

# **CS 5003: Parameterized Algorithms**

**Lectures 26-27**

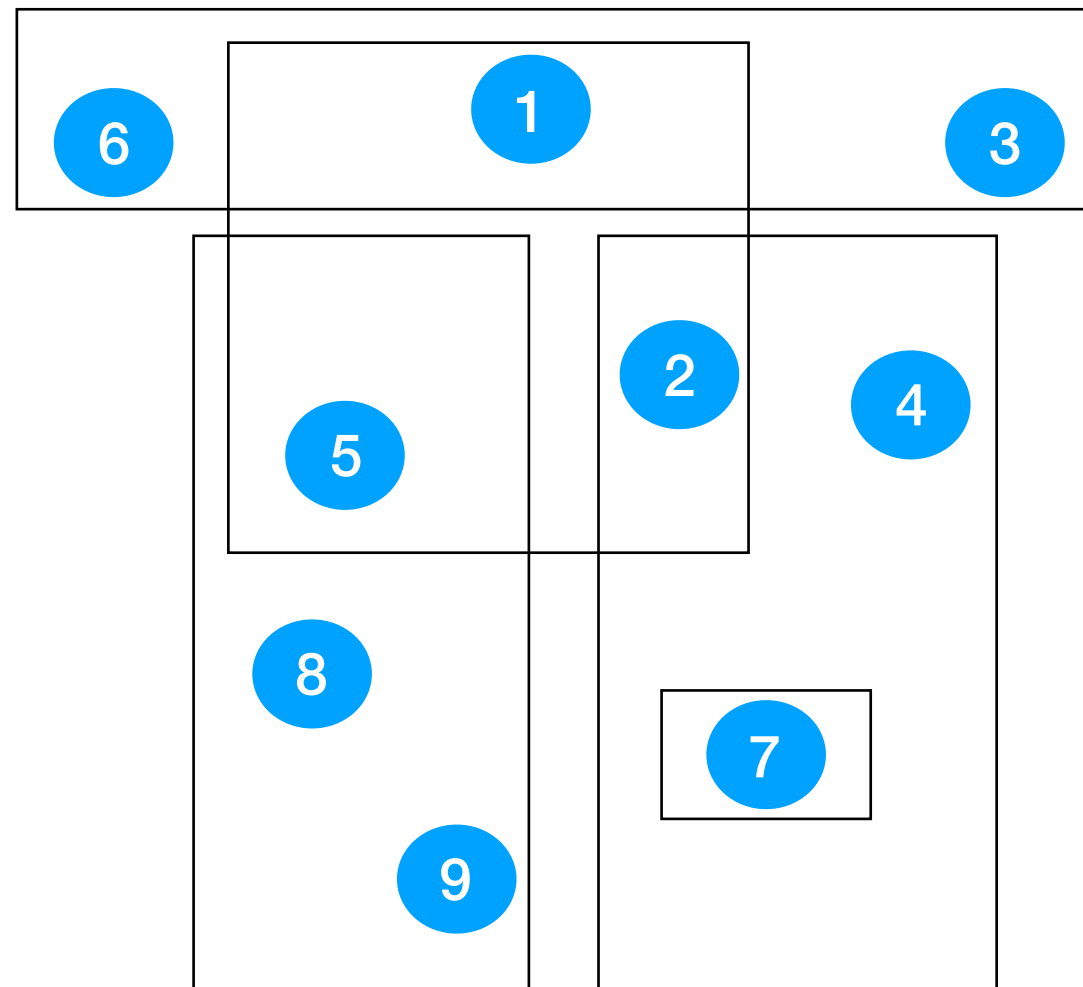
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**Reference Books: Parameterized Algorithms by Cygan et al. and Kernelization by Fomin et al.**

