

Computer Systems Principles

Data Representation- Bits and Bytes



Today's Class

1. Learn data representation in binary and hexadecimal
 - represent real numbers, negative numbers and fractions.
 2. Perform operations
 - addition, subtraction, multiplication, and division
- Announcements: Moodle Quiz 2 is up.

Group Task

- Form partners
- Using only the three symbols @ # & represent:
 - integers 0 – 10
 - represent movements:
 - 1) stand, 2) sit, 3) walk , 4)raise hand, 5) turn around 6) hop
- Using only your representation write down a series of commands and integers (raise arm - 3, turn - 2)
 - must have at least 5 commands

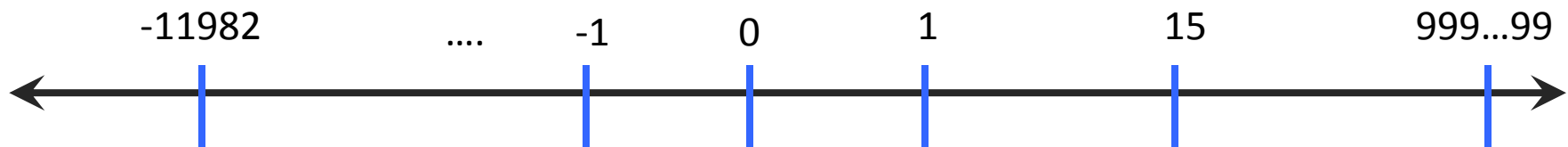
How do we think about numbers?

- Representing and reasoning numbers?
 - Computers store variables (data).
 - Data is typically composed of numbers and characters.
 - Want a representation that is sensible.

Decimal Number Systems



- Digits 0-9
- **Positional**
- Digits are composed to make larger numbers
 $12 = 1 \cdot 10^1 + 2 \cdot 10^0$
- **Base 10** representation

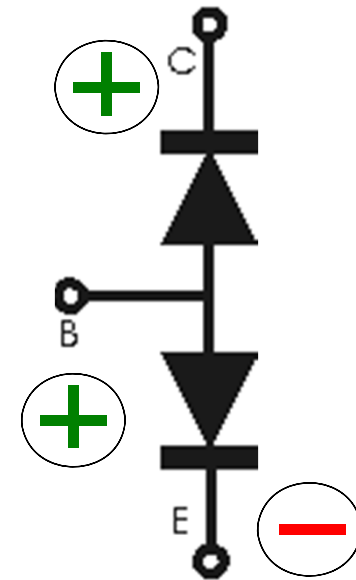
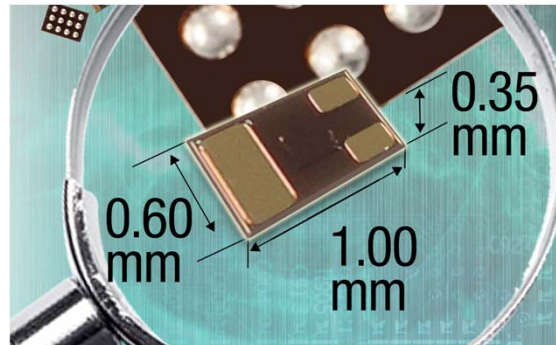


Given four positions: $[X X X X]_{10}$ what is the largest number you can represent?

- **9999_{10}**

Binary Number Systems

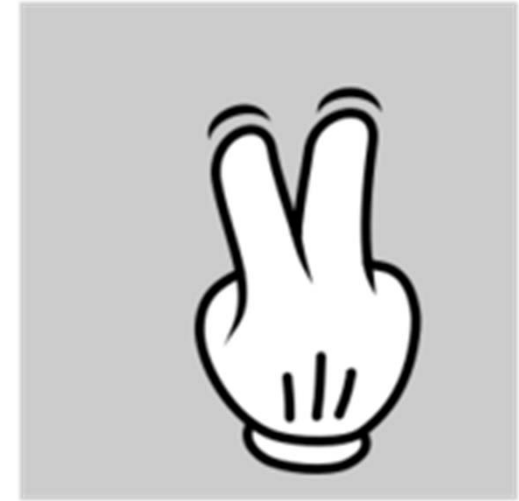
- Transistors: **On** or **Off**
- Optical: **Light** or **No light**
- Magnetic: **Positive** or **Negative**



Binary Digits: (BITS)

- Everything is bits
- Thinking in Base 2

Dec	0	1	2	3	4	...	10	...
Bin	0	01	10	11	100	...	1010	...



