CUBE: COMPACT PARTICLE-MESH N-BODY ALGORITHM

 ${
m Hao} ext{-}{
m Ran}~{
m Yu}^{1,2},~{
m Ue} ext{-}{
m Li}~{
m Pen}^{2,3,4,5},~{
m Xin}~{
m Wang}^2$

¹Kavli Institute for Astronomy and Astrophysics, Peking University, Beijing 100871, China
 ²Canadian Institute for Theoretical Astrophysics, University of Toronto, Toronto, ON M5H 3H8 Canada
 ³Dunlap Institute for Astronomy and Astrophysics, University of Toronto, Toronto, ON M5S 3H4, Canada
 ⁴Canadian Institute for Advanced Research, CIFAR Program in Gravitation and Cosmology, Toronto, ON M5G 1Z8, Canada
 ⁵Perimeter Institute for Theoretical Physics, Waterloo, ON N2L 2Y5, Canada
 Draft version May 29, 2017

Abstract

Cosmological large scale structure N-body simulations are computation-light but memory-heavy imbalanced problems in supercomputing. Traditional N-body simulation algorithms use at least 24-byte per particle in memory, where six 4-byte single precision real numbers keep track of each particle's phase space coordinate. This is usually the bottleneck of scaling-up of a N-body problem on modern supercomputers whose core/memory ratio is high. Here we present a parallel, memory-lite Particle-Mesh based N-body code, in which each particle can occupy as low as 6-byte memory. This is accomplished by storing only particles' relative location respect to a mesh-grid, and their velocity relative to a predictable distribution function. The remaining information is given by the ordering of particles, and a complimentary density field, which is negligible in memory space. Our numerical experiments show that this algorithm provides acceptable accuracy compared to traditional algorithm in cosmological simulations.

Draft version May 29, 2017 Preprint typeset using ETEX style emulateapj v. 12/16/11

CUBE: COMPACT PARTICLE-MESH N-BODY ALGORITH!

HAO-RAN YU^{1,2}, UE-LI PER^{2,3,4,5}, XIN WANG²

"Canadian Indicates for Astronomy and Astrophysics, Paking University, Beijing 10871, China Canada Tombarda Indicates for Astronomy and Astrophysics, Paking University, Beijing 10871, China Canada Tombarda Indicates for Astronomy and Astrophysics, University of Towarda, Concorda, ON MSS 3H4, Canada Tombarda Indicates for Astronomy and Astrophysics, University of Towarda, Concorda, ON MSS 3H4, Canada Canada Indicates for Astronomy of Towarda Indicates for To

Abstract

imbalanced problems in supercomputing. Tadditional N-body simulation algorithms use at least 24-byte per particle in memory, where six 4-byte single precision real numbers keep track of each particle's phase space coordinate. This is usually the bottleneck of scaling-up of a N-body problem Particle-Mesh Particle Mesh and the particle of the particl



1. INTRODUCTION

N-body simulation is a powerful tool to solve the highly nonlinear dynamic problems (Hockney & Eastwood 1988). It is widely used in cosmology to model the formation and evolution of the large scale structure (LSS). With the fast development of parallel supercomputers, we are able to simulate a system of more than a trillion (10¹²) N-body particles. To date the largest N-body simulation in application its ter "Tank" simulation run on the TianHe-2 supercomputer by costanulation run on the TianHe-2 supercomputer by costanulation run of the TianHe-2's costanulation (2013). It uses nearly 3: 10¹² particles to simulate the cold dark matter (CDM) and cosmic neutrino evolution through the cosmic age. However CDESPM and TianNt (2014) and the cold dark matter (CDM) and cosmic neutrino evolution through the cosmic age. However CDESPM and TianNt (2014) and the cold through the cosmic age. However CDESPM and TianNt (2014) and the cold through the cosmic and the cold through the cosmic and the cold through the cosmic and the cold through through the cold through through the cold through thr

2. ALGORITHM

The most memory consuming part of a N-body simulation is usually the phase space coordinates of each N-body particles = 24 bpp (4 single precision real numbers) must be used to store each particles' position and mean can use as low as 40 bpp in scrifficing the speed. This includes the phase coordinates (24 bpp) for particles includes the phase coordinates (24 bpp) for particles in physical domain and buffered region, a linked list (bpp), and a global coarse mesh and local fine mesh bpp), and a global coarse mesh and local fine mesh bpp. and a global coarse mesh and a global coarse mesh and