# Sunflower Lemma

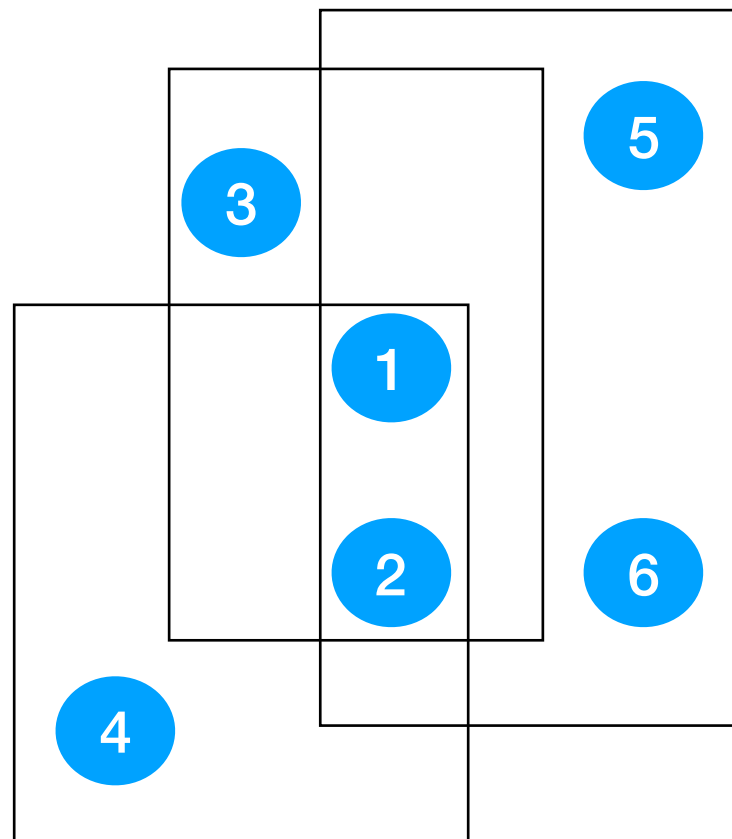
- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family of subsets of  $U$ :  $\mathcal{F} = \{S_1, S_2, \dots, S_m\}$



# Sunflower Lemma

- \* Universe  $U = \{1, 2, 3, \dots, n\}$ , Family of subsets of  $U$ :  $F = \{S_1, S_2, \dots, S_m\}$
- \* A sunflower with  $k$  petals and core  $Y$  is a collection  $C \subseteq F$  of sets s.t.
  - \*  $S \cap S' = Y$  for any two distinct  $S, S'$  in  $C$
  - \*  $S \setminus Y$  is non-empty for every  $S$  in  $C$

Sunflower

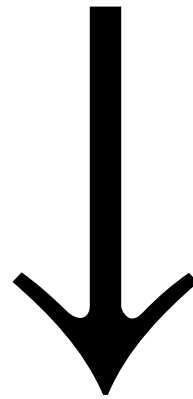


- \* 3 Petals  $\{3\}, \{4\}, \{5, 6\}$ 
  - \* Each petal is non-empty
- \* Core =  $\{1, 2\}$

- \* A set of pairwise disjoint sets is a sunflower with empty core

# Sunflower Lemma

- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family of subsets of  $U$ :  $F = \{S_1, S_2, \dots, S_m\}$  (no duplicates)
- \* Each set in  $F$  has size  $d$
- \*  $|F| > d! (k-1)^d$



- \*  $F$  has a sunflower with  $k$  petals that can be obtained in polynomial time

# Proof of Sunflower Lemma

## Induction on $d$

- \*  $d = 1$  (singletons)

