

Learning to Discover Efficient Mathematical Identities

Wojciech Zaremba^{1,3}, Karol Kurach^{2,3}, Rob Fergus^{1,4}

¹Courant Institute NYU; ²University of Warsaw; ³Google; ⁴Facebook

The Combinatorial Explosion Discrete search is a central issue in AI



From “Artificial Intelligence: A General Survey” by Professor Sir James Lighthill, 1973.

“...failure to recognize the implications of the **combinatorial explosion**. This is a general obstacle to the construction of a self-organising system on a large knowledge base which results from the explosive growth of any combinatorial expression, representing numbers of possible ways of grouping elements of the knowledge base according to particular rules, as the base’s size increases.”

Math & Theorem Proving

Contemporary theorem provers

- Requires search over all possible combinations of operators
- Intractable for all but simple proofs
- Automated Theorem Provers use heuristics

Yet (some) humans are able to prove theorems

- Have experience of related problems
- Known math “tricks”

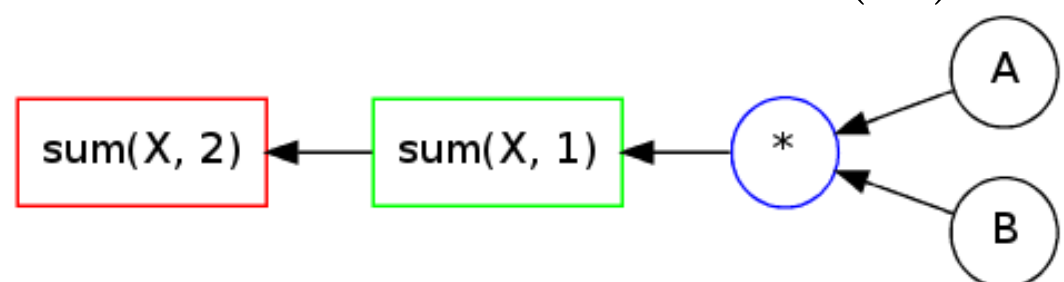
We focus on simpler problem: discovering identities in mathematics.

Toy Example

Consider two matrices A and B :

$$\sum_{i,k} (AB)_{i,k} = \sum_i \sum_j \sum_k a_{i,j} b_{j,k}$$

Naive computation takes $O(n^3)$:



Discovering Efficient Identities

Define a grammar G of operators Given some target expression T within the domain of G e.g. $\text{sum}(\text{sum}(A*B,2),1)$. Find an identical expression that has lower computational complexity i.e. avoids high complexity operators (like matrix multiplication).

| | | Polynomials of degree 1 |
|-------------------------|------------------------|-------------------------|
| Matrix-matrix multiply | $X * Y$ | |
| Matrix-vector multiply | $X * y$ | |
| Matrix-element multiply | $X .* Y$ | |
| Matrix transpose | X | |
| Column-sum | $\text{sum}(X,1)$ | |
| Row-sum | $\text{sum}(X,2)$ | |
| Column-repeat | $\text{repmat}(X,1,m)$ | |
| Row-repeat | $\text{repmat}(X,n,1)$ | |

$$A, A^T, \sum_i A_{i,:}, \sum_j A_{:,j}, \sum_{i,j} A_{i,j}, \dots$$

Polynomials of degree 2

$$A^2, (A^2)^T, AA^T, A^T A, \sum_i (AA^T)_{i,:}, \dots$$

Space grows exponentially fast.

Representing Symbolic Expressions

Target expression: $\text{sum}(\text{sum}(A*A,1),2)$ Use P copies of A
Representation of target is descriptor vector (length P)

- Each element is evaluation one copy
- Vector is of length P
- If descriptors match equivalent expressions.

Using is unstable, so use integers modulo q , where q is large prime.

Example: Taylor Series Approximation

Consider RBM partition function:

$$\sum_{v,h} \exp(v^T W h) = \sum_{k=0}^{\infty} \sum_{v,h} \frac{(v^T W h)^k}{k!}$$
$$v \in \{0,1\}^n$$
$$h \in \{0,1\}^m$$

1st term in Taylor series:

$$\sum_{k=0}^{\infty} \sum_{v,h} v^T W h = 2^{n+m-2} \sum_{i,j} W_{i,j}$$

2st term in Taylor series:

$$\sum_{k=0}^{\infty} \sum_{v,h} (v^T W h)^2 = 2^{n+m-4} [\sum_{i,j} W_{i,j}^2 + (\sum_{i,j} W_{i,j})^2 + \sum_i (\sum_j W_{i,j})^2 + \sum_j (\sum_i W_{i,j})^2]$$

Note that identity computes in $O(nm)$ versus $O(2^{n+m})$ for original expression.



6th term. Higher terms too complex for manual derivation.

