0.774	0.158	0.106	0.957	0.192	0.111	0.996	0.199	0.111	0.999	0.200
CNSS+CoreTree	0.137	0.496	0.215	0.198	0.662	0.305	0.788	0.502	0.613	0.895	0.821	0.856	0.923	0.972	0.947
CNSS+LowerBound	0.150	0.264	0.181	0.260	0.378	0.302	0.725	0.515	0.599	0.909	0.819	0.861	0.957	0.972	0.964
CNSS	0.137	0.491	0.211	0.167	0.701	0.267	0.728	0.508	0.579	0.880	0.820	0.849	0.910	0.972	0.940
Methods	Slashdot ($\mu = 1.5$)			Slashdot ($\mu = 2$)			Slashdot ($\mu = 3$)			Slashdot ($\mu = 4$)			Slashdot ($\mu = 5$)		
	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score
LTSS	0.099	0.663	0.173	0.116	0.823	0.203	0.130	0.973	0.229	0.133	0.998	0.234	0.132	0.999	0.234
EventTree	0.264	0.312	0.285	0.456	0.482	0.410	0.476	0.810	0.599	0.527	0.967	0.682	0.526	0.996	0.689
ColorCoding	0.077	0.713	0.140	0.112	0.797	0.197	0.145	0.959	0.251	0.172	0.997	0.293	0.196	1.000	0.327
NPHGS	0.153	0.539	0.238	0.186	0.731	0.296	0.215	0.949	0.351	0.223	0.996	0.364	0.222	1.000	0.363
CNSS+NoCalib.	0.063	0.693	0.116	0.074	0.822	0.136	0.086	0.965	0.159	0.089	0.997	0.164	0.089	0.999	0.165
CNSS+CoreTree+NoCalib	0.078	0.709	0.141	0.091	0.831	0.164	0.104	0.968	0.189	0.108	0.997	0.195	0.108	0.999	0.195
CNSS+CoreTree	0.157	0.529	0.242	0.192	0.720	0.303	0.591	0.672	0.629	0.886	0.829	0.857	0.905	0.971	0.937
CNSS+LowerBound	0.070	0.655	0.127	0.086	0.786	0.155	0.716	0.541	0.609	0.873	0.829	0.850	0.891	0.972	0.930
CNSS	0.159	0.528	0.245	0.193	0.722	0.304	0.587	0.678	0.629	0.872	0.829	0.850	0.901	0.971	0.935
Methods	DBLP ($\mu = 1.5$)			DBLP ($\mu = 2$)			DBLP ($\mu = 3$)			DBLP ($\mu = 4$)			DBLP ($\mu = 5$)		
	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score
LTSS	0.055	0.565	0.100	0.073	0.777	0.134	0.087	0.968	0.159	0.089	0.998	0.164	0.089	1.000	0.164
EventTree	0.081	0.357	0.132	0.119	0.538	0.194	0.174	0.834	0.288	0.198	0.972	0.329	0.203	0.998	0.340
ColorCoding	0.056	0.435	0.100	0.104	0.656	0.180	0.152	0.943	0.261	0.179	0.996	0.304	0.205	1.000	0.340
CNSS+CoreTree+NoCalib	0.022	0.623	0.043	0.028	0.775	0.054	0.034	0.955	0.066	0.035	0.994	0.068	0.035	0.995	0.068
CNSS+CoreTree	0.135	0.265	0.179	0.190	0.448	0.267	0.355	0.560	0.435	0.550	0.821	0.659	0.831	0.949	0.886

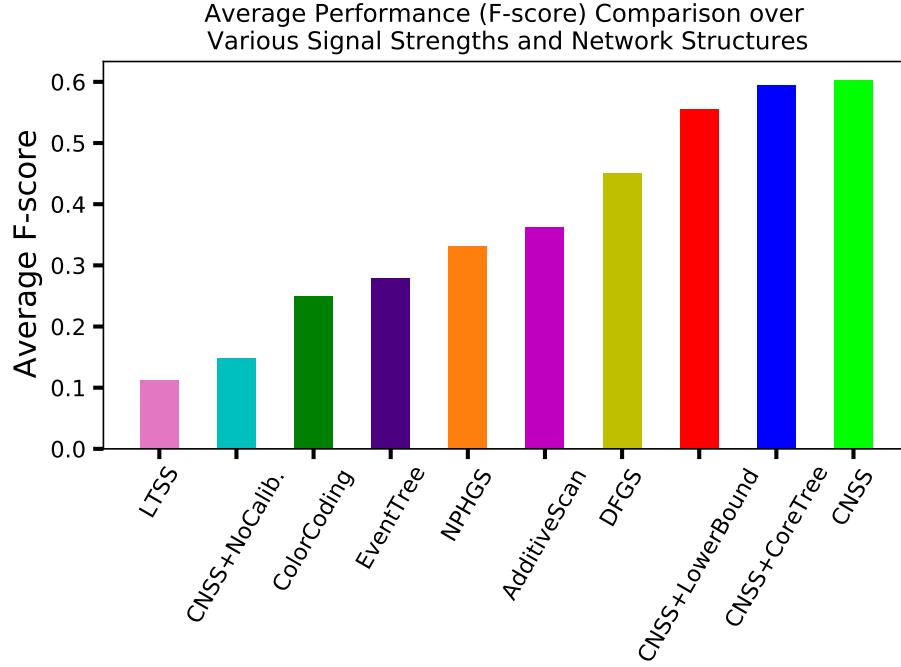


Figure 7: We compare the average performance (F-score) over various signal strengths and network structures to mimic the various real-world scenarios. We performed two-sample paired t-tests between each baseline and CNSS, and we find that all baselines have significantly lower performance than CNSS with $p\text{-value} < 0.05$.

