

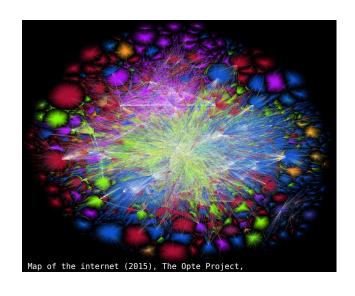
Computer Systems (Computer Networks)

David Marchant

Based on slides compiled by Marcos Vaz Salles, with adaptions by Vivek Shah and Michael Kirkedel Thomsen

Why study Computer Networks?

- How can we build networked applications?
- What are the protocols that power the Internet?
- How can we secure data transmission?





What should we learn in this portion of the course?



- Describe the design of application-layer protocols such as HTTP and DNS.
- Implement networked applications making use of sockets.
- Explain the mechanisms used by transport-layer protocols to achieve multiplexing, reliability, flow control, and congestion control.
- Describe network setups involving subnets, NAT, and LAN segments as well as related interconnection hardware such as routers, switches, and hubs.
- Explain the use of cryptography and operational measures to secure network protocols and applications.
- Build for performance by using all the hardware we have, both locally and remotely.
- Utilise events as basis for applications.
- De-mystify cloud computing and its impacts on data center networking

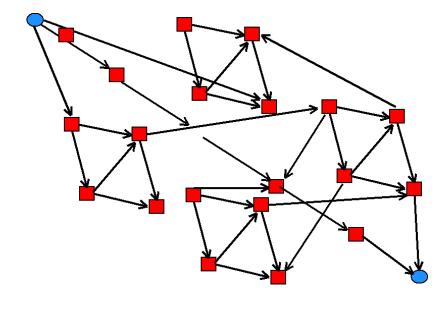


- How can we build networked applications?
 - Application-level protocols, e.g., HTTP and content delivery
 - Programming with sockets
 - Resolving names with DNS





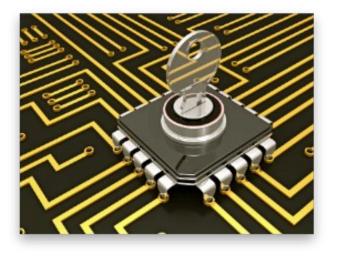
- What are the protocols that power the Internet?
 - UDP: Basic transport
 - TCP: Reliable and ordered transport with flow and congestion control
 - IP: Addressing, forwarding, routing
 - Ethernet et al: Physical transmission



Largest portion of the course



- How can we secure data transmission?
 - Cryptography
 - Authentication
 - HTTPS, IPSec: Securing protocols





- Not only theoretical knowledge, but also skills
 - Programming with sockets
 - Implementing protocols
 - Using network tools
 - Building distributed applications





References & Course Materials

- Book
 - Computer Networking, 8th ed., James F. Kurose and Keith W. Ross, Pearson, ISBN 13: 978-0-135-92873-8
- Ex-book
 - Computer Systems: A Programmers Perspective, 3rd ed., Randal Bryant and David O'Hallaron, ISBN 13: 978-0-134-09266-9
- Other references
 - Vast majority listed in the course schedule
 - Will keep updating them as we go



Acknowledgements

- Many of the slides in this course are based on or reproduce material kindly made available by Michael Freedman (Princeton), James Kurose & Keith Ross (RPI & NYU-Poly, textbook), Jerome Saltzer & M. Frans Kaashoek & Robert Morris (MIT), Randal E. Bryant & David R. Halloran (Computer Systems, textbook)
- Marcos Vaz Salles (Associate Professor, DIKU) for creating and compiling these slide sets



Questions so far?



What should we learn today?