1. INTRODUCTION

N-body simulation is a powerful tool to solve the highly nonlinear dynamic problems (Hockney & Eastwood 1988). It is widely used in cosmology to model the formation and evolution of the large scale structure (LSS). With the fast development of parallel supercomputers, we are able to simulate a system of more than

haoran@cita.utoronto.ca

a trillion (10^{12}) N-body particles. To date the largest N-body simulation in application is the "TianNu" simulation run on the TianHe-2 supercomputer by cosmological simulation code CUBEP3M (Harnois-Déraps et al. 2013). It uses nearly 3×10^{12} particles to simulate the cold dark matter (CDM) and cosmic neutrino evolution through the cosmic age. However CUBEP3M and TianNu did not fully use TianHe-2's computing power - 87% of the computing power on TianHe-2 are given by the Intel Xeon-Phi coprocessors while they occupy only 27% of the total memory. Although we accomplished preliminary works on developing native and offload modes of CUBEP3M on Xeon-Phi many integrated core (MIC) architecture, the memory problem is still the bottle neck of scaling up the N-body problem on the heterogeneous architectures. Lower the memory usage per particle will solve this problem, and will also help the N-body problem development on GPUs, which are even more massive parallelized and memory limited. In this paper we present a N-body algorithm which uses as low as 6 byte per particle (here after we use "bpp" referring to "byte per particle") and show that it can give accurate cosmological LSS simulations. The algorithm is presented in §2, and a comparison between this method and traditional method is shown in §3. Discussions and conclusions are in §4.

2. ALGORITHM

The most memory consuming part of a N-body simulation is usually the phase space coordinates of each N-body particles – 24 bpp (4 single precision real numbers) must be used to store each particles' position and velocity. A memory-lite parallel N-body code CUBEP3M can use as low as 40 bpp in sacrificing the speed. This includes the phase coordinates (24 bpp) for particles in physical domain and buffered region, a linked list (4 bpp), and a global coarse mesh and local fine mesh. We attempt to replace the coordinates and linked list

2 Yu et al.

24+4=28 bpp memory usage with a integer based storage, described as follows.

2.1. Particle location storage

Instead of storing the particle's global coordinates, we store its relative position respective to the coarse cell who contains the particle. We divide the coarse cell, in each dimension, evenly into $2^8 = 256$ bins, and we use a 1-byte integer (with range 0 to 255 or -128 to 127) to show which bin it locates in this dimension. Next, global locations of particles are given by cell-ordered format in memory space, and a complimentary number count of particle number in this coarse mesh (coarse density field) will give complete information of particle distribution in cells. Since the this field is usually coarser than particle distribution, it takes negligible memory. This coarse density field can be further compressed into 1-byte integer format, such that a 1-byte unsigned integer show the particle number in this coarse cell in range 0 to 255. In the densest cells (rarely happened) where there are > 255 particles, we can just write 255, and write the actual number as a 4-byte integer in another file. In a simulation with volume L^3 and N_c^3 coarse cells, particle positions are stored with a precision of $L/(256N_c)$, which is usually 10 times finer than the force resolution.

2.2. Particle velocity storage

Velocities, in each dimension, are compressed to 2-byte integers representing their index within a histogram containing 2^{16} equally spaced bins ranging from $-v_{\text{max}}$ to $v_{\rm max}$, where $v_{\rm max}$ is the maximum absolute particle velocity. $v_{\rm max}$ can be derived from either computation or prediction. We prefer the latter case. In cosmological simulations, velocity distribution of matter fluids can be predicted by linear theory, given a model and redshift z (time of the simulation). Then by any time step, we have a cumulative distribution function (CDF) F(|v|) of particles' velocity (see appendix A). We, instead, equally divide the y-axis of this CDF into 2^{16} bins, and the 2byte integer shows which velocity bin the particle locates, in each dimension. This step does finer sampling on lowvelocity end in velocity space and increases the simulation accuracy. This CDF may introduce velocity cutoffs, that particle cannot exceed some certain $v_{\text{max}}(z)$ given by the model. This process prevents the nonphysical scattering behaviours of particles, which do not exist in fluids. Practically, simulation results are very insensitive to the choice of velocity distribution. Even a top-hat distribution in range $-v_{\text{max}}$ to v_{max} give accurate results. One can just assume a Gaussian distribution and whose standard deviation is given by the model and redshift $\sigma(\theta,z)$. Proper choice of model only optimizes the velocity space sampling, but does not affect the physics.