Uniform[0,1] as usual. The results of these simulations are shown in Tables 7 and 8.

We observe that the results for piecewise constant p-values are highly consistent with those for Gaussian signals, demonstrating that it is miscalibration (not the shape of the signal) that is causing the uncalibrated methods to perform poorly. As we observed for the Gaussian signals, our proposed CNSS and CNSS+CoreTree outperformed all baselines and the uncalibrated CNSS+NoCalib by a wide margin in terms of detection accuracy and detection power, while the uncalibrated methods suffered from low detection power for low signal strengths, and low precision (and therefore low F-score) across all signal strengths. CNSS+LowerBounds consistently outperformed the uncalibrated scan and baseline methods for WikiVote across all signal strengths, and for CondMat for high signal strengths. For low signal strengths on CondMat, CNSS+LowerBounds achieved higher detection power and recall than the uncalibrated scan and baseline methods, but had lower precision and F-score.

Finally, we examined the mean and standard deviation of the selected value of α for each method across the 50 runs for each dataset and signal strength, noting that the signal was injected with a true α value of 0.01. These results are shown in Table 9. We observe that the uncalibrated scans and baseline NPSS methods fail to identify the true α value, instead consistently selecting the largest α value considered, i.e., $\alpha = 0.09$. In contrast, as the signal strength increases, CNSS, CNSS+CoreTree, and CNSS+LowerBounds are

all able to reliably identify the value, $\alpha = 0.01$, corresponding to the true injected signal.

D Case Studies

D.1 Black Lives Matter Event Detection in Twitter

In addition to the COVID-19 case study described in the main paper, we also conduct another case study using tweets with hashtag `#BlackLivesMatter`⁶ collected from August 8th, 2014 to August 31st, 2015 to discover events related to this social movement for racial justice during these 56 weeks. There are 442,077 unique tweets and 42,898 unique hashtags in total. This Twitter dataset contains account user names which may reveal personally identifiable information. However, we pre-processed the data to remove the user names, and only used the hashtags in each tweet and corresponding creation timestamps in our research. We generate a temporal graph with 100,123 nodes and 257,641 edges based on the mentioned hashtags in tweets during these 56 weeks, in which each node represents a mentioned hashtag in a particular week. We connect two nodes (two hashtags in a week) if they were co-mentioned in at least one tweet that week. We also add temporal edges between node $v_{i,t}$ (hashtag i in week t) and node $v_{i,t+1}$ (hashtag i in week $t+1$) if both nodes have non-zero mentioned counts in each week. In addition, we also add dummy nodes and edges to smooth

⁶<https://dataverse.harvard.edu/dataset.xhtml?persistentId=doi:10.7910/DVN/IQ525U&version=1.0>

Table 7: Detection Power Comparison on WikiVote and CondMat datasets, assuming piecewise constant p-values.

Methods	WikiVote ($q = 10$)	WikiVote ($q = 25$)	WikiVote ($q = 50$)	WikiVote ($q = 75$)	WikiVote ($q = 100$)
	Detection Power				
LTSS	0.12	0.18	0.4	0.64	0.86
EventTree	0.12	0.2	0.24	0.42	0.4
ColorCoding	0.02	0.12	0.98	1.0	1.0
NPHGS	0.0	0.0	0.0	0.0	0.0
AdditiveScan	0.04	0.12	0.52	0.98	1.0
CNSS+NoCalib	0.12	0.44	0.86	0.98	1.0
CNSS+CoreTree+NoCalib	0.08	0.4	0.8	0.98	1.0
CNSS+CoreTree	1.0	1.0	1.0	1.0	1.0
CNSS+LowerBound	1.0	1.0	1.0	1.0	1.0
CNSS	1.0	1.0	1.0	1.0	1.0

Methods	CondMat ($q = 10$)	CondMat ($q = 25$)	CondMat ($q = 50$)	CondMat ($q = 75$)	CondMat ($q = 100$)
	Detection Power				
LTSS	0.06	0.14	0.42	0.76	0.92
EventTree	0.06	0.1	0.28	0.38	0.54
ColorCoding	0.08	0.64	1.0	1.0	1.0
NPHGS	0.0	0.0	0.0	0.1	0.86
AdditiveScan	0.04	0.34	0.88	1.0	1.0
CNSS+NoCalib	0.12	0.36	0.84	0.98	1.0
CNSS+CoreTree+NoCalib	0.44	0.72	0.94	1.0	1.0
CNSS+CoreTree	1.0	1.0	1.0	1.0	1.0
CNSS+LowerBound	1.0	1.0	1.0	1.0	1.0
CNSS	1.0	1.0	1.0	1.0	1.0

Table 8: Detection Performance Results (Average Precision, Recall, and F-score) on WikiVote and CondMat datasets, assuming piecewise constant p-values. The bold number indicates that method has a significantly higher F-score than the other methods. Statistical significance is computed using paired t-tests ($p < 0.05$).