- Identify the key concepts in networking of protocols, layering, resource allocation, and naming.
- Describe what a protocol is and the main issues in protocol design.
- Explain the goal of layering in networked systems and the multiple layers in the Internet protocol suite.
- Explain how circuit switching and packet switching address the resource allocation problem in networks.
- Explain Time-Division Multiplexing (TDM) and Frequency-Division Multiplexing (FDM) and predict transmission time in circuit switched networks.
- Explain the advantages and challenges of packet switching and store-and-forward, and identify the components of transfer time, including processing, queueing, transmission, and propagation.
- Predict transfer time in such store-and-forward networks.
- Explain the notion of throughput in store-and-forward networks and predict throughput in specific scenarios.



Networking is Relevant



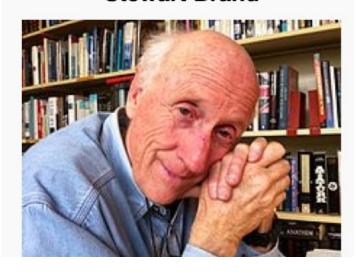
Information wants to be free because it has become so cheap to distribute, copy, and recombine... It wants to be expensive because it can be immeasurably valuable to the recipient. (1985)



Google news

WikipediA

Stewart Brand









What's the Internet: "nuts and bolts" view



PC



server



wireless laptop



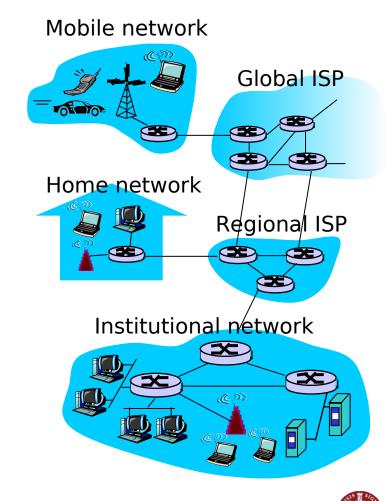
cellular handheld





- A network of networks
- Millions of connected computing devices:

 hosts = end systems
 - running network apps
- communication links
 - fiber, copper, radio, satellite
 - transmission rate = bandwidth
- routers: forward packets (chunks of data)





Key Concepts in Networking

Protocols

- Speaking the same language
- Syntax and semantics

Layering

- Standing on the shoulders of giants
- A key to managing complexity

Resource allocation

- Dividing scarce resources among competing parties
- Memory, link bandwidth, wireless spectrum, paths

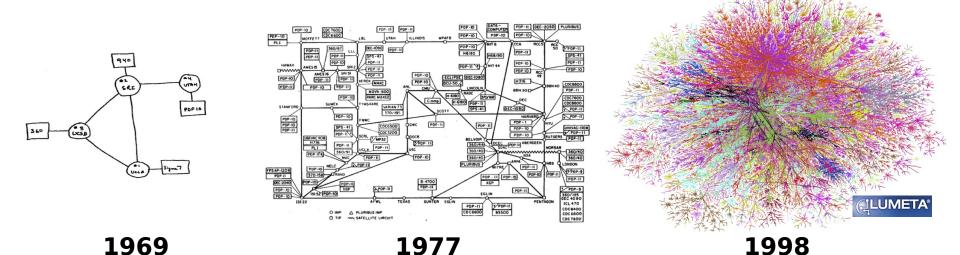
Naming

What to call computers, services, protocols, ...



Key Concepts in Networking

- Protocols
 - Speaking the same language
 - Syntax and semantics



All speak IPv4 "Internet Protocol version 4"



Protocol design is about tradeoffs

How should hosts and routers communicate?

- Standard protocol
- Fast: Machine readable in hardware at line rates

Browsers, web servers, and proxies?