2.3. Spacial decomposition

CUBE uses cubic decomposition structures. The global simulation volume is decomposed into nn³ cubic subvolumes, and each of these are assigned to a coarray $image^1$. Inside of an image, the sub-volume is further decomposed into nnt³ cubic tiles. Each tile is surrounded

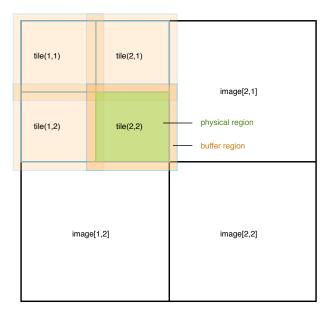


Fig. 1.— Spacial decomposition.

```
program CUBE
  call initialize
  call particle_initialize
  sync all
  call buffer_density
  call buffer_x
  call buffer_v
  do
    call timestep
    call update_x
    call buffer_density
    call buffer_x
    call PM ! update_v
    call buffer_v
    if(checkpoint_step) then
       call checkpoint
       if (final_step) exit
    endif
  enddo
  call finalize
endprogram
```

Fig. 2.— Overall structure of CUBE. by a buffer region which is usually ncb coarse cells thick. The buffer is designed for two reasons: (1) computing the fine mesh force, whose cut-off nforce_cutoff \leq ncb, and (2) collecting all possible particles travelling from a tiles buffer region to its center, physical region. Thus, a integer-based coarse mesh density is defined as

integer(1) rho_c(nex, nex, nex, nnt, nnt, nnt)[nn, nn, *]

where nex = nt + 2ncb covers the buffer region on both sides, nnt is the tile dimensions, and nn is the image $codimensions^2$. Particles' phase coordinates x and v are copied between tiles and images, when necessary. Figure 1 shows the spacial decomposition in a 2-dimensional analogy, with nn = 2 and nnt = 2.

2.4. Code overview

By using the previous compact format storage, we modified the CUBEP3M and utilized it in the outputs

¹ Images are the concept of computing nodes or MPI tasks in Coarray Fortran. We use this terminology in this paper.

 $^{^2}$ Coarray Fortran concept. Codimensions can do communications between images.

CUBE 3

TianNu simulation. On average, 9.125 bpp is used for all outputs and saves 62~% hard drive space on Tianhe-2. The compact version of CUBEP3M, named CUBE, accomplished the compact format in memory, and the code keeps only integer format for storing particles' phase space coordinates. CUBE is written in Coarray Fortran, where the coarray features are used to do MPI communications. Here we review the methodology of solving this integer based N-body problem. Like its predecessor CUBEP3M, CUBE uses two-level mesh grids for solving the Poisson equations. The advantages of this are described in Harnois-Déraps et al. (2013).

Figure 2 shows the overall structure of the code. Subroutine initialize creates necessary FFT plans and read in configuration files telling the program at which redshifts we need to do checkpoints, halofinds, or stop the simulation. Force kernels are also created or read in here. In particle_initialize, for each image, we read in initial positions and velocities in the compressed format. A initial condition generator is also written in the compressed format (see Appendix B). When this step is done by all images (synchronized by "sync all", which is equivalent to a mpi_barrier), the density fields rho_c are buffered between tiles (communications between images are needed) by subroutine buffer_density. Then, particles' phase space coordinates x and v are also buffered. by buffer_x and buffer_v respectively. In these steps, x and v's ordering is changed according to the updated rho_c.

