

CS 5003: Parameterized Algorithms

Lectures 8-9

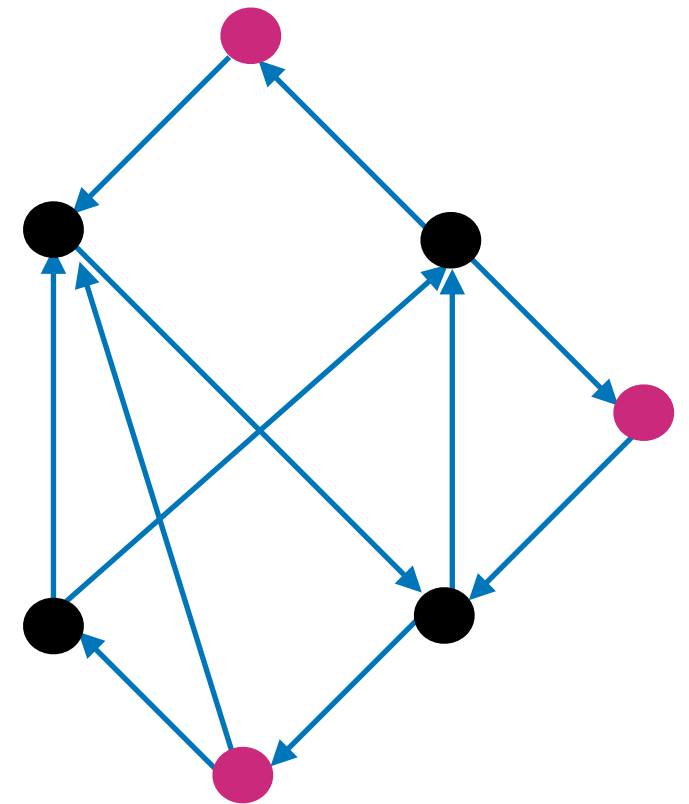
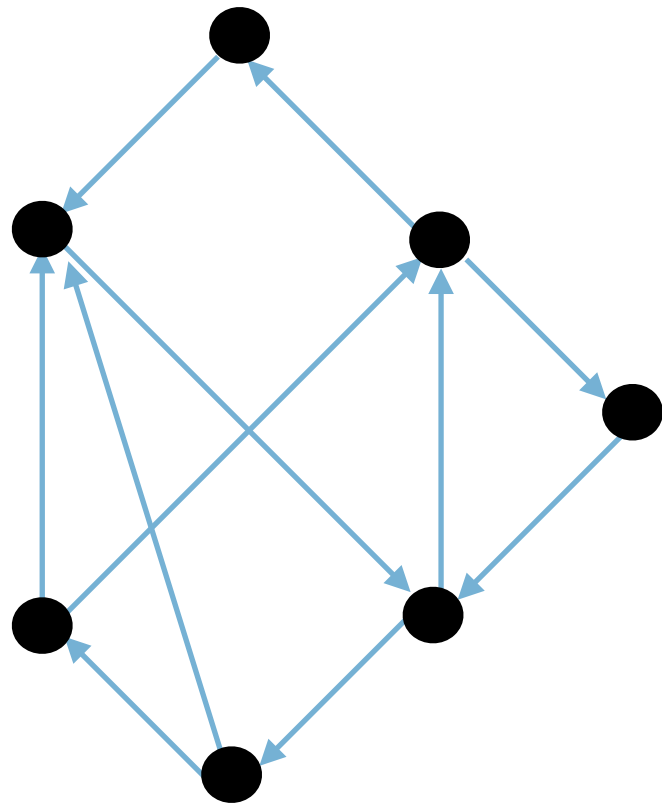
Krithika Ramaswamy

IIT Palakkad

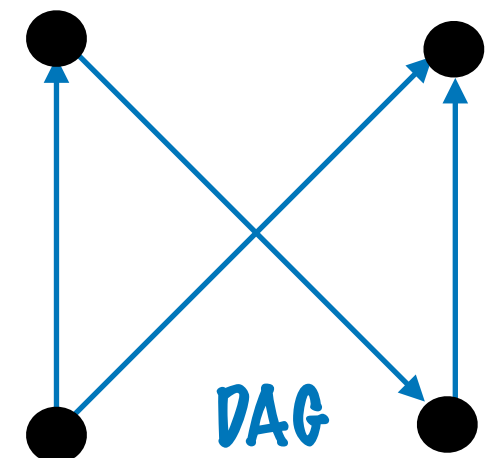
Reference Books: Parameterized Algorithms by Cygan et al. and Kernelization by Fomin et al.

Feedback Vertex Set

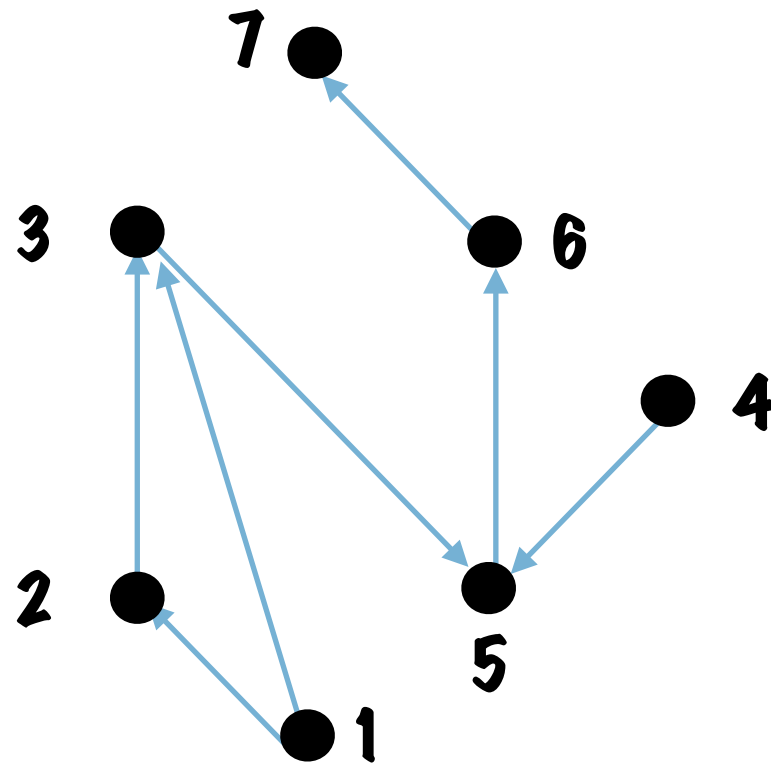
FVS - set of vertices that has at least one vertex of every directed cycle