- Can be slower: software readable
- Human readable
- Extensible and forward-compatible
- Not everybody might be familiar with extensions



IPv4 Packet



4-bit Header Version Length		8-bit Type of Service (TOS)	16-bit Total Length (Bytes)	
16-bit Identification			3-bit Flags	13-bit Fragment Offset
8-bit Time to Live (TTL)		8-bit Protocol	16-bit Header Checksum	
32-bit Source IP Address				
32-bit Destination IP Address				
Options (if any)				
Payload				

20-byte header



Example: HyperText Transfer Protocol

GET /courses/archive/spr09/cos461/ HTTP/1.1

Host: www.cs.princeton.edu

User-Agent: Mozilla/4.03

CRLF

Request

HTTP/1.1 200 OK

Date: Mon, 2 Feb 2009 13:09:03 GMT

Server: Netscape-Enterprise/3.5.1

Last-Modified: Mon, 21 Feb 2009 11:12:23 GMT

Content-Length: 42

CRLF

Site under construction

Response



Key Concepts in Networking

- Protocols
 - Speaking the same language
 - Syntax and semantics
- Layering
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 - A key to managing complexity



Layering = Functional Abstraction

- Sub-divide the problem
 - Each layer relies on services from layer below
 - Each layer exports services to layer above
- Interface between layers defines interaction
 - Hides implementation details
 - Layers can change without disturbing other layers