The number line



Binary Digits: (BITS)

Most significant bit \longrightarrow $\overset{7\ 6\ 5\ 4\ 3\ 2\ 1\ 0}{\underline{1}000111\underline{1}}$ \longleftarrow Least significant bit

Representation:

$$1 \times 2^7 + 0 \times 2^6 + \dots + 1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$

$$10001111 = 143$$

01011000 00010101 00101110 11100111

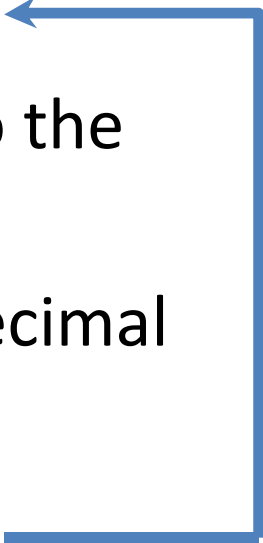
Represents: $0 \times 2^{31} + 1 \times 2^{30} + 0 \times 2^{29} + \dots + 1 \times 2^0$

- Sequence of eight bits: byte
- Sequence of four bits: nibble

Conversion: Decimal to Binary

Method 1:

We have a decimal number (x):

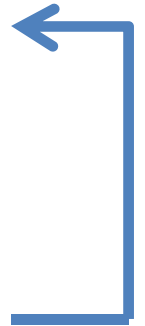
1. Highest **power of two** less than or equal to the decimal number (y)
 2. **Subtract** the power of two (y) from the decimal number (x) as $x = x - y$.
 3. If $x = 0$, we have our result! Else: **Repeat**
- 

Conversion: Decimal to Binary


Method 2:

We have a decimal number (x):

1. Divide (x) by 2 to obtain a **quotient and remainder**.
 - **Remainder** is 0 or 1.
2. If **quotient $\neq 0$** , assign (x) = quotient back to step 1.
3. If the **quotient = 0** proceed!
4. The sequence of remainders is the binary representation.



Example conversion by Method 2

quotient	remainder	
35	1	 read upward: $35_{10} = 100011_2$
17	1	
8	0	
4	0	
2	0	
1	1	
0		

Note: remainders always 0 or 1:
Easy to see by even versus odd

Hexadecimal Representation

Translate the binary to decimal?

1001000100001011001010001101001010001101

- Hexadecimal numbers use 16 digits: {0-9, A-F}
 - does not distinguish between upper and lower case!
- Encoding byte values using Octal numbers: use 8 digits {0 – 7}

DEC	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
HEX	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F

Hexadecimal Representation

- Bit patterns as base-16 numbers
- Convert binary to hexadecimal: by splitting into groups of **4 bits** each.
- If not a multiple of 4 (nibble):
 - pad the number **with leading zeros**.
- Example: $11\ 1100\ 1010\ 1101\ 1011\ 0011_2 = 3CADB3_{16}$

Bin	11	1100	1010	1101	1011	0011
Hex	3	C	A	D	B	3

iClicker Activity

Convert the decimal number 231 to binary and hexadecimal equivalents.

a) 1110 1111, F7

b) 1110 0110, E6

c) 1110 0111, E7

d) 0110 0111, 57

DEC	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
HEX	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F

Binary Addition

Rules

- $1+0 = 1$
- $1+1 = 10$
- $1+1+1 = 11$

	1	1	1	0		Carry in
		1	0	1	1	Augend (A)
+		1	1	1	0	Addend (B)
<hr/>						
	1	1	0	0	1	Sum

Binary Multiplication

Example:

$$\begin{array}{r}
 11 \\
 \times 13 \\
 \hline
 33 \\
 + 11- \\
 \hline
 143
 \end{array}$$

$$\begin{array}{r}
 1011 \\
 \times 1101 \\
 \hline
 1011 \\
 {}^10000- \\
 {}^{11}011-- \\
 + {}^{11}011-- \\
 \hline
 10001111
 \end{array}$$