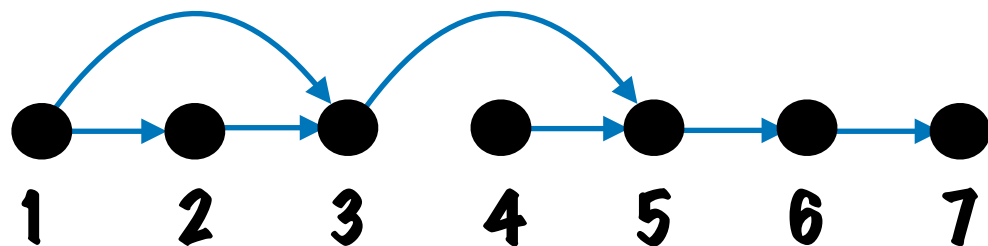
deletion
↓



Feedback Vertex Set

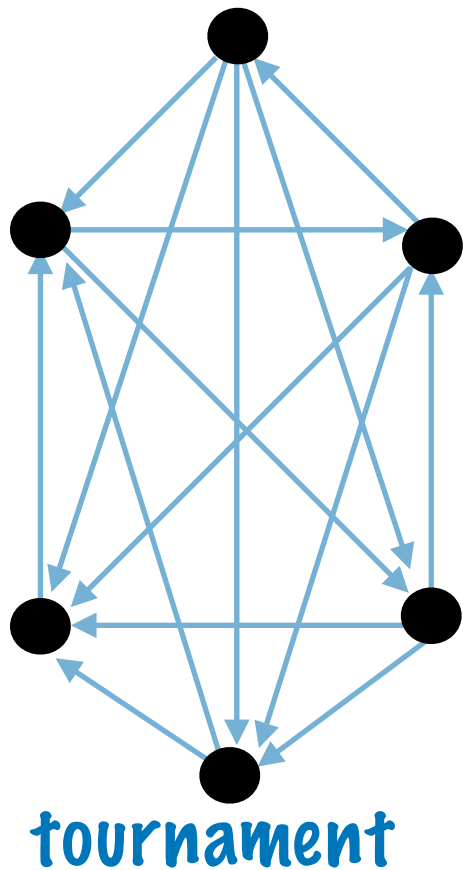


Lemma: A digraph is a DAG iff it has a topological ordering



Topological ordering

Feedback Vertex Set in Tournaments



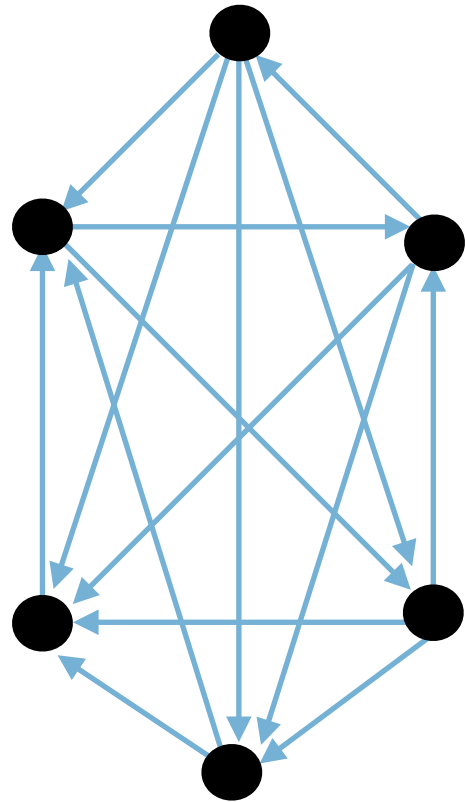
Feedback Vertex Set in Tournaments

Instance: A tournament T and an integer k

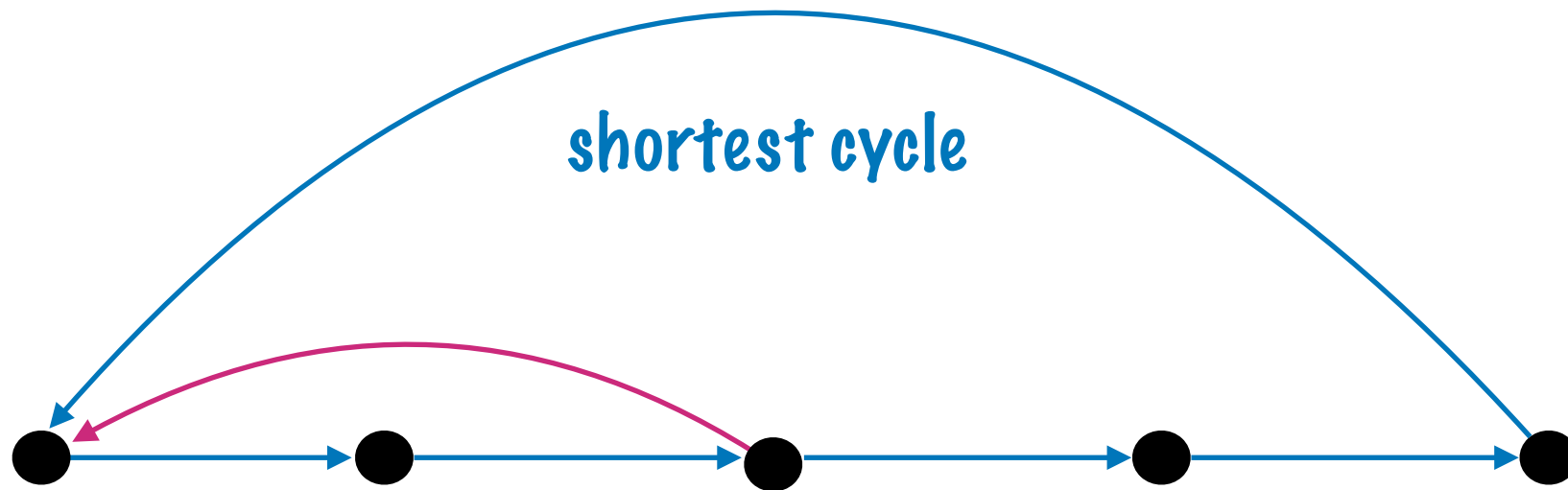
Question: Does there exist a feedback vertex set of T of size at most k ?

Lemma: Acyclic tournaments have unique topological ordering

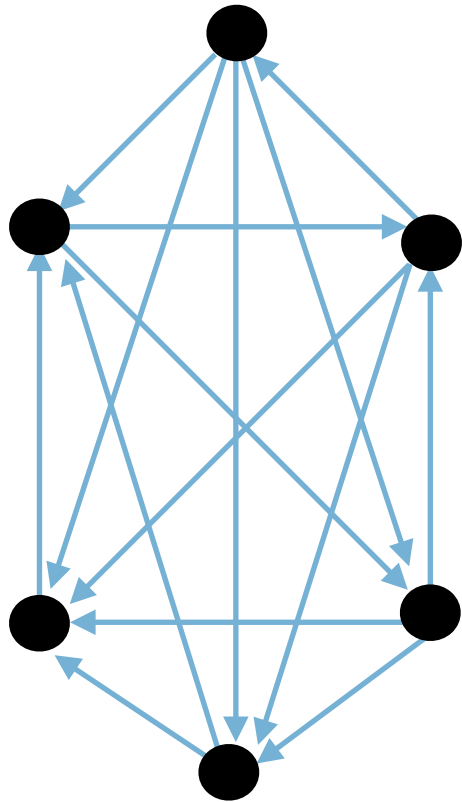
Feedback Vertex Set in Tournaments



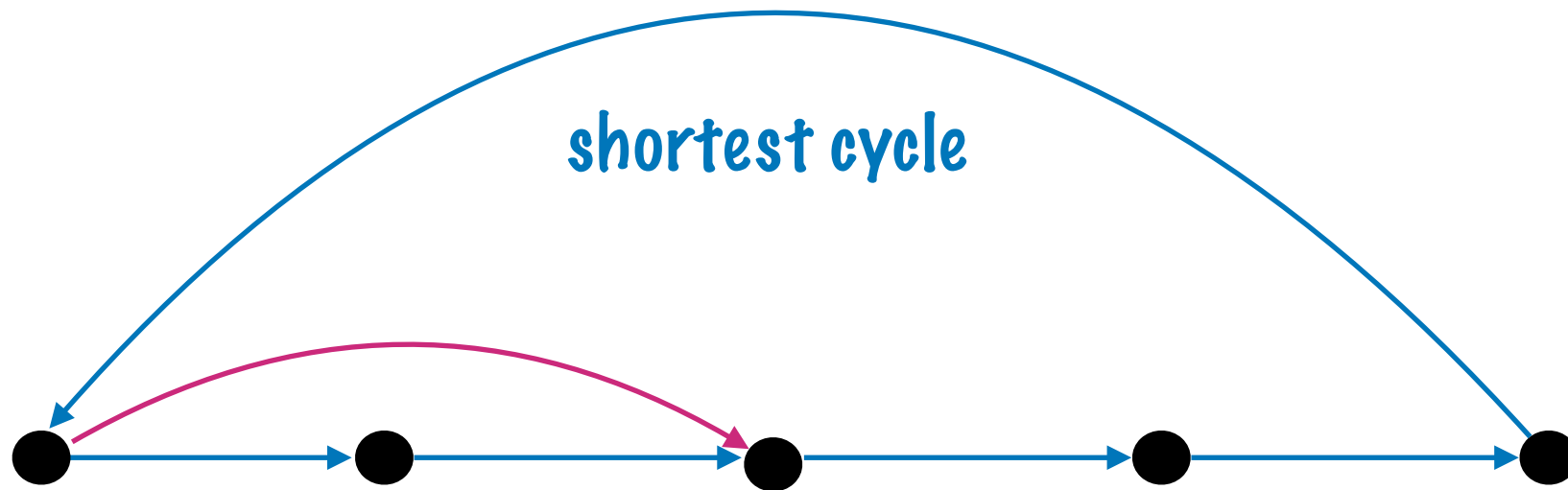
Lemma: A tournament is acyclic iff it has no triangle



Feedback Vertex Set in Tournaments



Lemma: A tournament is acyclic iff it has no triangle



Feedback Vertex Set in Tournaments

Lemma: A tournament is acyclic iff it has no triangle

Branching Algorithm?

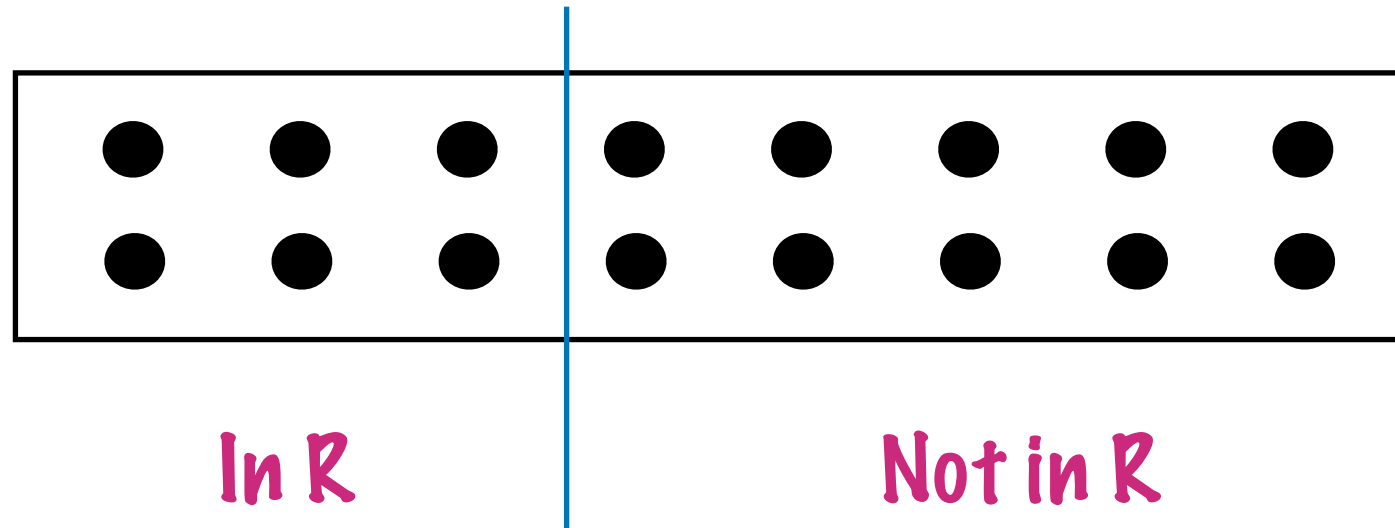
What are the recursive subproblems?

$O^*(3^k)$ algorithm

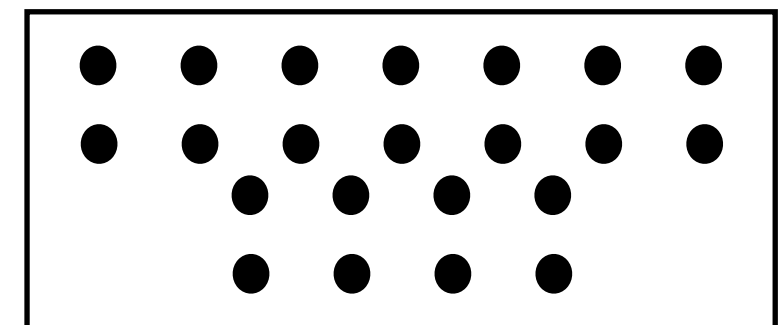
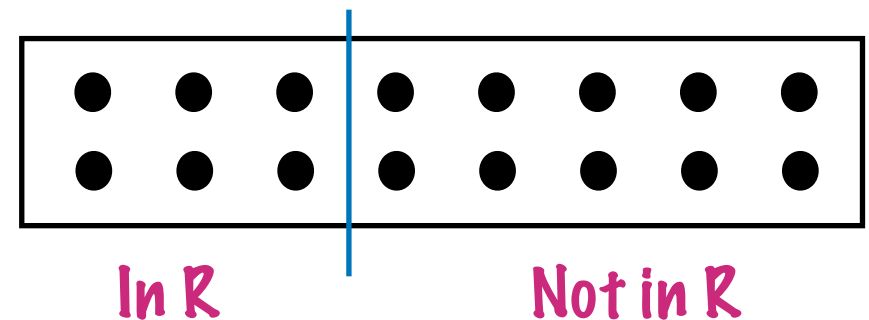
Theorem: FVST is FPT with respect to the solution size as parameter