Application

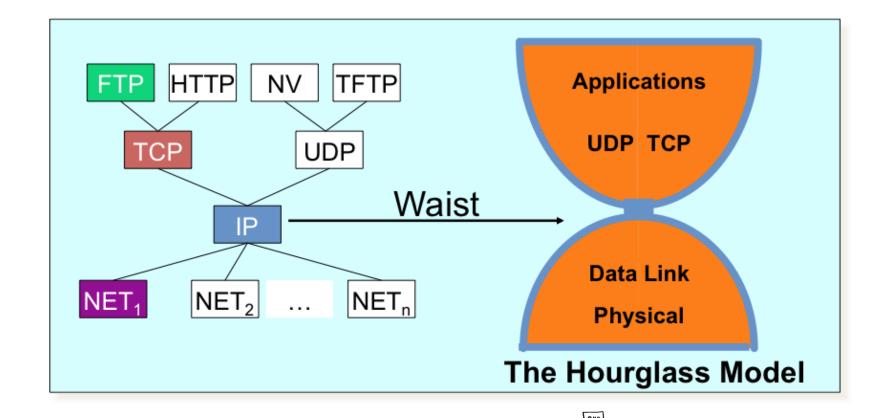
Application-to-application channels

Host-to-host connectivity

Link hardware



The Internet Protocol Suite



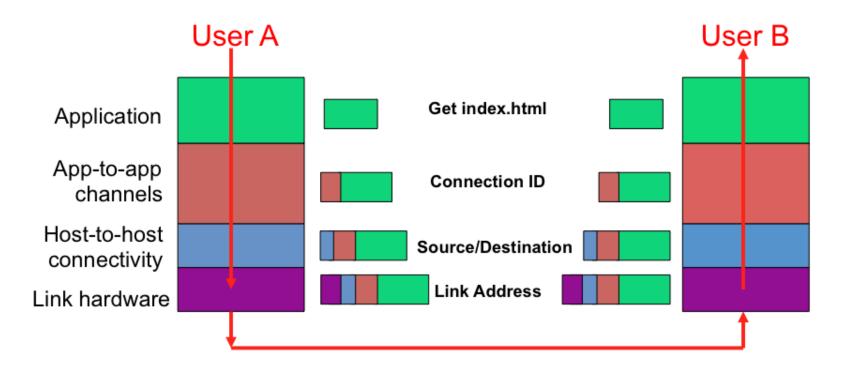
The waist facilitates interoperability



Source: Freedman

Layer Encapsulation in HTTP

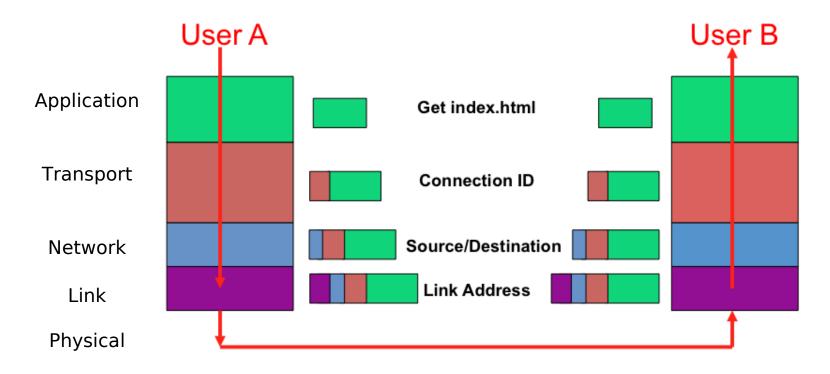






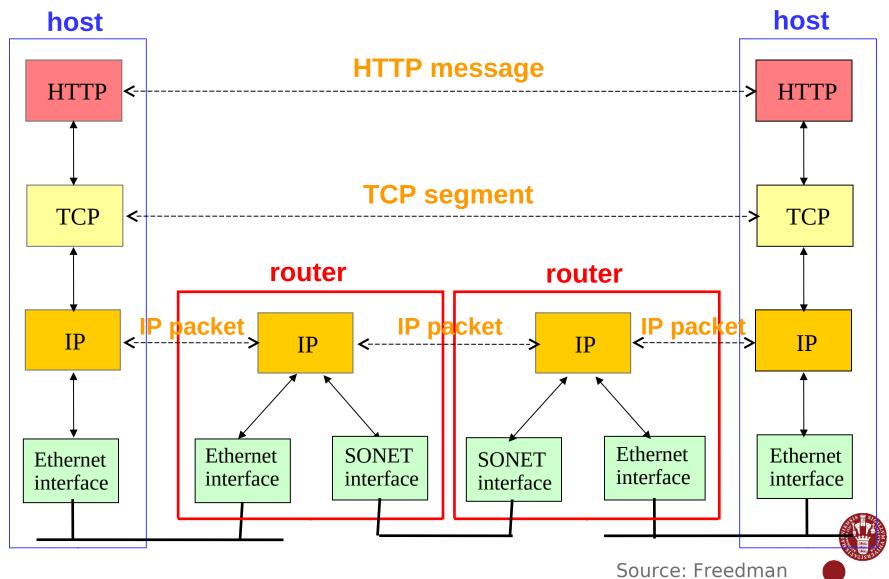
Layer Encapsulation in HTTP







IP Suite: End Hosts vs. Routers



Key Concepts in Networking

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Resource allocation

- Dividing scarce resources among competing parties
- Memory, link bandwidth, wireless spectrum, paths

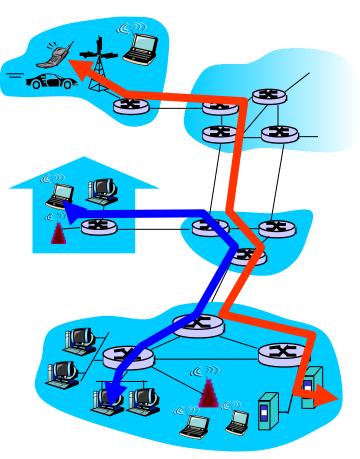
First Example: circuit vs. packet switching