- $1+0 = 1$
- $1+1 = 10$

Binary Multiplication

- Since we always multiply by either **0** or **1**, the **partial products** are always either 0000 or the multiplicand (in this example: **1011**).
- There are four partial products which are added to form the result.

iClicker Activity

Multiply the two numbers: 1101×1001

- a) 110 1001
- b) 100 1111
- c) 110 1101
- d) 111 0101

iClicker Activity

- Largest binary integer that can be stored in 3 bits?

a) 001

b) 100

c) 111

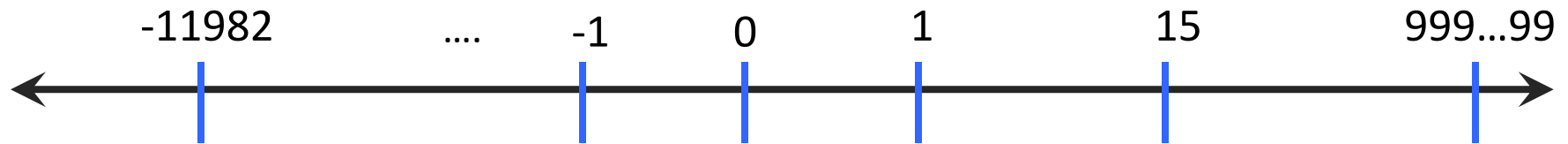
d) None of these

What about the smallest?

Unsigned Binary representation

- Issues:
 - complexity of arithmetic operations
 - negative numbers
 - maximum representable number

Negative Binary Numbers



- Which bit should we pick?
 - Left-most bit
 - Also called the **most significant bit**

Negative Binary: Sign-Value

To get the negative number, set the **sign bit to 1** and leave all the **other bits unchanged**.

81 =	0	0	0	0	0	0	0	0	0	1	0	1	0	0	0	1
-81 =	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0	1

Called **sign-magnitude** representation: sign bit + value

0 =	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0 =	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Two zeroes! Not fun

More difficulties ...

- It is not easy to add, subtract, or compare numbers in sign-magnitude format
- Everything get complicated by whether the signs of the numbers are the same or different, etc.
- Values do not fall in a natural order as compared with unsigned numbers

An alternative: Ones Complement

To get the negative of a number, **invert or complement** all the bits.

Complement = (\sim)

- $\sim 1 = 0$ and $\sim 0 = 1$

81 =	0	0	0	0	0	0	0	0	0	1	0	1	0	0	0	1
-81 =	1	1	1	1	1	1	1	1	1	0	1	0	1	1	1	0

This solves the “unnatural order” problem

0 =	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0 =	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1

But we still have two zeroes!

Two's Complement

To get the negative of a number, **invert** all the bits and then **add 1**.

If the addition causes a carry bit past the most significant bit, **discard the high carry**.

81 =	0	0	0	0	0	0	0	0	0	1	0	1	0	0	0	1
------	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

~81 =	1	1	1	1	1	1	1	1	1	0	1	0	1	1	1	0
-------	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

+1

-81 =	1	1	1	1	1	1	1	1	1	0	1	0	1	1	1	1
-------	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

Two's Complement

Do we have a unique representation of the value zero? **Yes!**

[illegible][illegible]

+1

[illegible]

Two's Complement: Advantages

- ✓ Unique representation of zero
- ✓ Exactly same method for addition, multiplication, etc. as unsigned integers **except** throw away the high carry (high borrow for subtraction).
- ✓ Also turns out to lead to reasonable simple electronic circuits: nice for building CPUs!

Two's Complement: Addition, Subtraction

$$45 - 14 = 45 + (-14) = 31$$

45 =	0	0	0	0	0	0	0	0	0	0	1	0	1	1	0	1
------	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

-14 =	1	1	1	1	1	1	1	1	1	1	1	1	0	0	1	0
-------	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

31 =	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1
------	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



ignore high carry 1!

