Computer Systems, B1-2 2023-24

Introduction

Michael Kirkedal Thomsen

DIKU, September 4, 2023

Course structure

Representing information as bits

Everything is bits
Bit-level manipulation

Integers

Overall outline

```
Week 36-37 Data representation and machine model
   Week 38 C programming
Week 39-40 Memory and operating systems
   Week 41 Concurrent and parallel programming
   Week 47 Fall break
Week 43-45 Computer networks - application and transport layer
   Week 46 No activities (reexam week)
   Week 47 Computer networks - security and efficiency
Week 48-51 Machine architecture
   Week 52 Christmas vacation
  Week 1-2 Computer networks and security - network and link layer
    Week 4 4-hour written exam
```

Lectures

- Mondays 13:15-15:00
- Wednesdays 10:15-12:00

Lecturers



Michael Kirkedal Thomsen: Course root, Networks and Security



Finn Schirmer Andersen: Computer Architecture



Teaching Material

- COD Computer Organization and Design (RISC-V Edition), David A. Patterson and John L. Hennesy, second edition, ISBN: 978-0-12-820331-6
 - KR Computer Networking: A Top-Down Approach, James F. Kurose and Keith W. Ross, Pearson, 8th and Global Edition, ISBN 13: 978-1-292-40546-0 (This book will not be used before December)-7th edition is also acceptable
 - JG Modern C, Jens Gustedt,
 https://hal.inria.fr/hal-02383654/document
- OSTEP Operating Systems: Three Easy Pieces, Remzi H. Arpaci-Dusseau and Andrea C. Arpaci-Dusseau,

https://pages.cs.wisc.edu/~remzi/OSTEP/

?? Some notes and book chapters that will be made available through the detailed course schedule

COD is (and KR will be) available at Academic Books at Panum (http://www.academicbooks.dk/) and Polyteknisk Boghandel at Biocenteret (http://www.polyteknisk.dk/).

TAs

TAs:

- Axel Kanne
- Christian Franck
- Jakob Holst Svenningsen
- Jóhann Utne
- Kjartan Martin Johannesen
- Lars Peter Jeppesen
- Lucas Haahr Yri

TAs will gladly help with

- Group members
- The right way to the administration
- A fellow student that can answer questions (or help find the answers)

Exercises and Assignment Cafées

Exercises

- Mondays 15:15-17:00
- Wednesdays 13:15-15:00

Exercises are only for posted exercises. Work on the exercises as they will prepare you both for the exam and assignments.

Cafés

- Wednesdays 15:15-17:00
- Fridays 13:15-15:00

Cafés are primarily for help with assignments.

Details: https://github.com/diku-compSys/compSys-e2022-pub. Also on Discord. See Absalon/Modules.

Groups

Size

2-3 student advised. 1 can be accepted but not recommended. More than 3 is only allowed is on special circumstances

- Sign up for classes with your group-mates on Absalon
- If you need one or more members
 - Look on announcements for details (use Absalon and Discord)
 - ► TA can help

Assignments

- There are 7 assignment in total during the course with deadline roughly every week or second week (all Sundays). The assignments will be evaluated with points.
- Assignments will be awarded zero to 4 points.
- You are required to achieve at least 50 % of the total number of points (equal to 12).
- Also we will require that you achieve points in each the of topics of the course to ensure that you have touched all parts of the curriculum.
- Assignments are made to be solved in groups of 2-3 students, but you can also do them alone.

Assignment rules

The Fundamental Principle of Group Assignments

Each group must make their own solution.

This means

- You can talk with other people about the assignments: Teachers, TAs, other students, etc.
- You cannot share written code with other groups.
- You are not allowed to use code that you did not write yourself without proper citation.
- You cannot share written text with other groups.
- You are not allowed to use text of material without proper citation
 - ► This also includes material provided on the course.
- You are only allowed to use ChatGPT and other AI as any other person.

Assignments vs. exercises

- Note! Both are equally important
- Assignments:
 - Seek to test learning goals that relates to implementation and development of computer systems.
 - Do not fully prepare you for the written exam.
- Exercises:
 - Help you understand the theoretical parts of the material.
 - Prepare you for part of the exam.