Network Core: Circuit Switching

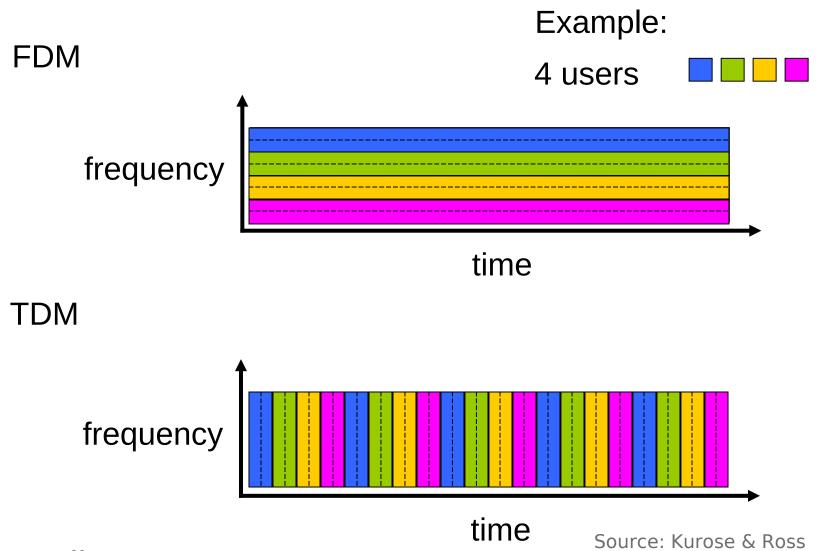
end-end resources reserved for "call"

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required





Circuit Switching: FDM and TDM





Numerical example

- How long does it take to send a file of 640,000 bits from host A to host B over a circuitswitched network?
 - all link speeds: 1536 Mbps
 - each link uses TDM with 24 slots/sec
 - 500 msec to establish end-to-end circuit
 - Note: $1 \text{ Mbps} = 10^6 \text{ bps}$
- Let's work it out!
- Possible answers
 - (a) 500 msec
 - (b) 500.4 msec
 - (c) 510 msec
 - (d) 1 sec



Numerical example

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- Let's work it out!
- Possible answers
 - (a) 500 msec
 - (b) 500.4 msec
 - (c) 510 msec (Assuming only 1 link)
 - (d) 1 sec



Source: Kurose & Ross



Network Core: Packet Switching

each end-end data stream divided into packets

- user A, B packets share network resources
- each packet uses full link bandwidth
- resources used as needed

Bandwidth division into 'pieces'

Dedicated allocation

Resource reservation

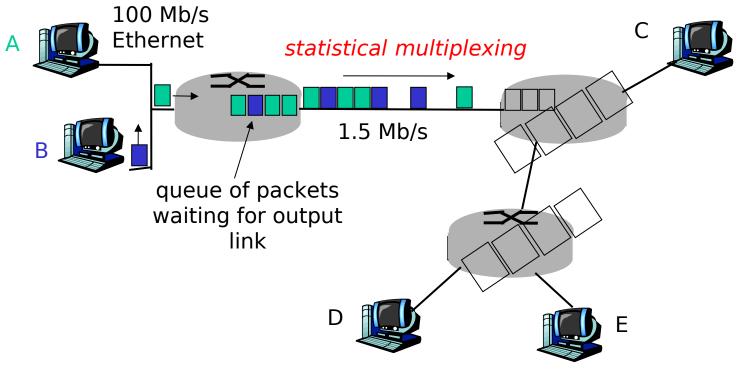
resource contention:

- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward: packets move one hop at a time
 - node receives complete packet before forwarding



Source: Kurose & Ross

Packet Switching: Statistical Multiplexing



- sequence of A & B packets has no fixed timing pattern
 - bandwidth shared on demand: statistical multiplexing.
- TDM: each host gets same slot in revolving TDM frame.



Packet switching versus circuit switching

Packet switching allows more users to use network!

Example:

- 1 Mb/s link
- each user:
 - 100 kb/s when "active"
 - active 10% of time

N users 1 Mbps link

•circuit-switching:

10 users, utilization?

•packet switching:

with 35 users, probability
 10 active at same time
 is less than .0004

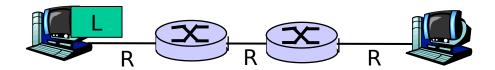
Q: How did we get value 0.0004?

A: $P(N=10) = C(35,10)*P(A)^{10*}(1-P(A))^{(35-10)}$ $P(N>10) \rightarrow \text{sum the above for } N=10..35$

Q: what happens if > 35 users?