Feedback Vertex Set in Tournaments

Suppose we have a $(k+1)$ -size solution S

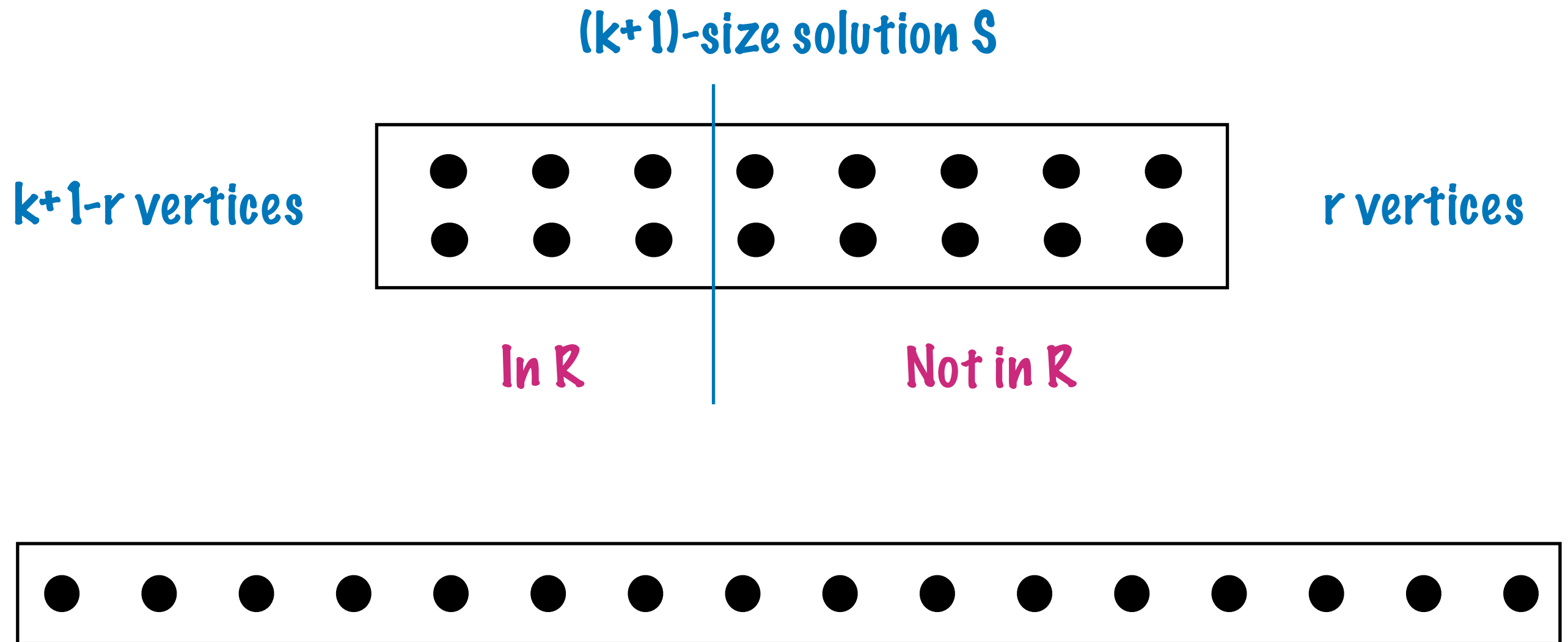


- * We want $\leq k$ size solution R
- * Suppose we know $S \cap R$
 - * If we don't know $S \cap R$, guess!
 - * 2^{k+1} choices



DAG (unique topo ordering)

Feedback Vertex Set in Tournaments

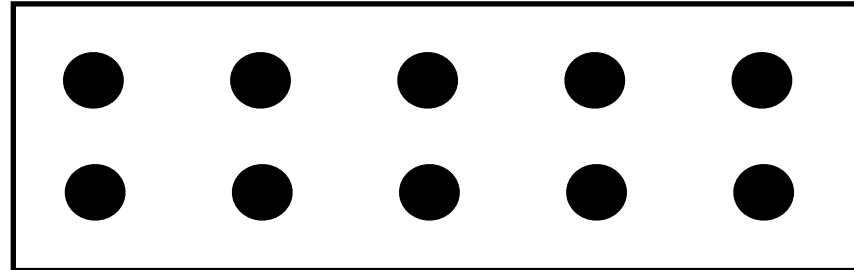


DAG with unique topological ordering

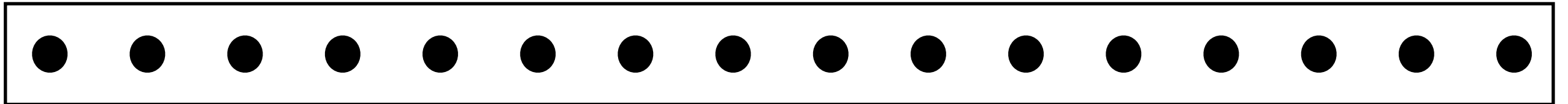
To find a set of $\leq r-1$ vertices here

Feedback Vertex Set in Tournaments

r -size solution



DAG (unique topo ordering)

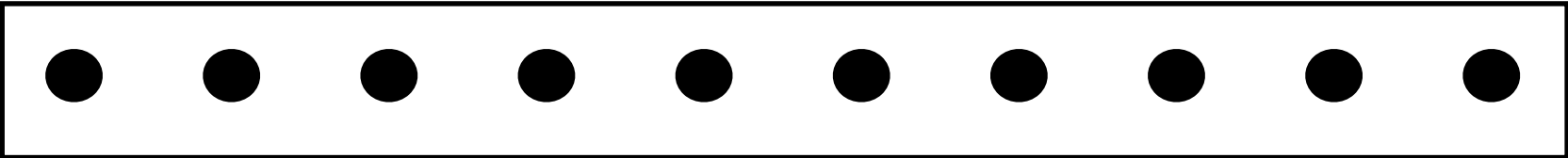


DAG (unique topo ordering)

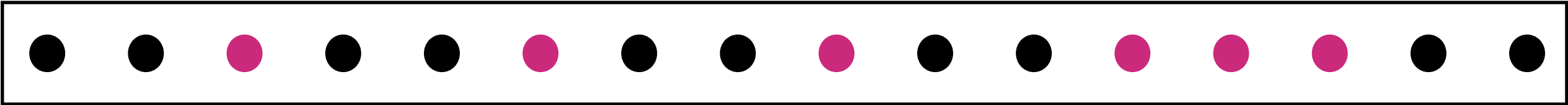
To find a set of $\leq r-1$ vertices here

Feedback Vertex Set in Tournaments

r-size solution



DAG (unique topo ordering)



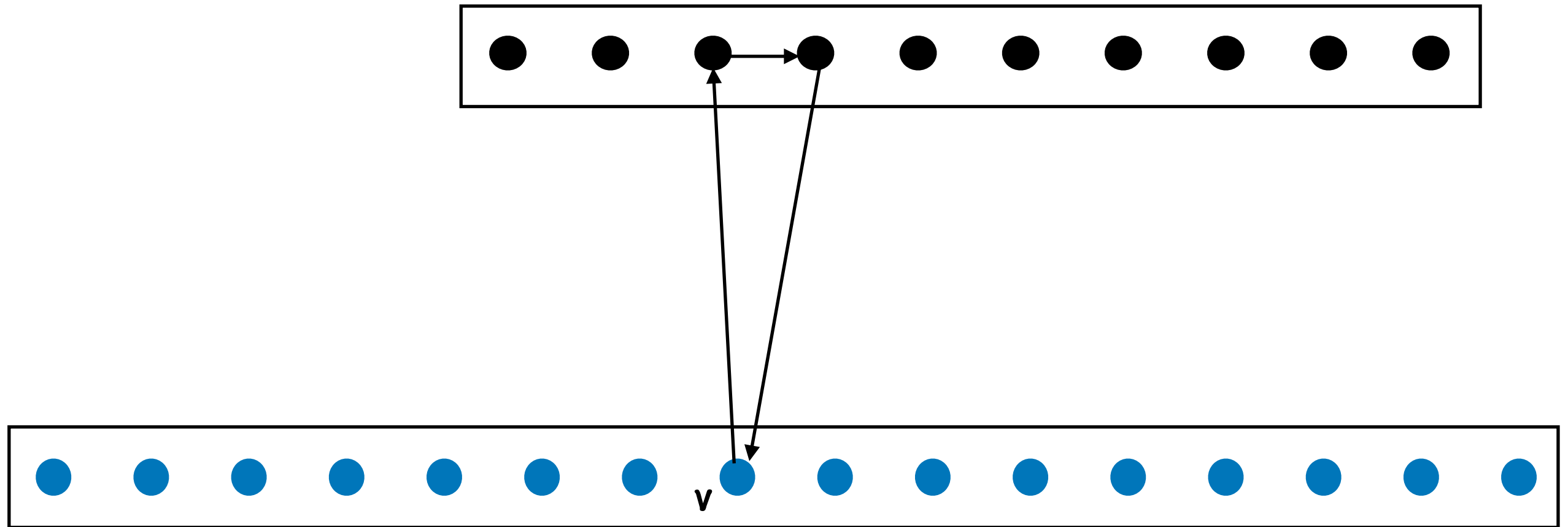
DAG (unique topo ordering)

Find a disjoint (r-1) size solution

Disjoint Compression

Feedback Vertex Set in Tournaments

Disjoint Compression



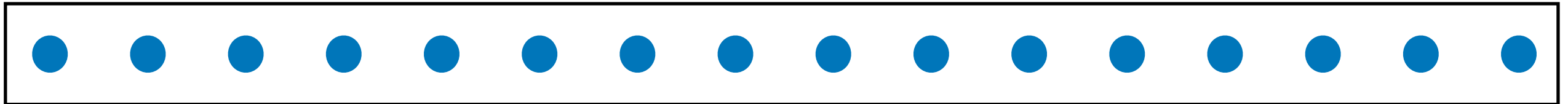
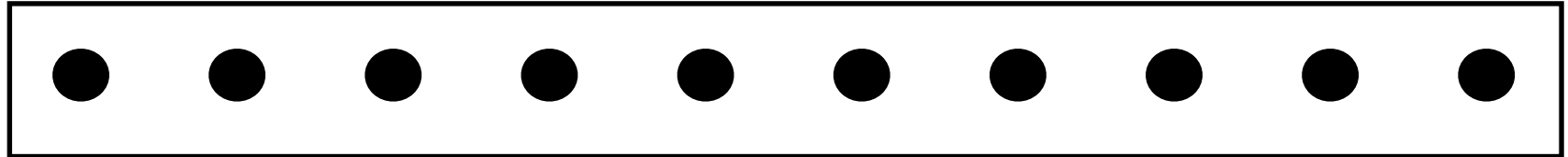
Add v into solution and reduce parameter by 1