Prior Over Computation Trees

Recall goal: find equivalent expressions to target i.e. descriptors match Restrict grammar to use operators with lower complexity than target If any match found then sure to be efficient w.r.t. target

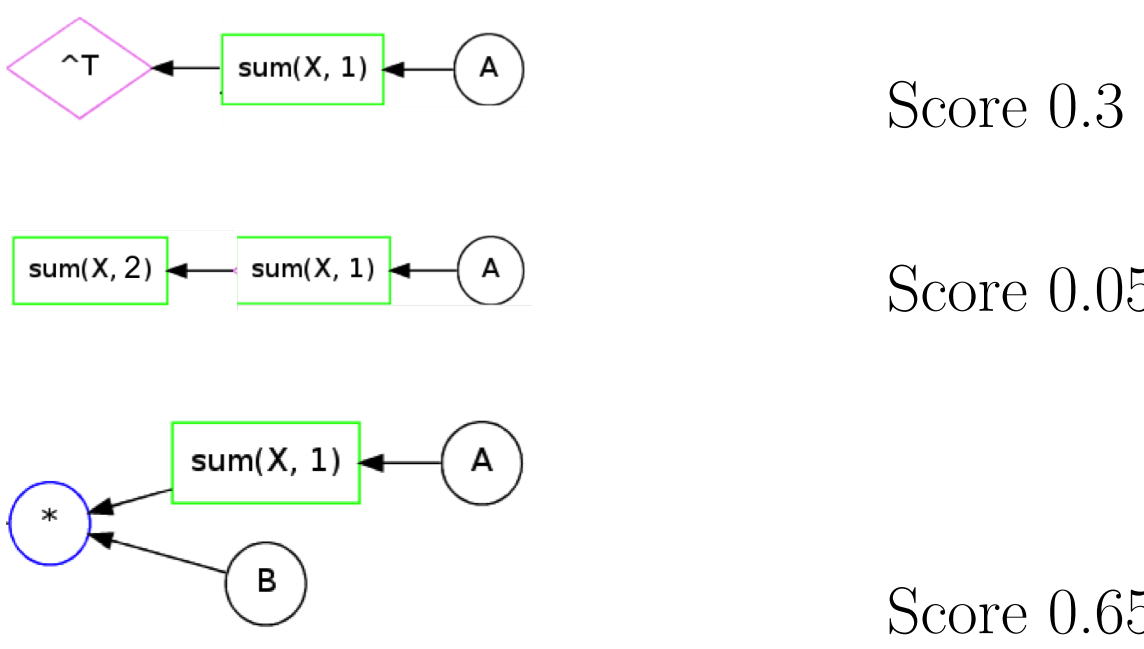
Want to learn a good **prior** over expressions

Scheduler picks potential new operators to append to current expression(s).

Example:



Scorer ranks each possibility (i.e. how likely they are to lead to the solution), using **prior**.



Sample new operator according to scorer probabilities

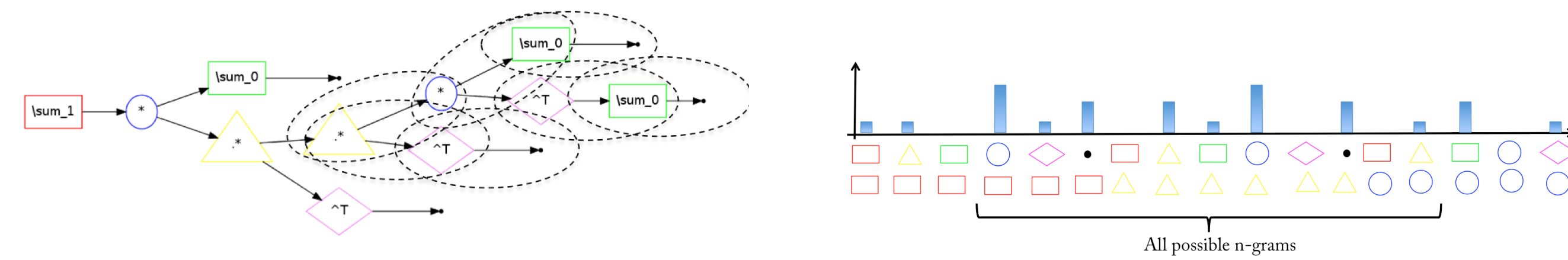
Scorer Strategy

- Naive: no prior, Just select randomly from all valid operators
- n -gram prior
- Recurrent Neural Network prior

Use curriculum learning approach. Start with easy targets, and uniform prior (low polynomial degree k). Refine prior using new solution

| | Target Expression | Efficient expression |
|---------|----------------------------|---|
| $K = 1$ | $\sum_i a_i$ | $\text{sum}(\text{sum}(A, 1), 2)$ |
| $K = 2$ | $\sum_{i,j} a_i a_j$ | $0.5 * (\text{sum}((\text{repmat}(\text{sum}(A', 1), 1, m) .* A), 2)) +$ $-0.5 * (\text{sum}((A .* A)', 1)))$ |
| $K = 3$ | $\sum_{i,j,k} a_i a_j a_k$ | $(-0.5 * (\text{sum}((\text{repmat}(\text{sum}((A' .* A'), 1), 1, m) .* A), 2)) +$ $1/3 * (\text{sum}(((A' .* A')' .* A), 2)) +$ $1/6 * (\text{sum}((A' .* \text{repmat}(\text{sum}((\text{repmat}(\text{sum}(A', 1), 1, m) .* A), 2), m, 1))), 1)))$ |