Source: Kurose & Ross (partial)

Packet-switching: store-and-forward



- takes L/R seconds to transmit (push out) packet of L bits on to link at R bps
- store and forward: entire packet must arrive at router before it can be transmitted on next link
- delay = 3L/R (assuming zero propagation delay)

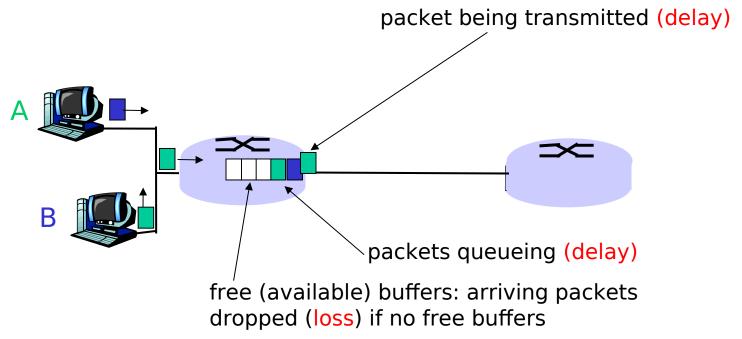
- Example:
 - L = 7.5 Mbits
 - R = 1.5 Mbps
 - transmission delay =
 - (a) 5 sec
 - (b) 10 sec
 - (c) 15 sec
 - (d) 20 sec

more on delay shortly ...



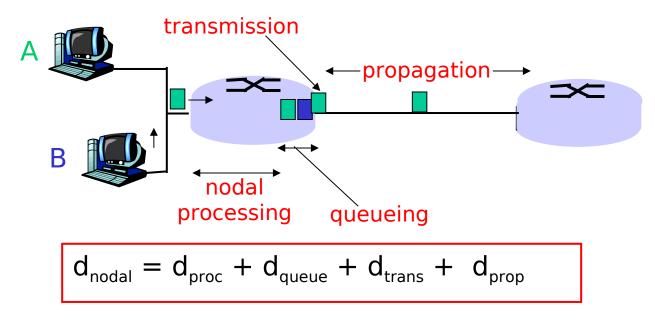
How do loss and delay occur?

- packets queue in router buffers
- packet arrival rate to link exceeds output link capacity
- packets queue, wait for turn





Four sources of packet delay



d_{proc}: nodal processing

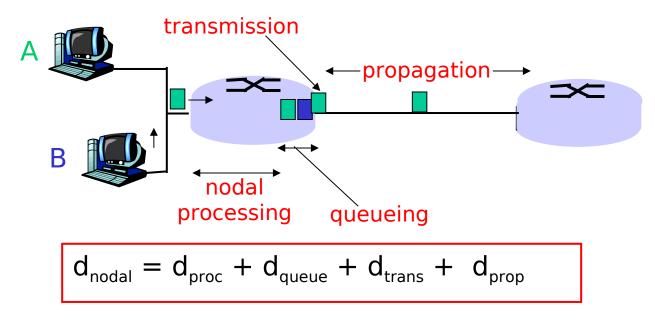
- check bit errors
- determine output link
- typically < msec

d_{queue}: queueing delay

- time waiting at output link for transmission
- depends on congestion level of router



Four sources of packet delay



d_{trans}: transmission delay

- L: packet length (bits)
- R: link bandwidth (bps)