Methods	WikiVote ($q = 10$)			WikiVote ($q = 25$)			WikiVote ($q = 50$)			WikiVote ($q = 75$)			WikiVote ($q = 100$)		
	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score									
LTSS	0.016	0.100	0.027	0.015	0.101	0.027	0.015	0.104	0.027	0.015	0.107	0.027	0.015	0.110	0.027
EventTree	0.058	0.006	0.010	0.083	0.008	0.015	0.064	0.010	0.017	0.060	0.010	0.017	0.049	0.010	0.016
ColorCoding	0.045	0.346	0.080	0.068	0.408	0.117	0.105	0.586	0.178	0.148	0.779	0.249	0.200	1.000	0.333
NPHGS	0.051	0.159	0.077	0.089	0.299	0.137	0.142	0.528	0.223	0.187	0.761	0.300	0.227	1.000	0.369
AdditiveScan	0.051	0.160	0.078	0.089	0.299	0.137	0.142	0.528	0.223	0.187	0.761	0.300	0.226	1.000	0.369
DFGS	0.052	0.159	0.078	0.092	0.292	0.139	0.150	0.501	0.228	0.184	0.684	0.288	0.209	0.820	0.329
CNSS+NoCalib	0.034	0.303	0.061	0.046	0.422	0.084	0.066	0.612	0.120	0.086	0.805	0.155	0.105	1.000	0.190
CNSS+CoreTree+NoCalib	0.035	0.301	0.063	0.049	0.423	0.088	0.070	0.614	0.125	0.090	0.806	0.162	0.110	1.000	0.198
CNSS+CoreTree	0.169	0.164	0.102	0.467	0.211	0.275	0.616	0.471	0.522	0.714	0.730	0.718	0.651	0.992	0.773
CNSS+LowerBound	0.068	0.174	0.096	0.127	0.312	0.177	0.274	0.534	0.355	0.404	0.760	0.526	0.485	1.000	0.653
CNSS	0.163	0.246	0.107	0.464	0.218	0.283	0.608	0.473	0.521	0.679	0.736	0.695	0.595	0.992	0.730

Methods	CondMat ($q = 10$)			CondMat ($q = 25$)			CondMat ($q = 50$)			CondMat ($q = 75$)			CondMat ($q = 100$)		
	Prec.	Rec.	F-Score	Prec.	Rec.	F-Score									
LTSS	0.009	0.092	0.017	0.009	0.093	0.017	0.009	0.095	0.017	0.009	0.097	0.017	0.009	0.099	0.017
EventTree	0.055	0.007	0.012	0.052	0.008	0.014	0.051	0.010	0.016	0.053	0.012	0.019	0.053	0.014	0.021
ColorCoding	0.036	0.120	0.055	0.076	0.228	0.113	0.139	0.460	0.213	0.204	0.730	0.318	0.273	1.000	0.429
NPHGS	0.052	0.061	0.055	0.103	0.154	0.122	0.193	0.398	0.258	0.261	0.697	0.379	0.316	1.000	0.480
AdditiveScan	0.048	0.076	0.059	0.100	0.182	0.128	0.184	0.433	0.257	0.247	0.715	0.367	0.298	1.000	0.459
DFGS	0.046	0.076	0.057	0.093	0.185	0.123	0.174	0.425	0.246	0.240	0.709	0.358	0.294	1.000	0.454
CNSS+NoCalib	0.021	0.289	0.039	0.030	0.418	0.056	0.044	0.614	0.081	0.056	0.799	0.105	0.070	1.000	0.131
CNSS+CoreTree+NoCalib	0.025	0.275	0.046	0.036	0.398	0.066	0.053	0.592	0.097	0.069	0.790	0.127	0.086	1.000	0.159
CNSS+CoreTree	0.049	0.235	0.052	0.294	0.213	0.130	0.584	0.314	0.350	0.703	0.530	0.577	0.836	0.836	0.829
CNSS+LowerBound	0.021	0.286	0.040	0.031	0.416	0.057	0.057	0.609	0.102	0.218	0.783	0.315	0.422	1.000	0.591
CNSS	0.072	0.300	0.064	0.372	0.200	0.192	0.524	0.394	0.390	0.766	0.631	0.677	0.893	1.000	0.943

Table 9: Chosen α values for each method on WikiVote and CondMat datasets, assuming piecewise constant p-values. Note that the EventTree method, unlike the nonparametric scan approaches, does not optimize over α .