n -gram

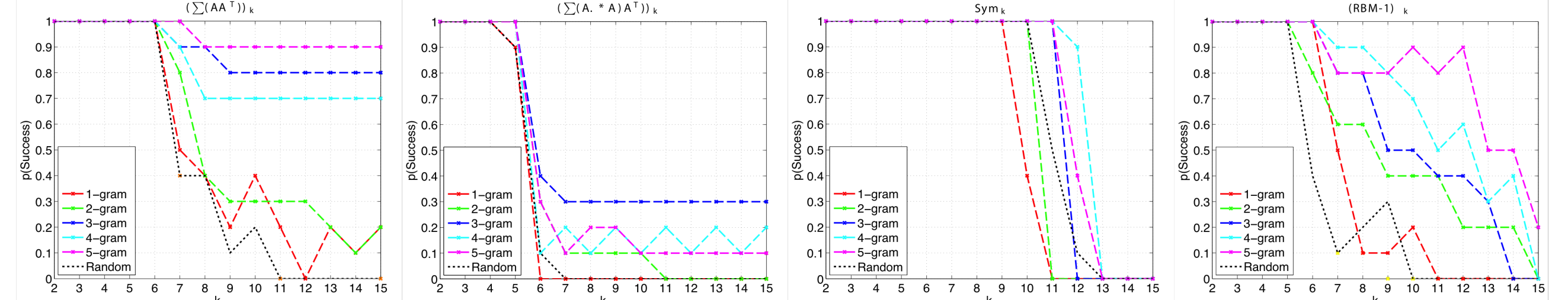


Experiments in discovering identities

5 families of expressions (vary degree k)

- Multiply-sum: $\sum (AA^T)_k$
- Element-wise multiply-sum: $\sum ((A .* A)A^T)_k$
- Symmetric polynomials: $\sum_{i,j,k} a_i a_j a_k$
- RBM-1: $\sum_{v \in \{0,1\}^n} (v^T A)^k$
- RBM-2: $\sum_{v \in \{0,1\}^n, h \in \{0,1\}^m} (v^T A h)^k$

Start with $k = 1$ and work up to $k = 15$. Time cut-off: 600 seconds Repeat 10 times, measure fraction successful

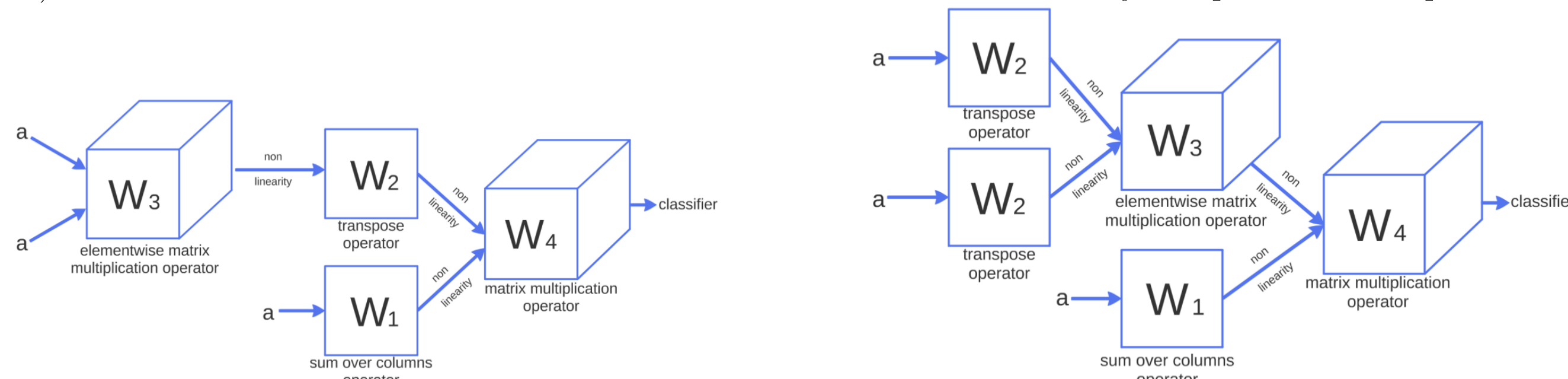


RBM-2

No scorer strategy able to get beyond $k = 5$ However, the $k = 5$ solution was found by the RNN consistently faster than the random strategy (100 +- 12 vs 438 +- 77 secs).

Recursive neural network strategy

n -gram has a shallow understanding of mathematical expressions (up to depth n). We trained recursive neural network to identify equivalent expressions.



| | Degree $k = 3$ | Degree $k = 4$ | Degree $k = 5$ | Degree $k = 6$ |
|-----------------------|----------------|------------------|------------------|------------------|
| Test accuracy | 100% \pm 0% | 96.9% \pm 1.5% | 94.7% \pm 1.0% | 95.3% \pm 0.7% |
| Number of classes | 12 | 125 | 970 | 1687 |
| Number of expressions | 126 | 1520 | 13038 | 24210 |