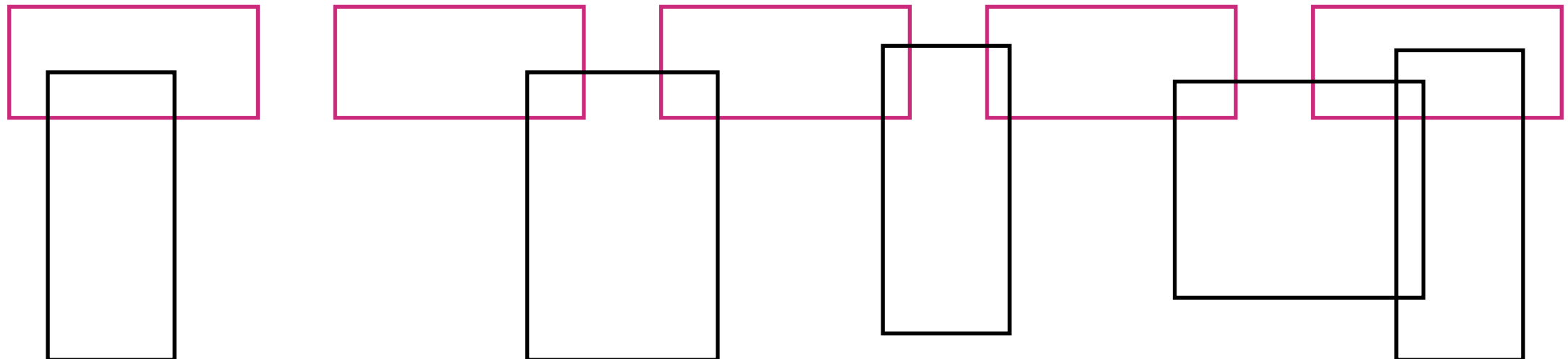
- \*  $d \geq 2$

- \*  $\{S_1, S_2, \dots, S_l\}$  maximal set of disjoint sets in  $F$

- \* If  $l \geq k$ ,  $\{S_1, S_2, \dots, S_l\}$  is the sunflower with  $k$  petals

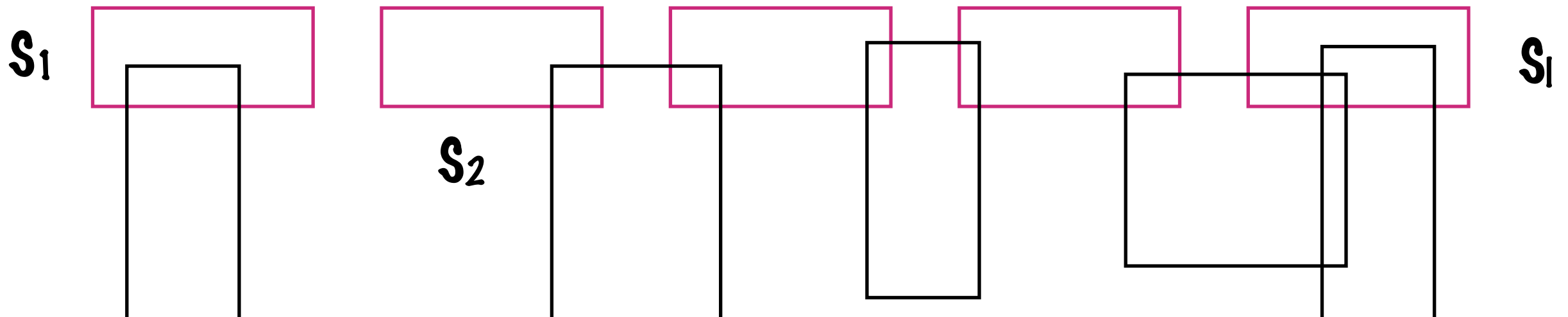
- \* Otherwise,  $l < k$

- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family of subsets of  $U$ :  $F = \{S_1, S_2, \dots, S_m\}$ 
  - \* No duplicates
- \* Each set in  $F$  has size  $d$ ,  $|F| > d! (k-1)^d$