Feedback Vertex Set in Tournaments

Disjoint Compression

DAG

X



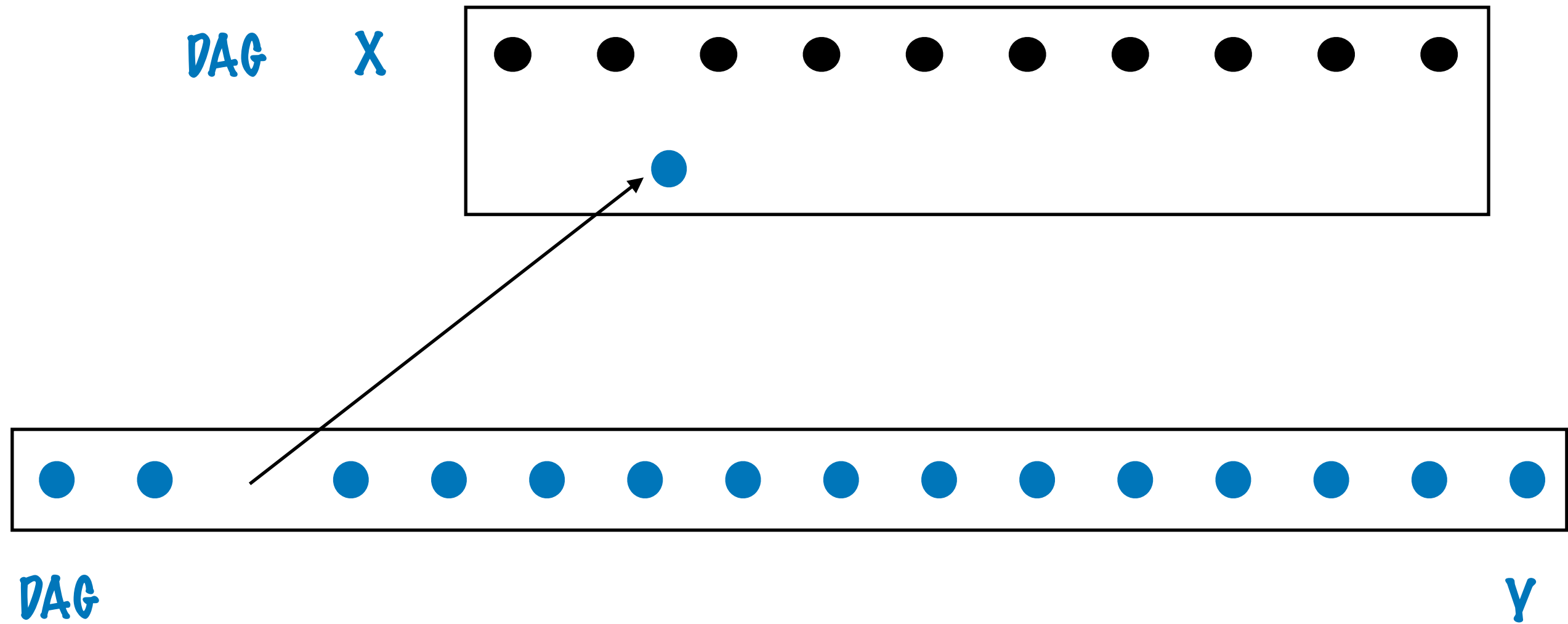
DAG

Y

Any triangle has two vertices from Y

Feedback Vertex Set in Tournaments

Disjoint Compression

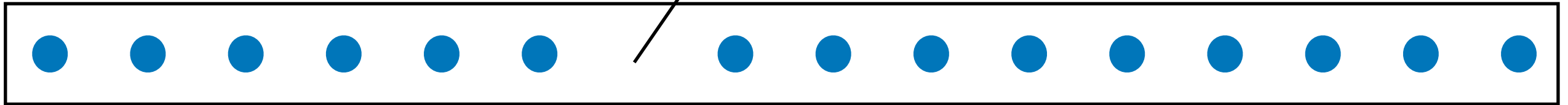
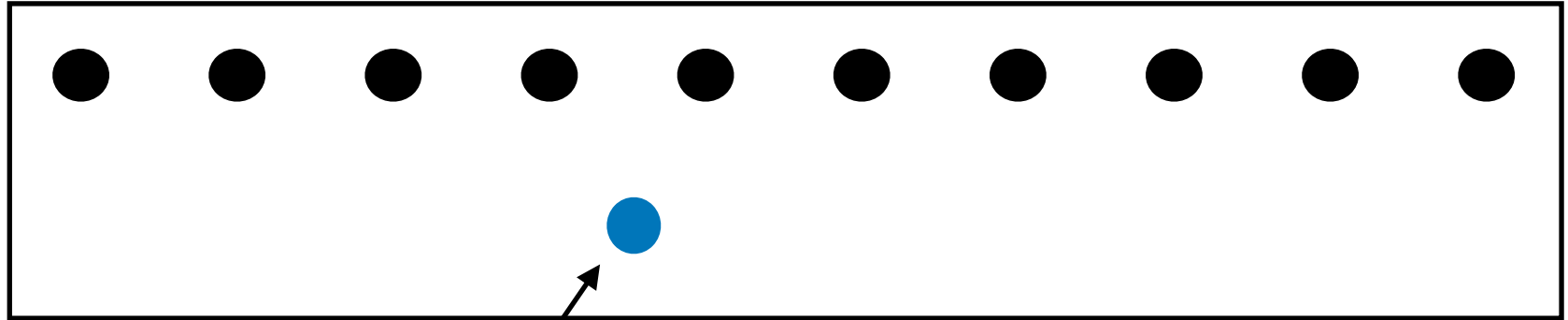