Tools

- RARS RISC-V simulator
- C compiler gcc (clang on macOS)
- C debugger gdb (lldb on macOS)
- Profiler Valgrind (perhaps leaks on macOS)
- You can also install all tools on you laptop
 - Linux: most available though apt
 - macOS: most available though Homebrew
 - ► Windows: Windows Subsystem for Linux
- Set up your tool chain
 - recommended using git to share code and reports in your group
 - Sign-up at GitHub today and apply for the Student Developer Pack
 - https://education.github.com/
- Tool-site is available on GitHub

Exam

- A 4-hour written exam; Jan 24 2022.
- The exam will be a BYOD-exam.
- The course syllabus is the exercises, assignments and reading material.
- Previous exams will available.

Questions?

Course structure

Representing information as bits Everything is bits

Bit-level manipulation

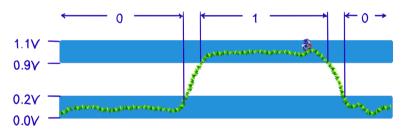
Everything is bits

- Each bit is 0 or 1.
- By interpreting sets of bits in various ways...
 - ...computers determine what to do.
 - ...represent and manipulate numbers, sets, strings data.

Why bits? Why not decimals? Could it have been some other way?

Everything is bits

- Why bits? Electronic implementation.
 - Easy to store with bistable elements.
 - Reliably transmitted on noisy and inaccurate wires (error correction).



- ... But there exist models that do not use bits.
 - ► The Soviet Setun computer used ternary *trits*.
 - Quantum computers use qubits that are in a superposition of the two states.
 - ...error correction is the main challenge here.

Binary numbers

Base 2 number representation.

- Represent 15213₁₀ as 11101101101101₂
- Represent 1.20₁₀ as 1.0011001100110011[0011]...₂
- Represent 1.5213×10^4 as $1.1101101101101_2 \times 2^{13}$

Machine numbers are of some finite size.

- If we use k bits to represent a number, only 2^k distinct values are possible.
- How we interpret those bits can vary.
- Why do we use finite-sized numbers?

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- How we interpret those bits can vary.
- Why do we use finite-sized numbers?
- ▶ A "k-bit machine" handles numbers of up to k bits "natively" (meaning fast).

Encoding byte values

Byte = 8 bits

- (Machine-specific, but is true for all mainstream machines.)
- 256 different values.
- Binary 00000000₂ to 11111111₂.
- Decimal 0₁₀ to 255₁₀.
- Hexadecimal 00₁₆ to FF₁₆.
 - Base 16 number representation.
 - ► Uses characters 0−9 and A−F.
 - ► In C we write FA1D37B₁₆ as
 - ▶ 0xFA1D37B
 - 0xfa1d37b (case does not matter)

Hex	Dec	Bin
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
В	11	1011
С	12	1100
D	13	1101
E	14	1110
F	15	1111

Example sizes of C types on various computers

C data type	Typical 16-bit	Typical 32-bit	Typical 64-bit	x86-64
char	1	1	1	1
short	1	2	2	2
int	2	4	4	4
long	4	4	8	8
int32_t	4	4	4	4
int64_t	8	8	8	8
float	4	4	4	4
double	8	8	8	8
long double	-	-	-	10
pointer	2	4	8	8

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Boolean algebra

Developed by George Boole in 19th century

- Algebraic representation of logic ("truth values").
- Encode true as 1 and false as 0.

And	
& 0	1
0 0	0
1 0	1
	-
Not	_
	_
Not	_

• These operations can be implemented with tiny electronic *gates*.

General boolean algebras

• The truth tables generalise to operate on *bit vectors*, applied elementwise.

	01101001	01101001		01101001		
&	01010101	01010101	^	01010101	~	01101001
	01000001	01111101		00111100		10010110

This is the form they take in programming languages such as C.

Bit-level operations in most C-like languages

Operations &, |, ~, ^ available in C.

- Apply to any integral type.
 - ► E.g. long, int, short, char...
- Interpret operands as bit vectors.
- Applied bit-wise.

Examples

- $\sim 0 \times 41 = 0 \times BE$
 - $\sim 01000001_2 = 101111110_2$
- $\sim 0 \times 00 = 0 \times FF$
 - $\sim 0000000002 = 1111111112$
- 0x69 & 0x55 = 0x41
 - \triangleright 01101001₂ & 01010101₂ = 01000001₂
- 0x69 & 0x55 = 0x7D
 - \triangleright 01101001₂ & 01010101₂ = 01111101₂

Shift operations

Left shift x << y</p>

- ► Shift bit-vector x left by y positions.
 - Throws away excess bits on the left.
 - Fills with zeroes on right.