Methods	WikiVote ($q = 10$)		WikiVote ($q = 25$)		WikiVote ($q = 50$)		WikiVote ($q = 75$)		WikiVote ($q = 100$)	
	Mean	SD	Mean	SD	Mean	SD	Mean	SD	Mean	SD
LTSS	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
EventTree	-	-	-	-	-	-	-	-	-	-
ColorCoding	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
NPHGS	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
AdditiveScan	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
DFGS	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
CNSS+NoCalib	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
CNSS+CoreTree+NoCalib	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
CNSS+CoreTree	0.040	0.038	0.014	0.010	0.014	0.012	0.011	0.006	0.010	0.001
CNSS+LowerBound	0.022	0.006	0.021	0.007	0.014	0.007	0.011	0.003	0.01	0
CNSS	0.040	0.038	0.014	0.010	0.014	0.012	0.011	0.006	0.01	0

Methods	CondMat ($q = 10$)		CondMat ($q = 25$)		CondMat ($q = 50$)		CondMat ($q = 75$)		CondMat ($q = 100$)	
	Mean	SD	Mean	SD	Mean	SD	Mean	SD	Mean	SD
LTSS	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
EventTree	-	-	-	-	-	-	-	-	-	-
ColorCoding	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
NPHGS	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
AdditiveScan	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
DFGS	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
CNSS+NoCalib	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
CNSS+CoreTree+NoCalib	0.09	0	0.09	0	0.09	0	0.09	0	0.09	0
CNSS+CoreTree	0.077	0.025	0.044	0.034	0.019	0.018	0.018	0.017	0.011	0.004
CNSS+LowerBound	0.087	0.006	0.087	0.006	0.076	0.022	0.039	0.036	0.011	0.004
CNSS	0.055	0.037	0.02	0.02	0.021	0.019	0.019	0.017	0.01	0

the temporal transition. For example, a dummy node $v_{i,t}$ is added into the graph if the hashtag i is not mentioned in week t but is mentioned in both week $t - 1$ and week $t + 1$. The temporal edges $(v_{i,t-1}, v_{i,t})$ and $(v_{i,t}, v_{i,t+1})$ are also added into the graph.

Some co-mentioned hashtags may not be relevant to each other. In order to detect a event with relevant hashtags, we remove the edges between any two hashtags in a week if the overlap coefficient ρ_{ij}^t of the edge $(v_{i,t}, v_{j,t})$ is smaller than a constant (i.e., $\rho_{ij}^t < 0.1$). Overlap coefficient is defined as

$$\rho_{ij}^t = \frac{\#\text{co-mention(tag}_i, \text{tag}_j\text{) in week } t}{(\#\text{tag}_i + \#\text{tag}_j - \#\text{co-mention(tag}_i, \text{tag}_j\text{)) in week } t}. \quad (16)$$

The processed graph includes 100,123 nodes and 137,984 edges. The p-value of each node in the graph is computed based on the rank of the expectation-based Poisson (EBP) statistic (Neill 2012) divided by the total number of nodes. For each node, we compute the EBP score as:

$$\text{EBP} = C \log \frac{C}{B} + B - C, \quad (17)$$

if $C > B$, and $\text{EBP} = 0$ otherwise, where C is the observed count (number of mentions of hashtag h in week w) and B is a baseline assuming independence of hashtag counts and time, i.e.,

$$B = \frac{(\#\text{ tweets in week } w)(\#\text{ tweets mentioning hashtag } h)}{\#\text{ total tweets}}. \quad (18)$$

We apply our CNSS method on this processed graph and discover one subgraph that consists of 3,294 nodes. By observing the hashtags, we find that it is a large subgraph connecting multiple events related to the Black Lives Matter social movement. One reason for this seems to be that hashtags related to certain events (such as the police-involved killings of Mike Brown in Ferguson, MO and Eric Garner in New York City) are used by the BLM movement not just at the time those events occurred, but as a rallying cry throughout the temporal duration of the data, perhaps to emphasize that these abuses have persisted throughout time and are all tied to the same underlying phenomena of societal injustice, inequity, and discrimination.

In order to narrow the focus of our detection method to a specific event and validate our algorithm, we decrease the value of α_{\max} from 0.09 to a smaller value 0.008 empirically, in where we could view the choice of α_{\max} as influencing the granularity of a detected event. In the end, we are able to obtain a significant event with much smaller hashtag cluster as shown in Figure 8. To be fair to the competing methods, we also attempt to shrink the granularity of detection by reducing α_{\max} for these methods as well. For LTSS, we shrink α_{\max} from 0.09 to 0.008, since it detects 6,024 nodes that maximize the BJ score with $\alpha_{\max} = 0.09$. With $\alpha_{\max} = 0.008$, LTSS detects 676 nodes that spread over 54 weeks of data. Therefore, we shrink the α_{\max} to 0.001 for LTSS and still detect 88 nodes that cross over 52 weeks and do not represent any single, specific event. For EventTree, the granularity is controlled by the normalization coefficient λ . We set

$\lambda \in \{0.001, \dots, 0.009, 0.01, \dots, 0.09\}$. The smallest detected subgraph or event has 130 nodes with $\lambda = 0.001$ that spread over 40 weeks, again failing to identify a single event that is localized in time.

As shown in Figure 8, the hashtags detected by CNSS correspond to the widespread protests related to two closely occurring events: the grand juries’ decisions not to indict the police officers responsible for the deaths of Eric Garner and Mike Brown. On July 17, 2014, Eric Garner died in the New York City borough of Staten Island after Daniel Pantaleo, a New York Police Department (NYPD) officer, put him in a prohibited chokehold while arresting him. On August 9, 2014, Mike Brown, an 18-year-old Black man, was fatally shot by a white Ferguson police officer in the city of Ferguson, Missouri. On November 24, 2014, the St. Louis County grand jury decided not to indict the police officer, and on December 4, 2014, a Richmond County grand jury decided not to indict Pantaleo. These decisions stirred massive public protests and rallies in Ferguson, New York City, and Seattle in the following weeks. As we can see in Figure 8, our detected subgraph clearly captures the emergence, the peak, and the end of this event using hashtags of tweets, while the lower volumes of hashtag mentions from continued references after these events are not included in the detected subgraph.