Two's Complement Range

- Representing negative and positive numbers in b bits:
 - Unsigned: 0 to $2^b - 1$
 - Signed: $-2^{(b-1)}$ to $2^{(b-1)} - 1$
- Example: range for 8 bits is:
 - Unsigned: 0 to $2^8 - 1 = 0 \dots 255$
 - Signed: $-2^{(7)}$ to $2^{(7)} - 1 = -128 \dots 127$
- Can think of the numbers as being arranged in a large circle
 - Signed and unsigned simply start at different places ...
 - ... and label differently the values with high bit of 1

Overflow & Underflow

- For an unsigned number: overflow happens when the **last carry cannot be accommodated**.
- For a signed number: overflow happens when the **most significant bit is not the same** as every bit to its left.
 - sum of two positive numbers is negative: overflow
 - sum of two negative numbers is positive: underflow
 - sum of a positive and negative number will never overflow!

Overflow: Example

Consider the 8-bit two's complement addition:

[illegible]

- The result should be +128, but the leftmost bit is 1, therefore **the result is -128!**
- **This is an overflow:** an arithmetic operation that should be positive gives a negative result.

Two's Complement Underflow

- Just as **overflow** is caused by a value near the **upper limit** of the range, **an underflow** is caused by values near the **lower limit** of the range.

iClicker Activity

- What is the result of the following 8-bit two's complement addition $1000\ 0000 - 1$?

- a) 0111 1111
- b) 1111 1111
- c) 0000 0000
- d) 0000 0001

Hint: Try converting it back to decimal and compare!

Underflow: Example

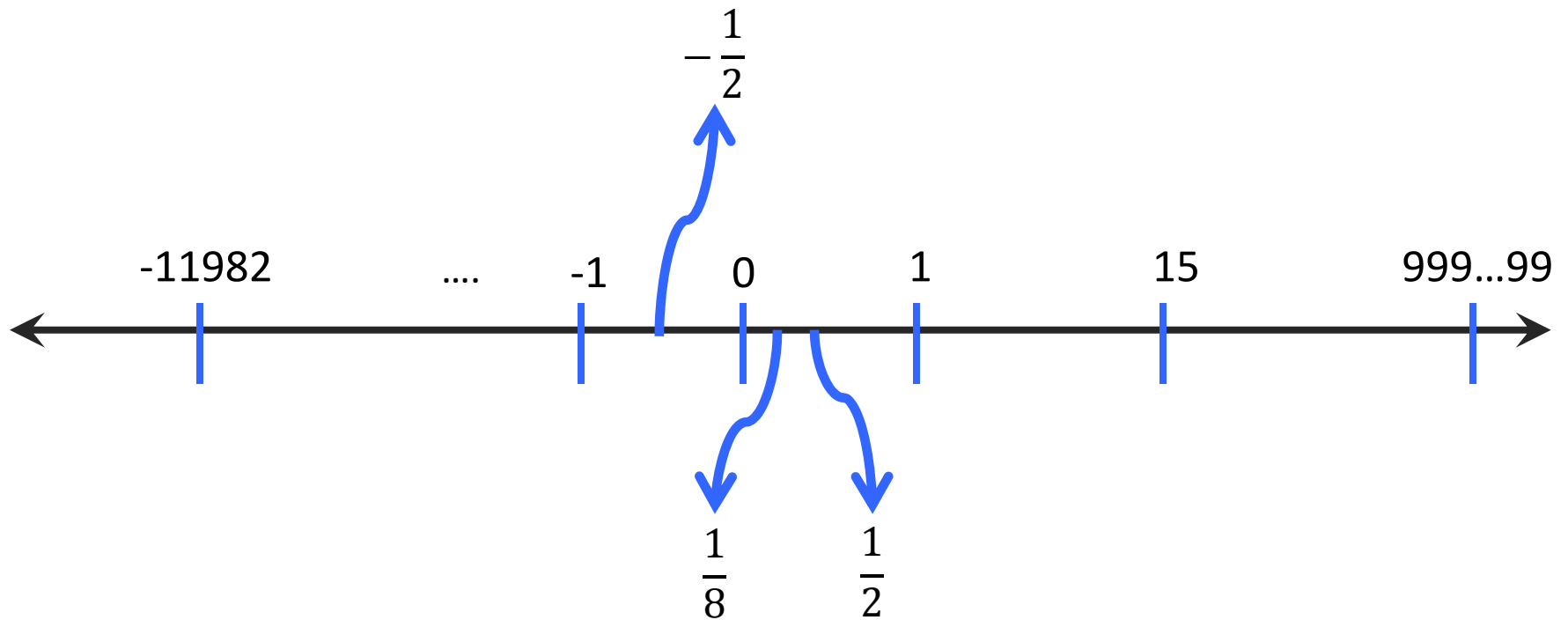
Consider the 8-bit two's complement addition:

$$\begin{array}{r} -128 \quad -1 \quad -1 \\ -128 \quad -1 \quad -1 \\ \hline 0 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \end{array}$$