Feedback Vertex Set in Tournaments

Disjoint Compression

DAG

X

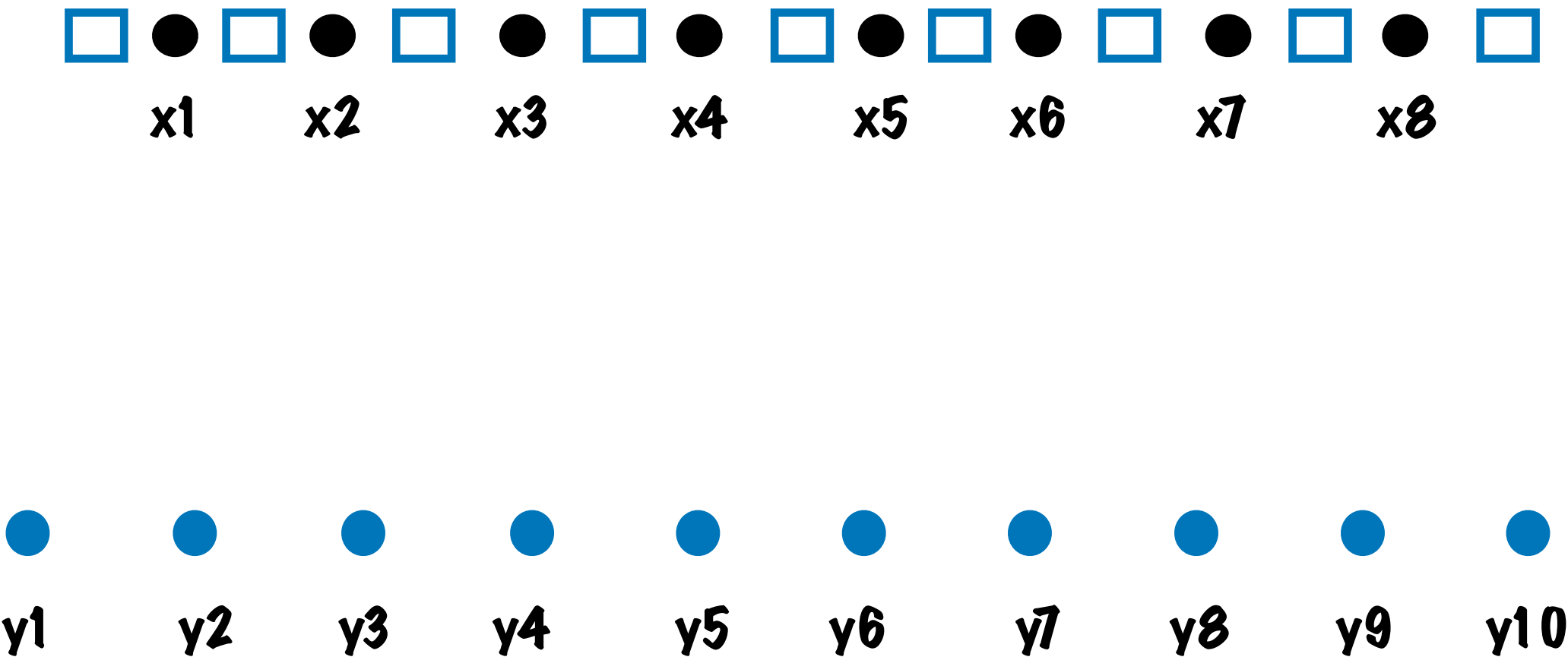


DAG

Y

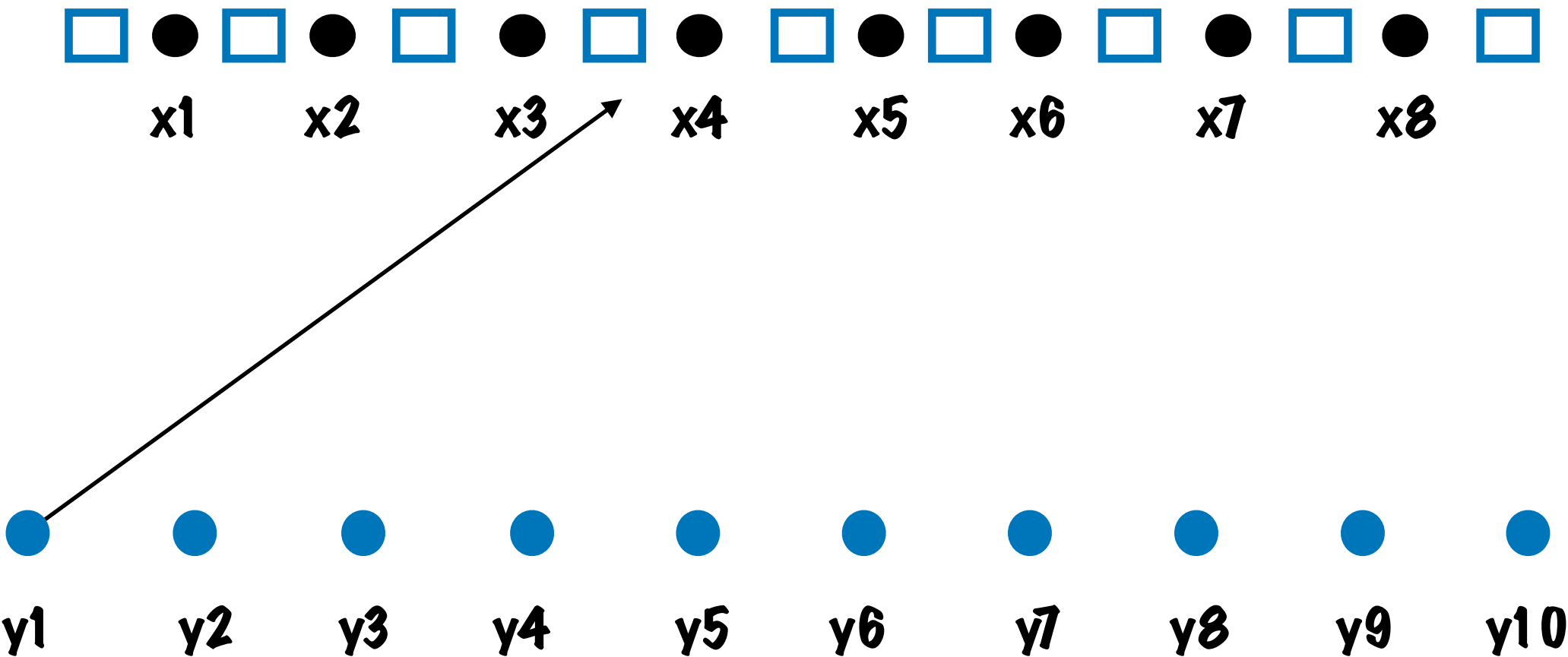
Feedback Vertex Set in Tournaments

Disjoint Compression



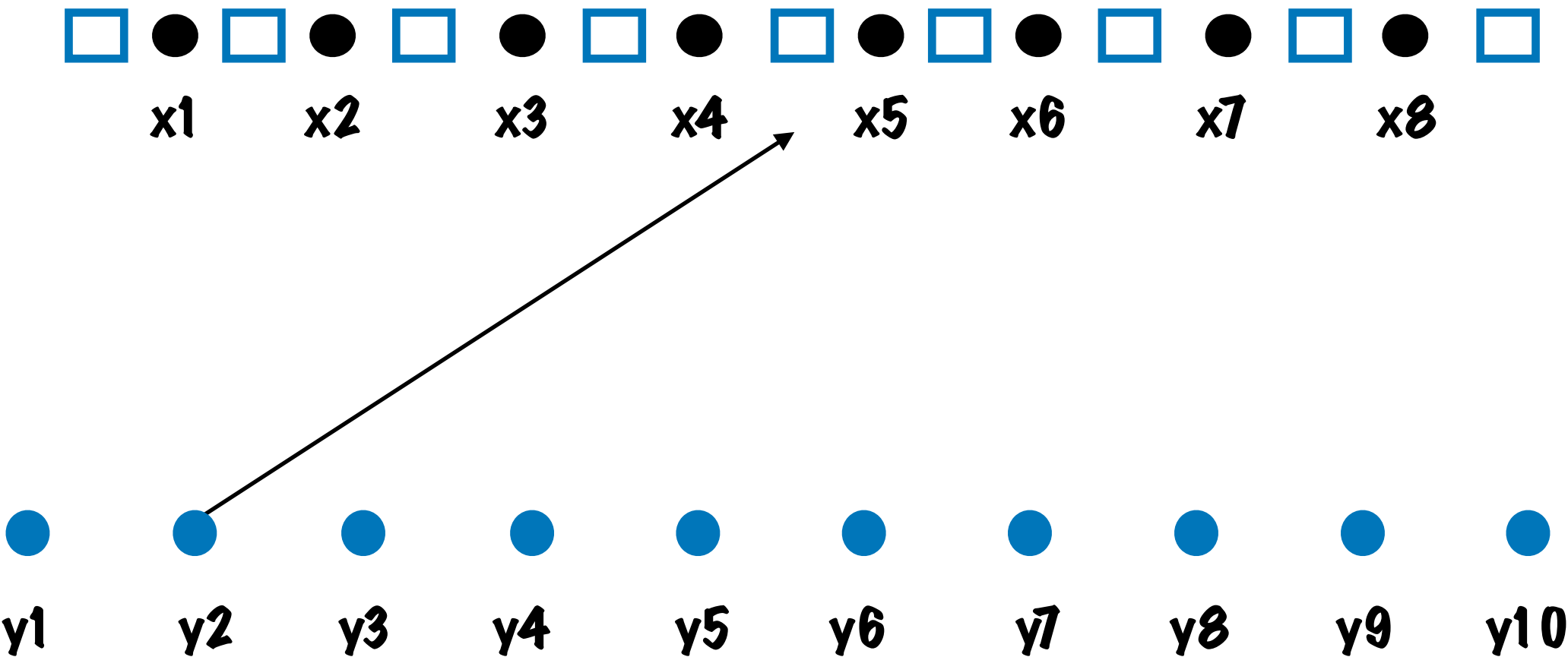
Feedback Vertex Set in Tournaments

Disjoint Compression



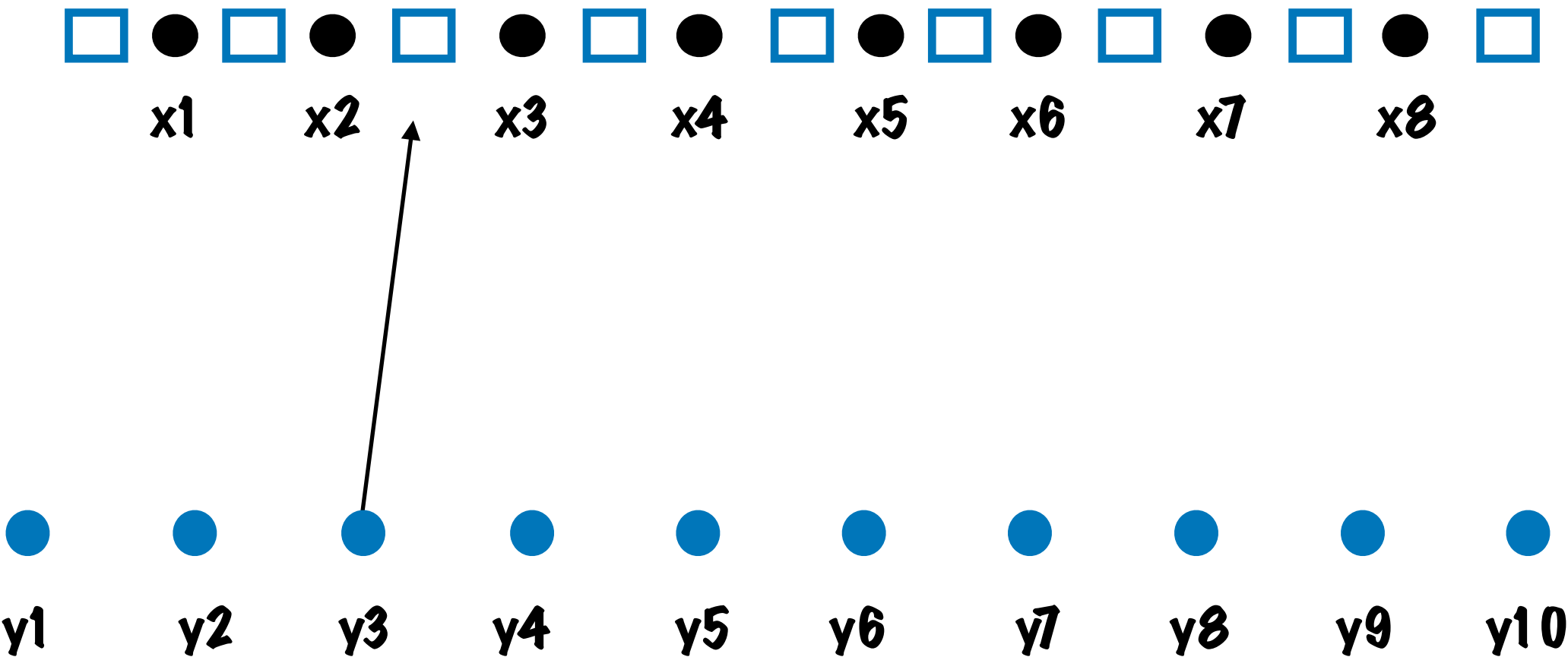
Feedback Vertex Set in Tournaments

Disjoint Compression



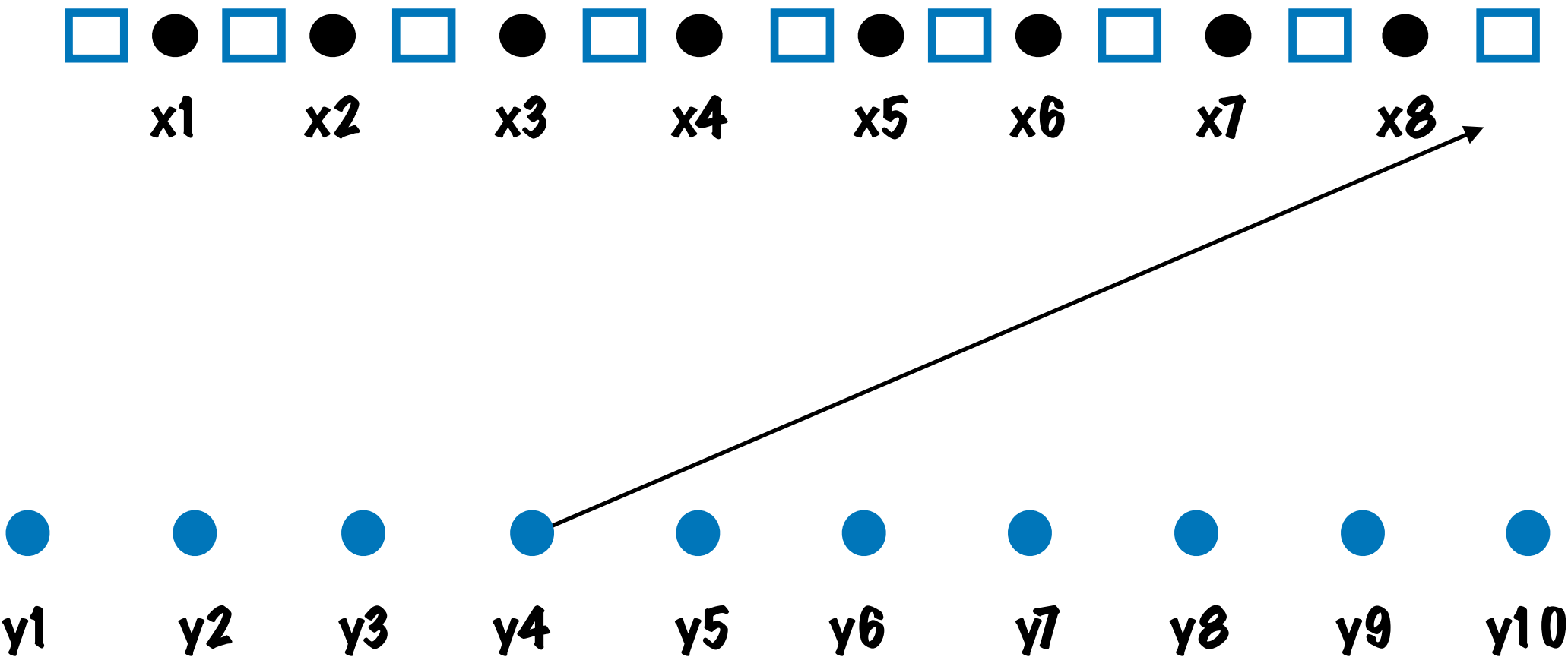