In contrast, as shown in Table 10 and Figure 9, EventTree and LTSS detect multiple small events indicated by multiple peaks in the Figure 9 across long periods. We annotate four peaks for EventTree. The first event is corresponding to the same event detected by CNSS in Figure 8. The second peak is around the Martin Luther King Jr. Day on January 19, 2015. The third event is about the death of Freddie Gray in Baltimore on April 19, 2015. The fourth event is about the death of Sandra Bland in Texas on July 13, 2015. We also annotate five peaks for LTSS. The first event is corresponding to the same event detected by CNSS in Figure 8. The second peak contains multiple unrelated hashtags, such as `#Nigeria`, `#Grammys`, `#Oscars`, and `#ReclaimMLK`. The third peak corresponds to the shooting of Tony Robinson on March 6, 2015, while the fourth and fifth peaks correspond to the deaths of Freddie Gray and Sandra Bland respectively.

These results demonstrate that, even when adjusting the significance threshold α for finer event granularity, the competing methods identify multiple individually anomalous hashtag-week combinations and fail to detect a coherent subgraph corresponding to a single event of interest.

D.2 COVID-19 Case Study

We now show visualizations of the top-1 detected subgraph for CNSS and two competing methods (LTSS and EventTree) in Figures 10, 11, and 12, respectively. We note that the overall spatial-temporal subgraphs identified by CNSS and EventTree are connected, though the set of spatial locations for any given time slice may not be connected. LTSS does not enforce any connectivity constraints.

The connected subgraph identified by CNSS clearly demonstrates the initial progression of the COVID-19 outbreak across the eastern United States between March-June

2020, with initial peaks in New York City and the northeastern U.S. that gradually spread into the southeastern U.S. and Texas. In contrast, the connected subgraph identified by EventTree and the subset identified by LTSS are dispersed over the entire country and do not clearly show the progression of the outbreak’s peak.

We used the following county adjacency data to build the spatial temporal network for the COVID-19 case study: <https://www.census.gov/geographies/reference-files/2010/geo/county-adjacency.html>, and statistics of the top-3 subgraphs detected by CNSS are presented in Table 11.