# Proof of Sunflower Lemma

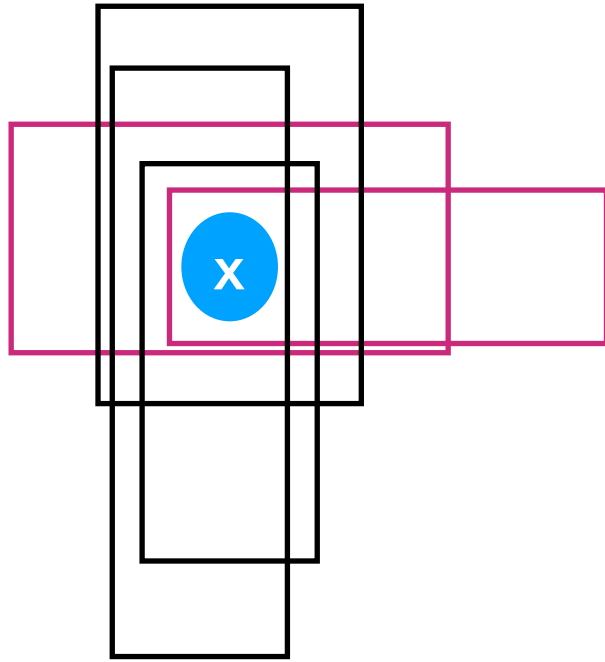
- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family of subsets of  $U$ :  $F = \{S_1, S_2, \dots, S_m\}$ 
  - \* No duplicates
- \* Each set in  $F$  has size  $d$ ,  $|F| > d! (k-1)^d$



Suppose each element of  $S$  is in less than  $|F|/|S|$  sets, then elements of  $S$  are in total  $< |F|$  sets but every set in  $F$  has an element

- \*  $S = S_1 \cup S_2 \cup \dots \cup S_i$  and  $|S| \leq d(k-1)$
- \* Every set in  $F$  has an element in  $S$
- \* There is an element  $x$  that is in  $\geq |F|/|S| > d! (k-1)^d / d(k-1)$  sets of  $F$

# Proof of Sunflower Lemma



- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family of subsets of  $U$ :  $F = \{S_1, S_2, \dots, S_m\}$ 
  - \* No duplicates
- \* Each set in  $F$  has size  $d$ ,  $|F| > d! (k-1)^d$

- \* There is an element  $x$  that is in  $\geq |F|/|S| > d! (k-1)^d / d(k-1)$  sets of  $F$
- \*  $F' =$  sets in  $F$  containing  $x$
- \*  $F''$  : obtained from  $F'$  by deleting  $x$  (no duplicates)
  - \*  $|F''| > (d-1)! (k-1)^{d-1}$
- \* By induction hypothesis,  $F''$  has a sunflower with  $k$  petals
  - \*  $\{S'_1, S'_2, \dots, S'_k\}$
- \*  $\{S'_1 \cup \{x\}, S'_2 \cup \{x\}, \dots, S'_k \cup \{x\}\}$  is a sunflower with  $k$  petals in  $F$

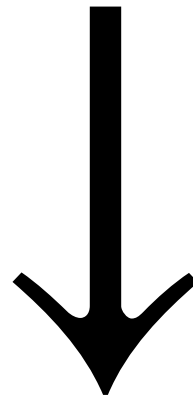
# Computing a Sunflower

- \* Universe  $U = \{1, 2, 3, \dots, n\}$
  - \* Family of subsets of  $U$ :  $F = \{S_1, S_2, \dots, S_m\}$ 
    - \* No duplicates
  - \* Each set in  $F$  has size  $d$ ,  $|F| > d! (k-1)^d$
- 
- \* Find a maximal set  $\{S_1, S_2, \dots, S_l\}$  of disjoint sets in  $F$
  - \* Let  $x$  be an element that is in a maximum number of sets in  $F$ 
    - \*  $x$  is in  $\geq |F|/|S|$  sets of  $F$
  - \*  $F'$  : sets in  $F$  containing  $x$
  - \*  $F''$  : obtained from  $F'$  by deleting  $x$  (no duplicates)
    - \*  $|F''| > (d-1)! (k-1)^{d-1}$  as  $|F|/|S| > d! (k-1)^d / d(k-1)$
  - \* Recurse on  $F''$  to find a sunflower  $\{S'_1, S'_2, \dots, S'_k\}$  with  $k$  petals
  - \*  $\{S'_1 \cup \{x\}, S'_2 \cup \{x\}, \dots, S'_k \cup \{x\}\}$  is a sunflower with  $k$  petals in  $F$