Feedback Vertex Set in Tournaments

Disjoint Compression



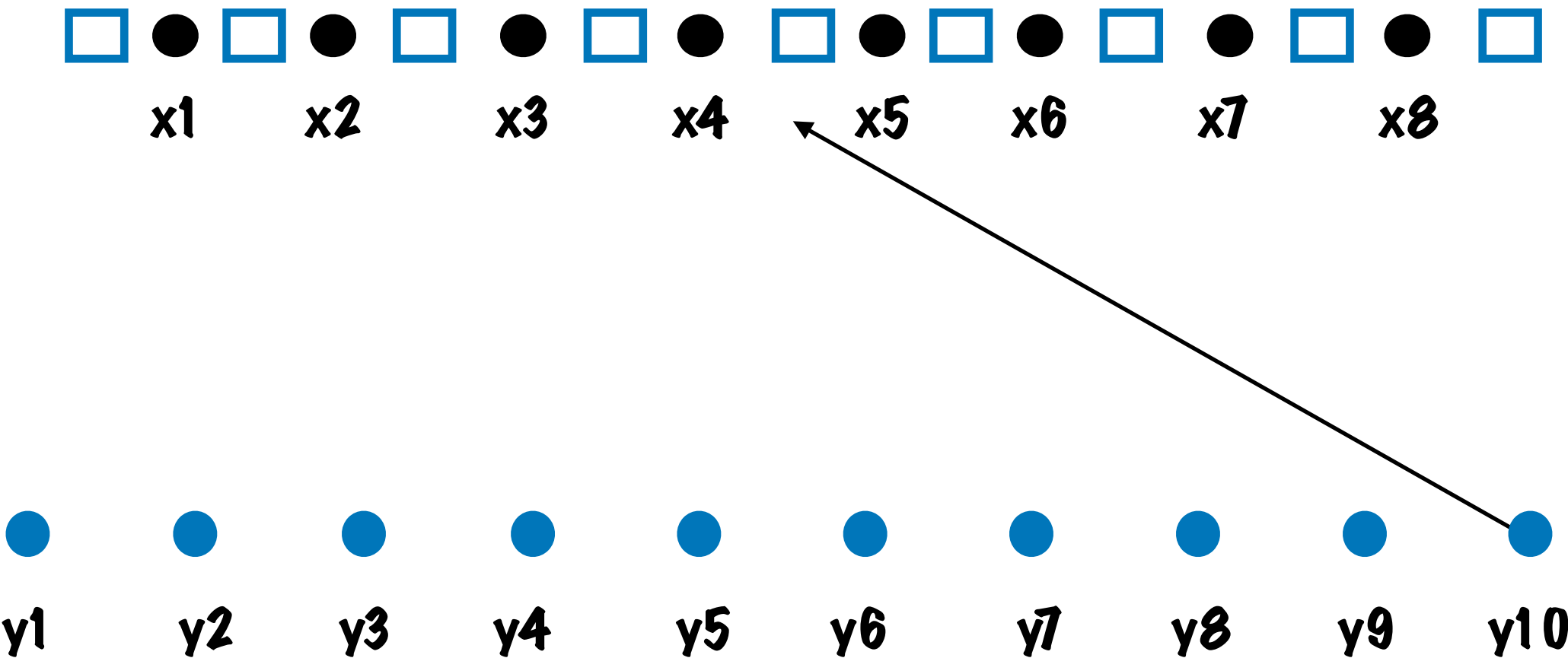
Feedback Vertex Set in Tournaments

Disjoint Compression



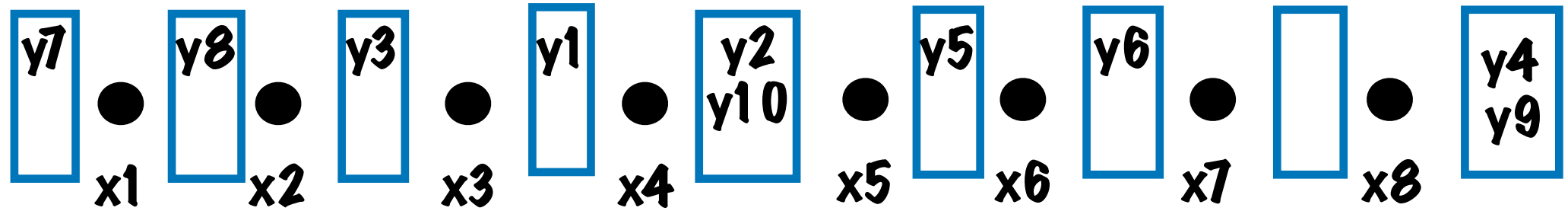
Feedback Vertex Set in Tournaments

Disjoint Compression



Feedback Vertex Set in Tournaments

Disjoint Compression



Feedback Vertex Set in Tournaments

Disjoint Compression



y1



y2



y3



y4



y5



y6



y7



y8



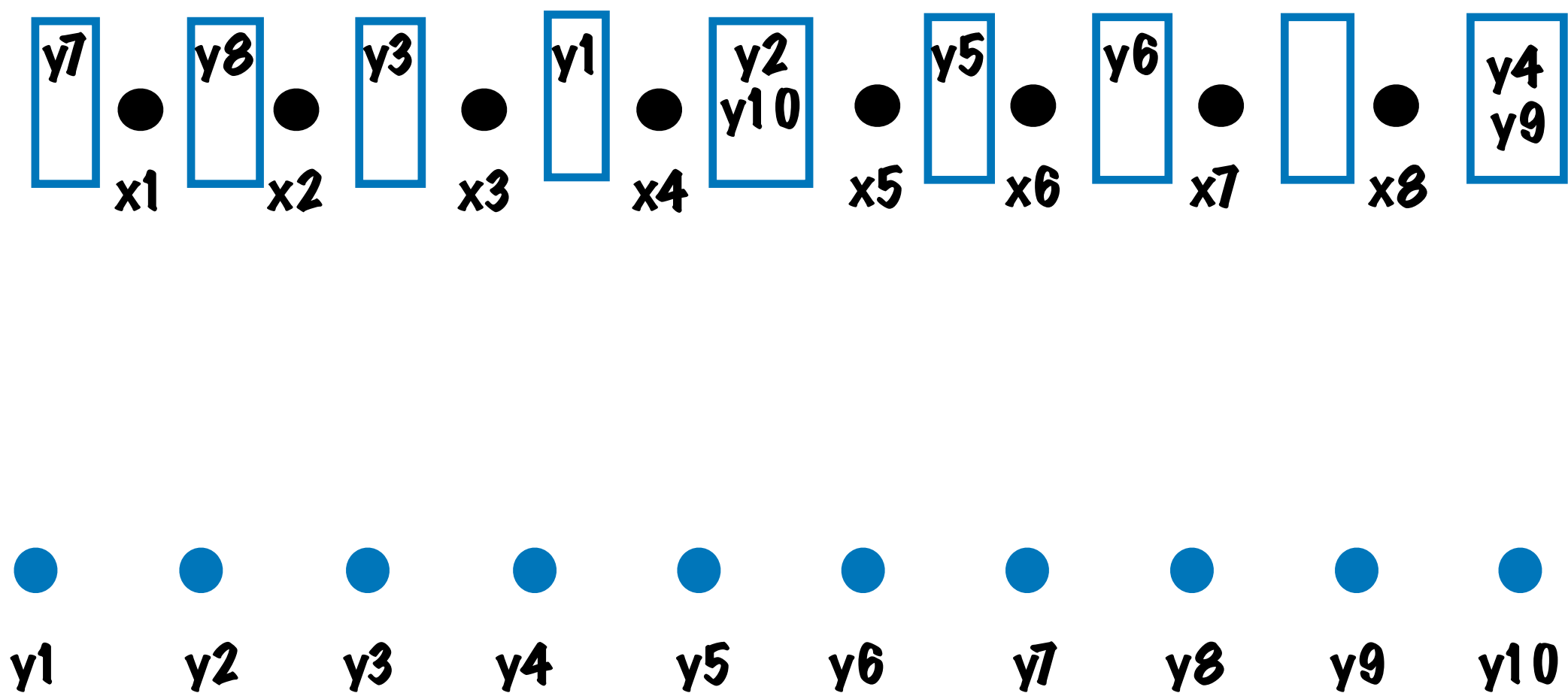