Right shift x >> y

- Shift bit-vector x right by y positions.
 - Throws away excess bits on the left.
- Logical shift: Fill with 0s on left.
- Arithmetic shift: Replicate most significant bit on left.

Undefined behaviour

Shifting a negative amount or by the vector size or more.

X		01100010
x << 3		00010000
x >> 2	(log)	00011000
x >> 2	(arith)	00011000
x		10100010
x << 3		10100010
	(log)	
x << 3		00010000

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Encoding integers

Suppose x_i is the *i*th bit of a *w*-bit word (with x_0 being the least significant bit).

Unsigned

Two's complement

$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

$$B2S(X) = -x_{w-1} \cdot 2^{w-1} + \sum_{i=0}^{w-2} x_i \cdot 2^i$$

$$int16_t x = 15213;$$

$$int16_t y = -15213;$$

	Decimal			Н	ex				Binary
Х	15213	3	В	5	D	0011	1011	0110	1101
У	-15213	С	4	9	3	1100	0100	1001	0011

Sign bit

- For 2's complement, most significant bit (x_{w-1}) indicates sign.
 - ▶ 0 for non-negative.
 - ▶ 1 for negative.

Two's complement encoding example

```
int16_t x = 15213; // 0011 1011 0110 1101 int16_t y = -15213; // 1100 0100 1001 0011
```

Weight	1	.5213		-15213	
1	1	1	1	1	
2	0	0	1	2	
4	1	4	0	0	
8	1	8	0	0	
16	0	0	1	16	
32	1	32	0	0	
64	1	64	0	0	
128	0	0	1	128	
256	1	256	0	0	
512	1	512	0	0	
1024	0	0	1	1024	
2048	1	2047	0	0	
4096	1	4096	0	0	
8192	1	8192	0	0	
16384	0	0	1	16384	
-32768	0	0	1	-32768	
Sum		15213		-15213	

Let's play a game

http://topps.diku.dk/compsys/integers.html

Numeric ranges, here for 16-bit signed and unsigned integers

Unsigned

Two's complement signed

Values for w = 16:

	Decimal	Hex	Binary		
UMax	65535	FFFF	1111 1111 1111 1111		
SMax	32767	7 F F F	0111 1111 1111 1111		
SMin	-32768	8 0 0 0	1000 0000 0000 0000		
-1	-1	FFFF	1111 1111 1111 1111		
0	0	0 0 0 0	0000 0000 0000 0000		

Values for different word sizes

	W							
	8	16	32	64				
UMax	255	65,535	4,294,967,295	18,446,744,073,709,551,615				
SMax	127	32,767	2,147,483,647	9,223,372,036,854,775,807				
SMin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808				

Observations

$$|\mathsf{SMin}| = \mathsf{SMax} + 1$$

 $|\mathsf{UMax}| = 2 \cdot \mathsf{SMax} + 1$

Note the assymetric range.

C Programming

- #include <limits.h>
- Declares constants, e.g:
 - ► ULONG_MAX
 - ► LONG_MAX
 - LONG_MIN
- Values are platform-specific.

Unsigned and signed numeric values (here w = 4)

X	DZU(X)	D23(X)
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

R211(v)

B)C(v)

Equivalence

Same encoding for non-negative values.

Uniqueness

- Every bit pattern represents distinct integer value.
- Each representable integer has unique bit encoding.
- ► The representation is **bijective**.

Can invert mappings

- \triangleright $U2B(x) = B2U^{-1}(x)$
 - Bit pattern for unsigned integer.
- \triangleright $S2B(x) = B2S^{-1}(x)$
 - Bit pattern for two's complement integer.

Main takeaways

- Distinguish between representation and interpretation.
- Low-level values do not describe their own structure.
- Everything is built in layers.
- A good computer scientist adds new, clean, layers of abstraction.
 - A bad one adds layers that hide without simplifying.
 - A terrible one adds layers that complicate and obfuscate.
- The point of this course is to show that there is no magic, only the work of careful people who put in a lot of effort.