In each iteration of the main loop, we firstly call timestep, where a increment of time dt is controlled by particles' maximum velocities, accelerations, and cosmic expansions. According to dt, subroutine update_x updates the positions of particles in a "gather" algorithm. x = x + v * dt is executed twice, first time to determine a updated density field on the tile, rho_c_new, and second time to generate a new particle list x_new and v_new on the tile. The reason for this repeated execution is the dependence of x_new and v_new's value and ordering on both the old x, v and the new rho_c_new. However, x = x + v * dt scales as o(N) and is computational inexpensive. Although x_new and v_new arrays take extra memory, they do not appear simultaneously for multiple tiles, and compared to x and v, the extra memory overhead is by factor of 1/nnt³. This requires the "gather" algorithm to move particles³ – as we synchronize the correct {rho_c, x, v} in the buffer region of each tile, and the particles in the extended (physical+buffer) region are a superset of particles locating in the physical region in the next time step, given that v * dt is not greater than the buffer depth. In update_x, {rho_c_new, x_new, v_new} can be set to collect only the physical region of the tile, with the rest of this subset discarded. At the end, {rho_c, x, v} is replaced by {rho_c_new, x_new, v_new}, and the memory of latter is freed up.

Next, buffer_density and buffer_x are called again to synchronize the updated particle positions in order that tile-based local fine forces are computed correctly.

A standard particle-mesh scheme is then performed in the PM subroutine. First, by looping over tiles, fine force is calculated and velocities are updated on the physical region of each tile. The memory overhead of fine-

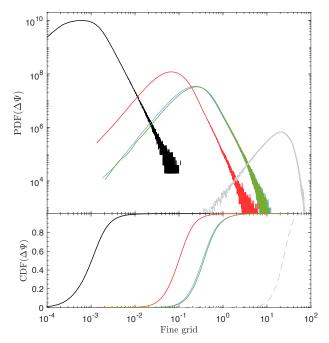


Fig. 3.— Displacement.

mesh arrays are well controlled by 1/nnt³. Next, global coarse force is solved by using a coarser (usually by factor of 4) mesh by dimensional splitting – a distributedmemory cubic decomposed 3D coarse density field is interpolated by particles, and we Fourier transform data in consecutive three dimensions with global transposition in between (known as the pencil decomposition). After the multiplication of force kernels, the inverse transform takes place to get the cubic distributed coarse force field, upon which velocities are updated again. The memory usage of coarse arrays are negligible. An optional particle-particle (PP) force can be called to increase the force resolution and the velocities are updated last time according to this. As discussed in §2.2, updating velocities use model predicted velocity distribution, thus v = v + dv is executed only once, scaling directly to the velocity distribution of the next time step d+dt. During each of these force calculations, maximum accelerations are collected serving for the timestep in the next iteration, controlling next dt.

By far, particle velocities in the physical regions of each tile are updated. Particle locations in the buffer regions has updated before PM and remained unchanged. So we simply call buffer_v again, such that the update_x in the next iteration will be done correctly.

If a desired redshift is reached, we call checkpoint and/or exit the main loop. The corresponding logical variables are controlled in timestep. Finally, in finalize subroutine we destroy all the FFT plans and finish up any timing or statistics taken in the simulation.

2.5. Memory layout

Here we summarize the memory usage in unit of bpp (byte per particle).

3. RESULTS

We use the same initial conditions, and use CUBEP3M and CUBE to simulate the large scale structure formation separately and compare their results.

³ In contrast, CUBEP3M uses a "scatter" algorithm.

4 Yu et al.

Results of izipx=1 Results of izipx=2

Fig. 4.— Power spectrum differential plot

4. DISCUSSION AND CONCLUSION

Cosmological simulation, PM method. PID. Summarize memory. Phi, GPU.

CUBE 5

APPENDIX

A. VELOCITY DISTRIBUTION FUNCTION

To optimize the storage of the velocity field, we need to estimate the velocity dispersion at each snapshot. At non-linear scale, the velocity distribution function is surely non-Gaussian. However, to the first order approximation, we simply assume it as Gaussian and characterized by the variance

$$\sigma_v = \frac{1}{3} (aHfD)^2 \int d^3q \frac{P_L(q)}{q^2} \tag{A1}$$

where a(z) is the scale factor, H(z) the Hubble parameter, D is the linear growth factor, $f = d \ln D/d \ln a$, and $P_L(k)$ is the linear power spectrum of density contrast at redshift zero.

B. INITIAL CONDITION GENERATOR

Thanks.

REFERENCES

Harnois-Déraps, J., Pen, U.-L., Iliev, I. T., Merz, H., Emberson,
J. D., & Desjacques, V. 2013, MNRAS, 436, 540, 1208.5098
Hockney, R. W., & Eastwood, J. W. 1988, Computer simulation using particles