y9



y10

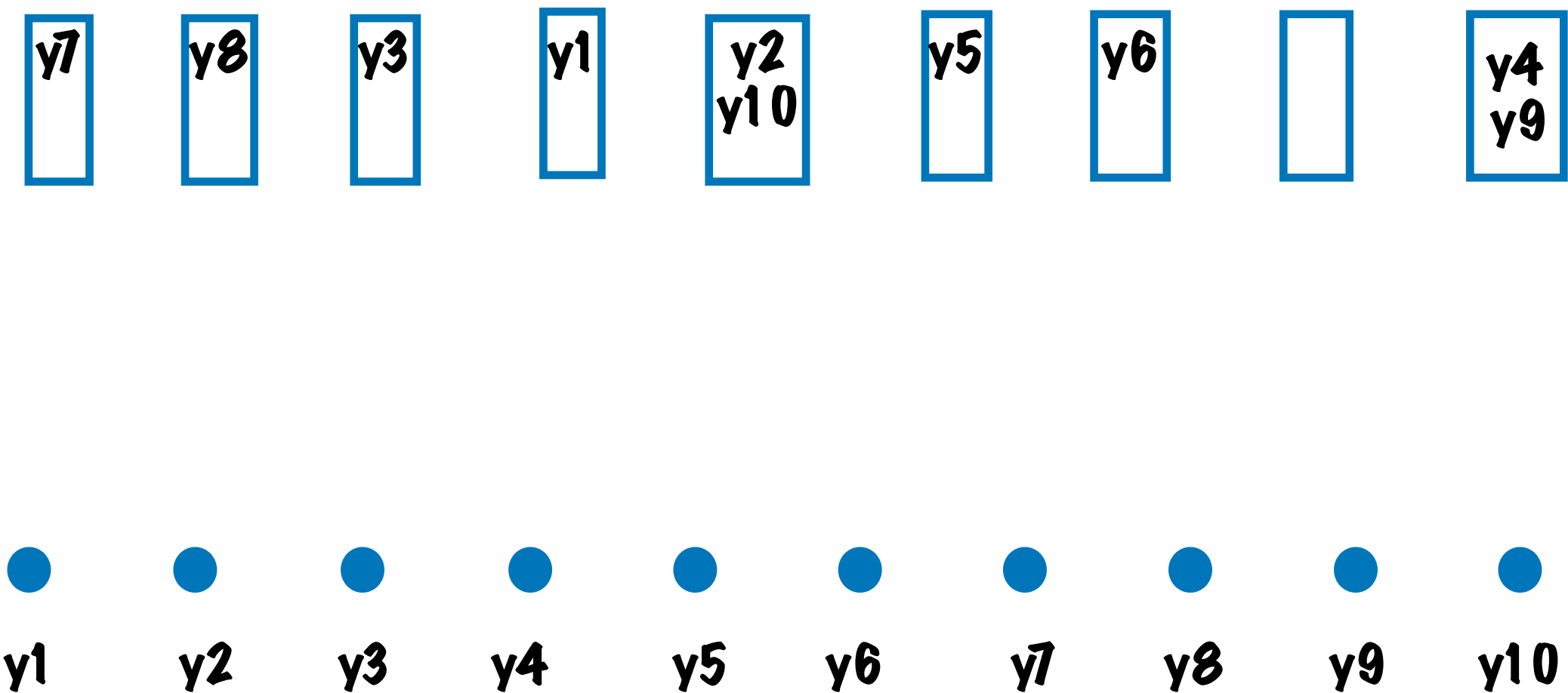
Feedback Vertex Set in Tournaments

Disjoint Compression



Feedback Vertex Set in Tournaments

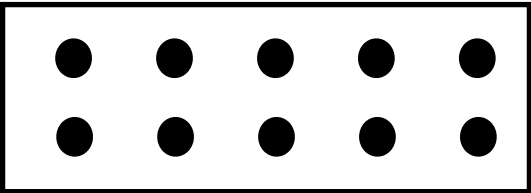
Disjoint Compression



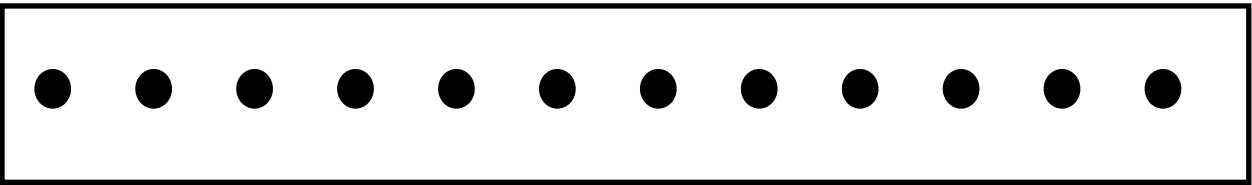
Feedback Vertex Set in Tournaments

Disjoint Compression

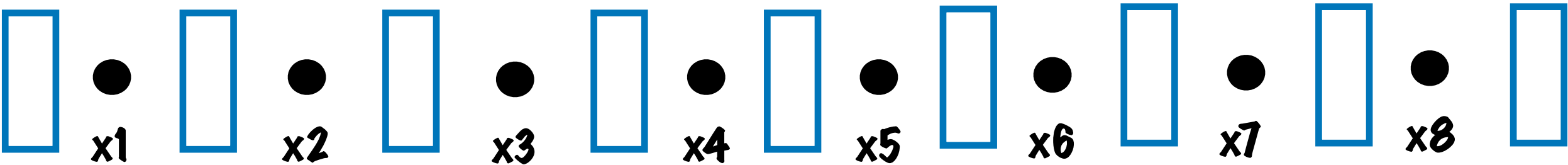
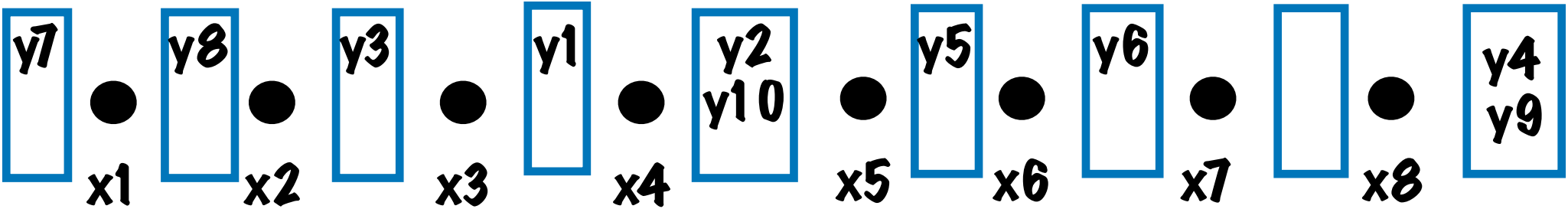
Consider solution DAG's topological order