•
$$d_{trans} = L/R$$

$$d_{trans} \text{ and } d_{prop}$$

$$very \text{ different}$$

d_{prop}: propagation delay

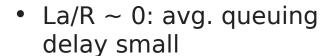
- d: length of physical link
- s: propagation speed in medium (~2x10⁸ m/sec)

$$d_{prop} = d/s$$

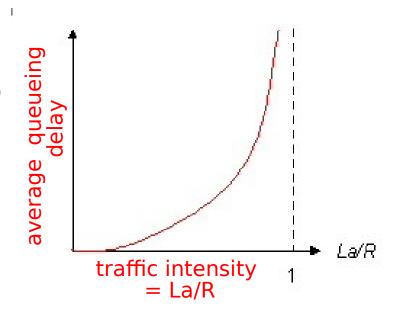


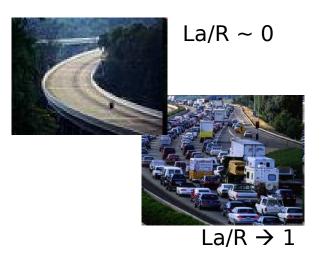
Queuing delay

- R: link bandwidth (bps)
- L: packet length (bits)
- a: average packet arrival rate



- La/R → 1: avg. queuing delay large
- La/R > 1: more "work" arriving than can be serviced, average delay infinite!



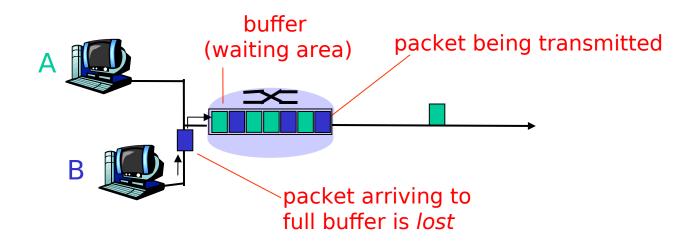




Source: Kurose & Ross

Packet Loss

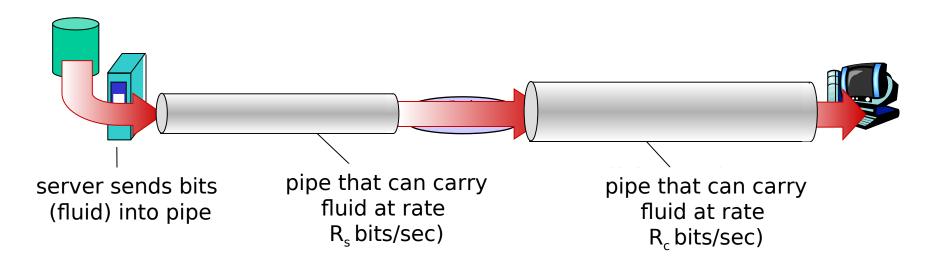
- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be re-transmitted by previous node, by source end system, or not at all





Throughput

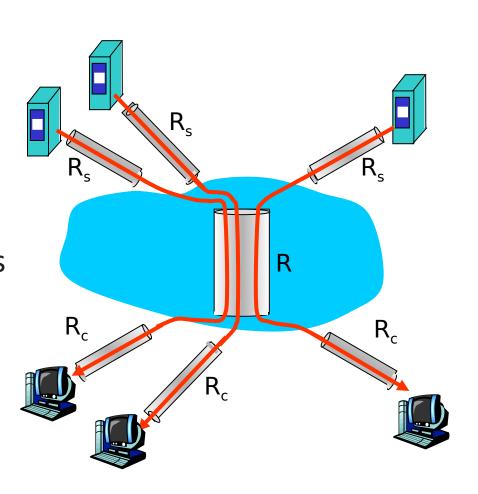
- throughput: rate (bits/time unit) at which bits transferred between sender/receiver
 - instantaneous: rate at given point in time
 - average: rate over longer period of time





Throughput: Internet scenario

- per-connection endend throughput: min(R_c, R_s, R/10)
- in practice: R_c or R_s is often bottleneck



10 connections (fairly) share backbone bottleneck link R bits/sec



Summary

- Internet is a network of networks
 - Inter-operability, power to the edges
- Key concepts in networking
 - Protocols, layers, resource allocation, and naming
- Circuit switching
 - FDM
 - TDM
- Packet switching
 - Store-and-forward
 - Delay, loss, throughput



What's next? What to program on the network?









What's next? How to program the network?