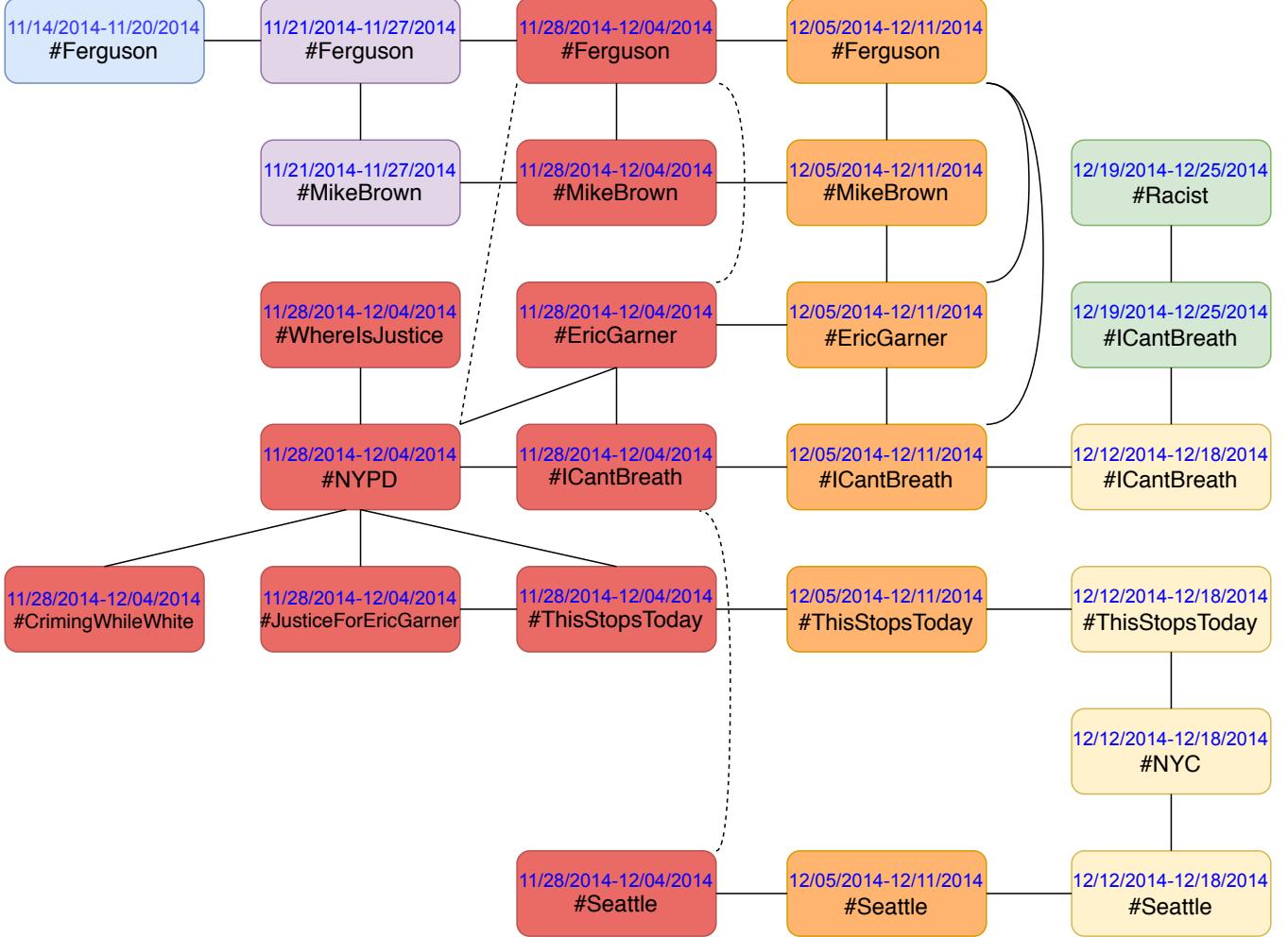


Figure 8: Detected Subgraph in #BlackLivesMatter Tweets. Different colors indicate different weeks. Dashed lines indicate crossing edges.

Table 10: BlackLivesMatter: Statistics of Detected Subgraphs by Different Methods

	EventTree	LTSS	CNSS
periods	11/21/2014-08/20/2015	08/29/2014-08/13/2015	11/14/2014-12/25/2014
# of weeks	40	52	6
# of nodes	130	88	25
# of hashtags	51	73	12

Table 11: COVID-19 Case Study: Statistics of Top-3 CNSS Detected Subgraphs. The calibrated BJ scores of these 3 subgraphs are higher than all 100 calibrated BJ scores under \mathcal{H}_0 .

	N	α	N_α	α'	Calibrated BJ-Score
1st Subgraph	4707	0.09	4702	0.843	774.249
2nd Subgraph	910	0.09	898	0.897	61.187
3rd Subgraph	100	0.03	100	0.708	34.595

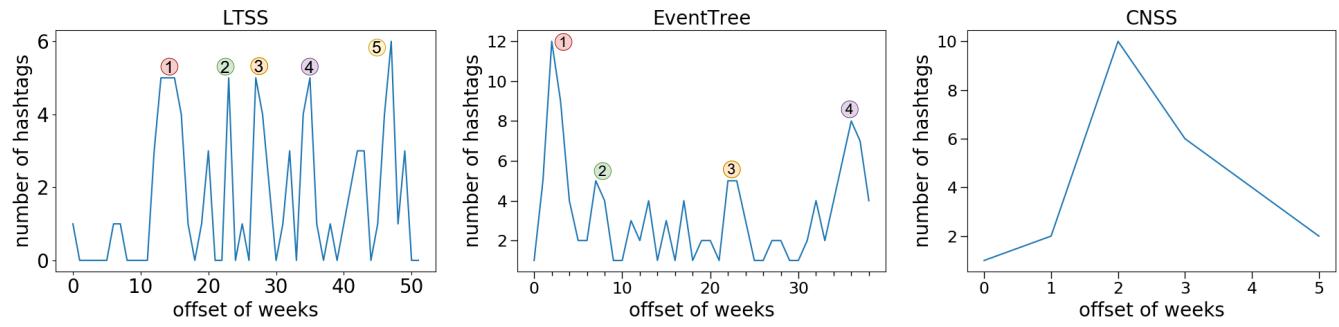


Figure 9: Detected Hashtags Distributions Over Time for LTSS, EventTree, and CNSS.