# Sunflower Lemma (Variant)

- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family of subsets of  $U$ :  $F = \{S_1, S_2, \dots, S_m\}$  (no duplicates)
- \* Each set in  $F$  has size at most  $d$
- \*  $|F| > d * d! (k-1)^d$



- \*  $F$  has a sunflower with  $k$  petals that can be obtained in polynomial time

**Hint:** There exists  $r \leq d$  s.t.  $|F_r| > d! (k-1)^d$  where  $F_r$  is the subset of  $F$  containing sets of size  $=r$

# d-Hitting Set

## Input:

- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family  $F = \{S_1, S_2, \dots, S_m\}$  of  $\leq d$  sized subsets of  $U$
- \* A non-negative integer  $k$

**Question:** Does there exist  $V \subseteq U$  with  $|V| \leq k$  s.t for each  $S$  in  $F$ ,  $S \cap V \neq \emptyset$ ?

## Some Common Hitting Sets

- \* Vertex Cover
- \* Feedback Vertex Set
- \* Odd Cycle Transversal
- \* Cluster Vertex Deletion

# d-Hitting Set

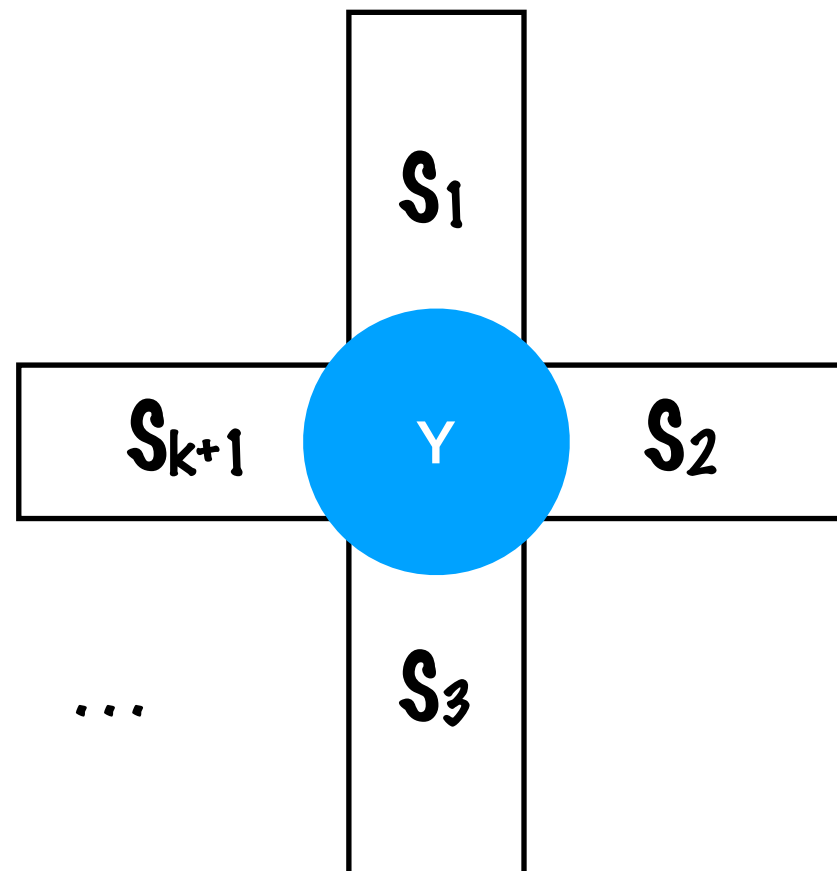
## Input:

- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family  $F = \{S_1, S_2, \dots, S_m\}$  of  $\leq d$  sized subsets of  $U$
- \* A non-negative integer  $k$

**Question:** Does there exist  $V \subseteq U$  with  $|V| \leq k$  s.t for each  $S$  in  $F$ ,  $S \cap V \neq \emptyset$ ?