$\leq k$ solution



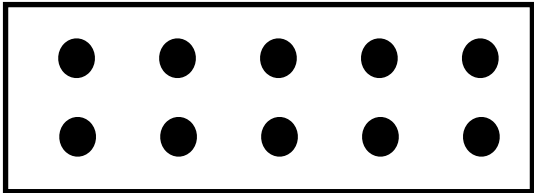
DAG



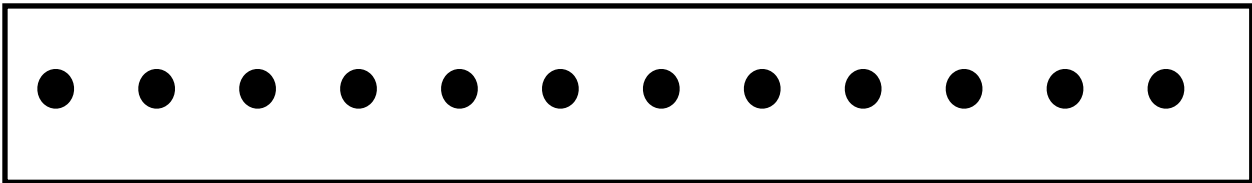
Feedback Vertex Set in Tournaments

Disjoint Compression

Consider solution DAG's topological order



$\leq k$ solution



DAG



y1



y2



y3



y4



y5



y6



y7



y8



y9



y10



x1



x2



x3



x4



x5



x6



x7



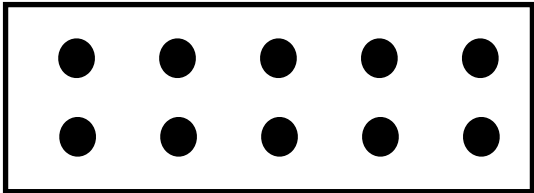
x8



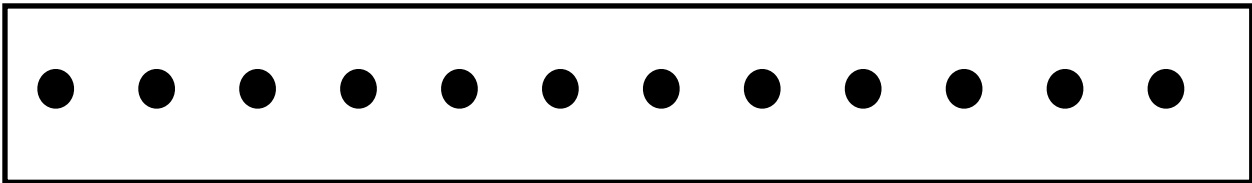
Feedback Vertex Set in Tournaments

Disjoint Compression

Consider solution DAG's topological order



$\leq k$ solution



DAG



y1

y2

y3

y4

y5

y6

y7

y8

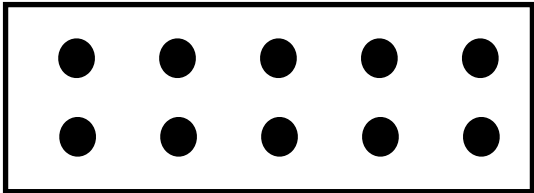
y9

y10

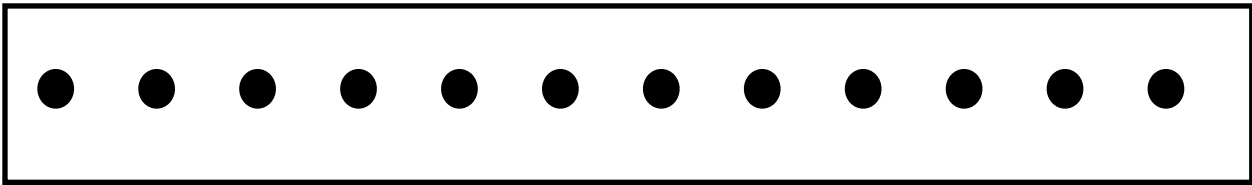
Feedback Vertex Set in Tournaments

Disjoint Compression

Consider solution DAG's topological order



$\leq k$ solution



DAG



y1



y2



y3



y4



y5



y6



y7



y8



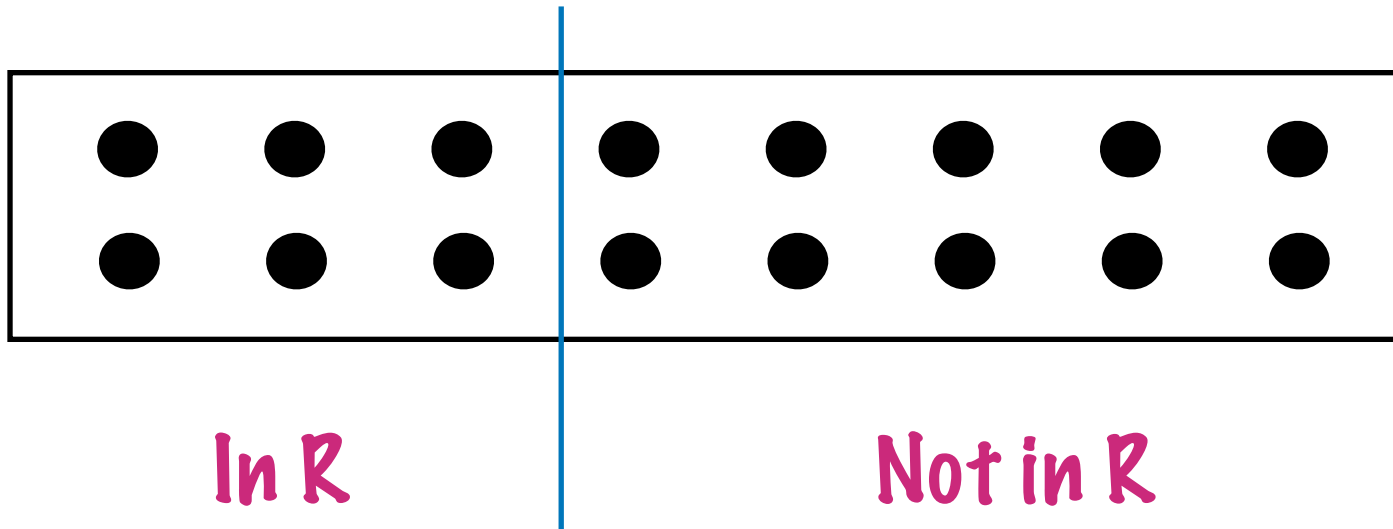
y9



y10

Find longest common subsequence

Feedback Vertex Set in Tournaments



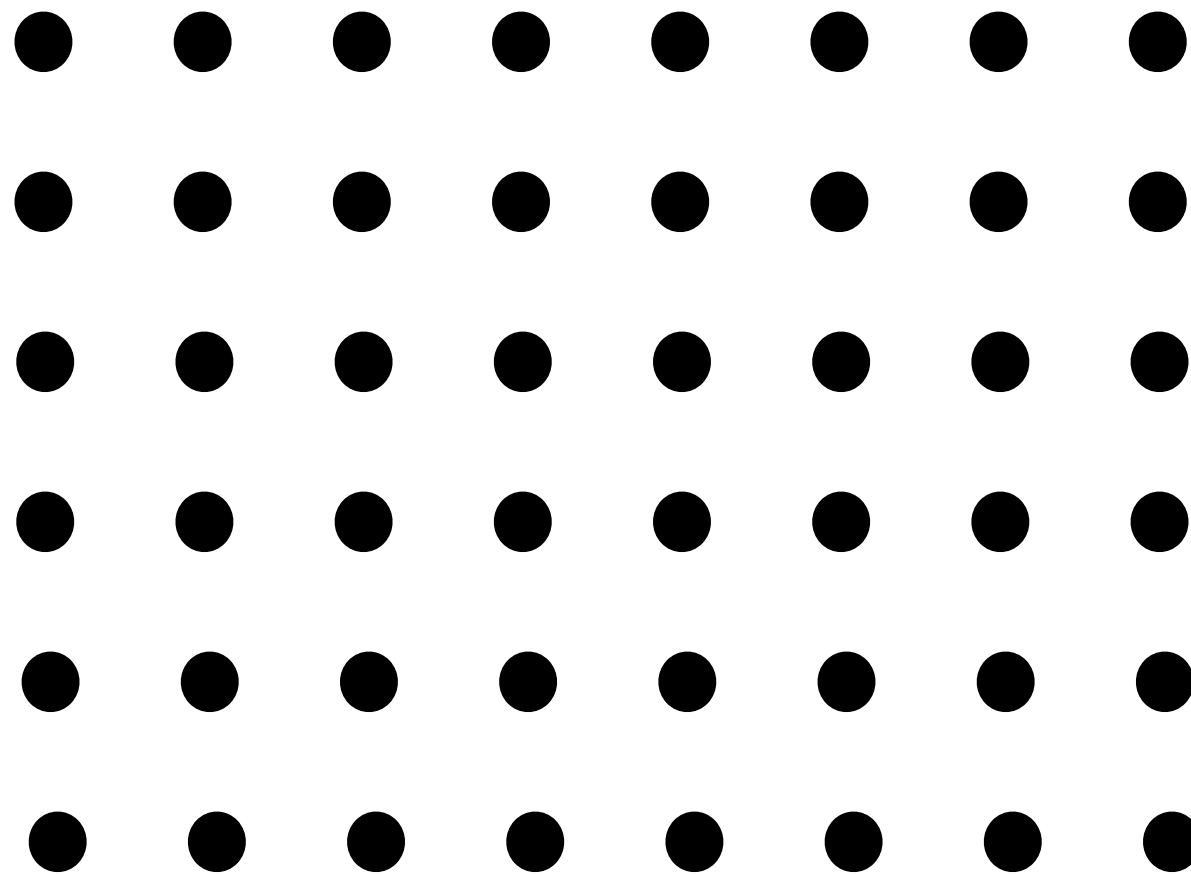
(k+1)-size solution S

- * We want $\leq k$ size solution R
- * Suppose we know $S \cap R$
 - * We don't know $S \cap R$, guess! (2^{k+1} choices)
 - * Solve Disjoint Compression
 - * Longest Common Subsequence (in polynomial time)

$O^*(2^k)$ algorithm

Feedback Vertex Set in Tournaments

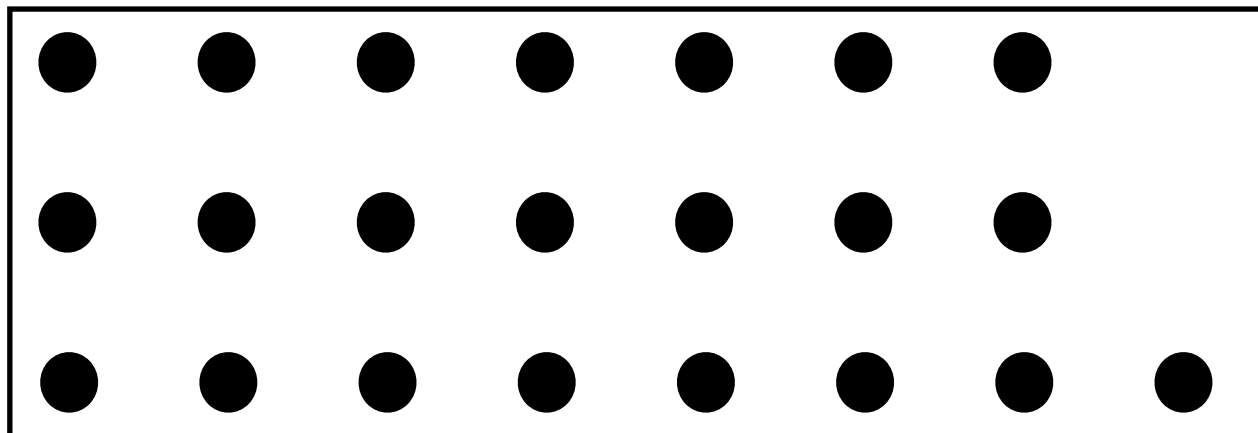
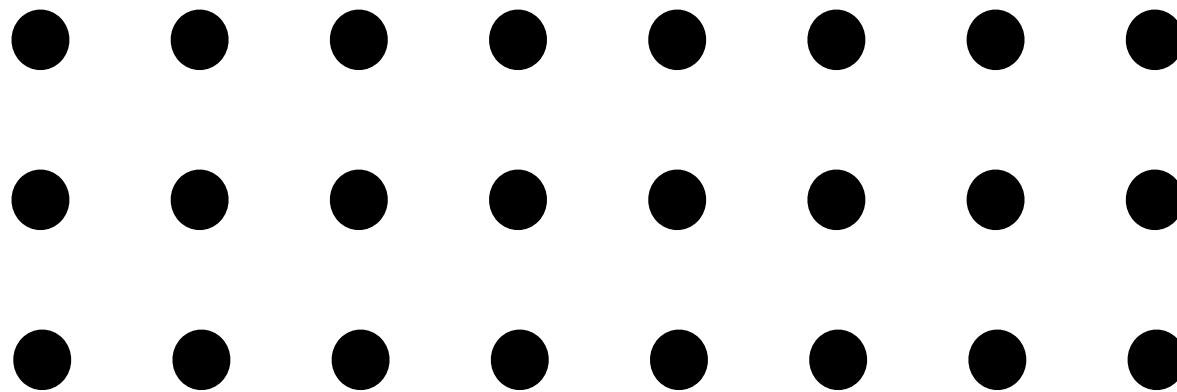
How to get a $(k+1)$ -size solution S ?



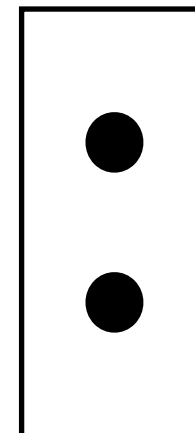
Tournament T

Feedback Vertex Set in Tournaments

Consider any $k+3$ vertices of T



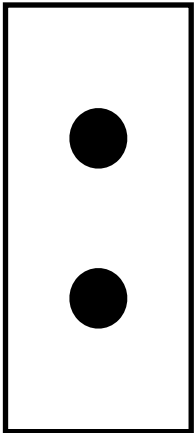
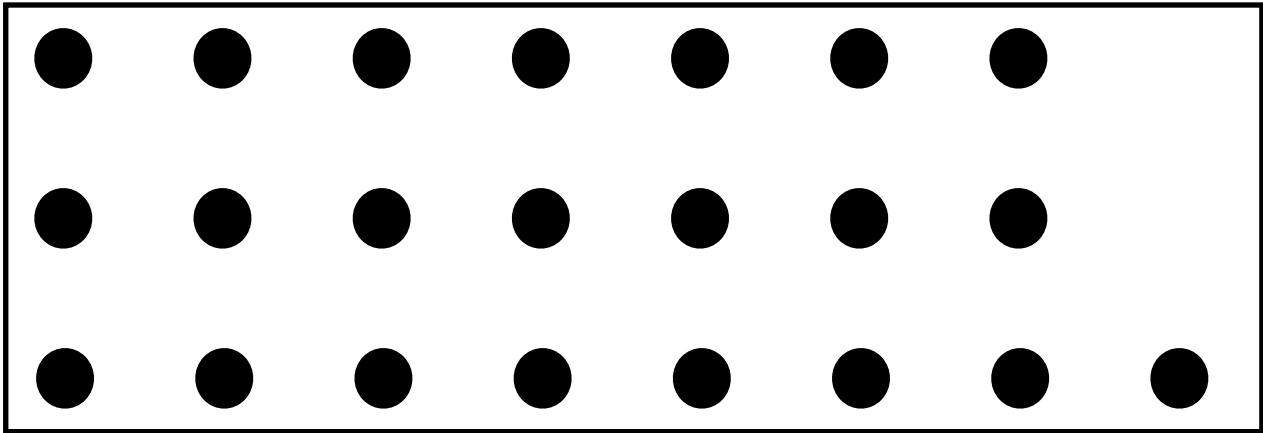
$k+1$ solution



DAG

Feedback Vertex Set in Tournaments

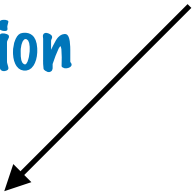
$k+1$ solution for subtournament on $k+3$ vertices



DAG

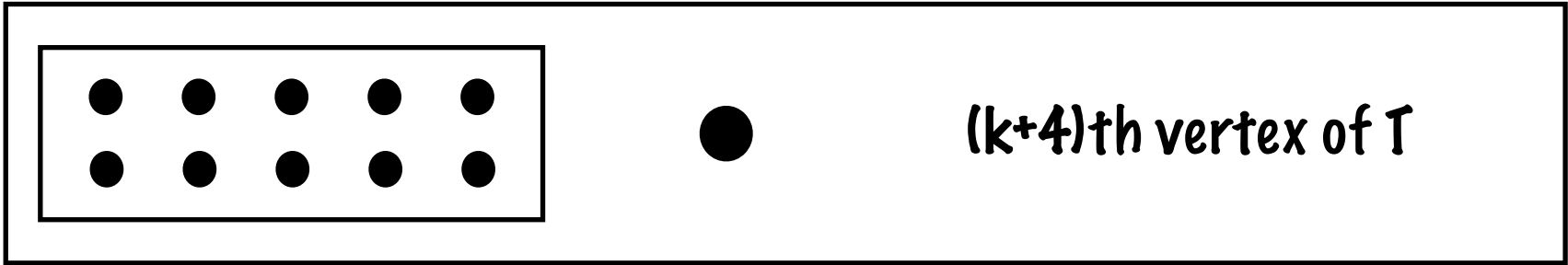
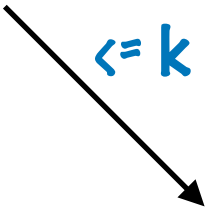
Compress in $O^*(2^k)$ time

No k solution



(T,k) is a no-instance

$\leq k$ solution



$k+1$ solution for subtournament on $k+4$ vertices

Iteratively Compress