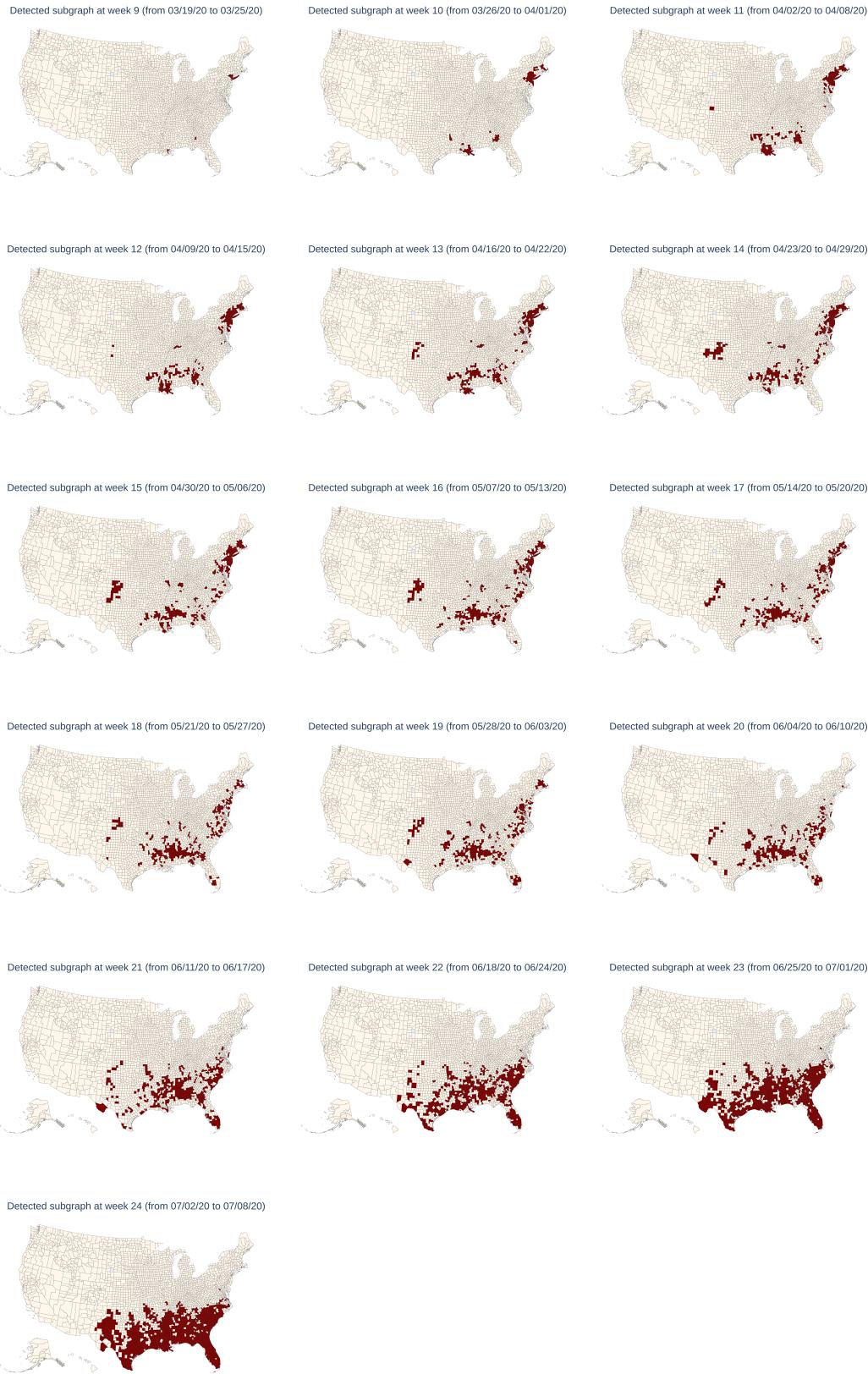


Figure 10: CNSS Top-1 Detected Spatial-Temporal Connected Subgraph on COVID-19 Dataset