- \* If  $|F| \leq d * d! k^d$  when something (d) is given as part of problem definition, treat it as c
  - \* Kernel with  $d * d! k^d$  sets and  $d^2 * d! k^d$  elements
- \* Otherwise
  - \* find a sunflower with  $(k+1)$  petals

# d-Hitting Set



forward dirn: Suppose there exist such a  $V$ ; a hitting set, then it clearly hits  $Y$ . reverse dirn: Hitting  $Y$  is equivalent to hitting all of

- \* If  $Y$  is empty then there is no  $(U, F)$  has no hitting set of size  $\leq k$
- \* Otherwise,
  - \* Any hitting set of size  $\leq k$  has a non-empty intersection with  $Y$
  - \* Delete  $S_1, S_2, \dots, S_{k+1}$  from  $F$  and add  $Y$  to  $F$  to get resultant instance  $(U', F', k)$

$(U, F, k)$  is a yes-instance iff  $(U', F', k)$  is a yes-instance

# d-Set Packing

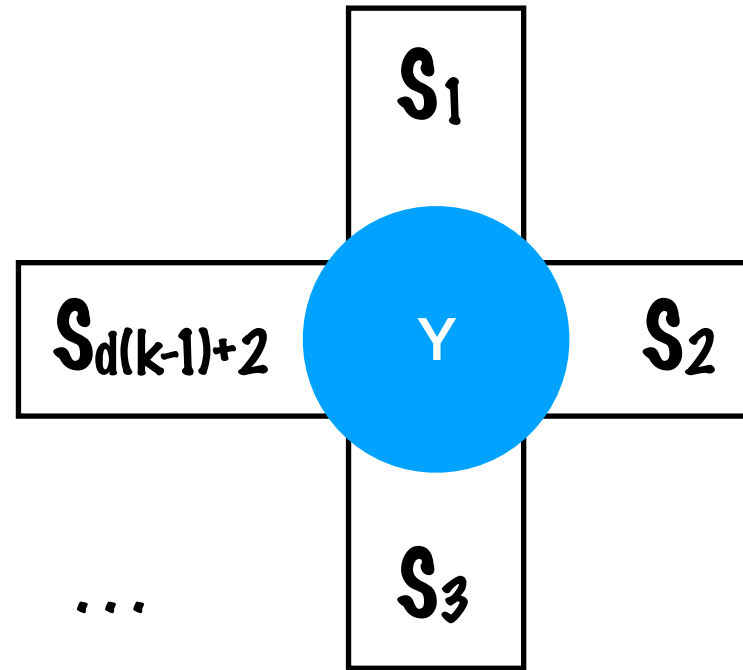
## Input:

- \* Universe  $U = \{1, 2, 3, \dots, n\}$
- \* Family  $F = \{S_1, S_2, \dots, S_m\}$  of  $\leq d$  sized subsets of  $U$
- \* A non-negative integer  $k$

**Question:** Does there exist  $F' \subseteq F$  of pairwise disjoint sets s.t.  $|F'| \geq k$ ?

- \* If  $|F| \leq d * d! (d(k-1)+1)^d$ 
  - \* Kernel with  $d * d! (d(k-1)+1)^d$  sets and  $d^2 * d! (d(k-1)+1)^d$  elements
- \* Otherwise
  - \* find a sunflower with  $d(k-1)+2$  petals

# d-Set Packing



- \* Delete  $S_1$  from  $F$  to get resultant instance  $(U', F', k)$
- \*  $(U, F, k)$  is a yes-instance iff  $(U', F', k)$  is a yes-instance

- \* Reverse direction is easy

- \* Forward direction

for forward dirn: Clearly if  $Y$  is empty then we already have our  $F'$ , o/w we can

- \* Suppose  $P$  is a  $k$ -set packing containing  $S_1$
- \*  $P \setminus \{S_1\}$  has  $\leq d(k-1)$  elements (set  $X$ )
- \* There is a set  $S_i$  that has no element from  $X$  for some  $2 \leq i \leq d(k-1)+2$
- \*  $P \setminus \{S_1\} \cup \{S_i\}$  is a  $k$ -set packing not containing  $S_1$