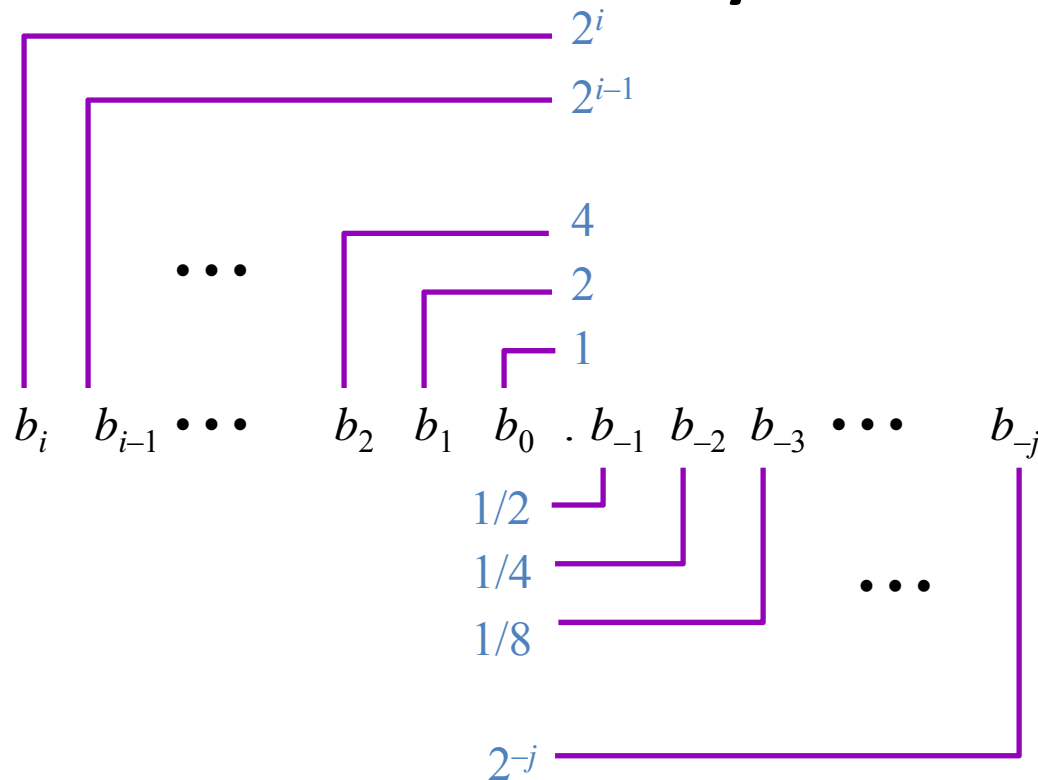
- The result should be -129, but the leftmost bit is 0, therefore **the result is +127!**
- **This is an underflow**: as an arithmetic operation that should be negative gives a positive result.

Fractional binary numbers

How do we represent fractions in binary?



Fractional Binary Numbers



- Representation

- Bits to right of “binary point” represent fractional powers of 2
- Represents the rational number:

$$\sum_{k=-j}^i b_k \cdot 2^k$$

IEEE Single Precision Floating point

- sign bit
- fraction
- exponent



- $(-1)^s \times M \times 2^E$
 1. Sign bit s determines whether number is negative or positive
 2. Fractional bit M normally a fractional value in range $[1.0, 2.0)$.
 3. Exponent E weights value by power of two

Next Class

- Lets represent Binary operations in C!
- Overflow conditions
 - How do we represent overflow?
 - What's the impact of an overflow?
 - How do we mitigate it?
- Readings posted on website