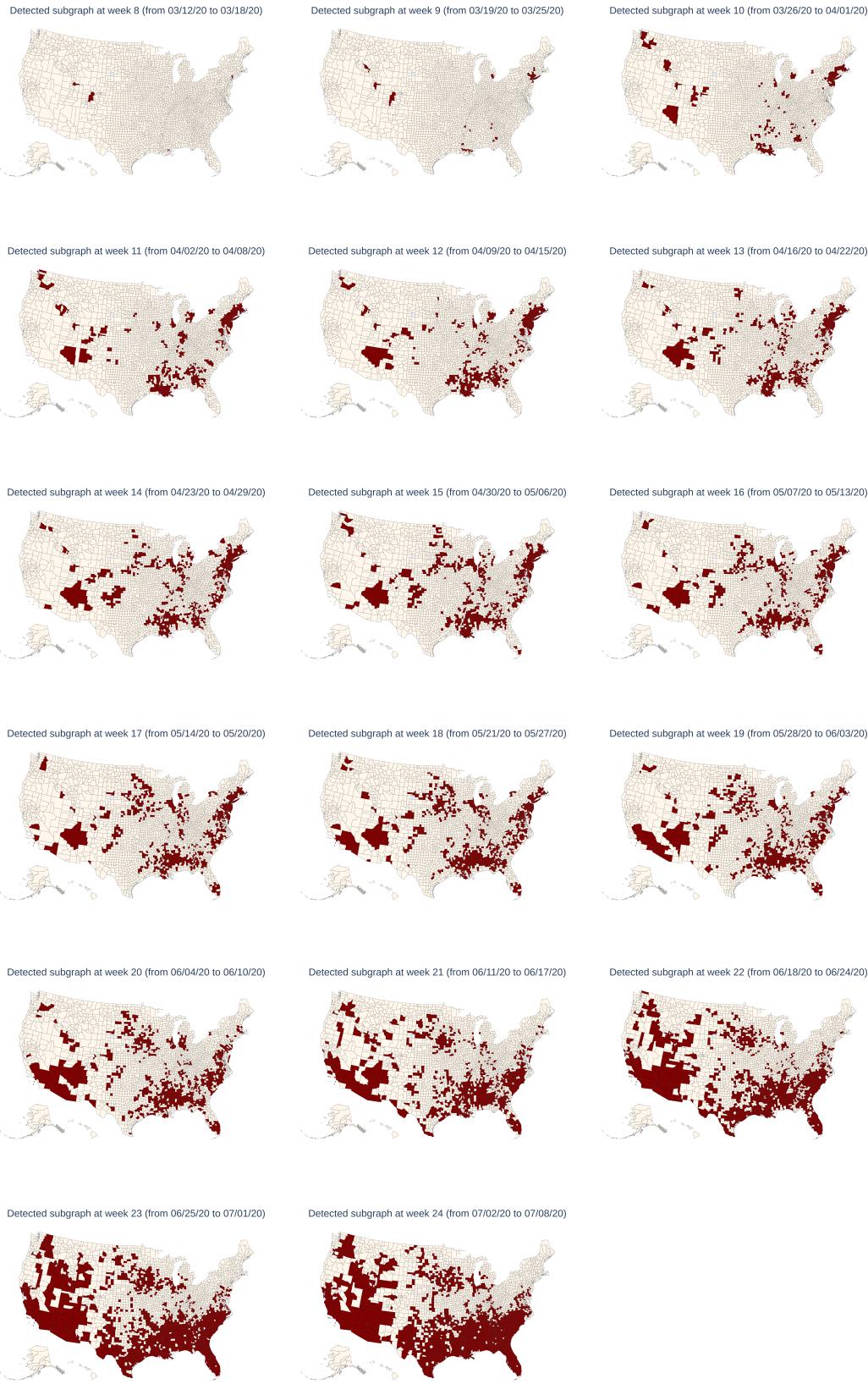


Figure 11: LTSS Top-1 Detected Spatial-Temporal Connected Subgraph on COVID-19 Dataset

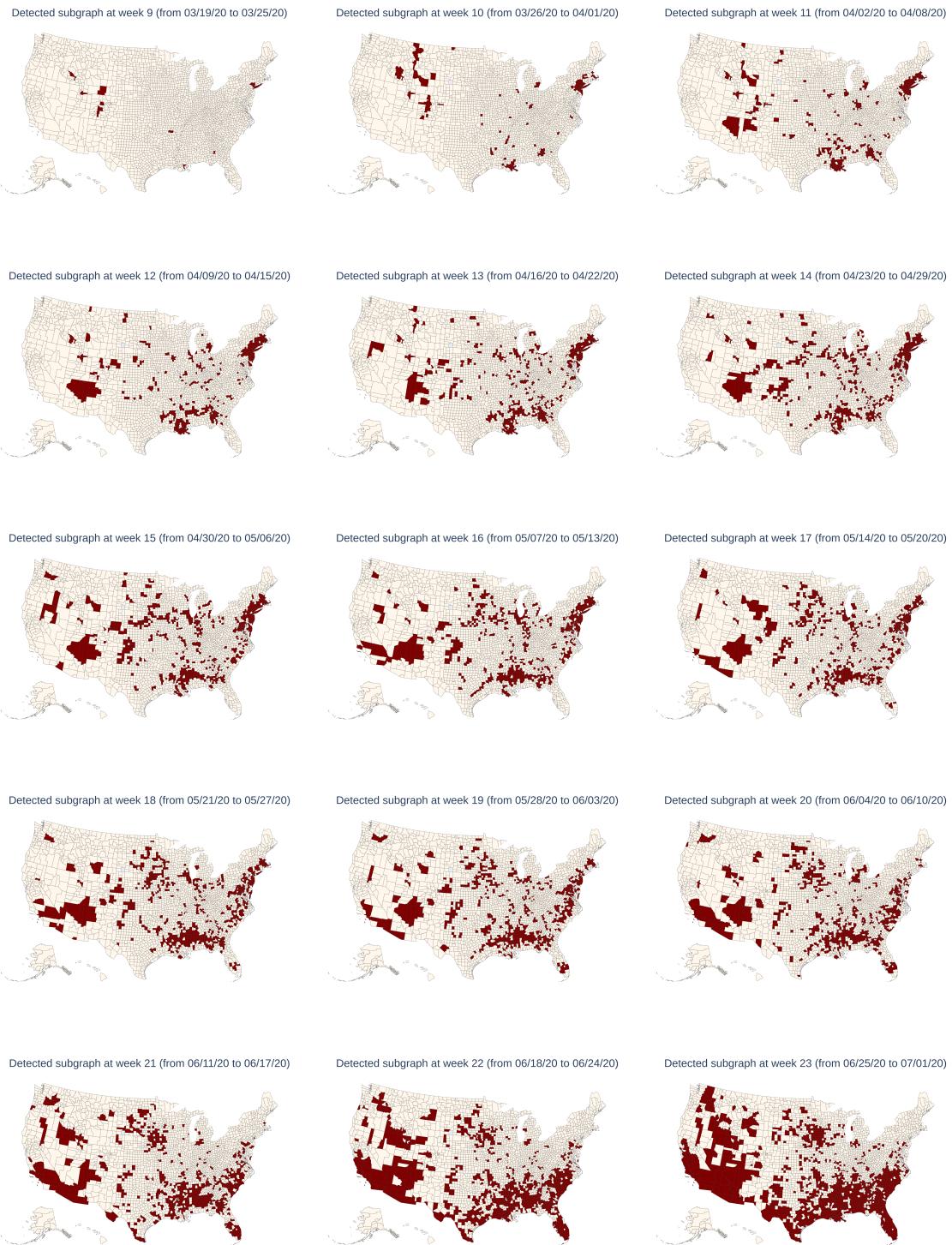


Figure 12: EventTree Top-1 Detected Spatial-Temporal Connected Subgraph on COVID-19 Dataset